Exponential Separation of Quantum and Classical One-Way Communication Complexity

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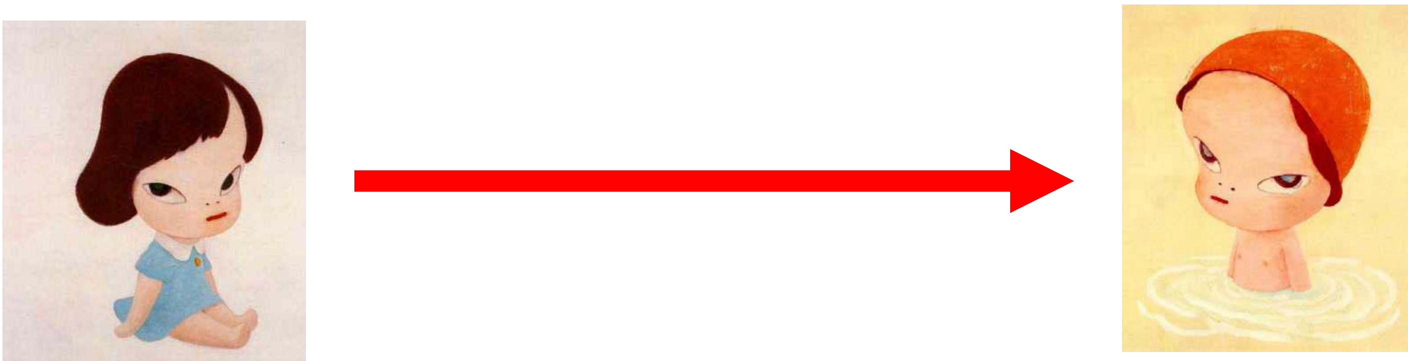
One-Way Communication Complexity



One-Way Communication Complexity

Spot a difference!

One-Way Communication Complexity

Are the images the same?

One-Way Communication Complexity

00111010001101...

Input **x**

Input **y**

Goal: Output P(**x**,**y**)

(minimum communication)

- Applications of Communication Complexity
  VLSI design, Boolean circuits, Data structures, Automata, Formulae size, Data streams, …

- Encoding/Compression scheme C(x), such that P(x,y)=g(C(x),y)

Quantum one-way communication complexity

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Input **y**

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(minimum communication)

Main question:

What is the relation between classical and quantum one-way communication?

Quantum one-way communication complexity

- Holevo's bound
  - We cannot compress information by using qubits.
    We need n qubits to transmit n classical bits.

- [Kremer95] defined a complete problem for boolean promise problems of logarithmic quantum communication complexity.

- [Raz 99] also considers the same problem. He gives an exponential separation for two-way communication.

Quantum one-way communication complexity

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Our result:

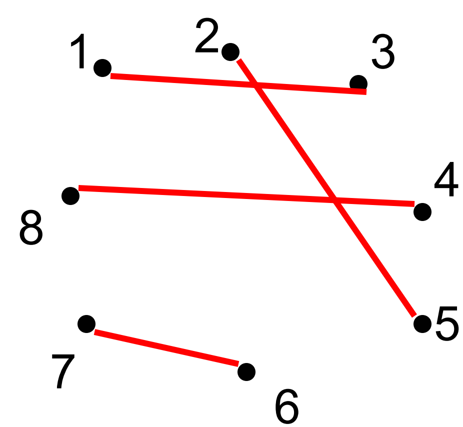
The first exponential separation of classical and quantum one-way communication complexity.

