# AN OVERVIEW OF MGMRES AND LAN/MGMRES METHODS FOR SOLVING NONSYMMETRIC LINEAR SYSTEMS 

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#### Abstract

We present an overview of the MGMRES and LAN/MGMRES iterative methods for solving large sparse linear systems.


## 1. Introduction

We begin with a brief discussion of background material on Idealized Generalized Conjugate Gradient (IGCG) methods and Krylov subspace methods. Following a review of the Generalized Minimum Residual (GMRES) method, we outline the MGMRES method, which is a modification of the GMRES method. Finally, we sketch a Lanczos-type procedure called the LAN/MGMRES method.

We consider linear systems of the form

$$
A u=b
$$

with true solution $\bar{u}=A^{-1} b$. Here $A$ is a large sparse nonsingular matrix of size $N \times N$. Recall that if we are given an arbitrary initial guess $u^{(0)}$ to be used in an iterative method, then the initial residual vector is $r^{(0)}=b-A u^{(0)}$. Iterative methods involve iterates $u^{(1)}, u^{(1)}, \ldots, u^{(n)}$ that hopefully converge to an approximation to the true solution; that is, the $n$th residual vector $r^{(n)}=b-A u^{(n)}$ is approximately the zero vector.

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## 2. Krylov Subspace and $\operatorname{IGCG}(Z)$ Methods

Let $Z$ be an auxiliary matrix for an iterative method such as $Z=I, Z=Y$, $Z=A^{T}$, or $Z=A^{T} Y$, for example. If $A$ is symmetric positive definite (SPD), then it can be shown that $Z=I$ for the conjugate gradient method and $Z=A^{T}$ for the conjugate residual method.

We state several important conditions for Krylov subspace methods and Idealized Generalized Conjugate Gradient IGCG(Z) methods.

## Condition I:

$$
u^{(n)}-u^{(0)} \in \mathcal{K}_{n}\left(r^{(0)}, A\right)=\operatorname{Span}\left\{r^{(0)}, A r^{(0)}, \ldots, A^{n-1} r^{(0)}\right\} .
$$

Here $\mathcal{K}_{n}\left(r^{(0)}, A\right)$ is the Krylov space associated with the initial residual vector $r^{(0)}$ and the matrix $A$.

Condition II (a) (Minimization condition): If $Z A$ is SPD, then

$$
\left\langle\left(u^{(n)}-\bar{u}\right),\left(u^{(n)}-\bar{u}\right)\right\rangle_{Z A}=\left\|u^{(n)}-\bar{u}\right\|_{Z A^{1 / 2}}^{2} \quad \text { is minimized. }
$$

Condition II (b) (Galerkin condition):

$$
\left\langle r^{(n)}, v\right\rangle_{Z}=0 \quad \text { for all } v \in \mathcal{K}_{n}\left(r^{(0)}, A\right)
$$

Here the $Z$-inner product is defined as $\langle x, y\rangle_{Z}=\langle Z x, y\rangle=y^{T} Z x$.
The minimization condition (Condition II (a)) can also be written as:

$$
\frac{1}{2}\left\langle u^{(n)}, u^{(n)}\right\rangle_{Z A}-\left\langle b, u^{(n)}\right\rangle_{Z A} \quad \text { is minimized. }
$$

Notice that if $Z=A^{T} Y$, where $Y$ is SPD, then $Z A=A^{T} Y A$ is SPD. It follows that Condition II (a) becomes $\left\langle r^{(n)}, r_{Y}^{(n)}=\left\|r^{(n)}\right\|_{Y^{\frac{1}{2}}}^{2}\right.$ is minimized.

