# Conditional Gradient Algorithms for Rank-One Matrix Approximations with a Sparsity Constraint 

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## Sparsity Constrained Rank-One Matrix Approximation $\equiv$ PCA

Principal Component Analysis solves

$$
\min \left\{\left\|A-x x^{T}\right\|_{F}^{2}:\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\} \Leftrightarrow \max \left\{x^{T} A x:\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\},\left(A \in \mathbb{S}_{+}^{n}\right)
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Sparse Principal Component Analysis solves

$$
\max \left\{x^{\top} A x:\|x\|_{2}=1,\|x\|_{0} \leq k, x \in \mathbf{R}^{n}\right\}, k \in[1, n] \text { sparsity }
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$\|x\|_{0}$ counts the number of nonzero entries of $x$

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## Difficulties:

(1) Maximizing a Convex objective.
(2) Hard Nonconvex Constraint $\|x\|_{0} \leq k$.

## Current Approaches:

(1) SDP Convex Relaxations [D'aspremont-El Ghaoui-Jordan-Lankcriet 07]
(2) Approximation/Modified formulations [Many....]

## Sparse PCA via Penalization/Relaxation/Approximation

The problem of interest is the difficult sparse PCA problem as is

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- Relaxed $I_{1}$-constrained PCA $\left(\|x\|_{1} \leq \sqrt{\|x\|_{0}}\|x\|_{2}, \forall x\right)$

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- Approximate-Penalized: Uses concave approximation of $\|x\|_{0}$

$$
\max \left\{x^{T} A x-s \varphi_{p}(\mid x \|):\|x\|_{2}=1\right\} \varphi_{p}(x) \simeq\|x\|_{0}, p \rightarrow 0^{+}
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- SDP-Convex Relaxation max $\left\{\operatorname{tr}(A X): \operatorname{tr}(X)=1, X \succeq 0,\|X\|_{1} \leq k\right\}$ Convex relaxations can be computationally expensive for very large problems and will not be discussed here.


## Quick Highlight of Simple Algorithms on "Modified Problems"

| Type | Iteration | Per-Iteration Complexity | References |
| :---: | :---: | :---: | :---: |
| $I_{1}$-constrained | $x_{i}^{j+1}=\frac{\operatorname{sgn}\left(\left(\left(A+\frac{\sigma}{2}\right) x^{j}\right)_{i}\right)\left(\left\|\left(\left(A+\frac{\sigma}{2}\right) x^{j}\right)_{i}\right\|-\lambda^{j}\right)_{+}}{\sqrt{\sum_{h}\left(\left\|\left(\left(A+\frac{\sigma}{2}\right) x^{j}\right)_{h}\right\|-\lambda^{j}\right)_{+}^{2}}}$ | $O\left(n^{2}\right), O(m n)$ | Witten et al. (2009) |
| $I_{1}$-constrained | $\begin{aligned} & x_{i}^{j+1}=\frac{\operatorname{sgn}\left(\left(A x^{j}\right)_{i}\right)\left(\left\|\left(A x^{j}\right)_{i}\right\|-s^{j}\right)_{+}}{\sqrt{\sum_{h}\left(\left\|\left(A x^{j}\right)_{h}\right\|-s^{j}\right)_{+}^{2}}} \text { where } \\ & s^{j} \text { is }(k+1) \text {-largest entry of vector }\left\|A x^{j}\right\| \end{aligned}$ | $O\left(n^{2}\right), O(m n)$ | Sigg-Buhman (2008) |
| ${ }^{1} 0$-penalized | $z^{j+1}=\frac{\sum_{i}\left[\operatorname{sgn}\left(\left(b_{i}^{T} z^{j}\right)^{2}-s\right)\right]_{+}\left(b_{i}^{T} z^{j}\right) b_{i}}{\left\\|\sum_{i}\left[\operatorname{sgn}\left(\left(b_{i}^{T} z^{j}\right)^{2}-s\right)\right]_{+}\left(b_{i}^{T} z^{j}\right) b_{i}\right\\|_{2}}$ | $O(m n)$ | Shen-Huang (2008), <br> Journee et al. (2010) |
| ${ }^{1} 0$-penalized | $x_{i}^{j+1}=\frac{\operatorname{sgn}\left(2\left(A x^{j}\right)_{i}\right)\left(\left\|2\left(A x^{j}\right)_{i}\right\|-s \varphi_{p}^{\prime}\left(\left\|x_{i}^{j}\right\|\right)\right)_{+}}{\sqrt{\sum_{h}\left(\left\|2\left(A x^{j}\right)_{h}\right\|-s \varphi_{p}^{\prime}\left(\left\|x_{h}^{j}\right\|\right)\right)_{+}^{2}}}$ | $O\left(n^{2}\right)$ | Sriperumbudur et al. (2010) |
| $I_{1}$-penalized | $\begin{aligned} y^{j+1} & =\underset{y}{\operatorname{argmin}}\left\{\sum_{i}\left\\|b_{i}-x^{j} y^{T} b_{i}\right\\|_{2}^{2}+\lambda\\|y\\|_{2}^{2}+s\\|y\\|_{1}\right\} \\ x^{j+1} & =\frac{\left(\sum_{i} b_{i} b_{i}^{T}\right) y^{j+1}}{\left\\|\left(\sum_{i} b_{i} b_{i}^{T}\right) y^{j+1}\right\\|_{2}} \end{aligned}$ |  | Zou et al. (2006) |
| $I_{1}$-penalized | $z^{j+1}=\frac{\sum_{i}\left(\left\|b_{i}^{T} z^{j}\right\|-s\right)+\operatorname{sgn}\left(b_{i}^{T} z^{j}\right) b_{i}}{\left\\|\sum_{i}\left(\left\|b_{i}^{T} z^{j}\right\|-s\right)+\operatorname{sgn}\left(b_{i}^{T} z^{j}\right) b_{i}\right\\|_{2}}$ | $O(m n)$ | Shen-Huang (2008), <br> Journee et al. (2010) |

