# Bi-parental product algorithm for coded waveform design in radar 

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#### Abstract

The problem of obtaining long sequences with finite alphabet and peaky aperiodic auto-correlation is important in the context of radar, sonar and system identification and is called the coded waveform design problem, or simply the signal design problem in this limited context. It is good to remember that there are other signal design problems in coding theory and digital communication. It is viewed as a problem of optimization. An algorithm based on two operational ideas is developed. From the earlier experience of using the eugenic algorithm for the problem of waveform design, it was realised that rather than random but multiple mutations, all the first-order mutations should be examined to pick up the best one. This is called Hamming scan, which has the advantage of being locally complete, rather than random. The conventional genetic algorithm for non-local optimization leaves out the anabolic role of chemistry of allowing quick growth of complexity. Here, the Hamming scan is made to operate on the Kronecker or Chinese product of two sequences with best-known discrimination values, so that one can go to large lengths and yet get good results in affordable time. The details of the ternary pulse compression sequences obtained are given. They suggest the superiority of the ternary sequences.


Keywords. Coded waveform design; global optimization; bi-parental products; Hamming scan.

## 1. Introduction

The term signal design has different meanings in coding theory, spread-spectrum communication, and radar. Here it is short for a coded waveform design problem in radar.

This involves obtaining sequences of finite lengths with small discrete alphabet, which are good approximations to white noise (Golay 1977). In particular, the alphabet ( $-1,0,1$ ) is considered. The sequences using them are called ternary sequences. Much of the earlier work deals with binary sequences with $(-1,1)$ as the alphabet. The goodness of the degree of approximation is measured here in terms of the discrimination. An alternative measure is the merit factor (Golay 1977).

To set the notation and to concretize the ideas, let

$$
\begin{equation*}
\mathbf{s}=\left[s_{0}, s_{1}, \ldots, s_{n-2}, s_{n-1}\right] \tag{1}
\end{equation*}
$$

be the sequence of length $n$, where the elements $s_{i}$ are from any of the two alphabets referred to above. Then

$$
\begin{equation*}
\rho(k)=\sum_{i=0}^{n-1-k} s_{i} s_{i+k}, \quad k=0,1, \ldots, n-1 \tag{2}
\end{equation*}
$$

is called the aperiodic auto-correlation of the sequence $\mathbf{s}$. The quantity,

$$
\begin{equation*}
D=\rho(0) /\left[\max _{k \neq 0}|\rho(k)|\right] \tag{3}
\end{equation*}
$$

is called the discrimination.
The problem of obtaining sequences with high discrimination is a very difficult one. It is so because whereas for the binary sequences with good periodic auto-correlation, analytical design procedures or construćtions (Baumert 1971; Golay 1983; Hoholdt et al 1985; Reddy \& Rao 1986; Jensen et al 1991) based on powerful number-theoretical results are available, for sequences with good aperiodic autocorrelations it has still to be basically a search. Further, the search cannot be exhaustive, because the number of sequences to be searched for grows combinatorially, as the length increases. Here, the emphasis is on obtaining as good sequences as possible with certain reasonably efficient procedures without claiming that even higher discrimination values are not possible. Golay's notion of sieves (Golay 1977) which restricts the search to a subset of sequences which can either be designed or searched efficiently and has properties which are desirable for enhancing the peakiness is a similar one. The sieves promise some good sequences without claiming that all good or even the best sequences are necessarily retained.

