

Comparative Correlatives as Anaphora to Differentials

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1. Comparative and Equative Correlatives in Romanian

The two goals of this paper are (i) to establish that there are comparative correlatives that are not comparative conditionals and that the semantics of such correlatives crucially involves a relation (possibly the identity relation) between differentials (against much of the previous literature, e.g., McCawley 1988, Beck 1997) and (ii) to argue that a unified analysis should be given for such non-conditional, differential-based comparative (and equative) correlatives and the more familiar, conditional-like comparative correlatives.

Correlatives are “biclausal topic-comment structures [. . .] [in which] the dependent clause introduces one or more topical referents to be commented on by the matrix clause, where each topical referent must be picked up by – correlated with – an anaphoric proform.” (Bittner 2001). Differentials, e.g., *2 cm* in the comparative *Gabby is 2 cm taller than Linus*, specify the difference between two measures, e.g., between Gabby’s and Linus’s heights.

1.1. Non-Conditional Comparative and Equative Correlatives

The first point is established by the existence of comparative correlatives with overt *than*-phrases (which are usually not acceptable in English) like the one in (1) below. Unless otherwise noted, all examples are in Romanian or English.

- (1) Cu cât e mai înalt fratele decât sora, (tot) cu
With how much is more tall brother.the than sister.the, (also) with
atât e mai înalt tatăl decât mama.
that much is more tall father.the than mother.the

‘The brother is taller than the sister by a certain amount and the father is taller than the mother by the same amount.’

Intuitively, sentence (1) is true iff (i) the brother is taller than the sister and the father is taller than the mother, that is, there is *no conditionality* (no ‘if the brother is taller than the sister . . .’ kind of interpretation), and (ii) the difference in height between the brother and the sister is the same as the difference in height between

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the father and the mother, that is, the correlative *equates* the two differentials under consideration (this is particularly clear if the particle *tot* is present).

These truth conditions are summarized by the formula in (2) below. Subscripts on terms indicate their type. The underlying logic is classical (many-sorted) type logic with three basic types: the usual t (truth values – the domain D_t of type t consists of $\{\mathbb{T}, \mathbb{F}\}$) and e (individuals – variables of type e : x, x', y, y', \dots) plus the type of degrees γ (from the Latin *gradus* – variables of type γ : d, d', \dots). Differentials are intervals, i.e., convex sets of degrees of type γt (variables of type γt : D, D', \dots). A formula of the form *tall*(x, d), where *tall* is a binary relation of type $\gamma(et)$, is interpreted as: the individual x is at least d -tall (tall at least to degree d). I take definite descriptions to be anaphoric (for simplicity): the brother, the sister, the father and the mother retrieve the contextually specified values for the variables x, x', y and y' respectively. Set-theoretic difference is symbolized as ‘\’.

$$(2) \quad \exists D_{\gamma t}(D \neq \emptyset \wedge D = \{d_{\gamma} : tall(x, d)\} \setminus \{d'_{\gamma} : tall(x', d')\}) \wedge \\ \exists D'_{\gamma t}(D' \neq \emptyset \wedge D' = \{d_{\gamma} : tall(y, d)\} \setminus \{d'_{\gamma} : tall(y', d')\}) \wedge D = D')$$

The equative correlative in (3) below is non-conditional in the same way: on its most salient reading, sentence (3) is true iff (i) Irina is (significantly) beautiful and (significantly) smart and (ii) the extent to which Irina is beautiful and the extent to which she is smart are (in some sense) equated / similar / comparable.¹

- (3) Pe cât e Irina de frumoasă, (tot) pe atât e de deșteaptă.
 PE how much is Irina DE beautiful, (also) PE that much is DE smart
 ‘Irina is beautiful to a certain, significant extent and she is smart to the same, equally significant extent.’

1.2. Conditional Comparative and Equative Correlatives

The idea that non-conditional, differential-based comparative correlatives and the more familiar, conditional-like comparative correlatives should receive a unified analysis is supported by the correlatives in (4) and (5) below, which have the same basic syntax – except that the *than*-phrase is missing in (4) – and the same morphology – a *wh*-indefinite in the topic clause, a.k.a. apodosis, and an anaphoric demonstrative in the comment clause, a.k.a. protasis – as sentence (1) above.

- (4) Cu cât e un avocat mai agresiv, cu atât e mai eficient.
 With how much is a lawyer more aggressive, with that much is more efficient
 ‘The more aggressive a lawyer is, the more efficient s/he is.’
- (5) Cu cât e un număr natural mai mare decât altul, (#tot) cu
 With how much is a number natural more great than another, (#also) with
 atât e pătratul lui mai mare decât pătratul celuiilalt.
 that much is square.the it.Gen more great than square.the other.one.Gen

¹The first meaning component is not necessarily present in equative correlatives and it is systematically absent in certain cataphoric equative correlatives (where the apodosis precedes the protasis). For example, as a criterion for estimating the age of a tree, *Un copac e pe atât de bătrîn pe cât e de gros* (A tree is as old as it is thick) applies to all trees, including the very young and thin ones.

‘The greater one natural number is (than another), the greater its square is (than the square of the other one).’

Moreover, the interpretations of (4) and (5) are very closely related to the interpretation of (1). Intuitively, sentence (4) has two salient interpretations – it can be paraphrased by either the conditional in (6a) or the one in (6b) below, as Beck (1997) points out with respect to similar examples in German. These two interpretations are not necessarily two distinct readings, since (6b) is ultimately just a refinement of (6a) that examines the aggressiveness and efficiency of lawyers at various times as opposed to a single, contextually salient temporal interval.

- (6) a. If a lawyer x is more aggressive than a lawyer y by a certain amount, then x is more efficient than y by a corresponding amount.
 b. If a lawyer x is more aggressive at time t than at time t' by a certain amount, then x is more efficient at t than at t' by a corresponding amount.

As (6a,b) above indicate, such conditional comparative correlatives involve a relation between differentials just as the non-conditional comparative correlative in (1) does. The truth conditions in (6a) are formalized in (7) below. The correspondence between the differential intervals D and D' is implicitly contributed by the relative scope of the two quantifiers $\forall D(D \neq \emptyset \dots \rightarrow \exists D'(D' \neq \emptyset \dots))$: any non-empty aggressiveness interval D induced by a pair of lawyers x and y is (somehow) related to a non-empty efficiency interval D' induced by the same pair of lawyers.

$$(7) \quad \forall x_e \forall y_e \forall D (lw(x) \wedge lw(y) \wedge D \neq \emptyset \wedge D = \{d : aggr(x, d)\} \setminus \{d' : aggr(y, d')\} \\ \rightarrow \exists D' (D' \neq \emptyset \wedge D' = \{d : eff(x, d)\} \setminus \{d' : eff(y, d')\}))$$

The interpretation of sentence (5) when the particle *tot* is present provides the strongest argument in favor of differential-based truth conditions for conditional comparative correlatives: (5) with *tot* is true iff, for any two natural numbers m and n such that m is greater than n , the difference $m - n$ is *identical* to the difference between their squares $m^2 - n^2$, i.e., $\forall m_e \in \mathbb{N} \forall n_e \in \mathbb{N} (m > n \rightarrow m - n = m^2 - n^2)$ (I take the set of natural numbers \mathbb{N} to be a subset of the domain of individuals D_e). Thus, the relation between differentials implicitly contributed by $\forall D(D \neq \emptyset \dots \rightarrow \exists D'(D' \neq \emptyset \dots))$ is now required to be the identity relation – which is why (5) with *tot* is intuitively false. And this falsity cannot be derived if the interpretation of (5) does not involve pairs of differential intervals that are required by *tot* to be identical. In contrast, (5) without *tot* is intuitively true because it just says that, for any natural numbers m and n such that $m > n$, their positive difference $m - n$ corresponds to a positive difference between their squares $m^2 - n^2$.² The truth conditions of (5) are formalized in (8) below. The particle *tot* contributes the last conjunct $D = D'$.

$$(8) \quad \forall m_e \forall n_e \forall D (m \in \mathbb{N} \wedge n \in \mathbb{N} \wedge D \neq \emptyset \wedge D = \{d : great(m, d)\} \setminus \{d' : great(n, d')\} \\ \rightarrow \exists D' (D' \neq \emptyset \wedge D' = \{d : great(m^2, d)\} \setminus \{d' : great(n^2, d')\} \wedge D = D'))$$

²Note that, for a given n , the correspondence is an increasing function $f(m - n) = m^2 - n^2$. In general, however, i.e., for varying n and m , the correspondence between differentials is neither increasing nor functional: (i) not increasing because, for example, $3 - 1 > 10 - 9$, but $3^2 - 1^2 < 10^2 - 9^2$; (ii) not functional because, for example, $2 - 1 = 10 - 9$, but $2^2 - 1^2 \neq 10^2 - 9^2$.

