An Adaptive Distance Vector Routing Algorithm for Mobile, Ad Hoc Networks

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Outline

- Introduction
- Current routing algorithms for MANETs
- Adaptive distance vector
- Performance analysis
- Conclusions

Introduction

- There are two category of routing algorithms
 - □ Proactive routing algorithms
 - Distance Vector routing, RIP, DSDV
 - □ On-demand routing algorithms
 - AODV, DSR
- A combination of proactive and on-demand routing algorithms
 - □ Adaptive Distance Vector (ADV)

Introduction

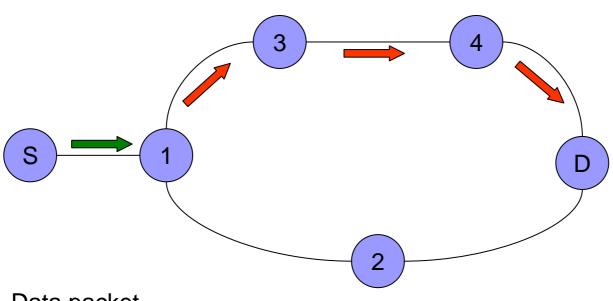
- This paper compare ADV with two ondemand protocols
 - □ Ad hoc On-demand Distance Vector (AODV)
 - □ Dynamic Source Routing (DSR)
- And a Distance Vector protocols
 - Destination-Sequenced Distance Vector (DSDV)

Current routing algorithms for MANETs

- Distance Vector Routing
 - The Routing Information Protocol (RIP) is an example of distance vector (DV) routing
 - Counting-to-infinity problem
 - □ Destination-Sequenced Distance Vector (DSDV)
 - DSDV solves counting-to-infinity problem in DV routing by attaching sequence numbers to routing entries.

Current routing algorithms for MANETs

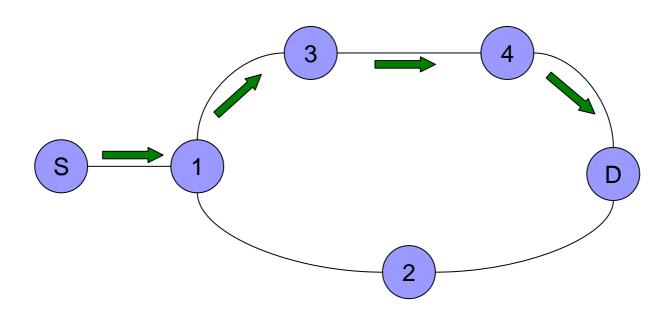
- On-demand Routing Algorithms
 - □ Ad hoc On-demand Distance Vector –AODV





Current routing algorithms for MANETs

□ Dynamic Source Routing - DSR



- The Adaptive Distance Vector (ADV) starts with DSDV
- Using the following methods to reduce the routing overhead
 - □ Varying the number of active routes maintained
 - □ Varying the frequency of routing updates

- Varying the number of active routes maintained
 - ADV advertises and maintains routes for active receivers only
 - □ A node is an active receiver if it is the receiver of any currently active connection.

- ☐ At the beginning of a new connection, the source floods with an *init-connection* control packet
- □ If destination node is not an active receiver already, broadcasting with a receiver alert packet.
- When a connection is to be closed, the source broadcasts network-wide an end-connection control packet
- If the destination node has no additional active connections, then it broadcasts a non-receiver alert control packet

- Varying the frequency of routing updates
 - A node should trigger an update under three conditions
 - if it has some buffered data packets due to lack of routes
 - if one or more of its neighbors make a request for fresh routes
 - it is a forwarding node that intends to advertise any fresh valid/invalid route to the destination so as to keep the route fresh.

- ☐ Then check trigger meter
 - If > TRGMETER_FULL, send full update
 - Else if > trigger threshold, send partial update
- If the number of neighbor changes exceeds a preset number, then the node categorizes the network as HIGH_SPEED or else as LOW_SPEED network.

- Sending routing updates
 - With every routing update entry, a node sends an expected response value of ZERO, LOW, MEDIUM or HIGH
 - □ The expected response values are determined using the following rules

- An expected response of HIGH is given when there are packets waiting for this route in the node buffers regardless of the speed of the network.
- In a HIGH SPEED network, an expected response of MEDIUM is given if this node is a forwarding node for the destination given in the routing entry.
- In a LOW SPEED network, an expected response of LOW is given if this node is a forwarding node for any of its neighbors to the routing entry's destination.
- If none of the above criteria apply then the expected response is set to ZERO.

