# Community Detection in Graphs through Correlation

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## **Community Detection**

### • What and why?







# **Related Work**

- Methods
  - Cut-based (M. Girvan and M. E. J. Newman, 2002)
  - Spectral-based (U. Luxburg, 2007)
  - Density-based (S. Mancoridis, et al., 1998)
  - Modularity-based (A. Clauset, et al., 2004)
  - Statistical-inference-based (M. E. J. Newman, 2013)

### • Goal

- More connections inside each community
- Fewer connections across different communities

## **Opportunities for improvement**

- Feature selection
  - Spectral-based
- Objective function
  - Cut-based
  - Modularity-based (Our focus in this paper)
  - Statistical-inference-based
- Search procedure
  - Greedy search
  - EM
  - Simulated annealing

# Major problem of modularity

- Resolution problem (Lancichinetti & Fortunato 2011)
  - $K_m$  is an m-clique
  - The detected communities are marked by circles with dash lines.
- Multi-resolution (Reichardt & Bornholdt 2006)
  - Further divide each detected community
  - Bias (Xiang et al. 2012)
    - Merge small communities
    - Split large communities



### Connection with itemset search

 Graph communities: number of internal edges is greater than expected under assumption of random partition

 Correlated itemsets: occur more than expected under the assumption of item independence

• Connection: modularity = leverage

### Correlated Itemsets (Duan, et al. 2014)

Given: itemset  $S = \{I_1, I_2, \dots, I_m\}$  with m items in a dataset with n transactions

- True probability:  $tp_s = P(S)$
- Expected probability:  $ep_s = \prod_{i=1}^m P(I_i)$
- Correlation measure: M<sub>S</sub> = f(tp<sub>S</sub>, ep<sub>S</sub>)
  Chi-square: (tp<sub>S</sub>-ep<sub>S</sub>)<sup>2</sup>/ep<sub>S</sub>
  Probability ratio / Lift: tp<sub>S</sub>/ep<sub>S</sub>
  Leverage: tp<sub>S</sub> ep<sub>S</sub>
  Likelihood ratio: (tp<sub>S</sub><sup>tp<sub>S</sub></sup>\*(1-tp<sub>S</sub>)<sup>1-tp<sub>S</sub></sup>/ep<sub>S</sub><sup>tp<sub>S</sub></sup>\*(1-ep<sub>S</sub>)<sup>1-tp<sub>S</sub></sup>)

### Correlated itemset example

- t1: Beef, Chicken, Milk
- t2: Beef, Cheese
- t3: Cheese, Boots
- t4: Beef, Chicken, Cheese
- t5: Beef, Chicken, Clothes, Cheese, Milk
- For the itemset {Beef, Chicken}

• 
$$tp = \frac{3}{5}, ep = \frac{4}{5} * \frac{3}{5}, Leverage = tp - ep = \frac{3}{25}$$

## **Modularity Function**

$$Q = \frac{1}{2m} \sum_{i \in G, j \in G} \left( w_{ij} - \frac{k_i * k_j}{2m} \right) * \delta(v_i, v_j)$$

- *n*: the number of nodes
- *m*: the number of links
- $w_{ij}$ : the edge between node *i* and *j*
- $k_i$ : the degree of node i
- $\delta(v_i, v_j)$ : the Kronecker delta function
  - $\delta(v_i, v_j) = I$  when  $v_i$  and  $v_j$  are in the same community
  - $\delta(v_i, v_j)$ =0 otherwise

For partition  $\{G_1, G_2, \dots, G_l\}$  on graph G

- $k_i$ : degree of node i
- $k_i^{internal}$ : number of nodes in the same group of node *i* that connect to node *i*
- For the green partition:
  - $\sum_{i \in G_p} k_i = (2+2+3) = 7$
  - $\sum_{i \in G_p} k_i^{internal} = (2+2+2) = 6$
  - Total number of edges: m = 10





For the green partition: (1)  $\sum_{i \in G_p} k_i = (2 + 2 + 3) = 7$ (2)  $\sum_{i \in G_p} k_i^{internal} = (2 + 2 + 2) = 6$ (3) Total number of edges: m = 10 The probability of the edge

- (1) Both ends in the green partition:  $6/20 = \sum_{i \in G_p} k_i^{internal} / 2m$
- (2) Started from the green partition:  $7/20 = \sum_{i \in G_p} k_i / 2m$
- (3) Ended in the green partition:  $7/20 = \sum_{i \in G_p} k_i / 2m$