The proposed truth conditions for conditional comparative correlatives are a refinement of the truth conditions in Beck (1997) (cf. p. 246:37''a, p. 252:51 and p. 253:52). The main difference is that the differential contributed by the protasis is available for anaphoric retrieval in the apodosis so that the protasis and apodosis differentials can be related (and, in certain cases, equated). We will see that this difference and the way in which the truth conditions are compositionally derived enable us to provide a unified analysis of both conditional and non-conditional comparative correlatives, with or without a *than*-phrase.

The proposed refinement of the truth conditions preserves an important result in Beck (1997), namely that comparative correlatives induce a correspondence between maximal degrees that is *monotonically increasing* and *non-functional* – and, just like in Beck (1997), these two properties do not need to be separately stipulated, but follow automatically from the truth conditions.

Consider, for example, (4) above (and its English counterpart) under the reading in (6a). As noted in Beck (1997: 257 et seqq), an intuitively attractive paraphrase of (4) involves a monotonically increasing correspondence between the maximal degrees of aggressiveness and the maximal degrees of efficiency attributable to lawyers, which covary in a ‘proportional’ way. Formally (cf. Beck 1997: 258,(63b)):

$$(9) \quad \exists F_{\gamma\gamma}(\text{MON}^\uparrow(F) \wedge \\ \forall x \forall d (lw(x) \wedge d = \text{MAX} \{d' : aggr(x, d')\} \rightarrow F(d) = \text{MAX} \{d' : eff(x, d')\})), \\ \text{where } \text{MON}^\uparrow(F) \text{ is true iff } \forall d \forall d' (d < d' \rightarrow F(d) < F(d')) \\ \text{and } \text{MAX} \{d' : \phi(d')\} := \iota d. \phi(d) \wedge \forall d' (\phi(d') \rightarrow d' \leq d)^3$$

Unlike Beck (1997) or the present proposal, this analysis does not involve pairs of lawyers. But, just like Beck (1997), it does not analyze comparative correlatives in terms of a relation between differentials: the correspondence F in (9) between the protasis and the apodosis is formulated in terms of degrees. Hence, this analysis also fails to generalize to (conditional or non-conditional) comparative correlatives with overt *than*-phrases and comparative correlatives that require identity of differentials. But, independently of these issues, Beck (1997) puts forth two arguments against approaches formulated in terms of degree-based correspondences like the one in (9) above – and it is important to see that they do not apply to the present approach, formulated in terms of differential-based correspondences.

The first argument is that the correspondence F has to be strictly monotonically increasing, otherwise we would incorrectly allow for functions F that are not order preserving, e.g., functions mapping two distinct aggressiveness degrees $d < d'$ to the same efficiency degree $F(d) = F(d')$ (lawyer x is more aggressive than – but as efficient as – lawyer y) or two aggressiveness degrees $d < d'$ to two

³The following variant of this formalization has been brought to my attention by Pranav Anand (p.c.): $\exists F_{\gamma\gamma}(\text{MON}^\uparrow(F) \wedge \forall x \forall d (lw(x) \wedge \text{MAX} \{d' : aggr(x, d')\} > d \rightarrow \text{MAX} \{d' : eff(x, d')\} > F(d)))$. This version, which builds on the informal proposal in McCawley (1988), is preferable because of its closeness to the surface form of (4) (or its English counterpart). Moreover, as far as I can see, it is not subject to the ‘functional’ problem raised in Beck (1997: 257 et seqq). But this version also fails to generalize to comparative correlatives with overt *than*-phrases and comparative correlatives that require identity of differentials.

inversely-ordered efficiency degrees $F(d') < F(d)$ (lawyer x is more aggressive – but less efficient – than lawyer y). The MON^\uparrow property needs to be separately stipulated in (9). In contrast, there is no need for such a stipulation in Beck’s or the present analysis. The reason is the same in both cases: requiring every non-empty differential D in the protasis to have a corresponding non-empty differential D' in the apodosis automatically ensures that every degree increase on the protasis scale is matched by a corresponding degree increase on the apodosis scale.

The second argument against approaches like (9) is that the correspondence between the two scales is not necessarily functional. Consider, for examples, the correlative (*In last year’s games*) *The warmer it got, the more goals Luise scored*, based on Beck (1997: 261,72). This correlative is intuitively true in a situation in which two games took place while the temperature was 25°C and Luise scored 3 and 4 goals during those games – as long as she scored 2 goals when the temperature was 20°C and 5 goals when the temperature was 30°C.

I take this argument to show that the degree-based correspondence intuitively established between the protasis and the apodosis is merely relational. Furthermore, this relation R needs to be strictly monotonically increasing in the following sense: $\forall d \forall d' \forall d'' \forall d''' (d < d' \wedge dRd'' \wedge d'Rd''' \rightarrow d'' < d''')$. Now, the relation between differentials implicitly contributed by the formula $\forall D (D \neq \emptyset \dots \rightarrow \exists D' (D' \neq \emptyset \dots))$ induces precisely this kind of monotonically increasing relation between degrees – and it induces it in a way that closely follows the overt morphosyntax of Romanian comparative correlatives like (1) and (5) above.

Further support for a unified analysis of conditional and non-conditional correlatives is provided by conditional equative correlatives like (10) below (interpreted roughly like (4)), which have the same morphology and syntax as non-conditional equative correlatives like (3) above.

(10) Pe cât e de agresiv un avocat, pe atât e de eficient.
 PE how much is DE aggressive a lawyer, PE that much is DE efficient
 ‘The aggressiveness of a lawyer is proportional to her/his efficiency.’

(11) #Pe cât de mare e un număr, pe atât de mare e pătratul lui.
 #PE how much DE great is a number, PE that much DE great is square.the it.Gen
 ‘#A number is equal to its square.’

As (11) above shows, the interpretation of such conditional equative correlatives is more constrained than the interpretation of their comparative correlative counterparts: (11) is not acceptable, in contrast to its counterpart without *tot* in (5) above, because it can only be interpreted as requiring the identity of two intervals. In this particular case, we falsely equate a number m and its square m^2 , for any number m .

2. Comparative Correlatives as Anaphora to Differentials

The main proposal is that the Romanian *atît* (that much) in (1), (3), (4), (5) and (10) above is anaphoric to differential intervals, i.e., *atît* is a proform in the degree

domain, and the *wh*-differential *cît* (how much) is an indefinite introducing a non-empty interval, anaphorically retrieved by *atît*.⁴ The idea that *atît* is an interval-based proform is supported by its anaphoric use in (12) below, which is similar to (1) above, and by its deictic use in (13) and its cataphoric use in (14) below.

- (12) Fratele e mai înalt decât sora cu 2 cm, iar tatăl e mai
 Brother.the is more tall than sister.the with 2 cm, and father.the is more
 înalt decât mama tot cu atât.
 tall than mother.the also with that much
 ‘The brother is 2 cm taller than the sister and the father is taller than the mother by the same amount.’
- (13) E atât de obosită.
 Is that much DE tired.f.sg
 ‘She is so tired.’
- (14) E atât de obosită încît o = doare capul.
 Is that much DE tired.f.sg that her.Acc = hurts head.the
 ‘She is so tired that she has a headache.’

