Design and FPGA Implementation of High Speed Vedic Multiplier

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ABSTRACT

Multiplication is an operation much needed in Digital Signal Processing for various applications. This paper puts forward a high speed Vedic multiplier which is efficient in terms of speed, making use of Urdhva Tiryagbhyam, a sutra from Vedic Math for multiplication and Kogge Stone algorithm for performing addition of partial products and also compares it with the characteristics of existing algorithms. The below two algorithms aids to parallel generation of partial products and faster carry generation respectively, leading to better performance. The code is written in Verilog HDL and implemented on Xilinx Spartan 3 and Spartan 6 FPGA kit using Xilinx ISE 9.1i. The propagation delay of the implemented architecture is obtained to be 28.699ns and 15.752ns respectively.

General Terms

Urdhva Tiryagbhyam Algorithm, Kogge Stone Algorithm, Vedic Multiplier.

Keywords

VM-Vedic Multiplier, KSA-Kogge Stone Adder, RCA-Ripple Carry adder, CLA-Carry look-ahead Adder, CSA-Carry Save Adder.

1. INTRODUCTION

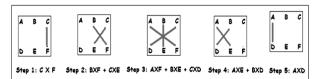
Vedic Mathematics is an ancient system of math practiced during Vedic age which was reconstructed by Jagadguru Swami Sri Bharati Krishna Tirthaji Maharaja[1] between 1911 and 1918 from certain Sanskrit manuscripts. It is perhaps the most refined and efficient mathematical system possible. One of such efficient technique has been employed to enhance the design of a multiplier. Multipliers are the key blocks of a Digital Signal processor. Multiplication is the key aspect, whereby improvement in computational speed of multiplication decreases the processing time of Digital Signal Processors. Convolution, Fast Fourier transforms and various other transforms make use of multiplier blocks.

A faster method for multiplication based on ancient Indian Vedic mathematics is studied in this paper. Among various methods of multiplications in Vedic mathematics, Urdhva Tiryagbhyam is efficient [2]. Urdhva Tiryagbhyam is a general multiplication formula applicable to all cases of multiplication. For addition of partial products in the multiplier Kogge Stone algorithm is used and realized. The code is written in Verilog HDL[3] and synthesized using Xilinx ISE 9.1i and implemented on Spartan 3 and 6 FPGA devices.

2. LITERATURE REVIEW

The algorithms and multiplier architecture was studied from [4, 5, 6, 7, 8, 9, 10, 11] and are represented below.

2.1 Urdhva Tiryagbhyam



Consider ABC as multiplicand and DEF as the multiplier. The steps of multiplication are descriptive in the figure above and the examples are solved below for better understanding. The intermediate carry generated is appended to the very next bit.

2.1.1 Illustration:

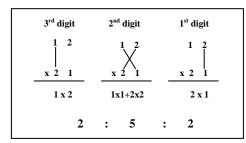


Figure 1: Illustration of decimal multiplication using Vedic technique

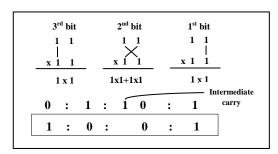


Figure 2: Illustration of binary multiplication using Vedic technique

2.1.2 How is it better than the conventional method?

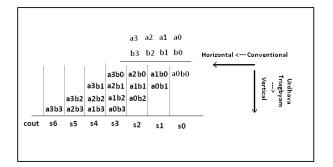


Figure 3: Difference between Conventional Multiplication and Vedic technique

In conventional method, partial products are summated only after every partial product is obtained. Whereas, in Vedic technique, partial products are obtained vertically as shown in the figure above and simultaneously once all the elements of a column are obtained, respective partial products are added. Hence, leads to advancement in speed over the conventional method.

2.2 Kogge Stone Algorithm

Kogge Stone algorithm was developed by Peter M. Kogge and Harold S. Stone and published in an IEEE seminar in 1973[6]. It generates carry in O (log n) time and is used in the industry for high performance arithmetic circuits considering it to be the fastest adder. Carries are computed faster using KSA [9, 10, 11] at the cost of increased area.

