Co-Simulation: A Survey

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Abstract

Modeling and simulation techniques are today extensively used both in industry and science. Parts of larger systems are, however, typically modeled and simulated by different techniques, tools, and algorithms. In addition, experts from different disciplines use various modeling and simulation techniques. Both these facts make it difficult to study coupled heterogeneous systems.

Co-simulation is an emerging enabling technique, where global simulation of a coupled system can be achieved by composing the simulations of its parts. Due to its potential and interdisciplinary nature, co-simulation is being studied in different disciplines but with limited sharing of findings.

In this survey, we study and survey the state-of-the-art techniques for co-simulation, with the goal of enhancing future research and highlighting the main challenges.

To study this broad topic, we start by focusing on discrete-event-based co-simulation, followed by continuous-time-based co-simulation. Finally, we explore the interactions between these two paradigms, in hybrid co-simulation.

To survey the current techniques, tools, and research challenges, we systematically classify recently published research literature on co-simulation, and summarize it into a taxonomy. As a result, we identify the need for finding generic approaches for modular, stable, and accurate, coupling of simulation units, as well as expressing the adaptations required to ensure that the coupling is correct.

1 Introduction

1.1 Motivation

Truly complex engineered systems that integrate physical, software, and network aspects are emerging, posing challenges in their design, operation, and maintenance.

The design of such systems, due to market pressure, has to be concurrent and distributed, that is, divided between different teams and/or external suppliers, each in its own domain and each with its own tools. Each participant develops a partial solution to a constituent system that needs to be integrated with all the other partial solutions. The later in the process the integration is done, the less optimal it is.

Innovative and optimal multi-disciplinary solutions can only be achieved through an holistic development process where the partial solutions developed independently are integrated sooner and more frequently, as each solution is refined. Furthermore, the traditional activities carried out at the partial solution level—such as requirements compliance check, or design space exploration—can be repeated at the global level, and salient properties spanning multiple constituent systems can be studied.

Modeling and simulation can improve the development of the partial solutions (e.g., see ), but falls short in fostering this holistic development process. To understand why, one has to observe that: (i) models of each partial solution cannot be exchanged or integrated easily, because these are likely developed by one of the many specialized tools deployed over the past 20 years; (ii) externally supplied models may have Intellectual Property (IP) that cannot be cheaply
disclosed to system integrators; and (iii) as solutions are refined, the system should be evaluated by integrating physical prototypes, software components, and even human operators, in what are denoted as Model/Software/Hardware/Human-in-the-loop simulations [7][185].

Consider now the interaction with, or operation of, a complex system. Such operation requires training, which, for safety or costs, may have to be conducted in a virtual environment. Developing a virtual environment is a difficult task [125] and reusing the models used in the development of the system allows the bulk of the effort to be redirected to where it is essential. Again, due to the aforementioned reasons, it may be difficult to obtain a single model of the whole system.

A high fidelity model of a system can also be used for maintenance of the system. Advanced sensory information, collected during the normal operation of the system, can be fed into a simulator to predict and prevent faults [111].

These are but a small sample of reasons for (and advantages of) being able to accurately compute the behavior of a coupled system. The fact that it should be carried out from a collection of interacting behaviors of the individual parts is what makes it a difficult challenge.

### 1.2 Co-simulation

Co-simulation consists of the theory and techniques to enable global simulation of a coupled system via the composition of simulators. Each simulator is broadly defined as a black box capable of exhibiting behaviour, consuming inputs and producing outputs. Examples of simulators include dynamical systems being integrated by numerical solvers [67], software and its execution platform [77], dedicated real-time hardware simulators (e.g., [128]), physical test stands (e.g., [276] Fig. 3), or human operators (e.g., [70] Fig. 24, [211] Fig. 6).

An alternative to co-simulation is co-modelling, were models are described in a unified language, and then simulated. There are advantages to this approach but each domain has its own particularities when it comes to simulation (e.g., see [67] [180] [267]) making it impractical to find a language and simulation algorithm that fits all.

As part of the systematic review that led to the current document (see section 6.1 for details), we took note of the approaches to co-simulation and the publications in applications of co-simulation. The approaches to co-simulation shaped the taxonomy in section 6.2 and the applications of co-simulation shows that in the last five years, co-simulation has been applied in many different engineering domains, as fig. 1 shows. In concrete, the publications are:

- **Automotive** - [4, 20, 26, 28, 32, 42, 49, 73, 81, 89, 144, 168, 181, 231, 258, 280, 286, 288]
- **HVAC** - [80, 89, 124, 202, 207, 272]
- **IC and SoC Design** - [227]
- **Maritime** - [209, 210]
- **Robotics** - [154, 215, 217, 257]

A closer look at the publications shows, however, that the average reported co-simulation scenario includes only two simulators, each a mock-up of a constituent system from a different domain. While this gives evidence that co-simulation enhances the development multi-domain systems, it is not yet up-to-par with the scale of Cyber-Physical Systems (CPSs). The unexplored potential is recognized in a number of completed and ongoing projects that address co-simulation (MODELISAR [1]).

[https://itea3.org/project/modelisar.html](https://itea3.org/project/modelisar.html)
DESTECS\textsuperscript{2}, INTO-CPS\textsuperscript{3}, ACOSAR\textsuperscript{4}, ACoRTA\textsuperscript{5}), and is one of the reasons why the Functional Mock-up Interface (FMI) Standard was created.

![Research publications of co-simulation applications over the past five years.](image)

**Figure 1:** Research publications of co-simulation applications over the past five years.

**Contribution.** We present a survey and a taxonomy, focused on the enabling techniques of co-simulation, as an attempt to bridge, relate, and classify the many approaches in the state of the art.

### 1.3 Need for the Survey

Despite the growing interest in the benefits and scientific challenges of co-simulation, to the best of our knowledge, no existing survey attempts to cover the heterogeneous communities in which it is being studied. The lack of such a survey means that the same techniques are being proposed independently with limited sharing of findings. To give an example, the use of dead-reckoning models is a well known technique in discrete event co-simulation \cite{164}, but only very recently it was used in a continuous time co-simulation approach \cite{240}.  

Our objective is to facilitate the exchange of solutions and techniques, highlight the essential challenges, and attain a deeper understanding of co-simulation.

Scientifically, co-simulation —multi-disciplinary by its very nature— mixes the following fields of research:

1. Numerical analysis – Accuracy and stability of the coupled system have to be studied \cite{10} \cite{56} \cite{71} \cite{118} \cite{119} \cite{121} \cite{140} \cite{214}.

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\textsuperscript{2}http://www.destecs.org/
\textsuperscript{3}http://into-cps.au.dk/
\textsuperscript{4}https://itea3.org/project/acosar.html
\textsuperscript{5}http://www.v2c2.at/research/ee-software/projects/acorta/
2. Differential Algebraic System Simulation – The composition of co-simulation units (SUs) is, in the most general sense, made through algebraic constraints [233, 238, 249].

3. Hybrid Systems – co-simulation scenarios, in the most general sense, are hybrid systems [53, 166, 186–188, 198, 272, 285].

4. Optimization – the heterogeneous capabilities of co-SUs pose interesting tradeoffs [52, 257].

5. Hierarchy – Systems of systems are hierarchical and the corresponding co-simulation scenarios should be hierarchical as well [105]. Compositionality properties of co-simulations becomes an interesting research challenge.

6. Formal Verification – The co-simulation orchestration algorithm, also known as the master, can be certified to be correct under certain assumptions about the co-simulation scenarios [52, 93, 108, 109, 193].

7. System Testing – Co-simulation can be used for exhaustively testing a set of black-box constituent systems, with a non-deterministic environment [162, 177].

8. Dynamic Structure Systems – Subsystems can have different dependencies depending on whom, and which level of abstraction, they interact with [20, 22, 255].

9. Multi-paradigm Modeling – Subsystems can have different models at different levels of abstraction [267]. The relationships between the multiple levels have to be known so that correct dynamic switching between levels abstraction can be made.

1.4 Outline

To help structure the characteristics of the simulators and how they interact, we distinguish two main approaches for co-simulation: Discrete Event (DE), described in section 3, and Continuous Time (CT), described in section 4. Both of these can, and are, used for the co-simulation of continuous, discrete, or hybrid coupled systems. We call Hybrid co-simulation, described in section 5, a co-simulation approach that mixes the DE and CT approaches. Section 6 summarizes the features provided by co-simulation frameworks, and classifies the state of the art with that taxonomy. Finally, section 7 concludes this publication. The section below provides the terminology used in the rest of the survey.

2 Modeling, Simulation, and Co-simulation

2.1 Dynamical Systems – Models of Real Systems

A dynamical system is a model of a real system (for instance a physical system or a computer system) characterized by a state and a notion of evolution rules. The state is a set of point values in a state space. The evolution rules describe how the state evolves over an independent variable, usually time. For instance, a traffic light system can be modeled as a dynamical system that can be in one state at a time. In this survey, we are focusing on timed formalisms (also called models of computation). Other formalisms, with no or only logical notion of time, are not discussed in this survey. For an overview of formalisms and models of computation, see [218] and [51].
of four possible states (red, yellow, green, or off). The evolution rules may dictate that it changes from red to green after some time (e.g., 60 seconds). Another example is a mass-spring-damper, modeled by a set of first order Ordinary Differential Equations (ODEs). The equations describe how the state—position and velocity of the mass—changes continuously over the simulated time. In contrast with the traffic light system, where the state cannot take an infinite number of different values over a finite duration of simulated time, the state of the mass-spring-damper can.

The behavior trace is the set of trajectories followed by the state (and outputs) of a dynamical system. For example, a state trajectory $x$ can be defined as a mapping between a time base $T$ and the set of reals $\mathbb{R}$, that is, $x : T \rightarrow \mathbb{R}$. fig. 2 shows a possible behavior trace for each of the example systems described before. In this example, the time base is $\mathbb{R}$.

![Figure 2: Examples of behavior traces.](image)

We refer to the time variable $t \in T$ as simulated time—or simply time, when no ambiguity exists—defined over a time base $T$ (typical the real numbers $\mathbb{R}$), as opposed to the wall-clock time $\tau \in WcT$, which is the time that passes in the real world. When computing the behavior trace of a dynamical system over an interval $[0, t]$ of simulated time, a computer takes $\tau$ units of wall-clock time that depend on $t$. $\tau$ can therefore be used to measure the run-time performance of simulators. fig. 3a highlights different kinds of simulation, based on the relationship between $\tau$ and $t$. In real-time simulation, the relationship between $t$ and $\tau$ is $t = \alpha \tau$, for a given $\alpha > 0$. In most cases $\alpha = 1$ is required, but making sure this is obeyed by the simulation algorithm is one of the main challenges in real-time simulation, and by extension, of co-simulation. In as-fast-as-possible—or analytical—simulation, the relationship between $\tau$ and $t$ is not restricted. Simulation tools that offer interactive visualization allow the user to pause the simulation and/or set a different value for $\alpha$.

Knowing when a dynamical system can be used to predict the behavior of a real system is crucial. The experimental frame describes, in an abstract way, a set of assumptions in which the behavior trace of the dynamical system can be compared with the one of the real system. By real system we mean either an existing physical system, or a system that does not yet exist. Validity is then the difference between the behavior trace of the dynamical system and the behavior trace of the real system, measured under the assumptions specified by the experimental frame. This is what conveys predictive power to dynamical systems. For example, Hooke’s law, in the mass-spring-damper system, can only be used to predict the reaction force of the spring for small deformations. For the traffic light dynamical system, the experimental frame includes the assumption that the transition from the red light to the green light is instantaneous. It is a valid assumption, provided that the executing platform in which the controller software runs, has enough computing power. A model is invalid when, within the experimental frame
assumptions, its behavior trace is so different than the one of the real system, that it cannot be used to predict properties of the real system.

In order to be practical, the behavior trace of a dynamical system has to highlight just the features of interest of the real system that are relevant for the tasks at hand [156]. In the traffic light model, the precise amount of wall-clock time a transition from red to green takes is unknown, but deemed small enough to be neglected. In the mass-spring-damper, Hooke’s law was chosen because the maximum displacement of the mass will not be large when the context in which the system will be used is taken into account.

Finally, we consider only those dynamical systems for which it is possible to obtain its meaning, i.e. the behavior trace, even if only an approximation.

2.2 Simulators – Computing the Behavior Trace

There are two generally accepted ways of obtaining the behavior trace of a dynamical system:

**Translational** Translate the dynamical system into another model, which can be readily used to obtain the behavior trace. Obtaining the analytical solution of the mass-spring-damper equations is an example of this approach. For instance, if the traffic light model is expressed in the Statechart formalism, it can be translated into a DEVS model, as done in [43], which can be used to obtain the behavior trace.

**Operational** Use of a solver – an algorithm that takes the dynamical system as input, and outputs a behavior trace. For the mass-spring-damper example, a numerical solver can be used to obtain an approximation of the behavior trace.

We focus on the latter.

A *simulator* (or solver) is an algorithm that computes the behavior trace of a dynamical system. If running in a digital computer, it is often the case that a simulator will only be able to approximate that trace. Two aspects contribute to the error in these approximations: inability to calculate a trajectory over a continuum, and the finite representation of infinitely small quantities. Simulators of discrete dynamical systems may also tolerate some inaccuracies in the behavior traces as well (e.g., if that brings a performance benefit). Fig. 3b shows an example approximation (dashed line) of the behavior trace (solid line) of the mass-spring-damper system, computed by the forward Euler method. Clearly, the trajectories differ.

In order to define what an *accurate simulator* is, or even be able to talk about error, we need to postulate that every dynamical system has an analytical behavior trace. The error can then be defined as the norm of the difference between the behavior trace produced by a simulator and the analytical trace. A simulator is accurate when the error is below a given threshold. Even if it is not possible to obtain the analytical behavior of every dynamical system, there are theoretical results that allow simulators to control the error they make. These techniques are applied to co-simulation in section 4.3. For the mass-spring-damper, and linear ODEs in general, the analytical trace follows a known structure [54]. For the traffic light, and timed statemachine models in general, the analytical behavior trace can be obtained with a sequential solver, that respects the causality of events. In short, validity is a property of a dynamical system whereas accuracy is a property of a simulator [67]. It is perfectly possible to have an accurate behaviour trace of a model that is invalid, and vice versa. For continuous time systems, the choice of an appropriate solver is important and should be made by domain experts [180][189].
2.3 Simulation Units - Mock-ups of Reality

In strict terms, a simulator is not readily executable: it needs a dynamical system and input trajectories, before being able to compute the behavior trace.

We use the term simulation unit (SU) to denote something that produces a behavior trace, when inputs are provided. A SU can be a composition of a simulator and a dynamical system, or it can be a real-world entity (with appropriate interface). Notice that, in contrast to a simulator, a SU only requires inputs to produce behavior.

A simulation is the behavior trace obtained with a SU. The correctness of a SU is dictated by the correctness of the simulation, which depends on the accuracy of the simulator and the validity of the dynamical system.

2.4 Compositional Co-simulation

As described in section 1, it is useful to obtain correct simulations of complex, not yet existing, systems as a combination of the behaviors of its constituent parts. Suppose each part is represented by a SU. Then these can be coupled via their inputs/outputs to produce a behavior trace of the coupled system. A co-simulation, a special kind of simulation, is the collection of combined simulations produced by the coupled SUs.

The SUs are independent black boxes. Hence, an orchestrator is necessary to couple them. The orchestrator controls how the simulated time progresses in each SU and moves data from outputs to inputs according to a co-simulation scenario. A co-simulation scenario is the information necessary to ensure that a correct co-simulation can be obtained. It includes how the inputs of each SU are computed from outputs, their experimental frames, etc.

Analogously to the simulator and SU concepts, the composition of a specific orchestrator with a co-simulation scenario, yields a co-SU, which is a special kind of SU, and a substitute of the real coupled system. It follows that a co-simulation is the simulation trace computed by a co-SU. This characterization enables hierarchical co-simulation scenarios, where co-SUs are coupled.
Co-simulation enables design decisions to be tried out in the model (what-if analysis), cheaply \[7\] early in the process, and possibly automatically \[117][92].

In this survey, we focus on the coupling techniques of black box SUs, where limited knowledge of the models and simulators is available. However, as will become clear in the following sections, the black box restriction has to be relaxed so that certain properties related to correctness can be ensured. Understanding what kind of information should be revealed and how IP can still be protected is an active area of research in co-simulation.

Most challenges in co-simulation are related to compositionality: if every SU \(S_i\) in a co-simulation scenario satisfies property \(P\), then the co-SU, with a suitable orchestrator, must also satisfy \(P\). The correctness is a property that should be compositional in co-simulation. Other properties include validity, or accuracy. It is an open research question to ensure that a co-simulator is compositional for a given set of properties. The following three sections provide an overview of the information and techniques being used throughout the state of the art, divided into three main approaches: discrete event (section 3), continuous time (section 4), and hybrid (section 5) co-simulation.

3 Discrete-Event-Based Co-simulation

The Discrete-Event-(DE)-based co-simulation approach describes a family of orchestrators and characteristics of simulation units (SUs) that are borrowed from the DE system simulation domain. We start with a description of DE systems, and then we extract the main concepts that characterize DE based co-simulation.

The traffic light is a good example of a DE system. It can be in one of the possible modes: red, yellow, green, or off. The off mode is often used by the police, which in some countries is characterized by a blinking yellow. Initially, the traffic light can be red. Then, after 60 seconds, it changes to green. Alternatively, before those 60 seconds pass, some external entity (e.g., a police officer) may trigger a change from red to off. The output of this system can be an event signaling its change to a new color. This example captures some of the essential characteristics of a DE dynamical system: reactivity – instant reaction to external stimuli (turning off by an external entity); and transiency – a DE system can change its state multiple times in the same simulated time point, and receive simultaneous stimuli. In the traffic light, transiency would happen if the light changes always after 0s (instead of 60s), or if the police officer would turn off and on the traffic light in the same instant.

These characteristics are embraced in DE based co-simulation, where the orchestrator acknowledges that SUs can change their the internal state and exchange values despite the fact that the simulated time is stopped.

3.1 DE Simulation Units

A DE SU is a black box that exhibits the characteristics of a DE dynamical system, but the dynamical system it stands for does not need to be a DE one. Furthermore, it is typical to assume that DE SUs communicate with the environment via time-stamped events, as opposed to signals. This means that the outputs of SUs can be absent at times where no event is produced.

\[7\] Another aspect to consider is the balance between insights gained and resources spent \[92].
We adapt the definition of the Discrete Event System Specification (DEVS) in [263] (originally proposed in [281]) to formally define a DE SU $S_i$, where $i$ denotes the reference of the SU:

$$S_i = (X_i, U_i, Y_i, \delta_i^{ext}, \delta_i^{int}, \lambda_i, t_{a_i}, q_i(0))$$

$$\delta_i^{ext} : Q_i \times U_i \rightarrow X_i$$

$$\delta_i^{int} : X_i \rightarrow X_i$$

$$\lambda_i : X_i \rightarrow Y_i \cup \{NaN\}$$

$$t_{a_i} : X_i \rightarrow \mathbb{R}_{\geq 0} \cup \{\infty\}$$

$$q_i(0) \in Q_i$$

$$Q_i = \{(x, e)|x \in X_i \text{ and } 0 \leq e \leq t_{a_i}(x)\}$$

where:

- $X_i, U_i,$ and $Y_i$ are the set of possible discrete states, input events, and output events, respectively;
- $\delta_i^{ext}(q_i, u_i) = x'_i$ is the external transition function that computes a new total state $(x'_i, 0) \in Q_i$ based on the current total state $q_i$ and an input event $u_i$;
- $\delta_i^{int}(x_i) = x'_i$ is the internal transition function that computes a new total state $(x'_i, 0) \in Q_i$ when the current total state is $(x_i, t_{a_i}(x_i)) \in Q_i$;
- $e$ denotes the elapsed units of time since the last transition (internal or external);
- $\lambda_i(x_i) = y_i \in Y_i \cup \{NaN\}$ is the output event function, invoked right before an internal transition takes place and $NaN$ encodes an absent value;
- $t_{a_i}(x_i) \in \mathbb{R}$ is the time advance function that indicates how much time passes until the next state change occurs, assuming that no external events arrive;
- $q_i(0)$ is the initial state.

The execution of a DE SU is described informally as follows. Suppose that the SU is at time $t_i \in \mathbb{R}_{\geq 0}$ and marks the current discrete state as $x_i$ for $e \geq 0$ elapsed units of time. Since $e \leq t_{a_i}(x_i)$, the total state is $(x_i, e) \in Q_i$. Let $t_n = t_i + t_{a_i}(x_i) - e$. If no input event happens until $t_n$, then at time $t_n$ an output event is computed as $y_i := \lambda_i(x_i)$ and the new discrete state $x_{i+1}$ is computed as $x_{i+1} = (\delta_i^{int}(x_i), 0)$. If, on the other hand, there is an event at time $ts < t_n$, that is, $u_i$ is not absent at that time, then the solver changes to state $x_{i+1} = (\delta_i^{ext}(x_i, e + ts - t_i), u_i), 0$ instead.

In the above description, if two events happen at the same time, both are processed before the simulated time progresses. Due to the transiency and reactivity properties, the state and output trajectories of a DE SU can only be well identified if the time base, traditionally the positive real numbers, includes a way to order simultaneous events, and simultaneous state changes. An example of such a time base is the notion of superdense time [166] [175], where each time point is a pair $(t, n) \in \mathbb{T \times N}$, with $\mathbb{T}$ typically being the positive real numbers and $\mathbb{N}$, called the index, is the set of natural numbers. In this time base, a state trajectory is a function $x_i : \mathbb{T \times N} \rightarrow V_{x_i}$, where $V_{x_i}$ is the set of values for the state, and an output/input trajectory is $u_i : \mathbb{T \times N} \rightarrow V_{u_i} \cup \{NaN\}$. Simultaneous states and events can be formally represented with incrementing indexes. See [53] for an introduction.

Eqs. (2) and (3) show instances of SUs represented in the adapted definition of DEVS.

---

8In the original DEVS definition, the initial state and the absent value in the output function are left implicit. Here we make them explicit, to be consistent with section 4. Note also that there are many other variants of DE formalisms. For instance, DE in hardware description languages (VHDL and Verilog) and actor based systems (for instance the DE director in Ptolemy II [218]).
A DE SU is passive: it expects some external coordinator to set the inputs and call the transition functions. This passivity enables an easier composition of SUs in a co-simulation, by means of a coordination algorithm, as will be shown later in section 3.2. Algorithm 1 shows a trivial orchestrator\(^9\) which computes the behavior trace of a single DE SU, as specified in eq. (1), that has no inputs. Remarks: \(tl\) holds the time of the last transition; and the initial elapsed time satisfies \(0 \leq e \leq ta_i(x_i(0))\).