The index $m=m\left(r^{(0)}, A\right)$ of $u^{(0)}$, with respect to $A$, is the largest integer $m$ such that the vectors $v^{(0)}, v^{(1)}, \ldots, v^{(m)}$ are linearly independent. For example, letting $v^{(0)}=r^{(0)}, v^{(1)}=A r^{(0)}, \ldots, v^{(m)}=A^{m} r^{(0)}$, it can be shown that

$$
\begin{aligned}
u^{(0)}-\bar{u} \in \mathcal{K}_{m+1}\left(r^{(0)}, A\right) & =\operatorname{Span}\left\{r^{(0)}, A r^{(0)}, \ldots, A^{m} r^{(0)}\right\} \\
& =\operatorname{Span}\left\{v^{(0)}, V^{(1)}, \ldots, v^{(m)}\right\}
\end{aligned}
$$

if $m \leq N-1$. Then $u^{(m+l)}=\bar{u}$ hopefully.
The $\operatorname{IGCG}(\mathrm{Z})$ method is $\left(n^{*}, u^{(0)}\right)$-computable if $n^{*} \leq m+1$ and if for all $n \leq n^{*}$ there exists a unique $u^{(n)}$ satisfying $u^{(n)}-u^{(0)} \in \mathcal{K}_{n}\left(r^{(0)}, A\right)$ and
$\left\langle Z r^{(n)}, v\right\rangle=0$ for all $v \in \mathcal{K}_{n}\left(r^{(0)}, A\right)$. Moreover, the $\operatorname{IGCG}(\mathrm{Z})$ method is $\left(n^{*}, u^{(0)}\right)$-computable if and only if the moment matrix $\Delta_{n^{*}}\left(Z A, r^{(0)}\right)$ is strongly regular. Here the moment matrix is given by

$$
\Delta_{n^{*}}\left(Z A, r^{(0)}\right)=\left[\begin{array}{ccc}
\left\langle v^{(0)}, v^{(0)}\right\rangle_{Z A} & \cdots & \left\langle v^{\left(n^{*}-1\right)}, v^{(0)}\right\rangle_{Z A} \\
\vdots & & \vdots \\
\left\langle v^{(0)}, v^{\left(n^{*}-1\right)}\right\rangle_{Z A} & \cdots & \left\langle v^{\left(n^{*}-1\right)}, v^{\left(n^{*}-1\right)}\right\rangle_{Z A}
\end{array}\right] .
$$

This matrix is strongly regular if all the principal submatrices are nonsingular, which means that for a matrix of order $n$, the $n$ submatrices of sizes $1 \times 1,2 \times$ $2, \cdots, n \times n$ in the top-left-hand corner are nonsingular.

In orthogonal implementations, there are two phases.
Phase I. Construct basis vectors $w^{(0)}, w^{(1)}, \ldots, w^{(n-1)}$ by orthogonalizing Krylov vectors with respect to $C$ :

$$
\left\langle w^{(i)}, w^{(j)}\right\rangle_{C}=0 \quad \text { for } i \neq j
$$

Here $C$ is usually SPD.
Phase II. Choose $c_{0}^{(n)}, c_{1}^{(n)}, \ldots, c_{n-1}^{(n)}$ so that the Galerkin condition $\left\langle Z r^{(n)}, w^{(i)}\right\rangle=$ 0 for $0 \leq i \leq n-1$ is satisfied. We have

$$
\begin{aligned}
u^{(n)} & =u^{(0)}+c_{0}^{(n)} w^{(0)}+\cdots+c_{n-1}^{(n)} w^{(n-1)} \\
& =u^{(0)}+W_{n-1} c^{(n)},
\end{aligned}
$$

where

$$
W_{n-1}=\left[W^{(0)} w^{(1)} \cdots w^{(n-1)}\right], \quad c^{(n)}=\left[c_{0}^{(n)} c_{1}^{(n)} \cdots c_{n-1}^{(n)}\right]^{T} .
$$

In Phase I, we have

$$
w^{(n)}=A w^{(n-1)}+\beta_{n, 0} w^{(0)}+\cdots+\beta_{n, n-1} w^{(n-1)} .
$$

Examples are as follows: $C=A^{T} Z$ corresponds to the $\operatorname{ORTHODIR}(Z)$ method, $C=A$ corresponds to the $\operatorname{ORTHORES}(Z)$ method, and $C=Y$ together with $Z=A^{T} Y$ corresponds to the $\operatorname{GMRES}\left(A^{T} Y\right)$ method when $Y$ is SPD. The latter method is really the GGMRES method. For the GMRES method, we have $C=I$.