There is extensive work on binary sequences (Barker 1953; Turyn 1963, 1968; Boehmer 1967; Baumert 1971; Golay 1972, 1977, 1982, 1983; Moharir 1975; Beenker et al 1985; Hoholdt et al 1985; Kerdock et al 1986; Reddy \& Rao 1986; Bernasconi 1987, 1988; Hoholdt \& Justesen 1988; Golay \& Harris 1990; Newmann \& Byrnes 1990; Jensen et al 1991; De Groot et al 1992), which however we do not attempt to review here. Beyond the Barker sequences (Barker 1953) of length 11 and 13, having the discrimination values of 11 and 13 , the four best binary sequences having discrimination values of $14,15,16$ and 17 (Kerdock et al 1986) have been obtained at the smallest lengths of 28, 51, 69 and 88. Complete enumeration is possible only up to small lengths (Kerdock et al 1986) and would not be regarded as too satisfactory, the ultimate aim being construction rather than search. Partial searches using the notion of sieves (Golay 1977) were extended to higher lengths. Starting with binary sequences with ideal periodic autocorrelation and searching for sequences with the best discrimination among their cyclic shifts, many good sequences
have been obtained. To cite some examples (Reddy \& Rao 1986), discriminations of 22.82, $26.85,36.86,41.45,41.77$ and 42.46 have been obtained at the lengths of $251,349,811$, 829,919 and 1019 , respectively.

Optimization techniques have been used for the signal design problem (Bernasconi 1987, 1988; De Groot et al 1992), but the design criterion has been different, viz., merit factor.

Ternary aperiodic sequences have been proposed and listed earlier with merit factor as the criterion of goodness (Moharir 1974, 1976; Moharir et al 1985). They were obtained using sieves such as terminal admissibility (Moharir 1976), skew-symmetry (Golay 1977), terminal admissibility for skew-symmetric sequences (Moharir et al 1985) etc. This earlier work establishing the superiority of ternary pulse compression sequences over binary sequences, if merit factor is accepted as a valid desideratum, is extended further in this paper with discrimination as a chosen criterion.

## 2. The algorithm and associated concepts

### 2.1 Hamming scan

An improved version of the genetic algorithm (Holland 1992; Michalewicz 1992) called eugenic algorithm was used (Singh et al 1996) to see whether better and longer ternary sequences could be obtained. Success was met with along both these directions. Yet it was not possible to go to very large lengths, as the search time requirement increased very fast. It was regarded that one should devise an optimization algorithm which is more efficient, even though possibly more suboptimal. The Hamming scan is one such algorithm. Genetic algorithms use random but possibly multiple mutations. Mutation is a term metaphorically used for a change in an element in the sequence. Thus, a single mutation is at a Hamming distance of one from the original sequence. The Hamming scan looks at all the Hamming -1 neighbours and picks up the one with the largest discrimination. If it is better than the original sequence with the chosen definition of goodness, the algorithm is recursively continued therefrom, as long as improvement is possible. Thus, an entirely probabilistic mechanism of mutation is replaced by a locally complete search. The Hamming scan is expedited, and hence, made applicable at large lengths, by not calculating the aperiodic autocorrelation of a Hamming neighbour ab initio, recognizing the fact that as only one element is different, only its different contributions need to be taken into account. Let the element $s_{j}$ be changed to $c_{j}$. As a result, let $\rho(k)$ change to $\rho^{\prime}(k)$. Then it can be shown that

$$
\begin{equation*}
\rho^{\prime}(0)=\rho(0)+\left(c_{j}^{2}-s_{j}^{2}\right) . \tag{4}
\end{equation*}
$$

and

$$
\begin{array}{ll}
\rho^{\prime}(k)=\rho(k)+\left(c_{j}-s_{j}\right) s_{j+k}+s_{j-k}\left(c_{j}-s_{j}\right), & k=1,2, \ldots, n-1 \\
& j=0.1, \ldots n-1 \tag{5}
\end{array}
$$

In (5), there are two correction terms. They have to be implemented with care. One way is to assume that $s_{p}$ is equal to zero if $p$ is outside $(0,1, \ldots, n-1)$. Alternatively, the correction term $s_{j-k}\left(c_{j}-s_{j}\right)$ is included only for $k=1,2, \ldots, n-j-1$ and the
correction term $s_{j-k}\left(c_{j}-s_{j}\right)$ is included only for $k=j+1, j+2, \ldots, n-1$. This idea is certainly trivial, but it has led to significantly increased efficiency, and hence, to search at longer lengths than would otherwise have been possible.

