



A note on Frobenius inner product and the m -distance matrices of a tree

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Abstract

The m -distance matrix D_m of a simple connected undirected graph has an important role in computing the distance matrix D of the graph from the powers of the adjacency matrix using Hadamard product. This paper shows that for an undirected tree T with diameter d , $\{D_0, D_1, \dots, D_d\}$ is an orthogonal basis for the space spanned by the binary equivalent matrices of the first $d + 1$ powers of the adjacency matrix of T and it gives an invertible conversion matrix for finding the m -distance matrix of T using Frobenius inner product on matrices.

Keywords

Adjacency matrix, Distance matrix, Binary matrix, Diameter, Hadamard product, m -distance matrix, Frobenius inner product, Frobenius norm.

AMS Subject Classification

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1. Introduction

Consider a simple, connected, undirected graph $G = (V, E)$ of order n with vertex set V and edge set E throughout this paper unless otherwise specified. Let A_G be the $n \times n$ adjacency matrix [1] of G . Then i_j th entry of A_G^m (m th power of A_G), represent the number of walks of length m between the vertices v_i and v_j of G .

Definition 1.1. ([2]) The distance matrix, $D = (d_{ij})$ of G is defined as,

$$d_{ij} = \begin{cases} d(v_i, v_j) & , \text{ if } i \neq j \\ 0 & , \text{ if } i = j \end{cases}$$

Where, $d(v_i, v_j)$ is the distance between the vertices v_i and v_j .

Definition 1.2. ([3]) The diameter of a graph G is the maximum distance between any two vertices of G and it is denoted by $Diam(G)$.

Definition 1.3. (Hadamard product [4]) Consider the vector space $\mathbb{R}^{m \times n}$ of all $m \times n$ real matrices over the real field \mathbb{R} . For $P, F \in \mathbb{R}^{m \times n}$, the Hadamard product \circ is a binary operation on $\mathbb{R}^{m \times n}$ defined by,

$$(P \circ F)_{ij} := (P)_{ij}(F)_{ij}, \quad \forall i, j.$$

Let $B = \{0, 1\}$ and let $B_{m \times n}$ denote the set of all binary matrices in $\mathbb{R}^{m \times n}$.

Then for $P, F \in B_{m \times n}$,

$$(P \circ F)_{ij} = \begin{cases} 1, & \text{ if } p_{ij} = 1 \text{ and } f_{ij} = 1 \\ 0, & \text{ otherwise} \end{cases}$$

Definition 1.4. (Frobenius inner product [5]) Consider the vector space $\mathbb{R}^{n \times n}$ over the field \mathbb{R} . Then the Frobenius inner product $\langle \cdot, \cdot \rangle_F : \mathbb{R}^{n \times n} \times \mathbb{R}^{n \times n} \rightarrow \mathbb{R}$ is defined by,

$$\langle P, Q \rangle_F = \sum_{i=1}^n \sum_{j=1}^n p_{ij}q_{ij} = \sum_{i=1}^n \sum_{j=1}^n (P \circ Q)_{ij} = Tr(PQ^T),$$

where $P, Q \in \mathbb{R}^{n \times n}$. The Frobenius norm $\|\cdot\|$ induced by this inner product on $\mathbb{R}^{n \times n}$ is, $\|P\|_F = (\langle P, P \rangle_F)^{\frac{1}{2}}$, $P \in \mathbb{R}^{n \times n}$. Then

the metric d_F induced by this norm on $\mathbb{R}^{n \times n}$ is $d_F(P, Q) = \|P - Q\|_F$, $P, Q \in \mathbb{R}^{n \times n}$.

Remark 1.5. (i) $\|P\|_F = (\langle P, P \rangle_F)^{\frac{1}{2}} = (Tr(PP^T))^{\frac{1}{2}}$, $P \in \mathbb{R}^{n \times n}$

(ii) Let $X, Y \in \mathbb{R}^{1 \times n}$. Then $\langle X, Y \rangle_F = \sum_{i=1}^n x_i y_i$ and

$$d_F(X, Y) = \|X - Y\|_F = \left(\sum_{i=1}^n (x_i - y_i)^2 \right)^{\frac{1}{2}}$$

which is the Euclidean distance between the vectors X and Y of \mathbb{R}^n .

