

Estimating the Largest Eigenvalue of a Positive Definite Matrix

By Dianne P. O'Leary, G. W. Stewart* and James S. Vandergraft

Abstract. The power method for computing the dominant eigenvector of a positive definite matrix will converge slowly when the dominant eigenvalue is poorly separated from the next largest eigenvalue. In this note it is shown that in spite of this slow convergence, the Rayleigh quotient will often give a good approximation to the dominant eigenvalue after a very few iterations—even when the order of the matrix is large.

Let A be a positive definite matrix of order n with eigenvalues $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_n > 0$ corresponding to the orthonormal system of eigenvectors x_1, x_2, \dots, x_n . In some applications, one must obtain an estimate of λ_1 without going to the expense of computing the complete eigensystem of A . A simple technique that is applicable to a variety of problems is the power method. Starting with a vector u_0 of Euclidean norm unity ($\|u_0\|_2 = 1$), one iterates as follows:

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1      : loop for  $k := 0, 1, 2, \dots$ 
1.1    :  $v_k := Au_k$ ;
1.2    :  $\rho_k := u_k^T v_k$ ;
1.3    :  $u_{k+1} := v_k / \|v_k\|_2$ ;
1      : end loop.
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The theory of the method is well understood (e.g., see [4]). If $\lambda_1 > \lambda_2$ and $x_1^T u_0 \neq 0$, then the vectors u_k converge linearly to x_1 at a rate proportional to $(\lambda_2/\lambda_1)^k$. The Rayleigh quotients ρ_k converge to λ_1 at a rate proportional to $(\lambda_2/\lambda_1)^{2k}$.

Convergence of the method can be hindered in two ways. First, if $x_1^T u_0$ is pathologically small compared to some of the numbers $x_i^T u_0$ ($i > 1$), then it will take many iterations for u_k to become a good approximation to x_1 . Second, if λ_2 is very near λ_1 , the final rate of convergence will be slow. We can do very little about the first problem, except to note that it is unlikely to occur with a randomly chosen starting vector u_0 . Moreover, if our object is to compute the eigenvector x_1 , the only way to accelerate slow convergence due to an unfavorable ratio λ_2/λ_1 is to use more elaborate methods, such as simultaneous iteration [2], [3] or Lanczos tridiagonalization [1]. However, if we are only interested in a rough approximation to λ_1 , it will

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be provided by ρ_k after a few iterations. The purpose of this note is to quantify this assertion with specific bounds.

To simplify the analysis, let the matrix A be scaled so that $\lambda_1 = 1$. If we set $w_i = x_i^T u_0$, then

$$(1) \quad u_0 = w_1 x_1 + w_2 x_2 + \dots + w_n x_n,$$

where w_1 is assumed to be nonzero. Then it is easy to verify that

$$(2) \quad \rho_k = \frac{w_1^2 + \sum_{i=2}^n w_i^2 \lambda_i^{2k+1}}{w_1^2 + \sum_{i=2}^n w_i^2 \lambda_i^{2k}} \equiv \frac{\nu_k}{\delta_k}.$$

The number ρ_k satisfies $0 < \rho_k \leq 1$, and it will be a good approximation to $\lambda_1 = 1$ if it is near one. Thus, we are led to investigate the worst case, when the eigenvalues of A are distributed to make ρ_k as small as possible. This distribution is easily determined by differentiating the expression (2) with respect to each λ_j and setting the results to zero. This gives

$$(2k + 1)w_j^2 \lambda_j^{2k} \delta_k - 2kw_j^2 \lambda_j^{2k-1} \nu_k = 0 \quad (j = 2, 3, \dots, n).$$

An inspection of (2) shows that if $w_j \neq 0$ then no minimum of ρ_k can have $\lambda_j = 0$. Hence, we take

$$(3) \quad \lambda_j = \frac{2k}{2k + 1} \frac{\nu_k}{\delta_k} = \frac{2k}{2k + 1} \rho_k \quad (j = 2, 3, \dots, n).$$

Thus, if $\tilde{\rho}_k$ denotes the minimum value of ρ_k , the minimizing distribution takes $\lambda_2, \lambda_3, \dots, \lambda_n$ equal and slightly less than $\tilde{\rho}_k$. We may obtain an equation for $\tilde{\rho}_k$ by substituting (3) into (2) to give

$$\tilde{\rho}_k = \frac{1 + \left(\frac{2k}{2k + 1}\right)^{2k+1} \tau^2 \tilde{\rho}_k^{2k+1}}{1 + \left(\frac{2k}{2k + 1}\right)^{2k} \tau^2 \tilde{\rho}_k^{2k}},$$

where

$$\tau^2 = (w_2^2 + w_3^2 + \dots + w_n^2)/w_1^2.$$

This expression may be simplified to the polynomial equation

$$(4) \quad f_k(\tilde{\rho}_k) \equiv \frac{c_k \tau^2}{2k + 1} \tilde{\rho}_k^{2k+1} + \tilde{\rho}_k - 1 = 0,$$

where

$$c_k = \left(\frac{2k}{2k + 1}\right)^{2k}.$$

The polynomial $f_k(\rho)$ is increasing for $\rho \geq 0$. Since $f_k(0) = -1$ and $f_k(1) > 0$, f_k has a unique zero $\tilde{\rho}_k \in (0, 1)$.