In Phase II, we have

$$
w^{(n)}=r^{(n)}+\alpha_{n, 0} w^{(0)}+\cdots+\alpha_{n, n-1} w^{(n-1)} .
$$

Examples are as follows: $C=Z A$ corresponds to the ORTHOMIN $(Z)$ method while $Z=A^{T}$ implies the conjugate residual method and $Z=I$ implies the conjugate gradient method.

## 3. Gmres Method

We now sketch the GMRES method of Saad and Schultz [6]. Let $Z=$ $A^{T} Y$, where $Y$ is a SPD matrix. Note that $Z A=A^{T} Y A$ is a SPD matrix. As mentioned above, Condition II (a) becomes $\left\langle Y r^{(n)}, r^{(n)}\right\rangle=\left\|r^{(n)}\right\|_{Y^{1 / 2}}$ is minimized.

In Phase I, we have

$$
\begin{aligned}
& \left\{\begin{array}{l}
\widehat{w}^{(0)}=r^{(0)} \\
w^{(0)}=\sigma_{0}-1 \\
\widehat{w}^{(0)}, \quad \text { where } \sigma_{0}=\left\langle Y \widehat{w}^{(0)}, \widehat{w}^{(0)}\right\rangle^{\frac{1}{2}}
\end{array}\right. \\
& \vdots \\
& \left\{\begin{array}{l}
\widehat{w}^{(n)}=A w^{(n-1)}+\beta_{n, 0} w^{(0)}+\cdots+\beta_{n, n-1} w^{(n-1)} \\
w^{(n)}=\sigma_{n}{ }^{-1} \widehat{w}^{(n)}, \quad \text { where } \sigma_{n}=\left\langle Y \widehat{w}^{(n)}, \widehat{w}^{(n)}\right\rangle^{\frac{1}{2}} .
\end{array}\right.
\end{aligned}
$$

Here

$$
\left\langle Y w^{(i)}, w^{(j)}\right\rangle= \begin{cases}1, & i=j \\ 0, & i \neq j .\end{cases}
$$

We have the basic relation

$$
A\left[w^{(0)} w^{(1)} \cdots w^{(n-1)}\right]=\left[w^{(0)} w^{(1)} \cdots w^{(n)}\right] H_{n}
$$

or

$$
A W_{n-1}=W_{n} H_{n}
$$

Here $H_{n}$ is an upper Hessenberg matrix of order $n$.
Example ( $n=2$ ):

$$
A\left[w^{(0)} w^{(1)}\right]=\left[w^{(0)} w^{(1)} w^{(2)}\right]\left[\begin{array}{rr}
-\beta_{1,0} & -\beta_{2,0} \\
\sigma_{1} & -\beta_{2,1} \\
0 & \sigma_{2}
\end{array}\right]
$$

Hence, we have

$$
A W_{1}=W_{2} H_{2} .
$$

In Phase II of the GMRES method, we have

$$
\begin{aligned}
u^{(n)} & =u^{(0)}+c_{0}^{(n)} w^{(0)}+\cdots+c_{n-1}^{(n)} w^{(n-1)} \\
& =u^{(0)}+W_{n-1} c^{(n)} .
\end{aligned}
$$

Consequently, from this equation we obtain

$$
\begin{aligned}
r^{(n)} & =b-A u^{(n)} \\
& =r^{(0)}-A W_{n-1} c^{(n)} \\
& =r^{(0)}-W_{n} H_{n} c^{(n)} \\
& =W_{n}\left(e^{(n+1)}-H_{n} c^{(n)}\right),
\end{aligned}
$$

using $A W_{n-1}=W_{n} H_{n}$ and $r^{(O)}=W_{n} e^{(n+1)}$, where $e^{(n+1)}=\left[\sigma_{n}, 0, \ldots, 0\right]_{n+1}^{T}$. Thus, we find