If we randomly select an edge from the doublydirected graph,

- The true probability of the edge in  $G_p$ :  $tp = \frac{\sum_{i \in G_p} k_i^{internal}}{2m}$
- Probability the edge started from  $G_p: \frac{\sum_{i \in G_p} k_i}{\sum_{j=2m} 2m}$
- Probability the edge ended in  $G_p: \frac{\sum_{j \in G_p} k_j^2}{2m}$
- The expected probability of the edge in  $G_p$  under the assumption of independence:

$$ep = \frac{\sum_{i \in G_p} k_i}{2m} * \frac{\sum_{j \in G_p} k_j}{2m}$$

For partition  $\{G_1, G_2, \dots, G_l\}$  on graph G

• 
$$Q = \frac{1}{2m} \sum_{i \in G, j \in G} \left( w_{ij} - \frac{k_i * k_j}{2m} \right) * \delta(v_i, v_j)$$

- We define  $Q_p$  as the partial modularity for the group p where

$$Q_p = \frac{1}{2m} \sum_{i \in G_p, j \in G} \left( w_{ij} - \frac{k_i * k_j}{2m} \right) * \delta(v_i, v_j)$$

• 
$$Q = \sum_{p=1}^{l} Q_p$$

• Partial modularity

$$\circ \quad Q_p = \frac{\sum_{i \in G_p} k_i^{internal}}{2m} - \frac{\sum_{i \in G_p} k_i}{2m} * \frac{\sum_{j \in G_p} k_j}{2m}$$

- If we randomly select an edge from the doubly-directed graph,
  - The true probability of the edge in  $G_p$ :

$$tp_p = \frac{\sum_{i \in G_p} k_i^{internal}}{2m}$$

• Expected probability of the edge in  $G_p$  under the assumption of independence:

$$ep_p = \frac{\sum_{i \in G_p} k_i}{2m} * \frac{\sum_{j \in G_p} k_j}{2m}$$

- Connecting correlation with modularity
  - $^\circ\,$  For a given partition  ${\it G}_p,$  partial modularity  $Q_p=tp_p-ep_p$
  - For a given itemset S, leverage =  $tp_s ep_s$
- For any correlation function f(tp, ep), the partial modularity function can be changed accordingly.

• Chi-square: 
$$\frac{(tp_s - ep_s)^2}{ep_s}$$

- Probability ratio / Lift:  $\frac{tp_s}{ep_s}$
- Leverage:  $tp_s ep_s$
- Likelihood ratio:  $\frac{tp_s^{tp_s}*(1-tp_s)^{1-tp_s}}{ep_s^{tp_s}*(1-ep_s)^{1-tp_s}}$

### Outline

- Basic concepts
  - Correlated itemset search
  - Modularity function
- Connecting modularity with leverage in correlated itemset search
- Upper bound analysis
- Evaluation

# Upper bound analysis

- Correlation Property
  - The correlation function monotonically increases with the decrease of *ep* when *tp* remains the same.
- Understanding the bias to the community size
  - Given a group  $G_p$  with fixed  $tp = \frac{\sum_{i \in G_p} k_i^{internal}}{2m}$
  - Partial modularity has the highest possible value when  $ep = \frac{\sum_{i \in G_p} k_i}{2m} * \frac{\sum_{j \in G_p} k_j}{2m}$  reaches its lowest value  $tp^2$



## Upper bound analysis

• The highest possible partial modularity:  $f(tp, ep = tp^2)$ 



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### Experiments

- Modify the objective function
- Greedy search (hierarchical clustering)
- Baseline:
  - Modularity-based methods (Leverage)
- Datasets:
  - Real life
  - Simulated with LFR model (Lancichinetti et al. 2008)
- Evaluation measures:
  - Rand Index (Rand1971), Jaccard, F-measure, Normalized mutual information (Danon 2005)



## Real life datasets

• Karate club: two equal size communities



• College football: 12 equal size communities



### **Real life datasets**

Data Set	Measure	NMI	Jaccard	RI	F-measure	DNC	ANC
	$\chi^2$	0.4852	0.2842	0.6453	0.4426	7	2
Karate	PR	0.3868	0.0945	0.5561	0.1728	14	2
	Leverage	0.6925	0.6833	0.8414	0.8118	3	2
	LR	0.5385	0.3958	0.6952	0.5671	5	2
	$\chi^2$	0.9141	0.7571	0.9793	0.8618	14	12
$\operatorname{Football}$	PR	0.6864	0.0829	0.9240	0.1531	55	12
	Leverage	0.6977	0.3622	0.8807	0.5317	6	12
	LR	0.9086	0.7897	0.9812	0.8825	12	12





(c) Leverage



### Simulated datasets

- Parameters:
  - Minimal community size: 50, 500, or 5000
  - Community structure ratio  $\beta$ : 5, 10, or 20





### Experiments

### • NMI when fixing Min-community-size



### Experiments

0

Fuzzy

Middle

 Number of detected groups when fixing <u>Min-community-size</u>



Clear

Chi-square

Actual

## Summary

- Connection between community detection and correlation search
- Conclusion
  - Modularity is good only when there are large and clear communities
  - Likelihood ratio is robust to any type of communities
  - Probability ratio partitions the whole graph into small communities with 2 or 3 objects



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