

Research Article

A Collection-Distribution Center Location and Allocation Optimization Model in Closed-Loop Supply Chain for Chinese Beer Industry

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Recycling waste products is an environmental-friendly activity that can result in manufacturing cost saving and economic efficiency improving. In the beer industry, recycling bottles can reduce manufacturing cost and the industry's carbon footprint. This paper presents a model for a collection-distribution center location and allocation problem in a closed-loop supply chain for the beer industry under a fuzzy random environment, in which the objectives are to minimize total costs and transportation pollution. Both random and fuzzy uncertainties, for which return rate and disposal rate are considered fuzzy random variables, are jointly handled in this paper to ensure a more practical problem solution. A heuristic algorithm based on priority-based global-local-neighbor particle swarm optimization (pb-glnPSO) is applied to ensure reliable solutions for this NP-hard problem. A beer company case study is given to illustrate the application of the proposed model and to demonstrate the priority-based global-local-neighbor particle swarm optimization.

1. Introduction

Due to resource scarcity and environmental concerns, responsible companies are beginning to pay attention to the future of the planet and the global environment. Recycling used products for remanufacturing is, therefore, becoming of greater importance in supply chain management, a move that can dramatically reduce carbon emissions [1]. Closedloop supply chain (CLSC) combines the forward supply chain with a reverse supply chain to cover the whole product life cycle [2], with the manufacturing of new products and the transportation to customers via distribution centers and retailers as the forward supply chain and recycling, sorting, disposal, and remanufacturing as the reverse supply chain. In recent years, the CLSC has received a great deal of academic and business attention because of the need to be socially responsible, global environmental concerns, and government legislation [3, 4], all of which have motivated companies to pay more attention to recycling to reduce costs and lessen their carbon footprint.

Facility location and allocation problems (FLAPs) have been widely studied. Subramanian et al. [5] developed

priority-based simulated annealing to solve a CLSC network design problem, in which the distribution center (DC) and the centralized return collection center (CC) were set. Amin and Zhang [6] presented facilities location model for manufacturing and remanufacturing plants and CLSC collection centers, which included demand and return uncertainties. Subulan et al. [7] developed a CLSC network design model for the lead/acid battery industry that considered both financial and collection objectives. CLSC network design in a competitive environment with price-dependent demand was examined by Rezapour et al. [4], in which the DC and CC were separately built. Zeballos et al. [8] proposed a model for a multiperiod CLSC design and planning problem with demand uncertainty that had ten echelons in which the DC and CC were considered. Wang et al. [9] developed a granular robust model for a two-stage waste-to-energy feedstock flow planning problem with uncertain capacity expansion costs. Tokhmehchi et al. [10] developed a hybrid approach to solve a closed-loop supply chain location and allocation problem that considered the minimization of total cost. Vahdani and Mohammadi [11] proposed capacitated bidirectional facilities for CLSC conduct distribution, in which a multipriority queuing system was studied. As a growing number of companies are now engaging in recycling activities due to economic and environmental concerns, distribution and collection activities using the same vehicle have been found to reduce carbon emissions and transportation costs as empty loads can be avoided. In this paper, to benefit company operation and reduce construction costs, a distribution center (DC) is combined with a collection center (CC) into a collection-distribution center (CDC). In practice, as the recycled product owners are usually at the same location as the potential new product buyer [12], a DC/CC combination has lower construction and operating expenses and can significantly reduce environmental pollution.

Ramkumar et al. [13] developed a multiechelon, multiperiod, multiproduct closed-loop supply chain network model which was solved using a genetic algorithm with fixed variables. Kaya and Urek [12] presented a facility locationinventory-pricing model without uncertainty to determine optimal facilities locations. Barz [14] proposed an optimization model for a two-stage capacitated facility location and allocation problem with additive manufacturing, in which all variables were certain. Jindal and Sangwan [15] developed a multiobjective model for a CLSC network design problem with the economic and environmental factors being fuzzy uncertain variable and the DC and CC were separate. Ramezani et al. [16] conducted research into a CLSC network design problem that only considered of fuzzy variables. In recent years, uncertainty has attracted more research attention [17-19]. Stochastic programing, robust optimization, and fuzzy set theory have been used to present uncertainty in FLAPs [20, 21]. Wang et al. [22] used prediction sets to solve an expansion planning problem for waste-to-energy (WtE) systems facing future waste supply uncertainty. Keyvanshokooh et al. [23] proposed a novel hybrid robust-stochastic programing (HRSP) approach to simultaneously model two different types of uncertainties by using stochastic scenarios for the transportation costs and polyhedral uncertainty sets for the demand and returns. However, the DC and the CC were separate and the collection disposal rate was a certain variable.

Uncertainties exist in both forward and reverse supply chains. However, the uncertainties in the reverse flow are higher than in the forward supply chain [7, 24, 25] as returned product quantity is generally seen as uncertain [23, 26]. Subjective uncertainties such as the decision maker's choices and environmental coefficients can be dealt with using fuzziness, while objective uncertainties such as unit transportation costs, product prices, and the quantity of unusable products can be dealt with using randomness. In this paper, to reflect the study problem, the return rate and disposal rate are considered fuzzy random variables, which are concurrently handled using triangular fuzzy numbers [7]. Based on the above, a model is formulated to determine the optimum CDC number and location and the allocation strategies for the different facility types.

As facilities location and allocation problems are seen as nonconvex, nondifferentiable, strongly NP-hard problems, a collection and distribution center location and allocation problem (CDCLAP) in a closed-loop supply chain under a fuzzy random environment is even more complicated. Several different methods have been used to solve NP-hard problems [27-29]. While particle swarm optimization (PSO) has been found to be generally effective [30–32], when the local optimal solution is found, the particle behavior in a basic PSO is directly influenced and therefore frequently falls into a local optimum [33–35]. Because of this problem, several advanced PSOs have been developed to more accurately solve supply chain management problems. Ai and Kachitvichyanukul [36] proposed a global-local-neighbor PSO which was found to be more effective. Based on this innovation, Xu et al. [33] proposed a fuzzy random simulation-based bilevel globallocal-neighbor particle swarm optimization (frs-bglnPSO). In this paper, a priority-based global-local-neighbor particle swarm optimization (pb-glnPSO) is applied to solve the CDCLAP.

In summary, this paper proposes a mathematical model to solve a collection-distribution center location and allocation problem in a closed-loop supply chain that considers economic and environmental factors and includes fuzzy random variables for return and disposal rates. The remainder of this paper is organized as follows. Section 2 presents the problem statement and model assumptions, after which a description of the model and its formulations is given in Section 3. The proposed hybrid solution based on the developed pb-glnPSO is described in Section 4 and case study is presented in Section 5 to illustrate model formulation and the proposed method. Finally, Section 6 gives conclusions and indications for future research extensions.

2. Research Problem Statement

In this paper, a company with factories in certain locations and retailers in different customer zones is considered. The company needs to decide the locations for their integrated collection and distribution centers (CDCs), at which both a used product collection network and a new product distribution network are to be jointly [12]. As CDCs reduce construction and transportation costs because the same vehicles are used for both distribution and recycling, in this paper, only CDCs are considered.

