

# A Detailed Study of the Generation of Optically Detectable Watermarks using the Logistic Map

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## Abstract

A digital watermark is a visible, or preferably invisible, identification code that is permanently embedded in digital media, to prove owner authentication and provide protection for documents. Given the interest in watermark generation using chaotic functions a detailed study of one chaotic function for this purpose is performed. In this paper, we present an approach for the generation of watermarks using the logistic map. Using this function, in conjunction with seed management, it is possible to generate chaotic sequences that may be used to create highpass or lowpass digital watermarks. In this paper we provide a detailed study on the generation of optically detectable watermarks and we provide some guidelines on successful chaotic watermark generation using the logistic map, and show using a recently published scheme, how care must be taken in the selection of the function seed.

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## 1 Introduction

Digital Watermarking provides a means of inserting additional information into digital documents [1]. This extra information may be used to provide a certain degree of protection to the document from malicious attack. This information is usually a piece of identifying information which may be used at a later date to prove ownership. Watermarking in contrast to steganography

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has the additional notion of resilience against attempts to remove the hidden data [2]. The original document owner can prove ownership of copies by proving that it contains their original watermark.

Numerous watermark generation techniques have been proposed to date, ranging from the use of personal logos to the use of pseudorandom sequences of numbers [1]. The use of chaotic functions for the generation of watermarks has also been proposed by Pitas *et al.* [3–5] where the Markov Maps and Bernoulli Maps are used as the chaotic functions. These types of watermark generation schemes require two values, the initial value and the function seed, in order to recreate the same watermark at a later stage. An advantage of these watermarks is the ability to analyse and control their spectral properties. Typically these watermarks are distributed over the full cover image using an additive or multiplicative technique.

Pitas *et al.* have described in detail the use of chaotic functions in watermark generation [3,5] and provide analysis of the functions and techniques used. A similar analysis for the logistic map does not exist and we believe that an understanding of it would be significant in the generation of chaotically generated watermarks. In this paper we deal with optical watermarks, which we consider to be watermarks which may be detected with optical correlation techniques. Previous watermarking schemes have proposed the use of optical devices such as photocopiers to retrieve the embedded watermarks [6,7] but no scheme has proposed the use of optical hardware to perform the task of watermark detection.

## 2 Relevant properties of the logistic map for the generation of optically detectable watermarks

In this paper, we utilise watermarks that are generated using a deterministic chaotic map. One motivation for using such chaotic functions to generate a watermark is that a single variable, seeding the chaotic function, will always result in the same output (mapping) when certain constraints or initial conditions are placed on the mapping. The primary advantage, however, is that it is possible to investigate the spectral properties of the resulting watermark [3–5]. The function used in this paper for the generation of sequences which are transformed into watermarks is the logistic difference equation, sometimes referred to as the logistic equation or mapping:

$$y_{n+1} = ay_n(1 - y_n) \tag{1}$$

where  $a$  is the ‘function seed’ and  $y_n$  is the current value of the mapping in time with an initial value  $y_0$  [8]. The properties of the logistic map have been

thoroughly studied and are well classified. The complicated but deterministic properties of the map make it ideally suited for watermark generation. The value  $y_{n+1}$  is dependant on its current value  $y_n$ . For low values of  $a$ ,  $y_n$  eventually converges to a single number as  $n \rightarrow \infty$ . When  $a = 3.0$ ,  $y_n$  no longer converges – it oscillates between two values. This characteristic change in behaviour is called a bifurcation, and can be seen in the bifurcation diagram shown in Fig. 1. Increasing the value of  $a$  ( $> 3.0$ ) causes  $y_n$  to oscillate between not two, but four values. As  $a$  increases,  $y_n$  goes through bifurcations of period  $2^n$  and eventually becomes chaotic. When the value of  $a \approx 3.55699$  (the accumulation point),  $y_n$  neither converges or oscillates – its value becomes random. For values of  $a > 3.5699$ , the behaviour is largely chaotic [8].