For the ternary sequences the actual implementation can be somewhat different. The elements in the sequence are $-1,0$ and 1 . Each one of them can mutate in two possible ways. The mutations $-1 \rightarrow 0,0 \rightarrow-1$ and $1 \rightarrow-1$ are considered to give the lower strand of Hamming neighbours and the other mutations $-1 \rightarrow 1,0 \rightarrow 1$ and $1 \rightarrow 0$ are regarded as giving the upper strand of Hamming neighbours. The best neighbours along these two strands were found separately. The idea is that if one strand gives improvement, the other strand may not even be considered in order to save time. As the results are certainly path-dependent and optimality is not guaranteed, efficiency is a valid determinant.

The Hamming scan yielded some better ternary sequences in reasonable time than were obtainable with the eugenic algorithm. However, the Hamming scan also became unaffordable at larger lengths.

### 2.2 Simon's principle through Kronecker and Chinese products

That is when Simon's principle (Koestler 1969; Simon 1981) suggested newer possibilities. It states that bigger systems evolve faster, when developed through the metastable intermediate subsystems, than if they are constructed $a b$ initio from the smallest components. In the present context, it took the form of using two sequences of the best discrimination values available and obtaining a sequence of much larger length from them, such that it already had better discrimination than would result from random choice. The actual mechanism is provided by bi-parental products (Moharir 1992) of two or more sequences. In particular, two products, viz., Kronecker product (Brewere 1978; Moharir 1992) and the Chinese product (Moharir 1977, 1992) are chosen. These products are said to be bi-parental because each element in the product depends exactly on one element each from the two component sequences.

The Kronecker product of two sequences

$$
\begin{equation*}
\mathbf{s}_{\mathbf{1}}=\left[s_{01}, s_{11}, \ldots, s_{(p-1) 1}\right], \quad \mathbf{s}_{\mathbf{2}}=\left[s_{02}, s_{12}, \ldots, s_{(q-1) 2}\right] \tag{6}
\end{equation*}
$$

of lengths $p$ and $q$ respectively, is a sequence $\mathbf{s}$ of length $p q$, defined as

$$
\begin{equation*}
\mathbf{s}_{\mathbf{k}}=\left(s_{01} \mathbf{s}_{\mathbf{2}}, s_{1 \mid \mathbf{s}_{\mathbf{2}}}, \ldots, s_{(p-1) \mid} \mathbf{s}_{\mathbf{2}}\right) \tag{7}
\end{equation*}
$$

The Kronecker product is not commutative.
The Chinese product of the two sequences of (6) is defined only when $p$ and $q$ are relatively prime, that is, when they do not have any common prime factor. The Chinese product is defined as

$$
\begin{equation*}
s_{\mathrm{Ck}}=s_{1 i} s_{2 j} \tag{8}
\end{equation*}
$$

where

$$
k=\left\{\begin{array}{l}
i \bmod p  \tag{9}\\
j \bmod q
\end{array}\right.
$$

and has the length $p q$. It is called the Chinese product because the solution of the congruence relation (9) is what the Chinese remainder theorem deals with. It is commutative.

Computationally, the Chinese product of the two sequences can be obtained easily by repeating the sequences $\mathbf{s}_{\mathbf{1}}$ and $\mathbf{s}_{\mathbf{2}}$ of relatively prime lengths $p$ and $q$ respectively, $q$ and $p$ times and then taking an element-wise (Schur) product (Moharir 1992).

The autocorrelation of $\mathbf{s}_{\mathbf{k}}$ can be expressed in terms of the autocorrelations of $\mathbf{s}_{\mathbf{1}}$ and $\mathbf{s}_{\mathbf{2}}$ as follows (Turyn 1968).

## Theorem 1.

$$
\begin{align*}
\rho_{s_{k}}\left(q k_{1}+k_{2}\right)= & \rho_{s_{1}}\left(k_{1}\right) \rho_{s_{2}}\left(k_{2}\right)+\rho_{s_{1}}\left(k_{1}+1\right) \rho_{s_{2}}\left(q-k_{2}\right) \\
& k_{1}=0,1 \ldots, p-1 ; \quad k_{2}=0,1, \ldots, q-1 . \tag{10}
\end{align*}
$$

The theorem shows that if the individual autocorrelations are good, so is the resultant autocorrelation, except at some lag values. It can further be shown that the discrimination of the Kronecker product of two sequences depends only on the discriminations of the component sequences. The attenuated minimum guarantee theorem below is important.

