

A Hybrid of Darboux's method and Singularity Analysis in Combinatorial Asymptotics

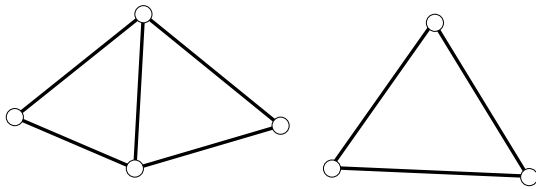
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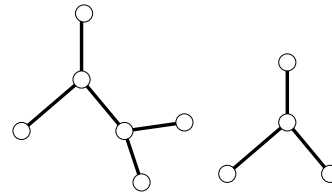
Motivations

- We consider **classes of objects**
Each object has a **size** (e.g. $\#(\text{vertices})$)

Graphs



Forests



Permutations

1 3 5 2 4

- Given a class \mathcal{C} , let c_n be the **number of objects of size n** in \mathcal{C}
- **Our aim:** find **automatic methods** for **estimating the coefficients c_n asymptotically**.

Generating functions

- Let \mathcal{C} be a class, with counting coefficients c_n .

OGF (unlabelled class): $C(z) = \sum c_n z^n$

EGF (labelled class): $C(z) = \sum_n c_n \frac{z^n}{n!}$

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- Dictionary for computing GF:

$\mathcal{C} = \mathcal{A} + \mathcal{B}$	$C(z) = A(z) + B(z)$
$\mathcal{C} = \mathcal{A} \times \mathcal{B}$	$C(z) = A(z) \cdot B(z)$
$\mathcal{C} = \text{Set}(\mathcal{A})$	$C(z) = \exp(A(z))$
$\mathcal{C} = \text{Cyc}(\mathcal{A})$	$C(z) = \log(1/(1 - A(z)))$

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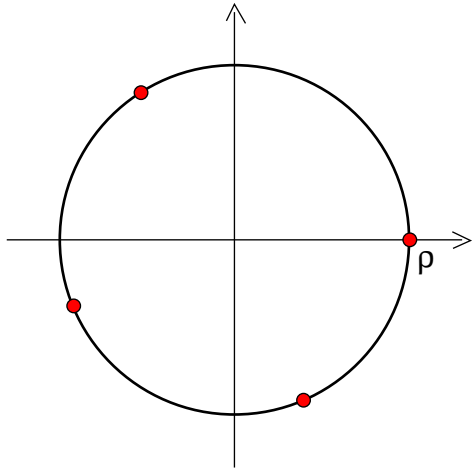
- Example:** permutations with **no fixed point**:

$$\mathcal{C} = \text{Set}(\text{Cyc}_{\geq 2}(\mathcal{Z}))$$

$$\Rightarrow C(z) = \exp \left(\log \left(\frac{1}{1-z} - z \right) \right) = \frac{e^{-z}}{1-z}.$$

Complex analysis

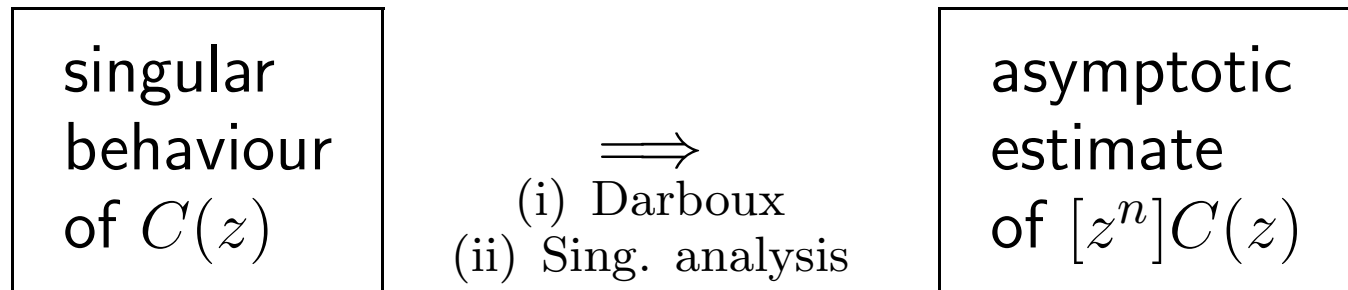
- Generating function $C(z)$ as a **complex function**



$$\rho = \limsup ([z^n]C(z))^{1/n}$$

$C(z)$ is **singular at ρ**
(Pringsheim)