(iii) For $P \in B_{n \times n}$,

$$d_F(P, 0) = \|P - 0\|_F = \|P\|_F = \left(\sum_{i=1}^n \sum_{j=1}^n p_{ij}^2 \right)^{\frac{1}{2}}$$

Then $(d_F(P, 0))^2$ represents the number of 1's in P . For $P, Q \in B_{n \times n}$,

$$d_F(P, Q) = \|P - Q\|_F = \left(\sum_{i=1}^n \sum_{j=1}^n (p_{ij} - q_{ij})^2 \right)^{\frac{1}{2}}$$

Then $(d_F(P, Q))^2$ represent the number of non identical entries in P and Q and the number of identical entries in P and Q is $n^2 - (d_F(P, Q))^2$

2. Preliminaries

Definition 2.1. ([6]) The function $\delta : \mathbb{R} \rightarrow B = \{0, 1\}$ is defined by

$$\delta(a) := \begin{cases} 0, & a \leq 0 \\ 1, & \text{otherwise} \end{cases}, a \in \mathbb{R}$$

Also, $\delta : \mathbb{R}^{m \times n} \rightarrow B_{m \times n}$ is defined by $\delta(D) = (\delta(D))_{ij} := (\delta(d_{ij}))$, $\forall i, j$ and $D \in \mathbb{R}^{m \times n}$. Then $\delta(A_G^m) \in B_{n \times n}$, $\forall m \in \{0, 1, 2, \dots\}$ and it is the equivalent binary matrix representation of A_G^m . Let $A_G^{(m)} = \delta(A_G^m)$. Then

$$(A_G^{(m)})_{ij} = \begin{cases} 1, & \text{if there exist a walk of length } m \\ & \text{from } v_i \text{ to } v_j \text{ in } G \\ 0, & \text{otherwise} \end{cases}$$

Also, for $C, F \in B_{m \times n}$, $\delta(C \circ F) = \delta(C) \circ \delta(F)$.

Definition 2.2. (m -distance matrix [6]) Consider the graph $G = (V, E)$ with n vertices. Then the m -distance matrix D_m of G is an $n \times n$ symmetric binary matrix defined by

$$(D_m)_{ij} := \begin{cases} 1, & \text{if } d(v_i, v_j) = m \\ 0, & \text{otherwise} \end{cases}$$

where, $d(v_i, v_j)$ is the distance between the vertices v_i and v_j .

Remark 2.3. (i) $D_m \in B_{n \times n}$, $\forall m = 0, 1, 2, \dots$

(ii) $D_0 = I_n$

(iii) $D_0 + D_1 + \dots + D_d =$

$$J_n = \begin{bmatrix} 1 & 1 & \dots & \dots & 1 \\ 1 & 1 & \dots & \dots & 1 \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & 1 & \dots & \dots & 1 \end{bmatrix} = \text{The matrix of ones}$$

(iv) $Diam(G) = \underset{m}{\text{Max}}\{m : D_m \neq 0\}$

(v) $D_i \circ D_j = \begin{cases} 0, & \text{if } i \neq j \\ D_i, & \text{if } i = j \end{cases}$, for $0 \leq i, j \leq d$.

i.e., $\{D_0, D_1, \dots, D_d\}$ is an orthogonal subset $\mathbb{R}^{n \times n}$ with respect to the binary operation \circ .

Theorem 2.4. ([6]) Let A_G be the adjacency matrix of $G = (V, E)$ with diameter d . Then $\beta_2 = \{D_0, D_1, \dots, D_d\}$ is a linearly independent subset of the vector space $\mathbb{R}^{n \times n}$.

Theorem 2.5. ([6]) Let A_G be the adjacency matrix of $G = (V, E)$ with diameter d . Then for $1 \leq m \leq d$, $D_m = A_G^{(m)} - \delta\left(\sum_{s=0}^{m-1} A_G^{(m)} \circ A_G^{(s)}\right)$, where D_m is the m -distance matrix of G .

3. Main Results

Theorem 3.1. Let A_G be the adjacency matrix of $G = (V, E)$. Let $d = diam(G)$. Then for $0 \leq K, m \leq d$,

$$D_k \circ A_G^{(m)} = \begin{cases} D_m, & \text{if } K = m \\ 0, & \text{if } m < K. \end{cases}$$

Proof. **Case(i):** $K = m$

$$\begin{aligned} (D_m \circ A_G^{(m)})_{ij} &= (D_m)_{ij} (A_G^{(m)})_{ij} \\ &= \begin{cases} 1, & \text{if } (D_m)_{ij} = 1 \text{ and } (A_G^{(m)})_{ij} = 1 \\ 0, & \text{otherwise} \end{cases} \\ &= \begin{cases} 1, & \text{if } d(v_i, v_j) = m \text{ and } (A_G^{(m)})_{ij} = 1 \\ 0, & \text{otherwise.} \end{cases} \end{aligned}$$

