AN ADAPTIVE STOCHASTIC GALERKIN METHOD FOR RANDOM ELLIPTIC OPERATORS

CLAUDE JEFFREY GITTELSON

ABSTRACT. We derive an adaptive solver for random elliptic boundary value problems, using techniques from adaptive wavelet methods. Substituting wavelets by polynomials of the random parameters leads to a modular solver for the parameter dependence of the random solution, which combines with any discretization on the spatial domain. In addition to selecting active polynomial modes, this solver can adaptively construct a separate spatial discretization for each of their coefficients. We show convergence of the solver in this general setting, along with a computable bound for the mean square error, and an optimality property in the case of a single spatial discretization. Numerical computations demonstrate convergence of the solver and compare it to a sparse tensor product construction.

INTRODUCTION

Stochastic Galerkin methods have emerged in the past decade as an efficient solution procedure for boundary value problems depending on random data; see [14, 32, 2, 30, 23, 18, 31, 28, 6, 5]. These methods approximate the random solution by a Galerkin projection onto a finite-dimensional space of random fields. This requires the solution of a single coupled system of deterministic equations for the coefficients of the Galerkin projection with respect to a predefined set of basis functions on the parameter domain.

A major remaining obstacle is the construction of suitable spaces in which to compute approximate solutions. These should be adapted to the stochastic structure of the equation. Simple tensor product constructions are infeasible due to the high dimensionality of the parameter domain in the case of input random fields with low regularity.

Parallel to but independently from the development of stochastic Galerkin methods, a new class of adaptive methods has emerged, which are set not in the continuous framework of a boundary value problem, but rather on the level of coefficients with respect to a hierarchic Riesz basis, such as a wavelet basis. Due to the norm equivalences constitutive of Riesz bases, errors and residuals in appropriate sequence spaces are equivalent to those in physically meaningful function spaces. This permits adaptive wavelet methods to be applied directly to a large class of equations, provided that a suitable Riesz basis is available.

©2013 American Mathematical Society

Reverts to public domain 28 years from publication

Received by the editor March 2, 2011 and, in revised form, September 24, 2011.

²⁰¹⁰ Mathematics Subject Classification. Primary 35R60, 47B80, 60H25, 65C20, 65N12, 65N22, 65N30, 65J10, 65Y20.

This research was supported in part by the Swiss National Science Foundation grant No. 200021-120290/1.

For symmetric elliptic problems, the error of the Galerkin projection onto the span of a set of coefficients can be estimated using a sufficiently accurate approximation of the residual of a previously computed approximate solution; see [8, 19, 16]. This results in a sequence of finite-dimensional linear equations with successively larger sets of active coefficients.

We use techniques from these adaptive wavelet methods to derive an adaptive solver for random symmetric elliptic boundary value problems. In place of wavelets, we use an orthonormal polynomial basis on the parameter domain. The coefficients of the random solution with respect to this basis are deterministic functions on the spatial domain.

Adaptive wavelet methods extend to this vector setting, and lead to a modular solver which can be coupled with any discretization of or solver for the deterministic problem. We consider adaptive finite elements with a residual-based a posteriori error estimator.

We review random operator equations in Section 1. In particular, we derive the weak formulation of such equations, construct orthonormal polynomials on the parameter domain, and recast the weak formulation as a bi-infinite operator matrix equation for the coefficients of the random solution with respect to this polynomial basis. We refer to [22] for further details.

A crucial ingredient in adaptive wavelet methods is the approximation of the residual. We study this for the setting of stochastic operator equations in Section 2. The resulting adaptive solver is presented in Section 3. We show convergence of the method, and provide a reliable error bound. Optimality properties are discussed in Section 4 for the special case of a fixed spatial discretization.

Finally, in Section 5, we apply the method to a simple elliptic equation. We discuss a suitable a posteriori finite element error estimator, and present numerical computations. These demonstrate the convergence of our solver and compare the adaptively constructed discretizations with the a priori adapted sparse tensor product construction from [5]; we refer to [21] for a comparison with other adaptive solvers. We discuss the empirical convergence behavior in the light of the theoretical approximation results in [11, 10].

1. Stochastic operator equations

1.1. **Pointwise definition.** Let $\mathbb{K} \in \{\mathbb{R}, \mathbb{C}\}$ and let V be a separable Hilbert space over \mathbb{K} . We denote by V^* the space of all continuous antilinear functionals on V. Furthermore, $\mathcal{L}(V, V^*)$ is the Banach space of bounded linear maps from V to V^* .

We consider operator equations depending on a parameter in $\Gamma := [-1, 1]^{\infty}$. Given

(1.1)
$$A: \Gamma \to \mathcal{L}(V, V^*)$$
 and $f: \Gamma \to V^*$,

we wish to determine

(1.2)
$$u: \Gamma \to V, \qquad A(y)u(y) = f(y) \quad \forall y \in \Gamma$$

Let $D \in \mathcal{L}(V, V^*)$ be the Riesz isomorphism, i.e., $\langle D \cdot, \cdot \rangle$ is the scalar product in V. We decompose A as

(1.3)
$$A(y) = D + R(y) \quad \forall y \in \Gamma$$

and assume that R(y) is linear in $y \in \Gamma$,

(1.4)
$$R(y) = \sum_{m=1}^{\infty} y_m R_m \qquad \forall y = (y_m)_{m=1}^{\infty} \in \Gamma;$$

e.g., as in [5, 6, 11, 10, 28]. Here, each R_m is in $\mathcal{L}(V, V^*)$. We assume $(R_m)_m \in \ell^1(\mathbb{N}; \mathcal{L}(V, V^*))$, and there is a $\gamma \in [0, 1)$ such that $||R(y)||_{V \to V^*} \leq \gamma$ for all $y \in \Gamma$. By [22, Proposition 1.2], this ensures existence and uniqueness of the solution of (1.1). For simplicity, we also assume that the sequence $(||R_m||_{V \to V^*})_{m=1}^{\infty}$ is nonincreasing.

1.2. Weak formulation. Let π be a probability measure on the parameter domain Γ with Borel σ -algebra $\mathcal{B}(\Gamma)$. We assume that the map $\Gamma \ni y \mapsto A(y)v(y)$ is measurable for any measurable $v \colon \Gamma \to V$. Then

(1.5)
$$\mathcal{A} \colon L^2_{\pi}(\Gamma; V) \to L^2_{\pi}(\Gamma; V^*) \ , \quad v \mapsto [y \mapsto A(y)v(y)] \ ,$$

is well defined and continuous. We assume also that $f \in L^2_{\pi}(\Gamma; V^*)$. The weak formulation of (1.2) is to find $u \in L^2_{\pi}(\Gamma; V)$ such that

(1.6)
$$\int_{\Gamma} \langle A(y)u(y), v(y) \rangle \, \mathrm{d}\pi(y) = \int_{\Gamma} \langle f(y), v(y) \rangle \, \mathrm{d}\pi(y) \qquad \forall v \in L^{2}_{\pi}(\Gamma; V)$$

The left term in (1.6) is the duality pairing in $L^2_{\pi}(\Gamma; V)$ of $\mathcal{A}u$ with the test function v, and the right term is the duality pairing of f with v. We follow the convention that the duality pairing is linear in the first argument and antilinear in the second.

By [22, Theorem 1.4], the solution u of (1.2) is in $L^2_{\pi}(\Gamma; V)$, and it is the unique solution of (1.6). In particular, the operator \mathcal{A} is boundedly invertible.

We define the multiplication operators

(1.7)
$$K_m \colon L^2_{\pi}(\Gamma) \to L^2_{\pi}(\Gamma) , \quad v(y) \mapsto y_m v(y) , \quad m \in \mathbb{N} .$$

Since y_m is real and $|y_m|$ is less than one, K_m is symmetric and has norm at most one.

By separability of V, the Lebesgue–Bochner space $L^2_{\pi}(\Gamma; V)$ is isometrically isomorphic to the Hilbert tensor product $L^2_{\pi}(\Gamma) \otimes V$, and similarly for V^* in place of V. Using these identifications, we expand \mathcal{A} as $\mathcal{A} = \mathcal{D} + \mathcal{R}$ with

(1.8)
$$\mathcal{D} \coloneqq \operatorname{id}_{L^2_{\pi}(\Gamma)} \otimes D$$
 and $\mathcal{R} \coloneqq \sum_{m=1}^{\infty} K_m \otimes R_m$.

This sum converges in $\mathcal{L}(L^2_{\pi}(\Gamma; V), L^2_{\pi}(\Gamma; V^*))$ due to the assumption that $(R_m)_m \in \ell^1(\mathbb{N}; \mathcal{L}(V, V^*))$.

Lemma 1.1. $\|\mathcal{R}\|_{L^{2}_{\pi}(\Gamma;V) \to L^{2}_{\pi}(\Gamma;V^{*})} \leq \gamma < 1.$

Proof. We note that, as in (1.5), $(\mathcal{R}v)(y) = R(y)v(y)$ for all $v \in L^2_{\pi}(\Gamma; V)$ and $y \in \Gamma$. Therefore, using the assumption $||R(y)||_{V \to V^*} \leq \gamma$,

$$\left\|\mathcal{R}v\right\|_{L^{2}_{\pi}(\Gamma;V^{*})}^{2} = \int_{\Gamma} \left\|R(y)v(y)\right\|_{V^{*}}^{2} d\pi(y) \leq \int_{\Gamma} \left\|R(y)\right\|_{V \to V^{*}}^{2} \left\|v(y)\right\|_{V}^{2} d\pi(y) ,$$

and the assertion follows using the assumption $||R(y)||_{V \to V^*} \leq \gamma$.

1.3. Orthonormal polynomial basis. In order to construct an orthonormal polynomial basis of $L^2_{\pi}(\Gamma)$, we assume that π is a product measure. Let

(1.9)
$$\pi = \bigotimes_{m=1}^{\infty} \pi_m$$

for probability measures π_m on $([-1,1], \mathcal{B}([-1,1]))$; see e.g. [4, Section 9] for a general construction of arbitrary products of probability measures. We assume that the support of π_m in [-1,1] has infinite cardinality.

For all $m \in \mathbb{N}$, let $(P_n^m)_{n=0}^{\infty}$ be an orthonormal polynomial basis of $L^2_{\pi_m}([-1,1])$, with deg $P_n^m = n$. Such a basis is given by the three term recursion $P_{-1}^m \coloneqq 0$, $P_0^m \coloneqq 1$ and

(1.10)
$$\beta_n^m P_n^m(\xi) \coloneqq (\xi - \alpha_{n-1}^m) P_{n-1}^m(\xi) - \beta_{n-1}^m P_{n-2}^m(\xi) , \quad n \in \mathbb{N} ,$$

with

(1.11)
$$\alpha_n^m \coloneqq \int_{-1}^1 \xi P_n^m(\xi)^2 \,\mathrm{d}\pi_m(\xi) \quad \text{and} \quad \beta_n^m \coloneqq \frac{c_{n-1}^m}{c_n^m} \,,$$

where c_n^m is the leading coefficient of P_n^m , $\beta_0^m \coloneqq 1$, and P_n^m is chosen as normalized in $L^2_{\pi_m}([0,1])$ with a positive leading coefficient.

We define the set of finitely supported sequences in \mathbb{N}_0 as

(1.12)
$$\Lambda \coloneqq \left\{ \nu \in \mathbb{N}_0^{\mathbb{N}} \; ; \; \# \operatorname{supp} \nu < \infty \right\} \; ,$$

where the support is defined by

(1.13)
$$\operatorname{supp} \nu \coloneqq \{ m \in \mathbb{N} ; \nu_m \neq 0 \} , \quad \nu \in \mathbb{N}_0^{\mathbb{N}} .$$

Then countably infinite tensor product polynomials are given by

(1.14)
$$\boldsymbol{P} \coloneqq (P_{\nu})_{\nu \in \Lambda} , \quad P_{\nu} \coloneqq \bigotimes_{m=1}^{\infty} P_{\nu_m}^m , \quad \nu \in \Lambda$$

Note that each of these functions depends on only finitely many dimensions,

(1.15)
$$P_{\nu}(y) = \prod_{m=1}^{\infty} P_{\nu_m}^m(y_m) = \prod_{m \in \text{supp } \nu} P_{\nu_m}^m(y_m) , \quad \nu \in \Lambda ,$$

since $P_0^m = 1$ for all $m \in \mathbb{N}$.

For example, by [22, Theorem 2.8], \boldsymbol{P} is an orthonormal basis of $L^2_{\pi}(\Gamma)$. By Parseval's identity, this is equivalent to the statement that the map

(1.16)
$$T: \ell^2(\Lambda) \to L^2_{\pi}(\Gamma) \ , \quad (c_{\nu})_{\nu \in \Lambda} \mapsto \sum_{\nu \in \Lambda} c_{\nu} P_{\nu} \ ,$$

is a unitary isomorphism. The inverse of T is

(1.17)
$$T^{-1} = T^* \colon L^2_{\pi}(\Gamma) \to \ell^2(\Lambda) \ , \quad g \mapsto \left(\int_{\Gamma} g(y) \overline{P_{\nu}(y)} \, \mathrm{d}\pi(y) \right)_{\nu \in \Lambda}$$

1.4. **Bi-infinite operator matrix equation.** We use the isomorphism T from (1.16) to recast the weak stochastic operator equation (1.6) as an equivalent discrete operator equation. Since T is a unitary map from $\ell^2(\Lambda)$ to $L^2_{\pi}(\Gamma)$, the tensor product operator $T_V := T \otimes id_V$ is an isometric isomorphism from $\ell^2(\Lambda; V)$ to $L^2_{\pi}(\Gamma; V)$. By definition, $w \in L^2_{\pi}(\Gamma; V)$ and $w = (w_{\nu})_{\nu \in \Lambda} \in \ell^2(\Lambda; V)$ are related by $w = T_V w$ if

(1.18)
$$w(y) = \sum_{\nu \in \Lambda} w_{\nu} P_{\nu}(y)$$
 or $w_{\nu} = \int_{\Gamma} w(y) \overline{P_{\nu}(y)} d\pi(y) \quad \forall \nu \in \Lambda$,

and either of these properties implies the other. The series in (1.18) converges unconditionally in $L^2_{\pi}(\Gamma; V)$, and the integral can be interpreted as a Bochner integral in V.

Let $\boldsymbol{A} \coloneqq T_V^* \mathcal{A} T_V$ and $\boldsymbol{f} \coloneqq T_V^* f$. Then $u = T_V \boldsymbol{u}$ for $\boldsymbol{u} \in \ell^2(\Lambda; V)$ with

$$(1.19) Au = f$$

since $u \in L^2_{\pi}(\Gamma; V)$ satisfies $\mathcal{A}u = f$.

