

Design and Implementation of a General-Purpose Median Filter Unit in CMOS VLSI

MUSTAFA KARAMAN, STUDENT MEMBER, IEEE, LEVENT ONURAL, MEMBER, IEEE, AND
ABDULLAH ATALAR, MEMBER, IEEE

Abstract—A general-purpose median filter unit configuration is proposed in the form of two single-chip median filters, one extensible and one real-time. The networks of the chips are pipelined and systolic at bit level and based on the odd/even transposition sorting. The chips are implemented in 3- μm standard CMOS by using full-custom VLSI design techniques. The exact median of elements, in a window size $w = 9$ with arbitrary word length L , can be found by using only one extensible median filter chip. The filter can be extended to arbitrary window size and word lengths by using many chips. For $w > 9$ with arbitrary L , the number of chips required to find the exact medians is no more than the smallest greater integer of $(w/9)^2$. Simulation results show that the extensible median filter chip can be clocked up to 40 MHz, and generate 30/ L megamedians per second. On the other hand, the real-time median filter chip can find the exact running medians of elements in a window of a fixed size $w = 9$ with $L = 8$. According to simulations, it can generate up to 50 megamedians per second with a 50-MHz clock. The chips can be used for the realization of various median filtering techniques. In this paper, the algorithms, VLSI implementations, and testing of the chips are presented together with some possible applications.

I. INTRODUCTION

IN MANY signal and image processing applications, it is necessary to smooth the noisy signals while at the same time preserving the edge information. The most commonly used smoothing techniques are linear filtering [1], averaging filtering [2], and median filtering [3]. The linear filters smooth the noisy signals but also the sharp edges. In addition, the impulsive noise components cannot be suppressed sufficiently by linear filtering, and digital implementation of the linear filters can be bulky and slow. The averaging filters have some undesirable features like outlier points that distort the filtered signal, and edge information loss. The median filters have proved to be good alternatives because they have some very interesting properties: 1) they can smooth the transient changes in signal intensity (e.g., noise); 2) they are very effective for removing the impulsive noises from the signals; 3) they can

preserve the edge information in the filtered signal; and 4) they can be implemented by using very simple digital nonlinear operations. Because of these properties of the median filters, they are frequently used in various signal and image processing applications, such as seismic signal processing, speech processing, computerized tomography, medical imaging, robotic vision, pattern recognition, peak detection, coding, and communication.

The median of a group, containing an odd number of elements, is defined as the middle element, when the elements of the group are sorted. A median filter finds the median of a number of elements at its input. In the standard median filtering applications, a window of size w , where w is odd, is moved along the sampled values of the signal or the image. For each position of the window, the median of the elements within the window is computed and then written as the output pixel located at the same position as the central element of the window. The median computed at this operation is called the *running* or the *moving* median. Since the size of the window is constant, the number of incoming elements is equal to the number of outgoing elements.

Detailed discussions of the theoretical analysis of the median filters and their properties can be found in [4]–[7]. The applications of the median filters in speech processing and in image processing were studied in [8]–[10]. Various median filtering techniques were presented in [11]–[17]. The median filtering was first proposed by Tukey in 1974 [3]; since that time, there have been various studies to develop fast off-line and on-line algorithms, especially for software implementation [18]–[23]. Most of the median filtering algorithms for hardware implementation are based on sorting [24]–[26].

It is possible to design high-performance hardware median filters by using the recent VLSI techniques. A good algorithm for VLSI implementation does not necessarily require the minimal computation but should possess the following properties [27]: the algorithm 1) can be implemented by using only a few simple units, 2) must have a simple and regular communication and control scheme, and 3) should employ parallelism and pipelining. A class of algorithms that have these properties is called the sys-

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The authors are with the Department of Electrical and Electronics Engineering, Bilkent University, 06572 Maltepe, Ankara, Turkey.
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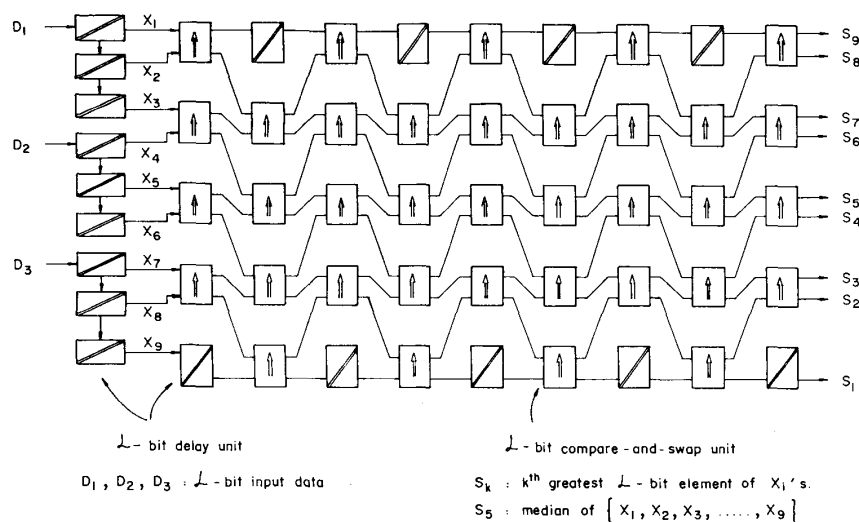


Fig. 1. Ordinary odd/even transposition sorting network for $w = 9$.

tolic algorithms. These algorithms for various computational problems have been discussed in [28]–[30].