Theorem 2. If the discriminations of the two sequences $\mathbf{s}_{\mathbf{1}}$ and $\mathbf{s}_{2}$ are $D_{1}$, and $D_{2}$ respectively, with $\min \left(D_{1}, D_{2}\right)=D_{\text {min }}$ and $\max \left(D_{1}, D_{2}\right)=D_{\text {max }}$, then the discrimination $D_{\mathrm{K}}$ of both their Kronecker products $\mathbf{s}_{\mathbf{1}} \times \mathbf{s}_{\mathbf{2}}$ and $\mathbf{s}_{\mathbf{2}} \times \mathbf{s}_{\mathbf{1}}$ is bounded as

$$
\begin{equation*}
D_{\min } \geq D_{\mathrm{K}} \geq \alpha D_{\min }, \quad \alpha=\left(D_{\max } /\left(1+D_{\max }\right)\right) \tag{11}
\end{equation*}
$$

where $\alpha$ may be called the attenuator.
The following has been shown by Moharir (1977).
Theorem 3. The periodic autocorrelation of the Chinese product of two sequences is a Chinese product of their periodic autocorrelations.

A sequence obtained by a bi-parental product (Moharir 1992) of two or more sequences can later be further improved by a Hamming scan. The time requirement comes down considerably because the algorithm begins with a good starting point. The importance of nonlocal optimizers such as the genetic algorithm is that the end result depends on the starting point only weakly, and ideally, not at all. But this statement is concerned with only the ultimate reachability of the algorithm, and the utility of a good starting point in determining the time requirements should not be underestimated. This is particularly so when good starting points can be designed or devised. Second, global optimization also is a goal and not a reality. In any case, the Hamming scan is not a global optimization algorithm. Therefore, if a large starting discrimination can be obtained by a simple procedure, in addition to the time saved, many lower local optima of which the Hamming scan possibly could not come out would also be avoided.

The anabolic bi-parental product scheme really works well as can be seen from table 1. It analyses the role of the Kronecker and the Chinese products in the proposed algorithm. Some ternary sequences with very good discrimination were chosen. They were obtained from the ternary sequences with the best merit factors listed by Golay (1977) by recursive Hamming scan to improve the discriminations. Then their Kronecker and Chinese products were obtained. The initial discriminations $D_{1}$ and $D_{2}$ and the discriminations $D_{\mathrm{K}}$ and $D_{\mathrm{C}}$

Table 1. Analysis of the efficacy of the bi-parental products, viz. Kronecker and Chinese, in designing ternary sequences with large discrimination. The lengths of the component sequences are $n_{1}$ and $n_{2}$. Their discriminations are $D_{1}$ and $D_{2}$. The discrimination of the product sequence is $D$. The bi-parental product efficiency is $\eta$ and the exponent is $\gamma$. Whether the bi-parental product $P$ is Kronecker or Chinese is indicated by $K$ and $C$ respectively.