The degree of randomness of any map  $F(x)$  can be seen from the variation of the Lyapunov Exponent for the map with respect to the value of  $a$ . We can calculate the Lyapunov Exponent,  $\lambda$ , for the logistic equation from [9]:

$$\lambda = \lim_{N \rightarrow \infty} \frac{1}{N} \log \left| \frac{d}{dx} F^n(y) \right|_{y_0} \quad (2)$$

which, on applying the chain rule

$$\left. \frac{dF^n}{dy} \right|_{y_0} = \prod_{i=0}^{N-1} \left. \frac{dF}{dy} \right|_{y_i} \quad (3)$$

reduces to

$$\lambda(y_0) = \lim_{N \rightarrow \infty} \frac{1}{N} \sum_{i=0}^{N-1} \log |F'(y_i)| \quad (4)$$

where  $y_i$  are points on the orbit. The map is dissipative for almost all  $y_0$  and gives rise to the same  $\lambda$  for a fixed value of  $a$ . For one-dimensional motion,  $\lambda$  is equal to the Kolmogorov-Sinai entropy (KSE). The KSE marks the transition from regular to chaotic behaviour with the notion that a positive KSE is a sufficient condition for chaos [10].

The Lyapunov Exponent is useful for distinguishing between the various types of orbits. When  $\lambda \leq 0$  the logistic map exhibits periodic behaviour, and when  $\lambda > 0$  the logistic map exhibits chaotic behaviour [11,12]. Fig. 1 illustrates the Lyapunov Exponent for the range  $[0, 4]$  for the logistic equation. Whenever an orbit contains  $y_n = \frac{1}{2}$ ,  $\lambda = -\infty$ , and every  $2^n$  periodic orbit has this as one of its periodic points.

For the logistic map there are two distinct regimes: (1): the periodic or bifurcation regime. Here  $0 < a < a_\infty$  and  $\lambda$  is negative, and, (2): the chaotic

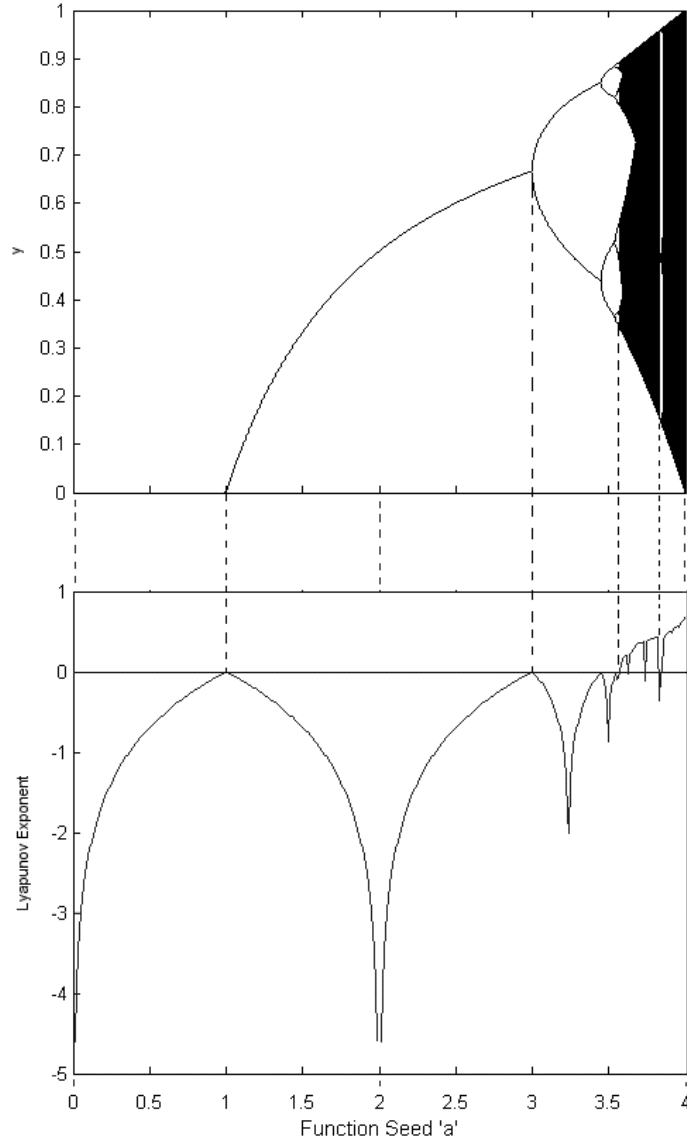


Fig. 1. The bifurcation diagram for the logistic map and the corresponding Lyapunov Exponent.

