

An integer arithmetic method to compute generalized matrix inverse and solve linear equations exactly

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Abstract. An algorithm that uses integer arithmetic is suggested. It transforms an $m \times n$ matrix to a diagonal form (of the structure of Smith Normal Form). Then it computes a reflexive generalized inverse of the matrix exactly and hence solves a system of linear equations error-free.

Keywords. Exact computation; integer arithmetic; generalized matrix inverse; linear equations; Smith diagonal form.

1. Introduction

Whether a matrix A over a complex field is singular square, or rectangular, it has always a generalized inverse (g -inverse) over the (complex) field. The true inverse exists only when A is nonsingular (i.e., a square matrix whose determinant is not zero). However, a g -inverse of an $m \times n$ matrix of rank r involves considerable errors if the r th order submatrices are near-singular. Further, the rank shown by the g -inverse may be less than the actual rank. In fact, identical are the pitfalls when a (square) near-singular matrix is inverted.

We present here a method that uses integer arithmetic to

- (i) transform an $m \times n$ integral matrix to a Smith Diagonal Form (defined later) without requiring to compute the greatest common divisors (GCDs) of the matrix elements as required in computing certain g -inverses (Hurt and Waid [4], Ben-Israel and Greville [2]),
- (ii) compute a reflexive g -inverse (Bowman and Burdet [3], Ben-Israel and Greville [2], Krishnamurthy and Sen [5]),
- (iii) obtain a solution vector x of $Ax=b$, b being 0 (null column vector) or not.

Since any computing system can represent only the rational numbers, we can, without any loss of generality, assume the inputs (here the matrix A and the right-hand-side column vector b) integral.

2. Definitions

2.1. Integral vector and integral matrix

- Let (a) K = the ring of integers $0, \pm 1, \pm 2, \dots$,
(b) K^m = the m dimensional vector space over K ,

- (c) $K^{m \times n}$ = the $m \times n$ matrices over K , and
 (d) $K_r^{m \times n}$ = the $m \times n$ matrices with rank r over K .

Any element of K^m is an integral vector. Any element of $K^{m \times n}$ is an integral matrix, and any element of $K_r^{m \times n}$ is an integral matrix of rank r .

2.2. Elementary row and column operations

A sequence of elementary row and column operations used here consists of

- (a) *Type 1.* Interchanging two rows (columns),
 (b) *Type 2.* Subtracting an integral multiple of one row (column) from another row (column), and
 (c) *Type 3.* Replacing a row (column) by an integral multiple of the row (column).

2.3. Elementary integral matrix

Any nonsingular matrix $P \in K^{m \times m}$ (any nonsingular matrix $Q \in K^{n \times n}$) which when pre- (post-) multiply a given matrix $A \in K_r^{m \times n}$, produces a combination of types 1, 2, and 3 operations.

2.4. Equivalent matrices

Two matrices $A, S \in K^{m \times n}$ are equivalent over K if there exist two elementary integral matrices $P \in K^{m \times m}$ and $Q \in K^{n \times n}$ such that $PAQ = S$.

2.5. Smith diagonal form

A matrix $S = (s_{ij}) \in K_r^{m \times n}$ is the Smith Diagonal Form (SDF) of $A \in K_r^{m \times n}$ if

- (a) $s_{ii} \neq 0, i = 1(1)r$; (b) $s_{ij} = 0$ otherwise, and (c) s_{ii} divides $s_{i+1, i+1}, i = 1(1)r-1$.

2.6. Generalized inverse of Smith diagonal form

The g -inverse of SDF S is the matrix $S^+ = (s_{ij}^+) \in K_r^{n \times m}$ if

- (a) $s_{ii}^+ = s_{ii}^{-1}, i = 1(1)r$, and (b) $s_{ij}^+ = 0$ otherwise.

3. The method

Let $A \in K_r^{m \times n}$ do not have any zero row or zero column.

Step 1. Computing SDF $S = (s_{ij}) \in K_r^{m \times n}$

- (i) Make the elements at positions (1, 1), (2, 2), ..., nonzero using type 1 operations if any element at any of the positions (1, 1), (2, 2), ..., is zero. Otherwise go to step 1(ii).
 (ii) Multiply (type 3 operations) second, third, ..., m th rows by $p_1, p_1 p_2, \dots, p_1 p_2 \dots p_{m-1}$, respectively, where p_i is the element at position (i, i) or p_i is a number so that the

elementary matrices P_1, P_2, \dots , where $P = \dots P_2 P_1$, are integral (see step 2, and examples in section 6) or p_i is a number so that P is integral and the remainder in any division is zero.

