

# An enhanced representation of time series which allows fast and accurate classification, clustering and relevance feedback

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## Abstract

We introduce an extended representation of time series that allows fast, accurate classification and clustering in addition to the ability to explore time series data in a relevance feedback framework. The representation consists of piecewise linear segments to represent shape and a weight vector that contains the relative importance of each individual linear segment. In the classification context, the weights are learned automatically as part of the training cycle. In the relevance feedback context, the weights are determined by an interactive and iterative process in which users rate various choices presented to them. Our representation allows a user to define a variety of similarity measures that can be tailored to specific domains. We demonstrate our approach on space telemetry, medical and synthetic data.

## 1.0 Introduction

Time series account for much of the data stored in business, medical, engineering and social science databases. Much of the utility of collecting this data comes from the ability of humans to visualize the *shape* of the (suitably plotted) data, and classify it. For example:

- Cardiologists view electrocardiograms to diagnose arrhythmias.
- Chartists examine stock market data, searching for certain shapes, which are thought to be indicative of a stock's future performance.

Unfortunately, the sheer volume of data collected means that only a small fraction of the data can ever be viewed.

Attempts to utilize classic machine learning and clustering algorithms on time series data have not met with great success. This, we feel, is due to the typically high dimensionality of time series data, combined with the difficulty of defining a similarity measure appropriate for the domain. For an example of both these difficulties, consider the three time series in Figure 1. Each of them contains 29,714 data points, yet together they account for less than .0001 % of the database from which they were extracted. In addition, consider the problem of clustering these three examples. Most people would group A and C

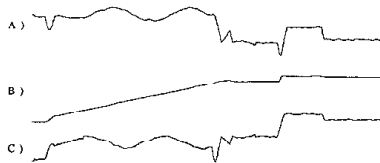


Figure 1: Examples of time series

together, with B as the outgroup, yet common distance measures on the raw data (such as correlation, absolute distance or squared error) group B and C together, with A as the outgroup.

What is needed is a representation that allows efficient computation on the data, and extracts higher order features. Several such representations have been proposed, including Fourier transformations (Faloutsos et al. 1994), relational trees (Shaw & DeFigueiredo, 1990) and envelope matching/R+ trees (Agrawal et al. 1995). The above approaches have all met with some success, but all have shortcomings, including sensitivity to noise, lack of intuitiveness, and the need to fine tune many parameters.

Piece-wise linear segmentation, which attempts to model the data as sequences of straight lines, (as in Figure 2) has innumerable advantages as a representation. Pavlidis and Horowitz (1974) point out that it provides a useful form of data compression and noise filtering. Shatky and Zdonik (1996) describe a method for fuzzy queries on linear (and higher order polynomial) segments. Keogh and Smyth (1997) further demonstrate a framework for probabilistic pattern matching using linear segments.

Although pattern matching using piece-wise linear segments has met with some success, we believe it has a major shortcoming. When comparing two time series to see

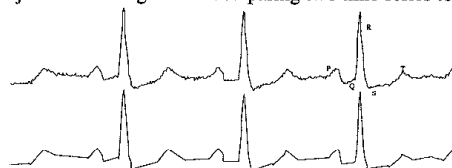


Figure 2: An example of a time series and its piece-wise linear representation

if they are similar, all segments are considered to have equal importance. In practice, however, one may wish to assign different levels of importance to different parts of the time series. As an example, consider the problem of pattern matching with electrocardiograms. If a cardiologist is attempting to diagnose a recent myocardial infarction (MI) she will pay close attention to the S-T wave, and downplay the importance of the rest of the electrocardiogram. If we wish an algorithm to reproduce the cardiologist's ability, we need a representation that allows us to weight different parts of a time series differently.

In this paper, we propose such a representation. We use piece-wise linear segments to represent the shape of a time series, and a weight vector that contains the relative importance of each individual linear segment.

## 2.0 Representation of time series

There are numerous algorithms available for segmenting time series, many of which were pioneered by Pavlidis and Horowitz (1974). An open question is how to best choose  $K$ , the 'optimal' number of segments used to represent a particular time series. This problem involves a trade-off between accuracy and compactness, and clearly has no general solution. For this paper, we utilize the segmentation algorithm proposed in Keogh (1997).

### 2.1 Notation

For clarity we will refer to 'raw', unprocessed temporal data as time series, and a piece-wise representation of a time series as a sequence. We will use the following notation throughout this paper. A time series, sampled at  $k$  points, is represented as an uppercase letter such as  $A$ . The segmented version of  $A$ , containing  $K$  linear segments, is denoted as a bold uppercase letter such as  $\mathbf{A}$ , where  $\mathbf{A}$  is a 5-tuple of vectors of length  $K$ .

$$\mathbf{A} \equiv \{AXL, AXR, AYL, AYR, AW\}$$

The  $i^{\text{th}}$  segment of sequence  $\mathbf{A}$  is represented by the line between  $(AXL_i, AYL_i)$  and  $(AXR_i, AYR_i)$ , and  $AW_i$ , which represents the segments weight. Figure 3 illustrates this notation.