Table : Cheap sparse PCA algorithms for modified problems.

## A Plethora of Models/Algorithms Revisited

All previous listed algorithms have been derived from various disparate approaches/motivations to solve modifications of SPCA:

- Nonsmooth reformulations
- Expectation Maximization
- Majoration-Mininimization techniques
- DC programming
- ... etc...

Q1: Are all these algorithms different? ...Any connection?

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\max \left\{x^{\top} A x:\|x\|_{2}=1,\|x\|_{0} \leq k, x \in \mathbf{R}^{n}\right\}
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Q2: Is is possible to derive a simple/cheap scheme to tackle directly the sparse PCA problem as is?

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## Answers

- All the previously listed algorithms are a particular realization of a "Father Algorithm": ConGradU (based on the well-known Conditional Gradient Algorithm)
- ConGradU CAN be applied directly to the original problem!


## The Conditional Gradient/Frank-Wolfe Algorithm

[Frank-Wolfe'56, Rubinov'64, Levitin-Polyak'66, Canon-Cullum' 68, Dunn'79,....]
\& Classic Conditional Gradient Algorithm solves

$$
\max \{F(x): x \in C\}
$$

- $F: \mathbf{R}^{n} \rightarrow \mathbf{R}$ is continuously differentiable
- $C$ is nonempty, convex compact subset of $\mathbb{R}^{n}$
via the following iteration for all $j \geq 0$ :

$$
x^{0} \in C, x^{j+1}=x^{j}+\alpha^{j}\left(p^{j}-x^{j}\right)
$$

with

$$
p^{j}=\operatorname{argmax}\left\{\left\langle x-x^{j}, \nabla F\left(x^{j}\right)\right\rangle: x \in C\right\}
$$

where $\alpha^{j} \in(0,1]$ is a stepsize (exact/or via line search).
© Here in SPCA :
$F$ is convex, possibly nonsmooth; (through equiv. reformulations)
$C$ is compact but nonconvex