If algorithm 1 is used to coordinate the execution of the traffic light SU in eq. (2), then the resulting behavior trace is the piecewise constant traffic light state \(x_1(t)\), together with the output events. The latter is represented as a trajectory \(y_i(t)\) that is mostly undefined (or absent), except for the single points where an output is produced, according to \(ta_1\).

---

**Algorithm 1:** Single autonomous DE SU orchestration.

**Data:** A \(S_i = (X_i, 0, Y_i, \delta_i^{ext}, \delta_i^{int}, \lambda_i, ta_i, (x_i(0), e_i))\).

\[
\begin{align*}
t_i &:= 0 ; \\
x_i &:= x_i(0) ; & \quad \text{Initial discrete state} \\
sl &:= -e_i ; & \quad \text{Account for initial elapsed time} \\
\text{while } \text{true do} & \\
  & t_i := tl + ta_i(x_i) ; & \quad \text{Compute time of the next transition} \\
  & y_i := \lambda_i(x_i) ; & \quad \text{Output} \\
  & x_i := \delta_i^{int}(x_i) ; & \quad \text{Take internal transition} \\
  & tl := t_i ; & \\
\text{end}
\end{align*}
\]

---

### 3.2 DE Co-simulation Orchestration

DEVS SUs communicate with their environment exclusively through inputs and outputs. DE co-simulation scenarios are comprised of multiple DE SUs (eq. (1)) coupled through output to input connections, which map output events of one SU to external events in other SU.

Consider the following DE SUs of a traffic light and a police office, respectively:

---

\(^9\)Algorithm 1 is based on [263] and is originally proposed in [281].
\[ S_1 = \langle X_1, U_1, Y_1, \delta_1^\text{ext}, \delta_1^\text{int}, \lambda_1, ta_1, q_1(0) \rangle \]
\[ X_1 = Y_1 = \{ \text{red}, \text{yellow}, \text{green}, \text{off} \} \]
\[ U_1 = \{ \text{toAuto}, \text{toOff} \} ; \quad q_1(0) = (\text{red}, 0) \]
\[ \delta_1^\text{ext}((x_1, e), u_1) = \begin{cases} \text{off} & \text{if } u_1 = \text{toOff} \\ \text{red} & \text{if } u_1 = \text{toAuto} \text{ and } x_1 = \text{off} \end{cases} \]
\[ \delta_1^\text{int}(x_1) = \begin{cases} \text{green} & \text{if } x_1 = \text{red} \\ \text{yellow} & \text{if } x_1 = \text{green} \\ \text{red} & \text{if } x_1 = \text{yellow} \end{cases} \]
\[ \lambda_1(x_1) = \begin{cases} 60 & \text{if } x_1 = \text{red} \\ 50 & \text{if } x_1 = \text{green} \\ 10 & \text{if } x_1 = \text{yellow} \\ \infty & \text{if } x_1 = \text{off} \end{cases} \]
\[ ta_1(x_1) = \begin{cases} 200 & \text{if } x_1 = \text{working} \\ 100 & \text{if } x_1 = \text{idle} \end{cases} \]

\[ S_2 = \langle X_2, U_2, Y_2, \delta_2^\text{ext}, \delta_2^\text{int}, \lambda_2, ta_2, q_2(0) \rangle \]
\[ X_2 = \{ \text{working}, \text{idle} \} \]
\[ U_2 = \emptyset \]
\[ Y_2 = \{ \text{toWork}, \text{toIdle} \} \]
\[ \delta_2^\text{ext}(x_2) = \begin{cases} \text{idle} & \text{if } x_2 = \text{working} \\ \text{working} & \text{if } x_2 = \text{idle} \end{cases} \]
\[ \delta_2^\text{int}(x_2) = \begin{cases} \text{toIdle} & \text{if } x_2 = \text{working} \\ \text{toWork} & \text{if } x_2 = \text{idle} \end{cases} \]
\[ \lambda_2(x_2) = \begin{cases} 200 & \text{if } x_2 = \text{working} \\ 100 & \text{if } x_2 = \text{idle} \end{cases} \]
\[ ta_2(x_2) = (\text{idle}, 0) \]

(2)

With the following remarks:
- The current state of the model in the definition of \( \delta_1^\text{ext} \) is \( q_1 = (x_1, e) \) with \( e \) being the elapsed time since the last transition.
- The output event function \( \lambda_1 \) is executed immediately before the internal transition takes place. It must then publish the next state instead of the current.

To model a scenario where the police officer interacts with a traffic light, the output events \( Y_2 \) have to be mapped into the external events \( U_1 \) of the traffic light SU (eq. (3)). In this example, if \( U_1 = \{ \text{toAuto}, \text{toOff} \} \) are the external input events handled by the traffic light SU, the mapping \( Z_{2,1} : Y_2 \to U_1 \) is defined by:

\[ Z_{2,1}(y_2) = \begin{cases} \text{toAuto} & \text{if } y_2 = \text{toIdle} \\ \text{toOff} & \text{if } y_2 = \text{toWork} \end{cases} \]

(4)

This way, if the police officer changes to working state at time \( t_n \), then the output signal \( y_2 := \text{toWork} \) will be translated by \( Z_{2,1} \) into an input event \( u_1 := \text{toOff} \) of the traffic light SU.

Based on the idea of abstract SUs [284], we formalize a DE co-simulation scenario with reference \( cs \) as follows:

\[ \langle U_{cs}, Y_{cs}, D, \{ S_d : d \in D \} \cup \{ I_d : d \in D \cup \{ cs \} \} \cup \{ Z_{i,d} : d \in D \wedge i \in I_d \} \rangle \]

(5)

where:
- \( U_{cs} \) is the set of possible input events, external to the scenario;
- \( Y_{cs} \) is the set of possible output events from the scenario to the environment;
- \( D \) is an ordered set of SU references;
• For each \( d \in D \), \( S_d \) denotes a DE SU, as defined in eq. (1);
• For each \( d \in D \cup \{cs\} \), \( I_d \subseteq (D \setminus \{d\}) \cup \{cs\} \) is the set of SUs that can influence \( S_d \), possibly including the environment external to the scenario \((cs)\), but excluding itself;
• For each \( i \in I_d \), \( Z_{i,d} \) specifies the mapping of events:
  \[
  Z_{i,d} : U_i \rightarrow U_d, \text{ if } i = cs \\
  Z_{i,d} : Y_i \rightarrow Y_d, \text{ if } d = cs \\
  Z_{i,d} : U_i \rightarrow U_d, \text{ if } i \neq cs \text{ and } d \neq cs
  \]
• Select : \( 2^D \rightarrow D \) is used to deterministically select one SU among multiple SUs ready to produce output events simultaneously, i.e., when at time \( t \), the set of SUs
  \[
  IMM(t) = \{d|d \in D \land q_d(t) = (x_d, ta_d(x_d))\}
  \]
  has more than one SU reference. This function is restricted to select one from among the set \( IMM(t) \), i.e., \( Select(IMM(t)) \in IMM(t) \).

The following co-simulation scenario \( cs \) couples the traffic light SU to the police officer SU:

\[
\langle \emptyset, Y_{cs}, \{1, 2\}, \{S_1, S_2\}, \{I_1, I_2, I_{cs}\}, \{Z_{2,1}, Z_{1,cs}\}, \text{Select} \rangle
\]

\[
Y_{cs} = Y_1; \quad I_1 = \{2\}; \quad I_2 = \emptyset; \quad I_{cs} = \{1\}; \quad Z_{1,cs}(y_1) = y_1
\]

where: \( S_1 \) is the traffic light SU and \( S_2 \) the police officer SU (eq. (3)); \( Y_1 \) is the output of \( S_1 \); \( Z_{2,1} \) is defined in eq. (4); and the omitted \( Z_{i,d} \) functions map anything to absent \( (NaN) \).

The Select function is particularly important to ensure that the co-simulation trace is unique. For example, consider the co-simulation scenario of eq. (7), and suppose that at time \( t_n \) both SUs are ready to output an event and perform an internal transition. Should the traffic light output the event and perform the internal transition first, or should it be the police office to do it first? In general, the order in which these output/transition actions are performed matters. The reason is that the way one SU reacts to the other SU’s output may be different, depending on the internal state of the former. In the example co-simulation scenario, the end result is always the same but this is not the general case.

Algorithm 2 illustrates the orchestrator of an autonomous (without inputs) DE co-simulation scenario\(^{10} \). It assumes that the co-simulation scenario does not expect external events, that is, all events that can affect the SUs are produced by other SUs in the same scenario. External output events are possible though. Remarks: \( t_{cs} \) holds the most recent time of the last transition in the scenario; \( e_d \) is the elapsed time of the current state \( q_d = (x_d, e_d) \) of \( S_d \); \( t_n \) is the time of the next transition in the scenario; \( i^* \) denotes the chosen imminent SU; \( I_{cs} \) is the set of SUs that can produce output events to the environment; \( y_{cs} \) is the output event signal of the scenario to the environment; and \( \{d \in D \land i^* \in I_d\} \) holds the SUs that \( S_{i^*} \) can influence.

fig. 4 shows the behavior trace of the traffic light in the co-simulation scenario of eq. (7).

Algorithm 2 is similar to algorithm 1: i) The time advance of the scenario \( ta_{cs} \) corresponds to the time advance of a single SU; ii) The output produced by the state transition is analogous to the \( \lambda \) function of a single SU; and iii) The output and state transition of child \( S_{i^*} \), together with the external transitions of the SUs influenced by \( S_{i^*} \), are analogous to the internal transition of a single SU. It is natural then that a co-simulation scenario \( cs \) as specified in eq. (4), can be

\(^{10}\text{Algorithm 2 is based on [263]}\)
Algorithm 2: Autonomous DE co-simulation scenario orchestration.

**Data:** A co-simulation scenario $cs = \langle \emptyset, Y_{cs}, D, \{S_d\}, \{I_d\}, \{Z_{i,d}\}, \text{Select} \rangle$.

$t_{cs} := 0$ ;

$x_i := x_i(0)$ for all $i \in D$ ; // Store initial discrete state for each unit

while true do

\[
\begin{align*}
ta_{cs} & := \min_{d \in D} \{ ta_d(x_d) - e_d \} ; \quad \text{// Time until the next internal transition} \\
tn & := t_{cs} + ta_{cs} ; \quad \text{// Time of the next internal transition} \\
i^* & := \text{Select}(\text{IMM}(tn)) ; \quad \text{// Get next unit to execute} \\
y_{i^*} & := \lambda_{i^*}(x_{i^*}) ; \\
x_{i^*} & := \delta_{i^*}(x_{i^*}) ; \quad \text{// Store new discrete state} \\
e_{i^*} & := 0 ; \\
\end{align*}
\]

if $i^* \in I_{cs}$ then

\[
\begin{align*}
y_{cs} & := Z_{i^*, cs}(y_{i^*}) ; \quad \text{// Compute output of the scenario} \\
\end{align*}
\]

end

for $d \in \{d | d \in D \land i^* \in I_d \}$ do

\[
\begin{align*}
u_d & := Z_{i^*, d}(y_{i^*}) ; \quad \text{// Trigger internal units that are influenced by unit $i^*$} \\
x_d & := \delta_{i^*}^{\text{ext}}((x_d, e_d + ta_{cs}), u_d) ; \\
e_d & := 0 ; \\
\end{align*}
\]

end

for $d \in \{d | d \in D \land i^* \notin I_d \}$ do

\[
\begin{align*}
e_d & := e_d + ta_{cs} ; \quad \text{// Update the elapsed time of the remaining units} \\
\end{align*}
\]

end

$t_{cs} := tn$ ; // Advance time

end

made to behave as a single DE SU $S_{cs}$. Intuitively, the state of $S_{cs}$ is the set product of the total states of each child DE SU; $ta_{cs}$ is the minimum time until one of the DE SUs executes an internal transition; the internal transition of $S_{cs}$ gets the output event of the imminent SU, executes the external transitions of all the affected SUs, updates the elapsed time of all unaffected SUs, and computes the next state of the imminent SU; the external transition of $S_{cs}$ gets an event from the environment, executes the external transition of all the affected SUs, and updates the elapsed time of all unaffected SUs, and updates the elapsed time
of all the unaffected SUs \[284\]. Formally:

\[ S_{cs} = \langle X_{cs}, U_{cs}, Y_{cs}, \delta_{cs}^{\text{ext}}, \delta_{cs}^{\text{int}}, \lambda_{cs}, t_{a_{cs}}, q_{cs}(0) \rangle \]

\[ X_{cs} = \times_{d \in D} Q_d \]

\[ q_{cs}(0) = (\times_{d \in D} q_i(0), \min_{d \in D} e_d) \]

\[ t_{a_{cs}}((\ldots, (x_d, e_d), \ldots)) = \min_{d \in D} \{ t_a(x_d) - e_d \} \]

\[ i^* = \text{Select}(\text{IMM}(t)) \]

\[ \lambda_{cs}(x_{cs}) = \begin{cases} Z_{i^*, cs}(y_{i^*}(tn)) & \text{if } i^* \in I_{cs} \\ \text{NaN} & \text{otherwise} \end{cases} \]

\[ \delta_{cs}^{\text{int}}(x_{cs}) = (\ldots, (x_d', e_d'), \ldots), \text{ for all } d \in D, \text{ where:} \]

\[ x_{cs} = (\ldots, (x_d, e_d), \ldots) \]

\[ (x_d', e_d') = \begin{cases} (\delta_{\text{int}}(x_d), 0) & \text{if } i^* = d \\ (\delta_{\text{int}}((x_d, e_d + t_{a_{cs}}(x_{cs})), Z_{i^*, d}(\lambda_{i^*}(x_{i^*})), 0) & \text{if } i^* \in I_d \\ (x_d, e_d + t_{a_{cs}}(x_{cs})) & \text{otherwise} \end{cases} \]

\[ \delta_{cs}^{\text{ext}}((x_{cs}, e_{cs}), u_{cs}) = (\ldots, (x_d', e_d'), \ldots), \text{ for all } d \in D, \text{ where:} \]

\[ x_{cs} = (\ldots, (x_d, e_d), \ldots) \]

\[ (x_d', e_d') = \begin{cases} (\delta_{\text{ext}}((x_d, e_d + e_{cs}), Z_{cs,d}(u_{cs})), 0) & \text{if } cs \in I_d \\ (x_d, e_d + e_{cs}) & \text{otherwise} \end{cases} \]
Remarks:
It is the Cartesian product of the total state of each child SU that makes the discrete state of the co-SU;
The elapsed times of each child SU are managed solely by the co-SU, whenever there is a transition (internal or external);
The external transition functions of each child are executed with the mapping of the events produced by the current state of the imminent child, and not the next one computed by \( (\delta^{out}_d(x_d), 0) \);
An internal transition of a child SU may cause an output event to the environment of the co-SU, if the child is connected to the output of the co-SU.

The same internal transition causes not only a change in the child discrete state, but also, due to its output event, may cause external transitions in other child SUs. This is not a recursive nor iterative process: at most one external transition will occur in all the affected child SUs; if any of the affected SUs becomes ready for an internal transition, it waits for the next internal transition invoked from the coordinator of the co-SU;

The resulting co-SU \( S_{cs} \) behaves exactly as a DE SU specified in eq. (1). It can thus be executed with algorithm 1 (in case of no inputs), or composed with other SUs in hierarchical co-simulation scenarios. Hierarchical co-simulation scenarios can elegantly correspond to real hierarchical systems, a natural way to deal with their complexity [149].

In summary, DE based co-simulation exhibits the following characteristics:

**reactivity:** A DE SU (analogously, a DE co-SU) has to process an event at the moment it occurs.

**transiency:** In both algorithm 2 and in a DE co-SU, the time advance \( ta_{cs} \) to the next imminent internal transition can be zero for successive iterations, so an orchestrator has to be able to tolerate the fact that simulated time may not advance for several iterations.

**predictable step sizes:** In a DE co-simulation scenario without inputs, the orchestrator, as shown in algorithm 2 can always predict the next simulated time step. In a scenario with inputs, if the environment provides the time of the next event, then the next simulated time step can be predicted too. For this to be possible, black box DE SUs have to be able to inform the orchestrator what their time advance is. This is not a trivial task for DE SUs that simulate continuous systems whose future behavior trace, especially when reacting to future inputs, is not easily predicted without actually computing it.

In the next subsection the main challenges in DE based co-simulation, and the requirements (or capabilities) their solutions impose in DE SUs, are made explicit.

### 3.3 Challenges

**Causality**
For the sake of simplicity, algorithm 2 is sequential. In a hierarchical co-SU, the imminent SU (closest to performing an internal transition) will be the one to execute, thus inducing that there is a global order in the events that are exchanged. This global order avoids causality violations but is too pessimistic. If an event \( y_1(t_1) \) causes another event —by changing the internal state of some other SU, which in turn changes its next output event— \( y_2(t_2) \), then \( t_1 \leq t_2 \), which is ok. However, the converse is not true: \( t_1 \leq t_2 \) does not necessarily imply that \( y_1(t_1) \) has caused \( y_2(t_2) \), which means that \( S_2 \) could execute before —in the wall-clock time sense— \( y_1(t_1) \) without violating causality, at least within a small window of simulated time. To see why, suppose that \( S_1 \)
and $S_2$ do not influence each other in the scenario. Then $y_2(t_2)$ would happen anyway, regardless of $y_1(t_1)$ occurring or not. Moreover, the co-simulation scenario holds information—the dependencies $\{I_d\}$—that can be used to determine who influences what.

A parallel optimistic orchestrator that takes $\{I_d\}$ into account is, in general, faster in the wall clock time sense, than a pessimistic, sequential one. However, most of these, the Time-warp algorithm being a well known example, require rollback capabilities of SUs. This is because SUs proceed to advance their own time optimistically, assuming that any other SUs will not affect them, until they are proven wrong by receiving an event which occurs before their own internal time. When that happens, the SU has to rollback to a state prior to the time of timestamp of the event that just arrived. This may in turn cause a cascade of rollbacks in other affected SUs. Moreover, in parallel optimistic DE co-simulation, any of the SUs in the scenario needs (theoretically) to support multiple rollbacks and have enough memory to do so for an arbitrary distant point in the past. This point in the past is limited in Time-warp by the Global Virtual Time (GVT). The GVT represents the minimum internal time of all SUs. By definition, no event that is yet to be produced (in wall-clock time) can have a timestamp smaller than the GVT.

We make a distinction between multiple rollback and single rollback capabilities. To support single rollback, a SU needs to store only the last committed state, thereby saving memory.

Causality is a compositionality property: if each child SU does not violate causality, then any orchestrator has to ensure that the causality is not violated when these SUs are coupled. Optimistic orchestration algorithms do so by requiring rollback capabilities from child SUs, whereas pessimistic algorithms do so at the cost of performance.

**Determinism and Confluence**

Determinism is also a compositional property. The Select function, in the co-simulation scenario definition of eq. (5), is paramount to ensure the compositionality of deterministic behavior. This function is used to ensure that a unique behavior trace can be obtained when the co-simulation scenario is executed by algorithm 2 or when it is turned into a co-SU, as in eq. (8). The alternative to the Select function is to ensure that all possible interleavings of executions always lead to the same behavior trace—this is known as confluence. Intuitively, if a co-SU is compositional with respect to confluence, then it is also compositional with respect to determinism.

Proving confluence is hard in general for black box DE co-simulation because it depends on knowledge about how the child SUs react to external events, which is potentially valuable IP. Parallel-DEVS is an approach, which leaves the confluence property to be satisfied by the modeler.

**Dynamic Structure**

Until now, the dependencies $\{I_d\}$, in eq. (5), have been assumed to be fixed over time. From a performance perspective, a static sequence of dependencies may be too conservative, especially if used to ensure causality in optimistic parallel co-simulation. To see why, consider that in a large scale simulation, there is a SU $S_1$ which may influence SU $S_2$ but only under a very specific set of conditions, which may not be verified until a large amount of simulated time has passed. A pessimistic co-SU assumes that $S_1$ may always affect $S_2$ and hence, tries to ensure that the simulated time of $S_2$ is always smaller than $S_1$, to minimize possible rollbacks. This incurs an unnecessary performance toll in the overall co-simulation because $S_1$ does not affect $S_2$ most of the time. This is where making $I_2$ dynamic can improve the performance of the co-simulation since...
the co-SU will know that most of the time, \( S_1 \) does not affect \( S_2 \). Dynamic structure co-simulation allows for \( \{I_d\} \) to change over time, depending on the behavior trace of the SUs. It can be used to study self-organizing systems \([256][20]\).

**Distribution**

Co-SUs whose child SUs are geographically distributed are common \([103]\). Interesting solutions like computation allocation \([194][261]\), bridging the hierarchical encapsulation \([262]\), and the use of dead-reckoning models \([164]\) have been proposed to mitigate the additional communication cost. Moreover, security becomes important, as pointed out, and addressed, in \([201]\).

4 Continuous-Time-Based Co-simulation

In the continuous time (CT) based co-simulation approach, the orchestrators’ and simulation units’ (SUs) behavior and assumptions are borrowed from the CT system simulation domain. We describe these below.

4.1 CT Simulation Units

A CT SU is assumed to have a state that evolves continuously over time. It is easier to get the intuitive idea of this by considering a SU of a CT dynamical system, such as a mass-spring-damper, depicted in the left hand side of fig. 5. The state is given by the displacement \( x_1 \) and velocity \( v_1 \) of the mass, and the evolution by:

\[
\begin{align*}
\dot{x}_1 &= v_1; \\
x_1(0) &= p_1; \\
m_1 \cdot v_1 &= -c_1 \cdot x_1 - d_1 \cdot v_1 + F_e \\
v_1(0) &= s_1
\end{align*}
\]

where \( \dot{x} \) denotes the time derivative of \( x \); \( c_1 \) is the spring stiffness constant and \( d_1 \) the damping coefficient; \( m_1 \) is the mass; \( p_1 \) and \( s_1 \) the initial position and velocity; and \( F_e \) denotes an external input force acting on the mass over time. The solutions \( x_1(t) \) and \( v_1(t) \) that satisfy eq. (9) constitute the behavior trace of the dynamical system. fig. 3b shows an example of such trace.

eq. (9) can be generalized to the state space form:

\[
\begin{align*}
\dot{x} &= f(x, u) \\
y &= g(x, u) \\
x(0) &= x_0
\end{align*}
\]

where \( x \) is the state vector, \( u \) the input and \( y \) the output vectors, and \( x_0 \) is the initial state.

A solution \([x(t), y(t)]^T\) that obeys eq. (10) is the behavior trace of the system. If \( f(x, u) \) is linear and time-invariant, an analytical form for \( x(t) \) can be obtained \([16]\). An analytical solution obtained by the application of mathematical identities is an example of a behavior trace obtained via the translational approach, described in section 2.2. Alternatively, the behavior trace can be computed.