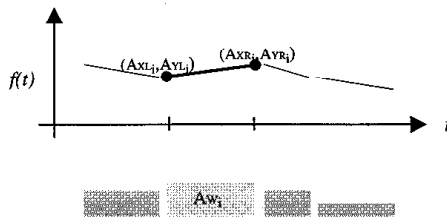


Figure 3: We represent times series by a sequence of straight segments, together with a sequence of weights (shown as the histogram) which contain the relative importance of each segment

After a time series is segmented to obtain a sequence, we initialize all the weights to one. Thereafter, if any of the weights are changed, the weights are renormalized such that the sum of the products of each weight with the length of its corresponding segment, equals the length of the entire sequence, so that the following is always true:

$$\sum_{i=1}^K W_i * (AXR_i - AXL_i) = AXR_k - AXL_k$$

## 2.2 Comparing time series

An advantage in using the piece-wise linear segment representation is that it allows one to define a variety of distance measures to represent the distance between two time series. This is important, because in various domains, different properties may be required of a distance measure. Figure 4 shows some of the various distortions one can encounter in time series.

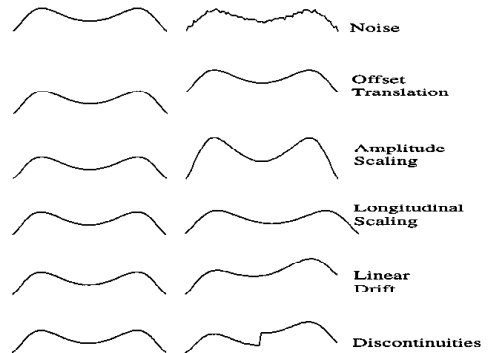


Figure 4: Some of the difficulties encountered in defining a distance measure for time series

It is possible to define distance measures for sequences that are invariant to any subset of the illustrated distortions. For the experiments in this paper we used the simple distance measure given below. This measure is designed to be insensitive to offset translation, linear trends and discontinuities.

$$D(A, B) = \sum_{i=1}^K AW_i * BW_i * |(AYL_i - BYL_i) - (AYR_i - BYR_i)|$$

It is convenient for notational purposes to assume that the endpoints of the two sequences being compared are aligned. In general, with real data, this is not the case. So we will have to process the sequences first, by breaking some segments at a point which corresponds to an endpoint in the other sequence. This can be done in  $O(K)$ .

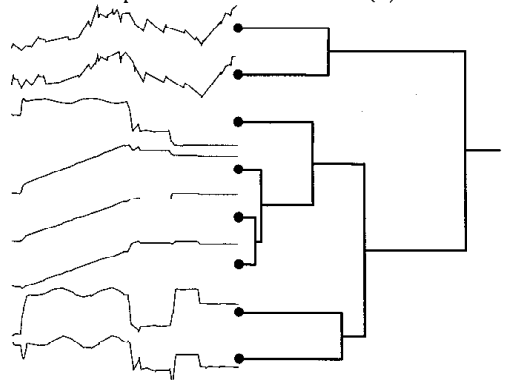


Figure 5: An example of time series hierarchically clustered using our distance measure

### 2.3 Merging time series

In this section we define an operation on sequences which we call 'merge'. The merge operator allows us to combine information from two sequences, and repeated application of the merge operator allows us to combine information from multiple sequences. The basic idea is that the merge operator takes two sequences as input, and returns a single sequence whose shape is a compromise between the two original sequences, and whose weight vector reflects how much corresponding segments in each sequence agree.

When merging sequences one may wish for one of the two input sequences to contribute more to the final sequence than the other does. To accommodate this, we associate a term called 'influence' with each of the input sequences. The influence term associated with a sequence  $S$  is a scalar, denoted as  $SI$ , and may be informally considered a 'mixing weight'. Where the influence term comes from depends on the application and is discussed in detail in sections 3.1 and 4.1 below.