## Maximizing a Convex function over a Compact Nonconvex set

ConGradU - Conditional Gradient with a Unit Step Size

$$
x^{0} \in C, x^{j+1} \in \operatorname{argmax}\left\{\left\langle x-x^{j}, F^{\prime}\left(x^{j}\right)\right\rangle: x \in C\right\}
$$

## Notes:

(1) Mangasarian (96) considered it for $C$ a polyhedral set.
(2) $F$ is not assumed to be differentiable and $F^{\prime}(x)$ is a subgradient of $F$ at $x$.
(3) The algorithm is useful when $\max \left\{\left\langle x-x^{j}, F^{\prime}\left(x^{j}\right)\right\rangle: x \in C\right\}$ is simple to solve

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## A Basic Convergence Result

(a) The sequence $F\left(x^{j}\right)$ is monotonically increasing and

$$
\lim _{j \rightarrow \infty} \gamma\left(x^{j}\right)=0, \text { where } \gamma(x):=\max \left\{\left\langle u-x, F^{\prime}(x)\right\rangle: u \in C\right\}
$$

(b) If $F$ is assumed continuously differentiable, then every limit point of the sequence $\left\{x^{j}\right\}$ converges to a stationary point.

## The Original $I_{0}$-constrained PCA via ConGradU

Applying ConGradU directly to

$$
\max \left\{x^{\top} A x:\|x\|_{2}=1,\|x\|_{0} \leq k, x \in \mathbf{R}^{n}\right\}
$$

results in the iteration

$$
x^{j+1}=\operatorname{argmax}\left\{x^{j T} A x:\|x\|_{2}=1,\|x\|_{0} \leq k\right\}, j=0,1, \ldots
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- Thus, the main step consists of maximizing a linear function on intersection of two nonconvex sets

$$
x \in C_{1} \cap C_{2} \text { with } C_{1}:=\left\{x:\|x\|_{2}=1\right\}, C_{2}:=\left\{x:\|x\|_{0} \leq k\right\}
$$

- It turns out that this problem is very simple!
- In fact, thanks to $C_{1}: x^{j+1}=\underset{x \in C_{1} \cap C_{2}}{\operatorname{argmin}}\left\|x-A^{T} x^{j}\right\|^{2}=P_{C_{1} \cap C_{2}}\left(A^{T} x^{j}\right) \ldots$ and...


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- Thanks to the "hard" constraint $C_{2}$...Projection on intersection "easy"...!

$$
P_{C_{1} \cap C_{2}}\left(A^{T} x^{j}\right) \equiv P_{C_{1}} \circ\left[P_{C_{2}}\left(A^{T} x^{j}\right)\right]
$$

## A Simple Key Result

A Simple Key Result Given $0 \neq a \in \mathbf{R}^{n}$,

$$
\begin{aligned}
& \max _{x}\left\{a^{T} x:\|x\|_{2}=1,\|x\|_{0} \leq k\right\}=\left\|T_{k}(a)\right\|_{2}, \text { with solution } x^{*}=\frac{T_{k}(a)}{\left\|T_{k}(a)\right\|_{2}} \\
& \qquad\left(T_{k}(a)\right)_{i}= \begin{cases}a_{i}, & \text { for } k \text { largest entries (in absolute values) of } a ; \\
0, & \text { otherwise. }\end{cases}
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$$

Definition $T_{k}: \mathbf{R}^{n} \rightarrow \mathbf{R}^{n}$ is the best $k$-sparse approximation of $a$

$$
T_{k}(a):=\operatorname{argmin}\left\{\|x-a\|_{2}^{2}:\|x\|_{0} \leq k\right\}
$$

Despite the nonconvex constraint, very easy to compute. In case $k$ largest entries are not uniquely defined, we select the smallest possible indices, with w.l.o.g, $a \in \mathbf{R}^{n}$ such $\left|a_{1}\right| \geq \ldots \geq\left|a_{n}\right|$.

Computing $T_{k}(\cdot)$ only requires determining the $k^{\text {th }}$ largest number of a vector of $n$ numbers which can be done in $O(n)$ time (Blum 73) and zeroing out the proper components in one more pass of the $n$ numbers.