regime, for which  $a_\infty < a \leq 4$  and  $\lambda$  is mostly positive. The periodic regime is characterised by  $a_n$  where the number of fixed points changes from  $2^{n-1}$  to  $2^n$  and the  $a_n$  scale according to Feigenbaum's  $\delta$  and  $\alpha$ , where

$$\delta = \lim_{n \rightarrow \infty} \frac{a_n - a_{n-1}}{a_{n+1} - a_n} = 4.669201 \quad (5)$$

and

$$\alpha = \lim_{n \rightarrow \infty} \frac{d_n}{d_{n+1}} = 2.502907 \quad (6)$$

where  $d_n$  are the distances of points in a  $2^n$  super cycle which is closest to the superstable point  $y = \frac{1}{2}$ , which is contained in every orbit.

In the chaotic regime, the chaotic intervals move together by inverse bifurcation as  $a$  is increased from  $a_\infty$  to 4 until the iterates become distributed over the whole interval  $[0, 1]$  at  $a = 4$ . Within the chaotic intervals are a-windows which are characterised by p-cycles ( $P = 3, 5, 6, \dots$ ) with successive bifurcations  $P, P.2^1, P.2^2, \dots$ . The corresponding  $a$  values scale as in Eq.5. There are also period tripling  $P.3^n$  and quadruplings  $P.4^n$ , etc. which occur at  $\bar{a}_n$  and give rise to constants  $\bar{\delta}$  and  $\bar{\alpha}$  which are again universal. For example,  $\bar{\delta} = 55.247$  and  $\bar{\alpha} = 9.277$  for  $P.3^n$  [13,14].

The logistic map maps the closed interval  $[0, 1]$  onto itself and examination of one such transformation (when  $a = 4$ ) illustrates this, as shown in Fig. 2 and may also be observed by examining Fig. 1 at  $a = 4$ . By choosing an initial value  $y_0 \in [0, 1]$  the successive states of the system are given by the trajectory [15]:

$$y_0, F(y_0), F^2(y_0) = F(F(y_0)), \dots \quad (7)$$

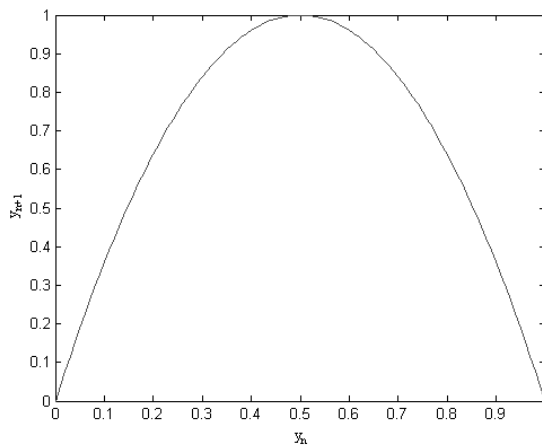


Fig. 2. Cross section  $y_{n+1}$  versus  $y_n$  when  $a = 4$ .

Sections 3 – 6 describe watermark generation, embedding and detection schemes along with unique properties of the logistic map which are useful in the generation of chaotic watermarks, in particular the choice of function seed.

### 3 Watermark Generation

When the logistic equation, shown in Eq.1 is seeded with a value  $3.5699 \leq a \leq 4.0$ , and iterated, chaotic behaviour is witnessed in general. It is this feature of

the logistic equation that is utilised to generate watermarks. The trajectories produced will differ greatly even for small differences in the seed value. The dense set of chaotic periodic orbits of the mapping are used for the generation of watermarks.