By definition, A is a boundedly invertible linear map from $\ell^2(\Lambda; V)$ to $\ell^2(\Lambda; V^*)$. It can be interpreted as a bi-infinite operator matrix

(1.20)
$$\boldsymbol{A} = [A_{\nu\mu}]_{\nu,\mu\in\Lambda}, \quad A_{\nu\mu} \colon V \to V^* ,$$

with entries

(1.21)
$$A_{\nu\nu} = D + \sum_{m=1}^{\infty} \alpha_{\nu_m}^m R_m , \quad \nu \in \Lambda ,$$
$$A_{\nu\mu} = \beta_{\max(\nu_m,\mu_m)}^m R_m , \quad \nu,\mu \in \Lambda , \quad \nu - \mu = \pm \epsilon_m$$

and $A_{\nu\mu} = 0$ otherwise, where ϵ_m denotes the Kronecker sequence with $(\epsilon_m)_n = \delta_{mn}$. If π_m is a symmetric measure on [-1, 1] for all $m \in \mathbb{N}$, then $\alpha_n^m = 0$ for all m and n, and thus $A_{\nu\nu} = D$. We refer to [22, 20] for details.

Similarly, the operator $\mathbf{R} \coloneqq T_V^* \mathcal{R} T_V$ can be interpreted as a bi-infinite operator matrix $\mathbf{R} = [R_{\nu\mu}]$ with $R_{\nu\nu} = A_{\nu\nu} - D$ and $R_{\nu\mu} = A_{\nu\mu}$ for $\nu \neq \mu$.

Let $\mathbf{K}_m = T^* K_m T \in \mathcal{L}(\ell^2(\Lambda))$. Due to the three term recursion (1.10),

(1.22)
$$(\boldsymbol{K}_m \boldsymbol{c})_{\mu} = \beta_{\mu_m+1}^m c_{\mu+\epsilon_m} + \alpha_{\mu_m}^m c_{\mu} + \beta_{\mu_m}^m c_{\mu-\epsilon_m} , \quad \mu \in \Lambda ,$$

for $\boldsymbol{c} = (c_{\mu})_{\mu \in \Lambda} \in \ell^2(\Lambda)$, where $c_{\mu} := 0$ if $\mu_m < 0$ for any $m \in \mathbb{N}$. Furthermore, $\boldsymbol{K}_m^* = \boldsymbol{K}_m$ and $\|\boldsymbol{K}_m\|_{\ell^2(\Lambda) \to \ell^2(\Lambda)} \leq 1$.

Using the maps $\boldsymbol{K}_m, \, \boldsymbol{R}$ can be written succinctly as

(1.23)
$$\boldsymbol{R} = \sum_{m=1}^{\infty} \boldsymbol{K}_m \otimes \boldsymbol{R}_m$$

with unconditional convergence in $\mathcal{L}(\ell^2(\Lambda; V), \ell^2(\Lambda; V^*))$. By Lemma 1.1,

(1.24)
$$\|\boldsymbol{R}\|_{\ell^2(\Lambda;V)\to\ell^2(\Lambda;V^*)} \le \gamma < 1$$

In particular, $\|\boldsymbol{A}\| \leq (1+\gamma)$ and $\|\boldsymbol{A}^{-1}\| \leq (1-\gamma)^{-1}$.

We also define the operator $\mathbf{D} := T_V^* \mathcal{D} T_V$. This is just the Riesz isomorphism from $\ell^2(\Lambda; V)$ to $\ell^2(\Lambda; V^*)$. By [22, Proposition 2.10],

(1.25)
$$(1-\gamma)\mathbf{D} \le \mathbf{A} \le (1+\gamma)\mathbf{D}$$
 and $\frac{1}{1+\gamma}\mathbf{D}^{-1} \le \mathbf{A}^{-1} \le \frac{1}{1-\gamma}\mathbf{D}^{-1}$.

In particular, using $\mathbf{A} = \mathbf{A}\mathbf{A}^{-1}\mathbf{A}$, we have

(1.26)
$$\frac{1}{1+\gamma} \boldsymbol{A} \boldsymbol{D}^{-1} \boldsymbol{A} \le \boldsymbol{A} \le \frac{1}{1-\gamma} \boldsymbol{A} \boldsymbol{D}^{-1} \boldsymbol{A} .$$

1.5. Galerkin projection. Let \mathcal{W} be a closed subspace of $L^2_{\pi}(\Gamma; V)$. The Galerkin solution $\bar{u} \in \mathcal{W}$ is defined through the linear variational problem

(1.27)
$$\int_{\Gamma} \langle A(y)\bar{u}(y), w(y) \rangle \, \mathrm{d}\pi(y) = \int_{\Gamma} \langle f(y), w(y) \rangle \, \mathrm{d}\pi(y) \qquad \forall w \in \mathcal{W} \, .$$

Existence, uniqueness and quasi-optimality of \bar{u} follow since \mathcal{A} induces an inner product on $L^2_{\pi}(\Gamma; V)$ that is equivalent to the standard inner product; see [22, Proposition 1.5].

For all $\nu \in \Lambda$, let W_{ν} be a finite dimensional subspace of V, such that $W_{\nu} \neq \{0\}$ for only finitely many $\nu \in \Lambda$. It is particularly useful to consider spaces \mathcal{W} of the form

(1.28)
$$\mathcal{W} \coloneqq \sum_{\nu \in \Lambda} W_{\nu} P_{\nu} \; .$$

The Galerkin operator on such a space has a similar structure to (1.20), with $A_{\nu\mu}$ replaced by its representation on suitable subpaces W_{ν} of V; see [22, Section 2].

2. Approximation of the residual

2.1. Adaptive application of the stochastic operator. We construct a sequence of approximations of \mathbf{R} by truncating the series (1.23). For all $M \in \mathbb{N}$, let

(2.1)
$$\boldsymbol{R}_{[M]} \coloneqq \sum_{m=1}^{M} \boldsymbol{K}_m \otimes R_m ,$$

and $\mathbf{R}_{[0]} \coloneqq 0$. For all $M \in \mathbb{N}$, let $\bar{e}_{\mathbf{R},M}$ be given such that

(2.2)
$$\left\|\boldsymbol{R} - \boldsymbol{R}_{[M]}\right\|_{\ell^{2}(\Lambda; V) \to \ell^{2}(\Lambda; V^{*})} \leq \bar{e}_{\boldsymbol{R}, M}.$$

For example, these bounds can be chosen as

(2.3)
$$\bar{e}_{\boldsymbol{R},M} \coloneqq \sum_{m=M+1}^{\infty} \|R_m\|_{V \to V^*} .$$

We assume that $(\bar{e}_{\boldsymbol{R},M})_{M=0}^{\infty}$ is nonincreasing and converges to 0, and also that the sequence of differences $(\bar{e}_{\boldsymbol{R},M} - \bar{e}_{\boldsymbol{R},M+1})_{M=0}^{\infty}$ is nonincreasing.

We consider a partitioning of a vector $\boldsymbol{w} \in \ell^2(\Lambda)$ into $\boldsymbol{w}_{[p]} \coloneqq \boldsymbol{w}|_{\Lambda_p}, p = 1, \ldots, P$, for disjoint index sets $\Lambda_p \subset \Lambda$. This can be approximate in that $\boldsymbol{w}_{[1]} + \cdots + \boldsymbol{w}_{[P]}$ only approximates \boldsymbol{w} in $\ell^2(\Lambda)$. We think of $\boldsymbol{w}_{[1]}$ as containing the largest elements of $\boldsymbol{w}, \boldsymbol{w}_{[2]}$ the next largest, and so on.

Such a partitioning can be constructed by the approximate sorting algorithm

(2.4)
$$\operatorname{BucketSort}[\boldsymbol{w}, \epsilon] \mapsto \left[(\boldsymbol{w}_{[p]})_{p=1}^{P}, (\Lambda_{p})_{p=1}^{P} \right]$$

which, given a finitely supported $\boldsymbol{w} \in \ell^2(\Lambda)$ and a threshold $\epsilon > 0$, returns index sets

(2.5)
$$\Lambda_p \coloneqq \left\{ \mu \in \Lambda \, ; \, |w_{\mu}| \in (2^{-p/2} \, \|\boldsymbol{w}\|_{\ell^{\infty}} \, , 2^{-(p-1)/2} \, \|\boldsymbol{w}\|_{\ell^{\infty}}] \right\}$$

and $\boldsymbol{w}_{[p]} \coloneqq \boldsymbol{w}|_{\Lambda_p}$; see [24, 3, 19, 16]. The integer P is minimal with

(2.6)
$$2^{-P/2} \|\boldsymbol{w}\|_{\ell^{\infty}(\Lambda)} \sqrt{\#\operatorname{supp} \boldsymbol{w}} \leq \epsilon \;.$$

By [19, Rem. 2.3] or [16, Prop. 4.4], the number of operations and storage locations required by a call of BucketSort[w, ϵ] is bounded by

(2.7)
$$\# \operatorname{supp} \boldsymbol{w} + \max(1, \lceil \log(\|\boldsymbol{w}\|_{\ell^{\infty}(\Lambda)} \sqrt{\# \operatorname{supp} \boldsymbol{w}}/\epsilon) \rceil) .$$

This analysis uses that every w_{μ} , $\mu \in \Lambda$, can be mapped to p with $\mu \in \Lambda_p$ in constant time by evaluating

(2.8)
$$p \coloneqq \left[1 + 2 \log_2 \left(\frac{\|\boldsymbol{w}\|_{\ell^{\infty}(\Lambda)}}{|w_{\mu}|} \right) \right] .$$

Alternatively, any standard comparison-based sorting algorithm can be used to construct the partitioning of w, albeit with an additional logarithmic factor in the complexity.

 $\mathtt{Apply}_{oldsymbol{R}}[oldsymbol{v},\epsilon]\mapstooldsymbol{z}$

$$\begin{split} \left[\cdot, (\Lambda_p)_{p=1}^P \right] & \longleftarrow \operatorname{BucketSort} \left[(\|v_{\mu}\|_V)_{\mu \in \Lambda}, \frac{\epsilon}{2\bar{e}_{\mathbf{R},0}} \right] \\ \text{for } p = 1, \dots, P \text{ do } v_{[p]} \leftarrow (v_{\mu})_{\mu \in \Lambda_p} \\ \operatorname{Compute the minimal } \ell \in \{0, 1, \dots, P\} \text{ s.t. } \delta \coloneqq \bar{e}_{\mathbf{R},0} \left\| \boldsymbol{v} - \sum_{p=1}^{\ell} \boldsymbol{v}_{[p]} \right\|_{\ell^2(\Lambda;V)} \leq \frac{\epsilon}{2} \\ \text{for } p = 1, \dots, \ell \text{ do } M_p \leftarrow 0 \\ \text{while } \sum_{p=1}^{\ell} \bar{e}_{\mathbf{R},M_p} \| \boldsymbol{v}_{[p]} \|_{\ell^2(\Lambda;V)} > \epsilon - \delta \text{ do} \\ \left[\begin{array}{c} q \leftarrow \operatorname{argmax}_{p=1,\dots,\ell}(\bar{e}_{\mathbf{R},M_p} - \bar{e}_{\mathbf{R},M_p+1}) \| \boldsymbol{v}_{[p]} \|_{\ell^2(\Lambda;V)} / \# \Lambda_p \\ M_q \leftarrow M_q + 1 \end{array} \right] \\ \boldsymbol{z} = (z_{\nu})_{\nu \in \Lambda} \leftarrow \mathbf{0} \\ \text{for all the } \mu \in \Lambda_p \text{ do} \\ \text{for all the } \mu \in \Lambda_p \text{ do} \\ \left[\begin{array}{c} \text{for } m = 1, \dots, M_p \text{ do} \\ w \leftarrow R_m v_{\mu} \\ z_{\mu + \epsilon_m} \leftarrow z_{\mu + \epsilon_m} + \beta_{\mu_m + 1}^m w \\ \text{ if } \mu_m \ge 1 \text{ then } z_{\mu - \epsilon_m} \leftarrow z_{\mu} - \epsilon_m + \beta_{\mu_m}^m w \\ \text{ if } \alpha_{\mu_m}^m \neq 0 \text{ then } z_{\mu} \leftarrow z_{\mu} + \alpha_{\mu_m}^m w \end{split} \right] \\ \end{array} \right]$$

The routine $\operatorname{Apply}_{R}[v, \epsilon]$ adaptively approximates Rv in three distinct steps. First, the elements of v are grouped according to their norm. Elements smaller than a certain tolerance are discarded. This truncation of the vector v produces an error of at most $\delta \leq \epsilon/2$.

Next, a greedy algorithm is used to assign to each segment $\boldsymbol{v}_{[p]}$ of \boldsymbol{v} an approximation $\boldsymbol{R}_{[M_p]}$ of \boldsymbol{R} . Starting with $\boldsymbol{R}_{[M_p]} = \boldsymbol{0}$ for all $p = 1, \ldots, \ell$, these approximations are refined iteratively until an estimate of the error is smaller than $\epsilon - \delta$.

Finally, the operations determined by the previous two steps are performed. Each multiplication $R_m v_{\mu}$ is performed just once, and copied to the appropriate entries of z.

Proposition 2.1. For any finitely supported $v \in \ell^2(\Lambda; V)$ and any $\epsilon > 0$, $\text{Apply}_{\mathbf{R}}[v, \epsilon]$ produces a finitely supported $z \in \ell^2(\Lambda; V^*)$ with

(2.9)
$$\# \operatorname{supp} \boldsymbol{z} \le 3 \sum_{p=1}^{\ell} M_p \# \Lambda_p$$

and

(2.10)
$$\|R\boldsymbol{v} - \boldsymbol{z}\|_{\ell^{2}(\Lambda; V^{*})} \leq \delta + \eta_{\boldsymbol{M}} \leq \epsilon , \quad \eta_{\boldsymbol{M}} \coloneqq \sum_{p=1}^{\ell} \bar{e}_{\boldsymbol{R}, M_{p}} \|\boldsymbol{v}_{[p]}\|_{\ell^{2}(\Lambda; V)} ,$$

where M_p refers to the final value of this variable in the call of Apply_R . The total number of products $R_m v_\mu$ computed in $\text{Apply}_R[v, \epsilon]$ is $\sigma_M := \sum_{p=1}^{\ell} M_p \# A_p$. Furthermore, the vector $\mathbf{M} = (M_p)_{p=1}^{\ell}$ is optimal in the sense that if $\mathbf{N} = (N_p)_{p=1}^{\ell}$ with $\sigma_N \leq \sigma_M$, then $\eta_N \geq \eta_M$, and if $\eta_N \leq \eta_M$, then $\sigma_N \geq \sigma_M$.