## $I_{0}$-constrained PCA via ConGradU

The iteration for ConGradU results in

$$
x^{j+1}=\operatorname{argmax}\left\{x^{j T} A x:\|x\|_{2}=1,\|x\|_{0} \leq k\right\}=\frac{T_{k}\left(A x^{j}\right)}{\left\|T_{k}\left(A x^{j}\right)\right\|_{2}}, j=0, \ldots
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- Convergence: Since the objective is continuously differentiable, by previous result, we have here that every limit point of the sequence $\left\{x^{j}\right\}$ converges to a stationary point.
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- Convergence: Since the objective is continuously differentiable, by previous result, we have here that every limit point of the sequence $\left\{x^{j}\right\}$ converges to a stationary point.
- Complexity: $O(k n)$ or $O(m n)$.
- The original $I_{0}$-constrained problem can be solved using ConGradU with the same complexity as when applied to solving modified problems!
- Penalized/modified problems require tuning a tradeoff penalty parameter to get the desired sparsity. This can be computationally very expensive, and is not needed in our scheme.


## Back to Q1 - ....All via ConGradU

- All currently known cheap schemes are particular realization of ConGradU
- Novel Schemes can be derived via ConGradU

All we need is a simple toolbox...

## Answer to Q1: A Simple ToolBox

All previously listed algorithms are particular realizations of ConGradU.

- Proposition 1 Given $a \in \mathbf{R}^{n}, s>0$,

$$
\max _{\|x\|_{2}=1}\left\{\langle a, x\rangle^{2}-s\|x\|_{0}\right\}=\sum_{i=1}^{n}\left(a_{i}^{2}-s\right)_{+}, x_{i}^{*}=\frac{a_{i}\left[\operatorname{sgn}\left(a_{i}^{2}-s\right)\right]_{+}}{\sqrt{\sum_{j=1}^{n} a_{j}^{2}\left[\operatorname{sgn}\left(a_{j}^{2}-s\right)\right]_{+}}} .
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$$

- Proposition 2 For $a \in \mathbf{R}^{n}, w \in \mathbf{R}_{++}^{n}$, and $W=\operatorname{diag}(w)$

$$
\max _{\|x\|_{2} \leq 1}\left\{\langle a, x\rangle-\|W x\|_{1}\right\}=\left\|S_{w}(a)\right\|, x^{*}=S_{w}(a) /\left\|S_{w}(a)\right\|_{2}
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$S_{w}(a)=(|a|-w)+\operatorname{sgn}(a)$. (Soft Threshold)

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- Proposition 3 Given $a \in \mathbf{R}^{n}$, we have
$\max \left\{\langle a, x\rangle:\|x\|_{2} \leq 1,\|x\|_{1} \leq k, x \in \mathbf{R}^{n}\right\}=\min \left\{\lambda k+\left\|S_{\lambda e}(a)\right\|_{2}: \lambda \in \mathbb{R}_{+}\right\}$
Moreover, if $\lambda$ solves the one-dimensional dual, then an optimal solution

$$
x^{*}(\lambda)=S_{\lambda e}(a) /\left\|S_{\lambda e}(a)\right\|_{2},\left(e \equiv(1, \ldots, 1) \in \mathbf{R}^{n}\right)
$$

## Nonsmooth Convex Reformulations

D'aspremont et al. (08), Journee et al. (10)
$1_{0}$-penalized PCA problem: $\max \left\{x^{\top} A x-s\|x\|_{0}:\|x\|_{2} \leq 1, x \in \mathbf{R}^{n}\right\}$
Exploiting $A$ PSD $A:=B^{T} B$ with $B \in \mathbf{R}^{m \times n}$, yields

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$$

The objective is neither concave nor convex. Using the simple fact $\|B x\|_{2}^{2}=\max _{\|z\|_{2} \leq 1}\left\{\langle z, B x\rangle^{2}\right\}$, the problem is equivalent to

$$
\max _{\|x\|_{2} \leq 1} \max _{\|z\|_{2} \leq 1}\left\{\langle z, B x\rangle^{2}-s\|x\|_{0}\right\}=\max _{\|z\|_{2} \leq 1} \max _{\|x\|_{2} \leq 1}\left\{\left\langle B^{T} z, x\right\rangle^{2}-s\|x\|_{0}\right\} .
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Exploiting $A$ PSD $A:=B^{T} B$ with $B \in \mathbf{R}^{m \times n}$, yields

$$
\max \left\{\|B x\|_{2}^{2}-s\|x\|_{0}:\|x\|_{2} \leq 1, x \in \mathbf{R}^{n}\right\}
$$