A chaotic binary sequence is quantified from the real sequence generated from iterating the logistic map using:

$$x_n = \begin{cases} 0 & y_n \leq 0.5 \\ 1 & y_n > 0.5 \end{cases} \quad (8)$$

This sequence is then converted into a 2D image which may be used as a watermark. Usually, watermarks are generated which are the same size as the image which will contain the watermark, as this ensures the watermark is spread across the entire image. A scanning technique, known as Peano Scanning [16], may be used to determine watermark pixel ordering in the 2D image and is preferable to the conventional Raster Scanning technique in that scanning is not predictable, and can produce many variations within the same image. A drawback with using Peano Scanning, however, is that it will only operate on images whose height and width are multiples of two. The scanning order always moves to a neighbouring pixel, but the pattern may appear in different orientations depending on the starting pixel within the scanning routine [17,18]. An example of a watermark generated using this technique is shown in Fig. 3(a). In this case, the logistic function was seeded with a value  $a = 4$  and an initial starting value of  $y_0 = 0.001$ .

Watermarks produced using chaotic functions are desirable, with only the initial starting value and seed needed to regenerate a sequence of numbers. This ensures that the sequence can only be reproduced if seeded with the corresponding initial starting value and also the function seed. Given that the logistic map is not invertible, there is added security over pseudorandom sequences. Chaotic watermarks are preferable to the more conventional pseudorandomly generated watermarks, in that their spectral properties can be controlled, which is discussed below.

An important property of the logistic map is that there are certain regions of the logistic map where chaotic behaviour is not observed. It should be obvious that no chaotic behaviour is observed for values of the seed less than the accumulation point, see Fig. 1. Beyond this point periodicity witnessed previously gives way to chaos. Within this chaotic region there are non chaotic regions present. These occur at the bifurcation zones of the chaotic region of the map, for example, sizeable regions exist when  $3.627 \leq a \leq 3.634$  (6 bifurcations  $-P.6^n$ ),  $3.739 \leq a \leq 3.744$  (5 bifurcations  $-P.5^n$ ) and  $3.829 \leq a \leq 3.856$  (3 bifurcations  $-P.3^n$ ).

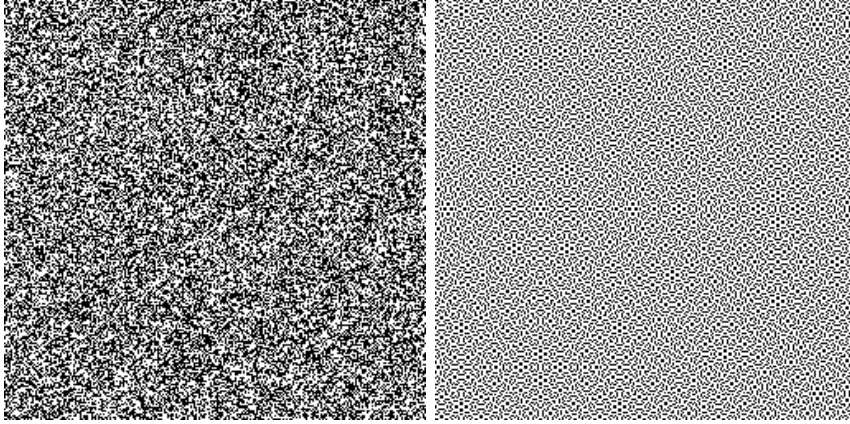


Fig. 3. Watermark generated by iterating the Logistic Map with an initial value of 0.001 and (a):  $a = 4$  (b):  $a = 3.83$ .

The latter of these regions, known as a period-tripling region, within this region the logistic map can be in one of three states independent of  $y_0$ . Using a seed for the function within this zone will result in the orbit of the map taking on only one of three values. The sequences produced are non chaotic, and also predictable and therefore are not suitable for watermark generation. A sample of a watermark generated when  $a = 3.83$  is shown in Fig. 3(b). In this case Peano Scanning is used to determine the scanning order of the pixels. This figure illustrate that non-chaotic predictable sequences are produced even in the chaotic region. It can be seen that there is a pattern developing within the watermark and this is due to the predictability of the map at this location. Similar results are found for each of the non-chaotic regions of the map, indicated previously.