But $d(v_i, v_j) = m \Rightarrow (A_G^{(m)})_{ij} = 1$

$$\begin{aligned} \therefore (D_m \circ A_G^{(m)})_{ij} &= (D_m)_{ij} (A_G^{(m)})_{ij} \\ &= \begin{cases} 1, & \text{if } d(v_i, v_j) = m \\ 0, & \text{otherwise} \end{cases} = (D_m)_{ij} \end{aligned}$$

That is, $(D_m \circ A_G^{(m)})_{ij} = D_m$.

Case(ii): $K > m$

$$(D_k \circ A_G^{(m)})_{ij} = (D_k)_{ij} (A_G^{(m)})_{ij}$$



If $(D_K)_{ij} = 1$, then $d(v_i, v_j) = K$.

That is, if there exist a shortest path of length K between v_i and v_j . But then this path cannot be traversed by any of the walk between v_i and v_j of length $< K$. Thus $(A_G^{(m)})_{ij} = 0$. ($\because K > m$)

$$\therefore (D_K)_{ij}(A_G^{(m)})_{ij} = 0$$

If $(D_K)_{ij} = 0$, then also $(D_K)_{ij}(A_G^{(m)})_{ij} = 0$.

Therefore, $(D_K \circ A_G^{(m)})_{ij} = 0, \forall i, j, \forall K > m$.

ie., $D_K \circ A_G^{(m)} = 0$ if $K > m$. Hence

$$D_K \circ A_G^{(m)} = \begin{cases} D_m, & \text{if } K = m \\ 0, & \text{if } m < K \end{cases}, \text{ for } K \leq d.$$

□

Theorem 3.2. Let A_G be the adjacency matrix of $G = (V, E)$. Let $d = \text{diam}(G)$. Then $\beta_1 = \{A_G^{(0)}, A_G^{(1)}, \dots, A_G^{(d)}\}$ is a linearly independent subset of the vector space $\mathbb{R}^{n \times n}$.

Proof. Suppose

$$c_0 \cdot A_G^{(0)} + c_1 \cdot A_G^{(1)} + \dots + c_d \cdot A_G^{(d)} = 0 \quad (3.1)$$

for some $c_0, c_1, \dots, c_d \in \mathbb{R}$. By taking Hadamard product \circ by D_d

$$\begin{aligned} (3.1) \Rightarrow D_d \circ (c_0 \cdot A_G^{(0)} + c_1 \cdot A_G^{(1)} + \dots + c_d \cdot A_G^{(d)}) &= 0 \\ c_0 \cdot (D_d \circ A_G^{(0)}) + c_1 \cdot (D_d \circ A_G^{(1)}) + \dots + c_d \cdot (D_d \circ A_G^{(d)}) &= 0 \\ \Rightarrow 0 + 0 + \dots + 0 + c_d \cdot (D_d \circ A_G^{(d)}) &= 0 (\because \text{Theorem 3.1}) \\ \Rightarrow c_d \cdot D_d &= 0 \Rightarrow c_d = 0, \text{ as } D_d \neq 0 \end{aligned}$$

$$(3.1) \Rightarrow c_0 \cdot A_G^{(0)} + c_1 \cdot A_G^{(1)} + \dots + c_{d-1} \cdot A_G^{(d-1)} = 0 \quad (3.2)$$

By taking Hadamard product \circ by D_{d-1} on both side of equation (3.2).

$$\begin{aligned} (3.2) \Rightarrow D_{d-1} \circ (c_0 \cdot A_G^{(0)} + c_1 \cdot A_G^{(1)} + \dots + c_{d-1} \cdot A_G^{(d-1)}) &= 0 \\ c_0 \cdot (D_{d-1} \circ A_G^{(0)}) + c_1 \cdot (D_{d-1} \circ A_G^{(1)}) + \dots & \\ + c_{d-1} \cdot (D_{d-1} \circ A_G^{(d-1)}) &= 0 \\ \Rightarrow 0 + 0 + \dots + 0 + c_{d-1} \cdot (D_{d-1} \circ A_G^{(d-1)}) &= 0 \\ \Rightarrow c_{d-1} D_{d-1} &= 0 \Rightarrow c_{d-1} = 0, \text{ as } D_{d-1} \neq 0 \end{aligned}$$

$$(3.2) \Rightarrow c_0 \cdot A_G^{(0)} + c_1 \cdot A_G^{(1)} + \dots + c_{d-1} \cdot A_G^{(d-1)} = 0$$

Continue the above process further a finite number of times, we get $c_0 = 0, c_1 = 0, \dots, c_d = 0$.