If \( f(x, u) \) is sufficiently differentiable, \( x \) can be approximated with a truncated Taylor series \([24][67]\):

\[
x(t + h) = x(t) + f(x(t), u(t)) \cdot h + O(h^2)
\]

where

\[
\O(h^{n+1}) = \max_i \left( \lim_{h \to 0} \frac{x^{n+1} \left( \zeta(t^*) \right)}{(n + 1)!} h^{n+1} \right) = \text{const} \cdot h^{n+1}
\]
denotes the order of the truncated residual term; \( t^* \in [t, t + h] \); and \( h \geq 0 \) is the micro-step size. eq. (11) is the basis of a family of numerical solvers that iteratively compute an approximated behavior trace \( \tilde{x} \). For example, the forward Euler method is given by:

\[
\begin{align*}
\tilde{x}(t + h) &:= \tilde{x}(t) + f(\tilde{x}(t), u(t)) \cdot h \\
\tilde{x}(0) &:= x(0)
\end{align*}
\] (12)

A CT SU is assumed to have a behavior that is similar to one of a numerical solver computing a set of differential equations. We reinforce that this does not restrict CT SUs to being mockups of CT systems, even though it is easier to introduce them as such. For example, a SU \( S_1 \) of the mass-spring-damper, using the forward Euler solver, can be written by embedding the solver (eq. (12)) into eq. (9):

\[
\begin{align*}
\tilde{x}_1(t + h_1) &:= \tilde{x}_1(t) + \tilde{v}_1(t) \cdot h_1 \\
\tilde{v}_1(t + h_1) &:= \tilde{v}_1(t) + \frac{1}{m_1} \cdot (-c_1 \cdot \tilde{x}_1(t) - d_1 \cdot \tilde{v}_1(t) + F_c(t)) \cdot h_1 \\
\tilde{x}_1(0) &:= p_1 \\
\tilde{v}_1(0) &:= s_1
\end{align*}
\] (13)

where \( h_1 \) is the micro-step size, \( F_c(t) \) is the input, and \([x(t + h), v(t + h)]^T\) is the output.

### 4.2 CT Co-simulation Orchestration

Consider now a second system, depicted in the right hand side of fig. 5. It is governed by the differential equations:

\[
\begin{align*}
\dot{x}_2 &= v_2 \\
m_2 \cdot \ddot{v}_2 &= -c_2 \cdot x_2 - F_c \\
F_c &= c_c \cdot (x_2 - x_c) + d_c \cdot (v_2 - \dot{x}_c) \\
x_2(0) &= p_2 \\
v_2(0) &= s_2
\end{align*}
\] (14)

where \( c_c \) and \( d_c \) denote the stiffness and damping coefficients of the spring and damper, respectively; \( x_c \) denotes the displacement of the left end of the spring-damper. Combining with the forward Euler solver, yields the following SU:

\[
\begin{align*}
\tilde{x}_2(t + h_2) &:= \tilde{x}_2(t) + \tilde{v}_2(t) \cdot h_2 \\
\tilde{v}_2(t + h_2) &:= \tilde{v}_2(t) + \frac{1}{m_2} \cdot (-c_2 \cdot \tilde{x}_2(t) - \tilde{F}_c(t)) \cdot h_2 \\
\tilde{F}_c(t) &= c_c \cdot (\tilde{x}_2(t) - x_c(t)) + d_c \cdot (\tilde{v}_2(t) - \dot{x}_c(t)) \\
\tilde{x}_2(0) &:= p_2 \\
\tilde{v}_2(0) &:= s_2
\end{align*}
\] (15)

where \( h_2 \) is the micro-step size, \( x_c \) and \( \dot{x}_c \) are inputs, and \( \tilde{F}_c \) the output. Suppose \( S_1 \) (eq. (13)) and \( S_2 \) are coupled, setting \( x_c = x_1, \dot{x}_c = v_1 \) and \( F_c = F_{e_1} \), so that the resulting co-simulation scenario represents the multi-body system in fig. 5.
In the co-modeling approach, the models in Equations 9 and 14 would be combined to get the following coupled model:

\[
\begin{align*}
\dot{x}_1 &= v_1 \\
m_1 \cdot \dot{v}_1 &= -c_1 \cdot x_1 - d_1 \cdot v_1 + F_c \\
\dot{x}_2 &= v_2 \\
m_2 \cdot \dot{v}_2 &= -c_2 \cdot x_2 - F_c \\
F_c &= c_c \cdot (x_2 - x_1) + d_c \cdot (v_2 - v_1)
\end{align*}
\]  

(16)

which can be written in the state space form (eq. (10)) as:

\[
\begin{bmatrix}
\dot{x}_1 \\
\dot{v}_1 \\
\dot{x}_2 \\
\dot{v}_2
\end{bmatrix} =
\begin{bmatrix}
0 & 1 & 0 & 0 \\
-c_1 + c_c & -d_1 + d_c & -c_c & -d_c \\
0 & 0 & m_2 & 0 \\
c_c & d_c & 0 & m_2
\end{bmatrix}
\begin{bmatrix}
x_1 \\
v_1 \\
x_2 \\
v_2
\end{bmatrix}
\]

(17)

The behavior trace of eq. (17) can be obtained either analytically, or with forward Euler solver (eq. (12)).

In CT based co-simulation, to overcome the fact that each SU’s micro-step sizes are independent, a communication step size \( H \) (also known as macro-step size or communication grid size) has to be defined. \( H \) marks the times at which the SUs exchange values of inputs/outputs.

Suppose a SU \( S_i \) is at time \( n \cdot H \), for some natural \( n \), and is asked by an orchestrator to execute until time \( (n + 1) \cdot H \). If \( S_i \) only gets its inputs valued at \( n \cdot H \), then extrapolation must be used to
get the inputs in any of the internal micro-steps of the SU. In other words, when time is \( n \cdot H + m \cdot h_i \), for \( m \leq \frac{H}{h_i} \) and micro-step size \( h_i \), an extrapolation function \( \phi_{u_i}(m \cdot h_i, u_i(n \cdot H), u_i((n-1) \cdot H), \ldots) \), built from input values known at previous communication time points, is used to approximate the value of \( u_i(n \cdot H + m \cdot h_i) \). Notice that \( m = \frac{H}{h_i} \) is allowed, even though, theoretically, the value of \( u_i((n+1) \cdot H) \) can be obtained from the environment. The reason for this becomes clear in section 4.3. Analogously, interpolation techniques have to be used when the orchestrator makes the input value available at time \((n+1) \cdot H\) but the SU is still at time \( n \cdot H \). For example, the input \( F_c \) of the SU described in eq. (13) can be defined as:

\[
F_c(n \cdot H + m \cdot h_1) := \phi_{F_c}(m \cdot h_1, F_c(n \cdot H), F_c((n-1) \cdot H), \ldots), \text{ for } m \leq \frac{H}{h_1}
\]  

Similarly, the inputs \( x_c \) and \( \dot{x}_c \) of the SU described in eq. (15) can be defined as:

\[
\begin{align*}
x_c(n \cdot H + m \cdot h_2) &:= \phi_{x_c}(m \cdot h_2, x_c(n \cdot H), x_c((n-1) \cdot H), \ldots) \\
\dot{x}_c(n \cdot H + m \cdot h_2) &:= \phi_{\dot{x}_c}(m \cdot h_2, \dot{x}_c(n \cdot H), \dot{x}_c((n-1) \cdot H), \ldots)
\end{align*}
\]

for \( m \leq \frac{H}{h_2} \).

In the simplest case, the extrapolations can be constant. In the coupled mass-spring-dampers, this means:

\[
\phi_{F_c}(t, F_c(n \cdot H)) = F_c(n \cdot H); \quad \phi_{x_c}(t, x_c(n \cdot H)) = x_c(n \cdot H); \quad \phi_{\dot{x}_c}(t, \dot{x}_c(n \cdot H)) = \dot{x}_c(n \cdot H)
\]

In the state of the art, input extrapolation approaches can be classified as: Constant; Linear; Polynomial; Extrapolated-Interpolation [52][55]; Context-aware [29][30]; and Estimated Dead-Reckoning Model [24][48][242]. See [55][10][238][8] for an overview of linear and higher order extrapolation techniques and how these affect the accuracy of the co-simulation trace.

The orchestrator for this co-simulation scenario, at a time \( t = n \cdot H \), gets the outputs of both SUs and computes their inputs. Then, each SU is instructed to compute its behavior trace until the next communication step size, at \( t = (n + 1) \cdot H \), making use of the extrapolating functions to get the inputs at each of the micro steps (Equations 18 and 19).

We are now ready to formally define the behavior of a CT SU \( S_i \):

\[
\begin{align*}
S_i = & (X_i, U_i, Y_i, \delta_i, \lambda_i, x_i(0), \phi_{U_i}) \\
\delta_i : & \mathbb{R} \times X_i \times U_i \to X_i \\
\lambda_i : & \mathbb{R} \times X_i \times U_i \to Y_i \text{ or } \mathbb{R} \times X_i \to Y_i \\
x_i(0) & \in X_i \\
\phi_{U_i} : & \mathbb{R} \times U_1 \times \ldots \times U_i \to U_i
\end{align*}
\]

where:

- \( X_i \) is the state vector space;
- \( U_i \) is the input vector space;
- \( Y_i \) is the output vector space;
- \( \delta_i(t, x_i(t), u_i(t)) = x_i(t + H) \) or \( \delta_i(t, x_i(t), u_i(t+H)) = x_i(t + H) \) is the function that instructs the SU to compute a behavior trace from \( t \) to \( t + H \), making use of the input extrapolation (or interpolation) function \( \phi_{U_i} \);
\begin{itemize}
  \item $\lambda_i(t, x_i(t), u_i(t)) = y_i(t)$ or $\lambda_i(t, x_i(t)) = y_i(t)$ is the output function; and
  \item $x_i(0)$ is the initial state.
\end{itemize}

For instance, the SU in eq. \eqref{eq:13} can be described as follows:

$$S_1 = \langle \mathbb{R}^2, \mathbb{R}, \mathbb{R}^2, \delta_1, \lambda_1, [p_1]_{s_1}, \phi_{F_c} \rangle$$

$$\delta_1(t, [\tilde{x}_1(t), \tilde{v}_1(t)], F_c(t)) = \begin{bmatrix} \tilde{x}_1(t + H) \\ \tilde{v}_1(t + H) \end{bmatrix}$$

$$\lambda_1(t, [\tilde{x}_1(t), \tilde{v}_1(t)]) = \begin{bmatrix} \tilde{x}_1(t) \\ \tilde{v}_1(t) \end{bmatrix}$$

(22)

where $[\tilde{x}_1(t + H), \tilde{v}_1(t + H)]^T$ is obtained by the iterative application of the SU in eq. \eqref{eq:13} over a finite number of micro-steps, making use of the extrapolation of $F_c$ (defined in eq. \eqref{eq:18}):

$$\begin{bmatrix} \tilde{x}_1(t + H) \\ \tilde{v}_1(t + H) \end{bmatrix} = \begin{bmatrix} \tilde{x}_1(t) \\ \tilde{v}_1(t) \end{bmatrix} + \begin{bmatrix} \dot{\tilde{x}}_1(t) \\ \dot{\tilde{v}}_1(t) \end{bmatrix} \cdot h + \begin{bmatrix} \tilde{x}_1(t + h) \\ \tilde{v}_1(t + h) \end{bmatrix} \cdot h + \ldots$$

A CT co-simulation scenario with reference $cs$ includes at least the following information\footnote{Please note that this formalization is related to the formalization proposed by \cite{52}, with the main differences: i) it is not designed to formalize a subset of the FMI Standard, ii) it accommodates algebraic coupling conditions, and iii) it does not explicitly define port variables.}:

\begin{equation}
\begin{align*}
&\langle U_{cs}, Y_{cs}, D, \{S_i : i \in D\}, L, \phi_{U_{cs}} \rangle \\
&L : (\Pi_{i \in D} Y_i) \times Y_{cs} \times (\Pi_{i \in D} U_i) \times U_{cs} \to \mathbb{R}^m
\end{align*}
\end{equation}

(23)

where $D$ is an ordered set of SU references, each $S_i$ is defined as in eq. \eqref{eq:21}, $m \in \mathbb{N}$, $U_{cs}$ is the vector space of inputs external to the scenario, $Y_{cs}$ is the vector space of outputs of the scenario, $\phi_{U_{cs}}$ a set of input approximation functions, and $L$ induces the SU coupling constraints, that is, if $D = \{1, \ldots, n\}$, then the coupling is the solution to $L(y_1, \ldots, y_n, y_{cs}, u_1, \ldots, u_n, u_{cs}) = 0$, where $0$ denotes the null vector. As an example, the co-simulation scenario representing the system of fig. \ref{fig:5} is:

$$cs = \langle \emptyset, 0, \{1, 2\}, \{S_1, S_2\}, L, \emptyset \rangle; \quad L = [x_c - v_1; \dot{x}_c - x_1; F_c - F_c]^T$$

(24)

where:

\begin{itemize}
  \item $S_1$ is the SU for the constituent system on the left (eq. \eqref{eq:22}), and $S_2$ is the SU for the remaining constituent system;
  \item $x_c, \dot{x}_c$ are the inputs of $S_2$, and $F_c$ is the input of $S_1$; and
  \item $x_1, v_1$ are outputs of $S_1$ and $F_c$ is the output of $S_2$.
\end{itemize}

Algorithm \ref{alg:3} summarizes, in a generic way, the tasks of the orchestrator related to computing the co-simulation of a scenario $cs$ with no external inputs. It represents the Jacobi communication approach: SUs exchange values at time $t$ and independently compute the trace until the next communication time $t + H$. The way the system in eq. \eqref{eq:25} is solved depends on the definition of $L$. In the most trivial case, the system reduces to an assignment of an output $y_j(t)$ to each input $u_i(t)$, and so the orchestrator just gets the output of each SU and copies it onto the input of some other SU, in an appropriate order. Concrete examples of algorithm \ref{alg:3} are described in \cite{25,102,151,105,86,118,57,274}.

\footnotetext{Please note that this formalization is related to the formalization proposed by \cite{52}, with the main differences: i) it is not designed to formalize a subset of the FMI Standard, ii) it accommodates algebraic coupling conditions, and iii) it does not explicitly define port variables.}
An alternative to the Jacobi communication approach is the Gauss-Seidel (a.k.a. sequential or zig-zag) approach, where an order of the SUs’ $\delta$ function is forced to ensure that, at time $t$, they get inputs from a SU that is already at time $t + H$. Gauss-Seidel approach allows for interpolations of inputs, which is more accurate, but hinders the parallelization potential. Examples are described in [11][25][10][257].

**Algorithm 3:** Generic Jacobi based orchestrator for autonomous CT co-simulation scenarios.

**Data:** An autonomous scenario $cs = (\emptyset, Y_{cs}, D = \{1, \ldots, n\}, \{S_i\}, L, \emptyset)$ and a communication step size $H$.

**Result:** A co-simulation trace.

$t := 0$;

$x_i := x_i(0)$ for $i = 1, \ldots, n$;

while true do

Solve the following system for the unknowns:

\[
\begin{align*}
  y_1 &= \lambda_1(t, x_1, u_1) \\
  \vdots \\
  y_n &= \lambda_n(t, x_n, u_n) \\
  L(y_1, \ldots, y_n, y_{cs}, u_1, \ldots, u_n) &= 0
\end{align*}
\]  

$x_i := \delta_i(t, x_i, u_i)$, for $i = 1, \ldots, n$; // Instruct each SU to advance

$t := t + H$; // Advance time

end

Similarly to DE based co-simulation, a CT co-simulation scenario, together with an orchestrator, should behave as a (co-)SU of the form of eq. (21), and thus be coupled with other SUs, forming hierarchical co-simulation scenarios: the state of the co-SU is the set product of the states of the internal SUs; the inputs are given by $U_{cs}$ and the outputs by $Y_{cs}$; the transition and output functions are implemented by the orchestrator; the communication step size $H$ used by the orchestrator is analogous to a SU’s micro-step sizes, and the input extrapolation function is $\phi_{U_i}$.

Algorithm [3] makes it clear that the SUs can be coupled with very limited information about their internal details. In concrete:

- The output $\lambda_i$ and state transition $\delta_i$ functions need to be executable but their internal details can remain hidden;
- the inputs $u_i$ need to be accessible;
- the state variables can be hidden. These are represented merely to illustrate that the internal state of the SU changes when executing $\delta_i$.

However, the blind coupling can lead to compositionality problems, as will be discussed in the sections below. The common trait in addressing these is to require more from the individual SUs: either more capabilities, or more information about the internal (hidden) dynamical system.
4.3 Challenges

Modular Composition – Algebraic Constraints

In the co-simulation scenario described in eq. (24), the coupling condition \( L \) translates into a set of assignments from outputs to inputs. This is because the inputs of the SU of the system in the left hand side of fig. 5 and the outputs of the SU of the system represented in the right hand side of the same picture can be connected directly, and vice versa. In practice, the SUs’ models are not created with a specific coupling pattern in mind and \( L \) can be more complex. As an example, adapted from [233], consider the system coupled by a massless rigid link, depicted in fig. 6. The first subsystem is the same as the one in the left hand side of fig. 5 and its SU is in eq. (13). The second constituent system is governed by the following differential equations:

\[
\dot{x}_3 = v_3 \\
m_3 \cdot \dot{v}_2 = -c_3 \cdot x_3 + F_c \\
x_3(0) = p_3 \\
v_3(0) = s_3
\] (26)

And the following SU:

\[
\ddot{x}_3(t + h_3) = \ddot{x}_3(t) + v_3(t) \cdot h_3 \\
\ddot{v}_3(t + h_3) = \ddot{v}_3(t) + \frac{1}{m_3} \cdot (-c_3 \cdot x_3(t) + F_c(t)) \cdot h_3 \\
\ddot{x}_3(0) = p_3 \\
\ddot{v}_3(0) = s_3
\] (27)

The input to \( S_3 \) is the coupling force \( F_c \), and the output is the state of the mass \([\ddot{x}_3, \ddot{v}_3]^T\). The input to \( S_1 \) is the external force \( F_e \) and the outputs are the state of the mass \([\ddot{x}_1, \ddot{v}_1]^T\). Recall eq. (13). There is clearly a mismatch. The outputs \([\ddot{x}_1, \ddot{v}_1]^T\) of the first SU cannot be coupled directly to the input \( F_c \) of the second SU, and vice versa. However, the massless link restricts the states and inputs of the two SUs to be the same. Whatever the input forces may be, they are equal and opposite in sign. Hence, any orchestration algorithm has to find inputs that ensure the coupling constraints are satisfied:

\[
L = [\ddot{x}_1(n \cdot H) - \ddot{x}_3(n \cdot H); \quad \ddot{v}_1(n \cdot H) - \ddot{v}_3(n \cdot H); \quad F_e(n \cdot H) + F_c(n \cdot H)]^T = 0
\] (28)

This problem has been addressed in [118][119][11][10][238][233][239]. The approach taken in [118] is worth mentioning because it defines a Boundary Condition Coordinator (BCC) which behaves as an extra SU, whose inputs are the outputs of the original two SUs, and whose outputs are \( F_e \) and \( F_c \). They show that the initial co-simulation scenario with the non-trivial constraint can be translated into a co-simulation, with a trivial constraint, by adding an extra SU. This is illustrated in fig. 7.

Transforming the co-simulation scenario to make it simpler marks an important step in separating the concerns of the orchestrator [112]. In fact, the newly created SU can be run with a smaller internal micro-step size, required to meet stability and accuracy criteria, as shown in [118]. In many of the solutions proposed (e.g., [11][10][238][233][239]), information about the rate of change
(or sensitivity) of outputs and states of each SU, with respect to changes in its inputs is required
to solve the non-trivial coupling condition. This information can be either provided directly as a
Jacobian matrix of the system and output functions, or estimated by finite differences, provided
that the SUs can be rolled back to previous states. A frequent characteristic of co-simulation: the
availability of certain capabilities from SUs can mitigate the lack of other capabilities.

To show why the sensitivity information is useful, one of the tasks of the BCC is to ensure that
\[ \tilde{x}_1 - \tilde{x}_3 \] is as close to zero as possible, by finding appropriate inputs \( F_e \) and \( F_c \). This is possible
since \( \tilde{x}_1 \) and \( \tilde{x}_3 \) are functions of the inputs \( F_e \) and \( F_c \), and \( -F_e = F_c \). So the constraint can be
written as
\[
g(F_e) = \tilde{x}_1(F_e) - \tilde{x}_3(-F_e) = 0 \tag{29}
\]

From one communication step to the next, \( g \) can be expanded with the Taylor series:
\[
g(F_e((n + 1) \cdot H)) = g(F_e(n \cdot H) + \Delta F_e) \approx g(F_e(n \cdot H)) + \frac{\partial g(F_e(n \cdot H))}{\partial F_e} \cdot \Delta F_e \tag{30}
\]

From a known input \( F_e(n \cdot H) \), Equations \( 29 \) and \( 30 \) can be combined to obtain the input \( F_e((n + 1) \cdot H) \) at the next communication step:
\[
g(F_e(n \cdot H) + \Delta F_e) \approx g(F_e(n \cdot H)) + \frac{\partial g(F_e(n \cdot H))}{\partial F_e} \cdot \Delta F_e \leftrightarrow \]
\[
g(F_e(n \cdot H)) = -\frac{\partial g(F_e(n \cdot H))}{\partial F_e} \cdot \Delta F_e \leftrightarrow \]
\[
\Delta F_e = -\left[ \frac{\partial g(F_e(n \cdot H))}{\partial F_e} \right]^{-1} \cdot g(F_e(n \cdot H)) \leftrightarrow \]
\[
F_e((n + 1) \cdot H) = F_e(n \cdot H) - \left[ \frac{\partial g(F_e(n \cdot H))}{\partial F_e} \right]^{-1} \cdot g(F_e(n \cdot H)) \]
with
\[
\frac{\partial g(F_e(n \cdot H))}{\partial F_e} = \frac{\partial \tilde{x}_1(F_e(n \cdot H))}{\partial F_e} + \frac{\partial \tilde{x}_3(-F_e(n \cdot H))}{\partial F_e} \tag{32}
\]
A simple orchestration algorithm will then perform the following steps, at each co-simulation step:
1. Let $\tilde{x}_1(nH), \tilde{x}_3(nH)$ be the current position outputs of the two SUs $S_1$ and $S_3$; 
2. Perform a co-simulation step with a known $F_e$, obtaining $\tilde{x}_1^p(nH), \tilde{x}_3^p(nH)$ as new outputs.
3. Rollback SUs to state $\tilde{x}_1(nH), \tilde{x}_3(nH)$;
4. Perform a co-simulation step with $F_e + \Delta F_e$, obtaining $\tilde{x}_1^d(nH), \tilde{x}_3^d(nH)$;
5. Approximate $\frac{\partial g(F_e(nH))}{\partial F_e}$ by finite differences and eq. (32);
6. Obtain a corrected $F_c$ by eq. (31);
7. Rollback SUs to state $\tilde{x}_1(nH), \tilde{x}_3(nH)$;
8. Perform the final co-simulation step with $F_c$;
9. Commit states and advance time;

As can be seen in fig. 8, this coupling cannot be carried out without errors: the constraint $g(F_e((n+1) \cdot H))$ cannot be accurately forced to zero at first try. Furthermore, finding initial conditions and initial inputs that satisfy Equations 9, 26, and 28 is very important and usually requires a fixed point iteration. The above algorithm could be changed to perform an arbitrary number of iterations, repeating steps 1–7 until $g(F_c((n+1) \cdot H))$ is close enough to zero. This would increase the accuracy but also increase the amount of computation.