#### 4 Probability Density

It is often useful to examine the frequency with which each state along a trajectory falls into a certain region of the state space. This is known as the probability density, or the invariant measure, and is calculated as follows: Assume that state space  $[0, 1]$  is divided into  $n$  discrete non-intersecting intervals, so the  $i^{th}$  interval is (neglecting the end point 1):

$$[(i - 1)/n, i/n) \quad i = 1, \dots, n \quad (9)$$

Next, an initial state  $y_0$  is chosen and a trajectory of length  $N$  ( $N \gg n$ ) calculated by:

$$y_0, F(y_0), F^2(y_0) = F(F(y_0)), \dots, F^N(y_0) \quad (10)$$

The fraction,  $f_i$ , of the  $N$  system states that are in the  $i^{th}$  interval is given by [15]

$$f_i = \frac{n}{N} \{\text{number of } F^j(y_0) \in [(i-1)/n, i/n), j = 1, \dots, N\} \quad (11)$$

Fig. 4(a) and (b) shows the calculated probability density when  $a = 4$  and when  $a = 3.83$  for 5000 iterations, and 200 possible states. It can be seen that in the case where  $a = 4$  there is symmetry with most states concentrated near 0, and near 1 with a minimum at  $\frac{1}{2}$ . In the case of  $a = 3.83$  it can be seen that the system can be in only one of three possible states. This is not surprising since when  $a = 3.83$  the logistic mapping is within a stable state, the period tripling state, referred to earlier.

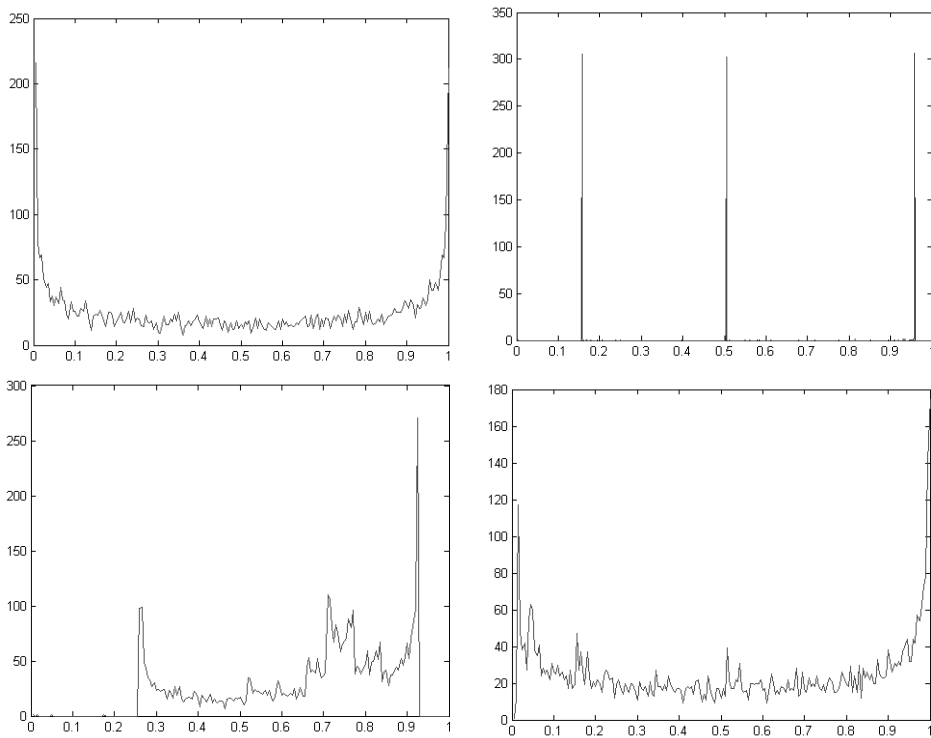


Fig. 4. Probability Density when (a):  $a = 4$  (top left), (b):  $a = 3.83$  (top right), (c):  $a = 3.7$  (bottom left), and (d):  $a = 3.99$  (bottom right).