$\Rightarrow \beta_1 = \{A_G^{(0)}, A_G^{(1)}, \dots, A_G^{(d)}\}$ is a linearly independent subset of $\mathbb{R}^{n \times n}$. □

Remark 3.3. $\beta_1 = \{A_G^{(0)}, A_G^{(1)}, \dots, A_G^{(d)}\}$ is a basis for span $\{A_G^{(0)}, A_G^{(1)}, \dots, A_G^{(d)}\}$.

Lemma 3.4. Let $G = (V, E)$ be a connected undirected graph of order $n (n > 1)$. Let d_{ij} denote the distance between v_i and v_j . Then there always exist at least one walk between v_i and v_j of length $d_{ij} + 2r, \forall r = 0, 1, 2, \dots$. That is,

$$(A_G^{(d_{ij}+2r)})_{ij} = 1, \forall r = 0, 1, 2, \dots$$

Proof. Since G is a connected undirected graph, there exist at least one path in between v_i and v_j . Let $P_{ij} : v_0 (= v_i) - v_1 - v_2 - \dots - v_{d_{ij}-1} - v_{d_{ij}} (= v_j)$ denote the shortest path between v_i and v_j of length d_{ij} . If we traverse back and forth once along the last edge in P_{ij} , then we get a walk $v_0 (= v_i) - v_1 - v_2 - \dots - v_{d_{ij}-1} - v_{d_{ij}} (= v_j) - v_{d_{ij}-1} - v_{d_{ij}} (= v_j)$ of length $d_{ij} + 2$ between v_i and v_j . Also If we traverse back and forth twice along the last edge in P_{ij} , then we get a walk $v_0 (= v_i) - v_1 - v_2 - \dots - v_{d_{ij}-1} - v_{d_{ij}} - v_{d_{ij}-1} - v_{d_{ij}} - v_{d_{ij}-1} - v_{d_{ij}} (= v_j)$ of length $d_{ij} + 4$ between v_i and v_j . If we proceed like this, then we always get a walk of length $d_{ij} + 2r, \forall r = 0, 1, 2, \dots$ between v_i and v_j in G . This walk will reflect as 1 in the ij^{th} entry of $A_G^{(d_{ij}+2r)}$.

$$\therefore (A_G^{(d_{ij}+2r)})_{ij} = 1, \forall r = 0, 1, 2, \dots$$

□

Theorem 3.5. Let $T = (V, E)$ be an undirected tree of order $n (n > 1)$. Let d_{ij} denote the distance between v_i and v_j . Then,

$$(A_T^{(m)})_{ij} = \begin{cases} 1, & \text{if } m = d_{ij} + 2r, \forall r = 0, 1, 2, \dots \\ 0, & \text{otherwise} \end{cases}$$

Proof. Let T_{ij} denote the unique path in T from v_i to v_j of distance d_{ij} . Consider a walk W_{ij} from v_i to v_j such that length of $W_{ij}, l(W_{ij}) > d_{ij}$. Now delete the edges of T_{ij} from W_{ij} . Let W_{ij}^1 be the remaining part of the walk W_{ij} . Then W_{ij}^1 should be disconnected. Otherwise W_{ij}^1 and some of the edges of T_{ij} together would form a cycle. But T has no cycle.

Let $W_{ij,1}, W_{ij,2}, \dots, W_{ij,s}$ be the component walks of W_{ij}^1 . Then $W_{ij,h} \cap T_{ij}, (1 \leq h \leq s)$ should be a single vertex $v_{w_{ij,h}}$, otherwise some of the edges of $W_{ij,h}$ and T_{ij} together form a cycle. Since each vertex $v_{w_{ij,h}}$ is a part of the walk $W_{ij}, W_{ij,h}$ should be either the single vertex $v_{w_{ij,h}}$ or a closed walk from $v_{w_{ij,h}}$ to $v_{w_{ij,h}}$.

But length of a closed walk in a tree is always even. So $l(W_{ij}^1) = \sum_{h=1}^s l(W_{ij,h})$ must be even.

