

## EXTREMAL PROBLEMS IN DISCRETE GEOMETRY

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Dedicated to Paul Erdős on his seventieth birthday

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In this paper, we establish several theorems involving configurations of points and lines in the Euclidean plane. Our results answer questions and settle conjectures of P. Erdős, G. Purdy, and G. Dirac. The principal result is that there exists an absolute constant  $c_1$  so that when  $\sqrt[n]{n} \cong t \cong \binom{n}{2}$ , the number of incidences between  $n$  points and  $t$  lines is less than  $c_1 n^{2/3} t^{2/3}$ . Using this result, it follows immediately that there exists an absolute constant  $c_2$  so that if  $k \cong \sqrt[n]{n}$ , then the number of lines containing at least  $k$  points is less than  $c_2 n^{2/k^3}$ . We then prove that there exists an absolute constant  $c_3$  so that whenever  $n$  points are placed in the plane not all on the same line, then there is one point on more than  $c_3 n$  of the lines determined by the  $n$  points. Finally, we show that there is an absolute constant  $c_4$  so that there are less than  $\exp(c_4 \sqrt[n]{n})$  sequences  $2 \cong y_1 \cong y_2 \cong \dots \cong y_t$  for which there is a set of  $n$  points and a set  $l_1, l_2, \dots, l_t$  of  $t$  lines so that  $l_j$  contains  $y_j$  points.

## 1. Introduction

Extremal problems in discrete geometry are among the most natural and most tantalizing areas of research in combinatorial mathematics. The reader is encouraged to consult P. Erdős' survey papers [2], [3], [4], [5] and W. Moser's summary [6] of recent progress on a wide range of problems in this area. In this paper, we concentrate on extremal problems in the Euclidean plane. We begin by establishing the following theorem limiting the number of incidences between a set of  $n$  points and a family of  $t$  lines in the Euclidean plane.

**Theorem 1.** *There exists a constant  $c_1$  so that if  $\mathcal{P}$  is a set of  $n$  points and  $\mathcal{L}$  is a family of  $t$  lines in the Euclidean plane, then the number of incidences between points in  $\mathcal{P}$  and lines in  $\mathcal{L}$  is at most  $c_1 n^{2/3} t^{2/3}$  whenever  $\sqrt[n]{n} \cong t \cong \binom{n}{2}$ .*

P. Erdős [5] had conjectured the conclusion of Theorem 1 in the special case  $t = n$ . From Theorem 1, we obtain as an immediate corollary the following theorem settling in the affirmative a conjecture of Erdős and Purdy [5].

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**Theorem 2.** *There exists a constant  $c_2$  so that if  $\mathcal{P}$  is a set of  $n$  points and  $\mathcal{L}$  is a family of  $t$  lines each containing at least  $k$  points from  $\mathcal{P}$  where  $k \cong \sqrt{n}$ , then  $t < c_2 n^2/k^3$ .*

A second corollary is the following result which is a partial solution to Dirac's conjecture [5]. (This result was also proved by J. Beck [1]).

**Theorem 3.** *There exists a constant  $c_3 > 0$  so that if  $\mathcal{P}$  is a set of  $n$  points, not all on the same line, and  $\mathcal{L}$  is the family of all lines determined by  $\mathcal{P}$ , then there exists at least one point in  $\mathcal{P}$  which belongs to more than  $c_3 n$  lines from  $\mathcal{L}$ .*

Let  $\mathcal{P}$  be a set of  $n$  points and let  $\mathcal{L} = \{l_1, l_2, \dots, l_t\}$  be a family of lines each containing at least two points from  $\mathcal{P}$ . For each  $j = 1, 2, \dots, t$ , let  $y_j$  count the number of points from  $\mathcal{P}$  which are on line  $l_j$ . We may then assume that the lines have been labelled so that  $y_1 \cong y_2 \cong \dots \cong y_t$ . Then  $\sum_{i=1}^t \binom{y_i}{2} \cong \binom{n}{2}$ . However, this condition is not sufficient to insure that an arbitrary nondecreasing sequence arises in this fashion. Let  $\mathcal{E}(n)$  count the number of distinct nondecreasing sequences  $y_1 \cong y_2 \cong \dots \cong y_t$  determined by all possible configurations of  $n$  points in the plane. We use Theorem 1 to prove the following result which also settles in the affirmative a conjecture of P. Erdős:

**Theorem 4.** *There exists a constant  $c_4$  so that  $\mathcal{E}(n) < 2^{c_4 \sqrt{n}}$  for all  $n \cong 1$ .*

## 2. The Covering Lemma

The principal tool used in establishing Theorem 1 is a covering lemma proved by the authors in [7]. We assume that a pair of coordinate axes has been chosen and when we use the term *square*, it will also be assumed that the sides are parallel to the coordinate axes. The following lemma asserts the existence of a family of squares covering a positive fraction of a set of  $n$  points with an additional restriction on the number of points contained in each square. We refer the reader to [7] for the proof of the lemma.