To merge the two sequences  $A$  and  $B$  with influence terms  $AI$  and  $BI$  respectively, we use the following algorithm that creates the sequence  $C$ :

```

if (AI * BI < 0) then      sign = -1
else sign = 1
end
mag = min(|AI|, |BI|) / max(|AI|, |BI|)
scale = max(max(AVL), (AVR)) - min(min(AVL), (AVR))
for i = 1 to K
  CXLi = AXLi
  CXRi = AXRi
  CYLi = ((AVLi * AI) + (BVLi * BI)) / (AI + BI)
  CYRi = ((AVRi * AI) + (BVRi * BI)) / (AI + BI)
  run = AXRi - AXLi
  rise = |(AVLi - BVLi) - (AVRi - BVRi)|
  d = (rise / run) * scale
  CWi = (AWi * BWi) * (1 + (sign * mag) / (1 + d))
end
CW = normalize(CW)

```

Table 1: The merge algorithm

### 2.4 Learning prototypes

Although the merge operator is designed to be a component in the more sophisticated algorithms presented below, it can, by itself, be considered a simple learning algorithm that creates a prototypical sequence. Creating a prototype solely from positive examples works in the following manner. We have a model sequence  $A$ , which is a typical example of a class of sequences. If we merge sequence  $A$  with sequence  $B$ , another member of the same class, the resultant sequence  $C$  can be considered a more general model for the class. In particular the differences in shape are *reduced* by averaging, and the weights for similar segments are *increased*.

In contrast, creating a prototype from both positive and negative examples uses a negative influence for the negative examples. As before, suppose we have a sequence  $A$ , which is an example of one class of sequences. However, suppose  $B$  is an example of a different class. If

we merge  $A$  with  $B$ , using a negative influence term for  $B$ , the resultant sequence  $C$  is a new prototype for  $A$ 's class where the differences in shape between  $A$  and  $B$  are *exaggerated* and the weights for similar segments are *decreased*.

The above maps neatly to our intuitions. If we are learning a prototype for a class of sequences from a set of positive examples, we want the shape learned to be an average of all the examples, and we want to increase the weight of segments that are similar, because those segments are characteristic of the class. If however, we are trying to learn from a negative example, we want to exaggerate the differences in shape between classes, and decrease the weight of similar segments, because segments that are similar across classes have no discriminatory power.

### 3.0 Relevance feedback

Relevance feedback is the reformulation of a search query in response to feedback provided by the user for the results of previous versions of the query. It has an extensive history in the text domain, dating back to Rocchio's classic paper (1971). However, no one has attempted explore time series databases in a relevance feedback framework, in spite of the fact that relevance feedback has been shown to significantly improve the querying process in text databases (Salton & Buckley, 1990). In this section we present a simple relevance feedback algorithm which utilizes our representation and we demonstrate it on a synthetic dataset.

Our relevance feedback algorithm works in the following manner. An initial query sequence  $Q$  is used to rank all sequences in the database (this query may be hand drawn by the user). Only the best  $n$  sequences are shown to the user. The user assigns influences to each of  $n$  sequences. A positive influence is given to sequences that the user approves of. Negative influences are given to sequences that the user finds irrelevant.

The relative magnitude of influence terms reflects how strongly the user feels about the sequences. So if a user "likes"  $S_i$  twice as much as  $S_j$ , he can assign influences of  $1/2$  or  $2/1$  etc. The sequences are then merged to produce a new query, and the process can be repeated as often as desired.

$$Q_{new} = \text{merge}(\{Q_{old}, Q_{old}I\}, \{S_1, S_1I\}, \{S_2, S_2I\}, \dots, \{S_n, S_nI\})$$

$$Q_{new}I = Q_{old}I + S_1I + S_2I + \dots + S_nI$$

### 3.1 Experimental results

To test the above algorithm, we conducted the following experiment. We constructed 500 "Type A", and 500 "Type B" time series, which are defined as follows:

- Type A:  $\text{Sin}(x^3)$  normalized to be between zero and one, plus Gaussian noise with  $\sigma = .1$   $-2 \leq x \leq 2$
- Type B:  $\text{Tan}(\text{Sin}(x^3))$  normalized to be between zero and one, plus Gaussian noise with  $\sigma = .1$   $-2 \leq x \leq 2$

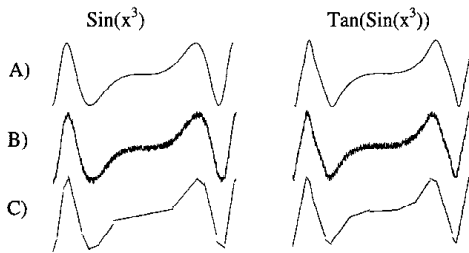


Figure 6: Synthetic data created for relevance feedback experiment

- A) The original time series.
- B) The original time series with noise added.
- C) The segmented version of the time series.