- **Asymptotic methods:**



- **Remark:** we can **assume $\rho = 1$** without loss of generality, using $[z^n]C(z) = \rho^{-n}[z^n]C(\rho \cdot z)$

Coefficients of basic functions

- A **log-power function** at 1 is a linear combination $\sigma(z)$ of functions of the form

$$(1 - z)^\alpha \log^k \left(\frac{1}{1 - z} \right), \quad \alpha \in \mathbb{R}, \quad k \in \mathbb{Z}_{\geq 0}$$

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- Coefficients have a **full asymptotic expansion**

Example:

$$[z^n] \frac{\log \left(\frac{1}{1-z} \right)}{\sqrt{1-z}} \sim \frac{\log n + \gamma + 2 \log 2}{\sqrt{\pi n}} - \frac{\log n + \gamma + 2 \log 2}{8\sqrt{\pi n^3}} + \dots$$

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- Applies for log-power with **finitely many singularities** $\zeta_1, \dots, \zeta_\ell$, using $[z^n] \sigma(z/\zeta_i) = \zeta_i^{-n} [z^n] \sigma(z)$

Darboux's method

- **Key-remark:** if $g(z)$ is \mathcal{C}_s -smooth on the closed unit disk, then

$$[z^n]g(z) = o(n^{-s}).$$

(from Cauchy's integral formula + int. by part)

- **Application:** given $C(z) = \sum_n c_n z^n$, **decompose** $C(z)$ as a sum

$$C(z) = \underbrace{\Sigma(z)}_{\text{log-power}} + \underbrace{g(z)}_{\mathcal{C}_s\text{-smooth}}.$$

Then $[z^n]C(z) = [z^n]\Sigma(z) + o(n^{-s})$.

Singularity analysis

- There holds the **transfer rule** (Flajolet-Odlyzko'90)

$$C(z) = O_{z \rightarrow \rho}(g(z)) \xrightarrow{\text{transfer}} [z^n]C(z) = O([z^n]g(z))$$

(+ analytic continuation conditions to check)

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$$C(z) = \Sigma(z) + g(z)$$

where $\Sigma(z)$ is a **log-power** and $g(z)$ is $O(z - \zeta_i)^\alpha$ at ζ_i .

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- **Remark:** $g(z)$ is $\mathcal{C}_{\lfloor \alpha \rfloor}$ -smooth, as $[z^n]g(z) = O(n^{-\alpha-1})$.
By **Darboux**, this **gives only** $[z^n]g(z) = o(n^{-\lfloor \alpha \rfloor})$

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- Singularity analysis:

$$C(z) \underset{z \rightarrow 1}{\sim} \frac{e^{-3/4}}{\sqrt{1-z}} \quad \Rightarrow \quad [z^n]C(z) \sim \frac{e^{-3/4}}{\sqrt{\pi n}}$$

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- Darboux's method:

$$C(z) = \frac{e^{-3/4}}{\sqrt{1-z}} + * \sqrt{1-z} + \underbrace{O((1-z)^{3/2})}_{\mathcal{C}_1\text{-smooth}}$$

$$\Rightarrow [z^n]C(z) \sim \frac{e^{-3/4}}{\sqrt{\pi n}} + \frac{*}{n^{3/2}} + o(n^{-1}) \sim \frac{e^{-3/4}}{\sqrt{\pi n}}$$

The Hybrid method

- Classical Darboux's method: **sum-decomposition** of $C(z) = \sum_n c_n z^n$ as

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- **Idea:** in such cases, decompose $C(z)$ as a **product**, where the **first factor** gathers the **most salient singularities**

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$$P(z) = \underbrace{\Sigma(z)}_{\text{log-power}} + \underbrace{R(z)}_{\mathcal{C}_t\text{-smooth}} \quad || \quad Q(z) = \underbrace{\overline{Q}(z)}_{\substack{\text{interpolation} \\ \text{polynomial} \\ \text{of order } c}} + \underbrace{S(z)}_{\substack{\text{high contact} \\ \text{(order } c\text{)} \\ \text{at } \zeta_1, \dots, \zeta_\ell}}$$