$$
\begin{aligned}
\left\langle Y r^{(n)}, r^{(n)}\right\rangle & =\left\langle Y W_{n}\left(e^{(n+1)}-H_{n} c^{(n)}, W_{n}\left(e^{(n+1)}-H_{n} c^{(n)}\right)\right\rangle\right. \\
& =\left\|e^{(n+1)}-H_{n} c^{(n)}\right\|_{2}^{2},
\end{aligned}
$$

since $W_{n}^{T} Y W_{n}=I_{n}$ and $Y$ is SPD.
Example ( $n=2$ ): Determination of $c^{(2)}$. The system

$$
H_{2} c^{(2)}=e^{(3)}
$$

has the form

$$
\left[\begin{array}{cc}
-\beta_{1,0} & -\beta_{2,0} \\
\sigma_{1} & -\beta_{2,1} \\
0 & \sigma_{2}
\end{array}\right]\left[\begin{array}{c}
c_{0}^{(2)} \\
c_{1}^{(2)}
\end{array}\right]=\left[\begin{array}{c}
\sigma_{2} \\
0 \\
0
\end{array}\right] .
$$

Using Givens rotations $Q=Q_{1} Q_{2}$ with $Q Q^{T}=I$, we have

$$
Q H_{2} c^{(2)}=Q e^{(3)},
$$

which is of the form

$$
\left[\begin{array}{cc}
\boxed{x} & \times \\
0 & \boxed{\times} \\
0 & 0
\end{array}\right]\left[\begin{array}{c}
c_{0}^{(2)} \\
c_{1}^{(2)}
\end{array}\right]=\left[\begin{array}{c}
\times \\
\times \\
\times
\end{array}\right]
$$

To get the least squares solution, we solve the first two equations for $c_{0}^{(2)}$ and $c_{1}^{(2)}$.

Note that the sum of the squares of the residuals are preserved:

$$
\begin{aligned}
\langle Q(b-A u), Q(b-A u)\rangle & =\left\langle(b-A u), Q^{T} Q(b-A u)\right\rangle \\
& =\langle b-A u, b-A u\rangle
\end{aligned}
$$

Some comparisons for orthogonal implementations. If the matrix $Z A$ is SPD (that is, if $Z=A^{T} Y$ for some SPD matrix $Y$ ), then the ORTHODIR
method converges but the ORTHOMIN and ORTHORES methods may breakdown. The ORTHODIR method is often numerically unstable and requires more work per iteration than the GMRES method. The $\operatorname{GMRES}\left(A^{T} Y\right)$ method, where $Y$ is a SPD matrix, is mathematically equivalent to the $\operatorname{ORTHODIR}\left(A^{T} Y\right)$ method, but requires less work per iteration and is more stable. The $\operatorname{GMRES}\left(A^{T} Y\right)$ method is widely used but the work per iteration increases as $n$ increases.

## 4. Mgmres Method

We now sketch the MGMRES method, which is a modification of the GMRES method. We assume $Y$ is symmetric and nonsingular (not necessarily SPD ). Also, we suppose that $Y A$ is symmetric,

In Phase I of the MGMRES method, we have

$$
\widehat{w}^{(n)}=A W^{(n-1)}+\beta_{n, n-1} w^{(n-1)}+\beta_{n, n-2} w^{(n-2)}
$$

and $\left\langle w^{(n)}, Y w^{(i)}\right\rangle=0$ for $0 \leq i \leq n-1$. Then we obtain

$$
w^{(n)}=\sigma_{n}{ }^{-1} \widehat{w}^{(n)}
$$

where $\sigma_{n}=\left|\left\langle Y \widehat{w}^{(n)}, \widehat{w}^{(n)}\right\rangle\right|^{1 / 2}$. Here the absolute value signs are used since the expression within them may be negative. Moreover, the process fails if $\sigma_{n}=0$. Next, we have

$$
W_{n}^{T} Y W_{n}=\operatorname{diag}( \pm 1, \pm 1, \ldots, \pm 1) \equiv D_{n}
$$

Here $D_{n}$ is a diagonal matrix with $\pm 1$ as diagonal entries. For the GMRES method, if $Y$ is a SPD matrix, then $D_{n}=\operatorname{diag}(1,1, \ldots, 1)$.