Twenty-five experimental runs were made. Each run consisted of the following steps. A coin toss decided whether Type A or Type B was to be the “target” shape (that is, the shape to be considered “relevant” for that particular experiential run). The initial query was made, and the quality of the ranked sequences was measured as defined below. The best 15 sequences were shown to the user, who then rated them by assigning influences that reflected how closely he thought they resembled the target shape. A new query was built and the search/rate process was repeated twice more.

We evaluated the effectiveness of the approach by measuring the average precision of the top 15 sequences, and the precision at the 25, 50 and 75 percent recall points. Precision ( $P$ ) is defined as the proportion of the returned sequences which are deemed relevant, and recall ( $R$ ) is defined as the proportion of relevant items which are retrieved from the database. These results are shown in Table 2.

In order to see if the ability to assign negative influence terms is helpful, we did the following. For each experimental run we also built queries which contained just the feedback from the sequences judged relevant. These results are shown in parentheses in Table 2.

	Initial Query	Second Query	Third Query
$P$ of top 15	.51	.91 (.68)	.97 (.72)
$P$ at 25% $R$	.52	.91 (.69)	.96 (.71)
$P$ at 50% $R$	.49	.89 (.66)	.95 (.69)
$P$ at 75% $R$	.51	.87 (.63)	.95 (.68)

Table 2: Results of relevance feedback experiments. The values recorded in parentheses are for the queries built just from positive feedback

As one might expect, the initial query (which does not have any user input) returns the sequences in essentially random order. The second query produces remarkable improvement, and the third query produces near perfect ranking. The queries built from just positive feedback do produce improvement, but are clearly inferior to the more general method. This demonstrates the utility of learning from both positive and negative instances.

## 4.0 Classification

Given our representation, an obvious approach to classification is to merge all positive instances in the training set, and use the resultant sequence as a template to which the instances to be classified are compared. This may work in some circumstances, but in some domains it may totally fail. The reason for the potential failure is that there may be two or more distinct shapes that are typical of the class.

To avoid this problem, an algorithm needs to be able to detect the fact that there are multiple prototypes for a given class, and classify an unlabeled instance to that class if it is sufficiently similar to *any* prototype. In the next section we describe such an algorithm, which we call CTC (Cluster, Then Classify).

### 4.1 Classification algorithm

Table 3 shows an outline of our learning algorithm. The input is  $S$ , the set of  $n$  sequences that constitute the training data. The output is  $P$ , a set of sequences, and a positive scalar  $\epsilon$ .

An unseen sequence  $U$  will be classified as a member of a class if and only if, the distance between  $U$  and at least one member of  $P$  is less than  $\epsilon$ .

The algorithm clusters the positive examples using the group average agglomerative method as follows. The algorithm begins by finding the distance between each negative instance in  $S$ , and its closest positive example. The mean of all these distances,  $neg-dis_1$ , is calculated. Next the distance between each positive instance in  $S$ , and the most similar positive example is calculated. The mean of all these distances,  $pos-dis_1$ , is calculated. The fraction  $q_1 = pos-dis_1 / neg-dis_1$  can now be calculated.

At this point, the two closest positive examples are replaced with the result of merging the two sequences, and the process above is repeated to find the fraction  $q_2 = pos-dis_2 / neg-dis_2$ . The entire process is repeated until a single sequence remains. Figure 7 shows a trace through this part of the algorithm for a dataset that contains just 4 positive instances. The set of sequences returned is the set for which  $q_i$  is minimized.

The  $\epsilon$  returned is  $(pos-dis_i + neg-dis_i) / 2$ .

Let  $P_i$  be the set of all positive instances in  $S$   
Cluster all sequences in  $P_i$

for  $i = 1$  to  $n$

$neg-dis_i =$  mean distance between all negative instances and their closest match in  $P_i$

$pos-dis_i =$  mean distance between all positive instances and their closest match in  $P_i$

$q_i = pos-dis_i / neg-dis_i$

Let  $A$  and  $B$  be the closest pair of sequences in  $P_i$

$C = merge([A, AI], [B, BI])$

$CI = AI + BI$

Remove  $A$  and  $B$  from  $P_i$

Add  $C$  to  $P_i$

end

let best equal the  $i$  for which  $q_i$  is minimized

return  $P_{best}$ ,  $(pos-dis_{best} + neg-dis_{best}) / 2$

Table 3: The CTC learning algorithm

## 4.2 Experimental results

To test the algorithm presented above we ran experiments on the following datasets.

- **Shuttle:** This dataset consists of the output of 109 sensors from the first eight-hours of Space Shuttle mission STS-067. The sensors measure a variety of phenomena. Eighteen of them are Inertia Movement Sensors, measured in degrees (examples are shown in Figures 1 and 7). The task is to distinguish these from the other 91 sensors.
- **Heart:** This dataset consists of RR intervals obtained from Holter ECG tapes sampled at 128 Hz (Zebrowski 1997). The data is in a long sequence that contains 96 ventricular events. We extracted the 96 events, and 96 other sections, of equal length, chosen at random.