Different two-dimensional applications require different window sizes and word lengths for median filtering. The window size varies from 3×3 to quite larger ones. Also the required speed of the filtering operation varies depending on the application. A median filter which is intended as a component of a general-purpose signal or image processor must meet these changing demands. Therefore such a filter must be 1) capable of operating at different word lengths, 2) extensible to larger window sizes, and 3) able to operate at a real-time rate which is about 30 megamedians per second for 1024×1024 frames with 25–30 frames per second. We propose a solution to this problem in the form of two single-chip median filter networks that are implemented using the full-custom VLSI design techniques. Both of the networks are bit-level, systolic, and pipelined odd/even transposition sorting networks that employ all advantages of the systolic algorithms and satisfy the listed constraints for a general-purpose median filter. The networks serve different purposes: one of them is designed for unlimited word length and extensibility to larger window sizes, and the other one is for real-time applications.

In the next section, first the ordinary odd/even transposition sorting network is summarized. Then, the networks of the extensible and the real-time median filter chips are described. Section III covers the VLSI feasibility of the networks, their implementations, and testing of the chips. Some possible applications of the chips are discussed in Section IV. Finally, the results and conclusions of this study are presented in Section V.

II. THE ALGORITHMS

A. The Odd/Even Transposition Sorting Network

The odd/even transposition sorting network presented in [24] is a pipelined modular structure consisting of w

compare-and-swap stages. A compare-and-swap stage has $(w-1)/2$ full-word (L -bit) parallel compare-and-swap units and an L -bit delay unit to delay the window element which is not compared at that stage. Each compare-and-swap unit operates on the *odd* and *even* pairs of the window elements alternately (Fig. 1). Odd pairs are in the form of X_n and X_{n+1} where n is odd, whereas n is even for even pairs. Each full-word compare-and-swap unit compares two L -bit elements at its inputs and interchanges them if necessary so that the larger element will be at the “top” and the smaller one will be at the “bottom” output. At the output of the last stage, the input elements will be sorted such that the largest will be at the “top” and the median will be in the middle.

If the odd/even transposition network described above is implemented in hardware as a single-chip median filter for a fixed window size w_0 , then one cannot use that chip to find the *exact* medians of a window element whose size is larger than w_0 . Also, there is no flexibility on the word length because the compare-and-swap unit can be implemented only for a fixed word length in this configuration. Furthermore, the compare-and-swap unit can complete its job after an L -bit comparison resulting in a large propagation delay. Although the throughput of the network is one L -bit median per clock (assuming a sequential logic implementation), it is not possible to clock this network at a rate required for real-time operation, because the clock period has to be equal at least to the time needed for a full-word comparison. All these results show that the odd/even transposition network in the configuration described above cannot satisfy any of the requirements for a general-purpose median filter: *extensibility* in window size, *unlimited* word length, and *real-time* operation capability.

However, the odd/even transposition network can be modified so that resultant structure(s) can satisfy our needs. Starting from this approach, two median filter networks are designed. These networks are bit-level, systolic, and pipelined odd/even transposition sorting networks.