Proof. The estimate (2.9) follows from the fact that each K_m has at most three nonzero entries per column; see (1.22). Since $\|\mathbf{R}\|_{\ell^2(A;V)\to\ell^2(A;V^*)} \leq \bar{e}_{\mathbf{R},0}$,

$$\left\| oldsymbol{R}oldsymbol{v} - oldsymbol{R}\sum_{p=1}^\ell oldsymbol{v}_{[p]}
ight\|_{\ell^2(\Lambda;V^*)} \leq ar{e}_{oldsymbol{R},0} \left\| oldsymbol{v} - \sum_{p=1}^\ell oldsymbol{v}_{[p]}
ight\|_{\ell^2(\Lambda;V)} = \delta \leq rac{\epsilon}{2} \; .$$

Due to (2.2) and the termination criterion in the greedy subroutine of Apply_R,

$$\sum_{p=1}^{\ell} \left\| \boldsymbol{R} \boldsymbol{v}_{[p]} - \boldsymbol{R}_{[M_p]} \boldsymbol{v}_{[p]} \right\|_{\ell^2(\Lambda; V^*)} \le \sum_{p=1}^{\ell} \bar{e}_{\boldsymbol{R}, M_p} \left\| \boldsymbol{v}_{[p]} \right\|_{\ell^2(\Lambda; V)} \le \epsilon - \delta.$$

For the optimality property of the greedy algorithm, we refer to the more general statement [20, Theorem 4.1.5]. $\hfill \Box$

2.2. Computation of the residual. We assume a solver for D is available such that for any $g \in V^*$ and any $\epsilon > 0$,

(2.11)
$$\operatorname{Solve}_{D}[g,\epsilon] \mapsto v , \quad \left\| v - D^{-1}g \right\|_{V} \le \epsilon$$

For example, $Solve_D$ could be an adaptive wavelet method (see e.g. [8, 9, 19]), an adaptive frame method (see e.g. [27, 12, 13]), or a finite element method with a posteriori error estimation; see e.g. [17, 25, 7].

Furthermore, we assume that a routine

is available to compute approximations $\tilde{f} = (\tilde{f}_{\nu})_{\nu \in \Lambda}$ of f with $\# \operatorname{supp} \tilde{f} < \infty$ and

(2.13)
$$\left\| \boldsymbol{f} - \tilde{\boldsymbol{f}} \right\|_{\ell^{2}(\Lambda; V^{*})} \leq \epsilon$$

for any $\epsilon > 0$.

The routine $\text{Residual}_{A,f}$ approximates the residual f - Av up to a prescribed relative tolerance.

 $\texttt{Residual}_{\boldsymbol{A},\boldsymbol{f}}[\epsilon,\boldsymbol{v},\eta_0,\boldsymbol{\chi},\boldsymbol{\omega},\boldsymbol{\alpha},\boldsymbol{\beta}]\mapsto [\boldsymbol{w},\eta,\boldsymbol{\zeta}]$

$$\begin{split} \zeta &\longleftarrow \chi \eta_0 \\ \textbf{repeat} \\ & \left[\begin{array}{c} \boldsymbol{g} = (g_{\nu})_{\nu \in \Lambda} &\longleftarrow \text{RHS}_{\boldsymbol{f}}[\beta(1-\alpha)\zeta] - \text{Apply}_{\boldsymbol{R}}[\boldsymbol{v}, (1-\beta)(1-\alpha)\zeta] \\ \boldsymbol{w} = (w_{\nu})_{\nu \in \Lambda} &\longleftarrow (\text{Solve}_D[g_{\nu}, \alpha\zeta(\# \operatorname{supp} \boldsymbol{g})^{-1/2}])_{\nu \in \Lambda} \\ \eta &\longleftarrow \|\boldsymbol{w} - \boldsymbol{v}\|_{\ell^2(\Lambda;V)} \\ \textbf{if } \zeta &\leq \omega\eta \text{ or } \eta + \zeta \leq \epsilon \text{ then break} \\ \zeta &\longleftarrow \omega \frac{1-\omega}{1+\omega}(\eta+\zeta) \end{split} \right] \end{split}$$

Proposition 2.2. For any finitely supported $\boldsymbol{v} = (v_{\nu})_{\nu \in \Lambda} \in \ell^2(\Lambda; V), \ \epsilon > 0, \ \eta_0 \geq 0, \ \chi > 0, \ \omega > 0, \ 0 < \alpha < 1 \ and \ 0 < \beta < 1, \ a \ call \ of \texttt{Residual}_{\boldsymbol{A},\boldsymbol{f}}[\epsilon, \boldsymbol{v}, \eta_0, \chi, \omega, \alpha, \beta]$ computes $\boldsymbol{w} \in \ell^2(\Lambda; V), \ \eta \geq 0 \ and \ \zeta \geq 0 \ with$ (2.14)

$$\left\| \eta - \| \boldsymbol{r} \|_{\ell^{2}(\Lambda; V^{*})} \right\| \leq \left\| \boldsymbol{w} - \boldsymbol{v} - \boldsymbol{D}^{-1} \boldsymbol{r} \right\|_{\ell^{2}(\Lambda; V)} = \left\| \boldsymbol{w} - \boldsymbol{D}^{-1} (\boldsymbol{f} - \boldsymbol{R} \boldsymbol{v}) \right\|_{\ell^{2}(\Lambda; V)} \leq \zeta,$$

where $\mathbf{r} = (r_{\nu})_{\nu \in \Lambda} \in \ell^2(\Lambda; V^*)$ is the residual $\mathbf{r} = \mathbf{f} - \mathbf{A}\mathbf{v}$, and ζ satisfies either $\zeta \leq \omega \eta$ or $\eta + \zeta \leq \epsilon$.

Proof. By construction,

$$\|\boldsymbol{g} - (\boldsymbol{f} - \boldsymbol{R}\boldsymbol{v})\|_{\ell^2(\Lambda; V^*)} \le (1 - \alpha)\zeta.$$

Furthermore, using $\|\boldsymbol{w} - \boldsymbol{D}^{-1}\boldsymbol{g}\|_{\ell^{2}(\Lambda;V)} \leq \alpha \zeta$,

$$\|\boldsymbol{w} - \boldsymbol{D}^{-1}(\boldsymbol{f} - \boldsymbol{R}\boldsymbol{v})\|_{\ell^{2}(\Lambda;V)} \leq \|\boldsymbol{w} - \boldsymbol{D}^{-1}\boldsymbol{g}\|_{\ell^{2}(\Lambda;V)} + \|\boldsymbol{g} - (\boldsymbol{f} - \boldsymbol{R}\boldsymbol{v})\|_{\ell^{2}(\Lambda;V^{*})} \leq \zeta.$$

The rest of (2.14) follows by triangle inequality with

$$\|\boldsymbol{r}\|_{\ell^{2}(\Lambda;V^{*})} = \|\boldsymbol{D}^{-1}\boldsymbol{r}\|_{\ell^{2}(\Lambda;V)}.$$

Remark 2.3. The tolerance ζ in **Residual**_{*A*,*f*} is initialized as the product of an initial estimate η_0 of the residual and a parameter χ . The update

(2.15)
$$\zeta \longleftarrow \omega \frac{1-\omega}{1+\omega} (\eta+\zeta) \eqqcolon \zeta$$

ensures a geometric decrease of ζ since if $\zeta > \omega \eta$, then

(2.16)
$$\zeta_1 = \omega \frac{1-\omega}{1+\omega} (\eta+\zeta) < \frac{1-\omega}{1+\omega} (\zeta+\omega\zeta) = (1-\omega)\zeta .$$

Therefore, the total computational cost of the routine is proportional to that of the final iteration of the loop. Furthermore, if $\zeta > \omega \eta$, then also

(2.17)
$$\zeta_1 = \omega \frac{1-\omega}{1+\omega} (\eta+\zeta) > \omega(1-\omega)\eta > \omega(\eta-\zeta) .$$

The term $\eta - \zeta$ in the last expression of (2.17) is a lower bound for the true residual $\|\mathbf{r}\|_{\ell^2(\Lambda; V_D^*)}$. In this sense, the prescription (2.15) does not select an unnecessarily small tolerance.

Finally, if $\zeta \leq 2\omega(1-\omega)^{-1}\eta$, then $\zeta_1 \leq \omega\eta$. If the next value of η is greater than or equal to the current value, this ensures that the termination criterion is met in the next iteration. For example, under the mild condition $\zeta \leq (1+4\omega-\omega^2)(1-\omega)^{-2}\eta$, we have $\zeta_1 \leq 2\omega(1-\omega)^{-1}\eta$. The loop can therefore be expected to terminate within three iterations.

Remark 2.4. In Residual_{A,f}, the tolerances of Solve_D are chosen such that the error tolerance $\alpha\zeta$ is equidistributed among all the nonzero indices of \boldsymbol{w} . This property is not required anywhere; Proposition 2.2 only uses that the total error in the computation of $\boldsymbol{D}^{-1}\boldsymbol{g}$ is no more than $\alpha\zeta$. Indeed, other strategies for selecting tolerances, e.g., based on additional a priori information, may be more efficient. Equidistributing the error among all the indices is a simple, practical starting point.

3. An adaptive solver

3.1. Refinement strategy. We use the approximation of the residual described in Section 2 to refine a Galerkin subspace $\mathcal{W} \subset L^2_{\pi}(\Gamma; V)$ of the form (1.28). For some approximate solution \boldsymbol{v} with $T_V \boldsymbol{v} \in \mathcal{W}$, let \boldsymbol{w} be the approximation of $\boldsymbol{D}^{-1}(\boldsymbol{f}-\boldsymbol{R}\boldsymbol{v})$ computed by Residual_{A,f}. We construct a space

(3.1)
$$\bar{\mathcal{W}} \coloneqq \sum_{\mu \in \Lambda} \bar{W}_{\mu} P_{\mu} \supset \mathcal{W} ,$$

with $\bar{W}_{\mu} \subset V$ finite-dimensional, such that \boldsymbol{w} can be approximated sufficiently in $\bar{\mathcal{W}}$. A simple choice is $\bar{W}_{\mu} \coloneqq W_{\mu} + \operatorname{span} w_{\mu}$, where $\mathcal{W} = \sum_{\mu} W_{\mu} P_{\mu}$.

We consider a multilevel setting. For each $\mu \in \operatorname{supp} \boldsymbol{w} \subset \Lambda$, let $W_{\mu} \cong W_{\mu}^{0} \subset W_{\mu}^{1} \subset \cdots$ be a scale of finite-dimensional subspaces of V such that $\bigcup_{i=0}^{\infty} W_{\mu}^{i}$ is dense in V. To each space, we associate a cost dim W_{μ}^{i} and an error $\|w_{\mu} - \Pi_{\mu}^{i}w_{\mu}\|_{V}^{2}$, where Π_{μ}^{i} denotes the orthogonal projection in V onto W_{μ}^{i} . In the construction of \overline{W} , we use a greedy algorithm to minimize the dimension of \overline{W} under a constraint on the approximation error of \boldsymbol{w} .

 $\texttt{Refine}_D[\mathcal{W}, \boldsymbol{w}, \epsilon] \mapsto [\bar{\mathcal{W}}, \bar{\boldsymbol{w}}, \varrho]$

for all the
$$\mu \in \operatorname{supp} \boldsymbol{w}$$
 do $j_{\mu} \leftarrow 0$
while $\sum_{\mu \in \operatorname{supp} \boldsymbol{w}} \left\| w_{\mu} - \Pi_{\mu}^{j_{\mu}} w_{\mu} \right\|_{V}^{2} > \epsilon^{2}$ do
 $\left\| \nu \leftarrow \operatorname{argmax}_{\mu \in \operatorname{supp} \boldsymbol{w}} \frac{\left\| \Pi_{\mu}^{j_{\mu}+1} w_{\mu} - \Pi_{\mu}^{j_{\mu}} w_{\mu} \right\|_{V}^{2}}{\dim(W_{\mu}^{j_{\mu}+1} \setminus W_{\mu}^{j_{\mu}})}$
for all the $\mu \in \operatorname{supp} \boldsymbol{w}$ do
 $\left\| \bar{W}_{\mu} \leftarrow W_{\mu}^{j_{\mu}} \\ \bar{w}_{\mu} \leftarrow \Pi_{\mu}^{j_{\mu}} w_{\mu} \right\|$
 $\varrho \leftarrow \left(\sum_{\mu \in \operatorname{supp} \boldsymbol{w}} \left\| w_{\mu} - \bar{w}_{\mu} \right\|_{V}^{2} \right)^{1/2}$

Proposition 3.1. If for every $\mu \in \operatorname{supp} \boldsymbol{w}$,

(3.2)
$$\frac{\left\|\Pi_{\mu}^{i+1}w_{\mu} - \Pi_{\mu}^{i}w_{\mu}\right\|_{V}^{2}}{\dim(W_{\mu}^{i+1} \setminus W_{\mu}^{i})} \geq \frac{\left\|\Pi_{\mu}^{j+1}w_{\mu} - \Pi_{\mu}^{j}w_{\mu}\right\|_{V}^{2}}{\dim(W_{\mu}^{j+1} \setminus W_{\mu}^{j})} \quad \forall i \leq j$$

then for any $\epsilon \geq 0$, a call of $\text{Refine}_D[W, w, \epsilon]$ constructs a space \overline{W} of the form (3.1) and $T_V \overline{w} \in \overline{W}$ satisfying

(3.3)
$$\varrho = \|\boldsymbol{w} - \bar{\boldsymbol{w}}\|_{\ell^2(\Lambda; V)} \le \epsilon$$

Furthermore, dim \overline{W} is minimal among all spaces of the form (3.1) with $\overline{W}_{\mu} = W^{i}_{\mu}$ and satisfying (3.3).

Proof. Equation (3.3) follows from the termination criterion in Refine_D . Convergence is ensured by (3.2) and $W^i_{\mu} \uparrow V$ for all μ . For the optimality property of the greedy algorithm, we refer to the more general statement [20, Theorem 4.1.5]. \Box

3.2. Adaptive Galerkin method. Let $\|\cdot\|_A$ denote the energy norm on $\ell^2(\Lambda; V)$, i.e., $\|v\|_A \coloneqq \sqrt{\langle Av, v \rangle}$. We assume that a routine

(3.4)
$$\operatorname{Galerkin}_{\boldsymbol{A},\boldsymbol{f}}[\mathcal{W},\tilde{\boldsymbol{u}}_0,\epsilon]\mapsto [\tilde{\boldsymbol{u}},\tau]$$

is available which, given a finite-dimensional subspace \mathcal{W} of $L^2_{\pi}(\Gamma; V)$ of the form (1.28), and starting from the initial approximation $\tilde{\boldsymbol{u}}_0$, iteratively computes $\tilde{\boldsymbol{u}} \in \ell^2(\Lambda; V)$ with $T_V \tilde{\boldsymbol{u}} \in \mathcal{W}$ and

$$\|\tilde{\boldsymbol{u}} - \bar{\boldsymbol{u}}\|_{\boldsymbol{A}} \le \tau \le \epsilon \;,$$

where $T_V \bar{u}$ is the Galerkin projection of u onto \mathcal{W} . An example of such a routine, based on a preconditioned conjugate gradient iteration, is given in [22].