**Lemma.** *Let  $r_1, r_2$  be integers with  $r_2 \cong 256 r_1$ , and let  $\mathcal{P}$  be any set of  $n$  points in the plane. Then there exists a family  $\mathcal{Q}$  of squares so that:*

1. *No point in the plane is in the interior of more than one square;*
2. *Each square contains at least  $r_1$  but no more than  $r_2$  points of  $\mathcal{P}$ ;*
3. *At least  $n/16$  of the points in  $\mathcal{P}$  are covered by the squares in  $\mathcal{Q}$ . ■*

The covering lemma will be used to localize the intersection patterns of lines. In order to apply the lemma, we take full advantage of the fact that linear transformations can be applied to any configuration of points and lines. Thus whenever we need to, we can rotate the configuration, and we can increase or decrease the slope of the lines in the configuration.

### 3. The Principal Theorem

In this section we present our principal theorem.

**Theorem 1.** *There exists a constant  $c_1$  so that if  $\mathcal{P}$  is a set of  $n$  points and  $\mathcal{L}$  is a family of  $t$  lines with  $\sqrt{n} \leq t \leq \binom{n}{2}$ , then the number of incidences between the points in  $\mathcal{P}$  and the lines in  $\mathcal{L}$  is less than  $c_1 n^{2/3} t^{2/3}$ .*

**Proof.** We will show that the conclusion holds when  $c_1 = 10^{60}$ . The proof will be by contradiction. We assume the theorem is false and choose a counterexample with  $n+t$  as small as possible. We then show that a large fraction of the points in  $\mathcal{P}$  are incident with a large number of lines from two subsets  $\mathcal{L}_1$  and  $\mathcal{L}_2$ . After a suitable linear transformation, the lines in  $\mathcal{L}_1$  will have slope very nearly  $+1$  and the lines in  $\mathcal{L}_2$  will have slope very nearly  $-1$ . The covering lemma will then be used to show that more pairs of lines cross where there is no point of  $\mathcal{P}$  than there are pairs of lines in  $\mathcal{L}$ .

Choose a coordinate system consisting of a horizontal line ( $x$ -axis) and a vertical line ( $y$ -axis). Then we may assume without loss of generality that all lines have slope in the interval  $(-1/2, 1/2)$ . (If this condition is not satisfied, we apply a linear transformation of the form  $(x, y) \rightarrow (x, \epsilon y)$  where  $\epsilon$  is a small positive number.)

We label the points  $p_1, p_2, \dots, p_n$  and the lines  $l_1, l_2, \dots, l_t$ . Without loss of generality, we may assume that if  $j_1 < j_2$ , then the slope of  $l_{j_1}$  is at least as large as the slope of  $l_{j_2}$ . For each  $i = 1, 2, \dots, n$ , we define the *degree* of  $p_i$ , denoted  $d_i$ , as the number of lines in  $\mathcal{L}$  which contain  $p_i$ . For each  $j = 1, 2, \dots, t$ , we define the *density* of  $l_j$ , denoted  $y_j$ , as the number of points in  $\mathcal{P}$  which are on the line  $l_j$ . We let  $I$  denote the total number of incidences between points and lines. Then

$$c_1 n^{2/3} t^{2/3} \leq I = \sum_{i=1}^n d_i = \sum_{j=1}^t y_j.$$

If  $l_{j_1}$  and  $l_{j_2}$  are lines for which there is some point  $p_i \in \mathcal{P}$  common to both lines, we say  $l_{j_1}$  and  $l_{j_2}$  have a *good* intersection. If  $l_{j_1}$  and  $l_{j_2}$  intersect where there is no point of  $\mathcal{P}$ , we say the lines have a *wasted crossing*. We let  $G$  denote the total number of good intersections and let  $W$  denote the total number of wasted crossings.

Clearly,  $G = \sum_{i=1}^n \binom{d_i}{2}$  and  $W \leq W + G \leq \binom{t}{2}$ . Thus:

$$\begin{aligned} \frac{t^2}{2} &\geq \binom{t}{2} \geq \sum_{i=0}^n \binom{d_i}{2} = \sum_{i=1}^n \frac{1}{2} d_i^2 - \frac{1}{2} \sum_{i=1}^n d_i \\ &\geq \frac{1}{2n} \left( \sum_{i=1}^n d_i \right)^2 - \frac{1}{2} \sum_{i=1}^n d_i = \frac{1}{2n} I^2 - \frac{1}{2} I \geq \frac{1}{3n} I^2 \geq \frac{c_1^2}{3} n^{1/3} t^{4/3}. \end{aligned}$$

In particular, it follows that  $t \geq \frac{c_1^3}{10} \sqrt{n}$ . At this stage, it is important to note that  $t$  is already known to be a reasonable distance away from the boundary value  $t = \sqrt{n}$ .

Let  $y_A$  denote the *average density* of a line in  $\mathcal{L}$ , i.e.,  $y_A=1/t$ . We show that every line has density at least as large as  $.6y_A$ . Suppose to the contrary that there are  $x_1t$  lines which have density less than  $.6y_A$  where  $0 < x_1 < 1$ . Let  $x_1 = \min \{x, 1/2\}$  and choose a set of  $x_1t$  lines each of which has density less than  $.6y_A$ . Since  $\binom{n}{2} > t - x_1t \cong t/2 \cong \sqrt{n}$ , we know that the number of incidences involving the remaining  $t - x_1t$  lines is less than  $c_1n^{2/3}(t - x_1t)^{2/3}$ . Meanwhile, the number of incidences involving the  $x_1t$  lines with small density is less than  $(x_1t)(.6y_A)$ . This requires:

$$(x_1t)(.6y) + c_1n^{2/3}(t - x_1t)^{2/3} > I.$$

Thus there is a number  $x_1$  with  $0 < x_1 \leq 1/2$  for which:

$$(1) \quad .6x_1 + (1 - x_1)^{2/3} > 1.$$

However, it is an easy exercise to show that there is no solution to this inequality in the desired range  $0 < x_1 \leq 1/2$ . The contradiction allows us to conclude that all lines have density at least as large as  $.6y_A$ . Since  $\sum_{j=1}^t \binom{y_j}{2} \leq \binom{n}{2}$ , we also know that