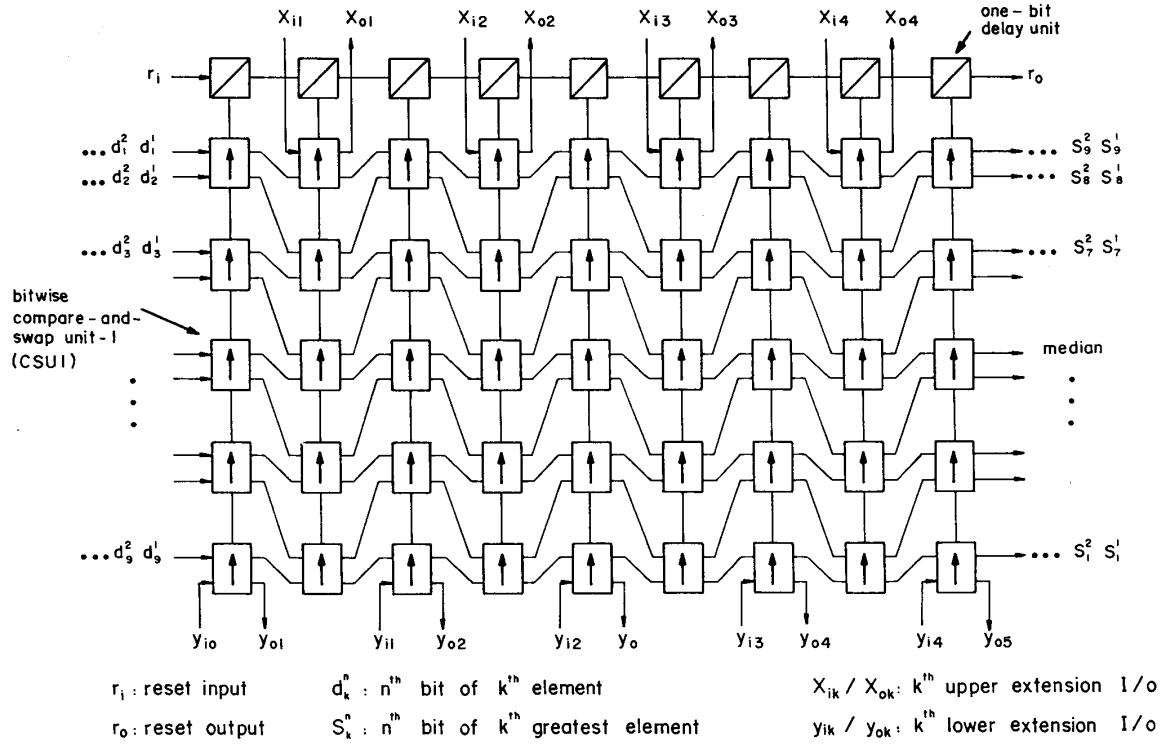


Fig. 2. Extensible median filter network.

They satisfy the listed requirements for a general-purpose median filter. One of the networks is designed for unlimited word length and extensibility to larger window sizes whereas the other one is for real-time applications.

B. The Extensible Median Filter Network

The extensible median filter network consists of nine compare-and-swap stages. A compare-and-swap stage consists of five *bitwise* compare-and-swap units. Each of these units compares two 1-b numbers at its inputs and interchanges them if necessary so that the larger one is at the "top." At the output of the last stage, the data will be sorted such that the largest will be at the top, and the median will be in the middle. At each clock, one bit from each word (total of nine bits) enters the network and one bit of the median is obtained at the output. The flow is from the most significant bits toward the least significant bits both at the input and the output. Due to serial bitwise data flow, this structure allows arbitrary word length L (see Fig. 2).

The bitwise compare-and-swap unit, CSU1, has two 1-b data inputs A_i and B_i , a reset input R , and two 1-b data outputs A_o and B_o . It also has two internal parameters S (1/0: swap/pass) and E (1/0: equal/not equal) which are updated and stored in the CSU1 (Fig. 3). Because of the feedbacks S and E , the best structure for CSU1 is a finite state machine in order to have a reliability of timing [31]. The CSU1 compares A_i and B_i , and either passes them unaltered or swaps them. Its legal operation states are

equal, *pass*, and *swap* states. The present state values of S and E show the cumulative subresult of the bit-wise comparison of the more significant bits while the next state values indicate the new subresult.

CSU1 is set to the *equal* state ($S = 0$, $E = 1$) by the reset signal R at the beginning of each computation cycle. This is the instant after the least significant bit of the previous cycle leaves and before the most significant bit of the next cycle enters the CSU1. Thus, the reset signal also flows through the stages of the network. For this purpose, a chain of the bit-wise delay units is used to reset the stages at appropriate time instants (see Fig. 2). During the computation, the CSU1 stays in the *equal* state ($S = 0$, $E = 1$) and passes the input data unaltered ($A_o = A_i$ and $B_o = B_i$) as long as the two input bits are equal as they flow in. However, it locks itself into the *pass* state when it first finds that $A_i > B_i$ and passes the inputs unaltered. On the other hand, it locks itself into the *swap* state when it first finds that $A_i < B_i$ and swaps the inputs, $A_o = B_i$ and $B_o = A_i$.

The *extension* I/O's ($x_{i/o}$'s and $y_{i/o}$'s, see Fig. 2) of the extensible median filter network are used to extend the filter to window sizes larger than $w = 9$. If the *upper* and *lower extension* inputs (x_i 's and y_i 's) are connected to logic ONE and ZERO, respectively, the corresponding CSU1's operate as delay units. If the network block diagram given in Fig. 2 is used as an array element to sort the elements of a window larger than 9, then these CSU1's operate as usual. As a result of this, the extensibility to indefinitely larger windows is achieved: in the normal mode, it is