For comparison purposes we evaluated 4 algorithms on the two datasets described above. CTC is the algorithm described in section 4.1. CTC<sub>uw</sub> is the same algorithm with the weight feature disabled (we simply hardcoded the weights to equal one). NN is a simple nearest neighbor algorithm that uses the raw data representation of the time series. An unlabeled instance is assigned to the same class as its closest match in the training set. We used absolute error as a distance measure, having empirically determined that it was superior to the other obvious candidates (i.e. squared error, correlation). NNS is the same algorithm as NN except it uses the sequence representation and the distance measure defined in section 2.2.

We ran stratified 10 fold cross validation once. All algorithms were trained and tested on exactly the same folds. The results are presented in table 4.

	CTC	CTC <sub>uw</sub>	NN	NNS	Default
Shuttle	100	96.0	82.1	89.3	83.5
Heart	85.6	69.1	59.8	64.6	50.0

Table 4: Experiment results of classification experiments

On both datasets the CTC algorithm performs the best. Its ability to outperform CTC<sub>uw</sub> seems to be a justification of our weighted representation. On the Shuttle dataset NN performs at the base rate. We surmise this is probably due to its sensitivity to discontinuities, which are a trait of this dataset. NNs ability to do significantly better supports this hypothesis.

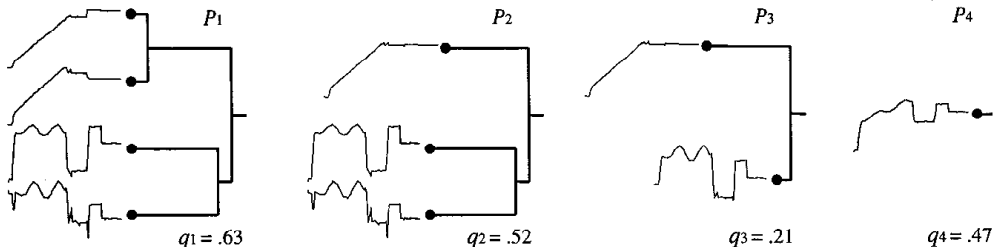


Figure 7: A trace through the CTC algorithm on a small dataset. The first set of prototypes  $P_1$ , are too specific and do not generalize to the test set. The final set  $P_4$  contains a single sequence that is too general, because it is trying to model two distinct shapes. The set  $P_3$  is the best compromise.

## 5.0 Related work

There has been no work on relevance feedback for time series. However, in the text domain there is an active and prolific research community. Salton and Buckley (1990) provide an excellent overview and comparison of the various approaches.

## 6.0 Conclusions

We introduced a new enhanced representation of time series and empirically demonstrated its utility for clustering, classification and relevance feedback.

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**Related Work.** Brin et al. (Brin, Motwani, & Silverstein 1997) used the chi-squared test to look for correlated associations, but did not take into account the number of hypotheses were being tested. Piatetsky-Shapiro (Piatetsky-Shapiro 1991) had a similar idea when he argued that a rule  $X \Rightarrow Y$  is not interesting if  $\text{support}(X \Rightarrow Y) \approx \text{support}(X) \times \text{support}(Y)$ , but again did not consider the number of hypotheses.

## 2. Statistical Significance of Association Rules

In this section, we discuss issues of statistical significance of a set of association rules. First, in Section 2.1, we discuss the statistical significance of a single rule. However, when we analyze many rules simultaneously, the significance test has to take into account the number of hypotheses being tested. This is the so-called multiple comparisons problem (Hochberg & Tamhane 1987). Typically the test statistics corresponding to the hypotheses being tested are not independent. It is important to observe that the number of hypotheses implicitly being tested may be much greater than the number of output rules; we give an upper bound for this number in Section 2.2. This bound may, in general, be too conservative. We offer a practical way of dealing with this problem in Section 2.3; the idea is to use resampling to determine a good acceptance threshold. Finally, in Section 2.4, we describe the computation of confidence intervals for the support and confidence of a rule.

### 2.1. Statistical Significance of a Single Association

We view a dataset consisting of  $n$  transactions as the realizations of  $n$  independent identically distributed random boolean vectors, sampled from the “real world” distribution. Let  $\pi^S$  denote the “real world” probability that a transaction contains a given itemset  $S$ . Thus, the number of transactions  $N^S$  in the sample (i.e., dataset) that contain  $S$  is a binomial random variable with success probability  $\pi = \pi^S$  and  $n$  trials.

