

# DDR-Distributed Dynamic Routing Algorithm for Mobile Ad Hoc Networks

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# Contents

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- Introduction
- Basic Idea
- DDR - Distributed Dynamic Routing  
Algorithm Description
- Conclusion and Future Work

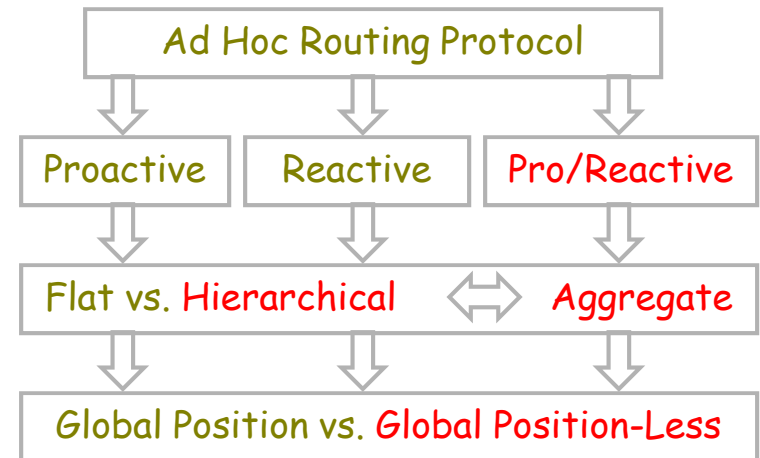
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Introduction, Basic Idea, DDR-Algorithm, Conclusion

# Introduction

🕒 A mobile ad hoc network is a set of wireless nodes forming dynamic autonomous networks.

## 📖 Routing



## 🔑 Critical key features of routing protocols:

🔒 Optimal Path:

🔒 #hops, most stable, ↓delay, ↓energy and ↓loss rate to Dest,

🔒 Fast adaptability to link changes,

🔒 Distributed operation,

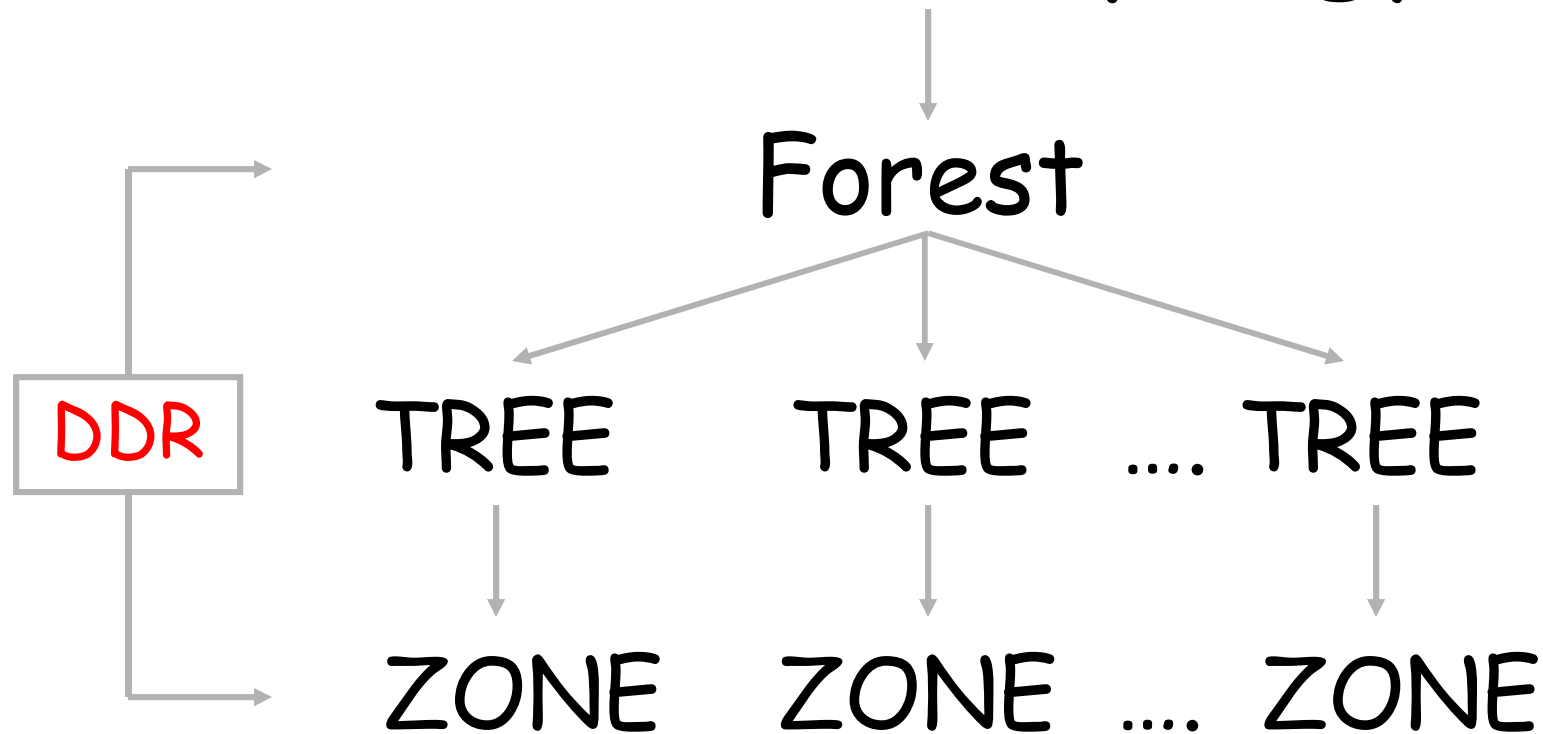
🔒 Loop avoidance.

