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An algorithm for the permanent of circulant matrices

Abstract

The permanent of an $n \times n$ matrix $A = (a_{ij})$ is the matrix function $(1) \text{ per } A = \sum_{\pi \in S_n} a_{1\pi(1)} \cdots a_{n\pi(n)}$ where the summation is over all permutations in the symmetric group, S_n . An $n \times n$ matrix A is a circulant if there are scalars a_0, \dots, a_{n-1} such that $(2) A = \sum_{j=0}^{n-1} a_j P^j$ where P is the $n \times n$ permutation matrix corresponding to the cycle $(12 \dots n)$ in S_n . In general the computation of the permanent function is quite difficult chiefly because it is not invariant under addition of a multiple of one row to another. Using the principle of "inclusion and exclusion", Ryser [6, p. 27] gave an expansion for the permanent. Also the Laplace expansion is available for the permanent [2, p. 20]. Neither of these methods are particularly efficient. In [4] Mine considered the permanents of matrices with entries either 0 or 1. Mine also studied tridiagonal circulants in [5]. Metropolis, Stein, and Stein [3] have given recurrence relations for evaluating the permanents of circulant matrices (2) where the first k scalars are 1 and the remaining ones are 0. Permanents of circulant matrices were also studied by Tinsley [7].

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AN ALGORITHM FOR THE PERMANENT OF CIRCULANT MATRICES

BY
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AND
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1. Introduction. The *permanent* of an $n \times n$ matrix $A = (a_{ij})$ is the matrix function

$$(1) \quad \text{per } A = \sum_{\pi \in S_n} a_{1\pi(1)} \cdots a_{n\pi(n)}$$

where the summation is over all permutations in the symmetric group, S_n . An $n \times n$ matrix A is a *circulant* if there are scalars a_1, \dots, a_n such that

$$(2) \quad A = \sum_{i=1}^n a_i P^{i-1}$$

where P is the $n \times n$ permutation matrix corresponding to the cycle $(12 \cdots n)$ in S_n . In general the computation of the permanent function is quite difficult chiefly because it is not invariant under addition of a multiple of one row to another. Using the principle of "inclusion and exclusion", Ryser [6, p. 27] gave an expansion for the permanent. Also the Laplace expansion is available for the permanent [2, p. 20]. Neither of these methods are particularly efficient. In [4] Minc considered the permanents of matrices with entries either 0 or 1. Minc also studied tridiagonal circulants in [5]. Metropolis, Stein, and Stein [3] have given recurrence relations for evaluating the permanents of circulant matrices (2) where the first k scalars are 1 and the remaining ones are 0. Permanents of circulant matrices were also studied by Tinsley [7].

2. The algorithm. If we consider the scalars as indeterminates over an underlying field every term of the permanent (1) of a circulant matrix (2) is a monomial in the scalars a_1, \dots, a_n . Our algorithm deletes appropriate monomials from the set of all n^n such monomials until only those appearing in the permanent remain. This is easily programmed because the monomials need only be considered one at a time and may be indexed by the n^n n -tuples chosen from $1, \dots, n$ and ordered lexicographically. It is convenient to state the algorithm in terms of these indices.

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Algorithm. If $I = (i_1, \dots, i_n)$ is an n -tuple with entries chosen from $1, \dots, n$ then discard I if

$$(i) \sum_{j=1}^n i_j \not\equiv 0 \pmod{n},$$

or if

$$(ii) i_{j+k} \equiv i_j - k \pmod{n} \text{ for any } k \text{ and } j = 1, \dots, n-1.$$

Condition (ii) excludes the occurrence of terms in the permanent of (2) with the following pattern

$$(3) \quad \dots a_i \underbrace{\dots}_{k-2 \text{ entries}} a_{i+k+1} \dots$$

where a_{i+n} is considered to be a_i if necessary. For example, if $n = 4$ condition (ii) of the algorithm discards a monomial whenever one of the following patterns occurs:

$$\begin{aligned} & \dots 14 \dots \dots 21 \dots \dots 32 \dots \dots 43 \dots \\ & \cdot 1^*3 \cdot, \cdot 2^*4 \cdot, \cdot 3^*1 \cdot, \cdot 4^*2 \cdot \\ & 1^{**}2, 2^{**}3, 3^{**}4, 4^{**}1. \end{aligned}$$

Condition (i) leaves the following 4-tuples:

1111	1214	1313	1412	<u>2114</u>	<u>2213</u>	2312	2411
1124	1223	<u>1322</u>	<u>1421</u>	<u>2123</u>	2222	<u>2321</u>	2424
<u>1133</u>	<u>1232</u>	<u>1331</u>	<u>1434</u>	<u>2132</u>	2231	2334	<u>2433</u>
<u>1142</u>	1241	1344	<u>1443</u>	<u>2141</u>	<u>2244</u>	<u>2343</u>	<u>2442</u>
<u>3113</u>	<u>3212</u>	<u>3311</u>	<u>3414</u>	4112	<u>4211</u>	<u>4314</u>	4413
3122	<u>3221</u>	<u>3324</u>	3423	<u>4121</u>	<u>4224</u>	<u>4323</u>	<u>4422</u>
3131	<u>3234</u>	3333	<u>3432</u>	4134	4233	<u>4332</u>	<u>4431</u>
<u>3144</u>	<u>3243</u>	3342	3441	<u>4143</u>	4242	<u>4341</u>	4444

Condition (ii) eliminates all of the above 4-tuples which are underlined.