The objective is neither concave nor convex. Using the simple fact
$\|B x\|_{2}^{2}=\max _{\|z\|_{2} \leq 1}\left\{\langle z, B x\rangle^{2}\right\}$, the problem is equivalent to

$$
\max _{\|x\|_{2} \leq 1} \max _{\|z\|_{2} \leq 1}\left\{\langle z, B x\rangle^{2}-s\|x\|_{0}\right\}=\max _{\|z\|_{2} \leq 1} \max _{\|x\|_{2} \leq 1}\left\{\left\langle B^{T} z, x\right\rangle^{2}-s\|x\|_{0}\right\} .
$$

Now, the inner minimization in $x$ can be solved (use P1):

$$
\max _{x \in \mathbf{R}^{n}}\left\{\|B x\|_{2}^{2}-s\|x\|_{0}:\|x\|_{2} \leq 1\right\}=\max _{z \in \mathbf{R}^{m}}\left\{\sum_{i=1}^{n}\left[\left\langle b_{i}, z\right\rangle^{2}-s\right]_{+}:\|z\|_{2} \leq 1\right\}
$$

where $b_{i} \in \mathbf{R}^{m}$ is the $i^{\text {th }}$ column of $B$.
Since the objective function $f(z):=\sum_{i}\left[\left\langle b_{i}, z\right\rangle^{2}-s\right]_{+}$is now clearly convex, we can apply ConGradU, recovering the alg. of Journee et al. (10).

## More Examples on NSO Reformulation

Similarly, for the $\iota_{1}$-penalized PCA problem one can show:

$$
\max \left\{x^{T} A x-s\|x\|_{1}:\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\}=\max _{z \in \mathbf{R}^{m}}\left\{\sum_{i=1}^{n}\left(\left|b_{i}^{T} z\right|-s\right)_{+}^{2}:\|z\|_{2} \leq 1\right\}
$$

We can now apply ConGradU to the convex objective $f(z)=\sum_{i}\left[\left|b_{i}^{T} z\right|-s\right]_{+}^{2}$, and for which our convergence results for the nonsmooth case hold true.

This recovers exactly the other algorithm of Journee et al. (2010).

## More Examples on NSO Reformulation

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## ConGradU is Very Flexible

 Tackling more general problems......
## A General Class of Problems

$$
\text { (G) } \quad \max _{x}\{f(x)+g(|x|): x \in C\}
$$

$f: \mathbf{R}^{n} \rightarrow \mathbf{R} \quad$ is convex,
$g: \mathbf{R}_{+}^{n} \rightarrow \mathbf{R}$ is convex differentiable and montonote decreasing $C \subseteq \mathbf{R}^{n} \quad$ is a compact set.

Here $|x|:=\left(\left|x_{1}\right|, \ldots,\left|x_{n}\right|\right)^{T}$; monotone decreasing means componentwise.

- Useful for handling penalized/approximate problems.
- Note: the composition $g(|x|)$ is not necessarily convex ...But after a simple transformation we can show that CondGradU can be applied to $(G)$, and produces the following simple scheme.


## A Simple Scheme for Solving (G)

(G) $\max _{x}\{f(x)+g(|x|): x \in C\}$

A-weighted $I_{1}$-norm maximization problem:

$$
x^{0} \in C, x^{j+1}=\operatorname{argmax}\left\{\left\langle a^{j}, x\right\rangle-\sum_{i} w_{i}^{j}\left|x_{i}\right|: x \in C\right\}, j=0, \ldots,
$$

where $w^{j}:=-g^{\prime}\left(\left|x^{j}\right|\right)>0$ and $a^{j}:=f^{\prime}\left(x^{j}\right) \in \mathbf{R}^{n}$.