The anaphora analysis enables us to capture the intuitive parallels between the interpretations of degree-based and individual-based correlatives. Consider, for example, the ‘singular’ / referential individual-based correlative in (15) below and the ‘plural’ / quantificational individual-based correlative in (16). Sentence (15) is parallel to (1) and (3): we refer to a single individual or a single pair of correlated intervals. Sentence (16) is parallel to (4), (5) and (10): we refer to a set of individuals or a set of pairs of correlated intervals.⁵

- (15) Care fată și = a = uitat ieri haina, pe aceea
 Which girl her.Dat = HAS = forgotten yesterday coat.the, PE that one
 o = caută tatăl ei.
 her.Acc = look for father.the her.Gen
 ‘The father of the girl that forgot her coat yesterday is looking for her.’
- (16) Pe care om l = a = interogat Securitatea, în acela nu
 PE which person him.Acc = HAS = interrogated security.the, in that one not
 mai am încredere.
 anymore have.1sg trust
 ‘I do not trust any person (whatsoever) that the secret police interrogated.’

Extending the investigation of anaphoric (and quantificational) parallels across domains initiated in Partee (1973, 1984) to encompass the degree domain is further supported by the following phenomena, exemplified in English:

⁴See Jespersen (1965) and den Dikken (2005) for related suggestions with respect to English comparative correlatives.

⁵See Dayal (1996), Bittner (2001) and Brasoveanu (2008) for more discussion of the semantics of individual-based correlatives.

- (17) Donkey anaphora:
- a. Every child that ate a lot of vanilla ice cream yesterday ate twice as much chocolate ice cream today.
 - b. Every farmer who owns a donkey beats it.
- (18) Quantificational subordination:
- a. Harvey eats a lot of vanilla ice cream at every convention, but Linus always eats twice as much chocolate ice cream.
 - b. Harvey courts a woman at every convention. She always comes to the banquet with him. (Karttunen 1976)
- (19) Modal subordination:
- a. Harvey might bring a lot of vanilla ice cream to the party. In which case Linus would get competitive and bring twice as much chocolate ice cream.
 - b. A wolf might come in. It would eat you first. (Roberts 1989)
- (20) Topicalization / Left Dislocation:
- a. As smart as Linus is, Gabby is even smarter.
 - b. Megan, I like (her).

Thus, I take the non-conditional comparative correlative in (1) above to relate *two cases*,⁶ one contributed by the protasis and the other by the apodosis. Each case features two heights and their differential and the two cases are related in terms of their respective differentials, which are equated. The conditional comparative correlatives in (4) and (5) are just a generalization of this basic pattern. They do not involve a single pair of cases related by means of their respective differentials, but involve multiple pairs of such cases. What is characterized in the literature (e.g., McCawley 1988, Beck 1997) as the *conditionality* of certain comparative correlatives is, under this analysis, just the fact that they correlate *sets* of cases.

The basic schema for the interpretation of comparative correlatives is provided in (21) below. The wh-morphology on the indefinite *cîț* contributes a maximality operator **max** with scope over the rest of the protasis ϕ . The **max** operator ensures that we take into consideration all cases, i.e., all sequences of objects, that satisfy the protasis. Each such case features a non-empty differential interval D , also contributed by the wh-indefinite – as indicated by the fact that D is a superscript on *cîț*. Then, the apodosis further elaborates on the set of cases contributed by the protasis: for each case / sequence, *atîț* anaphorically retrieves the differential D in that sequence – as indicated by the fact that D is a subscript on *atîț* – and relates it to a newly introduced differential D' that satisfies the rest of the apodosis ψ .

$$(21) \quad \begin{array}{l} cu \ cîț^D(\phi) \quad \quad \quad \text{—} \quad cu \ atîț_D^{D'}(\psi) \quad \rightsquigarrow \\ \mathbf{max}(D; D \neq \emptyset; \phi(D)) \quad \text{—} \quad D'; D' \neq \emptyset; \psi(D'); \mathcal{R}(D, D') \end{array}$$

⁶I use 'case' in the sense of Lewis (1975); (minimal) situations would work as well.

The value of the relational variable \mathcal{R} is contextually specified. I assume that the Cartesian product $D_{\gamma} \times D_{\gamma}$ is always contextually salient and can therefore be the default value assigned to \mathcal{R} – in which case we obtain the truth conditions in (7) above, where no particular constraints are placed on the relation between differentials. That is, if the value of \mathcal{R} is $D_{\gamma} \times D_{\gamma}$, the requirement $\mathcal{R}(D, D')$ is vacuously satisfied, so we can drop it and obtain truth-conditions of the form $\forall D(D \neq \emptyset \dots \rightarrow \exists D'(D' \neq \emptyset \dots))$.

The presuppositional particle *tot* (also) can require the relation \mathcal{R} to be the identity relation between intervals – in which case we obtain the stronger truth conditions in (8) above, which have the form $\forall D(D \neq \emptyset \dots \rightarrow \exists D'(D' \neq \emptyset \dots \wedge D = D'))$.

Finally, pragmatic factors can also force the value of \mathcal{R} to be the identity relation, e.g., in (1) above. In this case, the requirement $D = D'$ is present in the truth conditions whether the particle *tot* is present or not, as shown in (2) above.

The protasis and the apodosis of comparative correlatives are dynamically conjoined. This is a reflection of the fact that they are topic-comment structures: the protasis sets up a differential topic D that the apodosis comments on. Dynamic conjunction ensures that the cases introduced in the protasis, including the D -differentials, are passed on to the apodosis.

The **max** operator always contributes the maximal set of cases satisfying the protasis and requires this set of cases to be non-empty. When the set of cases contributed by the protasis is a singleton set, like in (1) above, we have a non-conditional comparative correlative. Because we have a singleton set of cases, the truth conditions of such non-conditional comparative correlatives are basically existential and have the form $\exists D(D \neq \emptyset \dots \wedge \exists D'(D' \neq \emptyset \dots \wedge D = D'))$ – hence the formula in (2) above. When the set of cases contributed by the protasis is a non-singleton set, like in (4) and (5) above, we have a conditional comparative correlative. Because the set of cases is not a singleton, the truth conditions of such conditional comparative correlatives are universal and have the form $\forall D(D \neq \emptyset \dots \rightarrow \exists D'(D' \neq \emptyset \dots))$ – hence the formulas in (7) and (8) above.

Besides accounting for the variable conditionality of comparative correlatives, dynamically conjoining the protasis and the apodosis – in that order – enables us to capture the fact that we can have anaphora from the latter to the former, e.g., *The more tyrannical α^x man is, the less his_x son likes him*, but not vice versa, e.g., *The more reckless his_x son is, the less α^x man likes him*.

Given a suitable framework, both individual-based correlatives like (15) and (16) above and run-of-the-mill conditionals like *If a wolf came in, it would / might eat you first* (example based on Roberts 1989) can be analyzed along the same lines as (21) above. The only difference is that the protasis and the apodosis are correlated in terms of the individuals or possible scenarios / possible worlds that they characterize instead of differential intervals; see Brasoveanu (2008) for an analysis of individual-based correlatives and Brasoveanu (2007) (building on proposals in Stone (1999) and Bittner (2001)) for an analysis of conditionals along these lines.

Because modalized conditionals correlate cases by means of the possible scenarios / possible worlds that they evoke and that can never be exhaustively char-

acterized up to singletonness (the set containing the actual world is the only modal singleton set that enters natural language interpretation – and it enters interpretation only deictically, never ‘descriptively’), conditionality is always part of the meaning of modalized conditionals, but not of individual-based or degree-based correlatives.

3. Anaphora to Differentials in Compositional Dynamic Semantics

The formal analysis of comparative correlatives presents us with two problems: (i) the variable conditionality problem and (ii) the compositionality problem that arises when there are no overt *than*-phrases, e.g., in (4) above.

We solve the conditionality problem by evaluating natural language expressions relative to – and letting them update – *sets* of sequences of objects (i.e., sets of cases in the sense of Lewis (1975)) instead of one sequence at a time.

We solve the compositionality problem by modeling such sequences of objects as *stacks* instead of total or partial variable assignments.⁷ The main motivation for using stacks is that we always add information to a stack (never overriding previous information) and we do this in an orderly manner, based on the particular position in the stack that the update targets. An important consequence for our analysis is that we can easily define a notion of stack ‘splitting’ (and re-concatenation), which enables us to ‘generate’ two independent variables out of the single indefinite *un avocat* (a lawyer) in (4) and compare their aggressiveness and efficiency.