These examples show that rollback capabilities are important. If a SU is a black box, then the rollback capability has to be provided by the SU itself and there is little that the orchestrator can do to make up for the lack of the feature. See [52] for an orchestrator that takes into account the existence of the rollback feature. If, on the other hand, the SU provides access to its state, and allows the state to be set, as in [38], then the orchestrator can implement the rollback by keeping
track of the state of the SU. Rollback also plays a key role when dealing with algebraic loops in the co-simulation scenario.

Finally, to explain why this subsection refers to modular composition of SUs, the example in fig. 9 makes explicit one of the problems in co-simulation: the “rigid” and protected nature of SUs can make their coupled simulation very difficult. To contrast, in a white box approach where the equations of both constituent systems are available, the whole system is simplified, with the two masses being lumped together, and their coupling forces canceling each other out. The simplified system is a lumped mass-spring-damper, which is easily solvable. Such an approach is common in acausal modeling languages, such as Modelica [1]. To be concrete, the coupled system is obtained by combining Equations 9, 26, and 28 and simplifying to:

\[
\begin{align*}
x_1 &= v_1 \\
(m_1 + m_3) \cdot v_1 &= -(c_1 + c_3) \cdot x_1 - d_1 \cdot v_1 \\
x_1(0) &= p_1 \\
v_1(0) &= s_1
\end{align*}
\]  

(33)

fig. 9 compares the behavior trace produced by algorithm 3 when applied to the co-simulation scenario described in eq. (24), with the analytical solution, obtained from the coupled model of eq. (17) (co-modelling). It is obvious that there is an error due to the extrapolation functions and the large communication step size $H = 0.1$. In the white-box approach, the same constituent system can be coupled to other systems in many different contexts, whereas in co-simulation it is possible to get around the modularity aspect, but at a cost.

**Algebraic loops**

Algebraic loops occur whenever there is a variable that indirectly depends on itself. To see how algebraic loops arise in co-simulation scenarios, recall (see eq. (21)) that the state evolution and
output of each SU \( S_i \) can be written as:

\[
\begin{align*}
    x_i(t + H) &= \delta_i(t, x_i(t), u_i(t)) \\
    y_i(t + H) &= \lambda_i(t, x_i(t + H), u_i(t + H))
\end{align*}
\] (34)

To simplify things, assume that the SUs are coupled by a set of assignments from outputs to inputs, i.e.,

\[
u_i(t + H) = y_j(t)
\] (35)

where \( u_i \) is the input of SU \( S_i \) and \( y_j \) the output of a SU \( S_j \), in the same co-simulation scenario.

With these definitions, it is easy to see that, depending on the coupling assignments of the co-simulation scenario, the output of a SU may depend on itself, that is,

\[
\begin{align*}
    y_i(t + H) &= \lambda_i(t, x_i(t + H), u_i(t + H)) \\
    u_i(t + H) &= y_j(t + H) \\
    y_j(t + H) &= \lambda_j(t, x_j(t + H), u_j(t + H)) \\
    u_j(t + H) &= y_k(t + H) \\
    \vdots \\
    u_z(t + H) &= y_i(t + H)
\end{align*}
\] (36)

We distinguish two kinds of algebraic loops in co-simulation [152]: the ones spanning just input variables, and the ones that include state variables as well. The first kind arises when the outputs of a SU depend on its inputs, while the second kind happens when implicit numerical solvers are used, or when the input approximating functions are interpolations. In the previous example, the first kind can be removed by replacing \( u_i(t + H) \) in eq. (34) by the corresponding extrapolation \( \phi_{u_i}(H, u_i(n \cdot H), u_i((n - 1) \cdot H), \ldots) \) which does not depend on \( u_i((n + 1) \cdot H) \), thus breaking the algebraic loop. As shown in [152][13] (and empirically in [25]), neglecting a loop can lead to a prohibitively high error in the co-simulation. Instead, fixed point iteration technique should be
used to solve algebraic loops. For those involving state variables, the same co-simulation step has
to be repeated until convergence, whereas for loops over inputs/outputs, the iteration just repeats
the evaluation of the output functions.

To see how algebraic loops involving state variables arise, suppose that, in the example above,
\( \phi_{u_i} \) is constructed from \( u_i((n + 1) \cdot H) \):

\[
\phi_{u_i}(m \cdot h_i, u_i((n + 1) \cdot H), u_i(n \cdot H), u_i((n - 1) \cdot H), \ldots)
\]

(37)

If an order can be imposed in the evaluation of the SUs that ensures \( u_i((n + 1) \cdot H) \) can be computed
from some \( \lambda_j(t, x_j((n + 1) \cdot H), u_j((n + 1) \cdot H)) \) that does not indirectly depend on \( u_i((n + 1) \cdot H) \),
then this approach — Gauss-Seidel — can improve the accuracy of the co-simulation, as shown in
[13] [11] [55] [140] [10]. Obviously, the execution of SU \( S_j \) has to start after SU \( S_i \) has finished and
its output \( \lambda_j(t, x_j((n + 1) \cdot H), u_j((n + 1) \cdot H)) \) can be evaluated. If the input \( u_j((n + 1) \) depends
indirectly on \( u_i((n + 1) \cdot H) \), then an algebraic loop exists. The output function \( \lambda_j(t, x_j((n + 1) \cdot H), u_j((n + 1) \cdot H)) \)
depends on the state of the SU at \( x_j((n + 1) \cdot H) \), which in turn can only be
obtained by executing the SU from time \( n \cdot H \) to \( (n + 1) \cdot H \), using the extrapolation of the input \( u_j, \phi_{u_i}(m \cdot h_i, u_j((n + 1) \cdot H, \ldots) \); any improvement in the input \( u_j((n + 1) \cdot H, \) means that the whole
cosimulation step has to be repeated, to get an improved \( x_j((n + 1) \cdot H) \) and by consequence, an
improved output \( \lambda_j(t, x_j((n + 1) \cdot H), u_j((n + 1) \cdot H)) \).

An orchestrator that makes use of rollback to repeat the co-simulation step with corrected
inputs is called dynamic iteration, waveform iteration, and strong or on ion coupling [13] [251]. If
the SUs expose their outputs at every internal micro-step, then the waveform iteration can be used
[167]. Strong coupling approaches are typically the best in terms of accuracy, but worst in terms
of performance. Approaches that do not perform any correction steps are the best in terms of
performance, but worst in accuracy. A variant that attempts to obtain the middle-ground is the
so-called semi-implicit method, where a fixed limited number of correction steps is performed. See
[238] [233] for examples of this approach.

In the current FMI Standard for co-simulation, it is not possible, in the step mode, to perform
a fixed point iteration on the output variables only. A workaround is to use a strong coupling
 technique. That is, rollback the SUs and repeat the co-simulation step, effectively treating the
algebraic loop as involving the state variables too.

Until here, we have assumed full knowledge of the models being simulated in each SU to explain
how to identify, and deal with, algebraic loops. In practice, with general black-box SUs, extra
information is required to identify algebraic loops. According to [52] [13] [35], a binary flag denoting
whether an output depends directly on an input is sufficient. A structural analysis, for example,
with Tarjan’s strong component algorithm [245], can then be performed to identify the loops.

Consistent Initialization of Simulators

The definition of a SU in eq. (21) assumes that an initial condition is part of the SU. However, as seen
in the example of fig. 6, the initial states of the SUs can be coupled by algebraic constraints, through
the output functions, which implies that the initial states of the SUs cannot be set independently of
the co-simulation in which they are used. For example, the constraint in eq. (28) has to be satisfied
for the initial states:

\[ \{ \tilde{x}_1(0), \tilde{v}_1(0), \tilde{x}_3(0), \tilde{v}_3(0) \} \].
In general, for a co-simulation scenario as defined in eq. (23), there is an extra coupling function $L_0$ that at the time $t = 0$, has to be satisfied. For example:

$$L_0(x_1(0), \ldots, x_n(0), y_1(0), \ldots, y_n(0), y_{cs}(0), u_1(0), \ldots, u_n(0), u_{cs}(0)) = 0$$

(38)

where:

- $x_i(0)$ denotes the initial state of $S_i$; and
- $L_0 : X_1 \times \ldots \times X_n \times Y_1 \times \ldots \times Y_n \times U_1 \times \ldots \times U_n \to \mathbb{R}^m$ represents the initial constraint, not necessarily equal to $L$ in eq. (23).

eq. (38) may have an infinite number of solutions – as in the case of the example in fig. 6 – or have algebraic loops. The initialization problem (or co-initialization) is identified in [38] and addressed in [105]. In the FMI Standard, there is a dedicated mode for the (possibly fixed point iteration based) search of a consistent initial state in all SUs.

**Compositional Convergence – Error Control**

The accuracy of a co-simulation trace is the degree to which it conforms to the real trace as described in section 2.2. Obtaining the real trace can be a challenge. Error — the difference between the co-simulation trace and the real trace — is then a measure of accuracy.

In the context of co-simulation of CT systems, the most accurate trace is the analytical solution to the coupled model that underlies the co-simulation scenario. For example, the coupled model in eq. (17), corresponding to the multi-body system in fig. 5, is implicitly created from the co-simulation scenario described in eq. (24). Fortunately, the analytical solution can be obtained for this coupled model because it forms a linear time invariant system. In practice, the analytical solution for a coupled model cannot be found easily. Calculating the error precisely is therefore impossible for most cases but getting an estimate in how it grows is a well understood procedure in numerical analysis.

In simulation, the factors that influence the error are [67]: the model, the solver, the micro-step size, and, naturally, the size of the time interval to be simulated. In co-simulation, the extrapolation functions introduce error in the inputs of the SUs, which is translated into error in the state/outputs of these, causing a feedback on the error that can increase over time. Intuitively, the larger the co-simulation step size $H$, the larger is the error made by the extrapolation functions.

For example, when the forward Euler solver (eq. (12)) is used to compute the approximated behavior trace of the dynamical system in eq. (10), in a single micro step, it is making an error in the order of

$$\left\| \frac{x(t) + f(x(t)) \cdot h + O(h^2)}{\text{by infinite Taylor series}} - \frac{x(t) + f(x(t)) \cdot h}{\text{by forward Euler}} \right\| = O(h^2)$$

Obviously, the order in the error made at one step $O(h^2)$, most commonly called the local error, depends on:

- $f$ having no unbounded derivatives – to see why, observe that if the derivative of $f$ is infinite, then the residual term cannot be bounded by a constant multiplied by $h^2$. Fortunately, since most CT dynamic systems model some real system, this assumption is satisfied.
The solver used – other solvers, such as the midpoint method, are derived by truncating higher order terms of the Taylor series. For the midpoint method, the local truncation error is $O(h^3)$;

- Naturally, the larger the micro step size $h$ is, the larger the local error $O(h^2)$ is.

The local error assumes that the solver only made one step, starting from an accurate point $x(t)$. To compute the approximate behavior trace, the only accurate point the solver starts from is the initial value $x(0)$. The rest of the trace is approximate and the error gets compounded over the multiple steps. For the forward Euler method, if there is a limit to how $f$ acts on deviations on its parameter $\tilde{x}(t) = x(t) + e(t)$ from the true parameter $x(t)$, that is, if

$$||f(x(t)) - f(x(t) + e(t))|| \leq \text{const} \cdot e(t)$$

and const $< \infty$, then the order of the total accumulation of error can be defined in terms of the micro-step size. This condition is called global Lipschitz continuity [57]. For the forward Euler solver, the total (or global) error is $O(h)$.

For a solver to be useful, it must be convergent, that is, the computed trace must coincide with the accurate trace when $h \to 0$ [273]. It means the error can be controlled by adjusting the micro step size $h$. The same concept of convergence applies to co-simulation but does, as the intuition suggests, decreasing the communication step size $H$ lead to a more accurate co-simulation trace? This cannot be answered yet in general co-simulation because the behavior of the coupled model induced by the coupling of SUs may not satisfy Lipschitz continuity.

According to [124][57][152][11][13], if the SUs are convergent and the coupled model induced by the scenario coupling conditions can be written in the state space form of eq. (10), then the co-SU induced by any of the Jacobi, Gauss-Seidel, or Strong coupling methods, is convergent, regardless of the polynomial extrapolation technique used. Presence of algebraic loops, or complex coupling constraints, are factors that may make it impossible to write the coupled model in state space form [10].

The local error vector, in a co-simulation, is defined as the deviation from the true trace after one co-simulation step $H$, starting from an accurate point.

$$x_1(t + H) - \tilde{x}_1(t + H)$$
$$\ldots$$
$$x_n(t + H) - \tilde{x}_n(t + H)$$
$$y_1(t + H) - \tilde{y}_1(t + H)$$
$$\ldots$$
$$y_n(t + H) - \tilde{y}_n(t + H)$$

(39)

where $\tilde{x}_i(t + H) = \delta_i(t, x_i(t), \phi_{ui}(t))$; $\tilde{y}_i(t + H) = \lambda_i(t, \tilde{x}_i(t + H), \phi_{ui}(t + H))$, and $x_i(t + H)$ and $y_i(t + H)$ are the true state vectors and outputs, respectively, for SU $S_i$.

For a convergent co-SU, some of the techniques traditionally used in simulation, have been applied in co-simulation to estimate the error during the computation:

**Richardson extrapolation:** This well-known technique is compatible with black-box SUs as long as these provide rollback and state saving/restore capabilities [14][13][105]. The essential idea is to get an estimate of the local error by comparing $[\tilde{x}_i(t + H), \tilde{y}_i(t + H)]^T$ with a less accurate point $[\bar{x}_i(t + H), \bar{y}_i(t + H)]^T$. The less accurate point can be computed by the same orchestrator but using a larger communication step size. We have seen that larger communication step sizes affect the accuracy so if the two points are not too far apart, it
means the communication step $H$ does not need to be changed. It is importance to notice that the less accurate point $[\tilde{x}_i(t + H), \tilde{y}_i(t + H)]^T$ has to be computed from the accurate starting point $[\tilde{x}_i(t), \tilde{y}_i(t)]^T$.

**Multi-Order Input Extrapolation:** The outputs of two different order input approximation methods are compared [57] [59].

**Milne’s Device:** Similar to the previous ones, but the extrapolation of the inputs is compared with its actual value, at the end of the co-simulation step. Iterative approaches such as the ones studied in [238] [233] [236] [11] [10] can readily benefit from this technique.

**Parallel Embedded Method:** This technique runs a traditional adaptive step size numerical method in parallel with the co-simulation [131]. The purpose is to piggy back in the auxiliary method, the decisions on the step size. The derivatives being integrated in each SU have to be either provided, or estimated.

**Conservation Laws:** The local error is estimated based on the deviation from a known conservation law. Extra domain knowledge about the coupling between SUs is required. For example, if the couplings form power bonds [208], then energy should be conserved across a co-simulation step. In practice there is always an error due to the usual factors. The magnitude of the energy residual at a start and at end of a co-simulation step serves as an estimate of the local error. This technique has been implemented and studied in [225]. It has the advantage that it may not require rollback functionalities.

**Embedded Solver Method:** If the individual SUs support adaptive step size, then the decisions made internally can be made public to help the orchestrator decide on the communication step size. To the best of our knowledge, there is no orchestrator proposed that performs this, but the FMI Standard allows SUs to reject too large communication step sizes [38] [52].

After the error is deemed too large by one of the above methods, the correction can be applied pessimistically (rolling back and repeating the same step) or optimistically (adapt the next step). To mitigate the overhead of a pessimistic approach, the corrections may be applied only to sensitive SUs, as carried out in [270].

**Compositional Stability**

In the previous section we have presented conditions in which an orchestration engine can reduce the communication step size to an arbitrarily small value in order to meet arbitrary accuracy. Theoretically, this is useful as it tells the orchestrator that by reducing the local error, it also reduces the global error. In practice, the communication step size cannot be reduced to an arbitrarily small value without facing performance and roundoff error problems. Performance because, for smaller communication step sizes, it takes more steps to compute a behavior trace over a given interval of time. Round-off accuracy because in a digital computer, real numbers can only be represented approximately. Computations involving very small real numbers incur a non-negligible round-off error. So that means that in practice convergence does not imply that arbitrary accuracy can be achieved. A better question is to analyze what happens to the global error, as the co-simulation trace is computed with a non-null communication step size $H$.

Suppose that the analytical solution to the coupled model induced by the co-simulation scenario eventually goes to zero. This is the case for the coupled multi-body system of fig. 5 described in eq. (17), provided that at least one of the constants $d_1$ or $d_2$ is positive non-zero. Intuitively, this means that the system will lose energy over time, until it eventually comes to rest.

Let $x_1(t)$ denote the analytical solution of the position the mass $m_1$ in the system, and let $\tilde{x}_1(t)$
be the solution computed by a co-SU. Then $e_{x_1}(t) = \|x_1(t) - \tilde{x}_1(t)\|$ denotes the global error at time $t$ made by the co-SU. If $\lim_{t \to \infty} x_1(t) = 0$, then $\lim_{t \to \infty} e_{x_1}(t) = \tilde{x}_1(t)$.

If the co-SU is convergent, then for an arbitrarily small $H \to 0$, $\lim_{t \to \infty} e_{x_1}(t) = 0$ will be arbitrarily small too. Since in practice we cannot take arbitrarily small $H$, we want to know whether there is some non-zero $H$ such that $\lim_{t \to \infty} \tilde{x}_1(t) = 0$, thus driving $e_{x_1}(t)$ to zero as well. If that is the case, then it means that, assuming the system will eventually come to rest, the co-SU will too. This property is called numerical stability.

Contrarily to convergence, numerical stability is a property that depends on the characteristics of the system being co-simulated. Numerical stability is always studied assuming that the system being co-simulated is stable. It makes no sense to show that the co-simulation trace will grow unbounded provided that the system does too. It is a comparison of two infinities. One of the ways numerical stability in co-simulation can be studied is by calculating the spectral radius of the error in the co-SU, written as an autonomous linear discrete system [56].

To give an example, recall that the coupled model induced by the co-simulation scenario described in eq. (24) can be written as:

$$
\begin{align*}
\dot{x}_1(t) &= A_1 x_1(t) + B_1 u_1(nH) \\
y_1(t) &= C_1 x_1(t) + D_1 u_1(nH) \\
\dot{v}_1(t) &= A_1 v_1(t) + B_1 u_1(nH) \\
y_2(t) &= C_2 x_2(t) + D_2 u_2(nH)
\end{align*}
$$

(40)

with the coupling conditions $u_1 = y_2$ and $u_2 = y_1$.

In order to write the co-simulation model as an autonomous linear discrete system, we have to write what happens at a single co-simulation step $t \in [nH, (n+1)H]$ when executed by the orchestrator presented in algorithm 3. Since the purpose is to analyze the stability of a co-SU, and not the stability of each of the SUs in the co-simulation, it is common to assume that the SUs compute the analytical trace of the system. This enables the study of the stability properties of the co-SU, starting from stable SUs.

From time $t \in [nH, (n+1)H]$, SU $S_1$ is computing the behavior trace of the following Initial Value Problem Ordinary Differential Equation (IVP-ODE):

$$
\begin{bmatrix}
\dot{x}_1(t) \\
\dot{v}_1(t)
\end{bmatrix} = A_1 \begin{bmatrix} x_1(t) \\
v_1(t) \end{bmatrix} + B_1 u_1(nH) 
$$

(41)

with initial conditions $[x_1(nH) \quad v_1(nH)]^T$ given from the previous co-simulation step. The term $u_1(nH)$ denotes the fact that we are assuming a constant extrapolation of the input in the interval $t \in [nH, (n+1)H]$. 
eq. (41) is linear and time invariant, so the value of \[
\begin{bmatrix}
x_1((n+1)H) \\
v_1((n+1)H)
\end{bmatrix}
\] can be given analytically as:
\[
\begin{bmatrix}
x_1((n+1)H) \\
v_1((n+1)H)
\end{bmatrix} = e^{A_1H} \begin{bmatrix} x_1(nH) \\ v_1(nH) \end{bmatrix} + \left( \int_{nH}^{(n+1)H} e^{A_1((n+1)H-\tau)} d\tau \right) B_1 u_1(nH)
\]  
(42)
or, replacing the integration variable with \( s = \tau - nH \),
\[
\begin{bmatrix}
x_1((n+1)H) \\
v_1((n+1)H)
\end{bmatrix} = e^{A_1H} \begin{bmatrix} x_1(nH) \\ v_1(nH) \end{bmatrix} + \left( \int_{0}^{H} e^{A_1(H-s)} ds \right) B_1 u_1(nH)
\]  
(43)
where \( e^X = \sum_{k=0}^{\infty} \frac{1}{k!} X^k \) is the matrix exponential.

Rewriting eq. (43) as a discrete time system gives us the computation performed by SU \( S_1 \) in a single co-simulation step, that is, the state transition function \( \delta_1 \):
\[
\begin{bmatrix}
x_1^{(n+1)} \\
v_1^{(n+1)}
\end{bmatrix} = e^{A_1H} \begin{bmatrix} x_1^{(n)} \\ v_1^{(n)} \end{bmatrix} + K_1 B_1 u_1^{(n)}
\]  
(44)
where \( z^{(n)} = z(nH) \).

At the end of the co-simulation step (\( t = (n+1)H \)) the output of the first SU, that is, its output function \( \lambda_1 \), is given by plugging in eq. (44) to the output \( y_1 \) in eq. (40):
\[
y_1^{(n+1)} = C_1 e^{A_1H} \begin{bmatrix} x_1^{(n)} \\ v_1^{(n)} \end{bmatrix} + C_1 K_1 B_1 u_1^{(n)}
\]  
(45)

Repeating the same procedure for the second SU, yields the state transition \( \delta_2 \) and output functions \( \lambda_2 \):
\[
\begin{bmatrix}
x_2^{(n+1)} \\
v_2^{(n+1)}
\end{bmatrix} = e^{A_2H} \begin{bmatrix} x_2^{(n)} \\ v_2^{(n)} \end{bmatrix} + K_2 B_2 u_2^{(n)}
\]  
(46)
\[
y_2^{(n+1)} = C_2 e^{A_2H} \begin{bmatrix} x_2^{(n)} \\ v_2^{(n)} \end{bmatrix} + (C_2 K_2 B_2 + D_2) u_2^{(n)}
\]
with \( K_2 = \int_{0}^{H} e^{A_2(H-u)} du \).