Hence, if $n = 4$ the permanent of (2) will be

$$\sum_{i=1}^4 a_i^4 + 2a_1^2 a_3^2 + 2a_2^2 a_4^2 + 4 \sum_{i=1}^4 a_i^2 a_{i+1} a_{i+3}.$$

Let R_n denote the set of n -tuples left by the algorithm. We remark that the n -tuples in R_n need not be formally distinct; e.g., 1313 and 3131 are both in R_4 . The number of formally distinct diagonal products in the permanent of an arbitrary circulant has been determined by Brualdi and Newman [1].

3. Proofs

THEOREM. *Let A be a circulant matrix (2) with scalars a_1, \dots, a_n . Then*

$$\text{per } A = \sum a_{i_1} \cdots a_{i_n}$$

where the summation is over all $(i_1, \dots, i_n) \in R_n$.

Proof. We are concerned with determining conditions for which $a_{i_1} \cdots a_{i_n}$ is a term of the permanent of the $n \times n$ matrix (2). Thus, a_{i_k} always denotes an element of the k th row of (2). The i th column of (2) is

$$\begin{bmatrix} a_i \\ a_{i-1} \\ \vdots \\ a_{i-n+1} \end{bmatrix}$$

where subscripts are taken modulo n . If the Laplace expansion along the first row is used to find $\text{per } A$ the entry a_{i-k+1} cannot be chosen from row k to appear in any monomial beginning with a_i . In any monomial of the permanent the pattern (3) cannot appear since we may expand along any row.

Therefore any (i_1, \dots, i_n) in R_n satisfies

$$i_{j+k} \neq i_j - k \quad \text{for } k = 1, \dots, n-1.$$

Again, subscripts are taken modulo n when necessary.

Write $i_{j+k} = i_j - k + x_{jk} \pmod{n}$ where $x_{jk} \neq 0$, $1 \leq x_{jk} \leq n-1$, and $k \neq 0$. We would like to show that $s \neq t$ implies $x_{js} \neq x_{jt}$.

Suppose $x_{js} = x_{jt}$. Then

$$x_{js} = i_{j+s} - i_j + s = i_{j+t} - i_j + t = x_{jt}.$$

Hence

$$i_{j+s} = i_{j+t} - (s - t),$$

but unless $s = t$

$$i_{j+s} = i_{j+t+(s-t)} \neq i_{j+t} - (s - t).$$

So assuming $x_{js} = x_{jt}$ leads to a contradiction. Hence the contrapositive is true and $s \neq t$ implies $x_{js} \neq x_{jt}$.

Step (i) is included in the algorithm because it is easy to implement. In fact, (ii) implies (i) as we now show:

$$\begin{aligned} \sum_{k=0}^{n-1} i_{j+k} &= i_j + \sum_{k=1}^{n-1} i_{j+k} = \left(i_j + \sum_{k=1}^{n-1} (i_j - k + x_{jk}) \right) \pmod{n} \\ &= \left(ni_j - \sum_{k=1}^{n-1} k + \sum_{k=1}^{n-1} x_{jk} \right) \pmod{n} \\ &= (ni_j - \frac{1}{2}n(n-1) + \frac{1}{2}n(n-1)) \pmod{n} \\ &\equiv 0 \pmod{n}. \end{aligned}$$

We have shown why the n -tuples mentioned in (i) and (ii) must be discarded. It remains to show that no more should be excluded. Condition (ii) says there are n choices for a_{i_1} , $n-1$ choices for a_{i_2} and in general $n-k+1$ choices for a_{i_k} . That is, condition (ii) does not eliminate exactly $n!$ terms. But there are $n!$ terms in the permanent so precisely the right number of monomials has been excluded.

4. Numerical results. Dr. Joan Cooper wrote a Fortran programme for our algorithm which was implemented on an ICL 1904A at the University of Newcastle, N.S.W., Australia. The following various 7×7 circulants were computed using 2.54 seconds of core time.

	First row of circulant matrix A						per A	row sum of $A = r$	per (A/r)
3	1	1	0	1	0	0	4416	6	0.0157750
1	1	1	0	0	0	0	31	3	0.0141747
1	1	0	0	0	0	0	2	2	0.0156250
1	1	1	1	1	1	1	5040	7	0.0061199
0	1	1	0	1	0	0	24	3	0.0109739
1	1	1	0	1	0	0	144	4	0.0087891
1	1	-1	0	0	0	0	1	1	1.0

We believe the algorithm is not shown to best advantage as most of the elapsed time is due to reading the 7-tuples of the example from disc.

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