$t \binom{.6y_A}{2} < n^2$ . A simple computation shows that  $t < 10^3 n^2 / c_1^6$ , and thus  $t$  is significantly less than the upper bound  $t = \binom{n}{2}$ .

Also note that if there are  $s$  lines in  $\mathcal{L}$  having the same slope, then  $s(.6y_A) \leq n$ . A simple computation shows  $s < 10^3 t / c_1^3$ , i.e., a very small fraction of the lines in  $\mathcal{L}$  belong to any class of parallel lines.

Let  $d_A$  denote the *average degree* of the points in  $\mathcal{P}$ , i.e.,  $d_A=I/n$ . We will show by a similar argument that all points in  $\mathcal{P}$  have degree at least as large as  $.6d_A$ . To the contrary, suppose there are  $x_1n$  points with degree less than  $.6d_A$  where  $0 < x_1 < 1$ . As before, let  $x_1 = \min \{x, 1/2\}$  and choose a set of  $x_1n$  points having degree less than  $.6d_A$ . Since  $\sqrt{n - x_1n} \leq t \leq \binom{n - x_1n}{2}$ , we know that the number of incidences involving the remaining  $n - x_1n$  points is less than  $c_1(n - x_1n)^{2/3}t^{2/3} = (1 - x_1)^{2/3}I$ . Since the number of incidences involving the  $x_1n$  points with small degree is less than  $(x_1n)(.6d_A)$ , this requires that  $x_1$  satisfy inequality (1) which we have already noted is impossible. The contradiction shows that every point in  $\mathcal{P}$  has degree at least as large as  $.6d_A$ .

The next step in the proof is to obtain a large subset  $\mathcal{P}^*$  of the points so that each point in  $\mathcal{P}^*$  is incident with a large number of lines from each of two sets  $\mathcal{L}_1, \mathcal{L}_2$  of lines. After a suitable linear transformation, all lines in  $\mathcal{L}_1$  will have slope in the interval  $(.99, 1.01)$  and all lines in  $\mathcal{L}_2$  will have slope in the interval  $(-1.01, -.99)$ . The tight control in the angles is necessary in order to apply the covering lemma. We will proceed in several stages.

First, let  $S_1 = \{l_j : 1 \leq j \leq t/2\}$  and let  $S_2 = \{l_j : t/2 < j \leq t\}$ . Then let  $\mathcal{P}_0$  denote the set of points in  $\mathcal{P}$  which are incident with at least  $d_A/100$  lines from  $S_1$  and at least  $d_A/100$  lines from  $S_2$ . Partition  $\mathcal{P} - \mathcal{P}_0$  into two subsets  $\mathcal{P}_1, \mathcal{P}_2$  where  $\mathcal{P}_1 = \{p \in \mathcal{P} - \mathcal{P}_0 : p \text{ is incident with more lines from } S_1 \text{ than with } S_2\}$  and  $\mathcal{P}_2 = \mathcal{P} - \mathcal{P}_0 - \mathcal{P}_1$ . For each  $i=0, 1, 2$ , let  $|\mathcal{P}_i| = x_i n$  where  $x_0 + x_1 + x_2 = 1$ . Also for each  $i$ , let  $I_i$  denote the number of incidences involving the points in  $\mathcal{P}_i$ . We claim that  $x_0 \geq .1$ . To the contrary, suppose that  $x_0 < .1$ . We proceed to a contradiction.

Consider the following inequalities:

$$(2) \quad \sqrt{x_0 n} \cong t \cong \binom{x_0 n}{2}.$$

$$(3) \quad \sqrt{x_1 n} \cong t/2 \cong \binom{x_1 n}{2}.$$

$$(4) \quad \sqrt{x_2 n} \cong t/2 \cong \binom{x_2 n}{2}.$$

If inequalities (2), (3), and (4) are satisfied, then we know:

$$I_0 < c_1(x_0 n)^{2/3} t^{2/3} = x_0^{2/3} I,$$

$$I_1 < (x_1 n)(d_A/100) + c_1(x_1 n)^{2/3}(t/2)^{2/3}, \quad \text{and}$$

$$I_2 < (x_2 n)(d_A/100) + c_1(x_2 n)^{2/3}(t/2)^{2/3}.$$

Since  $x_1^{2/3} + x_2^{2/3} \cong 2[(1-x_0)/2]^{2/3}$ , we conclude that:

$$I = I_0 + I_1 + I_2 < [x_0^{2/3} + (1-x_0)/100 + 2^{-1/3}(1-x_0)^{2/3}] I.$$

Thus  $x_0$  must be a solution to the inequality

$$(5) \quad 1 < x_0^{2/3} + (1-x_0)/100 + 2^{-1/3}(1-x_0)^{2/3}.$$

Again, it is an easy exercise to show that inequality (2) has no solutions in the interval  $(0, .1]$ . (However, it does have a solution in the interval  $(.1, .2)$ .) We may conclude that one or more of the inequalities (2), (3), and (4) does not hold. Since it is assumed that  $x_0 < .1$ , it is clear that at least one of (3) and (4) is valid. In view of the obvious symmetry, we may therefore assume that (4) is valid, and that at least one of (2) and (3) is invalid.