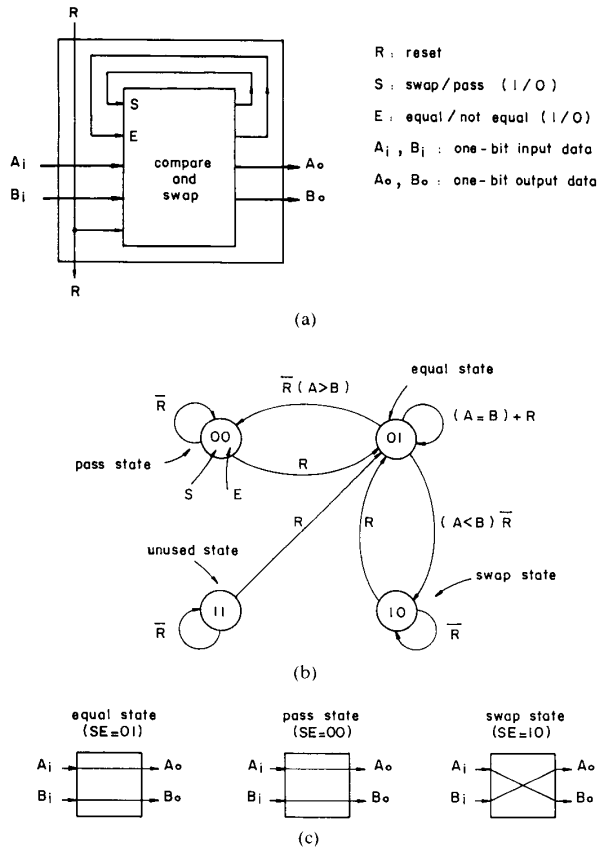


Fig. 3. Compare-and-swap unit-1 (CSU1): (a) block diagram, (b) state diagram, and (c) operations.

possible to compute the exact median of a group of 9 elements using only one network. For groups of w elements where $w > 9$, the number of networks required to find the exact medians is no more than the smallest greater integer of $(w/9)^2$, with appropriate interconnections of the extension and data I/O's.

The extensible network described above generates its outputs with a pipeline delay of $w + L$ clocks, and after the network is full, it finds one median per L clocks. The network can operate at a clock rate that is determined by the delay of one bit-wise compare-and-swap unit. Although the resulting speed can be sufficient for the real-time median filtering of 512×512 frames with $L < 5$, it is less than the real-time operation rate for 1024×1024 frames with $L > 1$.

C. The Real-Time Median Filter Network

A median filter with very high throughput can be designed by interconnecting an L number of pipelined odd/even transposition sorter blocks in parallel [26]. In each block, the compare-and-swap units (total of six) and delay units (total of four) at the upper and lower end parts of a sorter are eliminated because they are not necessary to find the median. The block diagram of the overall network

for $w = 9$ and $L = 8$ is shown in Fig. 4. In this network, the data enter in such a way that the most significant bits go to the first block, the second most significant bits to the second block, and so on. There are nine 8-b input data, and one 8-b median is obtained per clock. At every clock, three new elements enter the network, corresponding to the new elements of a sliding 3×3 window.

The bit-wise compare-and-swap unit, CSU2, used in this network is slightly different than CSU1 in such a way that the S and E parameters of CSU1 are external parameters in CSU2. These parameters are taken as inputs S_i and E_i (the cumulative subresult of the bit-wise comparison of the more significant bits) from the upper block, used, updated and sent out as S_o and E_o (the new cumulative subresult) to the lower block. Hence, CSU2 does not have any storage or the reset input signal. The functional description of CSU2 is given in Fig. 5.

During the computation, CSU2 compares the inputs A_i and B_i and passes or swaps them conditionally depending on S_i and E_i . If $E = 1$ and $S = 0$, it passes or swaps the inputs so the larger one goes to A_o and the smaller one to B_o . On the other hand, if $E = 0$ and $S = 0/1$, CSU2 passes/swaps the inputs, independent of the result of the comparison of the inputs. It also updates the S_i and E_i and sends them out as S_o and E_o for the comparison of the less significant bits.