A general illustration of the classical CDCLAP for a closed-loop supply chain is shown in Figure 1, with the CLSC framework shown in Loop 1. The CLSC framework has four echelons: factories, CDCs, retailers, and disposal centers [11]. The forward supply chain begins with new production, after which the finished products are transported from the factories to the retailers via the CDCs. In the reverse supply chain, returned products are collected and transported to the CDCs, where the recycled products are inspected, consolidated, and sorted into those that are available for remanufacturing, which are sent to the factories, and those that are unsuitable for remanufacturing, which are transported to the disposal centers [23]. A CDC supplies products to multiple retailers; however, retailer demand is fulfilled by only one production site. A CDC can also handle products from different factories and dispatch returned products to multiple factories for remanufacturing.



FIGURE 1: The closed-loop supply chain network.

In the CLSC examined in this paper, the retailers' demand is estimated based on preorders. However, the return rate is considered to be a fuzzy random variable as customers may not return the used product or the product may be broken. Consequently, the availability of recycled products is unsure because of unsure transportation and carrying losses. Another fuzzy random variable considered in this paper is the returned product disposal rate, which is dependent on the inspection and consolidation at the CDC.

The assumptions for the proposed problem investigation are as follows: (1) only one product in one period is considered; (2) all alternative CDC locations have been identified; (3) recycling a used product costs less than manufacturing a new one; (4) the CDCs and factories have capacity limits [37– 39]. As considering incapacitated facilities is an unrealistic assumption in many LAPs, many researchers have assigned a maximum capacity level to facilities to model more realistic decisions; (5) the factories', retailers', and disposal centers' locations are known; (6) new product and returned product storage are allowed at the CDCs [21].

The initial problem is deciding which CDCs to select from the candidate sites and which allocation strategies to select to minimize total CDC costs: operating costs, transportation costs, and transportation pollution costs, while also considering flow constraints, capacity limits, and retailer demand.

3. Modelling

In this section, the mathematical formulations are given for the CDCLAP in the CLSC and the notations are given in the Notations to facilitate the problem description.

3.1. Objective Functions. Based on the variables mentioned in the Notations, the objectives are to minimize total costs

and to minimize the environmental effects with the primary objective being minimizing total cost.

3.1.1. Economic Objective. In general, decision makers seek to minimize total costs, which are made up of transportation costs, fixed costs, and operating costs. The minimization objective can be described as

min F_1

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \tilde{\overline{a}}_{k}\right) \\ + \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \tilde{\overline{b}}_{i} \tilde{\overline{a}}_{k} Q_{ik} \\ + \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \tilde{\overline{a}}_{k} Q_{ik} \left(1 - \tilde{\overline{b}}_{i}\right) + \sum_{i=1}^{I} F_{i}^{c} X_{i} \\ + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} RV_{i}^{c} \tilde{\overline{a}}_{k}.$$
(1)

Equation (1) calculates the total cost, in which $\sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji}$ is the new product transport costs from the factories to the CDC, $\sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} (1 + \tilde{a}_{k})$ is the transport costs between the CDCs and the retailers, and $\sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \tilde{b}_{i} \tilde{a}_{k} Q_{ik}$ is the returned product delivery costs from the CDCs to the disposal centers. The returned product transportation costs from the CDCs to disposal centers are measured as $\sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \sum_{n=1}^{P} C_{ij}^{p} \tilde{a}_{k} Q_{ik} (1 - \tilde{b}_{i})$. The fixed costs for opening a new CDC are $\sum_{i=1}^{I} F_{i}^{c} X_{i}$. $\sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji}$ denotes the new product variable costs.

 $\sum_{i=1}^{I} \sum_{k=1}^{K} \mathrm{RV}_{i}^{c} \widetilde{\overline{a}}_{k} \text{ calculates the returned product operating costs.}$

As it is very difficult to handle objective functions with fuzzy random factors, Kruse and Meyer [40] demonstrated that the fuzzy expected value could be represented by a single fuzzy number. Based on the theory proposed by Heilpern [41], without a loss of generality, the expected value operator is used to convert the uncertain model into a deterministic model, which can then be used to transform the fuzzy random objective functions into their crisp equivalences, as shown in

min F_1

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right) \\ + \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
(2)
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right) \\ + \sum_{i=1}^{I} F_{i}^{c} X_{i} + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \text{RV}_{i}^{c} \text{EV}\left[\tilde{\tilde{a}}_{k}\right].$$

Note that $\text{EV}[\tilde{a}_k]$ or $\text{EV}[\tilde{b}_i]$ above represents two expected values: the first is used to convert the fuzzy random variables into fuzzy numbers based on Kruse and Meyer's 1987 theory, and the second is used to transform the fuzzy numbers into deterministic numbers based on Heilpern's 1992 theory.

3.1.2. Environmental Objective. The secondary objective is to minimize the transportation carbon emissions associated with the CLSC operations, an area that has attracted significant recent research attention [42]. The following expression represents the transportation carbon emissions between the CDCs and the factories, the CDCs and the retailers, and the CDCs and the disposal centers.

min F_2

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} \beta_{ij} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \beta_{ik} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right) + \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \beta_{in} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} + \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} \beta_{ij} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right).$$

$$(3)$$

 $\sum_{i=1}^{I} \sum_{j=1}^{J} \beta_{ij} P_{ji}$ is the environmental pollution caused by the transportation activities from the factories to the CDCs. $\sum_{i=1}^{I} \sum_{k=1}^{K} \beta_{ik} Q_{ik} (1 + \text{EV}[\tilde{a}_k])$ is the summation of the carbon footprints for transporting products between the CDCs and retailers. $\sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \beta_{in} \text{EV}[\tilde{b}_i] \text{EV}[\tilde{a}_k] Q_{ik}$ is the total carbon footprint from the CDCs to the disposal centers, and $\sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \sum_{n=1}^{N} \beta_{ij} \text{EV}[\tilde{\tilde{a}}_{k}]Q_{ik}(1 - \text{EV}[\tilde{\tilde{b}}_{i}])$ is the carbon footprint from the CDCs to factories for returned products.

3.2. Constraints. Note that as each CDC has its own capacity limit, it is unable to service goods beyond capacity; therefore, the capacity limit restriction can be written as follows:

$$\sum_{k=1}^{K} \operatorname{EV}\left[\widetilde{\widetilde{a}}_{k}\right] Q_{ik} + \sum_{j=1}^{J} P_{ji} \le \alpha_{i} \quad \forall i \in \Omega.$$
(4)

 \overline{a}_{ki} is a fuzzy random variable indicating the used product return rate transported from the retailer k to CDC i. Q_{ik} indicates the product quantity from CDC i to retailer k. P_{ji} indicates the product quantity transported from factory j to CDC i. α_i is the capacity of CDC i.

Within the capacity constraint, the factory is able to manufacture new products to meet retailer needs as well as deal with the returned products from the CDCs.

$$\sum_{k=1}^{K} D_{k} + \sum_{i=1}^{I} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\overline{b}}_{i}\right]\right)$$

$$\leq \sum_{i=1}^{J} \gamma_{j}.$$
(5)

 D_k is the demand from retailer k and γ_j is the capacity of factory j. $\sum_{i=1}^{I} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}[\tilde{a}_k] Q_{ik}(1 - \text{EV}[\tilde{b}_i])$ calculates the returned product quantity transported to factory j for remanufacturing.