*Introduction*, Basic Idea, DDR-Algorithm, Conclusion

# Basic Idea (1)

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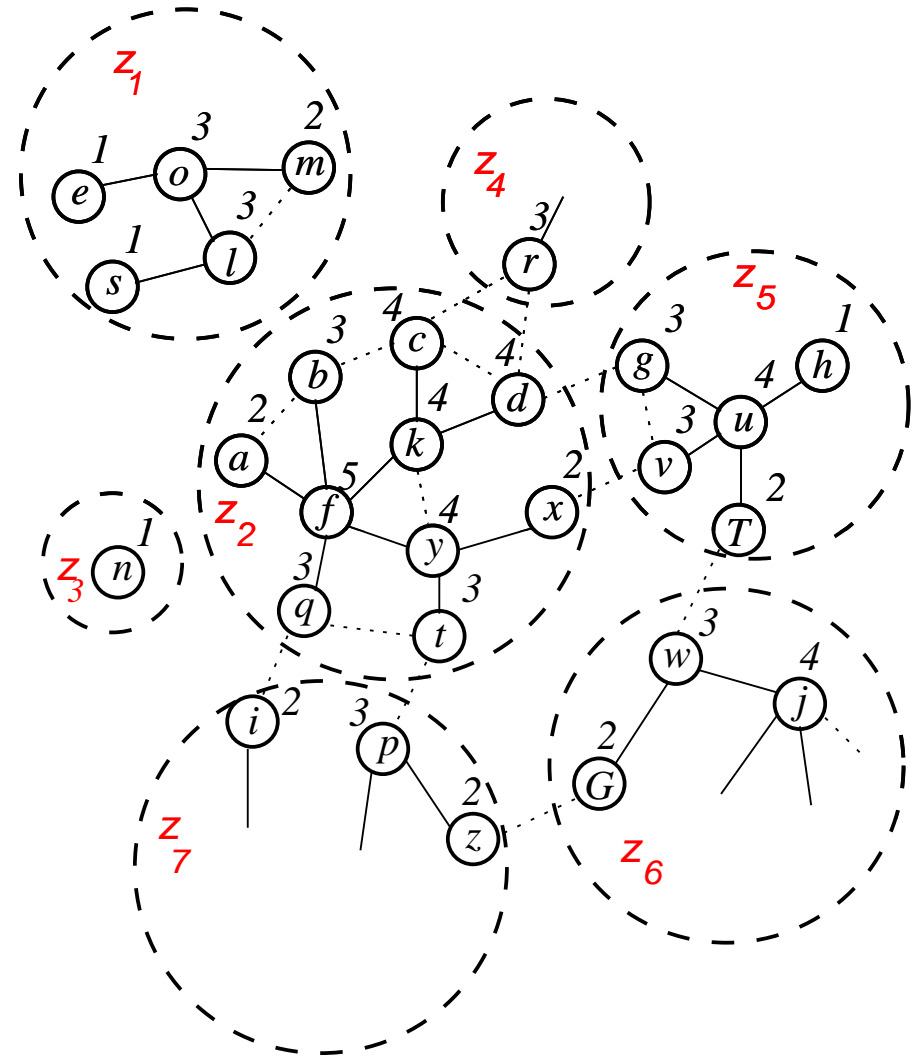
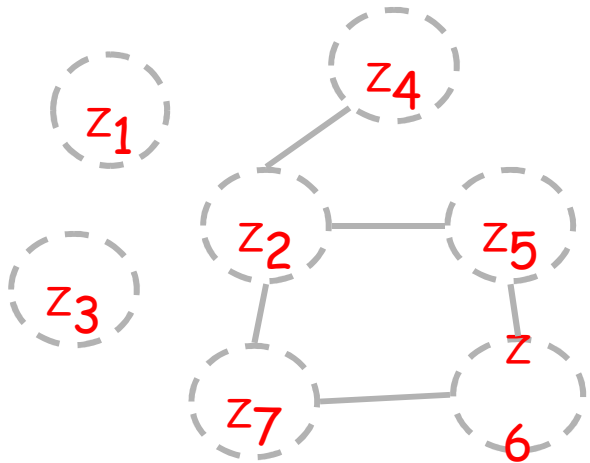
## Network Topology