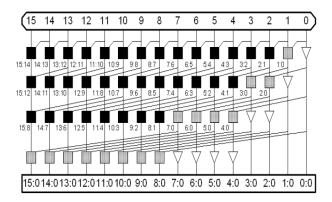


Figure 4: 16-bit Kogge stone adder network [8]

This is an attempt to apprehend the functioning of KSA in three distinct steps:

1. Pre processing

This step involves computation of generate and propagate signals corresponding to each pair of bits in A and B. These signals are given by the logic equations below:

$$P(i) = A(i) \oplus B(i)$$

$$G(i) = A(i) \cdot B(i)$$

2. Carry look ahead network

This is the block responsible for advancement in speed. This step involves computation of carries corresponding to each bit. It uses group propagate and generate as intermediate signals which are given by the logic equations below:

$$Pi: j = P(i:k+1) \cdot P(k:j)$$

Gi: j = G(i:k+1) + (P(i:k+1) \cdot G(k:j))

In the Figure 4, Black box represents the computation of both *Pi:j* and *Gi:j* whereas grey box represents the computation of Gi:j alone and White triangular objects are buffers.

3. Post processing

It involves computation of sum bits. Sum bits are computed by the logic given below:

$$S(i) = P(i) \oplus C(i-1)$$

3. PROPOSED MULTIPLIER

Proposed multiplier architecture of 2x2, 4x4, and 8x8 bit VM module are displayed below. The basic architecture was comprehended from the base paper [4] and modified to obtain the right output as well as gain speed. The major change adopted here in the architecture is that we have used Kogge stone algorithm to add partial products rather than RCA, CLA and CSA.

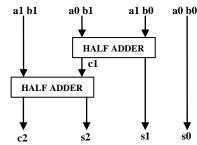


Figure 5: 2x2 Vedic Multiplier

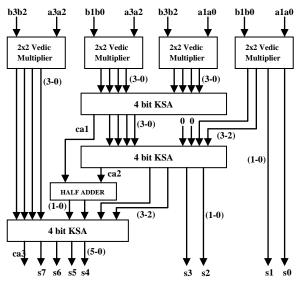


Figure 6: 4x4 Vedic Multiplier using KSA

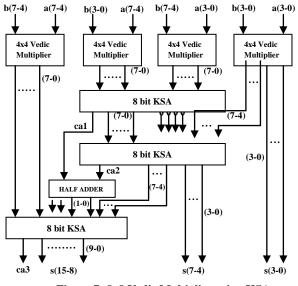


Figure 7: 8x8 Vedic Multiplier using KSA

4. RESULTS AND SIMULATIONS

The Verilog code of 8x8 Vedic multiplier was synthesized using Xilinx ISE 9.1i and was implemented on FPGA device xc3s400-5tq144 of SPARTAN 3 Family. The results are shown below. DIP switches are used as input devices and LEDs are used as output devices. Comparison of delays between 8x8 modified Vedic multipliers using RCA and KSA executed on xc3s700-afg484 and VM using RCA represented in the paper [4] are shown in Table.1.

Device Utilization Summary					
Logic Utilization	Used	Available	Utilization	Note(s)	
Number of 4 input LUTs	192	7,168	2%		
Logic Distribution					
Number of occupied Slices	105	3,584	2%		
Number of Slices containing only related logic	105	105	100%		
Number of Slices containing unrelated logic	0	105	0%		
Total Number of 4 input LUTs	192	7,168	2%		
Number of bonded <u>IOBs</u>	32	97	32%		
Total equivalent gate count for design	1,161				
Additional JTAG gate count for IOBs	1,536				