## A Simple Scheme for Solving (G)

$$
\begin{equation*}
\max _{x}\{f(x)+g(|x|): x \in C\} \tag{G}
\end{equation*}
$$

A-weighted $I_{1}$-norm maximization problem:

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x^{0} \in C, x^{j+1}=\operatorname{argmax}\left\{\left\langle a^{j}, x\right\rangle-\sum_{i} w_{i}^{j}\left|x_{i}\right|: x \in C\right\}, j=0, \ldots,
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```
where w}\mp@subsup{w}{}{j}:=-\mp@subsup{g}{}{\prime}(|\mp@subsup{x}{}{j}|)>0\mathrm{ and }\mp@subsup{a}{}{j}:=\mp@subsup{f}{}{\prime}(\mp@subsup{x}{}{j})\in\mp@subsup{\mathbf{R}}{}{n}
```

For penalized/approximate penalized SPCA, $C$ is a unit ball, and above admits a closed form solution thanks to $\mathbf{P} 2$ seen before:

$$
x^{j+1}=\frac{S_{w^{j}}\left(f^{\prime}\left(x^{j}\right)\right)}{\left\|S_{w^{j}}\left(f^{\prime}\left(x^{j}\right)\right)\right\|}, j=0, \ldots
$$

## Example I - A Novel Direct Approach for $l_{1}$-penalized SPCA via (G)

$$
\max \left\{x^{T} A x-s\|x\|_{1}:\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\},(s>0)
$$

Using our results, applying ConGradU reduces to

$$
x^{j+1}=\frac{S_{s e}\left(A_{\sigma} x^{j}\right)}{\left\|S_{s e}\left(A_{\sigma} x^{j}\right)\right\|_{2}}, e \equiv(1, \ldots, 1)
$$

and $S_{w}(a)=\underset{x}{\operatorname{argmin}}\left\{\frac{1}{2}\|x-a\|_{2}^{2}+\|W x\|_{1}\right\}=(|a|-w)_{+} \operatorname{sgn}(a)$.

- This approach can handle matrices $A$ that are not positive semidefinite (by taking $\left.\sigma>0, A_{\sigma}:=A+\sigma I_{n}\right)$.
- In fact, any other convex $f(\cdot)$ can be used!
- Allows for stronger convergence results than when applying the conditional gradient method to the nonsmooth equivalent reformulation.


## Example II: The Approximate $I_{0}$-penalized PCA Problem

$$
\max \left\{x^{\top} A x-s\|x\|_{0}:\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\},(s>0)
$$

- Approximations of the $I_{0}$ norm by some nicer continuous functions have been considered in various contexts, e.g., machine learning [Mangasarian (96), West (03)]; ... Compressed sensing [Borwein-Luke (11)] .
- Naturally emerged from very well-known mathematical approximations of the step and sign functions Bracewell (2000). Formally, we want to replace the problematic expression sgn $(|t|)$ by some nicer function

$$
\|x\|_{0}=\sum_{i=1}^{n} \operatorname{sgn}\left(\left|x_{i}\right|\right)=\lim _{p \rightarrow 0} \sum_{i=1}^{n} \varphi_{p}\left(\left|x_{i}\right|\right)
$$

where $\varphi_{p}: \mathbf{R}_{+} \rightarrow \mathbf{R}_{+}$is an appropriately chosen smooth concave functions, monotone increasing and normalized such that $\varphi_{p}(0)=0, \varphi_{p}^{\prime}(0)>0$.

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where $\varphi_{p}: \mathbf{R}_{+} \rightarrow \mathbf{R}_{+}$is an appropriately chosen smooth concave functions, monotone increasing and normalized such that $\varphi_{p}(0)=0, \varphi_{p}^{\prime}(0)>0$.

- The resulting approximate $I_{0}$-penalized PCA is in the form (G):

$$
\max \left\{x^{T} A x-s \sum_{i=1}^{n} \varphi_{p}\left(\left|x_{i}\right|\right):\|x\|_{2}=1, x \in \mathbf{R}^{n}\right\},(s>0, p>0)
$$

## Examples of Concave $\varphi_{p}(\cdot), p>0$ Approximations for $\|x\|_{0}$

(1) $\varphi_{p}(t)=(2 / \pi) \tan ^{-1}(t / p)$,
(2) $\varphi_{p}(t)=\log (1+t / p) / \log (1+1 / p)$,