# Basic Idea (2)

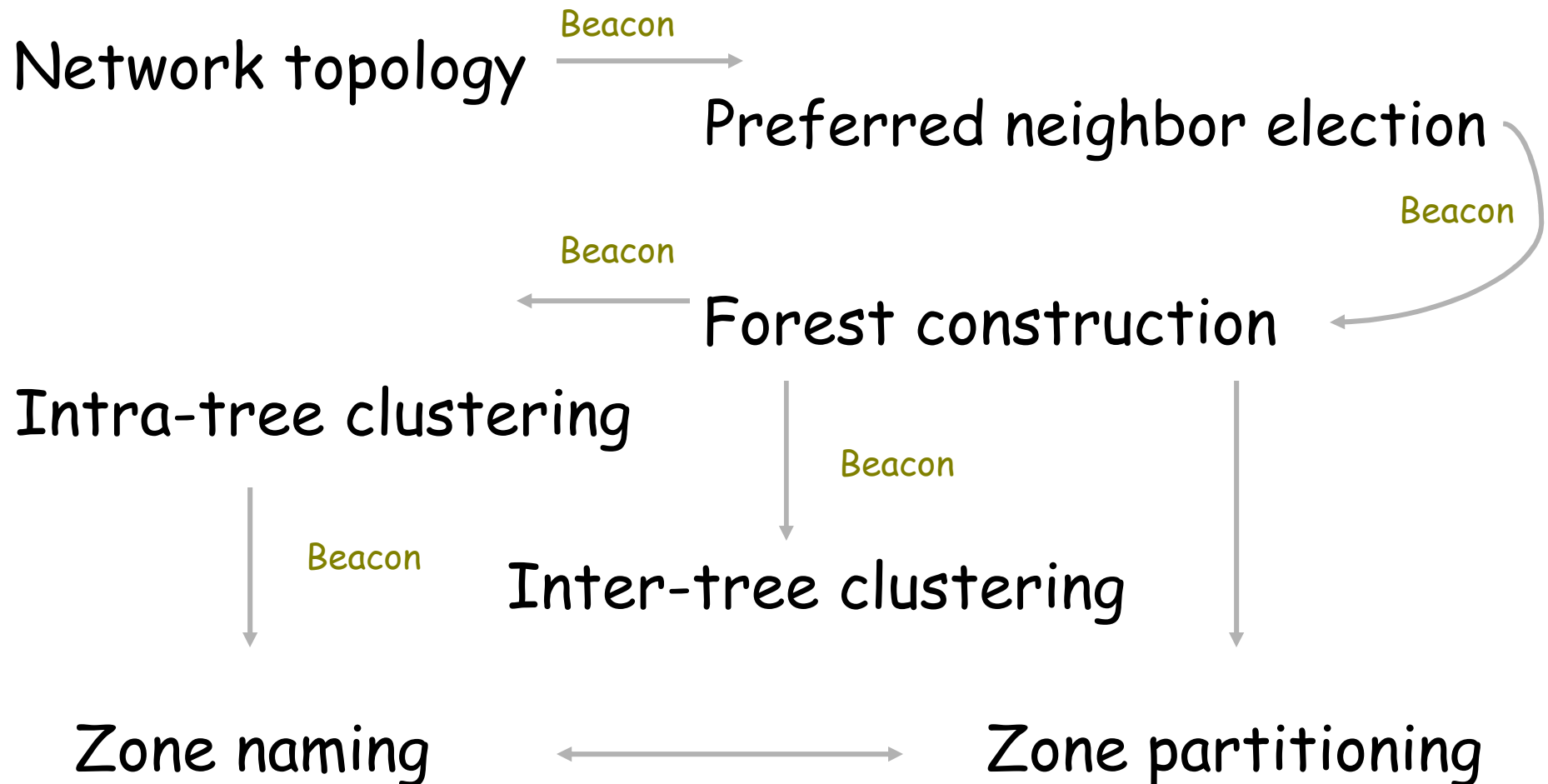
① Forest partitions the network into a set of non over-lapping dynamic zones.