All products for the retailers are processed through the CDCs. The recycled product quantity is always less than the product quantity transported from the factory to the CDC, which can be described as follows:

$$\sum_{j=1}^{J} P_{ji} \ge \sum_{k=1}^{K} \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik}.$$
(6)

 P_{ji} is a variable indicating the new product quantity transported from factory *j* to CDC *i*. $\tilde{a}_k Q_{ik}$ refers to the returned product quantity transported from retailer *k* to CDC *i*.

The products provided to the retailers should meet their demand.

$$\sum_{i=1}^{I} Q_{ik} \ge \sum_{k=1}^{K} D_k.$$
 (7)

 Q_{ik} is the product quantity CDC *i* sends to retailer *k*. The stochastic variable D_k is the retailer *k*'s demand based on order quantity.

The returned product quantity transported to the CDCs is more than the product quantity transported to disposal centers.

$$\sum_{k=1}^{K} \operatorname{EV}\left[\widetilde{\widetilde{a}}_{k}\right] Q_{ik} \geq \sum_{n=1}^{N} \sum_{k=1}^{K} \operatorname{EV}\left[\widetilde{\widetilde{b}}_{i}\right] \operatorname{EV}\left[\widetilde{\widetilde{a}}_{k}\right] Q_{ik}.$$
 (8)

 $EV[\tilde{\bar{a}}_k]Q_{ik}$ is the expression for the returned product quantity from retailer k to CDC i, and $EV[\overline{b}_i]EV[\overline{a}_k]Q_{ik}$ is the product quantity transported from CDC *i* to retailer *k*.

There should be at least one CDC and there should be no more CDCs than the specified upper limit.

$$1 \le \sum_{i=1}^{I} x_i, \tag{9}$$

$$\sum_{i=1}^{I} x_i \le U. \tag{10}$$

U is the upper limit for the number of CDCs, which is dependent on demand, returned product quantity, and fixed capacity constraints.

Each retailer must be serviced by only one CDC.

$$\sum_{i=1}^{I} y_{ik} = 1.$$
(11)

Since x_i and y_{ik} are binary variables, the following constraints are required:

$$x_i = \{0, 1\}, \quad \forall i \in \Omega,$$

$$y_{ik} = \{0, 1\}, \quad \forall i \in \Omega, \quad \forall k \in \Phi.$$
(12)

 x_i is a binary variable indicating whether a CDC is opened at point *i*. If location *i* is chosen to open a CDC, then x_i ; otherwise, $x_i = 0$. y_{ik} is a binary variable indicating whether retailer k is serviced by CDC i. If $y_{ik} = 1$, then retailer k is serviced by CDC *i*; otherwise, $y_{ik} = 0$.

3.3. Global Model. From the formulation above, the model for the CDCLAP with capacity, flow, and quantity constraints is developed with the aims of minimizing total costs and total transportation pollution with the primary objective of minimizing the total cost. In the CLSC, both new and returned products are considered. The product can be reproduced to save raw materials and reduce waste and pollution. In our model, all costs involved in the CDCLAP are considered as well as the influence of the transportation activity pollution. Fuzzy random theory is used to deal with the real world complex uncertainties and ensure more scientific decisions. Therefore, this CDC situation is closer to the real situation as it can deal with complicated practical problems. Finally, the global model is given:

min F_1

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right)$$

$$+\sum_{i=1}^{I} F_{i}^{c} X_{i} + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \mathrm{RV}_{i}^{c} \mathrm{EV}\left[\widetilde{\overline{a}}_{k}\right]$$

min F_2

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} \beta_{ij} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \beta_{ik} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \beta_{in} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} \beta_{ij} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right)$$

s.t.
$$\sum_{k=1}^{K} \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik} + \sum_{j=1}^{J} P_{ji} \le \alpha_{i} x_{i}$$

$$\forall i \in \Omega \ \forall j \in \Psi \ \forall k \in \Phi$$

$$\sum_{k=1}^{K} D_{k} + \sum_{i=1}^{I} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV} \left[\tilde{\tilde{a}}_{k} \right] Q_{ik} \left(1 - \text{EV} \left[\tilde{\tilde{b}}_{i} \right] \right)$$

$$\leq \sum_{j=1}^{J} \gamma_{j} \quad \forall j \in \Psi \; \forall k \in \Phi \; \forall n \in \Upsilon$$

$$\sum_{j=1}^{J} P_{ji} \geq \sum_{k=1}^{K} \text{EV} \left[\tilde{\tilde{a}}_{k} \right] \quad \forall i \in \Omega \; \forall j \in \Psi \; \forall k \in \Phi$$

$$\sum_{k=1}^{K} \text{EV} \left[\tilde{\tilde{a}}_{k} \right] Q_{ik} \geq \sum_{n=1}^{N} \sum_{k=1}^{K} \text{EV} \left[\tilde{\tilde{b}}_{i} \right] \text{EV} \left[\tilde{\tilde{a}}_{k} \right] Q_{ik}$$

$$\forall i \in \Omega \; \forall k \in \Phi \; \forall n \in \Upsilon$$

$$\forall i \in \Omega \ \forall k \in \Phi \ \forall n \in \Upsilon$$

$$\sum_{i=1}^{I} Q_{ik} \ge \sum_{k=1}^{K} D_k \quad \forall i \in \Omega \quad \forall k \in \Phi$$

$$1 \le \sum_{i=1}^{I} x_i \quad \forall i \in \Omega$$

$$\sum_{i=1}^{I} x_i \le U \quad \forall i \in \Omega$$

$$y_{ik} \le x_i \quad \forall i \in \Omega, \quad \forall k \in \Phi$$

$$\sum_{i=1}^{I} y_{ik} = 1 \quad \forall i \in \Omega \quad \forall k \in \Phi$$

$$x_i = \{0, 1\} \quad \forall i \in \Omega$$

$$y_{ik} = \{0, 1\} \quad \forall i \in \Omega, \quad \forall k \in \Phi.$$
(13)

3.4. Model Transformation. The two objectives are believed to be some similar ones to some extent. The minimization of the environmental objective requires low transportation carbon emissions which could be similar to the transportation costs

of the economic objective, which are mainly affected by transportation distance. Based on previous research [33, 43, 44], the secondary environmental impact reduction objective can be transformed into a constraint to reduce the complexity of the multiple-objective model (13). By the acceptable carbon emissions level, the decision makers are concerned much more about the total cost, namely, the primary objective. Suppose that there is a maximum average environmental impact level μ which is acceptable to the decision maker; therefore, the secondary objective can be transformed to be a constraint as follows:

$$F_{2} = \sum_{i=1}^{I} \sum_{j=1}^{J} \beta_{ij} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \beta_{ik} Q_{ik} \left(1 + \text{EV}\left[\tilde{\overline{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} \beta_{in} \text{EV}\left[\tilde{\overline{b}}_{i}\right] \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} \beta_{ij} \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\overline{b}}_{i}\right]\right)$$
$$\leq \mu.$$
(14)