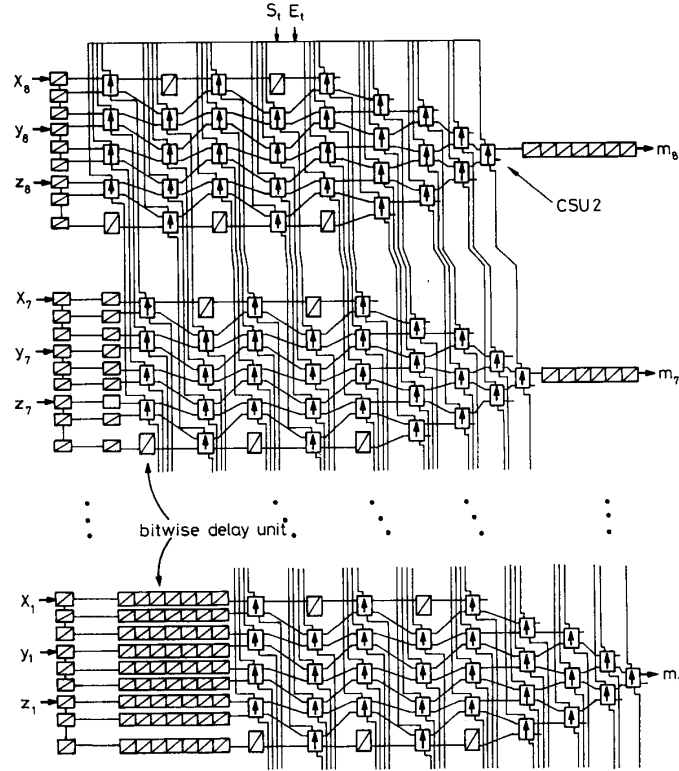
The computation at each stage of a lower block starts after the corresponding stage of the upper neighborhood block completes its operation so that the S_i 's and E_i 's are ready at the input of that stage. As a result of this, each stage of a lower block must operate one clock later compared to the corresponding stage of the upper block. In order to employ a proper timing, the pipeline delay units are used both at the input and output of the network (see Fig. 4).

The throughput of the network is one full-word median per clock. The clock period is determined by the delay of one CSU2. Recent VLSI technology allows the implementation of CSU2 at a speed larger than the real-time rate for the 1024×1024 frames with $L = 8$.

III. VLSI IMPLEMENTATION

A. VLSI Feasibility

Both the extensible and real-time median filter networks are regular arrays of bit-wise compare-and-swap and delay units. Also, their internal communication schemes are simple and regular. This makes the VLSI implementations easy and straightforward. The extensible median filter chip can be packaged in a 28-pin package with 2×1 multiplexed I/O's whereas the real-time median filter chip can be packaged in a 40-pin package without multiplexed I/O's. Furthermore, the testing of the VLSI chips may be easily accomplished by the *functional test* technique, since the operations of the cells can be selectively probed by issuing proper test vectors.



x_i, y_i, z_i : i^{th} bits of the inputs corresponding the new elements in a 3×3 sliding window.

m_i : i^{th} bit of the median

S_i, E_i : Test inputs for testing of the blocks individually.

Fig. 4. Real-time median filter network.

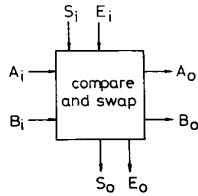


Fig. 5. Compare-and-swap unit-2 (CSU2). S_i and E_i indicate the cumulative subresult of the bit-wise comparison of more significant bits, while S_o and E_o indicate the new cumulative subresult.

$$S_o = S_i + E_i \quad (A_i < B_i)$$

$$E_o = E_i \quad (A_i = B_i)$$

$$A_o = \text{if } S_o \text{ then } B_i \text{ else } A_i$$

$$B_o = \text{if } S_o \text{ then } A_i \text{ else } B_i$$

One may choose to implement either the extensible median filter network or the real-time network for a larger w , but neither of the networks is preferable for a large w since the area is proportional to w^2 . We have chosen $w = 9$ because it is the minimum and the most commonly used window size in two-dimensional median filtering applications. In an attempt to increase the efficiency of the silicon area use, one may consider designing a single network that can operate as either one of the median filter networks presented above, depending on the state of an external signal. This may be achieved by designing a new compare-and-swap unit which is externally selectable to operate as

CSU1 or CSU2. Examination of this possibility did not give the desired result: the chip area turned out to be larger because of the decreased regularity and increased interconnections.

We have also examined the implementation of the two networks in the same chip. The I/O's of the two networks need to be multiplexed. Also, strong buffers are needed to drive the large wiring capacitive loads due to very long wiring paths from I/O pads to the network I/O's. All these increase the area and decrease the throughput of the networks significantly so that the real-time operation falls into a critical state. Also, while one of the networks operates the other one cannot be used due to multiplexed I/O's, i.e., one of the networks will be idle at all times. The implementation of the networks as two single chips avoids these undesired results.

B. The Chips

VLSI implementations of the algorithms consist of circuit designs and generation of layouts. Our chips are designed in $3\text{-}\mu\text{m}$ n-well double-metal CMOS, since CMOS technology possesses low power consumption, a well-

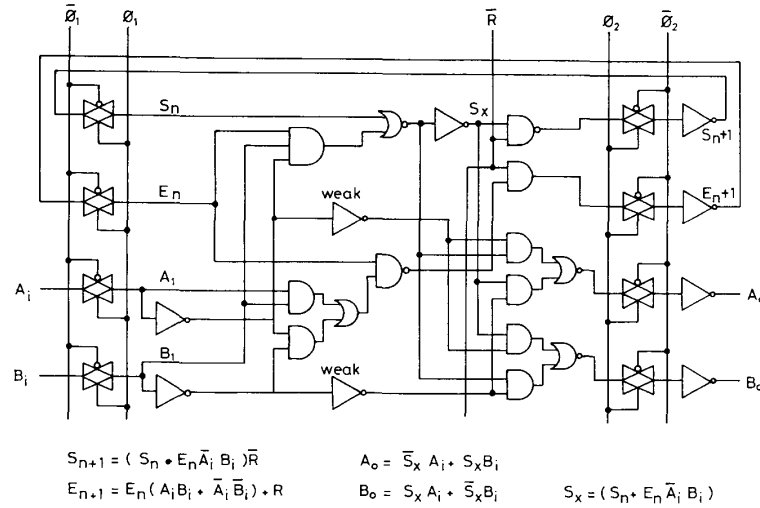


Fig. 6. Circuit diagram of the CSU1.