(iii) Make zeros of all the elements below position (1, 1) in the first column using type 2 operations.

(iv) If the element at position (2, 2) is not zero then make zeros of all the elements below position (2, 2) in the second column using type 2 operations. Otherwise interchange $p_2 p_3 \dots p_{s-1}$ times the second row and $1/(p_2 p_3 \dots p_{s-1})$ times the s th row whose second element is nonzero (types 1 and 3 operations).

Make zeros of all the elements below position (2, 2) in the second column (if these are not all zero). Continue the process for the elements at positions (3, 3), (4, 4), etc.

Note. Steps 1(i)–1(iv) make the first r rows non-null with all the elements below positions (1, 1), (2, 2), ..., etc. zero and other rows (($r+1$)st to m th rows) null.

(v) Multiply second, third, ..., r th columns by $q_1, q_1 q_2, \dots, q_1 q_2 \dots q_{r-1}$, respectively, where q_i is the element at position (i, i) or q_i is a number so that the elementary matrices Q_1, Q_2, \dots , where $Q = Q_1 Q_2, \dots$, are integral (see step 2 and examples in section 6) or q_i is a number so that Q is integral and the remainder in any division is zero. Multiply each of ($r+1$)st, ($r+2$)nd, ..., n th columns by $q_1 q_2 \dots q_{r-1}$.

(vi) Make zeros of all the elements after position (1, 1) in the first row using type 2 operations.

(vii) If the element at position (2, 2) is not zero then make zeros of all the elements after position (2, 2) in the second row using type 2 operations. Otherwise interchange $q_2 q_3 \dots q_{s-1}$ times the second column and $1/(q_2 q_3 \dots q_{s-1})$ times the s th column whose second element is nonzero (types 1 and 3 operations) and then make zeros of all the elements after position (2, 2) in second row (if these are not all zero). Continue the process for the elements at positions (3, 3), (4, 4), etc.

Note. Steps 1(i)–1(vii) give the SDF S . Also, in any divide operation denominator divides the numerator.

Step 2. Computing A^- .

(i) Compute the elementary matrix $P(Q)$ defined in sec. 2(iv), which is the product of all the elementary row (column) matrices, in the right order. Thus $PAQ = S$.

(ii) To obtain $A^- = QS^+ P$, compute

$$(i, j)\text{th element of } aA^- = \sum_{k=1}^r q_{ik} p_{kj} (a/s_{kk}),$$

$$i = 1(1)n, j = 1(1)m \text{ where } a = s_{rr}.$$

Step 3. Computing solution vectors. To solve $Ax = b$, x being rational,

(i) compute AaA^-b . If it is equal to ab , then solution exists. Otherwise, the system has no solution.

(ii) If $AaA^-b = ab$, then compute aA^-b .

Obtain $ax = aA^-b + ay - aA^-Ay$ for any $y \in K^n$.

Remarks

(i) *Computing g-inverse of a rational matrix.* Let B be an $m \times n$ nonintegral rational matrix. Then $A = \beta B$ is integral. Hence $B^- = \beta A^-$.

(ii) *Computing solution vector of a rational system.* Let $Cx=d$ be the nonintegral rational system. Then $Ex=f$, where $E=\gamma C$ and $f=\gamma d$, is an integral system.

4. Results

The method follows from the theorem and corollary below.

Theorem. Let $A \in K_r^{m \times n}$. Then A is equivalent over K to an SDF $S \in K_r^{m \times n}$. The proof follows from the construction of S described in step 1 of the method (Sec. 3).

Corollary. Let P and Q be elementary integral matrices and $PAQ=S$ be an SDF of $A \in K^{m \times n}$. Also, let $A^- = QS^+P$. Then

$$AA^-A = A, A^-AA^- = A^-.$$

Proof. $PAQ=S=SS^+S=PAQS^+PAQ=PAA^-AQ$. Hence $A=AA^-A$. A^-AA^- is proved similarly.

Note.