Hidden matching problem HM_n

Input: x 2 $\{0,1\}^n$

Input: a matching M on [n]

e.g. {(1,3),(2,5),(4,8),(6,7)}



Hidden matching problem HM_n

Input: x 2 $\{0,1\}^n$

Output:

$((i, j), x_i \oplus x_j),$
for $(i, j) \in M$

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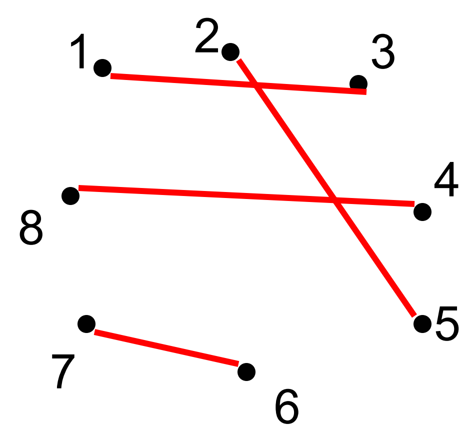

Complexity of HM_n

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$$((i,j), x_i \oplus x_j),$$
$$\text{for } (i,j) \in M$$

Theorem

- There exists a quantum algorithm with complexity $O(\log n)$
- Any randomized algorithm with public coins has complexity $\Omega(\sqrt{n})$

Quantum algorithm for HM_n

Let $M = \{(i_1, i_2), (i_3, i_4), \ldots, (i_{n-1}, i_n)\}$ be Bob's matching.

- Alice sends the state

$$\frac{1}{\sqrt{n}} \sum_{i=1}^{n} (-1)^{x_i} |i\rangle$$

- Bob measures in the basis

$$B = \{|i_1\rangle \pm |i_2\rangle, |i_3\rangle \pm |i_4\rangle, \ldots, |i_{n-1}\rangle \pm |i_n\rangle\}$$

and outputs $\begin{cases} ((j,k), 0) & \text{if he measures } |j\rangle + |k\rangle \\ ((j,k), 1) & \text{if he measures } |j\rangle - |k\rangle \end{cases}$

Quantum algorithm for HM_n

- Alice sends the state

$$\frac{1}{\sqrt{n}}\sum_{i=1}^{n}(-1)^{x_i}|i\rangle = \frac{1}{\sqrt{n}}(((-1)^{x_{i_1}}|i_1\rangle+(-1)^{x_{i_2}}|i_2\rangle)+\ldots+((-1)^{x_{i_{n-1}}}|i_{n-1}\rangle+(-1)^{x_{i_n}}|i_n\rangle))$$

- Bob measures in the basis

$$B = \{|i_1\rangle \pm |i_2\rangle, |i_3\rangle \pm |i_4\rangle, \ldots, |i_{n-1}\rangle \pm |i_n\rangle\}$$

- $Prob[\text{outcome is } |j\rangle+|k\rangle] = \frac{1}{2n}((-1)^{x_j}+(-1)^{x_k})^2$

  $Prob[\text{outcome is } |j\rangle-|k\rangle] = \frac{1}{2n}((-1)^{x_j}-(-1)^{x_k})^2$

- Bob can compute the XOR of a pair of the matching with prob. 1

HM_n and Other problems

Locally Decodable Codes

- The quantum algorithm relies on the property that we can compute efficiently the XOR of a pair of a matching from a uniform superposition.

- Same property was used in [K., deWolf] to prove a lower bound for 2-query Locally Decodable Codes.

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Locally Decodable Codes

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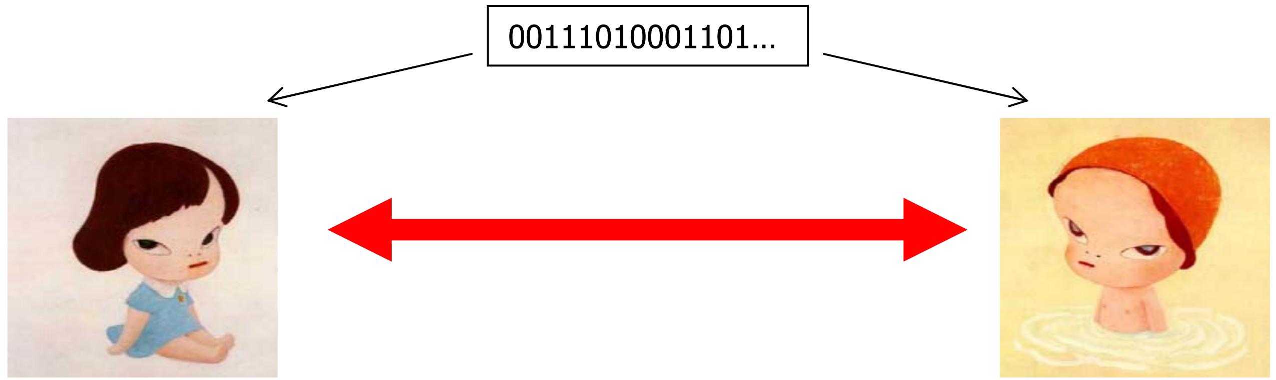
Complete Problems

- We can define a variant of Kremer's problem which is complete for non-boolean promise problems of logarithmic on-way quantum communication complexity.
- Our bounds extend to this problem.