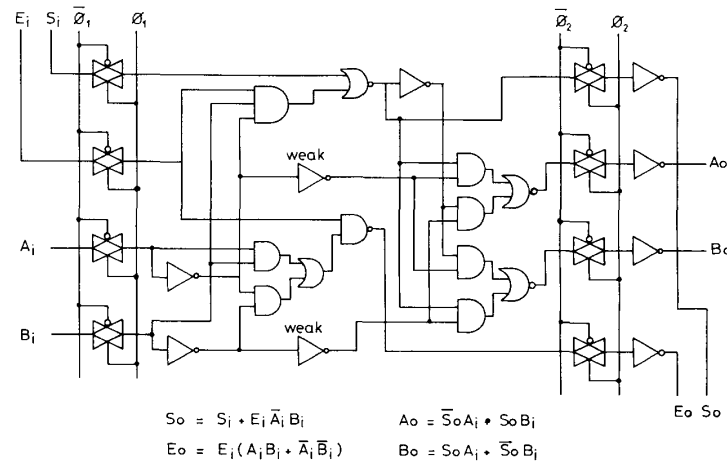


Fig. 7. Circuit diagram of the CSU2.

established design procedure, and a speed comparable to NMOS. Also, the noise immunity of CMOS is better than that of NMOS [32].

Due to pipelined structure and systolic data flow, a sequential circuit design is required with two-phase nonoverlapping clocking strategy. The circuit diagrams of the CSU1 and CSU2 are shown in Figs. 6 and 7. In these circuits, the inputs are latched in during $\phi_1 = 1$, and the computed outputs are sent out during $\phi_2 = 1$. Both of the circuits can function correctly even if the complement of ϕ_1 is used as ϕ_2 (see Fig. 8), but the operations of the circuits with two nonoverlapping clocks, ϕ_1 and ϕ_2 , will be more reliable. We have used standard (complementary) CMOS logic style because this logic style offers a more reliable timing, a rather easier mapping in layout, a lower power consumption, and a better noise immunity compared to the other styles. Furthermore, its area and speed performance is comparable with the other styles [31], [33].

In the layout generation step, first the floor plans of the chips are completed. Layouts of the networks are edited hierarchically starting from the compare-and-swap units (CSU1 and CSU2), which are the basic cells. Then circuit equivalents of the layouts are extracted and simulated. For editing and simulations, we have used the Berkeley CAD tools of the University of California [34] and the VLSI tools of the University of Washington [35] running under the Unix operating system on a SUN160 workstation. Layouts are edited with Magic. The timing and logic simulations are performed using the SPICE, Rn1, and Esim tools. After completing the layout editing and simulation tasks, the chip layouts are converted to CIF and GDS2 codes to be used in fabrication. The waveforms given in Fig. 8 are the samples for the SPICE simulation results of the CSU1, which are obtained at their maximum clock rates. The complete layouts of the chips are shown in Figs. 9 and 10.

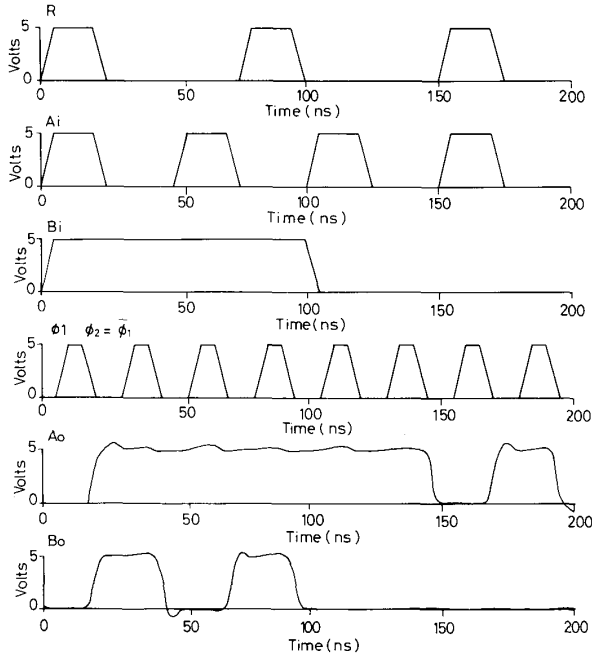


Fig. 8. Sample of the waveforms obtained by SPICE simulation of the CSU1.