**Hypothesis testing and minimum support.** The minimum support requirement can be cast in the hypotheses testing framework as follows. For example, suppose our minimum support requirement is 10%. For each itemset  $S$ , let  $H_0^S$  be the *null hypothesis* that  $\pi^S = 0.1$ , and let us test it against the *alternative hypothesis*  $H_1^S$  that  $\pi^S > 0.1$ . Let  $p^S$  be the fraction of transactions in the dataset that contain  $S$ . The test is to compare  $p^S$  with a threshold value  $p_0$  and reject  $H_0^S$

if and only if  $p^S \geq p_0$ . There are two kinds of possible errors (Cryer & Miller 1994):

Decision	Truth	
	$H_0^S$ true	$H_1^S$ true
Do not reject $H_0^S$	Correct decision	Type II error
Reject $H_0^S$	Type I error	Correct decision

The selection of  $p_0$  is determined by a bound on the desired probability of Type I error, which is called the *significance level*.

**P-values.** In general, the *p-value* of a test result is the probability of getting an outcome at least as extreme as the outcome actually observed; the p-value is computed under the assumption that the null hypothesis is true (Cryer & Miller 1994). In our example, the p-value corresponding to an observed fraction  $p^S$  is equal to the probability, under the assumption that  $\pi^S = 0.1$ , that the fraction of transactions that contain  $S$  is greater than or equal to  $p^S$ .

In order to compute the p-value, we use either the normal approximation, the Poisson approximation, or the exact binomial distribution, depending on the actual values of  $n$ , the minimum support requirement, and the observed support. For example, suppose  $n = 10,000$ , the minimum support is  $\pi = 0.1$  and the observed support is  $p = 0.109$ . We use the normal approximation. The mean is  $\pi = 0.1$ , and the standard deviation is  $\sqrt{(\pi(1-\pi)/n)} = \sqrt{(0.09/10000)} = 0.003$ . Since  $p$  is three standard deviations greater than  $\pi$ , the p-value is 0.0013.

**Testing Independence.** Consider an association rule  $S \Rightarrow T$ , where  $S$  and  $T$  are sets of items. As a null hypothesis we assume that  $S$  and  $T$  occur in transactions independently. Thus, under the null hypothesis,<sup>3</sup>  $\pi^{S \wedge T} = \pi^S \times \pi^T$ . As an alternative hypothesis, we can use the inequality  $\pi^{S \wedge T} > \pi^S \times \pi^T$ , which means that the conditional probability of  $T$  given  $S$  is greater than the probability of  $T$ .

If the values  $\pi^S$  and  $\pi^T$  are assumed to be known with sufficient accuracy, we can use the value  $\pi^S \times \pi^T$  to compute a p-value for  $S \Rightarrow T$ . This p-value corresponds to the probability, under the assumption that  $S$  and  $T$  are independent, that the empirical frequency of the set  $S \cup T$  will be greater than  $p^{S \wedge T}$ . Since we don't know the actual values of  $\pi^S$  and  $\pi^T$ , we use  $p^S$  and

<sup>3</sup>We denote the event that  $S \cup T$  is included in the transaction by  $S \wedge T$  since this is indeed the intersection of the events corresponding to  $S$  and  $T$ .

$p^T$  (the fractions of transactions that contain  $S$  and  $T$ , respectively) as estimates for  $\pi^S$  and  $\pi^T$ . The lower the p-value, the more likely it is that  $S$  and  $T$  are not independent.

## 2.2. Statistical Significance of a Set of Associations

Suppose we are testing  $k$  null hypotheses  $H_0^1, \dots, H_0^k$ , and denote by  $q^i$  the probability of rejecting  $H_0^i$  when it is true. The probability of rejecting at least one of the null hypotheses when they are all true is at most  $q^1 + \dots + q^k$ . Thus, if we wish the latter to be smaller than, say, 0.05, it suffices to determine thresholds for the individual tests so that  $q_i < 0.05/k$ . This bound may be very small if the number of hypotheses implicitly being tested is very large. Indeed, since under the null hypothesis the empirical p-value is distributed uniformly, when we test  $k$  true null hypotheses whose test statistics are independent random variables, the expected value of the smallest p-value is  $1/(k+1)$ , so in order to achieve a small probability of rejecting any true null hypothesis we would have to choose thresholds even smaller than that. Note that when we test independence of pairs in a set of, say, 10,000 items, then the value of  $k$  would be greater than  $10^7$ .

Obviously, if we wish to achieve a good probability of rejecting most “false discoveries”, we have to increase the probability of rejecting some true ones as well. In other words, when we attempt to discover more true rules, we also increase the risk of false discoveries (i.e., rejecting true null hypotheses). However, we would like to have an idea how many false discoveries there may be for any given threshold. We explain below how to compute an upper bound on the number of hypotheses that are implicitly tested; this number can be used to estimate the number of false discoveries for any given threshold.