Other models of communication complexity

- Two-way communication



- [Raz99] proved an exponential separation.
- The quantum protocol needs two rounds.

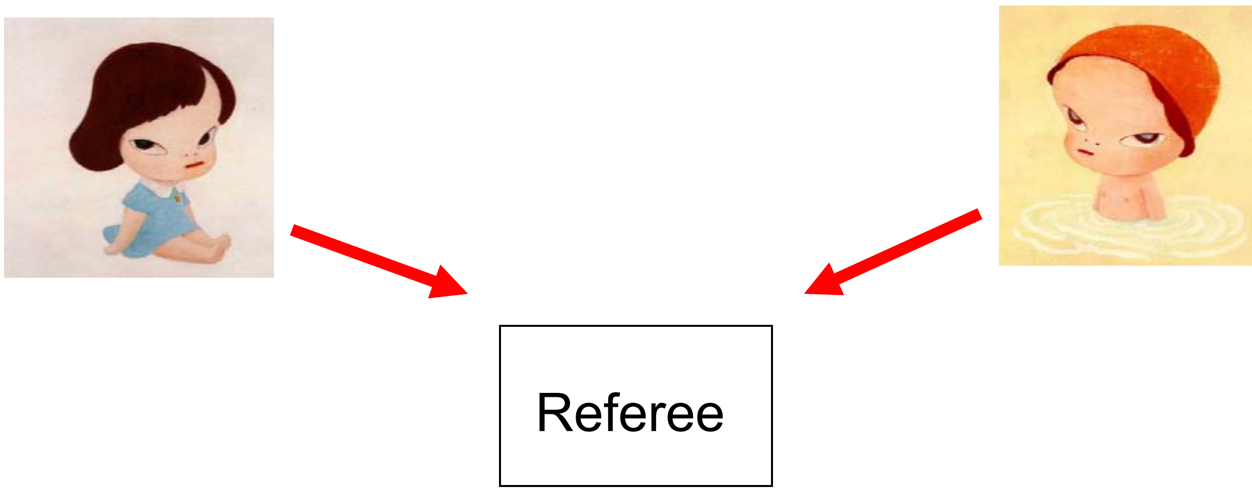
Other models of communication complexity

- Sampling model

- [ASTVW98] proved an exponential separation.
- The separation does not hold with public coins.

Other models of communication complexity

- Simultaneous Messages



- Quantum fingerprints [BCWdW01]
  The separation does not hold with shared public coins

Our problem provides the first exponential separation in the model of Simultaneous Messages with public coins.

Neat application [Harry Buhrman]

Hidden matching Problem as a non-locality game

- Using EPR pairs and NO communication, we can create correlations for which we need exponential classical communication even to approximate them!

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Lower bound for HM_n

- By Yao's Lemma we will construct a "hard" distribution over instances of HM_n and prove a distributional lower bound w.r.t. deterministic one-way protocols.