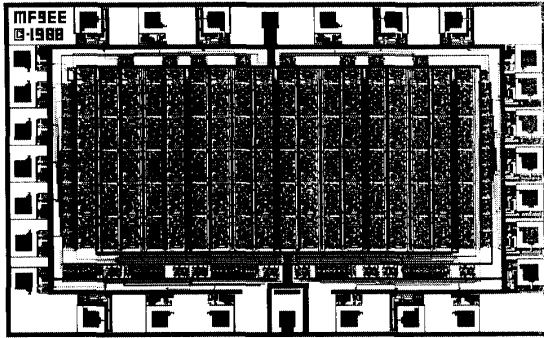


Fig. 9. Complete layout of the extensible median filter chip.

C. Testing

For the testing of the chips, the functional test method is used [36]–[38]. The testing operation is composed of the application of test vectors as inputs to the chips and the comparison of the outputs with expected results. The test vectors and the expected results are generated by the software tools written for these specific purposes.

The extensible median filter chip is tested by using a few hundred test vectors. If we apply the same element to all of the inputs, then every CSU1 will act in the *equal state*. On the other hand, the sorted inputs (the largest one is at the “top”) will cause all of the CSU1’s to operate in the *pass state*. Consequently, a set of inputs sorted in reverse order (the largest one is at the “bottom”) will make the CSU1’s operate in the *swap state*. In this way, the CSU1’s are tested for all operation states.

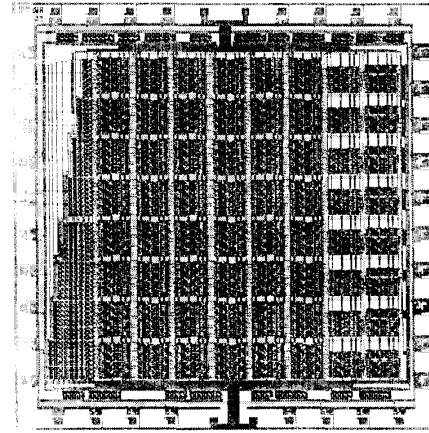


Fig. 10. Complete layout of the real-time median filter chip.

The real-time median filter chip is tested using 12 288 ($8 \times 3 \times 2^9$) test vectors. The S_i ’s and E_i ’s at the top of the network are connected to external test control inputs S_i and E_i , respectively (see Fig. 4). These test control signals will be used to isolate the operation of each block from the others. The k th block can be tested individually by applying logic ONE to all of the upper blocks’ inputs so that it receives the external S_i and E_i without any change. As a result, each block will be tested by using 1536 (3×2^9) test vectors: all combinations of one-bit 9 inputs times 3, where 3 comes from the legal values of the pair $S_i E_i$ (01, 00, 10).

IV. APPLICATIONS

The extensible and the real-time median filter chips can be selectively used in a general-purpose digital image or signal processor environment by means of the chip-enable signal that each chip has [39]. The exact medians of the elements, in a window size $w=9$ with arbitrary word length L , can be found by using only one extensible median filter chip. For $w > 9$ with arbitrary L , the number of required chips to find the exact medians is no more than the smallest greater integer of $(w/9)^2$. In this case the chips are interconnected similarly to the configuration shown in Fig. 11. The real-time median filter chip can find the exact running medians of the elements in a window of fixed size $w=9$ with fixed word length $L=8$ at the real-time rate. Furthermore, the extensible and/or real-time median filter can be used for the realization of many median filtering techniques, as outlined below.

- The extensible median filter is a favorable choice to realize *adaptive-length* median filters [16], [40], since one can change the window size from 3 to indefinitely large ones by using the extensible median filter chip(s) by applying the logic ZERO’s or ONE’s to unused inputs of the chip(s) appropriately.
- For the realization of *weighted* median filters [13], the extensible median filter can be used with a pipelined

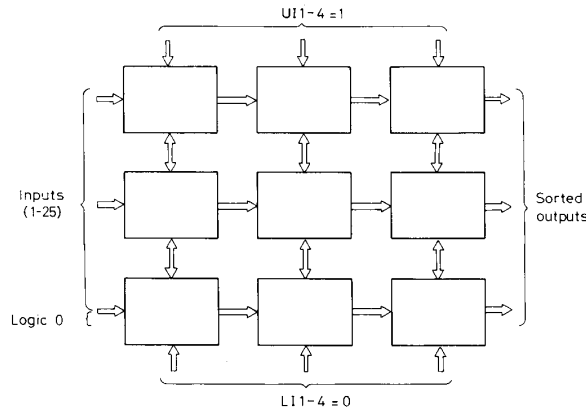


Fig. 11. Interconnections of the extensible median filter chips for $w = 25$.

multiplier to multiply the input data with the weight coefficients. Since all input data of the chip are entered to the chip directly at each move of the window, one can realize an adaptive weighted median filter by changing the weight coefficients at each position of the window on the frame.