Since the coupling conditions are \( u_1 = y_2 \) and \( u_2 = y_1 \), we can combine Equations 46, 45 and 41 into a single discrete time system:

\[
\begin{bmatrix}
x_1^{(n+1)} \\
x_2^{(n+1)} \\
y_1^{(n+1)} \\
y_2^{(n+1)}
\end{bmatrix} = e^{A_1H} \begin{bmatrix} 0 & 0 & K_1 B_1 \\ K_2 B_2 & 0 & 0 \\ 0 & C_2 K_2 B_2 + D_2 & 0 \\ 0 & C_2 e^{A_2H} & 0
\end{bmatrix} \begin{bmatrix}
x_1^{(n)} \\
x_2^{(n)} \\
y_1^{(n)} \\
y_2^{(n)}
\end{bmatrix}
\]  
(47)
The above system is stable if the behavior traces remain bounded (e.g., by going to zero) as \( n \to \infty \). This can be checked by observing whether the spectral radius \( \rho(A) < 1 \). For parameters \( m_1 = m_2 = c_1 = c_2 = d_1 = c_3 = d_c = 1, d_2 = 2 \), a communication step size of \( H = 0.001 \), \( \rho(A) = 0.9992 \), which means that the co-SU is stable. If the damping constant were \( d_c = 6.0E6 \), then the co-SU would be unstable (\( \rho(A) \approx 76.43 \)). A stable co-simulation is shown in fig. 10.

![Figure 10: Behavior trace of co-simulator described in eq. (47). Parameters are: \( m_1 = m_2 = c_1 = c_2 = d_1 = c_3 = d_c = 1, d_2 = 2, H = 0.001 \).](image)

Different coupling methods, and different approximation functions yield different stability properties. See [55][58][56][55] for the stability analysis of multiple coupling approaches and approximating functions. Stability of various co-SUs has been also studied in [153][236][140][118][10]. The rules of thumb drawn from these papers can be summarized as: (1) Co-simulators that employ fixed point iteration techniques typically have better stability properties; (2) Gauss-Seidel coupling approach has slightly better stability properties when the order in which the SUs compute is appropriate (e.g., the SU with the highest mass should be computed first [10]).

The main problem is that in co-simulation applied to industrial problems, the solvers and models may be coupled in a black box to protect IP, so there is little knowledge about the kind of solver and model being used and its stability properties. The best is then to always use iterative techniques that have been shown to have better stability properties. However, these techniques require rollback functionalities which can be difficult to support for certain SUs. Even if those functionalities are available, the cost of computing a co-simulation trace can be prohibitively high when compared with non-iterative approaches. This creates a paradox where industrial co-SUs should make use of iterative techniques but the performance toll may be too high.

**Compositional Continuity**

If a SU is a mock-up of a CT system, then it is reasonable to expect that its inputs are also continuous. As discussed in [225][55], the careless use of input extrapolations (e.g., constant extrapolation) may violate this assumption. Consider the point of view of a SU \( S_i \) in co-simulation. Throughout a co-simulation step \( t \in [nH, (n + 1)H] \) the input \( \phi_{ui}(t, u_i(nH)) = u_i(nH) \) is kept constant. At the next co-simulation step \( t \in [(n + 1)H, (n + 2)H] \), the input \( \phi_{ui}(t, u_i((n + 1)H)) = u_i((n + 1)H) \) may change radically if \( u_i((n + 1)H) \) is too far away from \( u_i(nH) \).
Any sudden change in the input to a CT SU may wreak havoc in the performance of its simulator, causing it to reduce inappropriately the internal micro step size, to reinitialize the solver [67], to discard useful information about the past (in multi-step solvers [8, 9]), and/or produce inaccurate values in its input extrapolation [206]. Furthermore, a discontinuity may be propagated to other SUs, aggravating the problem.

Most numerical methods assume that the input is a discretized version of a continuous trace. That means that, when a discontinuity occurs, SU $S_i$ cannot distinguish it from a very steep change in the continuous trace. The way traditional solvers deal with this behavior is to reduce the micro step size $h_i$ until the change is not so steep. This works with a continuous signal with a steep change, but does not work with a discontinuity: even if the micro-step size is reduced, the difference between $\lim_{t \to ((n+1)H)} \phi_{u_i}(t, u_i(nH)) = u_i(nH)$ and $\lim_{t \to ((n+1)H)} \phi_{u_i}(t, u_i((n + 1)H)) = u_i((n + 1)H)$ is still the same, as it depends on the communication step size $H$ and not on the micro step size $h_i$. The solver will reduce the micro step size until a minimum is reached, at which point it gives up and finally advantages the micro step [67].

Most of the times this gives acceptable results but has a huge performance toll: when the solver is repeatedly retrying a small micro-step size, it does not advance the simulated time. This means that a huge computational effort goes to waste until the solver finally gives up [66].

We defer the discussion of the correct ways to deal with discontinuities to co-simulation scenario where discontinuities are welcome, section 5. In continuous co-simulation scenarios, discontinuities should not occur.

A solution to avoid discontinuities in the input approximations is to use the extrapolated interpolation methods [55, 82]. These methods ensure at least that $\lim_{t \to ((n+1)H)} \phi_{u_i}(t, u_i(nH)) = \lim_{t \to ((n+1)H)} \phi_{u_i}(t, u_i((n + 1)H))$.

To give an example, we derive one possible linear extrapolated interpolation method for $\phi_{u_i}$ over the interval $t \in [nH, (n+1)H]$. Since $\phi_{u_i}$ is linear, then $\phi_{u_i}(t, u_i(nH), u_i((n - 1)H)) = b + a(t - nH)$, for some constants $a, b$. Let $\bar{u}_i(nH) = \phi_{u_i}(nH, u_i((n - 1)H), u_i((n - 2)H))$. To avoid discontinuities, we require that $\phi_{u_i}(nH, u_i(nH), u_i((n - 1)H)) = \bar{u}_i(nH)$. And we want that $\phi_{u_i}((n + 1)H, u_i(nH), u_i((n - 1)H)) = u_i(nH)$.

So putting these constraints together gives

$$
\phi_{u_i}(t, u_i(nH), u_i((n - 1)H)) = b + a(t - nH) \\
\bar{u}_i(nH) = \phi_{u_i}(nH, u_i((n - 1)H), u_i((n - 2)H)) \\
\phi_{u_i}(nH, u_i(nH), u_i((n - 1)H)) = \bar{u}_i(nH) \\
\phi_{u_i}((n + 1)H, u_i(nH), u_i((n - 1)H)) = u_i(nH)
$$

Solving this system for $\phi_{u_i}(t, u_i(nH), u_i((n - 1)H))$ gives:

$$
\phi_{u_i}(t, u_i(nH), u_i((n - 1)H)) = u_i((n - 1)H) + \frac{u_i(nH) - u_i((n - 1)H)}{H}(t - nH)
$$

### Real-time Constraints, Noise, and Delay

As introduced in section 2 the major challenge in real-time simulation is to ensure that a SU is fast-enough to satisfy the timing constraint $t = \alpha \tau$. In real-time co-simulation, this challenge gets aggravated due to the presence of multiple SUs, with different capabilities [241], and whose
internal workings are unknown. Furthermore, real-time co-simulation is often used when at least one of the SUs is a physical entity. This means that measurements may carry noise, and the extrapolation functions used in the other SUs have to be properly protected from that noise (e.g., using statistical techniques such as Kalman filtering [139][242]). Finally, the quality of the network is important, as the real-time SUs needs to receive their inputs in a timely manner. To mitigate this, the orchestration algorithm has to compensate for any delays in the receiving of data, and provide inputs to the real-time SU [240].

5 Hybrid Co-simulation Approach

Sections 3 and 4 described the essential characteristics and assumptions of simulation units (SUs) for each kind of co-simulation approach. When compared to a CT SU, whose state evolves continuously in time and whose output may have to obey to physical laws of continuity, a DE SU state can assume multiple values at the same time (transiency) and its output is discontinuous. For an orchestrator, a CT SU has some flexibility (safe for algebraic loops and complex coupling conditions) in deciding the parameters (e.g., step size or tolerance) of the co-simulation. In contrast, a DE SU has to get inputs and produce outputs at the precise time an event is supposed to occur, and there is no Lipschitz continuity conditions to help predict how a delay in the output of the DE SU can affect the overall co-simulation trace.

For example, in the SU of the mass-spring-damper system, eq. (22), with a constant extrapolation function, and running under the orchestrator in algorithm 3, the change in the input can only affect the output after at least $H$ units of time. For continuous time solvers in general, as can be seen for the explicit solver in eq. (12), a delayed response to the inputs is normal.

These differences are at the heart of many challenges in hybrid co-simulation scenarios.

5.1 Hybrid Co-simulation Scenarios

We do not give a formal definition of a hybrid co-simulation scenarios because that is related to finding an appropriate standard for hybrid co-simulation, which is a non trivial challenge (see section 5.2) [53]. Instead, we define it broadly as mixing the characteristics and assumptions of both kinds of SUs. These scenarios, together with an adequate orchestrator, can be used as mock-ups of hybrid systems [65][175][6][62]. A thermostat regulating the temperature in a room is a classical example [174]. The Continuous Time (CT) constituent system represents the temperature dynamics of the room, accounting for a source of heat (radiator). The Discrete Event (DE) part is a controller that turns on/off the radiator depending on the temperature.

The SU $S_1$ simulates the following dynamics:

$$\dot{x} = -\alpha (x - 30q) ; \quad x(0) = x_0$$

where $x$ is the output temperature in the room, $\alpha > 0$ denotes how fast the room can be heated (or cooled) down, and $q \in \{0,1\}$ is the control input that turns on/off the radiator. The SU $S_2$ simulates the statemachine shown in fig. 11, where one can think of the input event $\text{tooHot}$ as happening when $x(t) \geq 21$ and $\text{tooCold}$ when $x(t) \leq 19$. The output events $\text{off}$ and $\text{on}$ will assign the appropriate value to the input $q$ of $S_1$. Therefore, the temperature $x(t)$ is kept within a comfort region.
Clearly, the two SUs cannot just be coupled together via input to output assignments. Any orchestrator for this co-simulation scenario has to reconcile the different assumptions about the inputs and output of each SU.

- The CT SU expects a continuous input, whereas the output of the DE SU is an event signal.
- The output of the CT SU is a continuous signal, whereas the DE SUs expects an event signal as input.

The coupling of CT and DE black box SUs has been studied in the state of the art. In essence, two approaches are known, both based on adapting (or wrapping) the behavior of the SU:

**Hybrid DE** – adapt every CT SU as a DE SU, and use a DE based orchestrator;

**Hybrid CT** – wrap every DE SU to become a CT SU and use a CT based orchestrator.

According to the formalization that we have proposed for CT and DE SUs, the Hybrid DE approach, applied to the thermostat example may involve: adapting \( S_1 \) as a DE SU, \( S'_1 \), with a time advance that matches the size of the co-simulation step; and keeping track of the output of \( S_1 \) in order to produce an output event whenever it crosses the thresholds. Conversely, any output event from \( S_2 \) has to be converted into a continuous signal for the input \( q(t) \) of \( S_1 \). Other examples of Hybrid DE are described in [265][220][203][205][148][275][40][91][18][279][282][41][61][60][150][196][157].

The Hybrid CT, in our example, requires the adaptation of the DE \( S_2 \) as a CT SU that takes as input the temperature continuous signal, and internally reacts to an event caused by the crossing of the threshold. The output event of \( S_2 \) can be converted into a continuous signal \( q(t) \). Examples of the Hybrid CT include [106][219][162][76][240][90][253].

Regardless of the approach taken, the properties of the constituent systems have to be retained: the fact that an otherwise discontinuous signal becomes continuous as a result of a linear or higher order extrapolation may not respect the properties of the coupled system. Knowledge of the domain and the SUs is paramount to retain aforementioned properties.

A third option, compared to only using Hybrid CT or Hybrid DE, is to have different mechanisms of orchestrating the SUs depending on the semantic domain. For instance, in the actor modeling language Ptolemy II [218], an actor has many similarities to a SU. Instead of using either Hybrid CT or Hybrid DE, a so called Director block is used for a particular set of connected actors. In this context, the notion of superdense time is fundamental, as discussed in [53] and [74].

In the section below, different issues that arise in hybrid co-simulation will be described. These should be read in the light of hierarchical hybrid co-simulation scenarios, where compositionality is important.
5.2 Challenges

Semantic Adaptation and Model Composition

While a generic wrapper based on the underlying model of computation of the SU can be used, as done in [218][73], the realization of any of the approaches Hybrid DE or Hybrid CT depends on the concrete co-simulation scenario and the features of the SUs [46][192], as shown with the thermostat example. There is simply no best choice of wrappers for all scenarios. Even at the technical level, the manner in which the events or signals are sent to (or obtained from) the SU may need to be adapted [253]. To be concrete, the SU \( S_2 \) can assume that all events are communicated by encoding them in a single string signal, as opposed to having a different signal signal to denote different events. To account for this variability, the most common adaptations can be captured in a configuration language, as was done in [182][76], or in a specialization of a model of computation, as done in [157][190][210]. These approaches require that a person with the domain knowledge describes how the SUs can be adapted.

Our choice of wrapper for the Hybrid DE approach is meant to highlight another problem with the adaptations of SUs: the wrapper incorporates information that will ultimately have to be encoded in the software controller. As such, we argue that the need for sophisticated semantic adaptations should be smaller in later stages of the development of the components so that, for more refined models of the thermostat, the decision about when to turn off the radiator is not made by a wrapper of \( S_1 \).

Predictive Step Sizes and Event Location

In the Hybrid DE approach, the time advance has to be defined (recall eq. (1)). Setting it to whatever co-simulation step size \( H \) the orchestrator decides will work, but the adapted SU may produce many absent output events. Better adaptations have been proposed. In the thermostat example, \( S'_1 \) can propose a time advance that coincides with the moment that \( x(t) \) will leave the comfort region, thereby always being simulated at the relevant times.

Naturally, these approaches rely on information that may expose the IP of SUs. Others try to adaptively guess the right time advance by monitoring other conditions of interest, set over the own dynamics of the adapted SU, the most common approach being the quantization of the output space [283][41][147][148][204].

The capability to predict the time advance is also useful to enhance the performance/accuracy of CT based co-simulation, as shown in [52].

Locating the exact time at which a continuous signal crosses a threshold is a well known problem [14][285][42] and intimately related to guessing the right time advance for predicting the step size [61][105]. To address this, solutions typically require derivative information of the signal that causes the event, and/or the capability to perform rollbacks. In the thermostat example, a co-simulation that shows the output \( q \) of the controller changing from 0 to 1 at time \( t_e \) while the temperature of the room \( x \) actually crossed the comfort zone at \( t_e - k \), for \( k > 0 \), may not be accurate if \( k \) is too large. Note that \( k \) is a consequence of the decisions made in the orchestrator.

Discontinuity Identification

Until here, we have based our discussion in the knowledge of what kind of SUs comprise a co-simulation. In a general hierarchical co-simulation, a SU’s output may be an event signal coming from a wrapper of a CT SU, or vice-versa. In any case, at runtime, a signal is often represented
as a set of time-stamped points. Observing this sequence of points alone does not make it possible
to discern a steep change in a continuous signal, from a true discontinuity, that occurs in an event
signal $[166] [53] [58]$. Extra information is currently used: a) a formalization of time which
include the notion of absent signal, as proposed in $[246] [166] [53]$; or b) an extra signal can be used to
discern when a discontinuity occurs, as done in the FMI for Model Exchange $[58]$, even facilitating
the location of the exact time of the discontinuity; or c) symbolic information (e.g., Dirac impulses
$[79]$) that characterize a discontinuity can be included, as done in $[200] [114]$.

Discontinuity Handling

Once a discontinuity is located, how it is handled depends on the nature of the SUs and their
capabilities. If the SU is a mock-up of a continuous system then, traditionally, discontinuities in the
inputs should be handled by reinitializing the SU $[67]$. This step can incur a too high performance
cost, especially with multi-step numerical methods, and $[98]$ proposes an improvement for these
solvers. Furthermore, a discontinuity can cause other discontinuities, producing a cascade of re-
initializations. During this process, which may not finish, care must be taken to ensure that
physically meaningful properties such as energy distribution, are respected $[188]$.

Algebraic Loops, Legitimacy, and Zeno Behavior

Algebraic loops are non-causal dependencies between SUs that can be detected using feedthrough
information, as explained in section 4.3. In CT based co-simulation, the solution to algebraic loops
can be attained by a fixed point iteration technique, as covered in section 4.3. There is the possibility
that the solution to an algebraic loop will fail to converge. The result is that, if left unchecked, the
orchestrator would move an infinite number of input and output values between SUs, at the same
point in time.

In DE based co-simulation a related property is legitimacy $[254]$, which is roughly the undes-
irable version of the transiency property, explained in section 3. An illegitimate co-simulation
scenario will cause the co-simulation orchestrator to move an infinite number of events with the
same timestamp between SUs, never advancing time. Distance matrices, used to optimize parallel
optimistic approaches, as explained in $[103]$ and used in $[110]$, can be leveraged to detect statically
the presence of some classes of illegitimacy.

A similar behavior, but more difficult to detect is Zeno behavior. It occurs when there is
successively smaller intervals of time between two consecutive events, up to the point that the sum
of all these intervals is finite $[259]$. As shown in $[50]$, a simulator eventually fails to detect the
consecutive events. In particular, he advocates that the zeno behavior is a property of the model,
whereas the incorrectness is due to a simulation approximation error. However, while illegitimate
behaviors are not desired in pure DE co-simulation, Zenoness can be a desired feature in some hybrid
cosimulation scenarios (e.g., see $[47]$). We say in the theoretical sense because, for the purposes of
cosimulation, scenarios with Zenoness still have to be recognized and appropriate measures, such
as regularization $[137]$, have to be taken.

Stability under $X$

If a hybrid co-simulation represents a hybrid or switched system $[259]$, then it is possible that a
particular sequence of events causes the system to become unstable, even if all the individual con-
tinuous modes of operation are stable $[138]$ Example 1.1]. New analyses are required to identify
whether the CT SUs can yield unstable trajectories as a result of: 1. noisy inputs; 2. data quantization; 3. change of co-simulation orchestration [113]; 4. the events of wrapped DE SUs [115]; and, 5. delayed exchange of values.

**Theory of DE Approximated States**

In a pure DE based co-simulation, if round-off errors are neglected, the computed trajectories are essentially exact. To the best of our knowledge, only [284] addresses theoretically how the error in a discrete event system can be propagated. In CT based co-simulation however, error is an accepted and well studied and techniques exist to control it.

In Hybrid co-simulation, there is a need for analysis techniques that provide bounds on the error propagation in the DE SUs, when these are coupled to sources of error.

In addition, based on these analyzes, it should be possible for a DE SU to recognize that its error has exceeded a given tolerance, and measures should be taken to reduce that error. Having these techniques in place allows a hybrid co-simulation orchestrator to take appropriate measures (e.g., adapt the communication step size, etc.) the keep the error bounded in every SU.

**Standards for Hybrid Co-simulation**

While for CT co-simulation there is the Functional Mockup Interface (FMI) standard [38], and for DE co-simulation there is the High Level Architecture (HLA) [2] standard, as of the time of writing, both standards have limitations for hybrid co-simulation. References [240] [39] [106] [74] use/proposes techniques to perform CT simulation conforming to HLA. Recognizing that hybrid co-simulation is far from well studied, [53] proposes a set of idealized test cases that any hybrid co-SU, and underlying standard, should pass. In particular, it is important to have correct handling and representation of time, to achieve a sound approach for simultaneity.

Finally, even with a standardized interface, SUs have different capabilities: a fact that makes coding an optimal orchestration algorithm difficult. A possible approach to deal with this heterogeneity, proposed in [112], is to assume that all SUs implement the same set of features, code the orchestration algorithm for those features, and delegate to wrappers the responsibility of leveraging extra features (or mitigating the lack of). In the section below, these features are classified.

6 Classification and Applications

Having described the multiple facets of co-simulation, this section summarizes our classification and methodology, and applies it to a typical use case.

6.1 Methodology

To find an initial set of papers related to co-simulation, we used Google Scholar with the keywords “co-simulation”, “cosimulation”, “coupled simulation”, and collected the first 10 pages of papers. Every paper was then filtered by the abstract, read in detail, and its references collected. To guide our reading to the most influential papers, we gave higher priority to most cited (from the papers that we have collected).
We read approximately 30 papers to create the initial version of the taxonomy. Then, as we read new papers, we constantly revised the taxonomy and classified them.

After a while, new references did not cause revisions to the taxonomy, which prompted us to classify the collected papers in a more systematic fashion: all the papers that we collected from 2011 (inclusive) up to, and including, 2016 were classified. Two main reasons justify the last 5 years interval: limited time; and most of the papers refer to, and are based on, prior work. As a consequence, the classification would be very similar for many of the related references prior to 2011.

From the papers classified, those that report case studies where noted to create fig. 1. In total, 84 papers were read and classified.

6.2 Taxonomy

The taxonomy is represented as a feature model structured in three main categories, shown in fig. 12:

Non-Functional Requirements (NFRs): Groups concerns (e.g., performance, accuracy, and IP Protection) that the reference addresses.

Simulation unit (SU) Requirements (SRs): Features required/assumed from the SUs by the orchestrator described in the paper. Examples: Information exposed, causality, local/remote availability, or rollback support.

Framework Requirements (FRs): Features provided by the orchestrator. Examples: dynamic structure, adaptive communication step size, or strong coupling support.

Each main group is detailed in figs. 13 to 15. Abstract features denote concepts that can be easily detailed down but we chose not to, for the sake of brevity. Mandatory features are required for the activity of co-simulation, while optional are not.

6.3 Applications

To demonstrate how the taxonomy is used, we picked three examples from the state of the art: an industrial use case, a co-simulation framework, and a co-simulation standard.

6.3.1 An Industrial Application

The case study reported in applies co-simulation as part of the development of a controller for an exhaust gas recirculation water handling system. The purpose of this system is to clean and
recirculate exhaust gas to a ship engine intake manifold. The exhaust gas is cleaned by spraying water into it, and allowing the mixture to cool down and deposit in a receiving tank. Then, the (dirty) water is pumped to a water treatment center (externally developed) to be purified and reused.

The system is a representative example because: it includes parts that are developed by other departments (e.g., the ship engine) and external suppliers (e.g., the water treatment system); there are both continuous and discrete event dynamics (e.g., the control system is comprised of a state machine and a PI-Controller); and, quoting the authors, “to improve the control strategy of the WHS, a higher-fidelity model [of the systems interacting with the controller] should be used.”