Fig. 4(c) and (d) shows the cases where the above procedure was repeated with  $a = 3.7$  and with  $a = 3.99$ . It can be observed that in the case of  $a = 3.99$  there is a greater spread of states. There is a balance between the number of 0 pixels (white) and 1 pixels (black). In the case of  $a = 3.7$  it can be seen that the trajectory ends up in states greater than 0.5 substantially more times than in states less than 0.5. This means that in the watermark there will be more 1 pixels than 0 pixels, i.e. more white than black pixels. In the case of  $a = 3.99$  there is much less density variation in watermark produced than in the case of the watermark generated when  $a = 3.7$ . This means that the watermark



produced when  $a = 3.99$  is more lowpass than the watermark produced when  $a = 3.7$  and has shown to be the case [19].

There are, however, other stable states within this region which are not suitable for use in watermark generation. These stable states occur when a trajectory arrives at a fixed point of the logistic map, i.e. when  $y_{n+1} = y_n$ . Therefore, the fixed points may be obtained by solving

$$\begin{aligned} y[1 - a(1 - y)] &= y(1 - a + ay) \\ &= ay[y - (1 - a^{-1})] = 0 \end{aligned} \tag{12}$$

Thus

$$y_0 = 0 \quad \text{and} \quad y_0 = 1 - a^{-1} \tag{13}$$

For the case of  $a = 4$  the two fixed points which satisfy Eq. 13 are  $y_0 = 0$  and  $y_0 = \frac{3}{4}$  [15].

When the mapping is seeded with a value leading to a fixed point, the map is in a periodic state. Watermarks produced in this case are not suitable for watermark embedding as there is no variation in the pixels within the watermark (either all black or all white pixels). This would result in no variation in the image to be watermarked after watermark embedding and thus render detection impossible.

## 5 Watermark Embedding and Detection

Watermark embedding refers to the process of inserting the generated watermark robustly into some cover data, for example images, video or music. The concept of robustness refers to the ability to detect the watermark after common signal processing operations [20]. The generated watermark is usually embedded in the spatial domain or in some transform domain (DCT, DFT, DWT) in a multiplicative manner

$$m_n = j_n + \gamma j_n x_n \tag{14}$$

where  $m_n$  is the  $n^{th}$  pixel of the watermarked image  $M$ ,  $j_n$  is the  $n^{th}$  pixel of the cover image  $J$ ,  $x_n$  is the  $n^{th}$  pixel of the watermark  $X$ . The constant  $\gamma$  is known as the embedding factor and controls the watermark strength. The multiplicative watermark embedding rule relies on ‘‘Weber’s Law’’, which

indicates that the change in a stimulus that will be just noticeable is a constant ratio of the original stimulus [21], and may be expressed as

$$\frac{\Delta I}{I} = k \quad (15)$$

where  $I$  represents the initial intensity of a pixel,  $\Delta I$  the change in the intensity of the pixel.  $k$  signifies that the proportion on the left hand side remains constant despite the value of  $I$  changing. In Eq. 14,  $\gamma$  controls the trade-off between watermark visibility and watermark robustness within the image. The lower the value of  $\gamma$  the less noticeable a watermark is within an image and therefore does not noticeably alter the perceived quality of the image [20].

Watermark Detection is the process of determining whether an image is watermarked with a certain watermark. In some watermarking applications, the original, unwatermarked image is available during detection, this process is referred to as informed detection. However, detection must generally be performed without access to the original image, known as blind detection. The majority of watermark detection techniques use a form of correlation detection. Correlation has been shown to be an effective and relatively simple technique for comparison of signals, and thus is very popular in watermark detection systems. An example of a correlation detector given by

$$\rho = \frac{1}{3NM} \sum_{i=1}^N \sum_{j=1}^M I'(i, j)w(i, j) \quad (16)$$

where  $N$  is the image width,  $M$  is the image height,  $I'$  is the possibly watermarked image and  $w$  is the watermark sequence [22]. This correlation value is compared to a suitably selected threshold  $T$  and if the value of the correlation is greater than the threshold positive detection of the watermark is assumed.