Namely,

$$F_2 \le \mu. \tag{15}$$

As the economic costs for the primary objective and the secondary environmental objective are transformed into a constraint, the global model can be transformed into its corresponding equivalent model as follows:

min F_1

$$= \sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right)$$
$$+ \sum_{i=1}^{I} F_{i}^{c} X_{i} + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \text{RV}_{i}^{c} \text{EV}\left[\tilde{\tilde{a}}_{k}\right]$$
$$F_{2} \leq \mu$$

s.t. $F_2 \leq \mu$

$$\sum_{k=1}^{K} \operatorname{EV}\left[\widetilde{\widetilde{a}}_{k}\right] Q_{ik} + \sum_{j=1}^{J} P_{ji} \le \alpha_{i} x_{i}$$

$$\forall i \in \Omega \ \forall j \in \Psi \ \forall k \in \Phi$$

$$\begin{split} &\sum_{k=1}^{K} D_{k} + \sum_{i=1}^{I} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\overline{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\overline{b}}_{i}\right]\right) \\ &\leq \sum_{j=1}^{J} \gamma_{j} \quad \forall j \in \Psi \ \forall k \in \Phi \ \forall n \in \Upsilon \end{split}$$

$$\begin{split} \sum_{j=1}^{I} P_{ji} &\geq \sum_{k=1}^{K} \mathrm{EV}\left[\overline{a}_{k}\right] \quad \forall i \in \Omega \ \forall j \in \Psi \ \forall k \in \Phi \\ \sum_{k=1}^{K} \mathrm{EV}\left[\overline{\tilde{a}}_{k}\right] Q_{ik} &\geq \sum_{n=1}^{N} \sum_{k=1}^{K} \mathrm{EV}\left[\overline{\tilde{b}}_{i}\right] \mathrm{EV}\left[\overline{\tilde{a}}_{k}\right] Q_{ik} \\ &\forall i \in \Omega \ \forall k \in \Phi \ \forall n \in \Upsilon \\ \\ \sum_{i=1}^{I} Q_{ik} &\geq \sum_{k=1}^{K} D_{k} \quad \forall i \in \Omega \ \forall k \in \Phi \\ \\ 1 &\leq \sum_{i=1}^{I} x_{i} \quad \forall i \in \Omega \\ \\ \sum_{i=1}^{I} x_{i} &\leq U \quad \forall i \in \Omega \\ \\ \sum_{i=1}^{I} x_{i} &\leq U \quad \forall i \in \Omega, \ \forall k \in \Phi \\ \\ \\ \sum_{i=1}^{I} y_{ik} &= 1 \quad \forall i \in \Omega, \ \forall k \in \Phi \\ \\ \\ x_{i} &= \{0, 1\} \quad \forall i \in \Omega, \ \forall k \in \Phi. \end{split}$$

(16)

4. The Heuristic Algorithms Based on Pgln-PSO

Particle swarm optimization (PSO) is an evolutionary algorithm which simulates social behavior such as birds flocking and fish schooling [45]. Using a fixed population of individuals, the PSO searches the feasible zone to seek solutions, which are then updated to achieve an optimal solution. The particles [46], characterized by their position and velocity, are decided on by their flying experience, their discoveries, or the discoveries of their companions. They fly through the problem spaces following the currently optimum particles to find the best solution between the populations and the best solution for each population. Even though the PSO has been widely used to solve NP-hard problems [45, 47], in the basic PSO, as the particles in the swarm are weak, they tend to cluster rapidly toward the global best particle [31]. However, the global-local-neighbor particle swarm optimization (glnPSO) developed by Ai and Kachitvichyanukul [36] has been found to improve the weakness in the basic PSO. Xu et al. [43] proposed an even more advanced globallocal-neighbor particle swarm optimization with exchangeable particles (GLNPSO-ep). In this section, a prioritybased global-local-neighbor particle swarm optimization (pb-glnPSO) is proposed to solve the CDCLAP in the CLSC.

4.1. Notations for the pb-glnPSO. The basic elements of the PSO are particles, population, velocity, inertia weight, individual best, and global best. The notations needed for the pb-glnPSO are shown in the Notations.

4.2. Encoding and Decoding Algorithm. The decoding process is based on the priority-based encoding developed by Gen and Cheng and the priority-based decoding and encoding proposed by Gen et al. [48]. The priorities for the CDCs and the retailers are equal to the total number of retailers and CDCs. At each step, the CDC (retailer) with the highest priority is selected and connected to a retailer (CDC) under a minimum transportation cost constraint. Procedure 1 shows the decoding algorithm for the priority-based encoding and its trace table, with the priority-based encoding being random. The CDCLAP is solved in two stages [14]. In the first stage, the CDC location is chosen and the transportation between the CDCs and retailers calculated. In the second stage, the allocations between the factories and CDCs are dealt with.

4.3. The Fitness Value Function. The CDCLAP considered in this paper has a primary objective, the minimization of total costs, and a secondary objective, the minimization of environmental pollution, with the fitness value of each particle reflecting the objective value of the total cost. As the second objective has been transformed into a constraint, when the transportation carbon emissions are beyond the acceptable level [33, 43], a penalty function is assigned according to the actual situation. If the transportation carbon emissions are within the acceptable level $F_2 \leq \mu$, the fitness value function in the pb-glnPSO is as follows:

Fitness = min
$$\sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{p} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{p} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right)$$
$$+ \sum_{i=1}^{I} F_{i}^{c} X_{i} + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \text{RV}_{i}^{c} \text{EV}\left[\tilde{\tilde{a}}_{k}\right].$$
$$(17)$$

If the transportation carbon emissions are beyond the acceptable level, a penalty function is assigned and the fitness function is

Fitness = min
$$\sum_{i=1}^{I} \sum_{j=1}^{J} C_{ij}^{P} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} C_{ik}^{d} Q_{ik} \left(1 + \text{EV}\left[\tilde{\tilde{a}}_{k}\right]\right)$$
$$+ \sum_{i=1}^{I} \sum_{n=1}^{N} \sum_{k=1}^{K} C_{in}^{w} \text{EV}\left[\tilde{\tilde{b}}_{i}\right] \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik}$$
$$+ \sum_{i=1}^{I} \sum_{j=1}^{J} \sum_{k=1}^{K} \sum_{n=1}^{N} C_{ij}^{P} \text{EV}\left[\tilde{\tilde{a}}_{k}\right] Q_{ik} \left(1 - \text{EV}\left[\tilde{\tilde{b}}_{i}\right]\right) \quad (18)$$
$$+ \sum_{i=1}^{I} F_{i}^{c} X_{i} + \sum_{i=1}^{I} \sum_{j=1}^{J} V_{i}^{c} P_{ji} + \sum_{i=1}^{I} \sum_{k=1}^{K} \text{RV}_{i}^{c} \text{EV}\left[\tilde{\tilde{a}}_{k}\right]$$
$$+ C \left(\mu - F_{2}\right)$$

in which *C* is a large enough penalty factor.