**Upper Bound on the Number of Hypotheses.** The number of hypotheses that we are implicitly testing in the associations algorithm is typically much larger than just the number of frequent itemsets or the number of rules that are generated. To understand why, consider the set of frequent pairs. Let the null hypothesis  $H_0^{i,j}$  be that the items  $i$  and  $j$  are independent. To find the set of frequent pairs, the associations algorithm counts the cross-product of all the frequent items. Suppose there are 100 frequent items. Then there will be roughly 5000 pairs of items whose support is counted. If the algorithm throws away 4000 of these *at random* and tests  $H_0^{i,j}$  only for the remaining 1000 pairs, then only

1000 hypotheses have been tested. On the other hand, if the algorithm picks the 1000 pairs with the smallest p-values, then 5000 hypotheses have been tested. If the algorithm first considers the 1000 pairs with the highest support, and only then looks at p-values, then the actual number of hypotheses being tested is not readily available. In this case, we can use the number 5000 as an upper bound on the number of hypotheses.

We can extend this upper bound to include itemsets and rules with more than two items. Consider an itemset with three frequent items that does not include any frequent pairs. We do not need to include such an itemset while counting the number of hypotheses because this itemset clearly cannot have minimum support (in the dataset) and hence its properties are never examined by the algorithm. Hence for itemsets with three items, the number of hypotheses is less than the product of the number of frequent pairs times the number of frequent items. In fact, we can further bound the number of hypotheses to just those itemsets all of whose subsets are frequent. This number is exactly the number of candidates counted by current algorithms. By summing this over all the passes, we get

$$\begin{aligned} &\text{Number of Hypotheses} \\ &\leq \text{Number of Candidate Itemsets} \\ &\leq \text{Number of Frequent Itemsets} \times \text{Number} \\ &\quad \text{of Frequent Items.} \end{aligned}$$

where the set of candidate itemsets includes any itemset all of whose subsets are frequent. We emphasize that this is only an upper bound and may be much higher than necessary. Another serious problem is that even when all the null hypotheses are true, their test statistics are clearly not independent, thus prohibiting a direct calculation of appropriate thresholds. Below we present a practical solution to this problem.

## 2.3. Determining thresholds by resampling

Given the observed singleton frequencies  $p_i$  of the items, we generate a few synthetic data sets of transactions under a model where the occurrences of all the items are independent. Thus, the transactions are generated independently; for each transaction  $j$  and for item  $i$  we pick a number  $x_{ij}$  from a uniform distribution over  $[0, 1]$  and include  $i$  in  $j$  if and only if  $x_{ij} < p_i$ . Typically, we would generate 9 data sets, each consisting of 10,000 transactions. These numbers depend on the number of frequent items and minimum support rather than the number of transactions. We run the association rules algorithms on these data sets. Let  $v_{ij}$  denote the  $i$ th smallest p-value in dataset  $j$ . Let  $V_i$  denote the

Dataset	Supermarket	Dept. Store	Mail Order
Number of Customers	6200	Unknown	214,000
Number of Transactions	1.5 million	570,000	3 million
Items per Transaction	9.6	4.4	2.6
Min. Support (for exp.)	2%	1%	0.02%
Min. Conf (for exp.)	25%	25%	25%
# Frequent Items	201	283	2849
# Frequent Itemsets <sup>4</sup>	2541	943	10,173
# Candidates	30,000	42,000	4,090,000
# Rules	4828	1020	2479

Table 1: Dataset Characteristics

mean of the values  $v_{i1}, v_{i2}, \dots$ . The value  $V_i$  estimates the expectation of the  $i$ th smallest p-value when all the null hypotheses are true. So, we expect at most  $i$  false discoveries when we place the threshold at  $V_i$ . These estimates become useful when we wish to assess the quality of the set of rules we mine from the real data set. For example, if in the real data set we consider reporting all the rules with p-values smaller than some threshold  $t$ , and if  $V_i < t \leq V_{i+1}$ , then we expect no more than  $i$  of these rules to be false discoveries, since even in a purely synthetic data base where all the null hypotheses are true, no more than an expected number of  $i$  turn out to have such small p-values. As the value of  $t$  increases, more rules would be reported, but a larger number of them is expected to be false.

We tried this approach on three real-life datasets, whose characteristics are shown in Table 1. We present results for a specific minimum support and confidence for each rule; we got similar results for other values of support and confidence. Table 2 gives for each dataset the results of the simulation. We present results with three different random seeds to give an idea of the variation in p-values. For the supermarket and department store data, we also ran with three different data sizes: 1000, 10,000 and 100,000 transactions.<sup>5</sup> Notice that the average p-values are quite similar for the three data sizes.