Given the syntactically non-local, cross-clausal character of anaphora to differential intervals – which makes it similar to donkey anaphora –, the analysis is formalized in a dynamic semantics system. Following Muskens (1996), the system is couched in classical (many-sorted) type logic, which delivers Montague-style compositionality at sub-clausal level by the usual methods. In particular, we keep the three basic types introduced in section 1 above and add another basic type *s* that models stacks / sequences of individuals (variables of type *s*: *i, j, ...*).

We indicate the empty positions in a stack *i* by storing the dummy individual # (type *e*) there, which is the universal falsifier (# makes any lexical relation false).

$$(22) \quad \text{A stack of length } n: \begin{array}{|c|c|c|c|c|c|c|} \hline 0 & 1 & \dots & n-1 & n & n+1 & \dots \\ \hline \alpha_0 & \alpha_1 & \dots & \alpha_{n-1} & \# & \# & \dots \\ \hline \end{array}$$

The length of a stack *i*, abbreviated **lng**(*i*), is the leftmost position in which *i* stores an object different from the universal falsifier # plus 1, because the first position in a stack is the 0th position.⁸ For example, the stack in (22) above is of length *n*. We abbreviate the empty stack (of length 0) as *i*_∅ (that is, *i*_∅ := *ιj*.**lng**(*j*) = 0). From now on, we will omit the cells storing the universal falsifier # when we represent stacks.

⁷Following Bittner (2001), Nouwen (2003) and references therein.

⁸Stack length: **lng**(*i*) = 1 + *tn*. (*i*)_{*n*} ≠ # ∧ ∀*n'* > *n*((*i*)_{*n'*} = #) if ∃*n*((*i*)_{*n*} ≠ # ∧ ∀*n'* > *n*((*i*)_{*n'*} = #)); **lng**(*i*) = 0 if ∀*n*((*i*)_{*n*} = #); otherwise, **lng**(*i*) = #. The ‘otherwise’ case covers stacks of infinite length, for example, the stack storing the universal falsifier # at all odd-number positions 1, 3, 5, ... and objects different from # at the other positions.

Stack update, defined in (24), never overrides previously introduced anaphoric information. When we update the n^{th} position of a stack i , we obtain another stack j by shifting all the i -objects stored at positions greater than or equal to n by one position and introducing a new random object at position n . We can update a position n only if no previous position stores the universal falsifier $\#$.

(23) Projection functions: $(i)_n$ is the object stored at position n in stack i .

(24) $i[n]j := \forall m < n((j)_m = (i)_m \wedge (i)_m \neq \#) \wedge (j)_n \neq \# \wedge \forall m > n((j)_m = (i)_{m-1})$

For simplicity, I will only use discourse referents (drefs) for individuals u_0, u_1, \dots , which are functions of type se , and drefs for intervals (sets of degrees) $\Delta_0, \Delta_1, \dots$, which are functions of type $s(\gamma\tau)$. But we could just as easily add drefs for degrees $\delta_0, \delta_1, \dots$ of type $s\gamma$ and, in general, drefs for objects of arbitrary static types, e.g., for binary relations between intervals $\mathcal{R}_0, \mathcal{R}_1, \dots$ of type $s((\gamma\tau)((\gamma\tau)t))$. Drefs are type-driven projection functions over stacks. For example, $u_n := \lambda i_s.((i)_e)_n$: for any stack i , the individual $u_n(i)$ – or ui for short – is the n^{th} individual, i.e., the n^{th} object of type e stored in i . Similarly, $\Delta_n := \lambda i_s.((i)_{\gamma\tau})_n$: for any stack i , the interval $\Delta_n(i)$ – or $\Delta_n i$ for short – is the n^{th} interval, i.e., the n^{th} object of type $\gamma\tau$ stored in i . For any stack i and any static type τ that we have drefs for (in our case, $\tau \in \{e, \gamma\tau\}$), the function $(i)_\tau$ retrieves the sub-stack of i that stores the objects of type τ .⁹

We define the introduction of new drefs for individuals and intervals as shown in (25) below. All stacks i consist of the individual sub-stack $(i)_e$ and the interval sub-stack $(i)_{\gamma\tau}$. When we update a stack i with a dref u_n , we introduce a new random individual at the n^{th} position of $(i)_e$, while $(i)_{\gamma\tau}$ remains unchanged. When we update a stack i with a dref Δ_n , we introduce a new random interval at the n^{th} position of $(i)_{\gamma\tau}$, while $(i)_e$ remains unchanged.

(25) $i[u_n]j := (i)_e[n](j)_e \wedge (i)_{\gamma\tau} = (j)_{\gamma\tau}$ and $i[\Delta_n]j := (i)_{\gamma\tau}[n](j)_{\gamma\tau} \wedge (i)_e = (j)_e$

A set of five axioms ensures that the entities of type s actually behave as stacks.¹⁰

To solve the variable conditionality problem, we work with *plural* information states I, J, \dots , which are *sets* of stacks, i.e., terms of type st (van den Berg (1996) was the first to propose that natural language expressions should be evaluated relative to – and should dynamically update – sets of sequences of objects). Plural info states can be represented as matrices with stacks / sequences as rows:

⁹The initial position of type τ : $\mathbf{ini}_\tau(i) = \iota n.(i)_n \in D_\tau \wedge (i)_n \neq \# \wedge \neg \exists n' < n((i)_{n'} \in D_\tau \wedge (i)_{n'} \neq \#)$ if $\exists n((i)_n \in D_\tau \wedge (i)_n \neq \#)$; otherwise, $\mathbf{ini}_\tau(i) = \#$. The final position of type τ : $\mathbf{fin}_\tau(i) = \iota n.(i)_n \in D_\tau \wedge (i)_n \neq \# \wedge \neg \exists n' > n((i)_{n'} \in D_\tau \wedge (i)_{n'} \neq \#)$ if $\exists n((i)_n \in D_\tau \wedge (i)_n \neq \#)$; otherwise, $\mathbf{fin}_\tau(i) = \#$. The sub-stack of type τ : $(i)_\tau = i_\emptyset$ if $\mathbf{ini}_\tau(i) = \mathbf{fin}_\tau(i) = \#$; otherwise, $(i)_\tau = \iota j.\mathbf{lng}(j) = 1 + (\mathbf{fin}_\tau(i) - \mathbf{ini}_\tau(i)) \wedge \forall n < \mathbf{lng}(j)((j)_n = (i)_{n+\mathbf{ini}_\tau(i)})$.

¹⁰**Ax1** (stack identity in terms of projection functions) – $\forall i_s \forall i'_s (\forall n((i)_n = (i')_n) \rightarrow i = i')$; **Ax2** (stacks have finite length) – $\forall i_s (\exists n(\mathbf{lng}(i) = n))$; **Ax3** (the empty stack exists) – $\exists i_s (\mathbf{lng}(i) = 0)$; **Ax4** (only drefs for individuals and intervals; individuals before intervals and not interspersed with $\#$) – $\forall i_s \forall n((i)_n \in D_e \vee (i)_n \in D_{\gamma\tau}) \wedge \forall i'_s \forall n \forall n'((i)_n \in D_e \setminus \{\#\} \wedge (i)_{n'} \in D_{\gamma\tau} \rightarrow n < n') \wedge \forall i_s \forall n \forall n'((i)_n \neq \# \wedge (i)_{n'} \neq \# \rightarrow \forall m(n < m < n' \rightarrow (i)_m \neq \#)$; **Ax5** (enough stacks) – $\forall i_s \forall n \forall x_e (n \leq \mathbf{lng}((i)_e) \rightarrow \exists j_s (i[u_n]j \wedge u_n j = x)) \wedge \forall i_s \forall n \forall D_{\gamma\tau} (n \leq \mathbf{lng}((i)_{\gamma\tau}) \rightarrow \exists j_s (i[\Delta_n]j \wedge \Delta_n j = D))$.

$$(26) \quad \text{Plural Info State } I: \begin{array}{c} i_1 \\ i_2 \\ i_3 \end{array} \begin{array}{|c|c|c|c|c|} \hline u_0 & u_1 & \Delta_0 & \Delta_1 & \Delta_2 \\ \hline x_1 & y_1 & D_1 & D'_1 & D''_1 \\ \hline x_2 & y_2 & D_2 & D'_2 & D''_2 \\ \hline x_3 & y_3 & D_3 & D'_3 & D''_3 \\ \hline \end{array}$$

Such matrices are two-dimensional and enable us to encode two kinds of information: *values* and *structure*. Values are the sets of objects stored in the columns of the matrix, e.g., the dref u_0 stores a set of individuals $\{x_1, x_2, x_3\}$ relative to the plural info state I in (26) above because u_0 is assigned an individual by each stack / row in I . Structure, that is, dependencies / relations between sets of values, is encoded in the rows of the matrix: for each stack / row in I , the individual assigned to u_0 by that stack is correlated with the individual assigned to u_1 by the same stack (and, also, with the intervals assigned to Δ_0 , Δ_1 and Δ_2). Structure is what enables us to analyze comparative correlatives (and correlatives in general) as introducing sets of cases and simultaneously elaborating on each of them.