Reduced graph:  $G'=(V', E') \leq G=(V, E)$



# DDR - Algorithm

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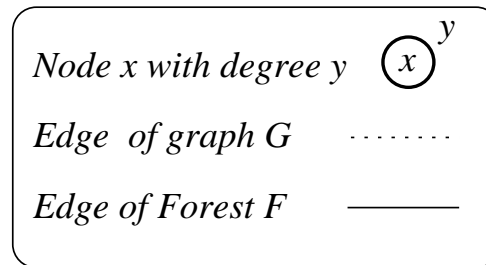
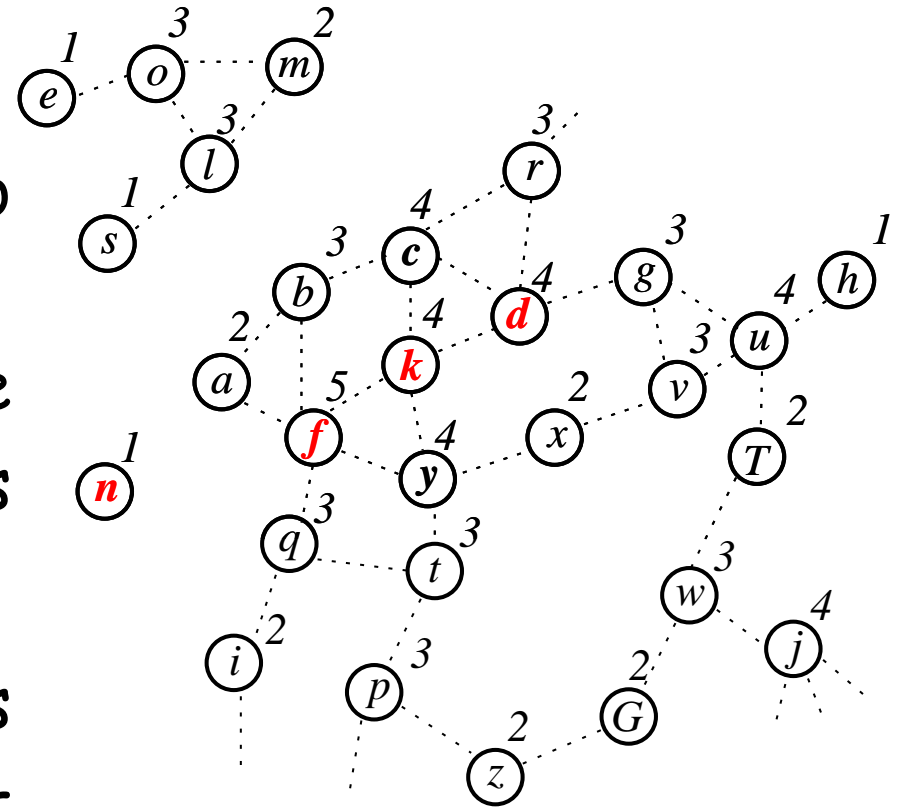
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Introduction, Basic Idea, *DDR-Algorithm*, Conclusion