The motivation for the use of an optical detection technique is that optical processing is faster than a similar digital technique due to the inherent parallelism of the system [23]. One technique to perform optical correlation is Matched Filtering [24,25] which has been shown to be an effective technique for watermark detection. For the optical detection system, the possibly watermarked image together with the watermark whose presence (or absence) in the image one wishes to determine are used as inputs to the detector. The watermark is used as the filter for the matched filter detector. The optical detection technique returns an output image which is the correlation between the possibly watermarked image and the watermark. The presence of the watermark in the image is represented by a sharp correlation peak in the centre of the output plane. This correlation peak is compared to a threshold to determine the presence or absence of the watermark. For more detail on the optical detection scheme employed one may consult earlier work by the authors [26,27,19].

## 6 Power Spectrum

Pitas *et al.* have used the power spectrum or power spectral density (PSD) to determine if a particular generated watermark sequence is lowpass or highpass in characteristics [3]. The power spectrum of a function is defined as the modulus-squared of the discrete Fourier Transform of some finite sampled section of it and is given by

$$P(u, v) = |FT[I(u, v)]|^2 \quad (17)$$

where  $P(u, v)$  is the power of a pixel  $(u, v)$ ,  $FT$  is the Fourier Transform and  $I(x, y)$  is the image pixel whose power spectrum one is calculating. The power spectrum of the logistic map is

$$P(y_n) = |FT[y_n]|^2 \quad (18)$$

Examination of the power spectrum of the logistic map in the chaotic region, has shown that the general trend is a decrease in the power spectrum with increasing seed values. This suggests that more lowpass watermarks are generated with higher values of the seed  $a$ , for example, the watermark generated when  $a = 3.7$  is more highpass in nature than a watermark generated when  $a = 3.9$ . [19], as can be seen in Fig. 5.

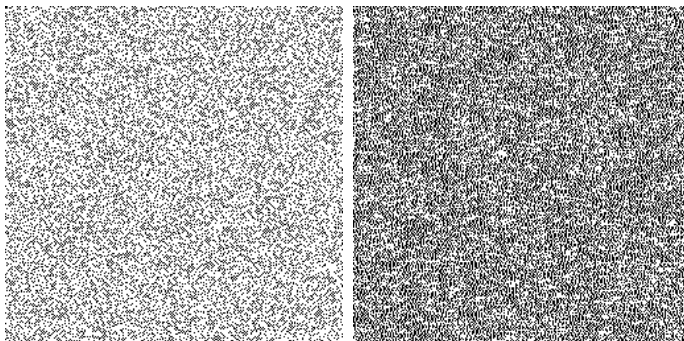


Fig. 5. Watermark produced when (a):  $a = 3.7$  and (b):  $a = 3.9$ .

The power spectrum of other chaotic maps, for example the Markov map, may be used solely to determine the spectral properties of the map. However, the PSD can not be used in isolation for the logistic map for certain values of  $a$ . One might expect that watermarks generated on either side of a periodic region within  $3.57 \leq a \leq 4$  would have similar spectral properties. However, by examining the PSD's and also the generated watermarks it can be seen that this is not the case. Fig. 6 shows watermarks generated by  $a = 3.828$  and  $a = 3.859$  (either side of the period tripling region). It can be seen that there is a difference in the density variations in the watermark pixels and little similarity as their PSD's suggest. Therefore, unlike other well behaved chaotic

maps the PSD alone may not be used to determine the lowpass and highpass properties of watermarks generated from the logistic map.

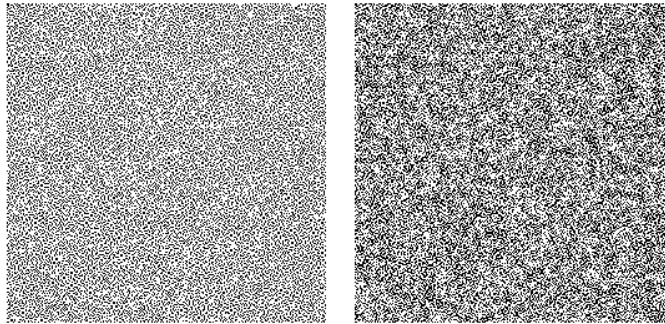


Fig. 6. Watermark produced when (a):  $a = 3.828$  and (b):  $a = 3.859$ .