In Phase II of the MGMRES method, we use the Galerkin condition

$$
W_{n-1}^{T}\left(Z r^{(n)}\right)=0
$$

Also, we have $Z=A^{T} Y, r^{(n)}=W_{n}\left(e^{(n+1)}-H_{n} c^{(n)}\right)$, and $\left\langle Z r^{(n)}, w^{(i)}\right\rangle=0$ for $0 \leq i \leq n-1$. So we obtain

$$
H_{n}^{T} W_{n}^{T} Y W_{n} H_{n} c^{(n)}=H_{n}^{T} W_{n}^{T} Y W_{n} e^{(n+1)}
$$

which implies that

$$
H_{n}^{T} D_{n} H_{n} c^{(n)}=H_{n}^{T} D_{n} e^{(n+1)} .
$$

If $D_{n}=I$, we get the normal equations

$$
H_{n}^{T} H_{n} c^{(n)}=H_{n}^{T} e^{(n+1)}
$$

Applying a sequence of Givens rotations, we form an upper triangular system

$$
Q H_{n}=\tilde{H}_{n}
$$

where $Q^{T} Q=I$.
Example ( $n=2$ ):

$$
H_{2}=\left[\begin{array}{rr}
-\beta_{1,0} & -\beta_{2,0} \\
\sigma_{1} & -\beta_{2,1} \\
0 & \sigma_{2}
\end{array}\right] \quad \Longrightarrow \quad Q H_{2}=\tilde{H}_{2}=\left[\begin{array}{cc}
\boxed{x} & \times \\
0 & \boxed{x} \\
0 & 0
\end{array}\right]
$$

where $Q=Q_{1} Q_{2}$.
Using $H_{n}=Q^{-1} \tilde{H}_{n}=Q^{T} \tilde{H}_{n}$ and $Q^{-1}=Q^{T}$, we obtain

$$
\tilde{H}_{n}^{T} Q D_{n} Q^{T} \tilde{H}_{n} c^{(n)}=\tilde{H}_{n}^{T} Q D_{n} e^{(n+1)}
$$

Letting

$$
\begin{aligned}
& z=Q D_{n} Q^{T} \tilde{H}_{n} c^{(n)} \\
& y=\tilde{H}_{n}^{T} Q D_{n} e^{(n+1)}
\end{aligned}
$$

we obtain

$$
\tilde{H}_{n}^{T} z=y .
$$

So our strategy is to first solve this system to get $z$ and then solve

$$
\tilde{H}_{n} c^{(n)}=Q D_{n}^{-1} Q^{T} z
$$

Example ( $n=2$ ): The system

$$
\tilde{H}_{2}^{T} z=y
$$

has the form

$$
\left[\begin{array}{ccc}
\boxed{x} & 0 & 0 \\
x & \boxed{x} & 0
\end{array}\right]\left[\begin{array}{l}
z_{1} \\
z_{2} \\
z_{3}
\end{array}\right]=\left[\begin{array}{l}
y_{1} \\
y_{2}
\end{array}\right] .
$$

And we obtain

$$
z=\left[\begin{array}{c}
z_{1} \\
z_{2} \\
0
\end{array}\right]+k\left[\begin{array}{l}
0 \\
0 \\
1
\end{array}\right]
$$

where $k$ is arbitrary.

$$
\tilde{H}_{2} c^{(2)}=Q D_{2}^{-1} Q^{T}\left\{\left[\begin{array}{c}
z_{1} \\
z_{2} \\
0
\end{array}\right]+k\left[\begin{array}{l}
0 \\
0 \\
1
\end{array}\right]\right\}=\left[\begin{array}{c}
z_{1}^{\prime} \\
z_{2}^{\prime} \\
0
\end{array}\right]
$$

for suitable $k$. (If $D_{2}=I_{2}$, let $k=0$.) Failure occurs if the third component of $Q D_{2}^{-1} Q^{T}\left[z_{1} z_{2} 0\right]^{T}$ is not zero and the third component of $Q D_{2}^{-1} Q^{T}\left[\begin{array}{lll}0 & 1\end{array}\right]^{T}$ is zero. For GMRES, $Q D_{2} Q^{T}=I$ and we let $k=0$.