Let $l(W_{ij}^1) = \sum_{h=1}^s l(W_{ij,h}) = 2r$, for some $r \in \{0, 1, 2, \dots\}$
 $\Rightarrow l(W_{ij}) = l(T_{ij}) + l(W_{ij}^1) = d_{ij} + \sum_{h=1}^s l(W_{ij,h}) = d_{ij} + 2r$, for some $r \in \{0, 1, 2, \dots\}$

\Rightarrow length of every walk from v_i to v_j must be of the form $d_{ij} + 2r$, for $r \in \{0, 1, 2, \dots\}$.

By Lemma 3.4, there always exist a walk between v_i and v_j of length $d_{ij} + 2r$, for all $r = 0, 1, 2, \dots$. So

$$(A_T^{(m)})_{ij} = \begin{cases} 1, & \text{if } m = d_{ij} + 2r, \forall r = 0, 1, 2, \dots \\ 0, & \text{otherwise} \end{cases}$$

□



Theorem 3.6. Let A_T be the adjacency matrix of a tree $T = (V, E)$ of order $n(n > 1)$. Let $d = \text{diam}(T)$. Then for $0 \leq K \leq d$,

$$(i) \quad D_K \circ A_T^{(K+2r)} = D_K.$$

$$(ii) \quad D_K \circ A_T^{(K+2r+1)} = 0, \forall r = 0, 1, 2, \dots$$

Proof.

$$\begin{aligned} (D_K \circ A_T^{(m)})_{ij} &= (D_K)_{ij} (A_T^{(m)})_{ij} \\ &= \begin{cases} 1, & \text{if } (D_K)_{ij} = 1 \text{ and } (A_T^{(m)})_{ij} = 1 \\ 0, & \text{otherwise} \end{cases} \end{aligned}$$

$$(D_K)_{ij} = 1 \Rightarrow d_{ij} = K.$$

But then $(A_T^{(K+2r)})_{ij} = 1, \forall r = 0, 1, 2, \dots$ (\because Theorem 3.5)

ie., $(D_K)_{ij} = 1 \Rightarrow (A_T^{(K+2r)})_{ij} = 1, \forall r = 0, 1, 2, \dots$

$$\begin{aligned} \therefore (D_K \circ A_T^{(K+2r)})_{ij} &= \\ \begin{cases} 1, & \text{if } K = d_{ij} \\ 0, & \text{otherwise} \end{cases} &= (D_K)_{ij}, \forall i, j \text{ and } r = 0, 1, 2, \dots \end{aligned}$$

$$\Rightarrow D_K \circ A_T^{(K+2r)} = D_K, \forall r = 0, 1, 2, \dots$$

By Theorem 3.5 there does not exist a walk of length $d_{ij} + 2r + 1$ between v_i and $v_j, \forall r = 0, 1, 2, \dots$

$$\Rightarrow (A_T^{(d_{ij}+2r+1)})_{ij} = 0, \forall i, j \text{ and } r = 0, 1, 2, \dots$$

$$\Rightarrow D_K \circ A_T^{(K+2r+1)} = 0 \forall r = 0, 1, 2, \dots \quad \square$$

Remark 3.7. Let A_T be the adjacency matrix of a tree $T = (V, E)$ of order $n(n > 1)$. Let $d = \text{diam}(T)$. Then for $0 \leq K \leq d$. For $m < K, A_T^{(m)} \circ D_K = 0$, (by Theorem 3.1).

For $m > K$,

$$(i) \quad D_K \circ A_T^{(K+2r)} = D_K.$$

$$(ii) \quad D_K \circ A_T^{(K+2r+1)} = 0. \forall r = 0, 1, 2, \dots, \text{ (by Theorem 3.6)}$$

$$\therefore A_T^{(m)} \circ D_K = \begin{cases} D_K, & m = K + 2r, \forall r = 0, 1, 2, \dots \\ 0, & \text{otherwise} \end{cases}$$

Theorem 3.8. Let A_T be the adjacency matrix of $T = (V, E)$ with $d = \text{diam}(G)$. Then $\beta_2 = \{D_0, D_1, \dots, D_d\}$ is an orthogonal basis for span $\{A_T^{(0)}, A_T^{(1)}, \dots, A_T^{(d)}\}$.