- Processing received updates
- the trigger meter is incremented by TRGMETER_HIGH, TRGMETER_MED or TRGMETER_LOW for an expected response of HIGH, MEDIUM or LOW respectively.

- To mitigate the effect of the periodic transmission of the updates, we have proposed some adaptive criteria that trigger routing updates based only on network load and mobility conditions.
- The overhead is reduced by varying the size and the frequency of the routing updates dynamically.
- We have shown using simulations that ADV outperforms on-demand protocols like AODV and DSR in many instances.

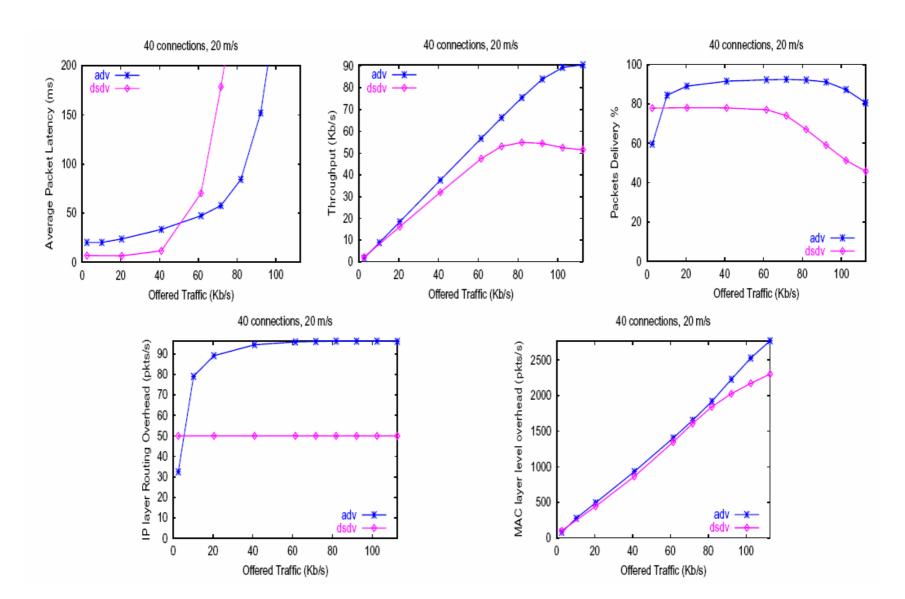
- ns-2 simulator with the CMU extensions
- CMU extensions include detailed implementations of IEEE 802.11 wireless LAN and ad hoc routing protocols DSDV, AODV, and DSR

- Mobility models
 - 50 nodes randomly placed on a 1000m x
 1000m field
 - 2. 100 nodes placed on a 2200m x 600m rectangular field
 - The nodes randomly choose a direction and speed and travel for a certain distance

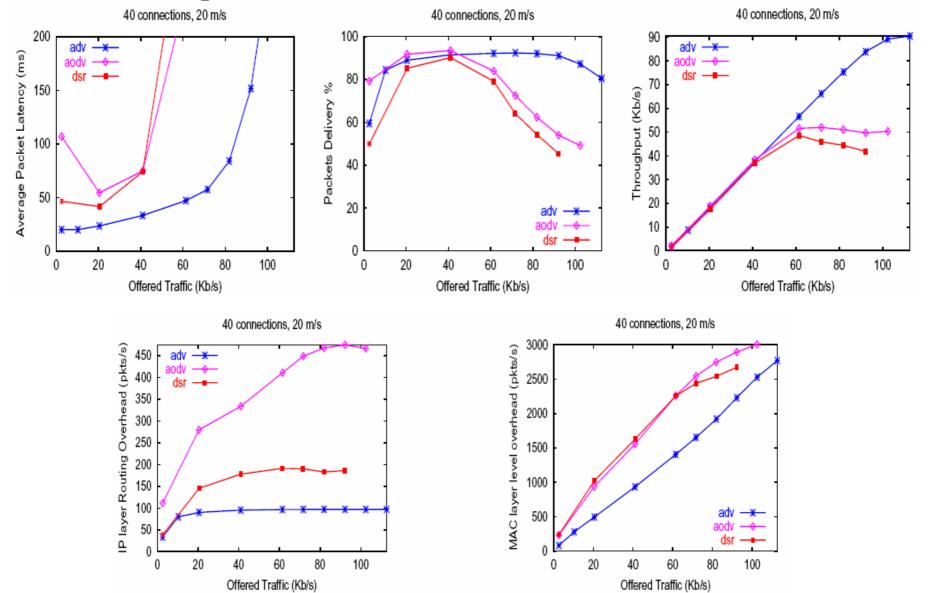
- Traffic load
 - ☐ CBR of 40 connections
- Performance metrics
 - □ average data packet latency
 - network throughput
 - percentage of data packets delivered
 - □ overheads of routing algorithms
 - IP layer
 - MAC layer