$$
\tilde{H}_{2} c^{(2)}=\left[\begin{array}{cc}
\boxed{\times} & \times \\
0 & \boxed{\times} \\
0 & 0
\end{array}\right]\left[\begin{array}{l}
c_{0}^{(2)} \\
c_{1}^{(2)}
\end{array}\right]=\left[\begin{array}{c}
z_{1}^{\prime} \\
z_{2}^{\prime} \\
0
\end{array}\right] .
$$

Finally, we solve for $c_{0}^{(2)}$ and $c_{1}^{(2)}$. Note this process might fail (if $z_{3}^{\prime} \neq 0$ ).

$$
Q D_{2}^{-1} Q^{T}\left[\begin{array}{c}
z_{1} \\
z_{2} \\
0
\end{array}\right]=\left[\begin{array}{c}
\times \\
\times \\
\times
\end{array}\right],
$$

and

$$
Q D_{2}^{-1} Q^{T}\left[\begin{array}{l}
0 \\
0 \\
1
\end{array}\right]=\left[\begin{array}{c}
\times \\
\times \\
0
\end{array}\right] .
$$

(Here $\times \neq 0$, which will not happen for the GMRES method since $D_{2}=I$.)

In the computation of the MGMRES methods, we assume $A$ is nonsingular, $Y$ is symmetric and nonsingular, $Z=A^{T} Y$, and $n^{*} \leq m$, which is the index of the $r^{(0)}$ vector. In Phase I, the MGMRES method is $\left(n^{*}, u^{(0)}\right)$ computable if and only if $\Delta_{n}\left(Y, r^{(0)}\right)$ is strongly regular. (This condition is not required for the $\operatorname{ORTHODIR}\left(A^{T} Y\right)$ method.) In Phase II, if $\Delta_{n}\left(Y, r^{(0)}\right)$ is strongly regular then the MGMRES method is $\left(n^{*}, u^{(0)}\right)$-computable if and only if $\Delta_{n}\left(A^{T} Y A, r^{(0)}\right)$ is strongly regular (that is, if the direct implementation of the $\operatorname{IGCG}\left(A^{T} Y\right)$ method is $\left(n^{*}, u^{(0)}\right)$-computable). (The $\operatorname{IGCG}\left(A^{T} Y\right)$ method is $\left(n^{*}, u^{(0)}\right)$-computable if and only if the $\operatorname{ORTHODIR}\left(A^{T} Y\right)$ method is ( $\left.n^{*}, u^{(0)}\right)$-computable.)

In a practical implication, if Phase I of the MGMRES method does not breakdown, and if the $\operatorname{IGCG}\left(A^{T} Y\right)$ method is $\left(n^{*}, r^{(0)}\right)$-computable, then SO is the MGMRES method.

## 5. Lan/mgmres Method

We now sketch a Lanczos-type method based on the MGMRES procedure. Consider the double system

$$
\left\{\begin{aligned}
A_{u} & =b \\
A^{T} \tilde{u} & =\tilde{b}
\end{aligned}\right.
$$

Here the second equation is called the shadow system for some $\tilde{b}$. We write the double system as

$$
\mathcal{A} \mathcal{U}=\mathcal{B},
$$

where

$$
\mathcal{A}=\left[\begin{array}{cc}
A & 0 \\
0 & A^{T}
\end{array}\right], \quad \mathcal{U}=\left[\begin{array}{c}
u \\
\tilde{u}
\end{array}\right], \quad \mathcal{B}=\left[\begin{array}{c}
b \\
\tilde{b}
\end{array}\right] .
$$