Proof. By Theorem 2.4, $\beta_2 = \{D_0, D_1, \dots, D_d\}$ is linearly independent. By Remark 2.3(v) β_2 is orthogonal also. So it is enough to prove that $A_T^{(m)}$ is a linear combination of D_0, D_1, \dots, D_d , for $0 \leq m \leq d$. By Remark 3.7,

$$\therefore A_T^{(m)} \circ D_K = \begin{cases} D_K, & m = K + 2r, \forall r = 0, 1, 2, \dots \\ 0, & \text{otherwise} \end{cases}$$

$\Rightarrow \sum_{K=0}^d A_T^{(m)} \circ D_K$ is a linear combination of D_0, D_1, \dots, D_d . i.e., $\sum_{K=0}^d A_T^{(m)} \circ D_K = \sum_{K=0}^d c_{m,K} D_K$, for $c_{m,0}, c_{m,1}, \dots, c_{m,d} \in B = \{0, 1\}$. Now

$$A_T^{(m)} = A_T^{(m)} \circ J_n = A_T^{(m)} \circ (D_0 + D_1 + \dots + D_d), (\because \text{Remark 2.3(iii)})$$

$$= \sum_{K=0}^d A_T^{(m)} \circ D_K = \sum_{K=0}^d c_{m,K} D_K$$

$\Rightarrow A_T^{(m)}$ is a linear combination of D_0, D_1, \dots, D_d .

For finding the scalars $c_{m,k}$, consider

$$\begin{aligned} \langle A_T^{(m)}, D_j \rangle_F &= \langle c_{m,0} \cdot D_0 + c_{m,1} \cdot D_1 + \dots + c_{m,d} \cdot D_d, D_j \rangle_F \\ &= c_{m,0} \cdot \langle D_0, D_j \rangle_F + c_{m,1} \cdot \langle D_1, D_j \rangle_F + \dots \\ &\quad + c_{m,d} \cdot \langle D_d, D_j \rangle_F \\ &= 0 + 0 + \dots + c_{m,j} \langle D_j, D_j \rangle_F + 0 + \dots + 0 \\ &= c_{m,j} \langle D_j, D_j \rangle_F (\because \text{Remark 2.3(v)}) \end{aligned}$$

$$\Rightarrow c_{m,j} = \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F}, j = 0, 1, \dots, d$$

the scalars, $c_{m,j} = \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F} \in B = \{0, 1\}, \forall 0 \leq m, j \leq d$,

So,

$$\begin{aligned} A_T^{(m)} &= c_{m,0} \cdot D_0 + c_{m,1} \cdot D_1 + \dots + c_{m,d} \cdot D_d \\ &= \frac{\langle A_T^{(m)}, D_0 \rangle_F}{\langle D_0, D_0 \rangle_F} \cdot D_0 + \frac{\langle A_T^{(m)}, D_1 \rangle_F}{\langle D_1, D_1 \rangle_F} \cdot D_1 + \dots \\ &\quad + \frac{\langle A_T^{(m)}, D_d \rangle_F}{\langle D_d, D_d \rangle_F} \cdot D_d \\ A_T^{(m)} &= \sum_{j=0}^d \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F} D_j \end{aligned}$$

$\Rightarrow \beta_2 = \{D_0, D_1, \dots, D_d\}$ is an orthogonal basis for span $\{A_T^{(0)}, A_T^{(1)}, \dots, A_T^{(d)}\}$. \square

Remark 3.9. (i) We have

$$\begin{aligned} A_T^{(m)} &= \sum_{j=0}^d \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F} D_j (\because \text{Theorem 3.8}) \\ A_T^{(m)} &= \frac{\langle A_T^{(m)}, D_d \rangle_F}{\langle D_d, D_d \rangle_F} D_d + \sum_{j=0}^{d-1} \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F} D_j \\ \frac{\langle A_T^{(m)}, D_d \rangle_F}{\langle D_d, D_d \rangle_F} D_d &= A_T^{(m)} - \sum_{j=0}^{d-1} \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F} D_j \end{aligned}$$

\Rightarrow The orthogonal basis for span $\{A_T^{(0)}, A_T^{(1)}, \dots, A_T^{(d)}\}$ obtained from the basis β_1 by Gram-Schmidt process is nothing but β_2 itself.

(ii) $A_T^{(m)} = \sum_{j=0}^d c_{m,j} D_j$, by Theorem 3.8, where

$$c_{m,j} = \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F}, m, j = 0, 1, \dots, d.$$



This can be written in the matrix form as given below,

$$\begin{bmatrix} A_T^{(0)} \\ A_T^{(1)} \\ \vdots \\ A_T^{(d)} \end{bmatrix} = \begin{bmatrix} c_{00} & c_{01} & \cdot & \cdot & \cdot & c_{0d} \\ c_{10} & c_{11} & \cdot & \cdot & \cdot & c_{1d} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ c_{d0} & c_{d1} & \cdot & \cdot & \cdot & c_{dd} \end{bmatrix} \begin{bmatrix} D_0 \\ D_1 \\ \cdot \\ \cdot \\ D_d \end{bmatrix} \quad (3.3)$$