4.4. Update. Based on the notations of pb-glnPSO mentioned in the Notations and the glnPSO proposed by Ai and Kachitvichyanukul [36], the inertia weight, velocity, and position are updated using the following equation:

$$\omega(\tau) = \omega(T) + \frac{\tau - T}{1 - T} \left[\omega(1) - \omega(T) \right], \tag{19}$$

$$v_{d}^{l}(\tau+1) = \omega(\tau) v_{d}^{l}(\tau) + c_{p}r_{1} \left[p_{ld}^{\text{best}}(\tau) - p_{d}^{l}(\tau) \right] + c_{g}r_{2} \left[p_{gd}^{\text{best}}(\tau) - p_{d}^{l}(\tau) \right] + c_{l}r_{3} \left[p_{gd}^{\text{best}}(\tau) - p_{d}^{l}(\tau) \right] + c_{n}r_{4} \left[p_{gd}^{\text{best}}(\tau) - p_{d}^{l}(\tau) \right],$$

$$p_{d}^{l}(\tau+1) = p_{d}^{l}(\tau) + v_{d}^{l}(\tau+1),$$
(21)

The glnPSO has been widely used in solving NP-hard facilities location and allocation problems.

4.5. Overall Process of the pb-glnPSO. In this paper, the glnPSO presented in Procedure 1 is used to solve a location and allocation problem. Due to uncertainties and environmental changes, a priority-based global-local-neighbor particle swarm optimization (pb-glnPSO) is proposed to solve this model. As the company pays close attention to economic costs, the environmental factor is dealt with as a constraint that has upper limits. The algorithmic details are as follows.

Step 1. Initialize *P* particles as a swarm: l = 1, 2, ..., L (the particle is the priority).

Step 2 (constraints check). If in the feasible region, go to Step 3; otherwise, return to Step 1.

Step 3. Calculate the fitness according to the decoding algorithm in Procedure 1.

Step 4. Update the particle positions and velocities.

Step 4.1. Acquire the expected value for Z from the above algorithm.

Step 4.2. For l = 1, 2, ..., L, decode each particle to an installment group. Calculate the fitness value of each particle and set the position of the *l*th particle as its personal best. The global best position is chosen from these personal best positions.

Step 4.3 (update pbest). For l = 1, 2, ..., L, if Fitness $(P_l) <$ Fitness (P_l^{best}) , $P_l^{\text{best}} = P_l$.

Step 4.4 (update gbest). For l = 1, 2, ..., L, if Fitness $(P_l) <$ Fitness $(P_q^{\text{best}}), P_q^{\text{best}} = P_l$.

Step 4.5 (update *l*best). For l = 1, 2, ..., L, among all *p*best of *M* neighbors around the *l*th particle, set the personal best which has the best fitness value as P_l^{Lbest} .

Input. Ω : set of CDCs, $i \in \Omega = \{1, 2, 3, \dots, I\}$; Φ : set of retailers, and $k \in \Phi = \{1, 2, 3, \dots, K\}$ D_k : demand of retailer $k, k \in \Phi$, α_i : the capability of the CDCs $i, i \in \Omega$; EV $[\tilde{a}_k]$: the return rate of retailer $k, k \in \Phi$ C_{ik}^d : unit transportation cost between the CDC *i* and retailer $k, i \in \Omega, k \in \Phi$; p(i + k): the priority settled, $i \in \Omega = \{1, 2, 3, ..., I\}, k \in \Phi = \{1, 2, 3, ..., K\}$ *Output*. Q_{ik} : the product quantity transported from CDC *i* to the retailer *k*. Q_{ki} : the product quantity transported from the retailer k to CDC i. $Step \ 1. \ Q_{ik} \leftarrow 0, i \in \Omega, k \in \Phi, Q_{ki} \leftarrow 0, i \in \Omega, k \in \Phi,$ *Step 2.* $t \leftarrow \arg \max p(l), l \in |\Omega| + |\Phi|$; select a node Step 3. If $t \in \Omega$, then $i^* \leftarrow t$; select a CDC, $k^* \leftarrow \arg\min C_{ik}^d \mid p(k) \neq 0, k \in \Phi$; select a retailer with the lowest cost else, k^*l ; select a retailer $i^* \leftarrow \arg\min C_{ik}^d \mid p(i) \neq 0, i \in \Phi$; select a CDC with the lowest cost Step 4. $Q_{ik} \leftarrow \min D_k(1 + \text{EV}[\tilde{\overline{a}}_k]), \alpha_i$: assign the available amount of units Update the availabilities on CDC (i^{\ast}) and retailer (k^{\ast}) $D_{k^*} = D_{k^*} - Q_{i^*k^*}, \alpha_i = \alpha_i - Q_{i^*k^*}$ Step 5. If $D_{k^*} = 0$ then $p_{k^*} = 0$ If $\alpha_{i^*} = 0$ then $p_{i^*} = 0$ Step 6. If p(|i| + k) = 0, $k \in \Phi$, then calculate transportation cost, find the chosen CDC and return, else go to Step 1.

PROCEDURE 1: Decoding of the priority for the location and allocation problem.

Step 4.6 (generate *n*best). For l = 1, 2, ..., L and d = 1, 2, ..., D, find p_{od} ensuring that the FDR takes a maximum value, and set p_{od} as P_l^{Lbest} .

Step 4.7. Update the position and the velocity of each *l*th particle using (20) and (21).

Step 4.8. Check whether the particles are beyond the mark. If $p_{ld} > P^{\max}$, $p_{ld} = P^{\max}$; otherwise, if $p_{ld} < P^{\min}$, then $p_{ld} = P^{\min}$.

Step 5. Based on the above calculation, replace the ranking vector using the new numbers.

Step 6. If the stopping criterion is met, stop; otherwise, $\tau = \tau + 1$ and return to Step 2.

The overall process can be clearly seen in Figure 2.

5. Case Study

5.1. Case Presentation. This model is based on a beer company in a developing country that bottles beer in glass bottles. The supply chain allows customers to return empty bottles to the retailers, which are then sent to the CDCs where they are inspected, consolidated, and sorted. After processing and disinfecting, the bottles are refilled and sold again. The company is now considering the construction of several CDCs to allow for bottle recycling as producing new bottles is far more expensive than recycling used bottles.

To illustrate the validity of the model and the usefulness of the solution method, the data needed to examine the CLSC performance for the four echelons is presented here. Based on the market analysis, ten alternative CDC coordinates are suggested, which are to be assessed based on location, capacity, fixed costs, new product variable costs (NPV cost), and recycled product variable costs (RPV cost). The assessments for the ten possible CDCs are shown in Table 1. Supermarkets and restaurants are considered to be the beer retailers with flexible demand. Table 2 presents the information regarding the retailers, factories, and disposal centers. It can be seen from Table 2 that k_1 to k_{30} represent 30 different retailers, j_1 to j_4 are the 4 different factories in different locations, each of which has a different capacity, and n_1 indicates the location and capacity of the disposal center. Therefore, 30 retailers, 4 factories, and 1 waste disposal center are considered in this study. The unit transportation costs and pollution are related to the distances between the facilities. The retailers' return rates, which are fuzzy random variables, are shown in Table 3.