We can select $Z$ as either of the following symmetric matrices

$$
\mathcal{Y}=\left[\begin{array}{ll}
0 & I \\
I & 0
\end{array}\right], \quad \mathcal{Y} \mathcal{A}=\left[\begin{array}{cc}
0 & A^{T} \\
A & 0
\end{array}\right]
$$

To apply the MGMRES method, let $Z=A^{T} Y$ where $A=\mathcal{A}$ and $Y=\mathcal{Y}$.
Related Lanczos methods are the $\operatorname{LANDIR}(\mathcal{Y})$ method, the $\operatorname{LANDIR}\left(\mathcal{A}^{T} \mathcal{Y}\right)$ method (equivalently, the LAN/MGMRES method), the LANMIN(Y) method (equivalently, the BCG method), the LANMIN $\left(\mathcal{A}^{T} \mathcal{Y}\right)$ method, the LANRES $(\mathcal{Y})$ method, and the LANRES $\left(\mathcal{A}^{T} \mathcal{Y}\right)$ method.

We discuss the motivation for the LAN/MGMRES method. Let $Z=A^{T} Y$ and $Y$ is a SPD matrix. The methods $\operatorname{ORTHODIR}(Y)$ and $\operatorname{ORTHODIR}\left(A^{T} Y\right)$ are more robust than the methods ORTHOMIN $(Y)$ and ORTHOMIN $\left(A^{T} Y\right)$, respectively, but they are often numerically unstable. The $\operatorname{GMRES}\left(A^{T} Y\right)$ method is mathematically equivalent to the $\operatorname{ORTHODIR}\left(A^{T} Y\right)$ method, but is more stable and requires less work per iteration.

Let

$$
\mathcal{Z}=\mathcal{A}^{T} \mathcal{Y}, \quad \mathcal{Y}=\left[\begin{array}{ll}
0 & I \\
I & 0
\end{array}\right]
$$

In theory, the methods $\operatorname{LANDIR}(\mathcal{Y})$ and $\operatorname{LANDIR}\left(\mathcal{A}^{T} \mathcal{Y}\right)$ are more robust than the ORTHOMIN $(\mathcal{Y})$ method (equivalently, the BCG method) and the method ORTHOMIN $\left(\mathcal{A}^{T} \mathcal{Y}\right)$, respectively, but they are often numerically unstable. The method LAN/MGMRES (equivalently, the $\operatorname{MGMRES}\left(\mathcal{A}^{T} \mathcal{Y}\right)$ method) is almost equivalent to the $\operatorname{LANDIR}\left(\mathcal{A}^{T} Y\right)$ method and is hopefully more stable. (However, an additional condition is needed so that Phase I of the LAN/MGMRES method can be carried out.)

We now outline Phase I of the LAN/MGMRES method. Let $u^{(0)}$ be arbitrary and compute $r^{(0)}=b-A u^{(0)}$. Let $\tilde{u}^{(0)}$ be arbitrary or set $\tilde{u}^{(0)}=u^{(0)}$ for the shadow system and compute $\tilde{r}^{(0)}=\tilde{b}-A^{T} \tilde{u}^{(0)}$. Then let

$$
\left\{\begin{array}{l}
\widehat{w}^{(0)}=r^{(0)} \\
\widehat{\tilde{w}}^{(0)}=\tilde{r}^{(0)} .
\end{array}\right.
$$

Next set

$$
s_{0}=2\left\langle\widehat{w}^{(0)}, \widehat{\tilde{w}}^{(0)}\right\rangle
$$

(The process fails if $s_{0}=0$.) Then set $\sigma_{0}=\sqrt{\left|s_{0}\right|}$ and compute

$$
\left\{\begin{array}{l}
w^{(0)}=\sigma_{0}{ }^{-1} \widehat{w}^{(0)} \\
\tilde{w}^{(0)}=\sigma_{0}{ }^{-1} \widehat{\widetilde{w}}^{(0)}
\end{array}\right.
$$

and

$$
\left\{\begin{array}{l}
\widehat{w}^{(n)}=A w^{(n-1)}+\beta_{n, n-1} w^{(n-1)}+\beta_{n, n-2} w^{(n-2)} \\
\widehat{\tilde{w}}^{(n)}=A^{T} \tilde{w}^{(n-1)}+\beta_{n, n-1} \tilde{w}^{(n-1)}+\beta_{n, n-2} \tilde{w}^{(n-2)} .
\end{array}\right.
$$