Suppose first that (2) does not hold, i.e., suppose that  $\binom{x_0 n}{2} < t$ . Then define

a number  $y_0$  by  $t = y_0^3 n^2 / 10$ . Then  $\sqrt{y_0 n} \cong t \cong \binom{y_0 n}{2}$  and  $x_0 n < y_0 n$ . Since  $t < 10^3 n^2 / c_1^3$ . Thus  $I_0 < c_1(y_0 n)^{2/3} t^{2/3} < 10^2 I / c_1^2 < 10^{-100} I$ .

Now suppose that (3) is valid. Following reasoning similar to our previous argument, we would conclude that  $I_1 + I_2 < c_1(x_1 n)^{2/3}(t/2)^{2/3} + c_1(x_2 n)^{2/3}(t/2)^{2/3} < 2^{-1/3} I$ . Thus  $I = I_0 + I_1 + I_2 < 10^{-100} I + 2^{-1/3} I$  which is clearly false. It must therefore be the case that when (2) is invalid, (3) is also invalid. However, in this case we know that  $x_1 n < y_0 n$  and thus  $I_1 < 10^{-100} I$ . Also, we know that  $I_2 < c_1 n^{2/3}(t/2)^{2/3} + n d_A / 100 = 2^{-2/3} I + 10^{-2} I$ . Together, these inequalities imply that  $I < 2 \cdot 10^{-100} I + 2^{-2/3} I + 10^{-2} I$  which is false.

So it remains to consider the case where (2) and (4) are valid but (3) is not. In this case, we know  $I_0 < x_0^{2/3} I$ ,  $I_1 < 10^{-100} I$ , and  $I_2 < 2^{-2/3} I + 10^{-2} I$ . These inequalities require that  $x_0$  satisfy the inequality:

$$1 < x_0^{2/3} + 10^{-100} + 2^{-2/3} + 10^{-2}.$$

However, there is no such number. This completes the proof of our claim that  $x_0 \cong .1$ .

Next, let  $m_0$  denote the maximum slope of the lines in  $S_2$  and let  $U_0 = \{l \in S_1: \text{slope}(l) > m_0\}$ . Also, let  $V_0 = S_2$ . Then, each of the points in  $\mathcal{P}_0$  is incident with at least  $d_A/200$  lines in  $U_0$  and at least  $d_A/200$  lines in  $V_0$ . Each line in  $U_0$  has slope greater than every line in  $V_0$ . Rotate the configuration so that the lines in  $U_0$  have positive slope and the lines in  $V_0$  have negative slope. To simplify the inductive construction which follows, we set  $M = 10^{10}$  and let  $n_0 = |\mathcal{P}_0|$ . We then observe that when  $i = 0$ , we have a configuration of points and lines satisfying the following properties.

1. A set  $\mathcal{P}_i$  of at least  $n_i$  points where  $n_i = .1n \left(1 - \frac{2}{M}\right)^i 2^{-i}$ .
2. A set  $U_i$  of lines with positive slope.
3. A set  $V_i$  of lines with negative slope.
4.  $|U_i| \leq t_i$  and  $|V_i| \leq t_i$  where  $t_i = t/2^{i+1}$ .
5. Each point in  $\mathcal{P}_i$  is incident with at least  $e_i$  lines from  $U_i$  and at least  $e_i$  lines from  $V_i$  where  $e_i = \frac{d_A}{200} \left(1 - \frac{2}{M}\right)^i$ .

We say that a point  $p \in \mathcal{P}_i$  is *closed* when it is incident with less than  $e_i/M$  lines from  $U_i$  with slope greater than  $1/2$ , and it is incident with less than  $e_i/M$  lines from  $V_i$  with slope less than  $-1/2$ . Dually, we say that  $p$  is *open* if it is incident with less than  $e_i/M$  lines from  $U_i$  with slope less than  $\sqrt{3}/2$  and it is incident with less than  $e_i/M$  lines from  $V_i$  with slope greater than  $-\sqrt{3}/2$ . We hope to find some linear transformation of the form  $(x, y) \rightarrow (x, my)$  so that there are at least  $n_i/M$  points in  $\mathcal{P}_i$  which are neither open nor closed. If we can find such a transformation, we terminate this inductive construction.

Suppose that there is no transformation which results in at least  $n_i/M$  points which are neither open nor closed. Then we choose a transformation for which the number of open points is as nearly equal to the number of closed points as possible. In this case, there are at least  $(n_i - 2n_i/M)/2$  open points and at least  $(n_i - 2n_i/M)/2$  closed points.