We combine the method $\text{Residual}_{A,f}$ for approximating the residual, Refine_D for refining the Galerkin subspace and $\text{Galerkin}_{A,f}$ for approximating the Galerkin projection, to an adaptive solver $\text{SolveGalerkin}_{A,f}$ similar to [8, 19, 16].

$$\begin{split} & \underbrace{\operatorname{SolveGalerkin}_{A,f}[\epsilon,\gamma,\chi,\vartheta,\omega,\sigma,\alpha,\beta] \mapsto u_{\epsilon}}_{\tilde{u}^{(0)} \longleftarrow 0} \{0\} \\ & \tilde{u}^{(0)} \longleftarrow 0 \\ & \delta_{0} \longleftarrow \sqrt{(1-\gamma)^{-1}} \|f\|_{\ell^{2}(\Lambda;V^{*})} \\ & \text{for } k = 0, 1, 2, \dots \text{ do} \\ & \left[\underbrace{[\boldsymbol{w}_{k},\eta_{k},\zeta_{k}] \leftarrow \operatorname{Residual}_{A,f}[\epsilon\sqrt{1-\gamma},\tilde{\boldsymbol{u}}^{(k)},\delta_{k},\chi,\omega,\alpha,\beta]}_{\delta_{k} \leftarrow (\eta_{k}+\zeta_{k})/\sqrt{1-\gamma}} \\ & \text{if } \min(\delta_{k},\bar{\delta}_{k}) \leq \epsilon \text{ then break} \\ & \left[\underbrace{\mathcal{W}^{(k+1)},\bar{\boldsymbol{w}}_{k},\varrho_{k} \right] \leftarrow \operatorname{Refine}_{D}[\mathcal{W}^{(k)},\boldsymbol{w}_{k},\sqrt{\eta_{k}^{2}-(\zeta_{k}+\vartheta(\eta_{k}+\zeta_{k}))^{2}}] \\ & \bar{\vartheta}_{k} \leftarrow (\sqrt{\eta_{k}^{2}-\varrho_{k}^{2}}-\zeta_{k})/(\eta_{k}+\zeta_{k}) \\ & \left[\underbrace{\tilde{\boldsymbol{u}}^{(k+1)},\tau_{k+1} \right] \leftarrow \operatorname{Galerkin}_{A,f}[\mathcal{W}^{(k+1)},\bar{\boldsymbol{w}}_{k},\sigma\min(\delta_{k},\bar{\delta}_{k})] \\ & \delta_{k+1} \leftarrow \tau_{k+1} + \sqrt{1-\bar{\vartheta}_{k}^{2}(1-\gamma)(1+\gamma)^{-1}}\min(\delta_{k},\bar{\delta}_{k}) \\ & u_{\epsilon} \leftarrow \widetilde{\boldsymbol{u}}^{(k)} \end{split}$$

3.3. Convergence of the adaptive solver. The convergence analysis of the method SolveGalerkin_{A,f} is based on [8, Lemma 4.1], which generalizes to our vector setting for Galerkin spaces \mathcal{W} of the form (1.28). Let $\Pi_{\mathcal{W}}$ denote the orthogonal projection in $\ell^2(\Lambda; V)$ onto $T_V^{-1}\mathcal{W}$, and let $\hat{\Pi}_{\mathcal{W}} \coloneqq D\Pi_{\mathcal{W}}D^{-1}$ be the orthogonal projection in $\ell^2(\Lambda; V^*)$ onto $DT_V^{-1}\mathcal{W} = T_V^*\mathcal{D}\mathcal{W}$.

Proposition 3.2. Let \mathcal{W} be as in (1.28), and $\vartheta \in [0, 1]$. Let $v \in \mathcal{W}$ with

(3.6)
$$\left\|\hat{\Pi}_{\mathcal{W}}(\boldsymbol{f}-\boldsymbol{A}\boldsymbol{v})\right\|_{\ell^{2}(\Lambda;V^{*})} \geq \vartheta \|\boldsymbol{f}-\boldsymbol{A}\boldsymbol{v}\|_{\ell^{2}(\Lambda;V^{*})}.$$

Then the Galerkin projection \bar{u} of u onto W satisfies

(3.7)
$$\|\boldsymbol{u} - \bar{\boldsymbol{u}}\|_{\boldsymbol{A}} \leq \sqrt{1 - \vartheta^2 \frac{1 - \gamma}{1 + \gamma}} \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}$$

Proof. Due to (3.6).

$$\begin{split} \|\bar{\boldsymbol{u}} - \boldsymbol{v}\|_{\boldsymbol{A}} &\geq \|\boldsymbol{A}\|^{-1/2} \|\boldsymbol{A}(\bar{\boldsymbol{u}} - \boldsymbol{v})\|_{\ell^{2}(\Lambda; V^{*})} \geq \|\boldsymbol{A}\|^{-1/2} \left\|\hat{\Pi}_{\mathcal{W}}(\boldsymbol{f} - \boldsymbol{A}\boldsymbol{v})\right\|_{\ell^{2}(\Lambda; V^{*})} \\ &\geq \|\boldsymbol{A}\|^{-1/2} \vartheta \|\boldsymbol{f} - \boldsymbol{A}\boldsymbol{v}\|_{\ell^{2}(\Lambda; V^{*})} \geq \|\boldsymbol{A}\|^{-1/2} \|\boldsymbol{A}^{-1}\|^{-1/2} \vartheta \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}} \ . \end{split}$$

By Galerkin orthogonality,

$$\|\boldsymbol{u} - \bar{\boldsymbol{u}}\|_{\boldsymbol{A}}^{2} = \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2} - \|\bar{\boldsymbol{u}} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2} \le (1 - \vartheta^{2} \|\boldsymbol{A}\|^{-1} \|\boldsymbol{A}^{-1}\|^{-1}) \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2}$$

The assertion follows using the estimates $\|\mathbf{A}\| \leq (1+\gamma)$ and $\|\mathbf{A}^{-1}\| \leq (1-\gamma)^{-1}$, which follow from (1.24).

Lemma 3.3. Let $\epsilon > 0$, $\chi > 0$ and $\alpha, \beta \in (0, 1)$. If $\vartheta > 0$, $\omega > 0$, and $\omega + \vartheta + \omega \vartheta \leq 0$ 1, then the space $\mathcal{W}^{(k+1)}$ in SolveGalerkin_{A,f} is such that

(3.8)
$$\left\|\hat{\Pi}_{\mathcal{W}^{(k+1)}}\boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})} \geq \bar{\vartheta}_{k} \left\|\boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})}$$

where $\mathbf{r}_k \coloneqq \mathbf{f} - \mathbf{A}\tilde{\mathbf{u}}^{(k)}$ is the residual at iteration $k \in \mathbb{N}_0$, and $\bar{\vartheta}_k \geq \vartheta$.

Proof. We abbreviate $\boldsymbol{z} \coloneqq \boldsymbol{w}_k - \tilde{\boldsymbol{u}}^{(k)}$. Due to $\zeta_k \leq \omega \eta_k$, the assumption $\omega + \vartheta + \omega \vartheta \leq \omega \eta_k$ 1 implies $\zeta_k + \vartheta(\eta_k + \zeta_k) \leq \eta_k$. Thus the tolerance in Refine_D is nonnegative. Since $\tilde{\boldsymbol{u}}^{(k)} \in \mathcal{W}^{(k)} \subset \mathcal{W}^{(k+1)}$, Proposition 3.1 implies

$$\varrho_k = \| \boldsymbol{w}_k - \bar{\boldsymbol{w}}_k \|_{\ell^2(\Lambda; V)} = \| \boldsymbol{w}_k - \Pi_{\mathcal{W}^{(k+1)}} \boldsymbol{w}_k \|_{\ell^2(\Lambda; V)} = \| \boldsymbol{z} - \Pi_{\mathcal{W}^{(k+1)}} \boldsymbol{z} \|_{\ell^2(\Lambda; V)} .$$

Consequently,

$$\|\Pi_{\mathcal{W}^{(k+1)}} \boldsymbol{z}\|_{\ell^{2}(\Lambda;V)}^{2} = \|\boldsymbol{z}\|_{\ell^{2}(\Lambda;V)}^{2} - \|\boldsymbol{z} - \Pi_{\mathcal{W}^{(k+1)}} \boldsymbol{z}\|_{\ell^{2}(\Lambda;V)}^{2} = \eta_{k}^{2} - \varrho_{k}^{2}.$$

Furthermore, since $\Pi_{\mathcal{W}^{(k+1)}}$ has norm one, Proposition 2.2 implies

$$\begin{split} \|\Pi_{\mathcal{W}^{(k+1)}} \boldsymbol{z}\|_{\ell^{2}(\Lambda;V)} &- \left\|\hat{\Pi}_{\mathcal{W}^{(k+1)}} \boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})} \leq \left\|\Pi_{\mathcal{W}^{(k+1)}} (\boldsymbol{z} - \boldsymbol{D}^{-1} \boldsymbol{r}_{k})\right\|_{\ell^{2}(\Lambda;V)} \\ &\leq \|\boldsymbol{z} - \boldsymbol{D}^{-1} \boldsymbol{r}_{k}\|_{\ell^{2}(\Lambda;V)} \leq \zeta_{k} \;. \end{split}$$

Combining these estimates, we have

$$\left\|\hat{\Pi}_{\mathcal{W}^{(k+1)}}\boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})} \geq \left\|\Pi_{\mathcal{W}^{(k+1)}}\boldsymbol{z}\right\|_{\ell^{2}(\Lambda;V)} - \zeta_{k} = \sqrt{\eta_{k}^{2} - \varrho_{k}^{2}} - \zeta_{k}$$

and (3.8) follows using $\|\boldsymbol{r}_k\|_{\ell^2(\Lambda; V^*)} \leq \eta_k + \zeta_k$. Finally, $\varrho_k^2 \leq \eta_k^2 - (\zeta_k + \vartheta(\eta_k + \zeta_k))^2$ implies $\sqrt{\eta_k^2 - \varrho_k^2} \ge \zeta_k + \vartheta(\eta_k + \zeta_k)$, and therefore $\bar{\vartheta}_k = (\sqrt{\eta_k^2 - \varrho_k^2} - \zeta_k)/(\eta_k + \zeta_k) \ge 1$ ϑ .

 $\begin{array}{l} \textbf{Theorem 3.4. } I\!\!f\,\epsilon>0,\,\chi>0,\,\vartheta>0,\,\omega>0,\,\omega+\vartheta+\omega\vartheta\leq 1,\,0<\alpha<1,\,0<\beta<1\\ and\,0<\sigma<1-\sqrt{1-\vartheta^2(1-\gamma)(1+\gamma)^{-1}},\,then\,\texttt{SolveGalerkin}_{\pmb{A},\pmb{f}}[\epsilon,\gamma,\chi,\vartheta,\omega,\sigma,\alpha,\beta] \end{array}$ constructs a finitely supported $\mathbf{u}_{\epsilon} \in \ell^2(\Lambda; V)$ with

$$\|\boldsymbol{u} - \boldsymbol{u}_{\epsilon}\|_{\boldsymbol{A}} \le \epsilon \;.$$

Moreover,

(3.10)
$$\sqrt{\frac{1-\gamma}{1+\gamma}\frac{1-\omega}{1+\omega}}\,\bar{\delta}_k \le \left\|\boldsymbol{u}-\tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}} \le \min(\delta_k,\bar{\delta}_k)$$

for all $k \in \mathbb{N}_0$ reached by SolveGalerkin_{A,f}.

Proof. Due to the termination criterion of $SolveGalerkin_{A,f}$, it suffices to show (3.10). For k = 0, since $\|\boldsymbol{u}\|_{\ell^{2}(\Lambda; V)} \leq \|\boldsymbol{A}^{-1}\|^{1/2} \|\boldsymbol{u}\|_{\boldsymbol{A}}$,

$$\left\|oldsymbol{u}- ilde{oldsymbol{u}}^{(0)}
ight\|_{oldsymbol{A}}^{2}=\left\|oldsymbol{u}
ight\|_{oldsymbol{\ell}^{2}\left(\Lambda;V
ight)}\leq\left\|oldsymbol{f}
ight\|_{\ell^{2}\left(\Lambda;V^{*}
ight)}\left\|oldsymbol{u}
ight\|_{\ell^{2}\left(\Lambda;V
ight)}\leq\delta_{0}\left\|oldsymbol{u}
ight\|_{oldsymbol{A}}$$

Let $\left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{A} \leq \delta_k$ for some $k \in \mathbb{N}_0$. Abbreviating $\boldsymbol{r}_k \coloneqq \boldsymbol{f} - \boldsymbol{A}\tilde{\boldsymbol{u}}^{(k)}$, using (1.26) then (2.14), we have

$$\left\|\boldsymbol{u}-\tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}} \leq \frac{1}{\sqrt{1-\gamma}} \left\|\boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})} \leq \frac{\zeta_{k}+\eta_{k}}{\sqrt{1-\gamma}} = \bar{\delta}_{k} \ .$$

If $\min(\delta_k, \bar{\delta}_k) > \epsilon$, then $\zeta_k \leq \omega \eta_k$ by Proposition 2.2. Due to Lemma 3.3, Proposition 3.2 implies

$$\|\boldsymbol{u} - \bar{\boldsymbol{u}}\|_{\boldsymbol{A}} \leq \sqrt{1 - \bar{\vartheta}_k^2 \frac{1 - \gamma}{1 + \gamma}} \min(\delta_k, \bar{\delta}_k) ,$$

where \bar{u} is the exact Galerkin projection of u onto $\mathcal{W}^{(k+1)}$. By (3.5), $\tilde{u}^{(k+1)}$ approximates $\bar{\boldsymbol{u}}$ up to an error of at most $\tau_{k+1} \leq \sigma \min(\delta_k, \bar{\delta}_k)$ in the norm $\|\cdot\|_{\boldsymbol{A}}$. It follows by triangle inequality that $\|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k+1)}\|_{\boldsymbol{A}} \leq \delta_{k+1}$.

To show the other inequality in (3.10), we note that for any $k \in \mathbb{N}_0$,

$$\left\|\boldsymbol{u}-\tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}} \geq \frac{1}{\sqrt{1+\gamma}} \left\|\boldsymbol{r}_{k}\right\|_{\ell^{2}(\Lambda;V^{*})} \geq \frac{\eta_{k}-\zeta_{k}}{\sqrt{1+\gamma}} = \sqrt{\frac{1-\gamma}{1+\gamma}} \frac{\eta_{k}-\zeta_{k}}{\eta_{k}+\zeta_{k}} \,\bar{\delta}_{k} \,,$$

and $(\eta_k - \zeta_k)(\eta_k + \zeta_k)^{-1} \ge (1 - \omega)(1 + \omega)^{-1}$. Finally, since

$$\delta_k \le \left(\sigma + \sqrt{1 - \vartheta^2 (1 - \gamma)(1 + \gamma)^{-1}}\right)^k \delta_0$$

and $\sigma + \sqrt{1 - \vartheta^2 (1 - \gamma)(1 + \gamma)^{-1}} < 1$ by assumption, the iteration does terminate.