- ADV vs. DSDV
- ADV vs. On-demand algorithms
 - ☐ Steady-state behavior
 - □ Transient state behavior
 - 10 new connections are initiated every 60 seconds

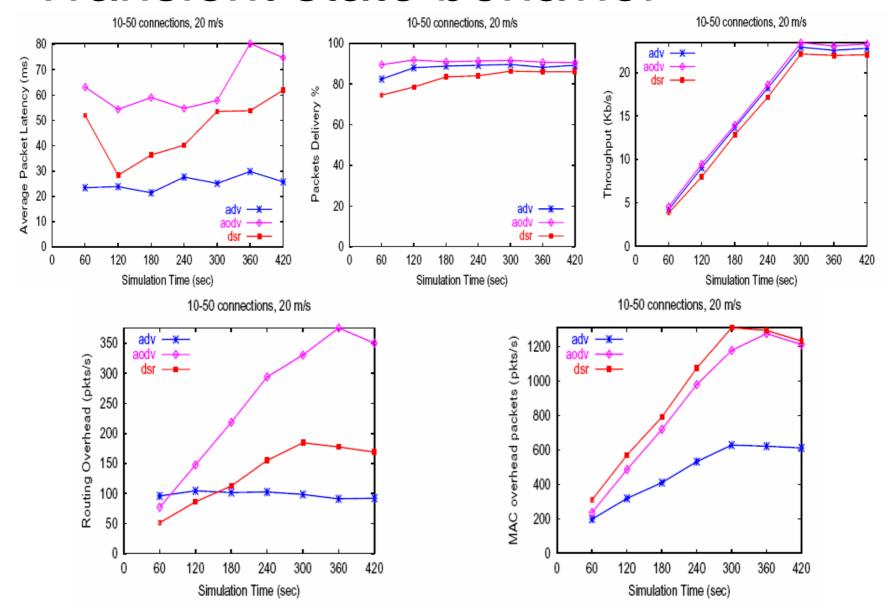
ADV vs. DSDV



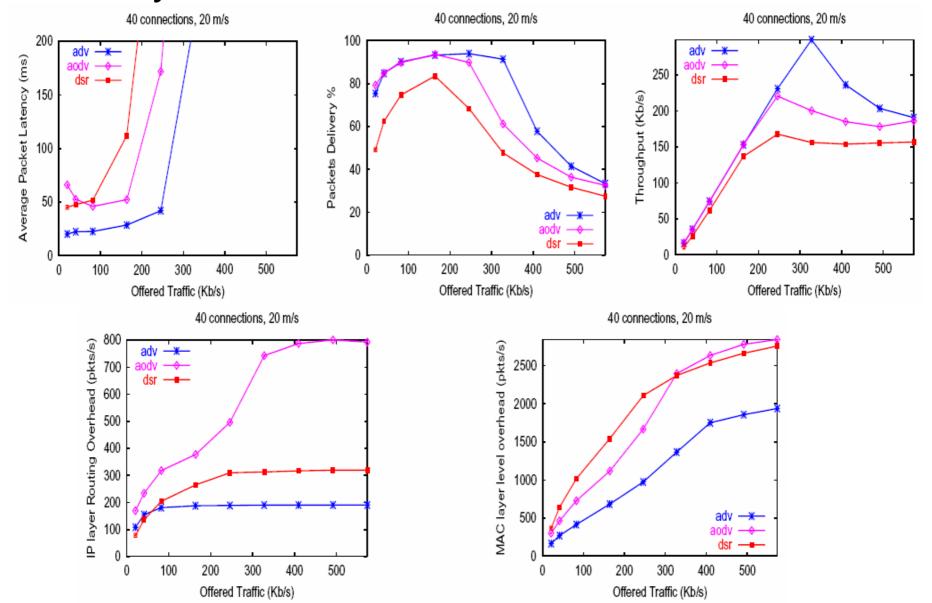
Steady-state behavior



Transient state behavior



Steady-state behavior with 100 node



Conclusions