In fact, thanks to the FMI Standard, its support by MATLAB/Simulink®, and to the INTO-CPS co-simulation framework, the authors were able to combine the behavior of higher fidelity models, with the behavior of the controller under development, simulated by an in-house C++ software application framework.

Through co-simulation, it was possible to reproduce and correct an issue that was previously encountered only during a (costly) Hardware-in-the-loop simulation with a physical engine test.
bench available at the MDT research center in Copenhagen.

This work is classified as highlighted in figs. 13 to 15.

6.3.2 A Framework

We next consider the work of [257], where an FMI based multi-rate orchestration algorithm is generated from a description of the co-simulation scenario. In the paper, the description language introduced can be reused in a tool-agnostic manner. The orchestration code generator analyzes the co-simulation scenario, and: a) identifies algebraic loops using I/O feedthrough information; b) separates the fast moving SUs from the slow moving ones, using the preferred step size information, and provides interpolation to the fast ones (multi-rate); and c) finds the largest communication step size that divides all step sizes suggested by SUs and uses it throughout the whole co-simulation. The algebraic loops are solved via successive substitution of inputs, storing and restoring the state of the SUs.

Based on these facts, [257] is classified as highlighted in figs. 13 to 15.

6.3.3 A Standard

The FMI standard for co-simulation, version 2.0 [97], can also be classified according to the assumptions it makes about the participating SUs. This is highlighted in fig. 14.

6.4 The State of the Art

The remaining state of the art is classified in Figs. 16-19. The raw data is available online\(^1\). The apparent lack of papers in the interval 2006-2009 is a consequence of our methodology (recall

\(^1\)http://msdl.cs.mcgill.ca/people/claudio/pub/Gomes2016bClassificationRawData/raw_data.zip
6.5 Discussion

Analyzing fig. 16, Accuracy is the most observed NFR, with 31 reports, followed by IP protection and Performance. The least observed NFRs are Fault tolerance, Hierarchy and Extensibility. Fault tolerance is especially important for long running co-simulations. One of the industrial partners of the INTO-CPS project has running co-simulations that takes a minimum of two weeks to complete. We argue that Extensibility (the ability to easily accommodate new features) should be given more importance: if an heterogeneous set of SUs participate in the same co-simulation scenario, the combination of capabilities provided (see fig. 14) can be huge. Thus, the orchestrator can either...
assume a common homogeneous set of capabilities, which is the most common approach, or can leverage the capabilities provided by each one. In any case, extensibility and hierarchy are crucial to address, and implement, new semantic adaptations.

As fig. 18 suggests, we could not find approaches that make use of the nominal values of state and output variables, even though these are capabilities supported in the FMI Standard, and are useful to detect invalid co-simulations. A-causal approaches are important for modularity, as explained in section 4.3, but these are scarce too.

As for the framework requirements, in fig. 19, the least observed features are dynamic structure co-simulation, interactive visualization, multi-rate, algebraic coupling, and partial/full strong coupling support. This can be explained by the fact that these features depend upon the capabilities of the SUs, which may not be mature.

Figs. 16 – 19 do not tell the full story because they isolate each feature. Feature interaction is a common phenomenon, and among many possible interactions, we highlight the accuracy concern, domain of the co-simulation, number of SUs supported, and IP protection. As can be seen from fig. 21, there is only one approach [157] that is both CT and DE based, up to any number of SUs. Accommodating the different CT and DE domains means that the approach assumes that the SUs can behave both as a CT and as a DE SU.

The concern with IP protection is evident in fig. 16 but the number of DE and CT based approaches that provide some support for it is small, as shown in fig. 20. Similarly, as fig. 22 suggests, accuracy does not show up a lot in the DE and CT approaches, for more than two SUs. Accuracy is particularly important in interactions between DE and CT SUs.

In general, from the observed classification, there is a lack of research into approaches that are both DE and CT based, and that leverage the extra features from the SUs.
7 Concluding Remarks

In this overview article, we show that there are many interesting challenges to be explored in co-simulation, which will play a key role in enabling the virtual development of complex heterogeneous systems in the decades to come. The early success can be attributed to a large number of reported applications. However, the large majority of these applications represent ad-hoc couplings between two simulators of two different domains (e.g., a network simulator with a power grid one, or a HVAC simulator with a building envelop one)

As systems become increasingly complex, the demand for co-simulation scenarios that are large, hierarchical, heterogeneous, accurate, IP protected, and so on, will increase.

This survey covers the main challenges in enabling co-simulation. To tackle such a broad topic,
we have covered two main domains—continuous-time- and discrete-event-based co-simulation—separately and then discussed the challenges that arise when the two domains are combined. A taxonomy is proposed and a classification of the works related to co-simulation in the last five years is carried out using that taxonomy.

From the challenges we highlight: semantic adaptation, modular coupling, stability and accuracy, and finding a standard for hybrid co-simulation. For early system analysis, the adaptations required to combine simulators from different formalisms, even conforming to the same standard, are very difficult to generalize to any co-simulation scenario.

One of the main conclusions of the classification is that there is lack of research into modular, stable, and accurate coupling of simulators in dynamic structure scenarios. This is where acausal approaches for co-simulation can play a key role. The use of bi-directional effort/flow ports can be a solution inspired by Bond-graphs [218], and there is some work already in this direction [225].

Finally, this document is an attempt to summarize, bridge, and enhance the future research in co-simulation, wherever it may lead us to.

Acknowledgment

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Table 1: Historical Perspective of Co-simulation.

<table>
<thead>
<tr>
<th>Time</th>
<th>Concept</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>&lt;80s</td>
<td>Single Formalism</td>
<td>The equations describing dynamic behavior are integrated together.</td>
</tr>
<tr>
<td>80s</td>
<td>Dynamic Iteration</td>
<td>Large circuits are decomposed into coupled constituent systems and dynamic iteration techniques are used [107][167][197][183][184][180].</td>
</tr>
<tr>
<td>90s</td>
<td>Multi-Formalism</td>
<td>Software and Hardware are developed and simulated concurrently [122][224][290][69] at multiple levels of abstraction [129][130][84]. Orchestration methods are explored in [91][100][63][254].</td>
</tr>
<tr>
<td>Late 90s and</td>
<td>Standard Interfaces</td>
<td>Recognized as key for co-simulation [291][155][213][127][243].</td>
</tr>
<tr>
<td>Early 2000s</td>
<td></td>
<td></td>
</tr>
<tr>
<td>2010s</td>
<td>IP Protection, X-in-the-loop, and Scale</td>
<td>Important to enhance industrial applicability of co-simulation [258][7][81][105][17][18][19].</td>
</tr>
</tbody>
</table>

A Historical Perspective of Co-simulation

This section provides an historical perspective that relates the major concepts in co-simulation to the time at which they are recognized in the studied state of the art, summarized in table 1.

A.1 One Formalism and Dynamic Iteration

Traditionally, the equations describing the dynamical behavior of large circuits were integrated together. These systems are sparsely coupled, reflecting the connections of the corresponding circuits, and many techniques were developed that take advantage of this structure [180].

The crucial idea that improved the simulation speed in up to two orders of magnitude is to decompose the large system into a set of coupled constituent systems and integrate them independently.

The decomposition of the circuit implies the definition of inputs and outputs for each of the resulting constituent systems. The coupling is then the assignment of outputs to inputs.

For a subsystem $S_i$, we call the subsystems, whose outputs are assigned to any of the inputs of $S_i$, for neighbor subsystems.

The essence of the dynamic iteration approach is to integrate each subsystem independently, for a period of time $T_n \to T_{n+1}$, using the extrapolated outputs of the neighbor subsystems as inputs [167][197][183][184].

Naturally, the fact that outputs are extrapolated introduces inaccuracies in the solution of the subsystem, so the integration can be repeated for the same period of time, with corrected outputs,
until some form of convergence criteria is met [134]. The extrapolated outputs of a subsystem $S_j$ can be corrected by collecting the outputs during the integration of $S_j$.

It is easy to see that this approach only requires communication between constituent systems at times $T_n$ and $T_{n+1}$ and that the integration of each subsystem can be done independently and in parallel [135], using any numerical method with any step size control policy. The signals exchanged are functions in the interval $[T_n, T_{n+1}]$.

The advantages of independent step size control policy become evident when one observes that many circuits have components that change at different rates. If the whole system were to be simulated, the simulation unit (SU) would have to use the smallest time step that ensures sufficient accuracy for the fastest changing component, which would be a huge waste of computational effort for the slow components. This is the similarity to multi-rate numerical methods [107].

To the best of our knowledge, dynamic iteration techniques and multi-rate numerical are the first to resemble co-simulation. The coordination software that implements these techniques expect any number of subsystems but assumes that the subsystems are all specific in the same formalism: differential equations.

### A.2 Two Formalisms: Digital and Analog Co-simulation

Co-simulation, in its modern definition, was applied to enable the virtual development of coupled software and hardware systems [122][224][290][69]. In this application domain, co-simulation decreases the need to build prototype board circuits to validate the composition of the software and the hardware part. It enables software and hardware to be developed and validated concurrently. To the best of our knowledge, this was one of the first uses of co-simulation in the modern sense. The co-simulation frameworks developed in this application domain typically assumed two SUs and two formalisms.

The hardware/software systems quickly became more complex and a new idea was introduced: use multiple models at different levels of abstraction of each subsystem. Simulations could be made arbitrarily faster in some intervals by solving the more abstract models, and arbitrarily accurate in other intervals, by solving the more detailed ones [84][227][179][146]. In the particular case of analog-digital co-simulation, each level of abstraction was solved by a different tool: a continuous time tool and a discrete event tool. The separation into continuous time and discrete event made the abstract synchronization problem and synchronization methods between SUs in these two domains were developed [91][100][63][254]. We could call these some of the first master algorithms.

### A.3 Multi-abstraction/Multi-Formalism Co-simulation

The heterogeneity aspect of co-simulation comes into play at this time: multiple formalisms can be used to describe the same subsystem at multiple levels of abstraction: state machines can describe a rough approximation of the modes, while differential equations can describe the detailed dynamics of the electronic circuit. Depending on the purpose of the co-simulation, a subsystem and its neighbors can be solved in detail, whereas subsystems that are “farther away” can be simulated with higher levels of abstraction [129][130]. For the domain of Hw/sw co-simulation, RTL and TLM classify the multiple abstraction levels of models [221][27] and switching between these multiple levels of abstraction have been studied in [143].

As the number and heterogeneity of simulation tools to be coupled increases, the need to provide a common interface to couple any number of tools is recognized in [291][155][213][127] and later in
In parallel with the previous advancements, co-simulation has also been in use for heterogeneous physical systems, such as automotive [163][118][153][231], railway [78][12] and HVAC [119][251], to name just a few. The common motivation is the fact that co-simulation enables specialized SUs to cooperatively simulate the system, with huge savings in time and cost, when compared to a monolithic modeling approach.

A.4 Black-box Co-simulation

Later, distributed and concurrent development processes, enabled by co-simulation, are studied and IP protection is identified as a desired characteristic [258][7] to enable suppliers and integrators to exchange co-SUs without having to disclose sensitive information and avoiding *vendor lock-in* contracts.

A.5 Real-time Co-simulation

Furthermore, co-simulation is used at every stage of the development process, from early system validation, to X-in-the-Loop co-simulation, bringing hard real-time constraints to the set of challenges [81].

A.6 Many SUs: Large Scale Co-simulation

More recently, with the acknowledgment that there is a need to be able to simulate even larger systems of systems, scale and distribution become inherent challenges in co-simulation [105][17][18][19].

B State of the Art in Co-simulation Frameworks

This section provides the detailed classification of each reference.

Ref 1. Trcka2010: Co-simulation for performance prediction of integrated building and HVAC systems – An analysis of solution characteristics using a two-body system [252].

*Summary.* They describe convergence and stability analysis for co-simulation.


*Summary.* They take a model which has been model checked, and map it to a more complex (stochastic) model, where stochastic analysis techniques can be applied. They describe a nice example, based on the railway domain. Heaters are used to prevent line switches from freezing.

It seems a great idea to apply their techniques to co-simulation itself. They use Contract Automata to perform the validation of the protocol. The formalism is quite elegant and allows for easy composition of many different simulators. They beauty is that we could use different automata for each kind of simulator participating in a co-simulation. Then we could synthesize a master algorithm that is considered optimal for that scenario.

Summary. They describe how to adapt traditional numerical methods to compute over approximations of the solutions.

They present an interesting problem called the oil-reservoir problem.


Summary. This paper gives many references to guaranteed integration techniques.

With these ideas, co-simulations could be computed with intervals of approximation, permitting reachability analyses.

The paper references other works which also are related to guaranteed integration.

Ref 5. Foster2017: Towards a UTP Semantics for Modelica [98].

Summary. They seem to formalize the event iteration system of the modelica language.

As far as I can see, it seems to be a formalization just for formalization’s sake. There is no convincing example of why it is useful to make such a formalization. Surely, the result is elegant.


Summary. They seem to propose a language to describe co-simulation toolchains?

I think the purpose of the language is to allow an easier integration of other tools into a given tool chain.

This work has been applied to two industrial use cases.


Summary. The paper is about the unique selling points of INTO-CPS.


Summary. This paper describes a control based scheme to compensate for extrapolation errors across signals in co-simulations. It is based on generalized energy concepts.

Summary. This work describes a co-simulation master in the context of the maritime industry.


Summary. This work describes a co-simulation between power system and network simulator.

Summary. This work describes a co-simulation approach that finds an appropriate co-simulation step size.


Summary. This work addresses co-simulation with real-time simulators.

Ref 13. Tomulik2011: Simulation of multibody systems with the use of coupling techniques: a case study [249].
Summary. This work discusses a co-simulation method for couplings with algebraic constraints. One of the results is that this kind of coupling should be done with many derivatives of the coupling variables.


Summary. This work describes the application of co-simulation to the power production domain.


Summary. This work describes a co-simulation approach.
Ref 16. Friedrich2011: Parallel Co-Simulation for Mechatronic Systems [102].

Summary. This work describes a co-simulation framework based on the Jacobi iteration scheme.

Ref 17. Gonzalez2011: On the effect of multirate co-simulation techniques in the efficiency and accuracy of multibody system dynamics [116].

Summary. This work deals with multi-rate co-simulation. Essentially, one of the simulators (the fast one) drives the simulation, while the slow one provides extrapolated inputs, to avoid excessive computation.
Ref 18. Nutaro2011: Designing power system simulators for the smart grid: Combining controls, communications, and electro-mechanical dynamics [203].

Summary. This work describes a tool that is formed by the coupling of a DEVS simulator with some other modules that wrap CT as DEVS simulators.


Summary. This work describes co-simulation approaches between two simulation tools. The main contribution is a semi-implicit method that applies a correction based on the jacobian of the subsystem’s coupling variables.
Ref 20. Pohlmann2012: Generating functional mockup units from software specifications [217].

Summary. This work describes an application of co-simulation to robotics.


Summary. This paper describes global error analysis for co-simulation, that takes into account sub-system solvers (instead of analytical solvers, as more commonly done).


Summary. This work present a coupling of the tools Crescendo and 20-sim.

Summary. This work introduces guidelines for the implementation of co-simulation between Matlab and SystemC. The case study is the JPEG Algorithm.


Summary. This work discusses co-simulation techniques for simulators coupled via algebraic constraints.

Summary. This work describes the MDPCosim framework.

- nfr: performance
- The decomposition of the system for co-simulation is done for performance reasons.
- nfr: parallelism
- nfr: accuracy
- sr: availability: local
- IPC communication is used.
- sr: causality: causal
- sr: info: derivatives: out
- sr: rollback: none
- sr: info: predict_step_sizes
- fr: results_visualization: post_mortem
- fr: alg_loop: explicit
- fr: coupling_model: io_assignments
- fr: domain: ct
- fr: communication_model: jacobi
- fr: sim_step_size: variable

The step size control approach is based on looking at the derivatives.

- fr: sim_rate: single
- fr: num_sim: three_more


Summary. Work describing another tool coupling.

- nfr: distribution
- sr: causality: causal
- sr: rel_time: analytic
- sr: rollback: none
- sr: availability: remote
- fr: results_visualization: post_mortem
- fr: alg_loop: explicit
- fr: coupling_model: io_assignments
- fr: domain: ct
- fr: num_sim: two
- fr: sim_rate: single
- fr: sim_step_size: fixed
- fr: communication_model: gauss_seidel

Ref 27. Al-Hammouri2012: A comprehensive co-simulation platform for cyber-physical systems [5].
Summary. The work describes the integration of two tools: Modelica, and NS-2.

Communication is done over named pipes.


Summary. This work describes the coupling of two tools: Matlab and NS-2. The coupling is done through HLA standard. The preliminary version of the tool is described in [222].


Summary. This work describes the coupling of two tools: Matlab and NS-2. The coupling is done through HLA standard.

**Summary.** This work describes a co-simulation between two tools in the power grid domain with matlab running the co-simulation.


**Summary.** This work describes the coupling between two tools: Overture and 20-sim.
Ref 32. Fitzgerald2013: A formal approach to collaborative modelling and co-simulation for embedded systems [96].

Summary. This work describes the coupling between two tools: Overture and 20-sim; already described in [93].


Summary. This work describes the integration of three simulators (ARGoS, NS-2 and NS-3) that can be used in co-simulation scenarios with two simulators.
Ref 34. Broman2013: Determinate Composition of FMUs for Co-simulation [52].

Summary. This work describes a master algorithm that ensures a determinate execution.

Ref 35. Benedikt2013: Guidelines for the Application of a Coupling Method for Non-iterative Co-simulation [33].

Summary. This work describes a co-simulation approach where energy information about the signals is used, and those errors are compensated in a corrector step.

Summary. The work describes a method for finding appropriate communication step sizes in cosimulations between LTI systems. Essentially, it provides rules of thumb to choose a communication step size based on the maximum instantaneous frequency of components.

Ref 37. Fuller2013: Communication simulations for power system applications [104].

Summary. This work describes a co-simulation between two co-simulation tools (ns-3 and GridLAB-D) for smart grid development.
Ref 38. Bombino2013: A model-driven co-simulation environment for heterogeneous systems [39].

Summary. This work describes the coupling between two simulation tools.


Summary. The approach described in this reference allows to arrange and process co-SUs, Modelica models and TinyOS applications. SysML is used to configure the co-simulation master. The coordination of simulators is done through the FMI standard.

Ref 40. Hafner2013: An Investigation on Loose Coupling Co-Simulation with the BCVTB [124].
Summary. This work discusses the consistency and stability of the Jacobi and Gauss-Seidel co-
simulation methods. Later, it presents a case study in HVAC systems.

Ref 41. Zhao2014: Co-simulation research and application for Active Distribution Network based on Ptolemy II and Simulink [289].

Summary. This work describes the co-simulation between Ptolemy II and Simulink.

Ref 42. Li2011c: VPNET: A co-simulation framework for analyzing communication channel effects on power systems [169].

Summary. This work describes the coupling of two simulation tools (VTB and OPNET) to achieve co-simulation.
The coordination is a sample discrete time system.

Ref 43. Awais2013b: Distributed hybrid simulation using the HLA and the Functional Mock-up Interface [19].

Summary. The main difference between this work and [17] is that this proposes a variable step size wrapper around CT components. The approach taken to do this is quantization.

Ref 44. Awais2013a: Using the HLA for Distributed Continuous Simulations [17].

Summary. This work addresses the need to adapt CT simulators as DE simulators, in order to be used in a hybrid co-simulation scenario that is fundamentally DE oriented.
Ref 45. Kuhr2013: FERAL - Framework for simulator coupling on requirements and architecture level [157].

Summary. They describe a framework that borrows many concepts from Ptolemy, but that is fundamentally event based co-simulation. It allows for the specialization of basic directors for the semantic adaptation of SUs.

Ref 46. Viel2014: Implementing stabilized co-simulation of strongly coupled systems using the Functional Mock-up Interface 2.0. [271].

Summary. This work describes the implementation of the method described in [10] in the context of the FMI standard.
Ref 47. Sicklinger2014: Interface Jacobian-based Co-Simulation [239].

*Summary.* Describes a co-simulation method that makes use of the Jacobian information for fixed point computations.


*Summary.* The work describes a co-simulation that integrates SystemC and CarSim.
Ref 49. Kounev2015: A microgrid co-simulation framework [50].

*Summary.* Describes the coupling of two simulators written in MATLAB and OMNeT++.

Ref 50. Bogomolov2015: Co-Simulation of Hybrid Systems with SpaceEx and Uppaal [39].

*Summary.* The orchestration algorithm is the one described in [52]. The work exploits the standard by allowing zero step transitions.
They abuse the FMI standard to be able to support state transitions.


Summary. Not many details are provided about the co-simulation orchestration. However, due to the fact that it is real-time, we can infer certain features.

Ref 52. Dols2016: Coupling the multizone airflow and contaminant transport software CONTAM with EnergyPlus using co-simulation [80].

Summary. The work described the coupling of the CONTAM and EnergyPlus tools to achieve HVAC simulation. The coupling is done through FMI. The coupling is done through the compiled binaries. The case study highlights the problems with an explicit method for co-simulation, even if the Gauss-seidel. Instabilities occur.

**Summary.** According to [28], this work explores variable step solvers.

- nfr: parallelism
- nfr: performance
- sr: info: causality: feedthrough
- sr: info: full_model
- sr: rel_time: fixed
- sr: rel_time: fixed_real_scaled_time_simulation
- sr: rollback: none
- sr: availability: local
- fr: standard: fmi
- fr: coupling_model: io_assignments
- fr: num_sim: three_more
- fr: domain: ct
- fr: sim_rate: single
- fr: sim_step_size: fixed
- fr: alg_loop: explicit
- fr: results_visualization: post_mortem
- fr: communication_model: jacobi
- fr: communication_model: gauss_seidel


**Summary.** This paper focuses on the parallelization of co-simulation. The approach is to start with a single model and partition it into multiple models, which are then executed in separate FMUs in parallel. The partitioning is important for accuracy reasons (e.g., break the algebraic loops at less sensitive variables).

- nfr: parallelism
- nfr: accuracy
- nfr: performance
- nfr: scalability
- sr: info: full_model
- sr: info: wcet
- sr: info: causality: feedthrough
- sr: causality: causal
- sr: rel_time: analytic
- sr: rollback: none
- sr: availability: local
- fr: communication_model: gauss_seidel
- fr: sim_step_size: fixed
- fr: sim_rate: single
- fr: standard: fmi
- fr: domain: ct
- fr: num_sim: three_more
It breaks the loops by establishing an order and delaying one of the variables in the loop.


Summary. The paper addresses the problem of performance in FMI co-simulation. The solution proposed is to go parallel. The parallelization approach is the same as the one presented in [145]. Since FMI does not enforce thread safety across multiple instances of the same FMU, the work presented ensures that these do not execute concurrently by using mutexes or changing the scheduling policy.