4. Optimality properties

4.1. A semidiscrete algorithm. The algorithm $SolveGalerkin_{A,f}$ is derived in Section 3 with arbitrary Galerkin subspaces of the form (1.28). We consider optimality properties of this method in the special case of a single spatial discretization, where a Galerkin subspace $\mathcal{W} \subset \ell^2(\Lambda; V)$ is fully determined by its set of active indices $\Xi \subset \Lambda$.

Since the spatial discretization is fixed throughout, only the part of the residual pertaining to the random part of the error needs to be computed to construct refinements. In particular, no adaptive solver is needed to invert D, making this a viable approach if no such solver is available, or whenever only a single spatial discretization is desired. It is not our intent to suggest that such spaces should generally be used in practice. The adaptive method $SolveGalerkin_{A,f}$ in its full generality has the potential to construct much sparser approximations of \boldsymbol{u} . However, the heuristic distribution of tolerances in $\text{Residual}_{A,f}$ precludes provable optimality statements in this setting; see Remark 2.4.

In this section, we think of the operator A from (1.1) as being already discretized in space, and V is, e.g., a finite element space. Thus, abstractly, we consider a semidiscrete version of the algorithm SolveGalerkin_{A,f}.

The Galerkin subspaces $\mathcal{W}^{(k)}$ have the form $\ell^2(\Xi^{(k)}; V)$ for finite sets $\Xi^{(k)} \subset \Lambda$. In the subroutine **Residual**_{A,f}, we assume that **Solve**_D inverts D exactly in V. The parameter α can thus be set to zero.

In the subsequent refinement step, $\Xi^{(k)}$ is augmented by sufficiently many elements of supp \boldsymbol{w}_k to represent \boldsymbol{w}_k to the desired accuracy. The method Refine_D reduces to ordering supp \boldsymbol{w}_k according to $\|\boldsymbol{w}_{k,\nu}\|_V$ and selecting the most important contributions.

In Galerkin_{A,f}, an iterative solver such as a conjugate gradient iteration is used to approximate the Galerkin projection of u onto $\ell^2(\Xi^{(k+1)}; V)$. Operations within V are assumed to be exact.

4.2. Optimal choice of subspaces. For $\boldsymbol{v} \in \ell^2(\Lambda; V)$ and $N \in \mathbb{N}_0$, let $P_N(\boldsymbol{v})$ be a best N-term approximation of \boldsymbol{v} , that is, $P_N(\boldsymbol{v})$ is an element of $\ell^2(\Lambda; V)$ that minimizes $\|\boldsymbol{v} - \boldsymbol{v}_N\|_{\ell^2(\Lambda; V)}$ over $\boldsymbol{v}_N \in \ell^2(\Lambda; V)$ with $\# \operatorname{supp} \boldsymbol{v}_N \leq N$. For $s \in (0, \infty)$, we define

(4.1)
$$\|\boldsymbol{v}\|_{\mathcal{A}^{s}(\Lambda;V)} \coloneqq \sup_{N \in \mathbb{N}_{0}} (N+1)^{s} \|\boldsymbol{v} - P_{N}(\boldsymbol{v})\|_{\ell^{2}(\Lambda;V)}$$

and

(4.2)
$$\mathcal{A}^{s}(\Lambda; V) \coloneqq \left\{ \boldsymbol{v} \in \ell^{2}(\Lambda; V) \; ; \; \|\boldsymbol{v}\|_{\mathcal{A}^{s}(\Lambda; V)} < \infty \right\} \; .$$

By definition, an optimal approximation in $\ell^2(\Lambda; V)$ of $\boldsymbol{v} \in \mathcal{A}^s(\Lambda; V)$ with error tolerance $\epsilon > 0$ consists of $\mathcal{O}(\epsilon^{-1/s})$ nonzero coefficients in V.

For any $\Xi \subset \Lambda$, let Π_{Ξ} denote the orthogonal projection in $\ell^2(\Lambda; V^*)$ onto $\ell^2(\Xi; V^*)$. The following statement is adapted from [19, Lemma 2.1] and [16, Lemma 4.1].

Lemma 4.1. Let $\Xi^{(0)}$ be a finite subset of Λ and $\mathbf{v} \in \ell^2(\Xi^{(0)}; V)$. If

(4.3)
$$0 \le \hat{\vartheta} < \sqrt{\frac{1-\gamma}{1+\gamma}}$$

and
$$\Xi^{(0)} \subset \Xi^{(1)} \subset \Lambda$$
 with
(4.4)
 $\#\Xi^{(1)} \leq \bar{c} \min \left\{ \#\Xi \; ; \; \Xi^{(0)} \subset \Xi \; , \; \|\Pi_{\Xi}(\boldsymbol{f} - \boldsymbol{A}\boldsymbol{v})\|_{\ell^{2}(\Lambda; V^{*})} \geq \hat{\vartheta} \, \|\boldsymbol{f} - \boldsymbol{A}\boldsymbol{v}\|_{\ell^{2}(\Lambda; V^{*})} \right\}$

for a
$$\bar{c} \geq 1$$
, then

(4.5)
$$\#(\Xi^{(1)} \setminus \Xi^{(0)}) \le \bar{c} \min\left\{ \#\hat{\Xi} \; ; \; \hat{\Xi} \subset \Lambda \; , \; \|\boldsymbol{u} - \hat{\boldsymbol{u}}\|_{\boldsymbol{A}} \le \tau \, \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}} \right\}$$

for $\tau = \sqrt{1 - \hat{\vartheta}^2 (1 + \gamma)(1 - \gamma)^{-1}}$, where $\hat{\boldsymbol{u}}$ denotes the Galerkin projection of \boldsymbol{u} onto $\ell^2(\hat{\boldsymbol{\Xi}}; V)$.

Proof. Let $\hat{\Xi}$ be as in (4.5) and $\check{\Xi} := \Xi^{(0)} \cup \hat{\Xi}$. Furthermore, let \hat{u} and \check{u} denote the Galerkin solutions in $\ell^2(\hat{\Xi}; V)$ and $\ell^2(\check{\Xi}; V)$, respectively. Since $\hat{\Xi} \subset \check{\Xi}$,

 $\|\boldsymbol{u} - \check{\boldsymbol{u}}\|_{\boldsymbol{A}} \le \|\boldsymbol{u} - \hat{\boldsymbol{u}}\|_{\boldsymbol{A}}$, and by Galerkin orthogonality,

$$\|\breve{\boldsymbol{u}} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2} = \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2} - \|\boldsymbol{u} - \breve{\boldsymbol{u}}\|_{\boldsymbol{A}}^{2} \ge (1 - \tau^{2}) \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2} = \hat{\vartheta}^{2} \frac{1 + \gamma}{1 - \gamma} \|\boldsymbol{u} - \boldsymbol{v}\|_{\boldsymbol{A}}^{2}$$

Therefore, using $\kappa(\mathbf{A}) = \|\mathbf{A}\| \|\mathbf{A}^{-1}\| \le (1+\gamma)(1-\gamma)^{-1}$,

$$egin{aligned} & \left\|\Pi_{\check{arepsilon}}(oldsymbol{f}-oldsymbol{A}oldsymbol{v})
ight\|_{\ell^2(\Lambda;V^*)} &= \left\|oldsymbol{A}(oldsymbol{\check{u}}-oldsymbol{v})
ight\|_{\ell^2(\Lambda;V^*)} &\geq \hat{artheta} \left\|oldsymbol{A}-oldsymbol{v}
ight\|_{\ell^2(\Lambda;V^*)} &= \hat{artheta} \left\|oldsymbol{A}-oldsymbol{v}
ight\|_{\ell^2(\Lambda;V^*)} &\geq \hat{artheta} \left\|oldsymbol{A}-oldsymbol{v}
ight\|_{\ell^2(\Lambda;V^*)} &= \hat{artheta} \left\|oldsymbol{A}-oldsymbol{A}-oldsymbol{v}
ight\|_{\ell^2(\Lambda;V^*)} &= \hat{artheta} \left\|oldsymbol{A}-oldsymbol{A}-oldsymbol{v}
ight\|_{\ell^2(\Lambda;V^*)} &= \hat{artheta} \left\|oldsymbol{A}-oldsymbo$$

By (4.4), $\#\Xi^{(1)} \leq \bar{c}\#\breve{\Xi}$ and, consequently,

$$\#(\Xi^{(1)} \setminus \Xi^{(0)}) \le \bar{c} \#(\check{\Xi} \setminus \Xi^{(0)}) \le \bar{c} \# \hat{\Xi} .$$

We use Lemma 4.1 to show that, under additional assumptions on the parameters, the index sets $\Xi^{(k)}$ generated by the semidiscrete version of SolveGalerkin_{A,f} are of optimal size, up to a constant factor.

Theorem 4.2. If the conditions of Theorem 3.4 are satisfied,

(4.6)
$$\hat{\vartheta} \coloneqq \frac{\vartheta(1+\omega) + 2\omega}{1-\omega} < \sqrt{\frac{1-\gamma}{1+\gamma}} ,$$

and $u \in \mathcal{A}^{s}(\Lambda; V)$ for an s > 0, then for all $k \in \mathbb{N}_{0}$ reached by SolveGalerkin_{A,f},

(4.7)
$$\#\Xi^{(k)} \le 2\frac{(\varrho/\tau)^{1/s}}{1-\varrho^{1/s}} \left(\frac{(1+\gamma)(1+\omega)}{(1-\gamma)(1-\omega)}\right)^{1/s} \left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\ell^2(\Lambda;V)}^{-1/s} \left\| \boldsymbol{u} \right\|_{\mathcal{A}^s(\Lambda;V)}^{1/s}$$

with
$$\rho = \sigma + \sqrt{1 - \vartheta^2 (1 - \gamma)(1 + \gamma)^{-1}}$$
 and $\tau = \sqrt{1 - \vartheta^2 (1 + \gamma)(1 - \gamma)^{-1}}$.

Proof. Let $k \in \mathbb{N}_0$, $\mathbf{r}_k = \mathbf{f} - \mathbf{A}\tilde{\mathbf{u}}^{(k)}$. Also, let $\boldsymbol{\varrho} = (\varrho_{\nu})_{\nu \in \Lambda}$, $\varrho_{\nu} \coloneqq \left\| w_{k,\nu} - \tilde{u}_{\nu}^{(k)} \right\|_V$ for the approximation $\mathbf{w}_k - \tilde{\mathbf{u}}^{(k)} = (w_{k,\nu} - \tilde{u}_{\nu}^{(k)})_{\nu \in \Lambda}$ of $\mathbf{D}^{-1}\mathbf{r}_k$ computed in Residual \mathbf{A}, \mathbf{f} , and let $\Delta \subset \operatorname{supp} \mathbf{w}_k$ denote the active indices selected by Refine_D.

We note that for $\alpha \coloneqq \omega + \vartheta + \omega \vartheta$, we have $\vartheta = \frac{\alpha - \omega}{1 + \omega}$ and $\hat{\vartheta} = \frac{\alpha + \omega}{1 - \omega}$. Let $\Xi^{(k)} \subset \bar{\Xi} \subset \Lambda$ satisfy $\|\Pi_{\bar{\Xi}} \boldsymbol{r}_k\|_{\ell^2(\Lambda; V^*)} \ge \hat{\vartheta} \|\boldsymbol{r}_k\|_{\ell^2(\Lambda; V^*)}$. Also, if $\tilde{\boldsymbol{u}}^{(k)}$ is used to refine the discretization, then the tolerance ϵ is not yet reached, and thus $\|\boldsymbol{\varrho}\|_{\ell^2(\Lambda)} - \|\boldsymbol{r}_k\|_{\ell^2(\Lambda; V^*)} \le \omega \|\boldsymbol{\varrho}\|_{\ell^2(\Lambda)}$ by Proposition 2.2. Therefore,

$$\begin{split} \hat{\vartheta} \| \boldsymbol{\varrho} \|_{\ell^{2}(\Lambda)} &\leq \hat{\vartheta} \| \boldsymbol{r}_{k} \|_{\ell^{2}(\Lambda; V^{*})} + \hat{\vartheta} \omega \| \boldsymbol{\varrho} \|_{\ell^{2}(\Lambda)} \\ &\leq \| \Pi_{\bar{\Xi}} \boldsymbol{r}_{k} \|_{\ell^{2}(\Lambda; V^{*})} + \hat{\vartheta} \omega \| \boldsymbol{\varrho} \|_{\ell^{2}(\Lambda)} \leq \| \Pi_{\bar{\Xi}} \boldsymbol{\varrho} \|_{\ell^{2}(\Lambda)} + (1 + \hat{\vartheta}) \omega \| \boldsymbol{\varrho} \|_{\ell^{2}(\Lambda)} \end{split}$$

and since $\hat{\vartheta} - (1 + \hat{\vartheta})\omega = \alpha$, it follows that $\|\Pi_{\bar{\Xi}}\boldsymbol{\varrho}\|_{\ell^2(\Lambda)} \geq \alpha \|\boldsymbol{\varrho}\|_{\ell^2(\Lambda)}$. By construction, Δ is a set of minimal cardinality with $\|\Pi_{\Delta}\boldsymbol{\varrho}\|_{\ell^2(\Lambda)} \geq \bar{\alpha} \|\boldsymbol{\varrho}\|_{\ell^2(\Lambda)}$ for $\bar{\alpha} \coloneqq \zeta_k \eta_k^{-1} + \vartheta(1 + \zeta_k \eta_k^{-1}) \leq \alpha$. Consequently, $\#(\Xi^{(k+1)} \setminus \Xi^{(k)}) \leq \#\Delta \leq \#\bar{\Xi}$. Since this holds for any $\bar{\Xi}$, using $\#\Xi^{(k)} \leq \bar{\Xi}$, it follows that

$$\#\Xi^{(k+1)} \le 2\min\left\{\#\bar{\Xi}\,;\,\Xi^{(k)} \subset \bar{\Xi} \subset \Lambda\,,\, \|\Pi_{\bar{\Xi}}\boldsymbol{r}_k\|_{\ell^2(\Lambda;V^*)} \ge \hat{\vartheta}\,\|\boldsymbol{r}_k\|_{\ell^2(\Lambda;V^*)}\right\}\,.$$

Lemma 4.1 implies

$$\#(\Xi^{(k+1)} \setminus \Xi^{(k)}) \le 2\min\left\{\#\hat{\Xi} \; ; \; \hat{\Xi} \subset \Lambda, \; \|\boldsymbol{u} - \hat{\boldsymbol{u}}\|_{\boldsymbol{A}} \le \tau \left\|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}}\right\}$$

with $\tau = \sqrt{1 - \hat{\vartheta}^2 (1 + \gamma)(1 - \gamma)^{-1}}$, where $\hat{\boldsymbol{u}}$ denotes the Galerkin projection of \boldsymbol{u} onto $\ell^2(\hat{\boldsymbol{\Xi}}; V)$.