# Preferred Neighbor Election

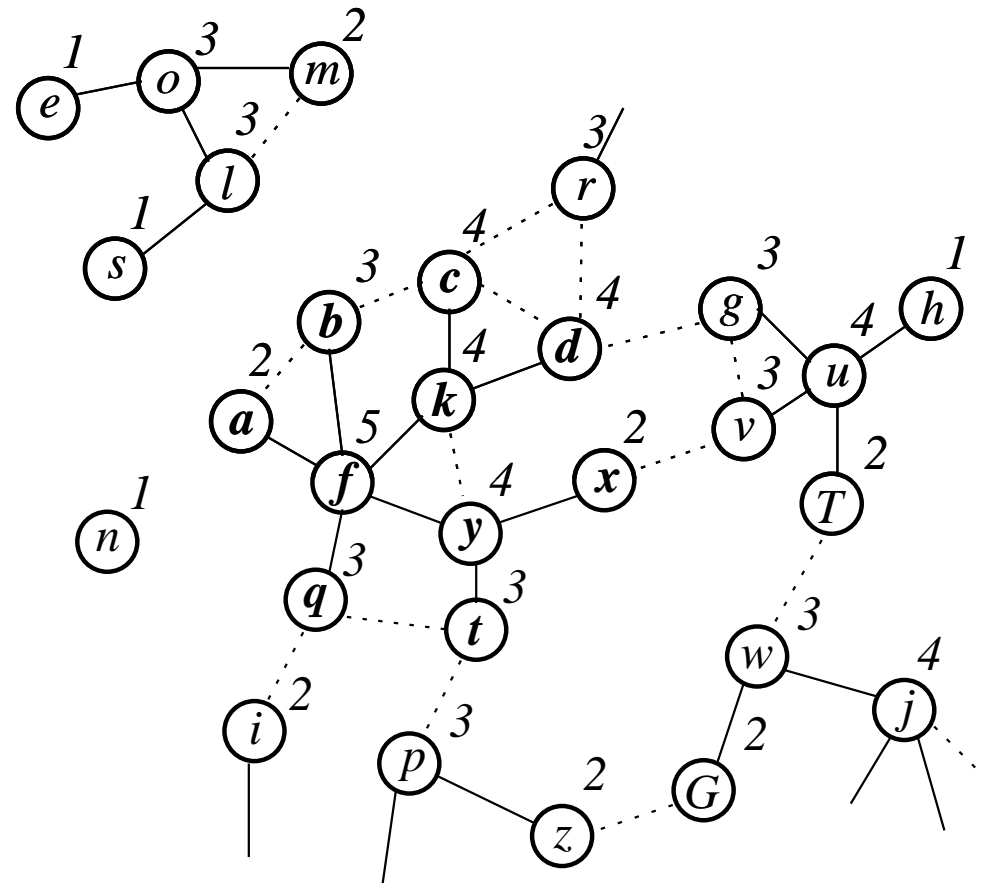
$$PN_x = \{y \mid \deg(y) = \text{Max}(\deg(N_x))\}$$

- ① If this set is empty, then no PN, e.g.  $PN_n = \{\emptyset\}$ .
- ② If this set has only one member, then this member is the elected PN, e.g.  $PN_k = \{f\}$ .
- ③ If more than one member has *max* degree, then we select the one with *max* ID number, e.g.  $PN_d = \{c, K\}$ .



# Forest Construction

① A forest is constructed by connecting each node to its preferred neighbor and vice versa.