| $n_{1}$ | $n_{2}$ | $n$ | $D_{1}$ | $D_{2}$ | P | D | $\eta$ | $\gamma$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 31 | 31 | 961 | 13.0000 | 13.0000 | K | 12.7547 | 0.9811 | 0.4963 |
| 9 | 33 | 297 | 7.0000 | 12.5000 | K | 7.0000 | 0.7483 | 0.4352 |
| 33 | 33 | 1089 | 12.5000 | 12.5000 | K | 12.5000 | 1.0000 | 0.5000 |
| 24 | 111 | 2664 | 20.0000 | 20.0000 | K | 19.7531 | 0.9877 | 0.4979 |
| 24 | 147 | 3528 | 20.0000 | 20.0000 | K | 19.6721 | 0.9836 | 0.4972 |
| 24 | 159 | 3816 | 20.0000 | 21.8333 | K | 20.0000 | 0.9571 | 0.4928 |
| 11 | 13 | 143 | 11.0000 | 13.0000 | K | 11.0000 | 0.9199 | 0.4832 |
|  |  |  |  |  | C | 8.9375 | 0.7474 | 0.4413 |
| 13 | 15 | 195 | 13.0000 | 13.0000 | K | 13.0000 | 1.0000 | 0.5000 |
|  |  |  |  |  | C | 12.0714 | 0.9286 | 0.4856 |
| 11 | 17 | 187 | 11.0000 | 11.0000 | K | 11.0000 | 1.0000 | 0.5000 |
|  |  |  |  |  | C | 6.7222 | 0.6111 | 0.3973 |
| 11 | 23 | 253 | 11.0000 | 19.0000 | K | 11.0000 | 0.7609 | 0.4488 |
|  |  |  |  |  | C | 7.7407 | 0.5354 | 0.3831 |
| 19 | 31 | 589 | 13.0000 | 13.0000 | K | 13.0000 | 1.0000 | 0.5000 |
|  |  |  |  |  | C | 12.0714 | 0.9286 | 0.4856 |
| 13 | 33 | 429 | 13.0000 | 12.5000 | K | 12.5000 | 0.9806 | 0.4961 |
|  |  |  |  |  | C | 10.1563 | 0.7967 | 0.4554 |
| 29 | 33 | 957 | 8.6667 | 12.5000 | K | 8.6667 | 0.8327 | 0.4609 |
|  |  |  |  |  | C | 8.5526 | 1.2141 | 0.5381 |
| 24 | 145 | 3480 | 20.0000 | 23.2000 | K | 20.0000 | 0.9285 | 0.4879 |
|  |  |  |  |  | C | 11.2621 | 0.5228 | 0.3944 |
| 24 | 155 | 3720 | 20.0000 | 22.1667 | K | 20.0000 | 0.9499 | 0.4916 |
|  |  |  |  |  | C | 19.4161 | 0.9221 | 0.4867 |
| 15 | 17 | 255 | 13.0000 | 11.0000 | K | 11.0000 | 0.9199 | 0.4832 |
|  |  |  |  |  | C | 11.0000 | 0.9199 | 0.4832 |
| 11 | 25 | 275 | 11.0000 | 10.5000 | K | 10.5000 | 0.9770 | 0.4951 |
|  |  |  |  |  | C | 10.5000 | 0.9770 | 0.4951 |
| 13 | 31 | 403 | 13.0000 | 13.0000 | K | 13.0000 | 1.0000 | 0.5000 |
|  |  |  |  |  | C | 13.0000 | 1.0000 | 0.5000 |
| 9 | 11 | 99 | 7.0000 | 11.0000 | K | 7.0000 | 0.7977 | 0.4480 |
|  |  |  |  |  | C | 8.5556 | 0.9750 | 0.4942 |
| 9 | 23 | 207 | 7.0000 | 19.0000 | K | 7.0000 | 0.6070 | 0.3979 |
|  |  |  |  |  | C | 8.3125 | 0.7208 | 0.4330 |

(Continued onfacing page)
of these two products are listed. It can be seen that impressive starting points for the Hamming scan are available. The purpose of table 1 is, however, different. It compares the two products as the bases of the bi-parental product signal design. In table 1 , six situations can be identified. In view of theorem 2, there is not much uncertainty about what the Kronecker product can achieve. For some entries in table 1, the component lengths are such that only the Kronecker product is defined and the Chinese product is not defined as the two lengths have a common factor. For these sets of component lengths, only the Kronecker product can be used. For the next set of entries, $D_{\mathrm{C}}$ is less than $D_{\mathrm{K}}$. For the next four sets of entries, $D_{\mathrm{C}}$ equals $D_{\mathrm{K}}$, lies between $D_{\mathrm{K}}$ and $D_{\max }$, equals $D_{\mathrm{max}}$, and exceeds