A sentence is interpreted as a Discourse Representation Structure (DRS), i.e. as a relation of type $(st)((st)t)$ between an input state I_{st} and an output state J_{st} . DRSs $\mathbb{D}, \mathbb{D}', \dots$ are pairs of the form **[new drefs|conditions]** and they relate an input state I and an output state J iff J differs from I at most with respect to the new drefs and J satisfies all the conditions. For example, $[u_0, \Delta_0 | lw \{u_0\}, tall \{u_0, \Delta_0\}] := \lambda I_{st}. \lambda J_{st}. I[u_0, \Delta_0] J \wedge lw \{u_0\} J \wedge tall \{u_0, \Delta_0\} J$ – that is, J differs from I at most with respect to the individuals assigned to u_0 and the intervals assigned to Δ_0 and all the u_0 -individuals stored in J are lawyers and they are Δ_0 -tall.

$$(27) \quad \mathbf{[new drefs|conditions]} := \lambda I_{st}. \lambda J_{st}. I[\mathbf{new drefs}] J \wedge \mathbf{conditions} J$$

Tests are DRSs that do not introduce new drefs: **[conditions]** $:= \lambda I_{st}. \lambda J_{st}. I = J \wedge \mathbf{conditions} J$. For example, $[tall \{u_0, \Delta_0\}] := \lambda I_{st}. \lambda J_{st}. I = J \wedge tall \{u_0, \Delta_0\} J$.

Conditions are sets of info states (sets of sets of stacks), i.e., they are terms of type $(st)t$. Moreover, conditions encoding lexical relations are distributive relative to the plural info states they accept: as shown in (28), (29) and (30) below, they accept an info state I iff they accept each stack $i \in I$. The first conjunct $I \neq \emptyset$ rules out the case in which the second conjunct $\forall i_s \in I(\dots)$ is vacuously satisfied.

$$(28) \quad lw \{u_n\} := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I(lw(u_n i))$$

$$(29) \quad eat \{u_m, u_n\} := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I(eat(u_m i, u_n i))$$

$$(30) \quad tall \{u_m, \Delta_n\} := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I(\Delta_n i = \{d_\gamma : tall(u_m i, d)\})$$

In the spirit of the Montagovian brace convention, the curly braces indicate that the static non-logical constants lw (type et), eat (type $e(et)$) and $tall$ (type $\gamma(et)$) apply to their dref arguments only after an index of evaluation (i.e., a stack i) is supplied.

A form of maximization over degrees is built into conditions like $tall \{u_m, \Delta_n\}$: for each stack i , the interval $\Delta_n i$ contains all the degrees d such that the individual $u_m i$ is at least d -tall. This will enable us to follow Beck (1997) and interpret the comparative morphology as (basically) relating two definite descriptions over degrees (see 43 below). I have chosen this formalization for simplicity and to facilitate

the comparison with Beck (1997) – but any analysis of comparatives that delivers a suitable interpretation for differentials would work as well.

We generalize new dref introduction to deal with plural info states and with the simultaneous introduction of an arbitrary finite sequence of drefs v_{m_1}, \dots, v_{m_n} (where v stands for u or Δ and m_1, \dots, m_n are natural numbers) as shown below.

- (31) $I[u_n]J := \forall i_s \in I(\exists j_s \in J(i[u_n]j)) \wedge \forall j_s \in J(\exists i_s \in I(i[u_n]j))$ and
 $I[\Delta_n]J := \forall i_s \in I(\exists j_s \in J(i[\Delta_n]j)) \wedge \forall j_s \in J(\exists i_s \in I(i[\Delta_n]j))$
- (32) Dynamic conjunction: $\mathbb{D}; \mathbb{D}' := \lambda I_{st}. \lambda J_{st}. \exists H_{st} (\mathbb{D}IH \wedge \mathbb{D}'HJ)$
- (33) Multiple dref introduction: $[v_{m_1}, \dots, v_{m_n}] := [v_{m_1}]; \dots; [v_{m_n}]$

We can now define the maximization operator **max** needed to capture the variable conditionality of comparative correlatives. As shown in (34) below, the result of updating an input state I with a DRS **max**(\mathbb{D}) is the maximal set J that we can get when we update I with the DRS \mathbb{D} : J stores all the stacks / cases that satisfy \mathbb{D} .

- (34) **max**(\mathbb{D}) := $\lambda I_{st}. \lambda J_{st}. DIJ \wedge \forall K_{st} (DIK \rightarrow K \subseteq J)$

Finally, we define an operator over stacks that removes the first individual in a stack i :¹¹ the stack i^{e-1} is the same as i except that the first position of type e has been dropped.¹² The operator is generalized to DRSs as shown in (35) below: a DRS $(\mathbb{D})^{e-1}$ updates an input state I by updating each stack $i \in I$ as if the first object of type e in that stack did not exist. The $(\dots)^{e-1}$ operator enables us to solve the compositionality problem posed by comparative correlatives without *than*-phrases.

- (35) $(\mathbb{D})^{e-1} := \lambda I_{st}. \lambda J_{st}. \exists \mathbb{R}_{s((st)t)} \neq \emptyset (I = \mathbf{Dom}(\mathbb{R}) \wedge J = \cup \mathbf{Ran}(\mathbb{R}) \wedge \forall k_s \forall L_{st} (\mathbb{R}kL \rightarrow \{((k)_e)_o\} = ((L)_e)_o \wedge \mathbb{D} \{k^{e-1}\} L^{e-1}))$ ¹³

The dynamic notion of truth is defined as usual, i.e., as existential ‘closure’.

- (36) Truth: a DRS $\mathbb{D}_{(st)((st)t)}$ is true relative to an info state I_{st} iff $\exists J_{st} (DIJ)$.

3.1. Non-Conditional Comparative Correlatives

Given the underlying type logic, we can define a Montague-style compositional interpretation procedure. This is largely determined by the types for the ‘saturated’ expressions, i.e., names and sentences, which we abbreviate as **e** and **t**. An extensional static framework identifies **e** with e and **t** with t . The translation of the noun *lawyer* is of type **et**, i.e., $et: lawyer \rightsquigarrow \lambda x_e. lw_{et}(x)$. The indefinite article *a* is of type **(et)((et)t)**, i.e., $(et)((et)t): a \rightsquigarrow \lambda X_{et}. \lambda X'_{et}. \exists x_e (X(x) \wedge X'(x))$.

¹¹This operator can be generalized to remove the first position storing a time, a world, an eventuality etc. – or to remove the first n positions storing individuals, times, worlds etc.

¹² $i^{e-1} = \iota j. (\mathbf{lng}((i)_e) = 0 \rightarrow \mathbf{lng}((j)_e) = 0) \wedge (\mathbf{lng}((i)_e) \neq 0 \rightarrow \mathbf{lng}((j)_e) = \mathbf{lng}((i)_e) - 1 \wedge \forall n < \mathbf{lng}((j)_e) ((j)_e)_n = ((i)_e)_{n+1}) \wedge (j)_n = (i)_n$.

¹³ $\mathbf{Dom}(\mathbb{R}) := \{k : \exists L(\mathbb{R}kL)\}$, $\mathbf{Ran}(\mathbb{R}) := \{L : \exists k(\mathbb{R}kL)\}$, $((L)_e)_o := \{((l)_e)_o : l \in L\}$ and $L^{e-1} := \{l^{e-1} : l \in L\}$.