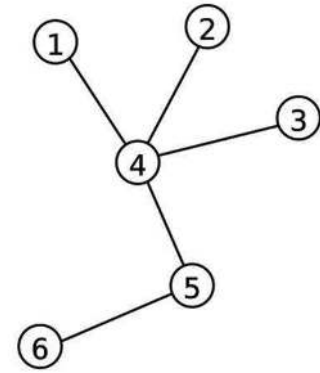


Figure 1

Let

$$Y = \begin{bmatrix} A_T^{(0)} \\ A_T^{(1)} \\ \vdots \\ A_T^{(d)} \end{bmatrix}, C = \begin{bmatrix} c_{00} & c_{01} & \cdot & \cdot & \cdot & c_{0d} \\ c_{10} & c_{11} & \cdot & \cdot & \cdot & c_{1d} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ c_{d0} & c_{d1} & \cdot & \cdot & \cdot & c_{dd} \end{bmatrix}, X = \begin{bmatrix} D_0 \\ D_1 \\ \cdot \\ \cdot \\ D_d \end{bmatrix}$$

Then (3.3) \Rightarrow

$$Y = CX \quad (3.4)$$

where, X, Y are $(d + 1) \times 1$ column matrices whose elements are $n \times n$ binary matrices and C is a $(d + 1) \times (d + 1)$ real matrix. But by Theorem 3.1, $c_{ii} = 1$ and $c_{ij} = 0$, when $i < j, (i, j = 0, 1, \dots, d)$. Which implies that C is a unit lower triangular matrix.

$\therefore C$ is invertible and C^{-1} is a lower triangular unit matrix with $|C| = 1$.

So, by multiplying C^{-1} on both sides of equation (3.4). Then (3.4) \Rightarrow

$$X = C^{-1}Y \quad (3.5)$$

Then (3.3) \Rightarrow

$$\begin{bmatrix} D_0 \\ D_1 \\ \cdot \\ \cdot \\ D_d \end{bmatrix} = \begin{bmatrix} c_{00} & c_{01} & \cdot & \cdot & \cdot & c_{0d} \\ c_{10} & c_{11} & \cdot & \cdot & \cdot & c_{1d} \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ c_{d0} & c_{d1} & \cdot & \cdot & \cdot & c_{dd} \end{bmatrix}^{-1} \begin{bmatrix} A_T^{(0)} \\ A_T^{(1)} \\ \cdot \\ \cdot \\ A_T^{(d)} \end{bmatrix}$$

$\Rightarrow C^{-1}$ and C are the conversion matrices for getting the orthogonal basis β_2 from β_1 and vice versa.

4. Illustration

Consider the following tree $T = (V, E)$. Then the adjacency matrix A_T of T is

$$A_T = \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{matrix} & \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix} \end{matrix}$$

Here the maximum distance $d = 3, A_T^{(0)} = I, A_T^{(1)} = \delta(A_T) = A_T$

$$\begin{aligned} A_T^2 &= \begin{bmatrix} 1 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 4 & 0 & 1 \\ 1 & 1 & 1 & 0 & 2 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}, \\ A_T^{(2)} = \delta(A_T^{(2)}) &= \begin{bmatrix} 1 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \end{bmatrix}, \\ A_T^{(2)} * A_T^{(1)} &= \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}, \\ A_T^3 &= \begin{bmatrix} 0 & 0 & 0 & 4 & 0 & 1 \\ 0 & 0 & 0 & 4 & 0 & 1 \\ 0 & 0 & 0 & 4 & 0 & 1 \\ 4 & 4 & 4 & 0 & 5 & 0 \\ 0 & 0 & 0 & 5 & 0 & 2 \\ 1 & 1 & 1 & 0 & 2 & 0 \end{bmatrix} \end{aligned}$$



$$A_T^{(3)} = \delta(A_T^{(3)}) = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 1 & 0 \end{bmatrix}$$

$$A_T^{(3)} * A_T^{(1)} = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix},$$

$$A_G^{(3)} * A_T^{(2)} = \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$A_T^{(1)} * I = 0$$

$$D_0 = A_T^{(0)} = I$$

$$D_1 = A_T^{(1)} - \delta(A_T^{(1)} * I) = A_T^{(1)} - 0 = A_T^{(1)}$$

$$(A_T^{(2)} * I + A_T^{(2)} * A_T^{(1)})$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix} + \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \end{bmatrix}$$