Table 1. (Continued)

| $n_{1}$ | $n_{2}$ | $n$ | $D_{1}$ | $D_{2}$ | P | D | $\eta$ | $\gamma$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 9 | 29 | 261 | 7.0000 | 8.6667 | K | 7.0000 | 0.8987 | 0.4740 |
|  |  |  |  |  | C | 8.2727 | 1.0621 | 0.5147 |
| 13 | 29 | 377 | 13.0000 | 8.6667 | K | 8.6667 | 0.8165 | 0.4571 |
|  |  |  |  |  | C | 11.2667 | 1.0614 | 0.5126 |
| 15 | 29 | 435 | 13.0000 | 8.6667 | K | 8.6667 | 0.8165 | 0.4571 |
|  |  |  |  |  | C | 11.6552 | 1.0980 | 0.5198 |
| 25 | 29 | 725 | 10.5000 | 8.6667 | K | 8.6667 | 0.9085 | 0.4787 |
|  |  |  |  |  | C | 9.1000 | 0.9539 | 0.4895 |
| 17 | 25 | 425 | 11.0000 | 10.5000 | K | 10.5000 | 0.9770 | 0.4951 |
|  |  |  |  |  | C | 11.0000 | 1.0235 | 0.5049 |
| 21 | 25 | 525 | 8.5000 | 10.5000 | K | 8.5000 | 0.8997 | 0.4765 |
|  |  |  |  |  | C | 10.5000 | 1.1114 | 0.5235 |
| 21 | 31 | 651 | 8.5000 | 13.0000 | K | 8.5000 | 0.8086 | 0.4548 |
|  |  |  |  |  | C | 13.0000 | 1.2367 | 0.5452 |
| 17 | 21 | 357 | 11.0000 | 8.5000 | K | 8.5000 | 0.8790 | 0.4716 |
|  |  |  |  |  | C | 13.3571 | 1.3814 | 0.5712 |
| 19 | 21 | 399 | 13.0000 | 8.5000 | K | 8.5000 | 0.8086 | 0.4548 |
|  |  |  |  |  | C | 13.8125 | 1.3140 | 0.5580 |
| 17 | 27 | 459 | 11.0000 | 8.3333 | K | 8.3333 | 0.8704 | 0.4693 |
|  |  |  |  |  | C | 11.4583 | 1.1968 | 0.5398 |
| 17 | 29 | 493 | 11.0000 | 8.6667 | K | 8.6667 | 0.8876 | 0.4738 |
|  |  |  |  |  | C | 13.0000 | 1.3314 | 0.5628 |
| 19 | 29 | 551 | 13.0000 | 8.6667 | K | 8.6667 | 0.8165 | 0.4571 |
|  |  |  |  |  | C | 14.0833 | 1.3268 | 0.5599 |
| 21 | 29 | 609 | 8.5000 | 8.6667 | K | 8.5000 | 0.9903 | 0.4977 |
|  |  |  |  |  | C | 9.2083 | 1.0729 | 0.5164 |
| 27 | 29 | 783 | 8.3333 | 8.6667 | K | 8.3333 | 0.9806 | 0.4954 |
|  |  |  |  |  | C | 9.4203 | 1.1085 | 0.5241 |
| 15 | 31 | 465 | 13.0000 | 13.0000 | K | 13.0000 | 1.0000 | 0.5000 |
|  |  |  |  |  | C | 13.5200 | 1.0400 | 0.5076 |
| 25 | 33 | 825 | 10.5000 | 12.5000 | K | 10.5000 | 0.9165 | 0.4821 |
|  |  |  |  |  | C | 13.1250 | 1.1456 | 0.5279 |
| 31 | 33 | 1023 | 12.5000 | 12.5000 | K | 12.5000 | 0.9806 | 0.4961 |
|  |  |  |  |  | C | 15.4762 | 1.2141 | 0.5381 |

$D_{\max }$. Thus, the Chinese product has the possibilities of providing better starting points than those provided by the Kronecker product, for the bi-parental product algorithm for signal design.