To continue the construction, we will either choose  $\mathcal{P}_{i+1}$  to be the set of open points or the set of closed points. With either choice, we note that  $|\mathcal{P}_{i+1}| \cong n_{i+1} = (.1n) \left(1 - \frac{2}{M}\right)^{i+1} 2^{-i-1}$ . The choice is determined by the distribution of lines. If less than half the lines in  $U_i \cup V_i$  have slope greater than 1 in absolute value, we choose then open points and set  $U_{i+1} = \{l \in U_i: \text{slope}(l) > 1\}$  and  $V_{i+1} = \{l \in V_i: \text{slope}(l) < -1\}$ . On the other hand, if at least half the lines have slope greater than 1 in absolute value, we choose the closed points and set  $U_{i+1} = \{l \in U_i: \text{slope}(l) < 1\}$  and  $V_{i+1} = \{l \in V_i: \text{slope}(l) > -1\}$ . In either case, the desired inequalities hold for  $t_{i+1} = t_i/2$ .

Now we pause to observe that the number of incidences  $J_i$  between the points in  $\mathcal{P}_i$  and the lines in  $U_i$  satisfies  $J_i \geq n_i e_i$ . Now suppose that  $i$  is relatively small, say  $i \leq 30$ . Then  $\left(1 - \frac{2}{M}\right)^i \geq \frac{1}{10}$  and  $2^i \leq M$ . Furthermore, we must have  $t_i \geq \sqrt{n_i}$ . For if  $t_i < \sqrt{n_i}$ , then  $J_i < c_1(n_i)^{2/3}(n_i^{1/2})^{2/3} = c_1 n_i$ . This in turn implies that  $c_1 > e_i = \frac{d_A}{200} \left(1 - \frac{2}{M}\right)^i \geq \frac{d_A}{2000} = \frac{c_1}{2000} n^{-1/3} t^{2/3} \geq \frac{c_1^3}{20\,000}$ , i.e.  $c_1^2 < 2 \cdot 10^4$  which is

false. On the other hand, we must also have  $t_i \cong \binom{n_i}{2}$ . For if  $\binom{n_i}{2} < t_i$ , then  $10^{-18} < \frac{2}{3 \cdot 10^2 \cdot 10^2 \cdot 10^3 2^i} \cong \frac{1}{3} \cdot n^2 (.1)^2 \left(1 - \frac{2}{M}\right)^{2i} 2^{-2i} \frac{2^{i+1}}{n^2} < \frac{1}{3} n_i^2 \frac{2^{i+1}}{t} \frac{10^3}{c_1^6} < \frac{\binom{n_i}{2} 10^3}{t_i} \frac{10^3}{c_1^6}$ . This requires  $c_1^6 < 10^{21}$  which is false.

Since  $\sqrt{n_i} \cong t_i \cong \binom{n_i}{2}$ , we know that  $J_i < c_1(n_i)^{2/3}(t_i)^{2/3}$ , and thus  $n_i e_i < c_1(n_i)^{2/3}(t_i)^{2/3}$ . Substituting, we obtain the inequality:

$$2^{i/3} \left(1 - \frac{2}{M}\right)^{4i/3} < \frac{200}{(.1)^{1/3} 2^{2/3}}$$

A simple calculation shows that this inequality fails if  $i \geq 30$ . So we conclude that our algorithm terminates after  $i$  steps where  $i < 30$ . We obtain a set of at least  $n_i/M$  points which are neither open nor closed.

Divide the quadrant  $[0^\circ, 90^\circ]$  into  $M$  sectors  $\alpha_1, \alpha_2, \dots, \alpha_M$  of equal size. Divide  $[-90^\circ, 0^\circ]$  into  $\beta_1, \beta_2, \dots, \beta_M$  similarly. Then for each point  $p$  which is neither open nor closed, we can choose a pair  $(\alpha_j, \beta_k)$  so that:

1.  $p$  is incident with at least  $e_i/M^2$  lines passing through  $p$  at an angle belonging to sector  $\alpha_j$  and with at least  $e_i/M^2$  lines passing through  $p$  at an angle belonging to sector  $\beta_k$ .
2. The angle between  $\alpha_j$  and  $\beta_k$  is at least  $30^\circ$  and at most  $150^\circ$ .

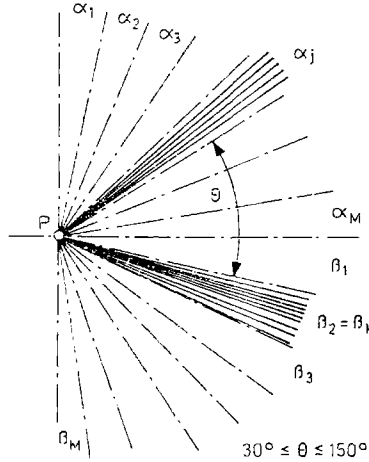


Fig. 1

It follows from the pigeonhole principle, that there is some pair  $(\alpha_{j_0}, \beta_{k_0})$  of sectors so that there is a subset  $\mathcal{P}^*$  containing at least  $n_i/M^3$  points which satisfy these two conditions for  $(\alpha_{j_0}, \beta_{k_0})$ . Rotate the configuration so that the  $x$ -axis is midway between these two sectors and then apply a linear transformation so that the sectors are centered on the lines with slope  $+1$  and  $-1$ . Note that  $e_i/M^2 \cong d_A/M^3$ .