Let $N \in \mathbb{N}_0$ be maximal with $\|\boldsymbol{u} - P_N(\boldsymbol{u})\|_{\ell^2(\Lambda;V)} > \tau(1+\gamma)^{-1/2} \|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)}\|_{\boldsymbol{A}}$, where $P_N(\boldsymbol{u})$ is a best N-term approximation of \boldsymbol{u} . By (4.1),

$$N+1 \le \|\boldsymbol{u} - P_N(\boldsymbol{u})\|_{\ell^2(\Lambda;V)}^{-1/s} \|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda;V)}^{1/s}$$

$$\le \tau^{-1/s} (1+\gamma)^{1/2s} \|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)}\|_{\boldsymbol{A}}^{-1/s} \|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda;V)}^{1/s}$$

For $\Xi_{N+1} \coloneqq \operatorname{supp} P_{N+1}(\boldsymbol{u})$, by maximality of N,

$$\|\boldsymbol{u} - \bar{\boldsymbol{u}}_{N+1}\|_{\boldsymbol{A}} \leq \|\boldsymbol{u} - P_{N+1}(\boldsymbol{u})\|_{\boldsymbol{A}}$$

$$\leq (1+\gamma)^{1/2} \|\boldsymbol{u} - P_{N+1}(\boldsymbol{u})\|_{\ell^{2}(\Lambda;V)} \leq \tau \left\|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}}$$

for the Galerkin solution $\bar{\boldsymbol{u}}_{N+1}$ in $\ell^2(\Xi_{N+1}; V)$, and thus

$$#(\Xi^{(k+1)} \setminus \Xi^{(k)}) \le 2(N+1) \le 2\tau^{-1/s} (1+\gamma)^{1/2s} \left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\boldsymbol{A}}^{-1/s} \left\| \boldsymbol{u} \right\|_{\mathcal{A}^{s}(\Lambda;V)}^{1/s}$$

Furthermore, by Theorem 3.4,

$$\left\|\boldsymbol{u}-\tilde{\boldsymbol{u}}^{(k)}\right\|_{\boldsymbol{A}}^{-1/s} \leq \left(\sqrt{\frac{1-\gamma}{1+\gamma}}\frac{1-\omega}{1+\omega}\,\bar{\delta}_k\right)^{-1/s}$$

We estimate the cardinality of $\Xi^{(k)}$ by slicing it into increments and applying the above estimates,

$$\begin{aligned} \#\Xi^{(k)} &= \sum_{j=0}^{k-1} \#(\Xi^{(j+1)} \setminus \Xi^{(j)}) \le 2\tau^{-1/s} (1+\gamma)^{1/2s} \|\boldsymbol{u}\|_{\mathcal{A}^{s}(\Lambda;V)}^{1/s} \sum_{j=0}^{k-1} \left\|\boldsymbol{u} - \tilde{\boldsymbol{u}}^{(j)}\right\|_{\boldsymbol{A}}^{-1/s} \\ &\le 2 \left(\frac{\tau (1-\gamma)^{1/2} (1-\omega)}{(1+\gamma)(1+\omega)}\right)^{-1/s} \|\boldsymbol{u}\|_{\mathcal{A}^{s}(\Lambda;V)}^{1/s} \sum_{j=0}^{k-1} \bar{\delta}_{j}^{-1/s} .\end{aligned}$$

By definition, $\delta_k \leq \varrho^{k-j} \bar{\delta}_j$. Therefore,

$$\sum_{j=0}^{k-1} \bar{\delta}_j^{-1/s} \le \delta_k^{-1/s} \sum_{j=0}^{k-1} \varrho^{(k-j)/s} = \delta_k^{-1/s} \sum_{i=1}^k \varrho^{i/s} = \frac{\varrho^{1/s} \delta_k^{-1/s}}{1-\varrho^{1/s}} \,.$$

The assertion follows using

$$(1-\gamma)^{1/2} \left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\ell^2(\Lambda;V)} \leq \left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\boldsymbol{A}} \leq \delta_k . \qquad \Box$$

4.3. **Complexity estimate.** We first cite an elementary result due to Stechkin connecting the order of summability of a sequence to the convergence of best *N*-term approximations in a weaker sequence norm; see e.g. [11, 15]. Note that, although it is formulated only for nonnegative sequences, Lemma 4.3 applies directly to, e.g., Lebesgue–Bochner spaces of Banach space valued sequences by passing to the norms of the elements of such sequences. Also, it applies to sequences with arbitrary countable index sets by choosing a decreasing rearrangement.

Lemma 4.3. Let $0 and let <math>\mathbf{c} = (c_n)_{n=1}^{\infty} \in \ell^2$ with $0 \le c_{n+1} \le c_n$ for all $n \in \mathbb{N}$. Then

(4.8)
$$\left(\sum_{n=N+1}^{\infty} c_n^q\right)^{1/q} \le (N+1)^{-r} \|\boldsymbol{c}\|_{\ell^p} , \quad r \coloneqq \frac{1}{p} - \frac{1}{q} \ge 0$$

for all $N \in \mathbb{N}_0$.

Proposition 4.4. Let s > 0. If either

(4.9)
$$\|R_m\|_{V \to V^*} \le s \delta_{\mathbf{R},s} (m+1)^{-s-1} \quad \forall m \in \mathbb{N}$$

or

(4.10)
$$\left(\sum_{m=1}^{\infty} \|R_m\|_{V \to V^*}^{\frac{1}{s+1}}\right)^{s+1} \le \delta_{\mathbf{R},s} ,$$

then

(4.11)
$$\|\boldsymbol{R} - \boldsymbol{R}_{[M]}\|_{\ell^2(\Lambda; V) \to \ell^2(\Lambda; V^*)} \leq \delta_{\boldsymbol{R}, s} (M+1)^{-s} \quad \forall M \in \mathbb{N}_0 .$$

Proof. By (1.23) and (2.1), using $\|\boldsymbol{K}_m\|_{\ell^2(\Lambda) \to \ell^2(\Lambda)} \leq 1$,

$$\|\boldsymbol{R} - \boldsymbol{R}_{[M]}\|_{\ell^{2}(\Lambda; V) \to \ell^{2}(\Lambda; V^{*})} \leq \sum_{m=M+1}^{\infty} \|R_{m}\|_{V \to V^{*}}$$

If (4.9) holds, then (4.11) follows using

$$\sum_{m=M+1}^{\infty} (m+1)^{-s-1} \le \int_{M+1}^{\infty} t^{-s-1} \, \mathrm{d}t = \frac{1}{s} \, (M+1)^{-s} \, .$$

If (4.10) is satisfied, then

$$\sum_{m=M+1}^{\infty} \|R_m\|_{V \to V^*} \le \left(\sum_{m=1}^{\infty} \|R_m\|_{V \to V^*}^{\frac{1}{s+1}}\right)^{s+1} (M+1)^{-s}$$

by Lemma 4.3.

Remark 4.5. If the assumptions of Proposition 4.4 are satisfied for all $s \in (0, s^*)$, then the operator \mathbf{R} is s^* -compressible with sparse approximations $\mathbf{R}_{[M]}$. In this case, \mathbf{R} is a bounded linear map from $\mathcal{A}^s(\Lambda; V)$ to $\mathcal{A}^s(\Lambda; V^*)$ for all $s \in (0, s^*)$; see [8, Prop. 3.8]. This carries over to the routine $\operatorname{Apply}_{\mathbf{R}}$ in that if $\mathbf{v} \in \mathcal{A}^s(\Lambda; V)$ and \mathbf{z} is the output of $\operatorname{Apply}_{\mathbf{R}}[\mathbf{v}, \epsilon]$ for an $\epsilon > 0$, then

(4.12)
$$\# \operatorname{supp} \boldsymbol{z} \lesssim \|\boldsymbol{v}\|_{\mathcal{A}^s(\Lambda;V)}^{1/s} \epsilon^{-1/s},$$

(4.13)
$$\|\boldsymbol{z}\|_{\mathcal{A}^{s}(\Lambda; V^{*})} \lesssim \|\boldsymbol{v}\|_{\mathcal{A}^{s}(\Lambda; V)}$$

with constants depending only on s and \mathbf{R} . Moreover, (4.12) is an upper bound for the total number of applications of operators R_m in $\text{Apply}_{\mathbf{R}}[\boldsymbol{v},\epsilon]$. This follows as in the scalar case (see e.g. [16, Prop. 4.6]), where the additional term $1 + \# \operatorname{supp} \boldsymbol{v}$ is only due to the approximate sorting of \boldsymbol{v} .

We make further assumptions on the routine RHS_{f} . If $f \in \mathcal{A}^{s}(\Lambda; V^{*})$ and \tilde{f} is the output of $\operatorname{RHS}_{f}[\epsilon]$ for an $\epsilon > 0$, then \tilde{f} should satisfy

(4.14)
$$\# \operatorname{supp} \tilde{\boldsymbol{f}} \lesssim \|\boldsymbol{f}\|_{\mathcal{A}^{s}(\Lambda; V^{*})}^{1/s} \epsilon^{-1/s} .$$

This is clearly satisfied for deterministic f, and is achieved for the right-hand sides of the form $\mathbf{R}\mathbf{w}$ for a finitely supported \mathbf{w} , stemming for example from inhomogeneous essential boundary conditions, by using $Apply_R$ to approximate this product. Note that if $u \in \mathcal{A}^{s}(\Lambda; V)$ and **R** is s^{*}-compressible with $s < s^{*}$, then also **A** is s^{*}compressible, and therefore $\|f\|_{\mathcal{A}^{s}(\Lambda;V^{*})} \lesssim \|u\|_{\mathcal{A}^{s}(\Lambda;V)}$.

Lemma 4.6. Under the conditions of Theorem 4.2,

(4.15)
$$\left\| \tilde{\boldsymbol{u}}^{(k)} \right\|_{\mathcal{A}^{s}(\Lambda;V)} \leq C \left\| \boldsymbol{u} \right\|_{\mathcal{A}^{s}(\Lambda;V)} \quad \forall k \in \mathbb{N}_{0} ,$$

with

(4.16)
$$C = 1 + \frac{2^{1+s}\varrho(1+\gamma)(1+\omega)}{\tau(1-\varrho^{1/s})^s(1-\gamma)(1-\omega)}$$

$$\rho = \sigma + \sqrt{1 - \vartheta^2 (1 - \gamma)(1 + \gamma)^{-1}}$$
 and $\tau = \sqrt{1 - \hat{\vartheta}^2 (1 + \gamma)(1 - \gamma)^{-1}}$

Proof. Let $k \in \mathbb{N}_0$. For any $N \geq \#\Xi^{(k)}, \left\| \tilde{\boldsymbol{u}}^{(k)} - P_N(\tilde{\boldsymbol{u}}^{(k)}) \right\|_{\ell^2(\Lambda \cdot V)} = 0$. For $N < \# \Xi^{(k)} - 1,$

$$\begin{split} \left\| \tilde{\boldsymbol{u}}^{(k)} - P_N(\tilde{\boldsymbol{u}}^{(k)}) \right\|_{\ell^2(\Lambda;V)} &\leq \left\| \tilde{\boldsymbol{u}}^{(k)} - \Pi_{\Xi_N} \tilde{\boldsymbol{u}}^{(k)} \right\|_{\ell^2(\Lambda;V)} \\ &\leq \left\| \boldsymbol{u} - \Pi_{\Xi_N} \boldsymbol{u} \right\|_{\ell^2(\Lambda;V)} + 2 \left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\ell^2(\Lambda;V)} \end{split}$$

where $\Xi_N \coloneqq \operatorname{supp} P_N(\boldsymbol{u})$, such that $\Pi_{\Xi_N} \boldsymbol{u} = P_N(\boldsymbol{u})$ and

$$\|\boldsymbol{u} - \boldsymbol{\Pi}_{\boldsymbol{\Xi}_N} \boldsymbol{u}\|_{\ell^2(\Lambda;V)} \leq (N+1)^{-s} \|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda;V)} .$$

Furthermore, Theorem 4.2 implies

$$\left\| \boldsymbol{u} - \tilde{\boldsymbol{u}}^{(k)} \right\|_{\ell^{2}(\Lambda; V)} \leq \frac{2^{s} \varrho(1 + \gamma)(1 + \omega)}{\tau(1 - \varrho^{1/s})^{s}(1 - \gamma)(1 - \omega)} (\# \Xi^{(k)})^{-s} \| \boldsymbol{u} \|_{\mathcal{A}^{s}(\Lambda; V)} ,$$

and $(N+1)^s < (\#\Xi^{(k)})^s$ by the definition of N. Consequently,

$$\left\|\tilde{\boldsymbol{u}}^{(k)}\right\|_{\mathcal{A}^{s}(\Lambda;V)} = \sup_{N \in \mathbb{N}_{0}} (N+1)^{s} \left\|\tilde{\boldsymbol{u}}^{(k)} - P_{N}(\tilde{\boldsymbol{u}}^{(k)})\right\|_{\ell^{2}(\Lambda;V)} \leq C \left\|\boldsymbol{u}\right\|_{\mathcal{A}^{s}(\Lambda;V)}$$

C from (4.16).

with C from (4.16).

Theorem 4.7. Let the conditions of Theorem 4.2 be satisfied. If (4.14) and the assumptions of Proposition 4.4 hold for all $s \in (0, s^*)$, then for any $\epsilon > 0$ and any $s \in (0, s^*)$, the total number of applications of D, $A_{\nu\nu}$ and D^{-1} in SolveGalerkin_{A,f}[$\epsilon, \gamma, \chi, \vartheta, \omega, \sigma, 0, \beta$] is bounded by $\|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda; V)}^{1/s} \epsilon^{-1/s}$ up to a constant factor depending only on the input arguments other than ϵ . The same bound holds for the total number of applications of R_m , $m \in \mathbb{N}$, up to an additional factor of $\max_{\mu \in \text{supp}} \boldsymbol{u}_{\epsilon} \# \text{supp} \mu$.