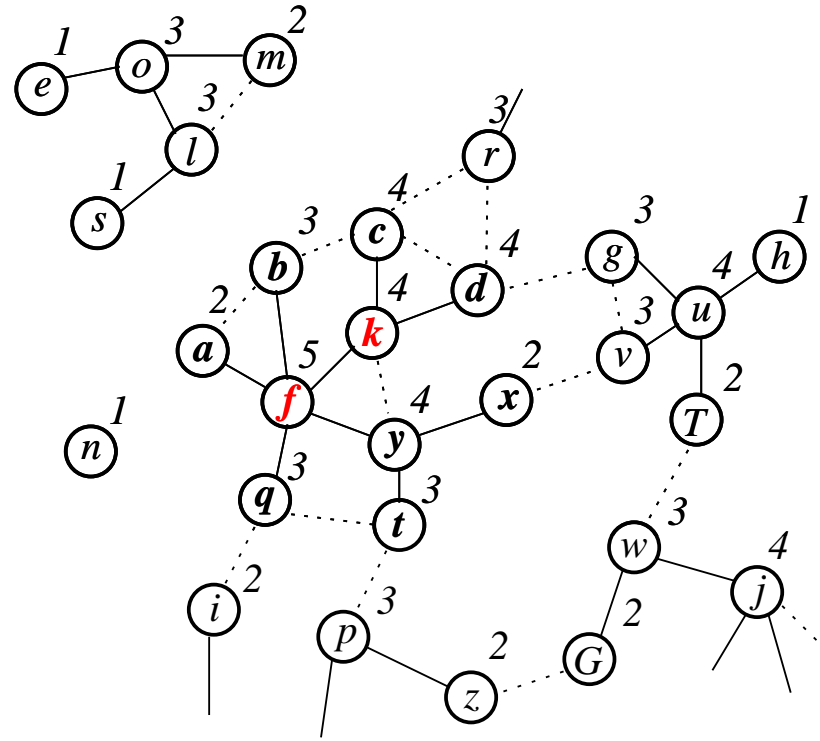
# Intra-tree Clustering (1)

## 👓 Node k:

- 1 PN = f then  $B = (ZID, k, 4, 1, f)$
- 2 Learned\_PN = c: d
- 3 if (PN not changed) then  
 $B = (ZID, k, 4, 0, c: d)$
- 4 Learned\_PN = a: b: q: y

## 👓 Node f:

- 1 PN = y then  $B = (ZID, f, 5, 1, y)$
- 2 Learned\_PN = a: b: q: y: k
- 3 if (PN not changed) then  
 $B = (ZID, f, 5, 0, a: b: q: y: k)$
- 4 Learned\_PN = c: d: x: t



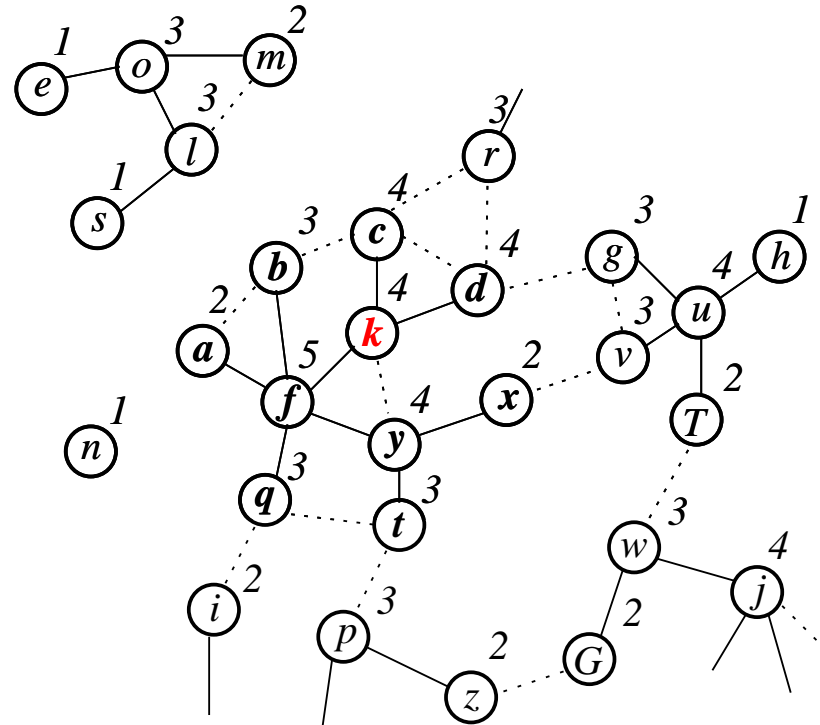
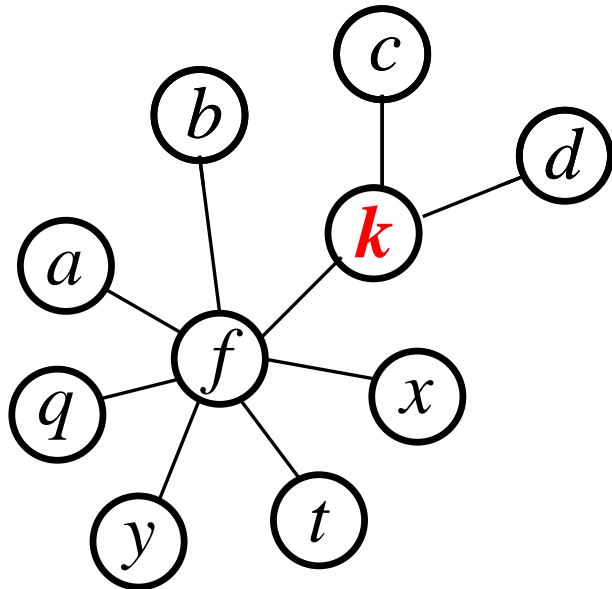
INTRA-ZONE TABLE OF NODES *k* AND *f*