So we now have a set  $\mathcal{P}^*$  of at least  $n/M^3$  points each which is incident with at least  $d_A/M^3$  lines from a set  $\mathcal{L}_1$  of lines each of which have slope in the interval  $(.99, 1.01)$  and with at least  $d_A/M^3$  lines from a set  $\mathcal{L}_2$  of lines each of which has slope in the interval  $(-1.01, -.99)$ .

Now we are ready to apply the covering lemma. We set  $r_2 = d_A/2M^3$  and  $r_1 = d_A/M^4$ . We then cover at least  $n/16M^3$  of the points in  $\mathcal{P}^*$  with squares containing at least  $r_1$  but no more than  $r_2$  points from  $\mathcal{P}^*$ . Now consider a square  $Q$  of this covering.  $Q$  is partitioned into four triangles by its diagonals. At least one of these triangles contains at least  $r_1/4$  points from  $\mathcal{P}^*$ . Choose one such triangle and call it  $T$ . For each point  $p \in T$ , there are at least  $d_A/M^3$  lines from  $\mathcal{L}_1$  which contain  $p$ . At most  $r_2$  of these lines pass through other points of  $\mathcal{P}^*$  which belong to  $Q$ . So there are at least  $d_A/M^3 - r_2 = d_A/2M^3$  lines in  $\mathcal{L}_1$  which contain  $p$  but do not contain any other point of  $\mathcal{P}^*$  inside  $Q$ . We can make a similar observation for the lines from  $\mathcal{L}_2$ .

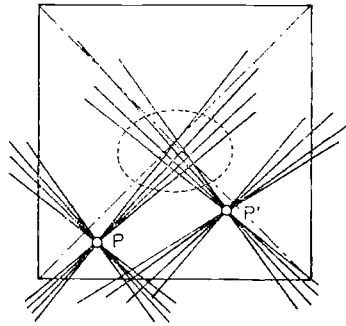


Fig. 2

Now suppose that  $p, p' \in T$  and that these lines not containing other points of  $\mathcal{P}^* \cap Q$  have been chosen for each of the two points. Then regardless of the location of  $p$  and  $p'$  inside  $T$ , there are at least  $d_A^2/4M^6$  wasted crossings produced by these lines.

Thus each  $T$  produces at least  $\binom{r_1/4}{2} d_A^2/4M^6$  wasted crossings. Since there are at least  $n/16M^3 r_2$  squares, we conclude that the total number of wasted crossings, which we denote by  $W$ , satisfies:

$$W \cong \frac{n}{16M^3 r_2} \binom{r_1/4}{2} \frac{d_A^2}{4M^6} > \frac{n}{M^{15}} d_A^3 \cong \frac{n}{M^{15}} c_1^3 n^{-1} t^2 \cong t^2,$$

since  $c_1^3 > 10^{150} = M^{15}$ . The inequality  $w > t^2$  is a contradiction since it implies that there are more wasted crossings than there are pairs of lines. With this observation, the proof of Theorem 1 is complete. ■

Before we close this section we comment that the special case of Theorem 1 when  $t = n$ , namely that the number of incidences between  $n$  points and  $n$  lines is at most  $c_1 n^{4/3}$  was conjectured by Erdős and Purdy [5].



### 4. Consequences of the Principal Theorem

We begin this section with the following result which follows immediately from Theorem 1.

**Theorem 2.** *There exists a constant  $c_2$  so that if  $2 \leq k \leq \sqrt{n}$  and  $\mathcal{L}$  is a family of  $t$  lines each containing at least  $k$  points from a set  $\mathcal{P}$  of  $n$  points, then  $t < c_2 n^2/k^3$ .*

**Proof.** Set  $c_2 = c_1^3$  where  $c_1$  is the absolute constant in Theorem 1. Then suppose that for some  $k$  with  $2 \leq k \leq \sqrt{n}$ , there is a family of  $t = c_1^3 n^2/k^3$  lines each containing at least  $k$  points from a set of  $n$  points. Then the number of incidences is at least  $c_1^3 n^2/k^2$ . Furthermore, since  $k \geq 2$ ,  $t \leq \binom{n}{2}$  and since  $k \leq \sqrt{n}$ ,  $t \geq \sqrt{n}$ . Thus,

$$\frac{c_1^3 n^2}{k^2} < c_1 n^{2/3} t^{2/3} \leq c_1 n^{2/3} \left(\frac{c_1^3 n^2}{k^3}\right)^{2/3} = \frac{c_1^3 n^2}{k^2}.$$

The contradiction completes the proof. ■

P. Erdős conjectured the result in Theorem 2 when  $k = \sqrt{n}$  and this special case was settled by the authors in [7]. Erdős also conjectured Theorem 2 when  $k = \sqrt{n}/2^u$ . In addition, Croft and Erdős also conjectured that for every  $\epsilon > 0$  and every  $k \geq 2$ , there exists a constant  $n(\epsilon, k)$  so that when  $n \geq n(\epsilon, k)$ , the number of lines each containing at least  $k$  points is less than  $\epsilon n^2/k^2$ . Theorem 2 shows that this last conjecture is also valid.