Proof. Let $k \in \mathbb{N}_0$; we consider the k-th iteration of the loop in SolveGalerkin_{A,f}. The routine Residual_{A,f} $[\epsilon \sqrt{1-\gamma}, \tilde{\boldsymbol{u}}^{(\Xi^{(k)})}, \delta_k, \chi, \omega, \beta]$ begins with $\#\Xi^{(k)}$ applications of D. Due to the geometric decrease in tolerances, the complexity of the loop in $\texttt{Residual}_{A,f}$ is dominated by that of its last iteration. By Remark 4.5 and Lemma 4.6, up to a constant factor, the number of applications of D^{-1} and R_m is bounded by $\|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda;V)}^{1/s} \zeta_k^{-1/s}$, and $\zeta_k \gtrsim \bar{\delta}_k$. Next, assuming the termination criterion of SolveGalerkin_{A,f} is not satisfied, the routine Galerkin_{A,f} $[\Xi^{(k+1)}, \boldsymbol{w}, \sigma \min(\delta_k, \bar{\delta}_k)]$ is called to iteratively approximate the Galerkin projection onto $\ell^2(\Xi^{(k+1)}; V)$. Since only a fixed relative error reduction is required, the number of iterations remains bounded. Therefore, the number of applications of D^{-1} and $A_{\nu\nu}$ is bounded by $\#\Xi^{(k+1)}$ and the total number of applications of $R_m, m \in \mathbb{N}$, is bounded by $2\bar{\lambda}(\Xi^{(k+1)})\#\Xi^{(k+1)}$, where $\bar{\lambda}(\Xi^{(k+1)})$ denote the average length of indices in $\Xi^{(k+1)}$; see [22, Proposition 3.5]. Since the sets $\Xi^{(k)}$ are nested, $\bar{\lambda}(\Xi^{(k+1)}) \leq \max_{\mu \in \text{supp} \boldsymbol{u}_{\epsilon}} \# \text{supp } \mu$. Furthermore, by Theorems 3.4 and 4.2, $\#\Xi^{(k+1)} \lesssim \|\boldsymbol{u}\|_{\mathcal{A}^s(\Lambda;V)}^{1/s} \bar{\delta}_{k+1}^{-1/s}$.

Let k be such that $\boldsymbol{u}_{\epsilon} = \tilde{\boldsymbol{u}}^{(k)}$. Due to the different termination criterion, the complexity of the last call of Residual_{A,f} can be estimated by $\|\boldsymbol{u}\|_{\mathcal{A}^{s}(\Lambda;V)}^{1/s} \zeta_{k}^{-1/s}$ with $\zeta_{k} \gtrsim \epsilon$. This bound obviously also holds for $\#\Xi^{(k)}$, and thus for the complexity of the final call of Galerkin_{A,f}.

Combining all of the above estimates, the number of applications of D^{-1} , D, $A_{\nu\nu}$ and R_m , $m \in \mathbb{N}$, in SolveGalerkin_{A,f} is bounded by

$$\| m{u} \|_{\mathcal{A}^{s}(\Lambda;V)}^{1/s} \left(\epsilon^{-1/s} + \sum_{j=0}^{k-1} ar{\delta}_{j}^{-1/s}
ight) \; .$$

Furthermore, $\bar{\delta}_{k-1} \ge \epsilon$, and using $\delta_{k-1} \le \varrho^{k-1-j}\bar{\delta}_j$,

$$\sum_{j=0}^{k-2} \bar{\delta}_j^{-1/s} \le \delta_{k-1}^{-1/s} \sum_{j=0}^{k-2} \varrho^{(k-1-j)/s} = \delta_{k-1}^{-1/s} \sum_{i=1}^{k-1} \varrho^{i/s} \le \delta_{k-1}^{-1/s} \frac{\varrho^{1/s}}{1-\varrho^{1/s}} + \frac{\rho^{1/s}}{1-\varrho^{1/s}} + \frac{\rho^{1/s}}{1-\rho^{1/s}} + \frac{\rho^{1/s}}{1-\rho^{1/s$$

where $\rho = \sigma + \sqrt{1 - \vartheta^2 (1 - \gamma)(1 + \gamma)^{-1}} < 1$. The assertion follows since $\delta_{k-1} \geq \epsilon$.

5. Computational examples

5.1. Application to isotropic diffusion. We consider the isotropic diffusion equation on a bounded Lipschitz domain $G \subset \mathbb{R}^d$ with homogeneous Dirichlet boundary conditions. For any uniformly positive $a \in L^{\infty}(G)$ and any $f \in L^2(G)$, we have

(5.1)
$$-\nabla \cdot (a(x)\nabla u(x)) = f(x) , \quad x \in G ,$$
$$u(x) = 0 , \quad x \in \partial G .$$

We view f as fixed, but allow a to vary, giving rise to a parametric operator

(5.2)
$$A_0(a): H_0^1(G) \to H^{-1}(G) , \quad v \mapsto -\nabla \cdot (a\nabla v) ,$$

which depends continuously on $a \in L^{\infty}(G)$.

We model the coefficient a as a bounded random field, which we expand as a series

(5.3)
$$a(y,x) \coloneqq \bar{a}(x) + \sum_{m=1}^{\infty} y_m a_m(x) \; .$$

Since a is bounded, a_m can be scaled such that $y_m \in [-1, 1]$ for all $m \in \mathbb{N}$. Therefore, a depends on a parameter $y = (y_m)_{m=1}^{\infty}$ in $\Gamma = [-1, 1]^{\infty}$. We define the parametric operator $A(y) \coloneqq A_0(a(y))$ for $y \in \Gamma$. Due to the linearity of A_0 ,

(5.4)
$$A(y) = D + R(y) , \quad R(y) \coloneqq \sum_{m=1}^{\infty} y_m R_m \qquad \forall y \in \Gamma$$

with convergence in $\mathcal{L}(H_0^1(G), H^{-1}(G))$, for

$$D \coloneqq A_0(\bar{a}) \colon H_0^1(G) \to H^{-1}(G) \ , \quad v \mapsto -\nabla \cdot (\bar{a}\nabla v) \ ,$$
$$R_m \coloneqq A_0(a_m) \colon H_0^1(G) \to H^{-1}(G) \ , \quad v \mapsto -\nabla \cdot (a_m \nabla v) \ , \quad m \in \mathbb{N} \ .$$

To ensure bounded invertibility of D, we assume there is a constant $\delta > 0$ such that

(5.5)
$$\operatorname{ess\,inf}_{x\in G} \bar{a}(x) \ge \delta^{-1} \,.$$

We refer, e.g., to [22, 20, 26] for further details.

5.2. A posteriori error estimation. Let the spaces W_{ν} from Section 1.5 be finite element spaces of continuous, piecewise smooth functions on meshes \mathfrak{T}_{ν} which contain at least the piecewise linear functions on \mathfrak{T}_{ν} . We assume that these meshes are compatible in the sense that for any $T_{\mu} \in \mathfrak{T}_{\mu}$ and $T_{\nu} \in \mathfrak{T}_{\nu}$, the intersection $T_{\mu} \cap T_{\nu}$ is either empty, equal to T_{μ} , or equal to T_{ν} . We denote the set of faces of \mathfrak{T}_{ν} by \mathfrak{F}_{ν} and define h_T and h_F as the diameters of $T \in \mathfrak{T}_{\nu}$ and $F \in \mathfrak{F}_{\nu}$, respectively.

In Residual_{*A*,*f*}, a generic solver Solve_{*D*} is used to approximate $D^{-1}g_{\nu}$ up to a prescribed tolerance. In the present finite element setting, this requires a reliable a posteriori error estimator to verify that the desired accuracy is attained.

The vector $\boldsymbol{g} = (g_{\nu})_{\nu \in A}$ is the approximation of $\boldsymbol{f} - \boldsymbol{R}\boldsymbol{v}$ computed with $\operatorname{RHS}_{\boldsymbol{f}}$ and $\operatorname{Apply}_{\boldsymbol{R}}$. For the call of $\operatorname{Residual}_{\boldsymbol{A},\boldsymbol{f}}$ inside $\operatorname{SolveGalerkin}_{\boldsymbol{A},\boldsymbol{f}}, \boldsymbol{v}$ is the approximate solution $\tilde{\boldsymbol{u}}^{(k)}$. Thus g_{ν} has the form

(5.6)
$$g_{\nu} = \tilde{f}_{\nu} - \sum_{i=1}^{k} \kappa_i R_{m_i} v_i ,$$

where \tilde{f}_{ν} is the approximation of f_{ν} generated by RHS_{f} , $v_{i} = v_{\mu_{i}}$ for some $\mu_{i} = \nu \pm \epsilon_{m_{i}}$ selected by Apply_{R} , and κ_{i} refer to the constants $\alpha_{\nu_{m_{i}}}^{m_{i}}$ and $\beta_{\max(\nu_{m_{i}},\mu_{m_{i}})}^{m_{i}}$ from (1.22). We abbreviate $\mathfrak{T}_{i} \coloneqq \mathfrak{T}_{\mu_{i}}$.

Standard error estimators have difficulties on faces of \mathfrak{T}_i that are not in the skeleton of \mathfrak{T}_{ν} since g_{ν} is singular on these faces. For all *i*, let \bar{v}_i be an approximation of v_i that is piecewise smooth on \mathfrak{T}_{ν} . Replacing g_{ν} by

(5.7)
$$\bar{g}_{\nu} \coloneqq \tilde{f}_{\nu} - \sum_{i=1}^{k} \kappa_i R_{m_i} \bar{v}_i$$

induces an error

(5.8)
$$\left\| D^{-1}g_{\nu} - D^{-1}\bar{g}_{\nu} \right\|_{V} \le \sum_{i=1}^{k} |\kappa_{i}| \left\| \frac{a_{m_{i}}}{\bar{a}} \right\|_{L^{\infty}(G)} \|v_{i} - \bar{v}_{i}\|_{V} \eqqcolon \mathrm{EST}_{\nu}^{P}$$

since

$$\begin{split} \sup_{\|z\|_{V}=1} \left| \int_{G} a_{m} \nabla v \cdot \nabla z \, \mathrm{d}x \right| &\leq \left\| \frac{a_{m}}{\bar{a}} \right\|_{L^{\infty}(G)} \sup_{\|z\|_{V}=1} \int_{G} |\bar{a} \nabla v \cdot \nabla z| \, \mathrm{d}x \\ &= \left\| \frac{a_{m}}{\bar{a}} \right\|_{L^{\infty}(G)} \|v\|_{V} \end{split}$$

for all $m \in \mathbb{N}$ and all $v \in H_0^1(G)$.

Let $\bar{w}_{\nu} \in W_{\nu}$ be the Galerkin projection of $D^{-1}\bar{g}_{\nu}$, i.e.,

(5.9)
$$\int_{G} \bar{a} \nabla \bar{w}_{\nu} \cdot \nabla z \, \mathrm{d}x = \int_{G} \tilde{f}_{\nu} z \, \mathrm{d}x - \sum_{i=1}^{k} \kappa_{i} \int_{G} a_{m_{i}} \nabla \bar{v}_{i} \cdot \nabla z \, \mathrm{d}x \qquad \forall z \in W_{\nu} .$$

Abbreviating

(5.10)
$$\sigma_{\nu} \coloneqq \bar{a} \nabla \bar{w}_{\mu} + \sum_{i=1}^{k} \kappa_{i} a_{m_{i}} \nabla \bar{v}_{i} ,$$

the residual of \bar{w}_{ν} is the functional

(5.11)
$$r_{\nu}(\bar{w}_{\nu};z) = \int_{G} \bar{g}_{\nu}z - \bar{a}\nabla\bar{w}_{\nu}\cdot\nabla z \,\mathrm{d}x = \int_{G} \tilde{f}_{\nu}z - \sigma_{\nu}\cdot\nabla z \,\mathrm{d}x \;, \quad z \in H^{1}_{0}(G) \;.$$

Due to the Riesz isomorphism,

(5.12)
$$\|D^{-1}\bar{g}_{\nu} - w_{\nu}\|_{V} = \sup_{z \in H_{0}^{1}(G) \setminus \{0\}} \frac{|r_{\nu}(\bar{w}_{\nu}; z)|}{\|z\|_{V}} \le \sqrt{\delta} \sup_{z \in H_{0}^{1}(G) \setminus \{0\}} \frac{|r_{\nu}(\bar{w}_{\nu}; z)|}{|z|_{H^{1}(G)}},$$

with δ from (5.5).

For all $T \in \mathfrak{T}_{\nu}$, let

(5.13)
$$R_{\nu,T}(\bar{w}_{\nu}) \coloneqq h_T \left\| \tilde{f}_{\nu} + \nabla \cdot \sigma_{\nu} \right\|_{L^2(T)}$$

where the dependence on \bar{w}_{ν} is implicit in σ_{ν} . Also, let

(5.14)
$$R_{\nu,F}(\bar{u}_{\nu}) \coloneqq h_F^{1/2} \| \llbracket \sigma_{\nu} \rrbracket \|_{L^2(F)} ,$$

where $\llbracket \cdot \rrbracket$ is the normal jump over the face $F \in \mathfrak{F}_{\nu}$. These terms combine to

(5.15)
$$\operatorname{EST}_{\nu}^{R}(\bar{w}_{\nu}) \coloneqq \left(\sum_{T \in \mathfrak{T}_{\nu}} R_{\nu,T}(\bar{w}_{\nu})^{2} + \sum_{F \in \mathfrak{F}_{\nu}} R_{\nu,F}(\bar{w}_{\nu})^{2}\right)^{1/2}$$

The following statement is a straightforward adaptation of the standard result from, e.g., [29, 25, 1] on reliability of residual error estimators.

Theorem 5.1. For all $z \in H_0^1(G)$,

(5.16)
$$|r_{\nu}(\bar{w}_{\nu};z)| \leq C \operatorname{EST}_{\nu}^{R}(\bar{w}_{\nu}) |z|_{H^{1}(G)}$$

with a constant C depending only on the shape regularity of \mathfrak{T}_{ν} .

Corollary 5.2. The Galerkin projection \bar{w}_{ν} from (5.9) satisfies

(5.17)
$$\left\| D^{-1}g_{\nu} - \bar{w}_{\nu} \right\|_{V} \leq \mathrm{EST}_{\nu}^{P} + \sqrt{\delta}C \, \mathrm{EST}_{\nu}^{R}(\bar{w}_{\nu})$$

for δ from (5.5) and C from Theorem 5.1.

Proof. The assertion follows by triangle inequality using (5.8), (5.12) and (5.16). \Box

5.3. Numerical computations. We consider as a model problem the diffusion equation (5.1) on the one-dimensional domain G = (0, 1). For two parameters k and γ , the diffusion coefficient has the form

(5.18)
$$a(y,x) = 1 + \frac{1}{c} \sum_{m=1}^{\infty} y_m \frac{1}{m^k} \sin(m\pi x), \quad x \in (0,1), \quad y \in \Gamma = [-1,1]^{\infty},$$

where c is chosen as

$$(5.19) c = \gamma \sum_{m=1}^{\infty} \frac{1}{m^k} ,$$

such that |a(y, x) - 1| is always less than γ . For the distribution of $y \in \Gamma$, we consider the countable product of uniform distributions on [-1, 1]; the corresponding family of orthonormal polynomials is the Legendre polynomial basis.

In all of the following computations, the parameters are k = 2 and $\gamma = 1/2$. A few realizations of a(y) and the resulting solutions u(y) of (5.1) are plotted in Figure 1.

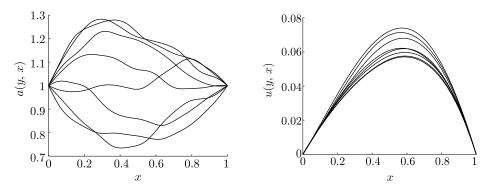


FIGURE 1. Realizations of a(y, x) (left) and u(y, x) (right).