In the present dynamic system, we let $\mathbf{t} := (st)((st)t)$, i.e., a sentence is interpreted as a DRS, and $\mathbf{e} := se$, i.e., a name is interpreted as a dref. The denotation of the noun *lawyer* is still of type \mathbf{et} , as shown in (37) below. The indefinite article *a* is still of type $(\mathbf{et})(\mathbf{et}\mathbf{t})$: as shown in (38), it introduces the dref u_n with which it is superscripted and predicates the restrictor and nuclear scope properties P and P' of this dref. Pronouns are translated as the Montagovian lift of the dref they anaphorically retrieve (type $(\mathbf{et}\mathbf{t})$, as shown in (39). Anaphoric definite articles receive the expected translation in (40) (type $(\mathbf{et})(\mathbf{et}\mathbf{t})$). We abbreviate the type of drefs for intervals as $\mathbf{\Gamma} := s(\gamma\mathbf{t})$ and translate the adjective *tall* as shown in (41) (type $\mathbf{\Gamma}(\mathbf{et})$). P, P', \dots are variables of type \mathbf{et} over dynamic properties of individuals, u, u', \dots are variables of type \mathbf{e} and Δ, Δ', \dots are variables of type $\mathbf{\Gamma}$.

$$(37) \quad \textit{lawyer} \rightsquigarrow \lambda u_{\mathbf{e}}. [lw\{u\}], \text{ i.e., } \textit{lawyer} \rightsquigarrow \lambda u_{\mathbf{e}}. \lambda I_{st}. \lambda J_{st}. I = J \wedge lw\{u\}J$$

$$(38) \quad a^{u_n} \rightsquigarrow \lambda P_{\mathbf{et}}. \lambda P'_{\mathbf{et}}. [u_n]; P(u_n); P'(u_n)$$

$$(39) \quad he_{u_n} / pro_{u_n} \rightsquigarrow \lambda P_{\mathbf{et}}. P(u_n)$$

$$(40) \quad the_{u_n} \rightsquigarrow \lambda P_{\mathbf{et}}. \lambda P'_{\mathbf{et}}. P(u_n); P'(u_n)$$

$$(41) \quad \textit{tall} \rightsquigarrow \lambda \Delta_{\mathbf{\Gamma}}. \lambda u_{\mathbf{e}}. [tall\{u, \Delta\}]$$

I take Romanian comparative correlatives to have roughly the same syntactic structure as the one proposed in Beck (1997). More precisely, I take them to be IP-adjunction structures¹⁴ in which the topic clause (the protasis) is an IP adjoined to the comment clause (the apodosis), which is also an IP – as shown in (42) below.

$$(42) \quad [_{IP} [_{IP} [Cu \hat{c}\hat{i}t \textit{mai}]^{\Delta} [_{I} e \Delta\text{-}\hat{i}nalt \textit{fratele}] [_{CP} [dec\hat{i}t]^{\Delta'} \textit{sora e } \Delta'\text{-}\hat{i}nalt\hat{a}]]] \\ [_{IP} [(tot) \textit{cu at\hat{i}t \textit{mai}}]^{\Delta''} [_{I} e \Delta''\text{-}\hat{i}nalt \textit{tat\hat{a}l}] [_{CP} [dec\hat{i}t]^{\Delta'''} \textit{mama e } \Delta'''\text{-}\hat{i}nalt\hat{a}]]]]$$

The wh-differential *cu c\hat{i}t* and the differential anaphor *cu at\hat{i}t* are in the Spec-IP of their respective clauses – since they are in complementary distribution with the subjects. I assume that the syntax of comparatives involves two instances of quantifying-in per comparative (see Schwarzschild (2008) and references therein for more discussion), adding up to four in a comparative correlative. The superscripted variables $\Delta, \Delta', \Delta''$ and Δ''' indicate where quantifying-in – hence, λ -abstraction – occurs. Finally, I follow Beck (1997: 248 et seqq) and take the comparative morpheme *mai* (more) to form a constituent with the differentials.

The meanings for *mai*, *cu c\hat{i}t* and *cu at\hat{i}t* in (43), (44) and (45) below are based on the syntax in (42) above. The proposed semantics for comparative correlatives, however, is compatible with other syntactic analyses (and frameworks). The wh-phrase *dec\hat{i}t* and the copula *e* are vacuous – they denote identity functions of the appropriate types: $(\mathbf{\Gamma}\mathbf{t})(\mathbf{\Gamma}\mathbf{t})$ for *dec\hat{i}t* and \mathbf{tt} (or $(\mathbf{et})(\mathbf{et})$) for the copula. P, P', \dots are variables of type $\mathbf{\Gamma}\mathbf{t}$ over dynamic properties of intervals and R is a variable of type $\mathbf{\Gamma}(\mathbf{\Gamma}(\mathbf{\Gamma}\mathbf{t}))$ over ternary dynamic relations between intervals.

$$(43) \quad \textit{mai} \rightsquigarrow \lambda \Delta_{\mathbf{\Gamma}}. \lambda \Delta'_{\mathbf{\Gamma}}. \lambda \Delta''_{\mathbf{\Gamma}}. [\Delta'' \neq \emptyset, \Delta'' = \Delta \setminus \Delta']^{15}$$

¹⁴Not CP-adjunction structures, because they can occur in attitude reports embedded under an overt complementizer, e.g., embedded under *Irina crede c\hat{a}...* (Irina believes that ...).

¹⁵ $\Delta'' \neq \emptyset := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I (\Delta'' i \neq \emptyset)$ and $\Delta'' = \Delta \setminus \Delta' := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I (\Delta'' i = \Delta i \setminus \Delta' i)$.

- (44) $cu\ c\hat{i}t^{\Delta_m, \Delta_{m'}, \Delta_{m''}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $\mathbf{max}([\Delta_m]; P(\Delta_m); [\Delta_{m'}]; P'(\Delta_{m'}); [\Delta_{m''}]; R(\Delta_m)(\Delta_{m'})(\Delta_{m''}))$
- (45) $cu\ at\hat{i}t_{\mathcal{R}, \Delta_{m''}}^{\Delta_n, \Delta_{n'}, \Delta_{n''}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $[\Delta_n]; P(\Delta_n); [\Delta_{n'}]; P'(\Delta_{n'}); [\Delta_{n''}]; R(\Delta_n)(\Delta_{n'})(\Delta_{n''}); [\mathcal{R}(\Delta_{m''}, \Delta_{n''})]$

The **max** operator in the meaning of $cu\ c\hat{i}t$ is contributed by its wh-morphology. In contrast, $cu\ 2\ cm$ or the null differential – translated in (46) and (47) below – do not have such an operator. The null differential, which I assume occurs in ‘bare’ comparatives like *Gabby is taller than Linus*, is, basically, just existential closure over the differential argument of the comparative morpheme *mai* (more).

- (46) $cu\ 2\ cm^{\Delta_m, \Delta_{m'}, \Delta_{m''}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $[\Delta_m]; P(\Delta_m); [\Delta_{m'}]; P'(\Delta_{m'}); [\Delta_{m''}]; R(\Delta_m)(\Delta_{m'})(\Delta_{m''}); [\mathbf{2cm}\{\Delta_{m''}\}]^{16}$
- (47) $null-differential^{\Delta_m, \Delta_{m'}, \Delta_{m''}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $[\Delta_m]; P(\Delta_m); [\Delta_{m'}]; P'(\Delta_{m'}); [\Delta_{m''}]; R(\Delta_m)(\Delta_{m'})(\Delta_{m''})$

We conjoin the matrix and *than* clauses of a comparative – in that order – because we have anaphora from the latter to the former, e.g., A^{u_0} *man was 2 cm taller than his_{u_0} son*, but not vice versa, e.g., $\#His_{u_0}$ *son was 2 cm taller than a^{u_0} man*.

The meaning of the anaphoric differential $at\hat{i}t$ in (45) above contains a relational dref \mathcal{R} , whose value is contextually supplied. Its default value is the Cartesian product over the domain of intervals, but the particle *tot* or various pragmatic factors can require it to be the identity relation over intervals (see the discussion in section 2 above). For simplicity, however, we will ignore this relational dref and work with the two translations in (48) and (49) below: (48) corresponds to the ‘Cartesian product \mathcal{R} ’ case, which reduces to no anaphora to differentials whatsoever; (49) corresponds to the ‘identity \mathcal{R} ’ case, which reduces to direct anaphora to a previously introduced differential interval.