Lower bound for HM_n

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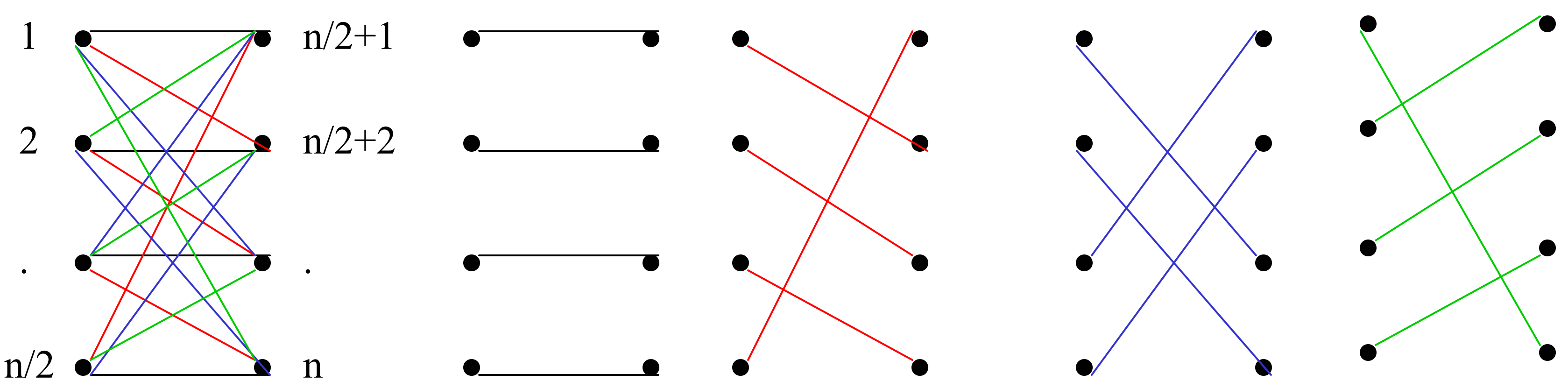
- Distribution of Alice's input: $x\ 2_R \{0,1\}^n$

- Distribution of Bob's input: $M\ 2_R\ M_n$

  M_n is any set of $\Omega(n)$ pairwise edge-disjoint matchings.



Lower bound for HM_n

Intuition :

Alice's message must contain information about at least one edge of each matching (Ω(n) edges).

(e.g. x_1© x_2 , x_2© x_3 , x_1© x_3 , …)

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Idea of Proof :

We prove that Alice cannot send the same message for too many inputs x.

Every matching imposes a linear constraint on x. (e.g. x_1© $x_2 = 0,\ldots$)

There are at least $\Omega(\sqrt{n})$ linearly independent constraints,

hence only $2^{n-\Omega(\sqrt{n})}$ x's can be mapped to the same message.

We need to take care of errors!!!