(i) *Integrality condition.* The A^- does not satisfy

$$A^-A \in K^{n \times n} \text{ and } AA^- \in K^{m \times m}$$

if (nontrivial) type 3 operations are used.

(ii) *Triangular matrices.* If type 1 operations are not needed then P and Q will be lower and upper triangular matrices, respectively. In such a case step 2(ii) can be written as

To obtain $A^- = QS^+P$, compute

$$(i, j)\text{th element of } aA^- = \sum_{k=i}^r q_{ik} p_{kj} (\alpha/s_{kk}), i \geq j$$

$$= \sum_{k=j}^r q_{ik} p_{kj} (\alpha/s_{kk}), i < j$$

$$i = 1(1)r, j = 1(1)r \text{ where } \alpha = s_{rr}.$$

5. Use of modular arithmetic

For exact computation the modular arithmetic (Adegbeyeni and Krishnamurthy [1], Rao *et al* [6]) can be used only when the integer arithmetic (Sen and Shamim [7]),

in a general purpose computing system, demands too long precision operands. The modular arithmetic offers the parallelism in computation but the total computation is about u times the computation needed by the integer arithmetic, u being the number of prime bases used.

6. Examples

$$(1) \quad A = \begin{bmatrix} 2 & 3 & 5 \\ 4 & 6 & 1 \\ 3 & 5 & 10 \end{bmatrix}$$

Here $p_1=2$, $p_2=6$, $p_3=10$.

Step 1(ii). Multiply second and third rows by 2 and 2×6 , respectively (type 3):

$$P_1 A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 12 \end{bmatrix} \quad A = \begin{bmatrix} 2 & 3 & 5 \\ 8 & 12 & 2 \\ 36 & 60 & 120 \end{bmatrix}$$

Step 1(iii). To reduce the first column elements below position (1, 1) zero premultiply $P_1 A$ by P_2 (type 2):

$$P_2 P_1 A = \begin{bmatrix} 1 & 0 & 0 \\ -4 & 1 & 0 \\ -18 & 0 & 1 \end{bmatrix} \quad P_1 A = \begin{bmatrix} 2 & 3 & 5 \\ 0 & 0 & -18 \\ 0 & 6 & 30 \end{bmatrix}$$

Step 1(iv). The element at position (2, 2) is zero. So interchange p_2 times the second row and $(1/p_2)$ times the third row (here $s=3$) whose second element is nonzero (types 1 and 3):

$$P_3 P_2 P_1 A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1/6 \\ 0 & 6 & 0 \end{bmatrix} \quad P_2 P_1 A = \begin{bmatrix} 2 & 3 & 5 \\ 0 & 1 & 5 \\ 0 & 0 & -108 \end{bmatrix}$$

Since the element below position (2, 2) in the second column is zero and $m=3$, we go to step 1(v).

Step 1(v). $q_1=2$, $q_2=1$, $q_3=-108$. Multiply second and third columns by 2, 2×1 , respectively (type 3):

$$P_3 P_2 P_1 A Q_1 = P_3 P_2 P_1 A \begin{bmatrix} 1 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & 2 \end{bmatrix} = \begin{bmatrix} 2 & 6 & 10 \\ 0 & 2 & 10 \\ 0 & 0 & -216 \end{bmatrix}$$

Step 1(vi). Make zeros of all the elements after position (1, 1) in the first row (type 2):

$$P_3 P_2 P_1 A Q_1 Q_2 = P_3 P_2 P_1 A Q_1 \begin{bmatrix} 1 & -3 & -5 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 2 & 10 \\ 0 & 0 & -216 \end{bmatrix}$$

Step 1(vii). Since the element at position (2, 2) is not zero, make zeros of all the elements (here one) after position (2, 2) in the second row (type 2):

$$P_3P_2P_1AQ_1Q_2Q_3 = P_3P_2P_1AQ_1Q_2 \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & -5 \\ 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} 2 & 0 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & -216 \end{bmatrix} = S \text{ (SDF)}$$

$$P = P_3P_2P_1 = \begin{bmatrix} 1 & 0 & 0 \\ -3 & 0 & 2 \\ -24 & 12 & 0 \end{bmatrix}$$

$$Q = Q_1Q_2Q_3 = \begin{bmatrix} 1 & -3 & 10 \\ 0 & 2 & -10 \\ 0 & 0 & 2 \end{bmatrix}, PAQ = S.$$

$$a = s_{33} = -216$$

$$aS^+ = \begin{bmatrix} -108 & 0 & 0 \\ 0 & -108 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$aA^{-1} = aA^- = QA^+P = \begin{bmatrix} -1320 & 120 & 648 \\ 888 & -120 & -432 \\ -48 & 24 & 0 \end{bmatrix}$$

$$(2) \quad A = \begin{bmatrix} -1 & 2 & 3 & 3 \\ 2 & 5 & 6 & 3 \\ -5 & -8 & -9 & -3 \end{bmatrix}.$$