- (48) $cu\ at\hat{i}t^{\Delta_n, \Delta_{n'}, \Delta_{n''}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $[\Delta_n]; P(\Delta_n); [\Delta_{n'}]; P'(\Delta_{n'}); [\Delta_{n''}]; R(\Delta_n)(\Delta_{n'})(\Delta_{n''})$
- (49) $(tot)\ cu\ at\hat{i}t_{\Delta_{m''}}^{\Delta_n, \Delta_{n'}} \rightsquigarrow \lambda R_{\Gamma(\Gamma t)}. \lambda P_{\Gamma t}. \lambda P'_{\Gamma t}.$
 $[\Delta_n]; P(\Delta_n); [\Delta_{n'}]; P'(\Delta_{n'}); R(\Delta_n)(\Delta_{n'})(\Delta_{m''})$

The comparative correlative in (1) above is indexed as shown in (50) below and it is compositionally translated as shown in (51). Given the definition of dynamic truth in (36), the translation derives the intuitively correct truth conditions in (2) above. The update proceeds as shown in (52): the anaphoric definites retrieve the contextually specified values for the drefs u_0 , u_1 , u_2 and u_3 . The differential $cu\ c\hat{i}t$ in the topic clause / protasis updates the context by introducing an interval dref Δ_2 that stores the difference between the brother’s height Δ_0 and the sister’s height Δ_1 . The comment clause / apodosis receives a parallel interpretation, except

¹⁶ $\mathbf{2cm}\{\Delta_{m''}\} := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I(\mathbf{2cm}(\Delta_{m''}i))$ and $\mathbf{2cm}$ is the set of 2 cm intervals (type (γt)).

that *cu atît* anaphorically retrieves the differential interval Δ_2 and equates it with the difference between the father's height Δ_3 and the mother's height Δ_4 .¹⁷

- (50) Cu cît ^{$\Delta_0, \Delta_1, \Delta_2$} e mai înalt fratele _{u_0} decât sora _{u_1} ,
(tot) cu atît ^{Δ_3, Δ_4} e mai înalt tatăl _{u_2} decât mama _{u_3} .
- (51) $\mathbf{max}([\Delta_0 | tall \{u_0, \Delta_0\}]; [\Delta_1 | tall \{u_1, \Delta_1\}]; [\Delta_2 | \Delta_2 \neq \emptyset, \Delta_2 = \Delta_0 \setminus \Delta_1]);$
 $[\Delta_3 | tall \{u_2, \Delta_3\}]; [\Delta_4 | tall \{u_3, \Delta_4\}]; [\Delta_2 \neq \emptyset, \Delta_2 = \Delta_3 \setminus \Delta_4]$
- (52)

u_0	u_1	u_2	u_3	topic clause / protasis						u_0	u_1	u_2	u_3	Δ_0	Δ_1	Δ_2
<i>bro</i>	<i>sis</i>	<i>fa</i>	<i>mo</i>							<i>bro</i>	<i>sis</i>	<i>fa</i>	<i>mo</i>	<i>D_{bro}</i>	<i>D_{sis}</i>	<i>D_{diff}</i>
comment clause / apodosis				u_0	u_1	u_2	u_3	Δ_0	Δ_1	Δ_2	Δ_3	Δ_4				
				<i>bro</i>	<i>sis</i>	<i>fa</i>	<i>mo</i>	<i>D_{bro}</i>	<i>D_{sis}</i>	<i>D_{diff}</i>	<i>D_{fa}</i>	<i>D_{mo}</i>				

3.2. Conditional Comparative Correlatives

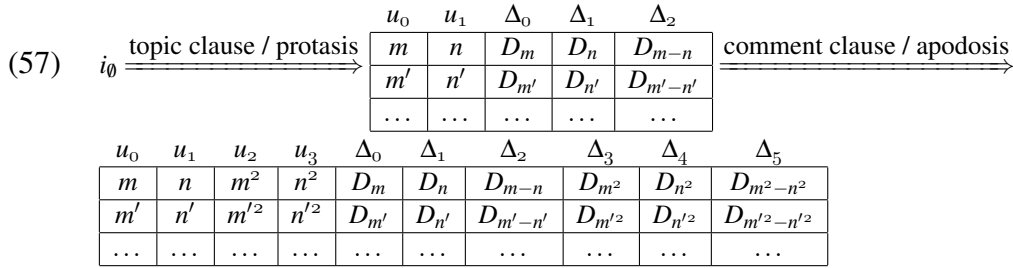
The comparative correlative with overt *than*-phrases in (5) above is indexed as shown in (55) below and it is compositionally translated as shown in (56) (the translation of the infelicitous variant with the particle *tot* present can be obtained if we use 49 above instead of 48). Once again, the translation derives the intuitively correct truth conditions. The only difference between the truth conditions derived based on (56) and the ones in (8) above (besides the final conjunct $D = D'$ in 8, which is due to the particle *tot*) is that the former have an additional existential requirement that there is at least one case satisfying the protasis, i.e., that there are at least two natural numbers such that one of them is greater than the other. This additional requirement is, in fact, part of the meaning of comparative correlatives (see Beck 1997: 253) and of correlative and quantificational structures in general, which are felicitous only if their domains are non-empty.

- (53) $the^{u_n} \rightsquigarrow \lambda P_{et}. \lambda P'_{et}. \mathbf{unique}^{u_n}(P(u_n)); P'(u_n)$ ¹⁸
- (54) $altul_{u_m}^{u_n} / an^{u_n} other_{u_m} nat. number \rightsquigarrow \lambda P_{et}. [u_n | u_n \neq u_m, u_n \in \mathbb{N}]; P(u_n)$ ¹⁹
- (55) Cu cît ^{$\Delta_0, \Delta_1, \Delta_2$} e un ^{$u_0$} număr natural mai mare decât altul ^{u_1} ,
cu atît ^{$\Delta_3, \Delta_4, \Delta_5$} e pătratul ^{$u_2$} lui ^{$u_0$} mai mare decât pătratul ^{u_3} celui ^{u_1} alt.
- (56) $\mathbf{max}([\Delta_0, u_0 | u_0 \in \mathbb{N}, great \{u_0, \Delta_0\}]; [\Delta_1, u_1 | u_1 \neq u_0, u_1 \in \mathbb{N}, great \{u_1, \Delta_1\}];$
 $[\Delta_2 | \Delta_2 \neq \emptyset, \Delta_2 = \Delta_0 \setminus \Delta_1]); \mathbf{unique}^{u_2}([square \{u_2, u_0\}]); [\Delta_3 | great \{u_2, \Delta_3\}];$
 $\mathbf{unique}^{u_3}([square \{u_3, u_1\}]); [\Delta_4 | great \{u_3, \Delta_4\}]; [\Delta_5 | \Delta_5 \neq \emptyset, \Delta_5 = \Delta_3 \setminus \Delta_4]$

¹⁷The equative correlative in (3) is analyzed in terms of anaphora to the interval obtained by subtracting the (contextual) standard of beauty from Irina's maximal degree of beauty. This interval is correlated with the one obtained by subtracting the contextual standard of smartness from Irina's maximal degree of smartness. Equative correlatives that do not involve reference to contextual standards, e.g., the correlative mentioned in footnote 1 above, correlate the intervals between the 0-points on the relevant scales and the maximal degrees under consideration.

¹⁸ $\mathbf{unique}^{u_n}(\mathbb{D}) := \lambda I_{st}. \lambda J_{st}. \exists \mathbb{R}_{s(st)} \neq \emptyset (I = \mathbf{Dom}(\mathbb{R}) \wedge J = \cup \mathbf{Ran}(\mathbb{R}) \wedge \forall k_s \forall L_{st} (\mathbb{R}kL \rightarrow ([u_n]; \mathbb{D}) \{k\} L \wedge |u_n L| = 1 \wedge \forall L'_{st} (([u_n]; \mathbb{D}) \{k\} L' \rightarrow u_n L' = u_n L))$, where $u_n L := \{u_n l : l \in L\}$.