Ref 56. Yamaura2016: ADAS Virtual Prototyping using Modelica and Unity Co-simulation via OpenMETA [278].

Summary. The co-simulation framework includes 4 tools. The communication between the tools is realized using OpenMeta. The work uses Unity for the modelling and simulation of the environment, allowing for live interaction. Communication is over UDP but there is no report on extra caution due to network delays and failures.
Ref 57. Camus2015: Combining DEVS with multi-agent concepts to design and simulate mult models of complex systems (WIP) [60].

Summary. This work is the preliminary description of [61].

Ref 58. Camus2016: Hybrid Co-simulation of FMUs using DEV & DESS in MECSYCO [61].

Summary. It proposes to use a FMU wrapper around DEV and DESS models, meaning that the co-simulation proceeds using a DE approach. It handles black box FMUs and the algorithm used to drive the co-simulation is the conservative parallel DEVS simulator. It requires that the FMU is able to perform rollback (through the use of state set and get).
Ref 59. Pedersen2016: FMI for Co-Simulation of Embedded Control Software [210].

Summary. The paper describes the adaptation of an embedded system to comply with FMI and thus interface with other FMUs. To validate the implementation, they run a co-simulation.


Summary. The paper proposes a co-simulation framework that takes into account network delays and compensates for that. It proposes to use cubic spline extrapolation to compensate for the delay but recognizes that if there are faults in the line (resulting in voltage drops), the derivatives used in the extrapolation assume gigantic proportions, thus wreaking havoc in the simulation. To address that, the framework employs an algorithm to detect discontinuities. The detection is simple: they check the derivative of the signal to see whether it exceeds a pre-determined empirically threshold. Basically, it looks for and Dirac delta. Figure 7 shows the effect of not handling a discontinuity.
Ref 61. Xie2016: Continuous-Mass-Model-Based Mechanical and Electrical Co-Simulation of SSR and Its Application to a Practical Shaft Failure Event [277].

**Summary.** Between two simulators. As it is explained in the paper, prior to co-simulation, the most common approach would be to run two simulations: one complete for one sub-system, and then another for the second sub-system, using the first as inputs. This is an open loop approach, whose results can be misleading due to ignoring the feedback loops. Each simulator advances in parallel and their communication is made with a barrier.

```plaintext
nfr: parallelism
fr: communication_model: jacob
fr: num_sim: two
fr: domain: ct
fr: sim_step_size: fixed
fr: results_visualization: post_mortem
fr: sim_rate: single
sr: availability: local
sr: causality: causal
sr: rel_time: analytic
sr: rollback: none
fr: alg_loop: explicit
fr: coupling_model: io_assignments
```


**Summary.** It describes an application of co-simulation in the distribution of energy in smart grids, supported by a real-time co-simulation framework. The simulators involved are the RTDS, which simulates the distribution network model, and the VVO Engine, coded in MATLAB.

```plaintext
sr: rel_time: fixed_real_scaled_time_simulation
fr: num_sim: two
sr: causality: causal
sr: rollback: none
sr: availability: local
fr: coupling_model: io_assignments
fr: sim_rate: single
fr: sim_step_size: fixed
fr: alg_loop: explicit
fr: communication_model: jacob
```

Ref 63. Schierz2012a: Co-simulation with communication step size control [229].
**Summary.** Describes a master algorithm. Does not allow for interpolation of inputs. Needs rollback. It touches upon accuracy, as it suggests an adaptive step size control mechanism. It does not address algebraic loops. It assumes that there is no feedthrough information.

---

Ref 64. Fourmigue2009: Co-simulation based platform for wireless protocols design explorations [69].

**Summary.** Application of co-simulation to wireless network development. One of the simulators is the actual Linux operating system, and the other is represents a wireless network protocol simulator.

---

Ref 65. Liu2001: Calculation of Wing Flutter by a Coupled Fluid-Structure Method [172].

**Summary.** A fully implicit method, dealing with parallelism.
Ref 66. Carstens2003: Coupled simulation of flow-structure interaction in turbomachinery [64].

Summary. Relates to the application of a co-simulation algorithm to the simulation of the deformation in the blades of a transonic compressor rotor under airflow. One of the simulators calculates deformation of the blades, while the other calculates the flow dynamics around the blades.

The communication of orchestration algorithm in use is shifted by half a step.

They highlight the need for it, because the computation of a rotor is just too expensive.

It seems that they perform the computation in separate computers.

Although it is a gauss seidel shifted in time.


Summary. Proposes to address the challenges in real-time co-simulation by using a model based coupling approach. The master has to keep track of two values for each packet of data: receiving time delay $t_r$ – the time it takes for a packet to reach the master from the simulator –, and sending time delay $t_s$ – the time it takes for a packet to leave the master and reach the simulator. When a sample is delayed, the master acts as a replacement for it. Basically, it is a dead reckoning model.
Ref 68. Benedikt2016: Automated configuration for non-iterative co-simulation [34].

Summary. Describes how a co-simulation master can configure some parameters throughout the co-simulation execution. This is the idea behind adaptive master algorithms.

Ref 69. Busch2011: An explicit approach for controlling the macro-step size of co-simulation methods [57].

Summary. Presents an approach to estimate the local truncation error caused by the extrapolations of inputs in a co-simulation. The sub-systems are assumed to make no error. It does not require rollback or re-initialization.

Categories:
- **Accuracy**: Because they study the global error and control the local error.
- **Performance**: They control the step size, which increases performance. And they study how to get an optimal step size.
In theory, they seem to support any communication model. In the paper they studied assuming the Jacobi.

**Ref 70. Quesnel2005**: DEVS coupling of spatial and ordinary differential equations: VLE framework [220].

*Summary.* Proposes a way to wrap a continuous time ODE simulator as a DEVS model. It requires that the state variables, and derivatives are available.

**Categories:**
- Hierarchy
- Open source
- Statevars
- Predict_step_sizes
- Causality: causal
- Domain: de
- Num_sim: three more
- Rel_time: analytic
- Sim_rate: multi
- Sim_step_size: variable
- Alg_loop: explicit
- Availability: local
- Results_visualization: post_mortem
- Coupling_model: io_assignments
- Communication_model: gauss_seidel

Any discrete event framework is by definition multi-rate.

Any discrete event framework is by definition in this category.

A discrete event framework is in this category as there is no extrapolation of inputs. Also, Gauss seidel does not violate the causality of inputs and outputs, because it sorts according to these dependencies. Events are processed to retain their causality.

**Ref 71. Arnold2014a**: Error analysis for co-simulation with force-displacement coupling [14].

*Summary.* Describes an FMI based master called SNiMoWrapper.

**Categories:**
- Accuracy
- Protection
- Causality: feedthrough
- Simulation: analytic
- Step_size: variable
- Alg_loop: explicit
- Availability: local
- Results_visualization: post_mortem
- Coupling_model: io_assignments
- Communication_model: gauss_seidel
- Domain: de
- Num_sim: three more
- Rel_time: analytic
- Sim_rate: multi
- Sim_step_size: variable

Summary. Studies the error control method known as Richard’s extrapolation.

Categories:
  - nfr: accuracy
    - Because they study the global error and control the local error.


Summary. Studies the convergence of the Gauss-Seidel dynamic iteration method and proposes a way to ensure it. The way to do it though, requires information from the model.

Categories:
Because they study the global error.

Summary. Proposes a predictor corrector master that evaluates the macro step twice and uses a perturbation on the inputs to get an estimate of the required partial derivatives. This approach is then generalized to multiple kinds of joints in the mechanical domain. A double pendulum, double mass-spring-damper and a slider crank mechanism are used as numerical examples.

Categories:

Ref 74. Schweizer2014: Semi-implicit co-simulation approach for solver coupling [234].

Ref 75. Schweizer2015d: Stabilized implicit co-simulation methods: solver coupling based on constitutive laws [237].

Summary. Proposes a master for co-simulation that requires the identification of power bonds between sub-systems. It assumes that the scenario is energy conserving and thus calculates the energy residual as an error to be minimized. The step size is adapted via a PI-Controller. When the step size is reduced, it is only on the next co-simulation step, so the method is explicit.

Ref 77. Busch2016: Continuous approximation techniques for co-simulation methods: Analysis of numerical stability and local error [55].

Summary. Analyses the stability and local error of multiple co-simulation approaches with multiple extrapolation approaches for the inputs. It considers Gauss-Seidel and Jacobi. It also talks about a method called the extrapolated interpolation method, which ensures no discontinuities at the inputs of the subsystems.
The method is explicit.

Ref 78. Arnold2010: Stability of Sequential Modular Time Integration Methods for Coupled Multibody System Models [10].

Summary. Studies stability of a gauss Seidel co-simulation method proposed in previous work: [11]. Based on that analysis, it proposes an implicit stabilization technique that uses Gauss-Seidel iteration. The resulting method is implicit but the equations that are being solved are linear.

Categories:


Summary. Describes a technique to deal with algebraically coupled sub-systems using a control theoretic approach. The highlights of this method are: it supports scenarios of arbitrary index; the boundary condition coordinator is seen as a co-SU (this is an elegant approach) and the method is explicit. The beauty of making the BCC as a co-SU, is that it can, just like any other sub-system be
run at a different rate and in the paper they show that by running it at a higher rate, the stability of the co-simulation increases.


Summary. Describes a technique to solve causal conflicts using a Boundary Condition Coordinator (BCC). Causal conflicts arise naturally from the coupling of different sub-systems and they are a relevant challenge that needs to be overcome in order to perform correct co-simulation. While in [118], the BCC requires the knowledge of the state variables of the simulations, in [119], some modifications are made to ensure that this information is not required.

Ref 81. Schweizer2016: Co-simulation method for solver coupling with algebraic constraints incorporating relaxation techniques [238].

Summary. A master algorithm capable of dealing with algebraic constraints is described. It requires the derivatives of the coupled variables to be available. It executes each communication step twice, being a semi-implicit method. It uses a predict step and a corrector step. The final corrected coupling variables are obtained by polynomial extrapolation and relaxation (to avoid instabilities).
Ref 82. Schweizer2015: Predictor/corrector co-simulation approaches for solver coupling with algebraic constraints [233].

Summary. Proposes a predictor corrector master that evaluates the macro step twice and uses a perturbation on the inputs to get an estimate of the required partial derivatives.

Ref 83. Schweizer2015a: Stabilized index-2 co-simulation approach for solver coupling with algebraic constraints [235].

Summary. A master algorithm capable of dealing with algebraic constraints is described. It requires the derivatives of the coupled variables to be available. It executes each communication step twice, being a semi-implicit method. It uses a predict step and a corrector step. The predictor step allows the method to estimate the sensitivity of the state variables with respect to the applied forces/torques.
Summary. This is a Phd Thesis. A linear extrapolation based master is proposed that is convergent and does not require fixed point iterations. Then, a modification to multi-step methods is proposed to increase their performance when executing in a co-simulation environment. This modification avoids the need to restart when dealing with discontinuities.

Categories:
- IP protection
- Platform independence
- Open source
- Rollback: None
- Causality: Causal
- Domain: CT
- Number of simulations: Three more
- Relative time: Analytic
- Time step size: Fixed
- Availability: Local
- Results visualization: Post mortem
- Communication model: Jacobi
- Coupling model: Algebraic constraints
- Standard: FMI

Ref 84. Andersson 2016: Methods and Tools for Co-Simulation of Dynamic Systems with the Functional Mock-up Interface [5].

Ref 85. Krammer 2015: Model-Based Configuration of Automotive Co-Simulation Scenarios [15].

Summary. The language is further developed in [15] with the addition of three novel diagrams to represent different aspects of the co-simulation configuration:

- Architectural – coupling of executable units;
- Tools – assignment of tools to models;
- Connections – connections (it was not clear what does this diagram do);
In addition, they define a couple of well formedness properties that can be checked more easily with the model-based approach. They give a brief summary of the tool ICOS.

Categories:

- nfr: config_reusability
- nfr: parallelism
- nfr: hierarchy
- nfr: extensibility
- fr: domain:ct
- fr: num:three_more
- sr: rel_time:analytic
- sr: availability :local
- fr: coupling_model:io_assignments
- fr: results_visualization :post_mortem
- sr: info:full_model

Ref 86. Galtier2015: FMI-Based Distributed Multi-Simulation with DACCOSIM [105].

Summary. DACCOSIM is able to perform Distributed simulations and multi-core simulations. The term “computation node” is used for a collection of FMU wrappers (which include an FMU) and a local master / global master. The FMU wrappers, and thereby not the masters, are responsible for passing outputs to connected inputs. This is to avoid bottlenecks. A component node contains a master and some FMUs, which are wrapped in so-called “FMU-wrappers”. The masters take responsibility of coordinated step sizes in case an FMU needs to roll back.

Because of the possibility of splitting the simulation over a cluster / multi-core and the focus on performance in the article. Additionally because of their use of variable step size

It is possible to create multiple co-simulation configuration files in the simulation configuration GUI. These can be stored and therefore reused.

Some level of IP protection because of FMI.

Because of the possibility of splitting the simulation over a cluster

Because a co-simulation can be executed on a cluster of computers

DACCOSIM is weak hierarchical because it has the notion of local and global masters.

The framework is considered to be scalable because of the multi-core and distributed architecture.

There are two versions of the DACCOSIM library. A cross-platform version relying on JAVA and a Windows version using C++ and QTronic SDK.
The article provides an example where the result of the co-simulation using DACCOSIM is compared to the simulation using Dymola and the results are very close to each other. Accuracy is ensured by each FMU examining its outputs and estimating how far they are from the exact value.

The framework is distributed under an open source license from January 2016.

The framework can perform a single rollback using the state variable serialization.

Because the framework is based on FMI for co-simulation it is considered to be causal.

The framework supports multiple formalisms because it is based on FMI for co-simulation.

The frameworks is capable of supporting many FMUs and thereby many SUs. DACCOSIM offers its own algorithm depending on global/local masters.

There is no mentioning of any other time models than this in the article.

The simulation rate is the same for all FMUs.

The framework uses Coordinated Variable Step

The framework uses Euler’s method and Richardson’s method. Whether this is default, parameterizable or fully customizable is unknown based on this article.

See Co-initialization bullet 2 in the article.

It is based on the FMI standard.

Ref 87. Fey1997: Parallel synchronization of continuous time discrete event simulators [91].

Summary. Presents two synchronization approaches, detailed in three different synchronization protocols, to coordinate simulation scenarios that include one discrete event simulator and one continuous time simulator. The discrete event simulator can implement any parallel simulation approach that we know, such as Time-Warp. This means that, even internally, the DE simulator can be forced to rollback due to straggler messages. The focus is on parallel approaches.
Ref 88. Acker2015: Generation of an Optimised Master Algorithm for FMI Co-simulation [257].

Summary. Essentially, this paper shows how a compiled approach increases the performance of the co-simulation. It also shows that, because there are so many decisions to be made when designing the master, a compiled approach allows for a more elegant, and specifically tailored master, to be generated.


Summary. The paper describes and compares two approaches to performing co-simulation of heterogeneous systems, namely the Functional Digital Mock-up (FDMU) and the Functional Mock-up Interface (FMI). Besides describing these approaches it also introduces the “FDMU framework”, a framework that implements the Functional Digital Mock-up approach. Furthermore, proposals are presented for combining FDMU and FMI approaches.
The FDMU approach is a tool-independent and web service-based framework built on the Web Service standards. It is capable of coupling different simulation tools and provide visualization based on CAD models.

FDMU consists of three main concepts: functional building blocks (FBB), wrappers, FDMU master, and FDMU Console. The functional building block can wrap geometric information (CAD Models), behavioral models, and a simulator tool. It is the responsibility of the wrappers to establish a connection between the different simulation tools and the FDMU Master Simulator. Finally, the FDMU master ensures correct communication between the simulators. The FDMU Console is the user's front-end.

**nfr: performance**
Communication overhead of a web service-based approach.

**nfr: ip protection**
Because of the web service-based approach IP protection should be possible.

**nfr: parallelism**
Because of the web service-based approach it is parallel by nature. It uses thread-safe queues and deadlock-free transmission of data.

**nfr: distribution**
Because of the web service-based approach it is easy distributable.

**nfr: scalability**
The distributed systems paradigm ensures scalability.

**nfr: platform independence**
web service-based approach.

**nfr: extensibility**
A new wrapper can be implemented.

**sr: causality : causal**
Every input of an FBB must have an appropriate output belonging to another FBB.

**fr: domain: ct**

**fr: coupling model: io assignments**

**fr: num sim: three more**

**fr: sim rate: multi**

**sr: availability : remote**

**fr: results visualization : live**

The framework provides an interactive 3D visualization based on CAD.

**fr: standard: fdmu**

---

Ref 90. Karner2010a: Heterogeneous co-simulation platform for the efficient analysis of FlexRay-based automotive distributed embedded systems [12].

**Summary.** Motivation: FlexRay is a wired network for automotive high-speed control applications and no solutions exist that simulates all parts of the network.

What: a co-simulation platform called TEODACS FlexRayExpert.Sim. The simulation approach used in the platform covers mechanics and all parts of the network from physical layer to application layer, which is not done by other solutions. The framework CISC SyAD is used to perform the microelectronics co-simulation, CarMaker/AVL InMotion for the mechanics, and they are bridged...
by TEODACS FlexRayEprt.Sim. The platform uses a very interesting approach to faster co-
simulations, namely the use of model switching, where a less detailed model replaces a more detailed
model in parts of the simulation.

The paper provides an overview of existing approaches such as transaction based modeling,
HDLs such as SystemC and Verilog, and cable harness and topology modeling along with why
these contain shortcomings to this domain. Furthermore, the paper provides some details of the
implementation of the models used in the co-simulation and showcases how the platform can analyse
a system with specific examples.

---


*Summary.* The paper describes MOKA, which is a framework for performing co-simulations and
creating FMUs using FMI 2.0 for co-simulation. The framework turns the creation of FMUs into
an object-oriented process by using C++. An FMU is created by inheriting one of the classes and
implementing virtual functions thereby avoiding writing boilerplate code. The implementation of
FMUs is realised by the concepts of FMUBlock, which is to be inherited, FMUPort, and FMU-
StateVariables. FMUBlock is to be extended by a concrete FMU slave and implements common
computation phase functions for slaves. It contains FMUPort for data exchange and FMUStateVa-
riables for state tracking during the simulation. The FMUPort classes provides the data exchange
interface of a slave. It abstracts the value references by automatically assigning a value reference
to the variable. The BaseStateVariable class also functions as base that is to be extended. It provi-
des virtual functions for state variable services. The StateVariable inherits from BaseStateVariable
and represents state variables for the slave. The framework also provides a template for the FMU
Master so the master code changes minimally for different scenarios.

The article exemplifies an application of the MOKA framework where two FMUs are used: The
bouncing ball and integer counter example from the QTronic SDK, where the bouncing ball has
been re-developed with MOKA.

In future work it is stated that development of a DSL in a current study, so that different
scenarios can be executed without altering the master code.
Ref 92. Wetter2010: Co-simulation of building energy and control systems with the Building Controls Virtual Test Bed [274].

Summary. Describes a co-simulation framework called Building Controls Virtual Test Bed (BCVTB) that can be used for real-time simulation and co-simulation. It is a modular extensible open-source platform to interface different simulation programs with each other. The intention is to give users the option to use the best tools suited to model various aspects of building energy and control systems, or use programs where they have expertise. The middleware to couple any number of simulation programs also provides libraries such that it can be extended. Furthermore, the paper describes how they gathered capabilities the framework should support. The framework is based on Ptolemy II, which is extended by some java packages. The simulator package adds functionality that allows an actor to perform system calls to start any executable on Windows, OSX or Linux. It simply starts a simulation program, sends input tokens to the simulation program, receives new values and sends them to its output port. Algorithms are also provided on how simulators are coupled. These are also exemplified with specific simulators. It is also described how to connect client programs. The article describes how the interfaces are created for simulink, matlab, modelica, and system cals. Furthermore, a specific example is presented.

In the paper, their explanation is focused on the CT domain.
**Ref 93. Neema2014**: Model-based integration platform for FMI co-simulation and heterogeneous simulations of cyber-physical systems [196].

**Summary.** The article concerns integrating FMI as an HLA federate and extending the Command and Control Wind Tunnel (C2WT) metamodel to include FMI-specifics. This enables the C2WT tool to use FMUs as part of a simulation. The C2WT tool describes as a multi-model integration platform that allows users to model and synthesize complex, heterogeneous, command and control simulations. The tool therefore has support for multiple simulation engines and an introduction to the tool is given in the paper. Furthermore, a case study on Vehicle Thermal Management using FMUs are presented and the tool compared to a simulation in a different environment. The work is sponsored by the US DoD.

Ref 94. Larsen2016c: Integrated Tool Chain for Model-Based Design of Cyber-Physical Systems [159].

**Summary.** This article presents an overview of the INTO-CPS project and thereby a Co-Simulation tool. The project concerns production of a well-founded tool chain for model-based design of CPSs, and therefore consists of a semantic foundation and several baseline tools such as Modelio, Overture, 20-sim, OpenModelica and RT-Tester. Furthermore, an application called the INTO-CPS application is the entry point for configuring co-simulations and uses the co-simulation orchestration engine (COE) to perform the actual simulations. This COE is based on the FMI standard. The entire tool chain and the semantic foundation are presented in this paper. This is related to [159].
Summary. This is an EU deliverable related to the INTO-CPS project. It contains the technical documentation of the INTO-CPS platform at the end of 2015 (the first year of the project). Part of this project is the Co-simulation Orchestration Engine (COE). This is related to [159].

Co-Simulation Scenario Categorization

This section describes each category and lists the references that belong to that category, classified in the previous section.

C.1 Non-Functional Requirements

C.1.1 Fault Tolerance

A co-simulation platform is fault tolerant if, for example, when one SU fails, other can take its place. This is particularly important for long running simulations. To be fault tolerant, certain features need to available: periodically store the state of SUs; record all inputs to each SU. If a
SU fails and the state is periodically stored, then the simulation can be paused while the state is restored in a new instance of the SU. The history of input values passed to each SU can be used to bring the SU to the current state.

References in this category:
- Fuller2013

C.1.2 Configuration Reusability

This category refers to the fact that frameworks can provide a way to configure co-simulation scenarios that can be reused. This means that the configuration is considered external to the execution of the co-simulation. External in this context means that the configuration can reused without altering the binaries for the co-simulation application.

If a tool/frame does not provide a way to reuse configurations for co-simulation, then it is a time-consuming, error-prone and non-trivial process to set up co-simulations [151].

References in this category:
- Wang2013
- Krammer2015
- Galtier2015
- Acker2015
- Aslan2015
- Wetter2010
- Neema2013

C.1.3 Performance

Performance is a relative measure: a co-simulation platform is performant when it is able to simulate a great deal in a short amount of time while needing little resources. This can be achieved by using variable step integration methods and signal extrapolation techniques. Parallelism also plays a role but mostly on the time aspect of performance.