$$= \begin{bmatrix} 1 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

$$D_2 = A_T^{(2)} - \delta(A_T^{(2)} * I + A_T^{(2)} * A_T^{(1)})$$

$$= \begin{bmatrix} 0 & 1 & 1 & 0 & 1 & 0 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 1 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \end{bmatrix}$$

$$(A_T^{(3)} * I + A_T^{(3)} * A + A_T^{(3)} * A^2) = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 \\ 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

$$D_3 = A_T^{(3)} - \delta(A_T^{(3)} * I + A_T^{(3)} * A_T^{(1)} + A_T^{(3)} * A_T^{(2)})$$

$$= \begin{bmatrix} 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

Let

$$Y = \begin{bmatrix} A_T^{(0)} \\ A_T^{(1)} \\ A_T^{(2)} \\ A_T^{(3)} \end{bmatrix}, X = \begin{bmatrix} D_0 \\ D_1 \\ D_2 \\ D_3 \end{bmatrix}$$

Then $Y = CX$, where

$$c_{m,j} = \frac{\langle A_T^{(m)}, D_j \rangle_F}{\langle D_j, D_j \rangle_F}, m, j = 0, 1, 2, 3$$

$$c_{00} = \frac{\langle A_T^{(0)}, D_0 \rangle_F}{\langle D_0, D_0 \rangle_F} = \frac{6}{6} = 1, c_{01} = c_{02} = c_{03} = 0$$

$$c_{10} = \frac{\langle A_T^{(1)}, D_0 \rangle_F}{\langle D_0, D_0 \rangle_F} = \frac{0}{6} = 0,$$

$$c_{11} = \frac{\langle A_T^{(1)}, D_1 \rangle_F}{\langle D_1, D_1 \rangle_F} = \frac{10}{10} = 1, c_{12} = c_{13} = 0$$

$$c_{20} = \frac{\langle A_T^{(2)}, D_0 \rangle_F}{\langle D_0, D_0 \rangle_F} = \frac{6}{6} = 1, c_{21} = \frac{\langle A_T^{(2)}, D_1 \rangle_F}{\langle D_1, D_1 \rangle_F} = \frac{0}{10} = 0,$$

$$c_{22} = \frac{\langle A_T^{(2)}, D_2 \rangle_F}{\langle D_2, D_2 \rangle_F} = \frac{14}{14} = 1, c_{23} = 0,$$

$$c_{30} = \frac{\langle A_T^{(3)}, D_0 \rangle_F}{\langle D_0, D_0 \rangle_F} = \frac{0}{6} = 0, c_{31} = \frac{\langle A_T^{(3)}, D_1 \rangle_F}{\langle D_1, D_1 \rangle_F} = \frac{10}{10} = 1,$$

$$c_{32} = \frac{\langle A_T^{(3)}, D_2 \rangle_F}{\langle D_2, D_2 \rangle_F} = \frac{0}{14} = 0, c_{33} = \frac{\langle A_T^{(3)}, D_3 \rangle_F}{\langle D_3, D_3 \rangle_F} = \frac{6}{6} = 1.$$

So,

$$C = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \end{bmatrix}, C^{-1} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ -1 & 0 & 1 & 0 \\ 0 & -1 & 0 & 1 \end{bmatrix}$$

Then $Y = CX$, also $X = C^{-1}Y$

$$Y = CX \Rightarrow$$

$$A_T^{(0)} = D_0$$

$$A_T^{(1)} = D_1$$

$$A_T^{(2)} = D_0 + D_2$$

$$A_T^{(3)} = D_1 + D_3$$

$$X = C^{-1}Y \Rightarrow$$

$$D_0 = A_T^{(0)}$$

$$D_1 = A_T^{(1)}$$

$$D_2 = (-1)A_T^{(0)} + A_T^{(2)}$$

$$D_3 = (-1)A_T^{(1)} + A_T^{(3)}$$



5. Conclusion

Generally, $\beta_2 = \{D_0, D_1, \dots, D_d\}$ is not an orthogonal basis for $\text{span} \{A_G^{(0)}, A_G^{(1)}, \dots, A_G^{(d)}\}$, for a simple connected undirected graph G with diameter d . But we proved that this is true for an undirected tree T with diameter d and also derived an invertible conversion matrix for computing one basis β_1 from the other basis β_2 and vice versa. Further study may be done on exploring some other connected undirected graphs having this property other than trees.

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