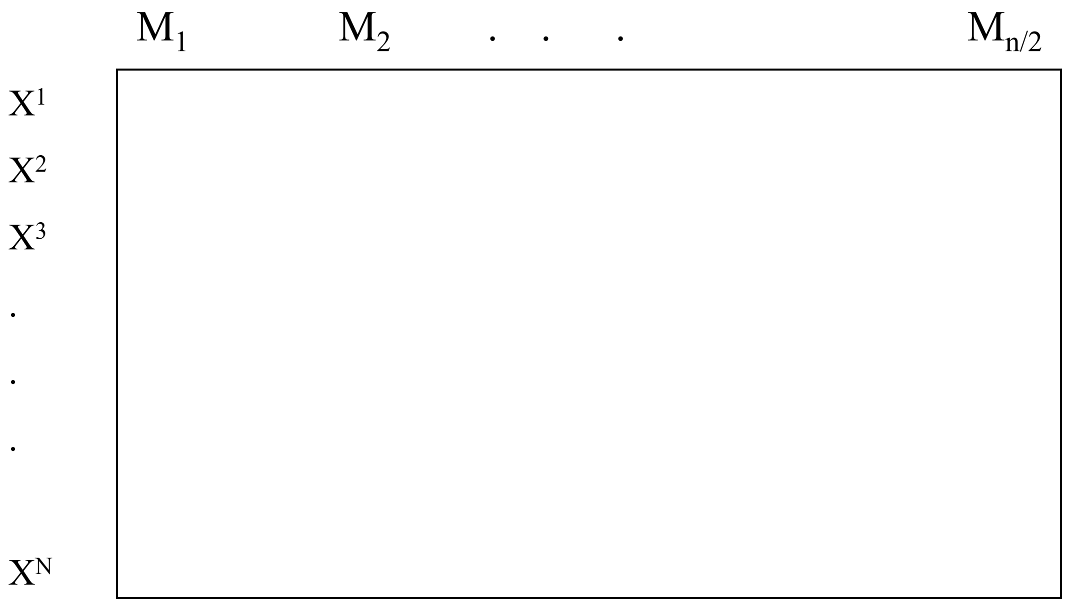
Lower bound for HM_n

$$
\begin{array}{c|ccccc|}
 & M_1 & M_2 & \cdots & & M_{n/2} \\
\hline
X^1 & & & & & \\
X^2 & & & & & \\
X^3 & & & & & \\
\vdots & & & ((i,j)\ ,\ x_i © x_j) & & \\
X^N & & & & & \\
\hline
\end{array}
$$

The Matrix

- At least a (1-δ) fraction of the entries are correct.
  - At least half the columns are (1-2δ)-"good".
  - At least half the rows are (1-2δ)-"good".

Step 1:Pick "good" rows corresponding to the same msg.



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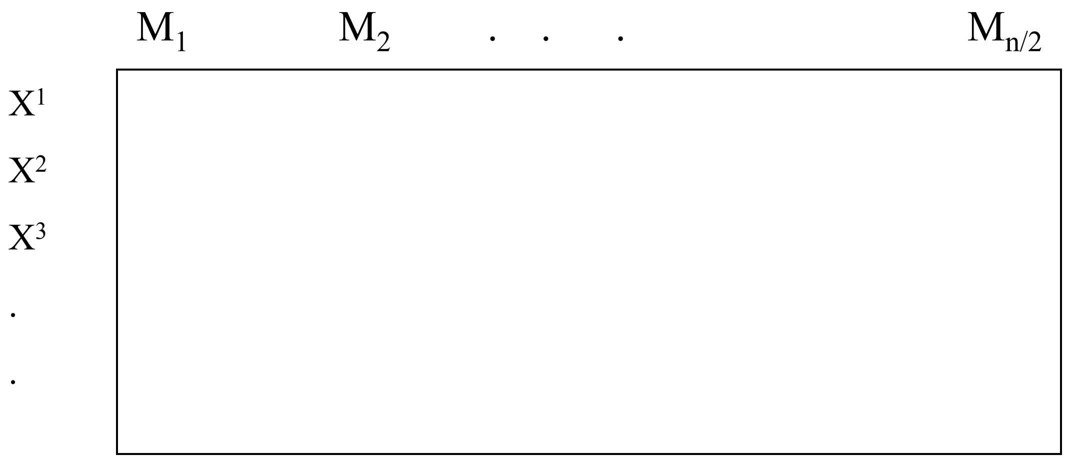

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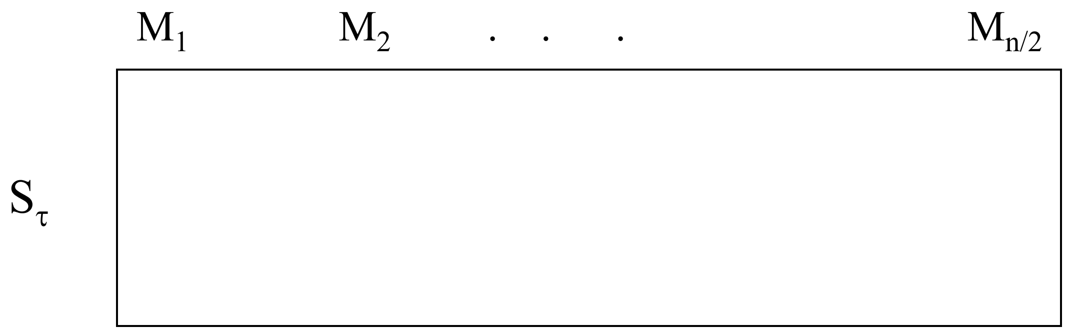
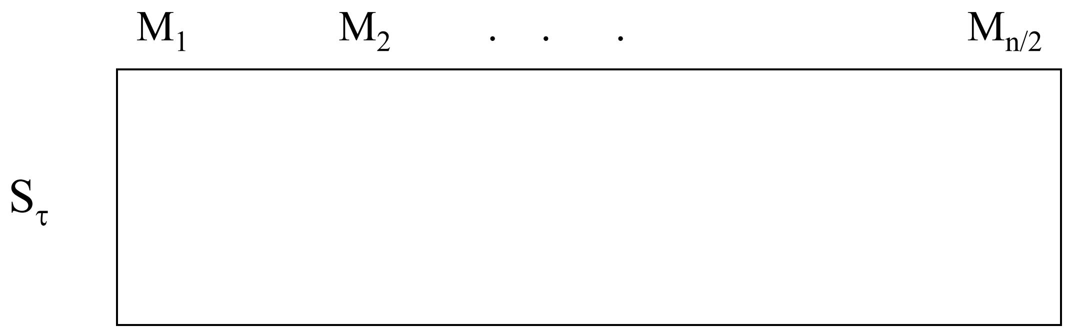