We estimated the smallest p-value for each dataset based on the conservative upper bound on the number of hypotheses that we derived in the previous section. There was more than a factor of 100 difference between the expected lowest p-value and the actual least p-value on all three datasets.

For the Supermarket data, only two rules (out of 4828 rules) had p-values higher than  $10^{-9}$ : their p-values

<sup>5</sup>For the mail order data, the minimum support was too low to get meaningful results with the first two data sizes. With 10,000 transactions, minimum support corresponds to just 2 transactions.

were .0037 and .0051. For the Department Store data, only nine rules (out of 1020 rules) had p-values higher than  $10^{-100}$ , and all their p-values were greater than 0.09. For the Mail Order data, none of the rules (out of 2479 rules) had p-values greater than  $10^{-40}$ . Hence the number of “false discoveries” was extremely small.

The reason for the extremely low number of false discoveries is that the support and confidence threshold already do an excellent job of pruning out most rules that are not statistically significant. For instance, consider a rule where the support of the consequent is 5%. For this rule to meet the minimum confidence constraint, the support (confidence) of this rule must be at least 5 times the expected support (confidence) assuming that the antecedent and consequent are independent. Hence, unless the minimum support was extremely low, this rule would have a very low p-value.

## 2.4. Confidence Intervals

Denote by  $B(k; n, s)$  the probability that a binomial random variable with success probability  $s$  and  $n$  trials will have a value greater than  $k$ . The p-value of a rule with observed frequency  $p$ , with respect to a desired support level of  $s$  is equal to  $B(np; n, s)$ . Let  $\pi$  denote the true frequency. The probability of the event  $\pi - x \leq p \leq \pi + y$  is the same as the confidence level of an interval of the form  $[p - y, p + x]$ . The symmetry of the normal approximation allows calculating confidence intervals based on the observed value  $p$ . If we construct for each rule a confidence interval of level 95%, then for each rule there is an a priori probability of 95% that the true frequency lies within the interval. This means that the expected proportion of the rules where the true frequency lies within the respective interval is 95%. With regard to constructing a confidence interval for the confidence of a rule, we can argue the following. In general, consider events  $E_1 \subset E_2$ . If  $[a, b]$  and  $[c, d]$  are confidence intervals of level  $1 - \epsilon$  for  $\pi(E_1)$  and  $\pi(E_2)$ , respectively, and if  $c > 0$ , then  $[a/d, b/c]$  is a confidence



Simulated Dataset	Number of Transactions	Expected Lowest p-value	Lowest p-value	Next Lowest p-value
Supermarket	1,000	3e-5	.0026, .0048, .0072	.0038, .0074, .0089
	10,000		.0030, .0044, .0064	.0049, .0110, .0140
	100,000		.0011, .0022, .0086	.0049, .0055, .0096
Dept. Store	1,000	2e-5	3e-5, .0025, .0025	.0010, .0027, .0029
	10,000		.0013, .0025, .0032	.0032, .0040, .0090
	100,000		.0002, .0021, .0045	.0006, .0022, .0090
Mail Order	100,000	2e-7	2e-5, 6e-5, .0002	7e-5, 8e-5, .0003

Table 2: Simulation Results

interval for  $p(E_1|E_2)$  with confidence level of at least  $1 - 2\epsilon$ .

These confidence intervals allow users to use associations rules predictively by giving them an idea of how much variance they can expect in the support and confidence of a rule in the future.

### 3. Conclusions

We looked at the issue of whether association rule algorithms produce many "false discoveries". It is straightforward to compute the statistical significance of a single rule. However, when looking at a set of rules, the significance test has to take into account the number of hypotheses being tested. We showed that the number of hypotheses implicitly being tested can be much greater than the number of output rules, and derived an upper bound for the number of hypotheses. Unfortunately deriving an acceptance threshold for the statistical significance test from this bound may be too conservative. We presented a novel approach of using resampling to determine the acceptance threshold for the significance test. The threshold value derived using this approach was typically more than 100 times greater than the threshold value derived from the upper bound.

We then used this threshold to evaluate the number of "false discoveries" on three real-life dataset. We found that less than 0.1% of the rules were false discoveries: the reason for this surprisingly low number is that the minimum support and confidence constraints already do an excellent job of pruning away the statistically insignificant rules. A bonus of this work is that the statistical significance measures we compute are a good basis for ordering the rules for presentation to users, since they correspond to the statistical "surprise" of the rule.

Finally, we derived confidence intervals for the support and confidence of an association rule, enabling users to use the rule predictively over future data.

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