References in this category:
- Hoepfer2011
- Faure2011
- Sun2011
- Friedrich2011
- Gonzalez2011
- Gunther2012
- Eyisi2017
- Riley2011
- Benedikt2013
- Sicklinger2014
- Kounev2015
- Bian2015
- BenKhaled2012
- BenKhaled2014
- Said2016
- Schierz2012a
- Liu2001
C.1.4 IP Protection

IP Protection deals with not requiring the models participating in the co-simulation to provide detailed structure, variables, etc. There are multiple levels of protection ranging from fully protected to not protected at all. A good IP protection enables component suppliers to provide the system integrators with detailed simulations of their components avoiding expensive lock-in contracts.

There are multiple techniques can be employed to ensure some degree of protection. For instance, making the models (and corresponding SUs) available as a web service is a possible solution. Another example is any framework that implements the FMI Standard [37][38], which allows models and SUs to be exported as a single functional unit, in binary format, that can be imported into a co-simulation.

References in this category:

- Hoepfer2011
- Sun2011
- Bastian2011a
- Friedrich2014
- Broman2013
- Fuller2013
- Wang2013
- Kuhr2013
- Viel2014
- Dols2016
- Said2016
- Camus2015
- Camus2016
- Benedikt2016
- Busch2011
- Arnold2014a
- Arnold2014
- Sadjina2016
- Gu2014
- Andersson2016
- Galtier2015
- Fey1997
- Acker2015
- Enge—Rosenblatt2011
- Karner2010a
C.1.5 Parallelism

A co-simulation framework is parallel when it makes use of multiple processes/threads to perform the co-simulation. This is typically in the same computer or same local network.

Techniques such as signal extrapolation help improve the speed-up gained from parallelism. Furthermore waveform relaxation techniques and the Jacobi iterations promote parallelism [180].

References in this category:
- Faure2011
- Bastian2011a
- Friedrich2011
- Gunder2012
- Awais2013b
- Awais2013a
- BenKhaled2012
- BenKhaled2014
- Said2016
- Yamaura2016
- Camus2015
- Camus2016
- Pedersen2016
- Oh2016
- Xie2016
- Lin2001
- Stettinger2014
- Krammer2015
- Galtier2015
- Fey1997
- Enge-Rosenblatt2011

C.1.6 Distribution

A co-simulation framework is parallel and distributed when it allows each SU to be remote, across a wide area network.

This is very important since suppliers can, instead of transferring the SUs in executable form across computers, can make them available over the web. This offers much more control over how the SUs are used.

The same techniques used in parallelism can be used to promote distribution, but fault tolerance is also important.

References in this category:
- Friedrich2011
- Nutaro2011
- Busch2012
- Quaglia2012
- Eyisi2012
- Riley2011
- Roche2012
- Fitzgerald2010
C.1.7 Hierarchy

A hierarchical co-simulation framework is able to abstract a co-simulation scenario as a black box SU. This is very intuitive and promotes abstraction of complex systems.

References in this category:
- Quesnel2005
- Krammer2015
- Galtier2015

C.1.8 Scalability

A co-simulation framework is scalable when it supports a large number of SUs. It is intimately related to performance and parallelism.

References in this category:
- Fuller2013
- BenKhaled2014
- Said2016
- Galtier2015
- Enge—Rosenblatt2011

C.1.9 Platform Independence

A co-simulation framework is platform independent when it works on multiple computing platforms. For this to be achieved, a platform independent language, such as Java, can be used to coordinate the simulation.

References in this category:
- Bastian2011a
- Eyisi2012
- Riley2011
- Fitzgerald2010
C.1.10 Extensibility

A co-simulation framework is extensible when it can be easily extended to support new kinds of SUs, with new kinds of capabilities. A higher level, domain specific, language can be used to specify the behaviour in a platform agnostic way. Code is then generated from this description. The hypothesis is that the high level description can be more easily extended to describe new behaviour and that the code generation process can be adapted accordingly.

References in this category:
- Kuhr2013
- Krammer2015
- Enge—Rosenblatt2011
- Wetter2010

C.1.11 Accuracy

A co-simulation is accurate when the error between the trace produced and the correct trace is minimal. This can be achieved by error control mechanisms.

References in this category:
- Hoepfer2011
- Tomulik2011
- Gonzalez2011
- Nutaro2011
- Busch2012
- Schmol2012
- Schierz2012
- Gunther2012
- Fitzgerald2010
- Fitzgerald2013
- Benedikt2013
- Benedikt2013b
- Hafner2013
- Viel2014
- Sicklinger2014
- BenKhaled2014
- Camus2015
- Camus2016
- Oh2016
- Schierz2012a
- Stettinger2014
C.1.12 Open source

We consider open source the frameworks that make available the source code under certain licenses that are not paid for in any way.

References in this category:
- Roche2012
- Quesnel2005
- Andersson2016
- Galtier2015
- Acker2015
- Wetter2010

C.2 Simulator Requirements

This sub section covers the taxonomy that focuses on individual simulators’ capabilities.

C.2.1 Information Exposed

Frequency of State  The instantaneous frequency of the state of the sub-system can be used to adjust the communication step size.

Frequency of Outputs  The instantaneous frequency of the output of the sub-system can be used to adjust the communication step size, as is done in [32].

References in this category:
- Benedikt2013b

Detailed Model  Simulators that make the equations of the dynamic system available fall into this category.

References in this category:
- Schmoll2012
- Hassaire2012
- Schierz2012
- Zhang2014
- BenKhaled2012
- BenKhaled2014
- Yamaura2016
**Nominal Values of Outputs**  This information indicates the order of magnitude of output signals.

**Nominal Values of State**  This information indicates the order of magnitude of state signals.

**I/O Signal Kind**  The kind of output signal helps the master algorithm understand what assumptions are in a signal.

References in this category:  
- **Kuhr2013**

**Time Derivative**

**Output**  References in this category:
- **Hoepfer2011**
- **Tomulik2011**
- **Gunther2012**
- **Schierz2012a**
- **Carstens2003**
- **Quesnel2005**

**State**  References in this category:
- **Hoepfer2011**

**Jacobian**

**Output**  References in this category:
- **Tomulik2011**
- **Bastian2011a**
- **Busch2012**
- **Schierz2012**
- **Viel2014**
- **Sicklinger2014**
- **Arnold2001**
- **Arnold2010**
- **Schweizer2016**
- **Schweizer2015**
- **Schweizer2015a**

**State**
**Discontinuity Indicator**  A discontinuity indicator is a signal that indicates the presence of a discontinuity in the output of the SU.

**Deadreckoning model**  A deadreckoning model is a function that can be used by other SUs to extrapolate the behavior of this SU.

**Preferred Step Size**  References in this category:
- Acker2015

**Next Step Size**  The next step size indicates the next communication time that is appropriate for the current simulator.

- References in this category:
  - Lin2011
  - Gunther2012
  - Eyisi2012
  - Riley2011
  - Broman2013
  - Kounev2015
  - Quesnel2005

**Order of Accuracy**  The order of accuracy can be used to determine the appropriate input extrapolation functions to be used in a co-simulation scenario.

**I/O Causality**  *Feedthrough*

- References in this category:
  - Broman2013
  - Bogomolov2015
  - BenKhaled2012
  - BenKhaled2014
  - Saidi2016
  - Benedikt2016
  - Busch2011
  - Arnold2014a
  - Arnold2014
  - Acker2015

*Propagation Delay*

The propagation delay indicates how many micro-steps have to be performed before a change in the input affects an output. In Simulink, this is the number of delay blocks in chain from an input to an output.

**Input Extrapolation**  This information denotes the kind of input extrapolation being performed by the SU.

- References in this category:
  - Vie2013
State Variables
References in this category:
- Hoepfer2011
- Quesnel2005
- Arnold2014
- Acker2015

Serialized State
References in this category:
- Bastian2011a
- Broman2013
- Viel2013
- Bogomolov2015
- Camus2015
- Camus2016
- Schierz2012a
- Galtier2015

Micro-Step Outputs
This information denotes the output of the SU, evaluated at each of the micro-steps.
References in this category:
- Benedikt2013
- Viel2014
- Arnold2001
- Arnold2010

WCET
This denotes the worst case execution time.
References in this category:
- Faure2011
- BenKhaled2014
- Saidi2016

C.2.2 Causality
Causal
References in this category:
- Pedersen2015
- Lin2011
- Hoepfer2011
- Tomulik2011
- Sun2011
- Bastian2011a
- Friedrich2011
- Gonzalez2011
- Nutaro2011
- Busch2012
- Schmoll2012
- Ni2012
- Hassairi2012
A-Causal

C.2.3 Time Constraints

Analytic Simulation References in this category:

- Pedersen2015
- Lin2011
- Hoepfer2011
- Tomulik2011
- Sun2011
- Friedrich2011
- Gonzalez2011
- Nutaro2011
- Busch2012
- Schmoll2012
- Ni2012
- Hassairi2012
- Schierz2012
- Quaglia2012
- Al-Hammouri2012
- Eyisi2012
- Riley2011
- Roche2012
- Fitzgerald2010
- Fitzgerald2013
- Kudelski2013
- Broman2013
Benedikt2013
Benedikt2013b
Fuller2013
Wang2013
Hahner2013
Zhao2014
Li2011c
Awais2013b
Awais2013a
Kuhr2013
Viel2014
Sicklinger2014
Zhang2014
Konne2015
Bogomolov2015
Dols2016
BenKhaled2014
Saidi2016
Yamaura2016
Camus2015
Camus2016
Oh2016
Xie2016
Scherz2012a
Fourmigue2009
Liu2001
Carstens2003
Benedikt2016
Busch2011
Quesnel2005
Arnold2014a
Arnold2014
Arnold2001
Schweizer2013
Schweizer2015d
Sadjina2016
Busch2016
Arnold2010
Gu2001
Gu2004
Schweizer2016
Schweizer2015
Schweizer2015a
Andersson2016
Krammer2015
Gaillier2015
Scaled Real Time Simulation

Fixed
A simulator is fixed scaled real time when it simulated time progresses according to a fixed linear relationship with the real time.

References in this category:
- Faure2011
- Bian2015
- BenKhaled2012
- Pedersen2016
- Manbachi2016
- Stettinger2014
- Wetter2010
- Neema2014

Dynamic
A simulator is dynamic scaled real time when the relation between the simulated time and the real time can be changed throughout the simulation.

References in this category:
- Bombino2013

C.2.4 Rollback Support

None  References in this category:
- Pedersen2015
- Lin2011
- Hoepler2011
- Faure2011
- Sim2011
- Bastian2011
- Friedrich2011
- Gonzalez2011
- Schmol2012
- Ni2012
- Schierz2012
- Gunther2012
- Quaglia2012
- Al–Hammouri2012
- Eyisi2012
- Riley2011
- Roche2012
Single rollback means that the SU is capable of saving a certain state in the past (simulated time) and revert to that state. Once reverted, the SU cannot revert further in the past.

References in this category:

<table>
<thead>
<tr>
<th>Author</th>
<th>Year</th>
</tr>
</thead>
<tbody>
<tr>
<td>Tomulik</td>
<td>2011</td>
</tr>
<tr>
<td>Nutaro</td>
<td>2011</td>
</tr>
</tbody>
</table>
Multiple rollback means that the SU is capable of saving a certain state in the past (simulated time) and revert to that state. Once reverted, the SU revert further into the past as many times as necessary.

C.2.5 Availability

Remote References in this category:

Friedlisch2011
Busch2012
Quaglia2012
Eyisi2012
Riley2011
Roche2012
Fitzgerald2010
Fitzgerald2013
Kudelski2013
Fuller2013
Bombino2013
Zhao2014
Awais2013b
Awais2013a
Bian2015
Dols2016
Yamaura2016
Oh2016
Lin2001
Carstens2003
Galtier2015
Enge–Rosenblatt2011
Local References in this category:

Pedersen2015
Lin2011
Hoepler2011
Faure2011
Tomulik2011
Sun2011
Bastian2011a
Gonzalez2011
Nutaro2011
Schmol2012
Ni2012
Hassairi2012
Schierz2012
Günther2012
Al-Hammouri2012
Broman2013
Benedikt2013
Benedikt2013b
Wang2013
Hafer2013
Li2011c
Awais2013b
Awais2013a
Kuhr2013
Viel2014
Sicklinger2014
Zhang2014
Kounev2015
Bogomolov2015
BenKhaled2012
BenKhaled2014
Saidi2016
Camus2015
Camus2016
Xie2016
Manbach2016
Schierz2012a
Fourmigue2009
Stettinger2014
Benedikt2010
Busch2011
Quesnel2005
Arnold2014a
Arnold2014
Arnold2001
C.3 Framework Requirements

C.3.1 Standard

High Level Architecture  References in this category:
- Eyisi2012
- Riley2011
- Awais2013b
- Awais2013a
- Neema2014

Functional Mock-up Interface  References in this category:
- Pedersen2015
- Sun2011
- Bastian2011a
- Broman2013
- Wang2013
- Awais2013b
- Awais2013a
- Kuhr2013
- Viel2014
- Bogomolov2015
- Dols2016
- BenKhaled2012
- BenKhaled2014
- Said2016
- Camus2015
- Camus2016
- Pedersen2016
- Schierz2012a
Functional Digital Mock-up

References in this category:

Enge—Rosenblatt2011

C.3.2 Coupling

Input/Output Assignments

References in this category:

Pedersen2015
Lin2011
Faure2011
Sun2011
Bastian2011a
Friedrich2011
Gonzalez2011
Nutaro2011
Busch2012
Schmolz2012
Ni2012
Hassairi2012
Gunther2012
Quaglia2012
Al—Hammouri2012
Eyisi2012
Riley2011
Roche2012
Fitzgerald2010
Fitzgerald2013
Kudelski2013
Broman2013
Benedikt2013
Benedikt2013b
Fuller2013
Bombino2013
Wang2013
Hafer2013
Zhao2014
Li2011c
Awais2013b
Awais2013a
Algebraic Constraints  References in this category:

Tomulik2011
Friedrich2011
Schierz2012
Viel2014
Sicklinger2014
Arnold2001
Schweizer2014
Schweizer2015d
Arnold2010
Gu2001
C.3.3 Number of Simulation Units

Two references in this category:

- Lin2011
- Sun2011
- Gonzalez2011
- Nutaro2011
- Busch2012
- Schmolz2012
- Ni2012
- Hassari2012
- Quaglia2012
- Al-Hammouri2012
- Eyisi2012
- Riley2011
- Roche2012
- Fitzgerald2010
- Fitzgerald2013
- Kudelski2013
- Benedikt2013
- Benedikt2013b
- Fuller2013
- Bombino2013
- Wang2013
- Zhao2014
- Li2011c
- Zhang2014
- Komek2015
- Bogomolov2015
- Bann2015
- Dols2016
- Pedersen2016
- Oh2016
- Xie2016
- Manbachi2016
- Fourmigue2009
- Liu2001
- Carstens2003
- Stettlager2014
- Schweizer2013
- Gu2001
Three or More References in this category:

Pedersen2015
Hoepler2011
Faure2011
Tomulik2011
Bastian2011a
Friedrich2011
Schierz2012
Gunther2012
Broman2013
Hafner2013
Awais2013b
Awais2013a
Knhr2013
Viel2013
Sicklinger2014
BenKhaled2012
BenKhaled2014
Saidi2016
Yamaura2016
Camus2015
Camus2016
Schierz2012a
Benedikt2016
Busch2011
Quesnel2005
Arnold2014a
Arnold2014
Arnold2001
Schweizer2015d
Sadjina2016
Busch2016
Arnold2010
Gu2004
Schweizer2016
Schweizer2015
Schweizer2015a
Andersson2016
Krammer2015
Galtier2015
Acker2015
Enge−Rosenblatt2011
Karner2010a
Aslan2015
C.3.4 Domain

CT References in this category:

Pedersen2015
Hoepler2011
Faure2011
Tomulik2011
Sun2011
Bastian2011a
Friedrich2011
Gonzalez2011
Busch2012
Schmolz2012
NI2012
Hassairi2012
Schierz2012
Gunther2012
Quaglia2012
Al-Hammouri2012
Roche2012
Fitzgerald2010
Fitzgerald2013
Broman2013
Benedikt2013
Benedikt2013b
Bombino2013
Wang2013
Halmer2013
Zhao2014
Li2011c
Kuhr2013
Viel2014
Sicklinger2014
Zhang2014
Bogomolov2015
Bian2015
Dols2016
BenKhaled2012
BenKhaled2014
Saidi2016
Yamaura2016
Pedersen2016
Oh2016
References in this category:

Lin2011
Nutaro2011
Al-Hammouri2012
Eyisi2012
Riley2011
Fitzgerald2010
Fitzgerald2013
Kulesi2013
Fuller2013
Awais2013b
Awais2013a
Kuhr2013
Zhang2014
Kounev2015
Bogomolov2015
C.3.5 Dynamic structure

References in this category:

Karner2010a

C.3.6 Co-simulation Rate

Single References in this category:

Pedersen2015
Lin2011
Hoepfer2011
Faure2011
Tomulik2011
Sun2011
Bastian2011a
Friedrich2011
Nutaro2011
Busch2012
Schmoll2012
Ni2012
Hassairi2012
Schierz2012
Gunther2012
Quaglia2012
Al-Hammouri2012
Eyisi2012
Riley2011
Roche2012
Fitzgerald2010
Fitzgerald2013
Kudelski2013
Broman2013
Benedikt2013
Benedikt2013b
Fuller2013
Bombine2013
Wang2013
Hainer2013
Zhao2014
Multi-rate co-simulation denotes that the framework distinguishes between slow and fast sub-systems and dimensions the communication step size accordingly, providing for interpolation/extrapolation of the slow systems.

References in this category:
Gonzalez2011
Awais2013b
Awais2013a
Kuhr2013
Camus2015
Camus2016
Busch2011
Quesnel2005
Arnold2014a
Arnold2014
Arnold2001
Schweizer2014
Schweizer2015d
Sadjina2016
Gu2001
Gu2004
Schweizer2016
C.3.7 Communication Step Size

Fixed References in this category:

- Pedersen2015
- Faure2011
- Tomulik2011
- Bastian2011a
- Friedrich2011
- Gonzaley2011
- Busch2012
- Schmoll2012
- Ni2012
- Hassairi2012
- Schierz2012
- Quaglia2012
- Roche2012
- Kudelski2013
- Benedikt2013
- Benedikt2013b
- Bombino2013
- Hafner2013
- Zhao2014
- Li2011c
- Awais2013a
- Vie2013
- Sicklinger2014
- Zhang2014
- Bian2015
- Dols2016
- BenKhaled2012
- BenKhaled2014
- Pedersen2016
- Oh2016
- Xie2016
- Manbachi2016
- Carstens2003
- Stettiager2014
- Arnold2014a
- Arnold2001
References in this category:

- Lin2011
- Hoepfer2011
- Sun2011
- Nutaro2011
- Gunther2012
- Al-Hammouri2012
- Eyisi2012
- Riley2011
- Fitzgerald2010
- Fitzgerald2013
- Broman2013
- Fuller2013
- Wang2013
- Awais2013b
- Kuhr2013
- Kounev2015
- Bogomolov2015
- Camus2015
- Camus2016
- Schierz2012a
- Benedikt2016
- Busch2011
- Quesnel2005
- Arnold2014
- Sadjina2016
- Galtier2015
- Neema2014
C.3.8 Strong Coupling Support

None – Explicit Method

References in this category:

- Pedersen2015
- Lin2011
- Hoepfer2011
- Faure2011
- Sun2011
- Friedrich2011
- Gonzalez2011
- Nutaro2011
- Busch2012
- Schmol2012
- Ni2012
- Hassairi2012
- Schierz2012
- Gunther2012
- Quaglia2012
- Al-Hammouri2012
- Eyisi2012
- Riley2011
- Roche2012
- Fitzgerald2010
- Fitzgerald2013
- Kudelski2013
- Broman2013
- Benedikt2013
- Benedikt2013b
- Fuller2013
- Bombino2013
- Wang2013
- Hainer2013
- Zhao2014
- Li2011c
- Awais2013b
- Awais2013a
- Zhang2014
- Kounev2015
- Bogomolov2015
- Bian2015
- Dols2016
- BenKhaled2012
- BenKhaled2014
- Saidi2016
- Yamaura2016
- Camus2015
- Camus2016
Partial – Semi-Implicit Method References in this category:

- Busch2012
- Schweizer2014
- Schweizer2015d
- Schweizer2016
- Schweizer2015
- Schweizer2015a

Full – Implicit Method References in this category:

- Tomulik2011
- Bastian2011a
- Busch2012
- Viel2014
- Sicklinger2014
- Liu2001
- Arnold2001
- Arnold2010
- Acker2015

C.3.9 Results Visualization

Postmortem The results are available after the simulation. References in this category:

- Pedersen2015
- Lin2011
- Hoepfner2011
- Faure2011
- Tomulik2011
Live References in this category:
- Hassairi2012
- Eyisi2012
- Riley2011
- Fitzgerald2010
- Fitzgerald2013
- Yamaura2016
- Pedersen2016
- Enge—Rosenblatt2011
- Neema2014

Interactive References in this category:
- Bombino2013
- Yamaura2016

C.3.10 Communication Approach

Jacobi References in this category:
- Pedersen2015
- Hoepfler2011
- Paure2011
- Tomulik2011
- Bastian2011a
- Friedrich2011
- Schmoll2012
- Schierz2012
- Gunther2012
- Kudelski2013
Gauss-Seidel  References in this category:
Lin2011
Hoepfer2011
Sun2011
Bastian2011a
Gonzalez2011
Nutaro2011
Busch2012
Ni2012
Hassairi2012
Schierz2012
Quaglia2012
Al–Hammouri2012
Eyisi2012
Riley2012
Roche2012
Fitzgerald2010
Fitzgerald2013
Benedikt2013
Benedikt2013b
Fuller2013
Bombino2013
Wang2013
Hainer2013
Zhao2014
Li2011c
Awais2013b
Awais2013a
Kuhr2013
Viel2014
Sicklinger2014
Kounev2015
Dols2016
BenKhaled2012
BenKhaled2014
Camus2015
Camus2016
Pedersen2016
Oh2016
Carstens2003
Stettinger2014
Quesnel2005
Arnold2014
Arnold2001
Busch2016
Arnold2010
Acker2015
Aslan2015

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D List of Acronyms

CPS  Cyber-Physical System
CT   Continuous Time
DE   Discrete Event
SU   Simulation Unit
DEVS Discrete Event System Specification
DTSS Discrete Time System Specification
FMI  Functional Mock-up Interface
FR   Framework Requirement
GVT  Global Virtual Time
IP   Intellectual Property
IVP  Initial Value Problem
NFR  Non-Functional Requirement
ODE  Ordinary Differential Equation
SR   Simulator Requirement