- Each row is $(1-2\delta)$-"good".
- Let S_τ be the set of x's that correspond to the most "popular" message τ.
- Number of Alice's distinct message ¸ $2^n / |S_\tau|$
- I need to bound the size of $|S_\tau|$!

Step 2: Pick "good" , independent columns



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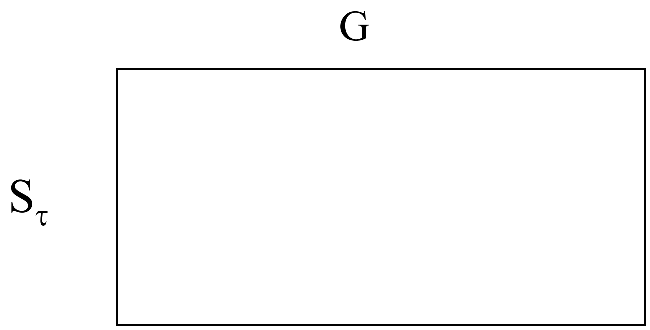

- Each column is (1-4δ)-"good". $|G| = \Omega(n)$
- All the rows of the matrix are the same.
- Each row contains $\Omega(n)$ entries of the form $((i,j),x_i$ © $x_j)$

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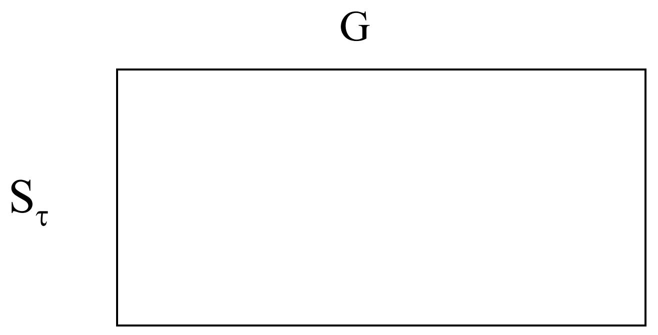

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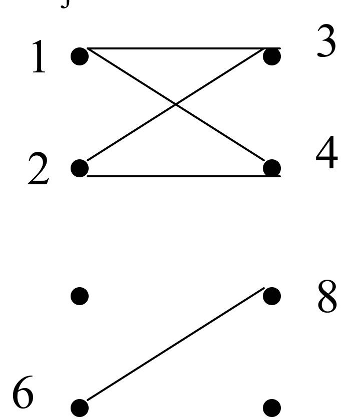
- Define the Graph G.

(e.g. x_1© x_3 , x_1© x_4 , x_2© x_4 , x_2© x_3 , x_6© x_8 …)



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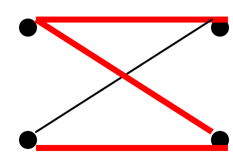
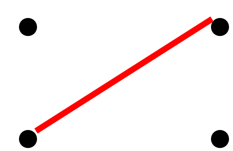
  $(\text{e.g. } x_1 \copyright x_3 ,\ x_1 \copyright x_4 ,\ x_2 \copyright x_4 ,\ x_2 \copyright x_3 ,\ x_6 \copyright x_8 \ \ldots)$



- There are $\Omega(n)$ edges)