The parameters of SolveGalerkin_{A,f} are set to $\chi = 1/8$, $\vartheta = 0.57$, $\omega = 1/4$, $\sigma = 0.01114$, $\alpha = 1/20$ and $\beta = 0$. These values do not satisfy the assumptions of Theorem 4.2; however, the method executes substantially faster than with parameters for which the theorem applies. All computations were performed in Matlab on a workstation with an AMD AthlonTM 64 X2 5200+ processor and 4GB of memory.

We consider a multilevel discretization in which the a posteriori error estimator from Section 5.2 is used to determine an appropriate discretization level independently for each coefficient. A discretization level j_{μ} , which represents linear finite elements on a uniform mesh with $2^{j_{\mu}}$ cells, is assigned to each index μ with the goal of equidistributing the estimated error among all coefficients. In particular, different refinement levels are used to approximate different coefficients u_{μ} .

In Figure 2, on the left, the errors are plotted against the number of degrees of freedom, which refers to the total number of basis functions used in the discretization, i.e., the sum of $2^{j_{\mu}} - 1$ over all μ . On the right, we plot the errors against an estimate of the computational cost. This estimate takes scalar products, matrix-vector multiplications and linear solves into account. The total number of each of these operations on each discretization level is tabulated during the computation, weighted by the number of degrees of freedom on the discretization level,

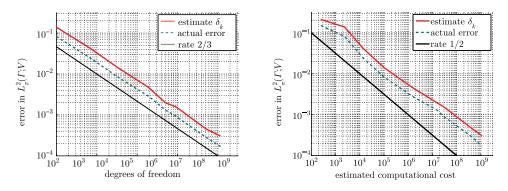


FIGURE 2. Convergence of SolveGalerkin_{A,f}.

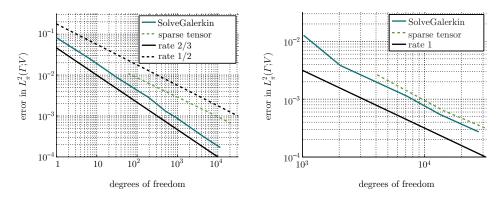


FIGURE 3. Comparison of SolveGalerkin_{A,f} and the sparse tensor construction, for a multilevel discretization (left) and with a fixed finite element mesh (right).

and summed over all levels. The estimate is equal to seven times the resulting sum for linear solves, plus three times the value for matrix-vector multiplications, plus the sum for scalar products. These weights were determined empirically by timing the operations for tridiagonal sparse matrices in Matlab.

The errors were computed by comparison with a reference solution, which has an error of approximately $5 \cdot 10^{-5}$. The plots show that the error bounds δ_k are good approximations of the actual error, and only overestimate it by a small factor.

We compare the discretizations generated adaptively by **SolveGalerkin**_{A,f} with the heuristic a priori adapted sparse tensor product construction from [5]. Using the notation of [26, Section 4], we set $\gamma = 2$ and $\eta_m = 1/(r_m + \sqrt{1 + r_m^2})$ for $r_m = cm^2/2$ and c from (5.19). These values are similar to those used in the computational examples of [5]. The coarsest spatial discretization used in the sparse tensor product contains 16 elements.

In order to isolate the stochastic discretization, we also consider a fixed spatial discretization, using linear finite elements on a uniform mesh of (0,1) with 1024 elements to approximate all coefficients. This mesh is sufficiently fine such that the finite element error is negligible compared to the total error. We refer to these simpler versions of the numerical methods as single level discretizations.

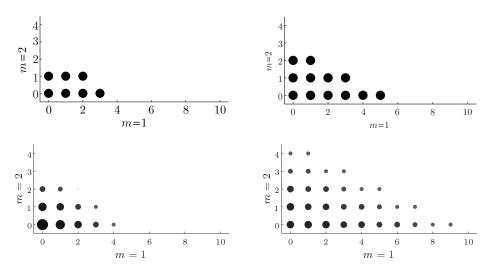


FIGURE 4. Slices of index sets generated by SolveGalerkin_{A,f} (left) and [5] (right) with single level discretization (top) and multilevel discretization (bottom). All sets correspond to the rightmost points in Figure 3. Active indices with support in $\{1, 2\}$ are plotted; the level of the finite element discretization is proportional to the radius of the circle.

The single level versions of SolveGalerkin_{A,f} and the sparse tensor method construct discretizations of equal quality, with only a slight advantage for the adaptive algorithm. However, with a multilevel discretization, SolveGalerkin_{A,f} converges faster than the sparse tensor method, with respect to the number of degrees of freedom. At least in this example, the adaptively constructed discretizations are more efficient than sparse tensor products.

As index sets $\Xi \subset \Lambda$ are infinite dimensional in the sense that they can contain indices of arbitrary length, they are difficult to visualize in only two dimensions. In Figure 4, we plot two-dimensional slices of sets generated by SolveGalerkin_{A,f} and the sparse tensor construction from [5]. We consider only those indices which are zero in all dimensions after the second, and plot their values in the first two dimensions. The upper plots depict index sets generated using single level discretizations; dots refer to active indices. The lower plots illustrate the discretizations generated with multilevel finite element discretizations. The radii of the circles are proportional to the discretization level.

The bottom two plots in Figure 4 illustrate differences between the discretizations generated by $SolveGalerkin_{A,f}$ and the sparse tensor construction. The former has many fewer active indices, but higher discretization levels for some of these. For example, the coefficient of the constant polynomial is approximated on meshes with 4096 and 256 elements, respectively. Also, while the sets constructed by sparse tensorization appear triangular in this figure, the adaptively generated index sets are somewhat more convex. All of the sets are anisotropic in the sense that the first dimension is discretized more finely than the second.

We use the convergence curves in Figures 2 and 3 to empirically determine convergence rates of SolveGalerkin_{A,f}. The convergence rate with respect to the

total number of degrees of freedom is 2/3, which is faster than the approximation rate of 1/2 shown in [11, 10]. It also compares favorably to the sparse tensor construction, which converges with rate 1/2. However, when considering convergence with respect to the computational cost, the rate of SolveGalerkin_{A,f} reduces to 1/2 also. We suspect that this is due to the approximation of the residual, which is performed on a larger set of active indices than the subsequent approximation of the Galerkin projection.

For the case of a single finite element mesh, [11, 10] show an approximation rate of 3/2, wheras we observe a rate of 1 for both SolveGalerkin_{A,f} and sparse tensorization. In principle, it is possible that SolveGalerkin_{A,f} does not converge with the optimal rate in this example, since the parameters used in the computations do not satisfy the assumptions of Theorem 4.2. Alternatively, due to large constants in the approximation estimates, the asymptotic rate may not be perceivable for computationally accessible tolerances.

CONCLUSION

The adaptive method SolveGalerkin_{A,f} efficiently constructs Galerkin spaces and approximations of the corresponding Galerkin projections for elliptic boudary value problems with random coefficients. It is proven to converge, and provides a reliable and efficient bound for the mean square error. In the case of a fixed spatial discretization, the Galerkin subspaces are shown to be optimal, and the algorithm has linear complexity with respect to the number of active polynomial modes, up to a logarithmic term in the computation of the Galerkin projection.

This solver has a modular structure, which allows any discretization of the spatial domain. For a model problem, we consider finite elements with a residual-based a posteriori error estimator. A minor modification of standard estimators is needed to account for finite element functions in the source term.

Numerical computations show that adaptively computed approximate solutions can be sparser than a sparse tensor product construction. Convergence with respect to the total number of degrees of freedom or the total computational cost agrees with or surpasses approximation estimates shown by nonconstructive means in the case of a multilevel spatial discretization.

References

- Mark Ainsworth and J. Tinsley Oden, A posteriori error estimation in finite element analysis, Pure and Applied Mathematics (New York), Wiley-Interscience [John Wiley & Sons], New York, 2000. MR1885308 (2003b:65001)
- Ivo M. Babuška, Raúl Tempone, and Georgios E. Zouraris, Galerkin finite element approximations of stochastic elliptic partial differential equations, SIAM J. Numer. Anal. 42 (2004), no. 2, 800–825 (electronic). MR2084236 (2005h:65012)
- A. Barinka, Fast evaluation tools for adaptive wavelet schemes, Ph.D. thesis, RWTH Aachen, March 2005.
- 4. Heinz Bauer, *Wahrscheinlichkeitstheorie*, Fifth ed., de Gruyter Lehrbuch. [de Gruyter Textbook], Walter de Gruyter & Co., Berlin, 2002. MR1902050 (2003b:60001)
- Marcel Bieri, Roman Andreev, and Christoph Schwab, Sparse tensor discretization of elliptic SPDEs, SIAM J. Sci. Comput. 31 (2009/10), no. 6, 4281–4304. MR2566594
- Marcel Bieri and Christoph Schwab, Sparse high order FEM for elliptic sPDEs, Comput. Methods Appl. Mech. Engrg. 198 (2009), no. 37-40, 1149–1170. MR2500242 (2010g:65205)
- Peter Binev, Wolfgang Dahmen, and Ronald A. DeVore, Adaptive finite element methods with convergence rates, Numer. Math. 97 (2004), no. 2, 219–268. MR2050077 (2005d:65222)

- Albert Cohen, Wolfgang Dahmen, and Ronald A. DeVore, Adaptive wavelet methods for elliptic operator equations: convergence rates, Math. Comp. 70 (2001), no. 233, 27–75 (electronic). MR1803124 (2002h:65201)
- 9. _____, Adaptive wavelet methods. II. Beyond the elliptic case, Found. Comput. Math. 2 (2002), no. 3, 203–245. MR1907380 (2003f:65212)
- Albert Cohen, Ronald DeVore, and Christoph Schwab, Analytic regularity and polynomial approximation of parametric and stochastic elliptic PDE's, Anal. Appl. (Singap.) 9 (2011), no. 1, 11–47. MR2763359
- Albert Cohen, Ronald A. DeVore, and Christoph Schwab, Convergence rates of best N-term Galerkin approximations for a class of elliptic sPDEs, Found. Comput. Math. 10 (2010), no. 6, 615–646. MR2728424
- Stephan Dahlke, Massimo Fornasier, and Thorsten Raasch, Adaptive frame methods for elliptic operator equations, Adv. Comput. Math. 27 (2007), no. 1, 27–63. MR2317920 (2008c:65366)
- Stephan Dahlke, Thorsten Raasch, Manuel Werner, Massimo Fornasier, and Rob Stevenson, Adaptive frame methods for elliptic operator equations: the steepest descent approach, IMA J. Numer. Anal. 27 (2007), no. 4, 717–740. MR2371829 (2008i:65239)
- Manas K. Deb, Ivo M. Babuška, and J. Tinsley Oden, Solution of stochastic partial differential equations using Galerkin finite element techniques, Comput. Methods Appl. Mech. Engrg. 190 (2001), no. 48, 6359–6372. MR1870425 (2003g:65009)
- Ronald A. DeVore, Nonlinear approximation, Acta Numerica, 1998, Acta Numer., vol. 7, Cambridge Univ. Press, Cambridge, 1998, pp. 51–150. MR1689432 (2001a:41034)
- Tammo Jan Dijkema, Christoph Schwab, and Rob Stevenson, An adaptive wavelet method for solving high-dimensional elliptic PDEs, Constr. Approx. 30 (2009), no. 3, 423–455. MR2558688
- Willy Dörfler, A convergent adaptive algorithm for Poisson's equation, SIAM J. Numer. Anal. 33 (1996), no. 3, 1106–1124. MR1393904 (97e:65139)
- Philipp Frauenfelder, Christoph Schwab, and Radu Alexandru Todor, *Finite elements for elliptic problems with stochastic coefficients*, Comput. Methods Appl. Mech. Engrg. **194** (2005), no. 2-5, 205–228. MR2105161 (2005i:65186)
- Tsogtgerel Gantumur, Helmut Harbrecht, and Rob Stevenson, An optimal adaptive wavelet method without coarsening of the iterands, Math. Comp. 76 (2007), no. 258, 615–629 (electronic). MR2291830 (2008i:65310)
- Claude Jeffrey Gittelson, Adaptive Galerkin methods for parametric and stochastic operator equations, Ph.D. thesis, ETH Zürich, 2011, ETH Dissertation No. 19533.
- 21. _____, Adaptive stochastic Galerkin methods: Beyond the elliptic case, Tech. Report 2011-12, Seminar for Applied Mathematics, ETH Zürich, 2011.
- 22. _____, Stochastic Galerkin approximation of operator equations with infinite dimensional noise, Tech. Report 2011-10, Seminar for Applied Mathematics, ETH Zürich, 2011.
- Hermann G. Matthies and Andreas Keese, Galerkin methods for linear and nonlinear elliptic stochastic partial differential equations, Comput. Methods Appl. Mech. Engrg. 194 (2005), no. 12-16, 1295–1331. MR2121216 (2005j:65146)
- A. Metselaar, Handling wavelet expansions in numerical methods, Ph.D. thesis, University of Twente, 2002. MR2715507
- Pedro Morin, Ricardo H. Nochetto, and Kunibert G. Siebert, *Data oscillation and convergence of adaptive FEM*, SIAM J. Numer. Anal. **38** (2000), no. 2, 466–488 (electronic). MR1770058 (2001g:65157)
- Christoph Schwab and Claude Jeffrey Gittelson, Sparse tensor discretization of highdimensional parametric and stochastic PDEs, Acta Numerica, Acta Numer., vol. 20, Cambridge Univ. Press, Cambridge, 2011, pp. 291–467. MR2805155
- Rob Stevenson, Adaptive solution of operator equations using wavelet frames, SIAM J. Numer. Anal. 41 (2003), no. 3, 1074–1100 (electronic). MR2005196 (2004e:42062)
- Radu Alexandru Todor and Christoph Schwab, Convergence rates for sparse chaos approximations of elliptic problems with stochastic coefficients, IMA J. Numer. Anal. 27 (2007), no. 2, 232–261. MR2317004 (2008b:65016)
- R. Verfürth, A Review of a Posteriori Error Estimation and Adaptive Mesh-Refinement Techniques, Teubner Verlag and J. Wiley, Stuttgart, 1996.

- Xiaoliang Wan and George Em Karniadakis, An adaptive multi-element generalized polynomial chaos method for stochastic differential equations, J. Comput. Phys. 209 (2005), no. 2, 617–642. MR2151997 (2006e:65007)
- Multi-element generalized polynomial chaos for arbitrary probability measures, SIAM J. Sci. Comput. 28 (2006), no. 3, 901–928 (electronic). MR2240796 (2007d:65008)
- Dongbin Xiu and George Em Karniadakis, The Wiener-Askey polynomial chaos for stochastic differential equations, SIAM J. Sci. Comput. 24 (2002), no. 2, 619–644 (electronic). MR1951058 (2003m:60174)

Seminar for Applied Mathematics, ETH Zurich, Rämistrasse 101, CH-8092 Zurich, Switzerland

 $Current \ address:$ Department of Mathematics, Purdue University, 150 N. University Street, West Lafayette, Indiana 47907

E-mail address: cgittels@purdue.edu