- A pair of the extensible or the real-time median filter chips can be used as a *selective median* filter [17] together with an external control logic consisting of two full-word subtractors and a full-word comparator.
- Either the extensible or the real-time median filter chip can be used as a *line-recursive* median filter [16] by loading the window elements from the frame appropriately.
- The extensible median filter chips can be used for the realization of *multilevel* median filters [15] together with a reasonable external hardware.
- The chips can be used for the realization of *separable* median filters [12] without any external hardware.

V. RESULTS AND CONCLUSIONS

Two single-chip median filters are designed to form a general-purpose median filter unit. One of the chips is designed for unlimited word length and extensibility to larger window sizes whereas the other one is for real-time video applications. The chips are implemented in 3- μ m double-metal CMOS by using full-custom VLSI design techniques. The performance characteristics of the chips are summarized in Table I. These characteristics of the chips are comparable with those of the chips in the literature, for example [41]. The extensible and real-time median filter chips with/without a reasonable external hardware can be used for the realization of various median filtering techniques.

The main contributions of this study are the architecture of the extensible median filter and its VLSI implementation. Another achievement of this study is a fast-running VLSI median filter that can meet the demands for real-time applications both today and in the near future. Further-

TABLE I
CHARACTERISTIC PARAMETERS OF THE CHIPS

Parameter	The Extensible Median Filter Chip	The Real-Time Median Filter Chip
Die size	$2.8 \times 4.5 \text{ mm}^2$	$7.2 \times 6.9 \text{ mm}^2$
Transistor count	5,000	22,000
Transistors/ mm^2	400	450
Max. clock frequency	40 MHz	50 MHz
Max. throughput	30/L mega medians/sec	50 mega medians/sec
Pin number	28	40
Max. power dissipation	250 mW	800 mW

more, it is concluded that a general-purpose median filter unit can be formed by selectively using the chips in full-scale general-purpose digital signal or image processor environments.

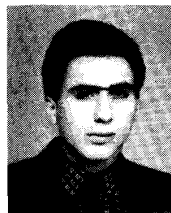
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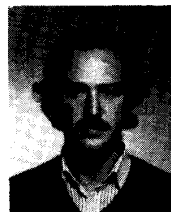
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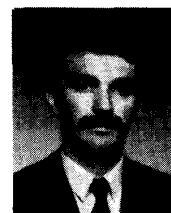
Mustafa Karaman was born in Balıkesir, Turkey, in 1964. He received the B.S. degree from the Middle East Technical University, Ankara, Turkey, and the M.S. degree from Bilkent University, Ankara, Turkey, in 1986 and 1988, respectively, both in electrical and electronics engineering. He is studying for his Ph.D. degree in the Department of Electrical and Electronics Engineering at Bilkent University, where he has been a Research Assistant since 1986. His present research interests include VLSI signal processing, computer-aided design tools for VLSI, and circuit simulation.



Levent Onural (S'82-M'85) was born in Izmir, Turkey, in 1957. He received the B.S. (high honors) and M.S. degrees in electrical engineering from the Middle East Technical University, Ankara, Turkey, in 1979 and 1981, respectively, and the Ph.D. degree in electrical and computer engineering from the State University of New York at Buffalo in 1985. He was a Fulbright Scholar between 1981 and 1985.

After a Research Assistant Professor position at the Electrical and Computer Engineering Department of the State University of New York at Buffalo, he joined the Electrical and Electronics Engineering Department of Bilkent University, Ankara, Turkey, where he is an Associate Professor at present. His present research interests are in signal processing, image processing, holography, fractal models, and signal processing hardware.

Dr. Onural is a member of SPIE.



Abdullah Atalar (M'88) was born in Gaziantep, Turkey, in 1954. He received the B.S. degree from the Middle East Technical University, Ankara, Turkey, in 1974, and the M.S. and Ph.D. degrees from Stanford University, Stanford, CA, in 1976 and 1978, respectively, all in electrical engineering. His thesis work was on reflection acoustic microscopy.

From 1978 to 1980 he was a first Post-Doctoral Fellow and later an Engineering Research Associate at Stanford University, where he continued his work on acoustic microscopy. For eight months he was with Hewlett Packard Laboratories, Palo Alto, CA, engaged in photoacoustics research. In 1980 he joined the Middle East Technical University as an Assistant Professor. From 1982 to 1983 he was on leave from the University with Ernst Leitz Wetzlar, Wetzlar, West Germany, where he was involved in the development of the commercial acoustic microscope. He is presently an Associate Professor and Chairman of the Electrical and Electronics Engineering Department at Bilkent University, Ankara, Turkey. His current research interests include computer-aided design in electrical engineering, acoustic imaging, and linear acoustics.