It is useful to define some quantitative measures for the bi-parental products in the context of the objective function chosen. The Kronecker efficiency $\eta_{\mathrm{K}}$ (for discrimination) is defined as

$$
\begin{equation*}
\eta_{\mathrm{K}}=D_{\mathrm{K}} /\left(D_{1} D_{2}\right)^{1 / 2} \tag{12}
\end{equation*}
$$

The Kronecker exponent $\gamma_{\mathrm{K}}$ is defined by

$$
\begin{equation*}
D_{\mathrm{K}}=\left(D_{1} D_{2}\right)^{\gamma_{k}} \quad \text { or } \quad \gamma_{\mathrm{K}}=\log D_{\mathrm{K}} / \log \left(D_{1} D_{2}\right) \tag{13}
\end{equation*}
$$

Table 1 shows values of $\eta_{\mathrm{K}}$ and $\gamma_{\mathrm{K}}$. It can be seen that they lie in narrower ranges than is suggested by theorem 2 . A starting sequence of a particular large length can be obtained by
factoring the length differently. In such cases two thumb rules could be used for guidance. the Chinese efficiency $\eta_{\mathrm{C}}$ and the Chinese exponent $\gamma_{\mathrm{C}}$ can be defined by replacing $D_{\mathrm{K}}$ by $D_{\mathrm{C}}$ in (10) and (11), where $D_{\mathrm{C}}$ is the discrimination obtained by using a Chinese product. Table 1 also shows $\eta_{\mathrm{K}}$ and $\gamma_{\mathrm{K}}$. The ranges of these are larger than those of $\eta_{\mathrm{K}}$ and $\gamma_{\mathrm{K}}$. The extension of the ranges on the higher sides must be viewed as advantageous.

### 2.3 Hamming scan and Simon's principle

Our interest was in obtaining ternary sequences with good discrimination values. The Kronecker product is not commutative. But the discriminations of $\mathbf{s}_{\mathbf{1}} \times \mathbf{s}_{\mathbf{2}}$ and $\mathbf{s}_{\mathbf{2}} \times \mathbf{s}_{\mathbf{1}}$ are the same. The Chinese product may have a superior, equal or inferior discrimination. Subsequently, however, they may evolve differently under a Hamming scan, so that the initial advantage of a superior discrimination may not last. When starting sequences of a particular length can be obtained by bi-parental products using different factors, further progress under a Hamming scan is also different. These are indicators that a Hamming scan does not give the global optimum. Let the ratio of the discrimination eventually obtained by recursively using a Hamming scan to the discrimination of the starting sequence obtained by abi-parental product be called the Hamming gain for that product. It is not as constant as the bi-parental product efficiency or the exponent. Let both the components in the Kronecker product be Barker sequence of length 11 . Then the discrimination of the resultant sequence of length 121 cannot be improved by the Hamming scan, meaning that the Hamming gain is just unity. But such situations are rather rare. Variability of the Hamming gain makes it difficult to use a good starting point as a very dependable criterion and offers no good thumb rules. One may make a choice on the basis of the average size of the improvement in the first few Hamming scans. But that is more a way of trying to minimize the computational effort than trying to estimate beforehand what discriminations may be achievable, as the improvement is not uniform over successive Hamming scans and the number of successful scans is also not known beforehand. It has been observed that the Hamming gain for the Chinese product is frequently less than that for the Kronecker product. It may also be noted that a sequence obtained by Hamming scan operating on a bi-parental product of two sequences, can later be used as a component in the same or different bi-parental product. All these possibilities have not been exhausted.

However, one more idea has been found to be very useful. After taking a bi-parental product of two good sequences, if the merit factor is improved by the recursive Hamming scan for some time and then the discrimination is taken to be the objective function to be

Table 2. The number of lengths $N_{T}$ at which ternary sequences exceeding the discrimination threshold of $D_{T}$ have been obtained.

| $D_{T}$ | $N_{T}$ | $D_{T}$ | $N_{T}$ | $D_{T}$ | $N_{T}$ |
| :--- | ---: | ---: | ---: | ---: | ---: |
| 48 | 2 | 44 | 7 | 40 | 27 |
| 36 | 47 | 34 | 75 | 32 | 119 |
| 30 | 170 | 28 | 242 | 26 | 304 |
| 24 | 344 | 22 | 367 | 20 | 389 |
| 18 | 414 | 16 | 434 | 14 | 464 |

Table 3. A list of ternary sequences, with discrimination values greater than 36, in the descending order of the discrimination.

