

An Adaptive Transport Protocol for Multimedia Communication

Dane Dwyer, Sungwon Ha, Jia-Ru Li, Vaduvur Bharghavan University of Illinois at Urbana-Champaign

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Outline

- Introduction
- HPF
 - Window management
 - Application interface
- Measurements
- Conclusions



Introduction

Motivation

- User traffic become increasingly multimedia oriented and *heterogeneous*
- Neither of TCP and UDP is well suited for multimedia flows



Introduction (cont.)

Goal

- Support multiple interleaved reliable and unreliable data sub-streams
- Decouple the congestion control and reliability mechanisms
- Use application-defined priorities for the link scheduler to drop low-priority packets during congestion



HPF

- HPF (Heterogeneous Packet Flow)
 - Window management
 - Application interface



Window management

• Similar with TCP

 Use or enhance the mechanisms that are provided by TCP for flow control, reliability, and sequencing

• Difference

- Congestion control mechanism
- Support interleaved reliable and unreliable packet sub-streams

- Flow control, reliability, and sequencing
 - Packet may have been either a low or a high priority packet
 - —low priority: the dropped packet be ignored
 - —high priority: should be retransmitted
 - Each packet needs to identify the previous high priority packet (called *prevHIGH*)



Example

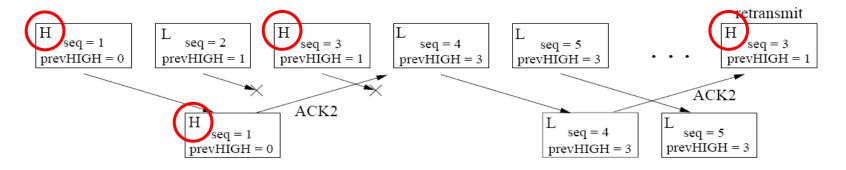


Figure 1. Example of acknowledgements in HPF. Only high priority packets that are lost will be retransmitted. When the receiver gets the retransmitted high priority packet with seq=3, it will send ACK6 (rather than ACK2). The lost low priority packet with seq=2 is implicitly acknowledged and will not be retransmitted.



- Congestion control
 - The sender estimates the congestion window based on the fraction of received packets in the current window



- Compute the fraction of received packets
 - Use a congestion estimation window
 - The sender maintains epoch and cwnd
 - The receiver maintains sender.epoch and sender.cwnd
 - Each packet has two fields packet.epoch and packet.cwnd



- if packet.epoch == sender.epoch received packet counter ++
- if packet.epoch < sender.epoch the congestion information is ignored
- if packet.epoch > sender.epoch
 the receiver starts a new estimation window received packet counter = 0
 sender.epoch = packet.epoch
 sender.cwnd = packet.cwnd



- Updating the congestion window
 - if cong_ack.fraction > α cwnd = f_I (cwnd)
 - if cong_ack.fraction $\leq \alpha$ cwnd = f_D (ack.cwnd, ack.fraction)
 - if a timeout occursssthresh = cwnd/2

 α : (1- random packet loss probability)



Application interface

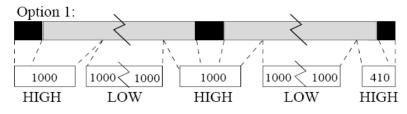
- The application must to signal the priority level
 - Maximize throughput
 - Merge reliable and unreliable data into a packet
 - Maximize adaptation
 - Merge only 'like-priority' data bytes
 - Loss based
 - Based on the fraction of received packets

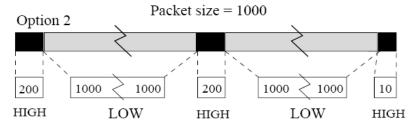


Application interface (cont.)

Sender

- 1 send(sock, buffer, 200, HIGH);
- 2 send(sock, buffer, 4000, LOW);
- 3 send(sock, buffer, 200, HIGH);
- 4 send(sock, buffer, 4000, LOW);
- 5 send(sock, buffer, 10, HIGH);

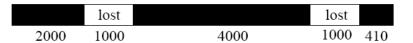




Receiver

Option 1

- 1 read(sock, buffer, 3000) returns 2000, last read status = 1
- 2 read(sock, buffer, 3000) returns 1000, last read status = -1
- 3 read(sock, buffer, 3000) returns 3000, last read status = 1
- 4 read(sock, buffer, 3000) returns 1000, last read status = 1
- 5 read(sock, buffer, 3000) returns 1000, last_read_status = -1
- 6 read(sock, buffer, 3000) returns 410, last read status = 1



Option 2

	lost		lost	
1200	1000	3200	1000	2010



Measurements

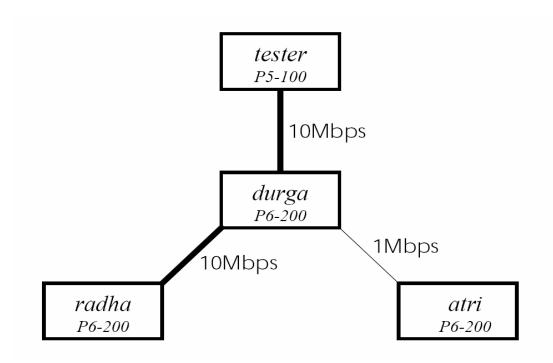


Figure 3. The experimental testbed configuration used for the performance tests.

Measurements (cont.)

		High:Low	Packets	Improvement
	Protocol	Ratio	Dropped	vs. TCP
_	TCP		0%	
	HPF	1:0	0%	-4.0%
		7:3	1.2%	2.2%
		5:5	5.1%	16.7%
		3:7	5.5%	21.3%

Table 1. The performance of HPF vs TCP at various priority ratios and bursty UDP traffic

Measurements (cont.)

	High:Low	Packets	Improvement
Protocol	Ratio	Dropped	vs TCP
TCP		0%	
HPF	1:0	0%	-4.3%
	7:3	1.1%	8.0%
	5:5	5.3%	29.7%
	3:7	5.7%	32.5%

Table 2. The performance of HPF vs TCP at various priority ratios with multiple concurrent streams



Conclusions

- Support multiple interleaved reliable and unreliable data sub-streams
- Decouple the congestion control and reliability mechanisms
- Use application-defined priorities for the link scheduler to drop low-priority packets during congestion