Multiply second and third rows by 1 and 1×5 , respectively (type 3):

$$P_1A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 5 \end{bmatrix} A = \begin{bmatrix} -1 & 2 & 3 & 3 \\ 2 & 5 & 6 & 3 \\ -25 & -40 & -45 & -15 \end{bmatrix}$$

To reduce the first column elements below position (1, 1) zero pre-multiply P_1A by P_2 (type 2):

$$P_2P_1A = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 5 \end{bmatrix} P_1A = \begin{bmatrix} -1 & 2 & 3 & 3 \\ 0 & 9 & 12 & 9 \\ 0 & -90 & -120 & -90 \end{bmatrix}$$

To reduce the second column element below position (2, 2) zero pre-multiply P_2P_1A by P_3 (type 2):

$$P_3P_2P_1A = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 10 & 1 \end{bmatrix} P_2P_1A = \begin{bmatrix} -1 & 2 & 3 & 3 \\ 0 & 9 & 12 & 9 \\ 0 & 0 & 0 & 0 \end{bmatrix}.$$

Multiply second, third, and fourth columns by 1, 1×9 , 1×9 , respectively (type 3):

$$P_3 P_2 P_1 A Q_1 = P_3 P_2 P_1 A \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 9 & 0 \\ 0 & 0 & 0 & 9 \end{bmatrix} = \begin{bmatrix} -1 & 2 & 27 & 27 \\ 0 & 9 & 108 & 81 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

To reduce the first row elements after position (1, 1) zero post-multiply $P_3 P_2 P_1 A Q_1$ by Q_2 (type 2):

$$P_3 P_2 P_1 A Q_1 Q_2 = P_3 P_2 P_1 A Q_1 \begin{bmatrix} 1 & 2 & 27 & 27 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} = \begin{bmatrix} -1 & 0 & 0 & 0 \\ 0 & 9 & 108 & 81 \\ 0 & 0 & 0 & 0 \end{bmatrix}$$

To reduce the second row elements after position (2, 2) zero post-multiply $P_3 P_2 P_1 A Q_1 Q_2$ by Q_3 (type 2):

$$P_3 P_2 P_1 A Q_1 Q_2 Q_3 = P_3 P_2 P_1 A Q_1 Q_2 \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & -12 & -9 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} \\ = \begin{bmatrix} -1 & 0 & 0 & 0 \\ 0 & 9 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix} = S$$

$$P = P_3 P_2 P_1 = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -5 & 10 & 5 \end{bmatrix}, \quad Q = Q_1 Q_2 Q_3 = \begin{bmatrix} 1 & 2 & 3 & 9 \\ 0 & 1 & -12 & -9 \\ 0 & 0 & 9 & 0 \\ 0 & 0 & 0 & 9 \end{bmatrix}$$

$$\alpha = s_{22} = 9, \quad \alpha S^+ = \begin{bmatrix} -9 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

$$\alpha A^- = Q \alpha S^+ P = \begin{bmatrix} -5 & 2 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}$$

Note that $PAQ = S$, and P and Q are lower and upper triangular matrices, respectively since type 1 operations are not used.

$$A \alpha A^- = \begin{bmatrix} 9 & 0 & 0 \\ 0 & 9 & 0 \\ 9 & -18 & 0 \end{bmatrix}$$

If $b = (7 \ 16 \ -25)^t$ then $AaA^{-1}b = ab$ and hence solutions exist. If $b = (8 \ 16 \ -25)^t$ then $AaA^{-1}b \neq ab$ and hence solutions do not exist. For $b = (7 \ 16 \ -25)^t$ the solution vector $ax = (-3 \ 30 \ 0 \ 0)^t$.

References

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