Equation (4) allows us to give a qualitative description of the behavior of $\tilde{\rho}_k$. The positive zero of (4) decreases as the leading coefficient increases. Now c_k is well

behaved, having e^{-1} as a limit as $k \rightarrow \infty$ and satisfying

$$.367 < e^{-1} \leq c_k \leq \frac{4}{9} < .445.$$

The number τ^2 can be written

$$\tau^2 = \tan^2 \theta,$$

where θ is the angle between u_0 and x_1 . Thus, it reflects how good an approximation u_0 is to x_1 . In particular, if all the w_i are equal, then $\tau^2 = n - 1$. As τ^2 grows, the approximation becomes poorer and $\tilde{\rho}_k$ decreases. On the other hand, as k increases, the leading coefficient of f_k decreases and $\tilde{\rho}_k$ increases.

To obtain some quantitative results on the behavior of the power method, suppose that we wish to estimate λ_1 to within a tolerance α ; that is, for $0 < \alpha < 1$ we wish the Rayleigh quotient ρ_k to satisfy $\rho_k \geq \alpha \lambda_1$. Then we must choose k so that

$$(5) \quad \frac{c_k \tau^2}{2k + 1} \alpha^{2k+1} + \alpha - 1 < 0.$$

This inequality simplifies if we set

$$x = -(2k + 1) \ln \alpha,$$

so that

$$(6) \quad \frac{e^{-x}}{x} < \frac{\alpha - 1}{c_k \tau^2 \ln \alpha}.$$

The following is a table of values of k satisfying (5) for $\alpha = .9$ and various values of τ^2 .

$\tau^2 + 1$	5	10	15	25	50	75	100	250	500	1000
k	4	5	7	8	10	12	13	16	19	21

As we pointed out earlier, if the components of u_0 along the eigenvectors of A are all equal, then $\tau^2 + 1 = n$. In this case the table gives k as a function of the order of A . What is significant is the extremely slow growth of k ; a tenfold increase in $\tau^2 + 1$ increases k by about eight. The effect is even more marked when α is taken to be 0.5, as it might be when one wants a rough estimate of the magnitude of λ_1 . In this case $k = 2$ for $\tau^2 + 1 = 100$ and $k = 3$ for $\tau^2 + 1 = 1000$.

We may get a crude approximation to k as a function of τ^2 by making the approximation

$$c_k \cong e^{-1}, \quad \frac{1 - \alpha}{\ln \alpha} \cong 1$$

and deleting the $1/x$ term in (6). This gives

$$x > \ln \tau^2 - 1$$

or

$$(7) \quad k > \frac{\ln \tau^2 - 1}{2 |\ln \alpha|} - \frac{1}{2},$$

where \gtrsim indicates that the inequality is only approximate. Thus, the number of iterations required to obtain a given accuracy in λ_1 increases as $\ln \tau/|\ln \alpha|$.

If one wishes to apply the power method with a fixed number of iterations, one must estimate τ^2 , which may not be easy to do. The slow growth of k with τ^2 suggests that underestimating τ^2 , even by an order of magnitude, will not affect things very much; a few extra iterations will wipe out the effect. This view is reinforced by the fact that our analysis assumes the worst possible distribution of eigenvalues.

We can make a probabilistic estimate of τ^2 , if we assume that the components of u_0 are chosen to be independent normal random variables with mean zero and variance σ^2 . In this case $\tau^2/(n-1)$ has an F distribution with $n-1$ and 1 degrees of freedom. It follows that at least 90% of the time, $\tau^2 \leq 64(n-1)$. If the approximation (7) is to be believed, this will add approximately $-2/\ln \alpha$ iterations to those required for $\tau^2 = (n-1)$.

An alternative to fixing the number of iterations, is to terminate the process when ρ_k satisfies a convergence criterion such as $|\rho_k - \rho_{k-1}| \leq (1-\alpha)\rho_k$. The results of this note suggest that if α is not too stringent, say $\alpha \leq .9$, then iteration will stop after a very few iterations.

Department of Computer Science
University of Maryland
College Park, Maryland 20742

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