5.2. Maximum Generation and Population Size. With an increase in the maximum iterations, computer time may increase or the optimal result may improve. For the pb-glnPSO in this case study, the maximum generation is set at T and the population size is set at N; therefore, the maximum iteration is $T \times N$. In the test, N is set from 20 to 60 with a step-length of 10, and T is set from 100 to 500 with a step-length of 100, while $c_p = c_g = c_l = c_n = 2$, from which 25 different groups are obtained, and then pb-glnPSO was run 20 times for each group.

Figures 3 and 4 show the test results, the computing time, and the optimal results. For the horizontal ordinate TN (e.g., "0–5" represents five different groups) [49], when N = 20, T increases from 100 to 500 with a step-length of 100, with the remainder following the same analogy. From Figure 3, it can be seen that the maximum iterations significantly influenced the computing time. When the population size N was the same, there was a positive correlation between the maximum generation and the computing time. For the optimal result, the maximum generation had an obvious influence on the result, as shown in Figure 4. As can be seen, the results

Node	Location	Capability	Fixed cost	NPV cost	RPV cost	$\tilde{\overline{b}}_i$	Parameters $\overline{\rho}$
1	(23, 23)	900	12300	0.01	0.05	$(0.18, \overline{\rho}_1, 0.25)$	$\overline{\rho}_1 \sim N(0.21, 0.02)$
2	(25, 35)	550	12100	0.02	0.06	$(0.23, \overline{\rho}_2, 0.28)$	$\overline{\rho}_2 \sim N(0.25, 0.02)$
3	(34, 29)	1050	15600	0.01	0.05	$(0.14, \overline{\rho}_3, 0.24)$	$\overline{\rho}_3 \sim N(0.18, 0.04)$
4	(32, 25)	650	11300	0.01	0.07	$(0.16, \overline{\rho}_4, 0.22)$	$\overline{\rho}_4 \sim N(0.18, 0.03)$
5	(35, 37)	1050	17800	0.01	0.05	$(0.25, \overline{\rho}_5, 0.30)$	$\overline{\rho}_5 \sim N(0.28, 0.02)$
6	(36, 31)	1050	22400	0.01	0.06	$(0.17, \overline{\rho}_6, 0.26)$	$\overline{\rho}_6 \sim N(0.22, 0.03)$
7	(29, 28)	1050	16300	0.02	0.07	$(0.15, \overline{\rho}_7, 0.23)$	$\overline{\rho}_7 \sim N(0.20, 0.02)$
8	(18, 21)	800	14900	0.01	0.06	$(0.19, \overline{\rho}_8, 0.28)$	$\overline{\rho}_8 \sim N(0.24, 0.03)$
9	(29, 23)	1100	26500	0.01	0.06	$(0.12, \overline{\rho}_9, 0.22)$	$\overline{\rho}_9 \sim N(0.17, 0.04)$
10	(35, 26)	1050	2200	0.02	0.05	$(0.17, \overline{\rho}_{10}, 0.23)$	$\overline{\rho}_{10} \sim N(0.20, 0.02)$

TABLE 1: CDCs information (unit: 1×10^2 RMB).



FIGURE 2: The heuristic algorithms based on pb-glnPSO.

were better when *T* was from 300 to 500 and *N* was from 40 to 60. Thus, for the above *T* and *N* group, the optimal results, the standard deviation of 20 optimal results, and the average computing time have been given in Table 4 for better analysis. From Table 4, it can be seen that the result was stable and reliable and that any further increase in the maximum iterations resulted in a longer computing time but no improvement in the optimal result, with the best result being when N = 50 and T = 400. Therefore, the suitable values for *T* and *N* in the pb-glnPSO for this case were found to be 400 and 50.

5.3. Sensitivity Analysis on the Parameters. To find the best solution to the proposed model, a series of experiments were conducted, all of which were performed using MATLAB 7.0

on a workstation with an Intel(R) Core(TM)i7, a 2.59 GHz clock pulse with 8 GB memory, and a Windows 10 operating system. A sensitivity analysis was performed to examine the effectiveness and behavior of the proposed algorithm, as shown in Table 5. Several parameters were changed: the population size N, the maximum generation T, and the acceleration constants c_p , c_g , c_l , and c_n . After trying various values for the population size and maximum generations, the results were found to be better when T was from 300 to 500 and N was from 40 to 60. The different fitness values obtained using the pb-glnPSO with the different parameters N, T, c_1 , and c_2 are shown in Table 5.

As can be seen from Table 5, when the parameters c_p , c_g , c_l , and c_n increased, the fitness value improved except for $c_p = c_q = c_l = c_n = 2.5$ with the same generation

TABLE 2: Retailers, factories, and disposal center.

Node	Location	Demand
k_1	(27, 28)	50
k_2	(30, 19)	60
k_3	(32, 22)	40
k_4	(37, 16)	80
k_5	(23, 29)	30
k_6	(27, 17)	40
j_1	(13, 22)	920
k_7	(33, 26)	80
k_8	(34, 32)	40
k_9	(37, 22)	100
k_{10}	(17, 22)	90
k_{11}	(26, 39)	60
k_{12}	(38, 26)	40
j_2	(31, 44)	530
<i>k</i> ₁₃	(38, 34)	50
k_{14}	(36, 25)	70
k_{15}	(41, 19)	40
k_{16}	(27, 33)	30
k_{17}	(25, 39)	20
k_{18}	(38, 37)	40
j_3	(32, 15)	850
k ₁₉	(36,27)	50
k_{20}	(39, 28)	60
k_{21}	(25, 31)	70
k ₂₂	(29, 35)	90
k ₂₃	(18, 29)	50
k_{24}	(18, 14)	60
j_4	(42, 31)	940
k_{25}	(35, 11)	80
k_{26}	(23, 33)	50
k ₂₇	(36, 37)	60
k_{28}	(28, 26)	40
k ₂₉	(25, 24)	30
k_{30}	(32, 19)	80
<i>n</i> ₁	(18, 47)	800

and population size, with the fitness value increased from $c_p = c_g = c_l = c_n = 2$ to $c_p = c_g = c_l = c_n = 2.5$. Therefore, when $c_p = c_g = c_l = c_n = 2$, the result was found to be optimal. For *T*, given the same c_p, c_g, c_l , and c_n and population size, it was found that when *T* was 400, the fitness value was better than for any other generation. Finally, for *N*, the results improved as the population size increased and it was found to be optimal when *N* is 50. The most effective and efficient results were gained with *T* at 400, *N* at 50, and $c_p = c_q = c_l = c_n = 2$.

5.4. Result Analysis. In this section, the pb-glnPSO is conducted to solve the model using the above data. The parameters for the problem were set as follows: population size: popsize = 50; maximum generation: maxGen = 400; inertia weight: $\omega(1) = 1$ and $\omega(T) = 0.1$; acceleration constant:

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TABLE 3: Retailers' return rate.