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4. Adjust the parameters: with $t = \lfloor \frac{s+\lfloor a \rfloor}{2} \rfloor$, $c = \lfloor \frac{s-\lfloor a \rfloor}{2} \rfloor$,

we obtain $[z^n]C(z) = [z^n]\Sigma(z)\bar{Q}(z) + o(n^{-t})$

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Theorem: $\Sigma(z)\overline{Q}(z)$ is the sum of the radial singular expansions of order t of $C(z)$ at $\zeta_1, \dots, \zeta_\ell$.

2. Compute the radial expansions of $C(z)$ of order t
 (all computations within the disk of convergence)

Example (computation steps)

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1. **exp-log** transformation gives

$$C(z) = (1 + z) \exp\left(\sum_{\ell \geq 1} \frac{(-1)^{\ell-1}}{\ell} (A_\ell(z))\right),$$

where $A_\ell(z) = \text{Li}_\ell(z^\ell) - z^\ell$, $\text{Li}_\ell(t) = \sum_{n \geq 1} \frac{t^n}{n^\ell}$ (**polylog.**)

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2. $\text{Li}_\ell(t)$ is **C_s -smooth for $\ell \geq s + 2$** , so prod.-decomp. is

$$C(z) = \underbrace{(1+z) \exp\left(\sum_{\ell \leq s+1} A_\ell(z)\right)}_{P(z)} \cdot \underbrace{\exp\left(\sum_{\ell \geq s+2} A_\ell(z)\right)}_{Q(z)}$$

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3. **Singularities of $P(z)$** are roots of unity till order $s+1$,
 \Rightarrow **compute the radial expansions of $C(z)$** at these roots,
 using the **singular expansions of polylogarithms**
 (Zagier-Cohen'91, Flajolet'99)

Example (results)

Proposition: The probability that a permutation is made of cycles of distinct lengths admits a **full asymptotic expansion** of the form

$$f_n \sim e^{-\gamma} + \frac{e^{-\gamma}}{n} + \frac{e^{-\gamma}}{n^2} (-\log n - 1 - \gamma + \log 2) \\ + \frac{1}{n^3} \left[e^{-\gamma} (\log^2 n + c_{3,1} \log n + c_{3,0}) + 2(-1)^n + \Re \left(\frac{18\Gamma(\frac{2}{3})\omega^{-n}}{\Gamma(\frac{1}{6} + \frac{i\sqrt{3}}{6})\Gamma(\frac{1}{2} - \frac{i\sqrt{3}}{6})} \right) \right] \\ + \sum_{r \geq 4} \frac{P_r(n)}{n^r},$$

where $c_{3,1}$ and $c_{3,0}$ are explicit constants. Each $P_r(n)$ is a **polynomial** of degree $r - 1$ in $\log n$ with coefficients that are **periodic functions of n** with period $D(r) = \text{lcm}(2, 3, \dots, r)$.

Scope of applications

- The hybrid method **typically applies** for a function $C(z)$ of the form

$$C(z) = \prod_{n \geq 1} (1 + c_n z^n), \text{ with } c_n \sim n^{-\alpha}, \alpha \geq 1$$

$Q_{-1} = \prod_{n \geq 1} \left(1 + \frac{z^n}{n}\right)$	Perm. dist. cycle lengths	$\propto 1$
$Q_{-3/2} = \prod_{n \geq 1} \left(1 + \frac{z^n}{n^{3/2}}\right)$	Forest dist. comp. sizes	$\propto n^{-3/2}$
$Q_{-2} = \prod_{n \geq 1} \left(1 + \frac{z^n}{n^2}\right)$	Perm. same cycle type	$\propto n^{-2}$

- Outside of scope:**
 - Partitions dist. summands: $\prod_{n \geq 1} (1 + z^n)$
 - Set partitions dist. sizes: $\prod_{n \geq 1} \left(1 + \frac{z^n}{n!}\right)$