NID	Learned_PN
f	a, b, q, y, t, x
c	-
d	-

NID	Learned_PN
y	x, t
k	c, d
b, a, q	-

# Intra-tree Clustering (2)

VIEW OF NODE **k**  
ABOUT ITS TREE



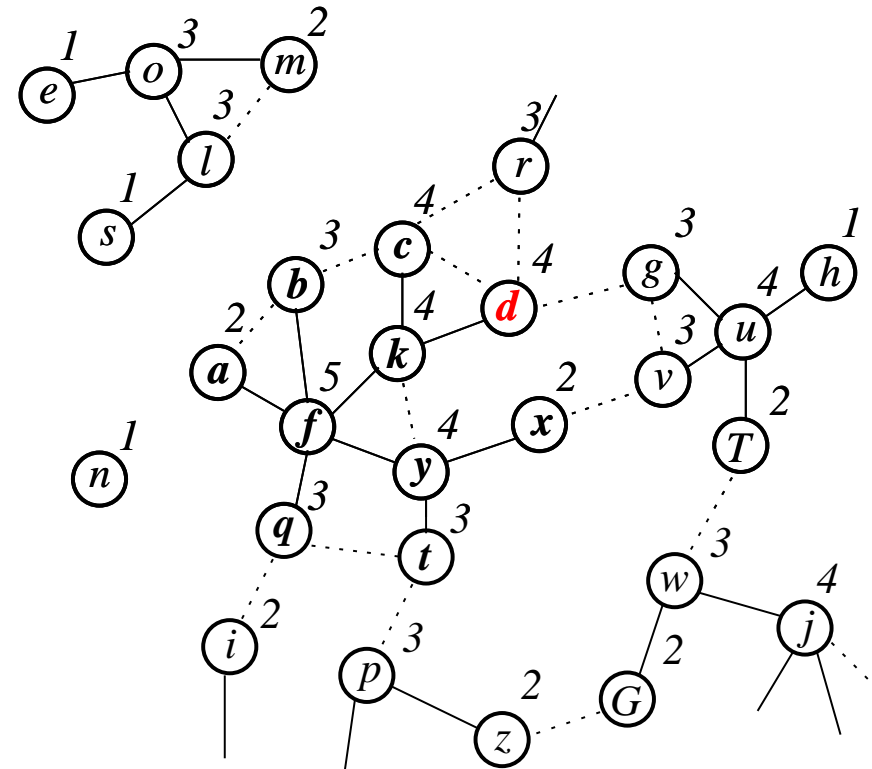
INTRA-ZONE TABLE OF NODES **k** AND **f**

NID	Learned_PN
f	a, b, q, y, t, x
c	-
d	-

NID	Learned_PN
y	x, d
k	c, d
b, a, q	-

# Inter-tree Clustering

- 1 Either a node can succeed to add some nodes to its intra-zone table.
- 2 Otherwise, it puts the remaining nodes in its inter-zone table.