(3) $\varphi_{p}(t)=(1+p / t)^{-1}$,
(9) $\varphi_{p}(t)=1-e^{-t / p}$. A nice feature:it also lower bounds $I_{0}$, $\sum_{i=1}^{n} \varphi_{p}\left(\left|x_{i}\right|\right) \leq\|x\|_{0}, \quad \forall x \in \mathbf{R}^{n}$.



Figure : The left plot $\varphi_{p}(t)$ for fixed $p=.05$. The right plot how concave approximation $1-e^{-t / p}$ converges to the indicator function as $p \rightarrow 0$.

## Some Simulations - Random Matrices -[For more see the paper]

- Our goal is to solve very large sparse PCA problems. The largest dimension we approach is $n=50000$.
- However, the ConGradU algorithm applied to $I_{0}$-constrained PCA has a very cheap $O(m n)$ iterations and is limited only by storage of a data matrix.
- Thus, on larger computers, extremely large-scale sparse PCA problems (much larger than those solved even here) are also feasible.


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- Thus, on larger computers, extremely large-scale sparse PCA problems (much larger than those solved even here) are also feasible.
- We here consider random data matrices $F \in \mathbf{R}^{m \times n}$ with $F_{i j} \sim N(0,1 / m)$.
- The experiments consider $n=10(m=6)$ and $n=5000,10000,50000$ (each with $m=150$ ), each using 100 simulations.
- We consider $I_{0}$-constrained PCA with $k=2, \ldots, 9$ for $n=10$ and $k=5,10, \ldots, 250$ for the remaining tests.
- The svdTime is the time required to compute the principal eigenvector of $F^{\top} F$ which is used to compute an initial solution for $I_{0}$-constrained PCA.
- Comparison of ConGradU: with $I_{0}, l_{1}$ penalized version(GPower of Journee et al.) and EM for $l_{1}$-constrained.


## Average Time to Produce Sparse Eigenvectors of $F^{\top} F$

$$
A=F^{T} F \text { with } F \in \mathbf{R}^{m \times n} \text { with } F_{i j} \sim N(0,1 / m)
$$






## Summary and Extensions

Problem structures beneficially exploited to build one very simple scheme ConGradU:

- Encompasses all currently known cheap methods for sparse PCA..and more..
- Can be applied just as easily to solve the original $I_{0}$-constrained problem
- All of the cheap algorithms give similar performance. When desired sparsity is known, our novel scheme appears as the cheapest
- Caveat: None of currently known algorithms provide certificate/bounds to global optimality for the original SPCA.


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Our tools can be easily used to produce novel simple algorithms for tackling directly other similar problems, (details in our paper). For example:
(1) Sparse Singular Value Decomposition:

$$
\max \left\{x^{\top} B y:\|x\|_{2}=1,\|y\|_{2}=1,\|x\|_{0} \leq k_{1},\|y\|_{0} \leq k_{2}\right\}
$$

(2) Sparse Canonical Correlation Analysis:

$$
\max \left\{x^{T} B^{T} C y: x^{T} B^{T} B x=1 y^{T} C^{T} C y=1,\|x\|_{0} \leq k_{1},\|y\|_{0} \leq k_{2}\right\}
$$

(3) Sparse PCA with other convex objectives $f(\cdot)$ or/and additonal "simple" constraints:

$$
\max \left\{f(x):\|x\|_{2}=1,\|x\|_{0} \leq k, x \in \mathcal{C}\right\}
$$

## For More Details, Results....

R. Luss and M. Teboulle. Conditional Gradient Algorithms for Rank-One Matrix Approximations with a Sparsity Constraint.

SIAM Review, (2013). In Press

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## Thank you for listening!