Figure 8: Device Utilization Summary

All values displaye	d in nanose	econds (r	13)		
Timing constraint: Total number of p	aths / dest	tination	ports:		
	Delay: 28.669ns (Levels of Logic = 18)				
Source:	Source: b<0> (PAD)				
Destination:	Destination: m<15> (PAD)				
Data Path: b<0>	:o m<15>				
		Gate	Net		
				Logical Name (Net Name)	
IBUF: I->0				b 0 IBUF (b 0 IBUF)	
LUT4: I0->0	Z	0.479	1.040	m01/r01/gp[0]/Mxor y Result1 (m01/g<0>)	
LUT4: I0->0	2	0.479	1.040	m01/r02/gp[0]/x1 (m01/r02/g<0>)	
LUT4: I0->0	4	0.479	1.074	m01/r02/g10/x1 (m01/r02/gij<0>)	
LUT4: I0->0	3	0.479	0.941	m01/r02/g30/x1 (m01/c2)	
LUT4: I1->0	2	0.479	1.040	m01/r03/gp[2]/Mxor_y_Result1 (m01/r03/p<2>)	
LUT4: I0->0	2			m01/r03/g20/x12 (m01/r03/g20/x_map4)	
LUT4: I0->0	1			m01/r03/Mxor_s<3>_Result1_F (N746)	
MUXF5:I0->0	3			m01/r03/Mxor_s<3>_Result1 (w<7>)	
LUT3:10->0	4			r02/g30/x (r02/g1j<4⊳)	
LUT3: I0->0	2			r02/g70/x11 (r02/N0)	
LUT4: I3->0	2			r02/Mxor_s<6>_Result1 (i<6>)	
LUT3:12->0	2			r03/g30/x_SW0 (r03/gij<2>)	
LUT3:11->0	2			r03/g60/x1_SW0 (r03/gij<4⊳)	
LUT3: I1->0	3			r03/g70/x1_SW0 (N194)	
LUT4: I3->0	1			r03/Mxor_s<7>_Result1_SW1 (N738)	
LUT4:11->0	1		0.681	r03/Mxor_s<7>_Result1 (j<7>)	
OBUF: 1->0		4.909		m_15_OBUF (m<15>)	
Total				3ns logic, 15.546ns route)	

Figure 9: Timing Details

4.1 Outputs of Simulation:

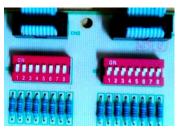


4.2 Outputs of Implementation:

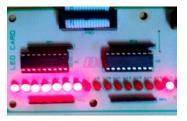
1.	$\underline{(11111111)2*(11111111)2}$
1.	- (111111100000001)2

= (111111100000001)2

Inputs:

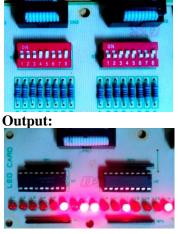


Output:



 $2. \quad \frac{(00111001)2*(00011001)2}{=(0000010110010001)2}$

Inputs:



4.3 Comparison

Table 1. Comparison of delay produced by 8x8 VM using RCA and KSA adders.

using KCA and KSA adders.				
Xc3s700-4fg484				
DELAY(nS)				
[a]. 28.27				
[b]. 27.586				
[c]. 26.178				
0.684				
2.092				
1.408				

Table 2. Comparison of delay produced by 8x8 VM on SPARTAN 3 and SPARTAN 6 device.

DEVICE	PROGRAM	DELAY(nS)
xc3s400-5tq144	8X8 Vedic	28.699
	Multiplier	
xc6slx75t-3fgg676	8x8 Vedic	15.752
	Multiplier	

Table 3. Comparison of delay produced by various 8x8 multipliers.

multipliers.			
DEVICE	xc3s50a-5tq144		
PROGRAM	DELAY(nS)		
Array Multiplier[5]	32.01		
Booth Multiplier [5]	29.549		
VM8x8KSA	23.644		

5. CONCLUSION

By comparing tables 1, 2 and 3 we conclude that the proposed technique of multiplication using UrdhvaTiragbyam algorithm and Kogge Stone algorithm causes less latency when compared to available techniques in literature. The proposed technique when implemented for 8x8 bit multiplication, the delay is found to be 28.699ns on SPARTAN 3 and 15.752 ns on SPARTAN 6. The adoption of KSA algorithm for higher bit size multipliers will further show improvement in speed[9]. Further, higher speeds can be achieved by making use of pipelining and parallel processing techniques. This work will increase awareness of Vedic mathematics techniques in the field of engineering and delivers high performance in DSP Processors.

6. AUTHORS

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