  There exists a forest of size $\Omega(\sqrt{n})$



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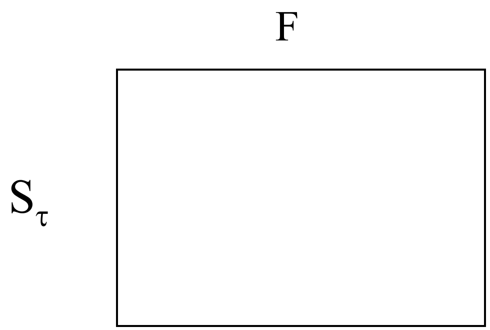
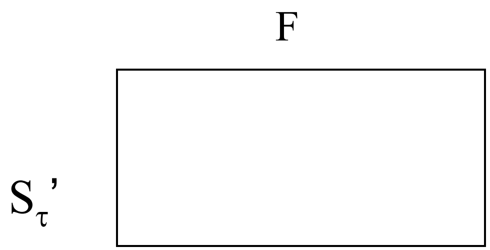


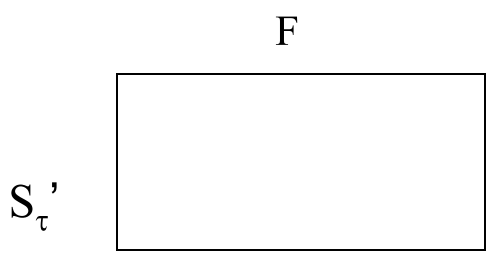
- The columns in F are independent and $|F| = \Omega(\sqrt{n})$

Pick "good" rows again!



- The columns in F are independent and |F|= $\Omega(\sqrt{n})$
- Each row is (1-8δ)-"good"

Lower bound for HM_n



- The rows correspond to inputs mapped to the same message.
- The columns correspond to independent edges, $|F|=\Omega(\sqrt{n})$
- In each row, (1- 8δ) fraction of the entries are correct.

Lower bound for HM_n

- How many x's can be mapped to the same message?
- n variables and a set F of $\Omega(\sqrt{n})$ independent linear constraints.
  - There are $2^{n-\Omega(\sqrt{n})}$ solutions.
- We also need to count all x's that satisfy a set of constraints which agrees with F on at least a (1-8δ) fraction.
  - There are $2^{H_2(8\delta)\Omega(\sqrt{n})}$ such sets of constraints.
- Total number of x's mapped to the same message:

$$|S_\tau| \cdot 2^{n-(1-H_2(8\delta))\Omega(\sqrt{n})}$$

Lower bound for HM_n

- Total number of x's mapped to the same message:

$$|S_\tau| \cdot \; 2^{n-(1-H_2(8\delta))\Omega(\sqrt{n})}$$

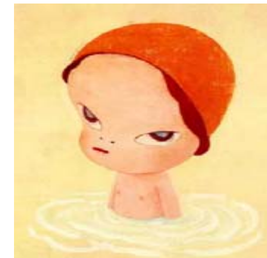
- Size of Alice's message = $\log(2^n / |S_\tau|) = \Omega(\sqrt{n})$

Theorem:

The one-way randomized communication complexity of HM_n is $\Theta(\sqrt{n})$

Upper bound: It's sufficient for Alice to send $O(\sqrt{n})$ random bits of x.

Boolean Hidden Matching Problem



Input: x 2 $\{0,1\}^n$

Output:
$\begin{cases} \text{0 if w is correct} \\ \text{1 if w is wrong} \end{cases}$

Input: a matching M on [n],
w 2 $\{0,1\}^{n/2}$

Theorem

- There exists a quantum algorithm with complexity $O(\log n)$
- Any linear randomized algorithm with public coins has complexity $\Omega(n^{1/3})$

Open problems

- Work in progress
  - Boolean HM_n : extend the lower bound to general randomised protocols.
  - Provide a separation between quantum one-way and classical two-way communication.

- Open problems
  - One-way communication complexity of total functions
  - Simultaneous Messages
  - Quantum advice: BQP/poly vs. BQP/qpoly
  - Quantum proofs: QMA vs. QCMA