Now we have

$$
\left\langle w^{(n)}, \tilde{w}^{(n-1)}\right\rangle=\left\langle w^{(n)}, \tilde{w}^{(n-2)}\right\rangle=0 .
$$

And set $s_{n}=2\left\langle\widehat{w}^{(n)}, \widehat{\widetilde{w}}^{(n)}\right\rangle$. (Process fails if $s_{n}=0$.) Set $\sigma_{n}=\sqrt{\left|s_{n}\right|}$. Finally, we have

$$
\left\{\begin{aligned}
w^{(n)} & =\sigma_{n}{ }^{-1} \widehat{w}^{(n)} \\
\tilde{w}^{(n)} & =\sigma_{n}{ }^{-1} \widehat{\widetilde{w}}^{(n)}
\end{aligned}\right.
$$

We now outline Phase II of the LAN/MGMRES method just for the nonshadow system. We have

$$
\begin{aligned}
u^{(n)} & =u^{(0)}+c_{0} w^{(0)}+c_{1} w^{(1)}+\cdots+c_{n-1}^{(n)} w^{(n-1)} \\
& =u^{(0)}+W_{n-1} c^{(n)} \\
& =\tilde{u}^{(n)}+\widetilde{W}_{n-1} c^{(n)} .
\end{aligned}
$$

The last equation is for the shadow system. Here

$$
\begin{aligned}
W_{n-1} & =\left[\begin{array}{llll}
w^{(0)} & w^{(1)} & \cdots & w^{(n-1)}
\end{array}\right] \\
\widetilde{W}_{n-1} & =\left[\begin{array}{llll}
\tilde{w}^{(0)} & \tilde{w}^{(1)} & \cdots & \tilde{w}^{(n-1)}
\end{array}\right], \\
c^{(n)} & =\left[\begin{array}{llll}
c_{0}^{(n)} & c_{1}^{(n)} & \cdots & c_{n-1}^{(n)}
\end{array}\right]^{T} .
\end{aligned}
$$

So

$$
H_{n}^{H} D_{n} H_{n} c^{(n)}=H_{n}^{T} D_{n} e^{(n+1)} .
$$

Example ( $n=2$ ):

$$
\begin{gathered}
H_{2}=\left[\begin{array}{rr}
-\beta_{1,0} & -\beta_{2,0} \\
\sigma_{1} & -\beta_{2,1} \\
0 & \sigma_{2}
\end{array}\right], \quad D_{2}=\left[\begin{array}{ccc}
d_{1} & 0 & 0 \\
0 & d_{2} & 0 \\
0 & 0 & d_{3}
\end{array}\right], \quad\left(d_{i}= \pm 1\right), \\
c^{(2)}=\left[\begin{array}{c}
c_{0}^{(2)} \\
c_{1}^{() 2}
\end{array}\right], \quad e^{(3)}=\left[\begin{array}{c}
\sigma_{2} \\
0 \\
0
\end{array}\right] .
\end{gathered}
$$

Use Givens rotations to find $Q$ with $Q Q^{T}=I$ and apply it to

$$
Q H_{2}=\tilde{H}_{2}=\left[\begin{array}{cc}
\boxed{x} & \times \\
0 & \boxed{x} \\
0 & 0
\end{array}\right]
$$

Solve for $c^{(2)}$ in

$$
\tilde{H}_{2}^{T} Q D_{2} Q^{T} \tilde{H}_{2} c^{(2)}=\tilde{H}_{2}^{T} Q D_{2} e^{(3)} .
$$

This process may fail. However, if Phase I is computable, then Phase II is computable if and only if LAN/IGCG $\left(\mathcal{A}^{T} \mathcal{Y}\right)$ is computable.

Additional details on the methods sketched in this paper can be found in $[1,2]$.

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