Node	$\tilde{\overline{a}}_k$	Parameters $\overline{\varsigma}$
k_1	$(0.28, \overline{\varsigma}_1, 0.33)$	$\bar{\varsigma}_1 \sim N(0.31, 0.02)$
k_2	$(0.52, \overline{\varsigma}_2, 0.62)$	$\bar{\varsigma}_2 \sim N(0.56, 0.02)$
k_3	$(0.32, \overline{\varsigma}_3, 0.38)$	$\bar{\varsigma}_3 \sim N(0.36, 0.03)$
k_4	$(0.62, \overline{\varsigma}_4, 0.73)$	$\bar{\varsigma}_4 \sim N(0.67, 0.04)$
k_5	$(0.55, \overline{\varsigma}_5, 0.62)$	$\bar{\varsigma}_5 \sim N(0.58, 0.02)$
k_6	$(0.65, \overline{\varsigma}_6, 0.72)$	$\overline{\varsigma}_6 \sim N(0.69, 0.02)$
k_7	$(0.73, \overline{\varsigma}_7, 0.81)$	$\overline{\varsigma}_7 \sim N(0.78, 0.03)$
k_8	$(0.72, \bar{\varsigma}_8, 0.78)$	$\overline{\varsigma}_8 \sim N(0.75, 0.03)$
k_9	$(0.75, \overline{\zeta}_9, 0.82)$	$\overline{\varsigma}_9 \sim N(0.80, 0.04)$
k_{10}	$(0.34, \overline{\varsigma}_{10}, 0.38)$	$\overline{\varsigma}_{10} \sim N(0.36, 0.02)$
k_{11}	$(0.42, \overline{\varsigma}_{11}, 0.48)$	$\overline{\varsigma}_{11} \sim N(0.46, 0.02)$
k_{12}	$(0.46, \overline{\varsigma}_{12}, 0.50)$	$\overline{\varsigma}_{12} \sim N(0.48, 0.04)$
<i>k</i> ₁₃	$(0.65, \overline{\varsigma}_{13}, 0.70)$	$\overline{\varsigma}_{13} \sim N(0.67, 0.04)$
k_{14}	$(0.62, \overline{\varsigma}_{14}, 0.68)$	$\overline{\varsigma}_{14} \sim N(0.64, 0.03)$
k_{15}	$(0.72, \overline{\varsigma}_{15}, 0.78)$	$\overline{\varsigma}_{15} \sim N(0.75, 0.04)$
k_{16}	$(0.74, \overline{\varsigma}_{16}, 0.79)$	$\overline{\varsigma}_{16} \sim N(0.77, 0.04)$
k_{17}	$(0.70, \overline{\varsigma}_{17}, 0.75)$	$\overline{\varsigma}_{17} \sim N(0.73, 0.04)$
k_{18}	$(0.75, \overline{\varsigma}_{18}, 0.80)$	$\overline{\varsigma}_{18} \sim N(0.78, 0.02)$
<i>k</i> ₁₉	$(0.68, \overline{\varsigma}_{19}, 0.75)$	$\bar{\varsigma}_{19} \sim N(0.72, 0.03)$
k_{20}	$(0.55, \overline{\varsigma}_{20}, 0.65)$	$\bar{\varsigma}_{20} \sim N(0.61, 0.02)$
k_{21}	$(0.35, \overline{\varsigma}_{21}, 0.45)$	$\overline{\varsigma}_{21} \sim N(0.39, 0.04)$
k ₂₂	$(0.35, \overline{\varsigma}_{22}, 0.42)$	$\overline{\varsigma}_{22} \sim N(0.38, 0.03)$
k_{23}	$(0.55, \overline{\varsigma}_{23}, 0.60)$	$\overline{\varsigma}_{23} \sim N(0.58, 0.03)$
k_{24}	$(0.52, \overline{\varsigma}_{24}, 0.62)$	$\overline{\varsigma}_{24} \sim N(0.56, 0.04)$
k_{25}	$(0.62, \overline{\varsigma}_{25}, 0.72)$	$\overline{\varsigma}_{25} \sim N(0.66, 0.04)$
k_{26}	$(0.69, \overline{\varsigma}_{26}, 0.75)$	$\overline{\varsigma}_{26} \sim N(0.72, 0.03)$
k_{27}	$(0.29, \overline{\varsigma}_{27}, 0.35)$	$\overline{\varsigma}_{27} \sim N(0.32, 0.03)$
k ₂₈	$(0.40, \overline{\varsigma}_{28}, 0.46)$	$\overline{\varsigma}_{28} \sim N(0.43, 0.02)$
k ₂₉	$(0.42, \bar{\varsigma}_{29}, 0.48)$	$\bar{\varsigma}_{29} \sim N(0.45, 0.03)$
k ₃₀	$(0.50, \overline{\varsigma}_{30}, 0.58)$	$\overline{\varsigma}_{30} \sim N(0.54, 0.04)$

 $c_p = c_g = c_l = c_n = 2$. After running the program 20 times, the most satisfactory solution was found. Figure 6

	N = 40		N = 50			N = 60			
	T = 300	T=400	T=500	T = 300	T=400	T=500	T = 300	T=400	T=500
Optimal result	82589.8	82434.4	82669.5	82102.4	81920.1	82634.9	82434.4	82189.2	82308.9
Standard deviation	512.485	627.973	407.732	565.617	306.345	527.466	543.704	526.525	678.494
Computing time	5.600	5.938	6.699	6.623	8.384	9.228	8.929	9.855	10.690

TABLE 4: pb-glnPSO result (unit: 1×10^2 RMB).

TABLE 5: Sensitivity analysis (unit: 1×10^2 RMB).	
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$c_1 = c_2$	N = 40				N = 50			N = 60		
	T = 300	T = 400	T = 500	T = 300	T = 400	T = 500	T = 300	T = 400	T = 500	
0.5	84449.3	83447.3	83243.5	82889.4	82766.2	83859.6	82879.9	82830.6	83233.8	
1	83256.4	83125.5	83128.3	82679.3	82459.0	83199.5	82596.8	82448.8	82614.4	
1.5	83201.8	82765.3	82815.6	82486.5	82564.3	83041.9	82498.6	82257.8	82430.6	
2	82589.8	82434.4	82669.5	82102.4	81920.1	82634.9	82434.4	82189.2	82308.9	
2.5	83844.7	83102.9	82782.4	82370.5	82198.5	82792.7	82448.8	82469.9	82377.3	





FIGURE 5: The distribution strategy.

shows the specific objective values found by the pb-glnPSO in different iterations and shows the reductions in the total costs. The results are presented in Figure 5. From the calculations, at least 3 CDCs could satisfy all markets. Alternative CDC positions 4, 5, and 8 should be chosen. CDC 4 can send products to markets 2, 3, 4, 6, 7, 9, 14, 15, 25, and 30. CDC 5 can transport products for 1, 8, 11, 12, 13, 16, 17, 18, 19, 20, 21, 22, and 27 and retailers 5, 10, 23, 24, 26, 28, and 29 can be serviced by CDC 8. The total cost was found to be 8.192 million RMB, in which the fixed costs were 4.4 million RMB, the transportation costs were 3.776 million RMB, and the operating costs were 0.016 million RMB.

5.5. Algorithm Comparison. To better illustrate the effectiveness of the proposed algorithm, a brief comparison between the pb-glnPSO, glnPSO, and an immune algorithm (IM) is given in this section. The glnPSO is a well-respected evolutionary algorithm that has been successfully implemented in a variety of engineering and combinatorial problems. The IM is also being widely used to solve facilities location problems.