INTER-ZONE TABLE OF NODE **d**

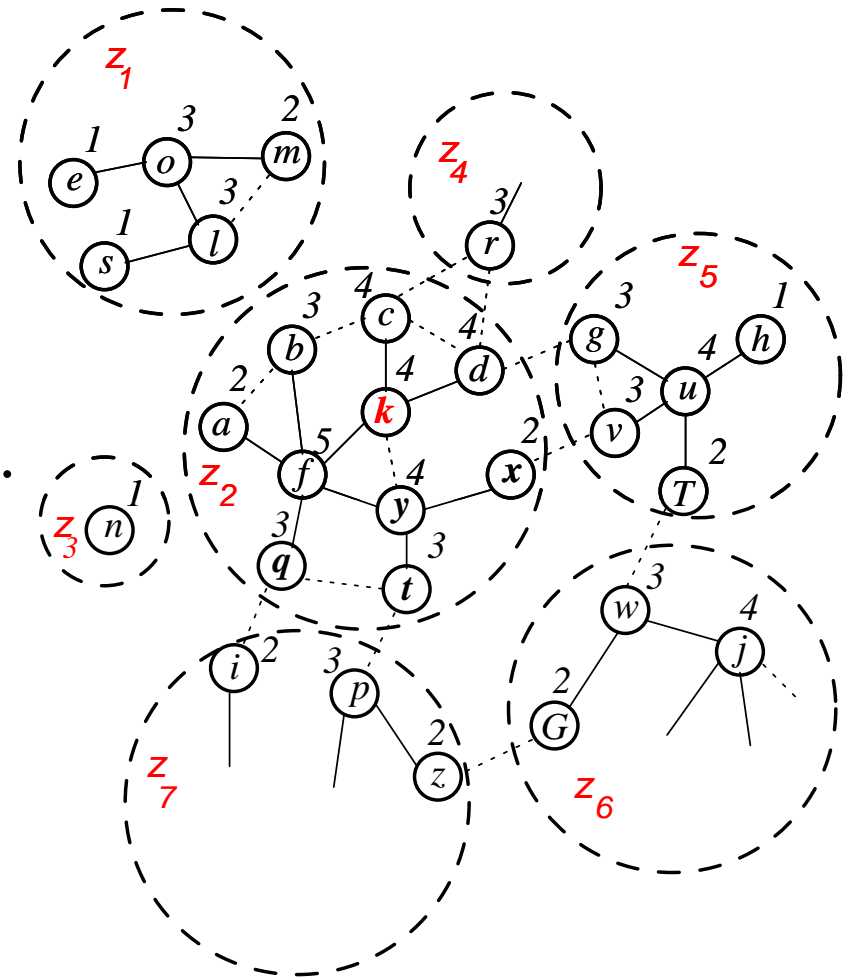
NID	NZID	Z_Stability
r	z4	++
g	z5	++

# Zone Naming

- 1 Select  $q$  highest ID # in intra-zone table, where  $q = \lfloor \frac{n}{d} \rfloor$ .
- 2 Compute a hash function on each selected ID # separately.
- 3 Concatenate all the hashed ID #.

👉 Node  $k$  (for  $n=12$  &  $d=4$ ):

- 1  $q = 4 \Rightarrow$  selected nodes:  $y, x, t, q$
- 2  $h(y)|h(x)|h(t)|h(q)$
- 3  $Z_2 = y'x't'q'$



# Conclusion and Future Work

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🔑 DDR algorithm is:

- 🔓 Simple,
- 🔓 Loop-free,
- 🔓 Distributed,
- 🔓 Bandwidth-efficient.

🔓 Routing protocol description with both numerical and performance analysis.

🔓 Mobile agent over ad hoc networks.

# Preliminary Definitions

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- Beacon

<i>ZID</i>	<i>NID</i>	<i>NID_DEG</i>	<i>MY_PN</i>	<i>PN</i>
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- Intra-zone table

<i>NID</i>	<i>Learned_PN</i>
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- Inter-zone table

<i>GNID</i>	<i>NZID</i>	<i>Z_Stability</i>
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