**Theorem 3.** *There exists a constant  $c_3 > 0$  so that if  $\mathcal{P}$  is a set of  $n$  points not all on the same line and  $\mathcal{L}$  is the family of all lines determined by  $\mathcal{P}$ , then there is at least one point in  $\mathcal{P}$  which belongs to more than  $c_3 n$  of the lines in  $\mathcal{L}$ .*

**Proof.** We show that the theorem holds for  $c_3 = 10^{-7} c_2^{-6}$  where  $c_2$  is the absolute constant in Theorem 2. Suppose to the contrary that  $\mathcal{P} = \{p_1, p_2, \dots, p_n\}$  is a set of  $n$  points not all on the same line,  $\mathcal{L} = \{l_1, l_2, \dots, l_t\}$  is the family of all lines determined by, but that no point belongs to more than  $c_3 n$  lines in  $\mathcal{L}$ . For each  $j = 1, 2, \dots, t$ , let  $y_j$  count the number of points on line  $l_j$ . For each  $i = 1, 2, \dots, n$ , let  $d_i$  count the number of lines containing  $p_i$ . Then  $d_i \leq c_3 n$  for each  $i$ . We may then write the following inequalities.

$$\sum_{j=1}^t y_j^2 > \sum_{j=1}^t \binom{y_j}{2} = \binom{n}{2} > \frac{4n^2}{10}, \quad \text{and}$$

$$\sum_{j=1}^t y_j = \sum_{i=1}^n d_i \leq c_3 n^2.$$

Next, let  $M = 100c_3^2$ . Split the sum  $\sum_{j=1}^t y_j^2$  into four terms:

$$S_1 = \sum_{2 \leq y_j < M} y_j^2; \quad S_2 = \sum_{M \leq y_j < \sqrt{n}} y_j^2; \quad S_3 = \sum_{\sqrt{n} \leq y_j < n^{2/3}} y_j^2; \quad \text{and} \quad S_4 = \sum_{n^{2/3} \leq y_j} y_j^2.$$

First, we show that  $S_1 < n^2/10$ . To the contrary suppose that  $S_1 \geq n^2/10$ . For each  $j=2, 3, \dots, M-1$  suppose there are  $s_j$  lines containing  $j$  points. Then  $\sum_{j=2}^{M-1} s_j j^2 \geq \frac{n^2}{10}$ . It follows that there is some  $j$  with  $2 \leq j \leq M-1$  so that  $s_j \geq \frac{n^2}{10M^3}$ . These  $s_j$  lines account for at least  $\frac{n^2}{10M^3} \cdot j > \frac{n^2}{10M^3} \geq \frac{n^2}{10^7 c_2^6} \geq c_3 n^2$  incidences which is a contradiction. Thus  $S_1 < n^2/10$ .

We next show that  $S_2 < n^2/10$ . For each  $i \geq 0$  with  $M2^i < \sqrt{n}$ , let  $t_i$  count the number of lines containing at least  $M2^i$  but less than  $M2^{i+1}$  points. Then  $t_i < c_2 n^2 / M^3 2^{3i}$ . Since these lines contain at most  $M \cdot 2^{i+1}$  points, it follows that  $S_2 < \sum_{i=0}^{\infty} \frac{c^2 n^{3i}}{M^3 2^{3i}} (M \cdot 2^{i+1})^2 = \frac{8c_2 n^2}{M} < \frac{n^2}{10}$ .

Next, we show that  $S_3 < n^2/10$ . Suppose there are  $t_0$  lines containing at least  $\sqrt{n}$  but less than  $n^{2/3}$  points, and that  $t_0 > 0$ . Then  $t_0 \leq c_2 \sqrt{n}$  so  $S_3 < c_2 \sqrt{n} \cdot n^{4/3} = c_2 n^{11/6}$ . We claim that  $c_2 n^{11/6} < n^2/10$ , i.e.,  $10c_2 < n^{1/6}$ . To see that this claim is valid, suppose that  $10c_2 \geq n^{1/6}$ . Then  $10^6 c_2^6 \geq n$  so  $c_3 (10^6 c_2^6) = 1 \geq c_3 n \geq d_i$  for each  $i=1, 2, \dots, n$ . This is a contradiction since each  $d_i$  is a positive integer. This completes the proof of our claim that  $c_2 n^{11/6} < n^2/10$ . We conclude that  $S_3 < n^2/10$ .

Since  $S_1 + S_2 + S_3 + S_4 > 4n^2/10$ , we are left to conclude that  $S_4 > n^2/10$ . For each  $i \geq 0$ , let  $u_i$  count the number of lines which contain at least  $n^{2/3} 2^i$  but less than  $n^{2/3} 2^{i+1}$  points. We claim that  $u_i \leq 2n^{1/3} 2^{-i}$  for each  $i \geq 0$ . To see that this is true, suppose that  $L_1, L_2, \dots, L_{u_i}$  are lines each containing at least  $n^{2/3} 2^i$  points. Then for each  $j=1, 2, \dots, u_i$  there are  $n^{2/3} 2^i - j + 1$  points on  $L_j$  which do not belong to any  $L_k$  with  $1 \leq k < j$ . Therefore  $\sum_{j=0}^{u_i} n^{2/3} 2^i - j + 1 \leq n$ . Summing, we obtain (being generous)  $u_i \leq 2n^{1/3} 2^{-i}$ .

On the other hand, we observe that  $u_i = 0$  whenever  $n^{2/3} 2^i > c_3 n$ . To see that this is true, observe that if there is a line containing more than  $c_3 n$  points, then any point not on the line has more than  $c_3 n$  lines passing through it.