To establish the solution quality for the pb-glnPSO, it was compared with the glnPSO and the IM. Each run time for the pb-glnPSO, glnPSO, and IM were run 20 times using the same data. For a fair comparison between the groups, the population size was set at 50 and the maximum generation at 400. In the glnPSO, the acceleration constant was designed as $c_p = c_g = c_l = c_n = 2$ and the inertia weight was $\omega(1) = 1$ and $\omega(T) = 0.1$. For the IM algorithm, the crossover probability was 1 and the mutation probability was 0.1.

From Figure 7, it can be seen that the pb-glnPSO outperformed both the glnPSO and the IM, and, as the glnPSO converged faster, it had a better result than the IM. This demonstrated that a better solution can be obtained using the glnPSO, and an even better result can be obtained using the



FIGURE 6: The Pgln-PSO iterative process.



FIGURE 7: The iterative process of pb-glnPSO, glnPSO, and IM.

pb-glnPSO. The blue profile shows the convergence for the best in history for the pb-glnPSO. It can be seen from Figure 7 that, as the programs ran, the results become stable for the pbglnPSO and glnPSO after about the 150th generation, while the IM became stable after the 175th generation. As shown in Figure 7, the best solution for the pb-glnPSO was superior to, more stable than, and had the smallest CPU run time compared to the other algorithms (Table 6), with the IM having the highest run time.

6. Conclusion

Economic development has resulted in many environmental pollution problems, the seriousness of which has encouraged

TABLE 6: Results of the pb-glnPSO, glnPSO, and IM.

Item	pb-glnPSO	glnPSO	IM
Best result	81920.1	82884.5	83315.6
Worst result	82434.9	83826.6	84585.5
Average result	82231.1	83467.6	83852.8
Difference between the best and the worst	514.81	942.10	1269.90
Difference between the average and the best	510.986	992.734	864.315
Standard deviation	203.82	583.08	537.27
CPU time	8.6796	10.9764	14.5000

people to recycle and reuse products. To examine this problem and seek appropriate solutions, a mathematical model for a collection-distribution center location and allocation problem in a closed-loop supply chain under a fuzzy random environment was presented for the beer industry in China. For this problem, a new model was formulated, in which the decision makers sought to minimize costs and pollution under capacity and quantity constraints. To more accurately represent actual production situations, the return rate and disposal rate were considered to be fuzzy random variables. A heuristic algorithm, the pb-glnPSO, was then applied to solve the problem. Based on the proposed priority, the distribution and collection activity were shown to satisfy retailer demand and reduce costs and pollution after the CDCs start operations. After calculation, the best solution was determined and the advantages of the algorithm were illustrated. The proposed model and method can be applied to the location and allocation of CDCs in the beer industry to improve supply chain management. The model was shown to assist in generating retailer demand and dealing with the returned products, which could benefit company recycling and reuse policies. At the same time, the transportation costs and pollution were reduced because of the reductions in losses from empty loads.

Notations

Sets

- Ω: Set of CDCs, $Ω = \{1, 2, 3, ..., I\}$
- Ψ : Set of factories, $\Psi = \{1, 2, 3, \dots, J\}$
- $\Phi: \text{ Set of retailers, } \Phi = \{1, 2, 3, \dots, K\}$
- Y: Set of disposal centers, $\Upsilon = \{1, 2, 3, ..., N\}$

Indices and Parameters

- *i*: Alternative location position for the CDCs, $i \in \Omega = \{1, 2, 3, ..., I\}$
- *j*: Known position of the factories, $j \in \Psi = \{1, 2, 3, \dots, J\}$
- *k*: Known position of the retailers, $k \in \Phi = \{1, 2, 3, \dots, K\}$
- *n*: Known disposal center, $n \in \Upsilon = \{1, 2, 3, \dots, N\}$

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- U: The upper limit of the CDCs
- The demand of retailer k D_k :
- α_i : The capability of CDC *i*
- The capability of factory *j* γ_i :
- P_{ii} : Product quantity from factory *j* to CDC *i*

- $\begin{array}{l} Q_{ik}: & \text{Product quantity from CDC } i \text{ to retailer } k\\ \overline{\tilde{a}}_{k}: & \text{Product quantity from CDC } i \text{ to retailer } k\\ \overline{\tilde{b}}_{i}: & \text{The product return rate from retailer } k\\ \overline{\tilde{b}}_{i}: & \text{The product disposal rate at CDC } i\\ F_{i}^{c}: & \text{The fixed costs of the CDC } i\\ V_{i}^{c}: & \text{The variable costs of the CDC } i \text{ for a new} \end{array}$ product unit
- RV_i^c : The variable cost of the CDC *i* triage for a returned product unit
- C_{ii}^p : Unit transportation cost between CDC i and factory *j*
- C_{ik}^{d} : Unit transportation cost between CDC *i* and retailer k
- C_{in}^{ω} : Unit transportation cost between CDC *i* and disposal center *n*
- β_{ii} : Environmental impact of transportation between CDC *i* and factory *j*
- β_{ik} : Environmental impact of transportation between CDC *i* and retailer *k*
- β_{in} : Environmental impact of transportation between CDC *i* and disposal center *n*
- The environmental impact level accepted by μ: decision makers.

Decision Variables

- x_i : A binary variable indicating whether point *i* is chosen. If point *i* is chosen, then $x_i = 1$; else, $x_i = 0$
- y_{ik} : It indicates whether retailer k is served by CDC *i*. If *i* is chosen, then $y_{ik} = 1$; else, $y_{ik} = 0$

Notations for pb-glnPSO

τ:	Iteration index, $\tau = 1, 2, \dots, T$
<i>l</i> :	Particle index, $l = 1, 2, \ldots, L$
$v_{ld}(\tau)$:	Velocity of the <i>l</i> th particle at the <i>d</i> th
	dimension in the $ au$ th iteration
p_{ld}^{best} :	Personal best position
p_{ld}^{Lbest} :	Local best position
c_p :	Personal best position acceleration
Γ	constant
<i>cl</i> :	Local best position acceleration constant
P^{\max} :	Maximum position value
P_l :	Velocity vector of <i>l</i> th particle
P_l^{best} :	Vector personal best position of <i>l</i> th
-	particle
P_l^{Lbest} :	Vector local best position of <i>l</i> th particle
Fitness(P_l):	Fitness value of P_l
d:	Dimension index, $d = 1, 2,, D$
ω_{τ} :	Inertia weight in τ th iteration
$p_d^l(\tau)$:	Position of the <i>l</i> th particle at the <i>d</i> th
	dimension in the τ th iteration

p_{ad}^{best} :	Global best position
p_{ld}^{Nbest} :	Near-neighbor best position
c_q :	Global best position acceleration constant
c_n :	Near-neighbor best position acceleration
	constant
P^{\min} :	Minimum position value
V_l :	Position vector of <i>l</i> th particle
P_q^{best} :	Vector global personal best position
r_1, r_2, r_3, r_4 :	Uniform distributed random number within
	[0, 1].

Conflicts of Interest

The authors declare that they have no conflicts of interest.

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