The watermarks generated using the logistic function have also been “attacked” using JPEG compression and noise addition. JPEG is a lossy compression technique designed to exploit known limitations of the human visual system, notably the fact that small changes in the image will go unnoticed when viewed by a human. Images processed by an optical system are often degraded by some random errors – this degradation is referred to as noise and may occur during image capture, transmission or processing [28]. Our studies have shown [27,19,29] that, in general, as the watermarks used had more lowpass characteristics, determined using the PSD, increased robustness to JPEG compression was observed. This is what is expected as lowpass watermarks have increased robustness with respect to image distortions that have lowpass characteristics (filtering, nonlinear filtering such as median filtering, lossy compression etc.) [30]. The maximum level of noise corruption which the detector can detect the embedded watermark was determined for numerous seed values of the logistic map in [27]. It was found that lowpass watermarks generated in the chaotic region of the logistic function perform much better in the presence of noise than highpass watermarks.

As a case study let us look at a previous study in “Chaos, Solitons and Fractals” by Dawei *et al.*, [31], which used the logistic map in the generation and embedding of watermarks. Watermark embedding takes place locally in the wavelet domain and watermark detection is performed using a correlation technique between the watermarked pixels and the watermark pixels. In this study the logistic map is used in two instances. The logistic map is used to select unique pixel blocks which will form a subimage. It is to this new subimage that the watermark is embedded before being replaced in the original image. Blocks of pixels of size  $8 \times 8$  are selected from the original image based on the output from the iteration of the logistic mapping as shown in Fig. 7(a). The logistic map is also used to generate the watermark in a similar format to that presented in this paper in Eq. 8. The chaotic watermark is embedded into part of the subbands of the wavelet domain representation of the image, after

a three level wavelet decomposition. This watermarking algorithm is novel, efficient and is successfully robust against attacks.

However, if the user of this algorithm is unaware of the properties of the logistic map, watermark generation and embedding fails. When the map is seeded with a value which results in the map entering a periodic orbit, the algorithm breaks down as can be seen, for example, in Fig. 7(b). In this case the logistic map was seeded with  $a = 3.83$ , for use in the selection of pixel blocks, and it can be seen that only three blocks of pixels are selected for use as the subimage to be watermarked. Similar results are obtained for any of the periodic regions of the logistic map if a periodic orbit is generated. If the logistic map is seeded with a value which results in the orbit being periodic, for the purpose of watermark generation, similar watermarks are produced to those shown in Fig. 3(b). Users must be fully aware of the properties of the logistic map before using this robust watermarking application.

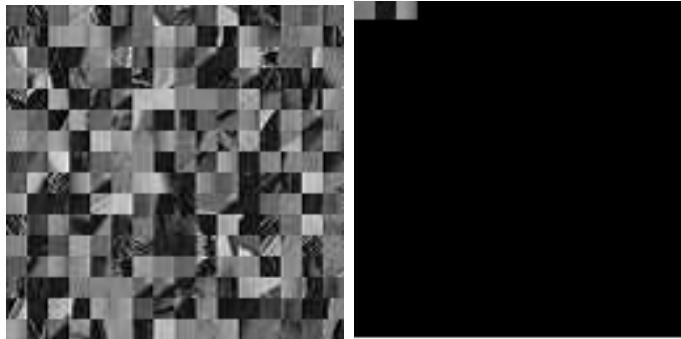


Fig. 7. Sample subimage generated from sequences of the logistic map where (a):  $a = 3.924$  and (b):  $a = 3.83$ .

## 7 Conclusion

In this paper, we have provided a detailed study of the generation of watermarks using the logistic map. The relevant properties of the logistic map were presented for the generation of optically detectable watermarks. It was shown how care must be taken in the selection of the function seed in the generation of chaotic watermarks. Reliance on the power spectrum alone is not enough to determine the spectral properties of the watermarks generated from the logistic map, unlike other well behaved chaotic maps, and has been presented here. A case study was presented based on another watermarking algorithm using the logistic map. We show that although this technique is an efficient robust watermarking technique but fails if the user is careless in the selection of the function seed of the logistic map.

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