Now let  $\alpha$  be the least integer for which  $n^{2/3} 2^{\alpha+1} > c_3 n$ . Then  $2^\alpha \leq c_3 n^{1/3}$ . It follows that we must have  $\sum_{i=0}^{\alpha} u_i n^{1/3} 2^{2i+\alpha} \leq S_4 > n^2/10$ , and thus  $S_4 \leq \sum_{i=0}^{\alpha} 8 \cdot n^{5/3} 2^i = 8 \cdot n^{5/3} (2^{\alpha+1} - 1) \leq 16 \cdot n^{5/3} 2^\alpha \leq 16 \cdot c_3 n^2$ . This is a contradiction since it requires  $c_3 > 1/160$ . With this observation, the proof is complete. ■

Theorem 3 is a partial solution to an old problem of Dirac [5] who conjectured that there is always a point on at least  $(n/2) - c$  lines. It also shows that  $n$  points not all on the same line determine at least  $c_3 n$  different angles. Corrádi, Hajnal and Erdős conjecture that  $n$  points not on the same line determine at least  $n - 2$  different angles.

Finally, we turn our attention to the following problem. Let  $\mathcal{E}(n)$  denote the number of distinct nondecreasing sequences  $y_1 \leq y_2 \leq \dots \leq y_t$  for which there is a set  $\mathcal{P}$  of  $n$  points and a family  $\mathcal{L} = \{l_1, l_2, \dots, l_t\}$  of  $t$  lines so that  $l_j$  contains  $y_j$  points from  $\mathcal{P}$  for each  $j=1, 2, \dots, t$ .

**Theorem 4.** *There exists a constant  $c_4$  so that  $\mathcal{E}(n) < 2^{c_4 \sqrt{n}}$  for all  $n \geq 1$ .*

**Proof.** We show that the theorem is valid when  $c_4 = 3c_2$  where  $c_2$  is the absolute constant in Theorem 2. To count the number of sequences associated with configura-

tion of points and lines, we first choose the value of  $t$ , the number of lines, as an integer not exceeding  $\binom{n}{2}$ . For each  $i=1, 2, 3, \dots, \log n$ , we then let  $s_i$  denote the number of lines which are to contain at least  $2^i$  but less than  $2^{i+1}$  points. The nonnegative numbers  $s_1, s_2, \dots, s_{\log n}$  sum to  $t$  so the number of different choices for the sequence of  $s_i$ 's is at most  $\binom{t+\log n-1}{\log n-1}$ . For each  $i$ , we choose a sequence  $y_j, y_{j+1}, \dots, y_{j+s_i-1}$  of nondecreasing integers with  $2^i \leq y_j \leq y_{j+s_i-1} < 2^{i+1}$ . Clearly, the number of such sequences is at most  $\binom{s_i+2^i}{2^i}$ . Therefore, we can write:

$$\mathcal{E}(n) \leq \binom{n}{2} \binom{t+\log n-1}{\log n-1} \prod_{i=1}^{\log n} \binom{s_i+2^i}{2^i}.$$

Now

$$\binom{n}{2} \binom{t+\log n-1}{\log n-1} < n^2 \binom{n^2+\log n}{\log n} < n^2 \cdot n^{3 \log n} < 2^{c_2 \sqrt{n}}$$

Now let  $i$  be an integer for which  $2^i < 2\sqrt{n}$ . Then  $s_i < 10c_2 n^2 / 2^{3i}$  and  $10c_2 n^2 / 2^{3i} \leq 100 \cdot 2^i$ . Therefore,

$$\binom{s_i+2^i}{2^i} < \binom{\frac{10c_2 n^2}{2^{3i}} + 2^i}{2^i} < \left( \frac{100c_2 n^2}{2^{4i}} \right)^{2^i}.$$

It follows that

$$\prod_{i: 2^i < 2\sqrt{n}} \binom{s_i+2^i}{2^i} < \prod_{i: 2^i < 2\sqrt{n}} \left( \frac{100c_2 n^2}{2^{4i}} \right)^{2^i} < 2^{c_2 \sqrt{n}}.$$

On the other hand, consider an integer  $i$  for which  $2\sqrt{n} \leq 2^i \leq n$ . In this case, we have the trivial inequality  $s_i < 2n/2^i < 2^i$ . Thus

$$\binom{s_i+2^i}{2^i} < \binom{\frac{2n}{2^i} + 2^i}{\frac{2n}{2^i}} < \left( \frac{100 \cdot 2^{2i}}{n} \right)^{\frac{2n}{2^i}}.$$

It follows that

$$\prod_{i: 2\sqrt{n} \leq 2^i \leq n} \binom{s_i+2^i}{2^i} < \prod_{i: 2\sqrt{n} \leq 2^i \leq n} \left( \frac{100 \cdot 2^{2i}}{n} \right)^{\frac{2n}{2^i}} < \prod_{j=0}^{\infty} (1000 \cdot 2^{2j})^{\sqrt{n} 2^{-j}} < 2^{c_2 \sqrt{n}}.$$

Combining these inequalities, we obtain the desired result  $\mathcal{E}(n) < 2^{3c_2 \sqrt{n}}$ . ■

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