| Length | Discrimi- <br> nation | Length | Discrimi- <br> nation | Length | Discrimi- <br> nation | Length | Discrimi- <br> nation |
| ---: | :--- | ---: | :--- | ---: | :--- | ---: | :--- |
| 1564 | 48.9048 | 924 | 42.4667 | 1584 | 40.2414 | 1001 | 37.6087 |
| 1131 | 48.2222 | 649 | 42.4615 | 590 | 40.0000 | 592 | 37.5455 |
| 1073 | 47.3125 | 975 | 41.7059 | 582 | 40.0000 | 555 | 37.4000 |
| 591 | 44.6667 | 1080 | 41.3889 | 584 | 39.8000 | 596 | 37.2727 |
| 599 | 44.4444 | 819 | 41.1538 | 1650 | 39.7667 | 593 | 36.9091 |
| 1632 | 44.2400 | 1188 | 40.9474 | 1680 | 39.6000 | 851 | 36.8125 |
| 1221 | 44.0588 | 1305 | 40.8889 | 1088 | 38.5714 | 1121 | 36.5789 |
| 1610 | 43.5385 | 726 | 40.8571 | 925 | 38.5625 | 1608 | 36.5556 |
| 1700 | 43.3214 | 623 | 40.7000 | 625 | 38.1818 | 1617 | 36.1389 |
| 1147 | 43.1579 | 858 | 40.6471 | 1071 | 37.9524 | 1053 | 36.1200 |
| 1518 | 43.0000 | 1750 | 40.6333 | 1173 | 37.7500 | 767 | 36.0000 |
| 999 | 42.5294 | 792 | 40.3846 | 1464 | 37.6429 |  |  |

increased by the same procedure, superior discrimination values frequently result, than if the discrimination was increased from the start.

## 3. Results

Only the sequence having the largest discrimination obtained at any length has been retained. The numbers of lengths at which various discrimination thresholds have been exceeded are shown in table 2. Thus the discrimination thresholds of 32,26 and 18 have been exceeded at more than 100,300 and 400 lengths respectively. The details about the sequences having the discrimination values of greater than 36 are shown in table 3 .

It is seen that as the length increases it is easier to reach higher discrimination values. This is a simple consequence of theorem 2 . To see this point, assume that the work of obtaining ternary sequences with good discrimination values is conducted up to the length of 160 . Then, the best sequences are of lengths $24,111,145,147,155$ and 159 , and have discriminations of $20,20,23,20,22.1667$ and 21.8333 respectively. That is, they all have discriminations exceeding 20. Then, theorem 2 implies that a Kronecker product of any two of them must have a discrimination exceeding 19.05. In general, the actual discrimination achieved is frequently much better than $\alpha D_{\min }$, as can be seen from table 1. Hamming scan can then raise the discrimination even further. It is very rare that the Hamming gain is just 1 . Thus, the highest discrimination obtained up to any length can almost always be exceeded at some higher lengths.

## 4. Conclusion

The bi-parental product algorithm has given ternary sequences with very good discriminations. The success of the algorithm indicates that the locally complete search may be preferable to the Monte Carlo search which depends rather excessively on the efficacy of chance in the matter of optimization. Whereas chance is a useful ally while dealing with combinatorially complex optimization problems, where feasible, it should be helped by
design procedures (Singh et al 1996) which can minimise the search effort. The anabolic role of bi-parental products which permitted going to large lengths quickly is particularly noteworthy in this regard. The algorithm has taken the problem of aperiodic signal design one step closer to that of the periodic signal design. The latter has two features. One of them is the availability of the regular construction procedures. That goal is still far away for the former problem. The second feature is that the discovery of a new good sequence means automatic construction of good sequences at many larger lengths. The bi-parental product procedure in the algorithm gives that feature to the aperiodic signal design problem also, though in a weaker sense, as the Hamming scan, which includes choice, is still to be performed.

To the best of our knowledge, the ternary sequences obtained here are the best. Yet, the bi-parental product algorithm is not a global optimisation algorithm. Therefore, there could be procedures which can improve upon the sequences obtained here.

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