¹⁹ $u_n \neq u_m := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I (u_n i \neq u_m i)$ and $u_n \in \mathbb{N} := \lambda I_{st}. I \neq \emptyset \wedge \forall i_s \in I (u_n i \in \mathbb{N})$.



We finally turn to the comparative correlative without overt *than*-phrases in (4) above. Compositionally deriving the meaning of such correlatives is problematic because, just as the other comparative correlatives we considered, they involve a comparison in both the protasis and the apodosis, but the second term of comparison is missing in both cases – and, as Beck (1997) notes, it is intuitively provided by reusing the first term of comparison, i.e., by using the two overt clauses twice.

Formally, the problem with this ‘double use’ idea is that both static and dynamic approaches find it difficult to introduce the variable contributed by the protasis indefinite *a lawyer* twice, make the two occurrences of the variable semantically independent and anaphorically retrieve them as two independent variables with the same pronoun in the apodosis, which, to this end, is also used twice.

This is ultimately the reason for the non-compositionality of the account in Beck (1997), where we need to postulate the existence of pairs of individual / temporal / world variables (as the case may be; cf. p. 248:42 and p. 252:49a,b) in the meaning of the correlative particles, then we need to coindex these variables with a covert adverb of quantification that takes scope over the entire correlative structure and, finally, we need these variables to satisfy two predicates obtained by covert λ -abstraction over the indefinite in the protasis of (4) and the pronoun in the apodosis of (4). That is, most of the work in deriving the correct truth conditions is done not by the meanings of the lexical items and the syntax-based way they are put together, but by postulating covert pairs of variables and a variety of covert operators inserted in particular structural positions that are suitably coindexed with both the covert variables and the overtly present indefinites and pronouns.

I take the basic intuition behind the ‘double use’ proposal to be that the protasis clause is used not once, but twice to update the context, so the context changes brought about by this clause are recorded twice (the same goes for the apodosis clause) – and I will use the $(\dots)^{e-1}$ operator to formalize this intuition. The ‘double update’ analysis is just this: the first update is executed as usual, while the second update with the same clause is executed in the scope of a $(\dots)^{e-1}$ operator. The operator temporarily removes the first e -position (contributed by the first update) and reattaches it after the second update is executed. The idea is that we treat the second execution of the update as if it was the first one, i.e., we interpret the subordinate *than*-clause as if it temporarily was the matrix clause. This is analogous to the way we manipulate world variables when we interpret clauses embedded under modal operators: the world variables contributed by the modal operators are treated as if they were the actual world and, given this temporary assumption, we interpret

the embedded clauses as if they were matrix clauses.²⁰

The two $(\dots)^{e-1}$ operators (one in the protasis, the other in the apodosis) are contributed by the two correlated differentials. Their translations, related in a reflexivization-like way to the translations in (44) and (48) above, are provided in (58) and (59) below. The proposal that, cross-linguistically, the presence vs. absence of *than*-phrases in comparative correlatives is licensed by the lexical properties of the correlated differentials is supported by the fact that Russian has two such pairs of differentials, *naskol'ko... nastol'ko...* and *čem... tem...*, the first of which is used for comparative correlatives with *than*-phrases, while the second is used for comparative correlatives without *than*-phrases.²¹

$$(58) \quad cu \hat{c}\hat{i}\hat{t}^{\Delta_m, \Delta_{m'}, \Delta_{m''}, e-1} \rightsquigarrow \lambda R_{\Gamma}(\Gamma(\Gamma t)) . \lambda P_{\Gamma t} . \\ \mathbf{max}([\Delta_m]; P(\Delta_m); [\Delta_{m'}]; (P(\Delta_{m'}))^{e-1}; [\Delta_{m''}]; R(\Delta_m)(\Delta_{m'})(\Delta_{m''}))$$

$$(59) \quad cu \hat{a}\hat{t}\hat{i}\hat{t}^{\Delta_n, \Delta_{n'}, \Delta_{n''}, e-1} \rightsquigarrow \lambda R_{\Gamma}(\Gamma(\Gamma t)) . \lambda P_{\Gamma t} . \\ [\Delta_n]; P(\Delta_n); [\Delta_{n'}]; (P(\Delta_{n'}))^{e-1}; [\Delta_{n''}]; R(\Delta_n)(\Delta_{n'})(\Delta_{n''})$$

The comparative correlative in (4) is indexed as shown in (60) below and it is compositionally translated as shown in (61). The translation derives the intuitively correct truth conditions in (7) above – with the addition of the existential requirement that there are at least two lawyers with different degrees of aggressiveness.

$$(60) \quad Cu \hat{c}\hat{i}\hat{t}^{\Delta_0, \Delta_1, \Delta_2, e-1} e \text{ un}^{u_0} \text{ avocat mai agresiv, cu } \hat{a}\hat{t}\hat{i}\hat{t}^{\Delta_3, \Delta_4, \Delta_5, e-1} e \text{ p r o }_{u_0} \text{ mai eficient.}$$

$$(61) \quad \mathbf{max}([\Delta_0, u_0 | lw \{u_0\}, aggr \{u_0, \Delta_0\}]; [\Delta_1]; ([u_0 | lw \{u_0\}, aggr \{u_0, \Delta_1\}])^{e-1}; \\ [\Delta_2 | \Delta_2 \neq \emptyset, \Delta_2 = \Delta_0 \setminus \Delta_1]); \\ [\Delta_3 | eff \{u_0, \Delta_3\}]; [\Delta_4]; (eff \{u_0, \Delta_4\})^{e-1}; [\Delta_5 | \Delta_5 \neq \emptyset, \Delta_5 = \Delta_3 \setminus \Delta_4]$$

$$(62) \quad i_{\emptyset} \xrightarrow{\text{topic clause / protasis}} \begin{array}{|c|c|c|c|c|} \hline & u_0 & u_1 & \Delta_0 & \Delta_1 & \Delta_2 \\ \hline x & y & D_x & D_y & D_{x-y} \\ \hline x' & y' & D_{x'} & D_{y'} & D_{x'-y'} \\ \hline \dots & \dots & \dots & \dots & \dots \\ \hline \end{array} \xrightarrow{\text{comment clause / apodosis}}$$

²⁰Another way to formalize the ‘double update’ analysis is to follow Bittner (2001) and model contexts not as single stacks (or sets thereof), but as pairs of stacks, one storing the foregrounded / topical drefs, the other one storing the backgrounded drefs. The ‘double use’ of a clause in comparative correlatives like (4) is, once again, double update: we update the topical stack the first time around and, then, we reuse the same protasis clause to update the background stack. We thus obtain two independent drefs (a foregrounded and a backgrounded one) that can be anaphorically retrieved by the apodosis pronoun in a parallel ‘double update’ way.

²¹The Russian counterpart of (1) is: *Naskol'ko brat vyše čem sestra, nastol'ko otec vyše čem mat'* (in-how-much brother.Nom taller than sister.Nom, insomuch father.Nom taller than mother.Nom). Expressing the two *than*-phrases by means of the Genitives *sestry* (sister.Gen) and *materi* (mother.Gen) is also possible. The Russian counterpart of (4) is: *Čem advokat agresivnee, tem on effektivnee*. (than lawyer.Nom more-aggressive that.Instr he more-effective). The data, however, is more complicated, since there are two Russian counterparts of (5), one of them with a *than*-phrase in the protasis, i.e., *Čem bol'she odno natural'noe čislo čem drugoe, tem bol'she (i) ego kvadrat*. (than bigger one natural number.Nom than other.Nom, that.Instr bigger (and/EMPH) its square.Nom) (expressing the *than*-phrase by means of the Genitive *drugogo* (other.Gen) is also possible) and the other one without any *than*-phrase, i.e., *Čem bol'she natural'noe čislo, tem bol'she (i) ego kvadrat*. (than bigger natural number.Nom, that.Instr bigger (and/EMPH) its square.Nom).

u_0	u_1	Δ_0	Δ_1	Δ_2	Δ_3	Δ_4	Δ_5
x	y	D_x	D_y	D_{x-y}	D'_x	D'_y	D'_{x-y}
x'	y'	$D_{x'}$	$D_{y'}$	$D_{x'-y'}$	$D'_{x'}$	$D'_{y'}$	$D'_{x'-y'}$
...

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