

Flight Gate Assignment with a Quantum Annealer

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Knowledge for Tomorrow



Flight Gate Assignment

A day at Frankfurt Airport

- about 1300 aircraft movements (arrival and departure)
- more than 90% are passenger flights
- more than 170000 passengers
- about 60% transfer passengers
- 278 gates



Passenger Flows

- F, G sets of flights and gates

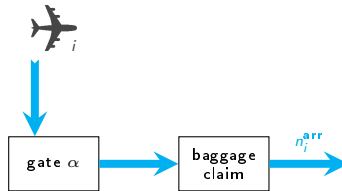


gate α

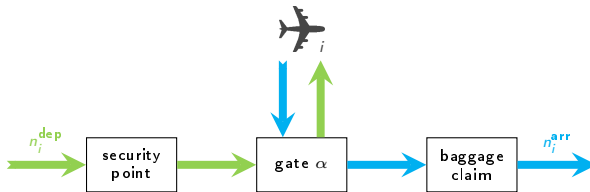


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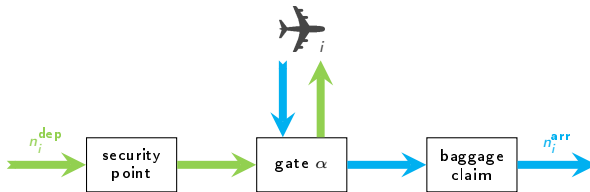
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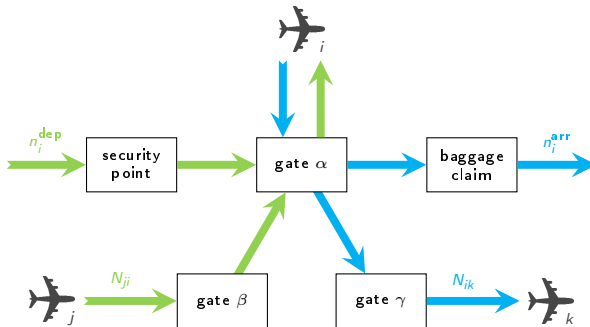
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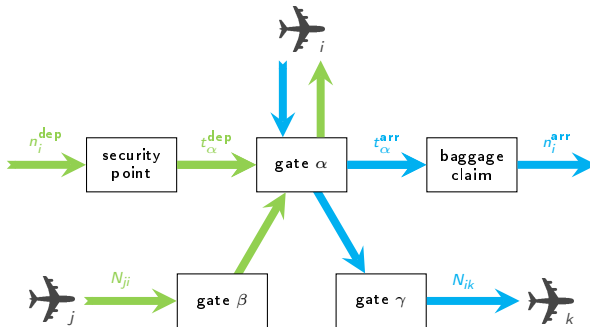
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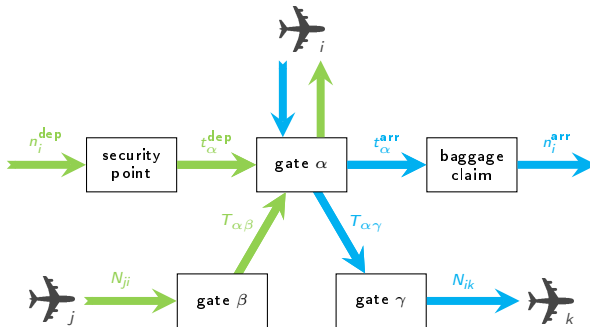
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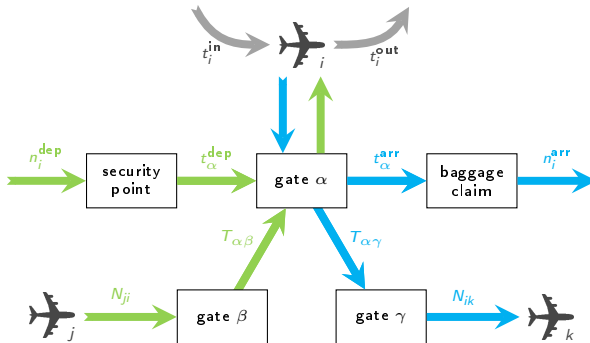
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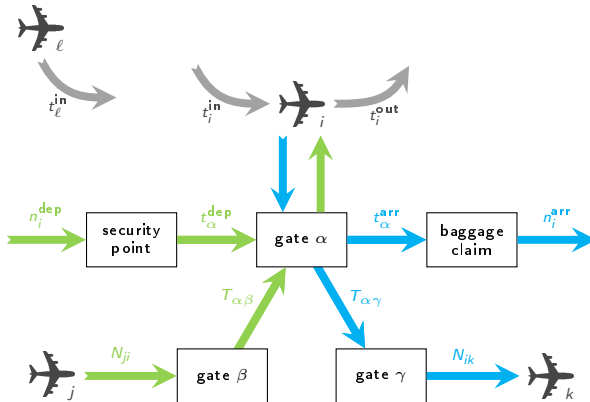
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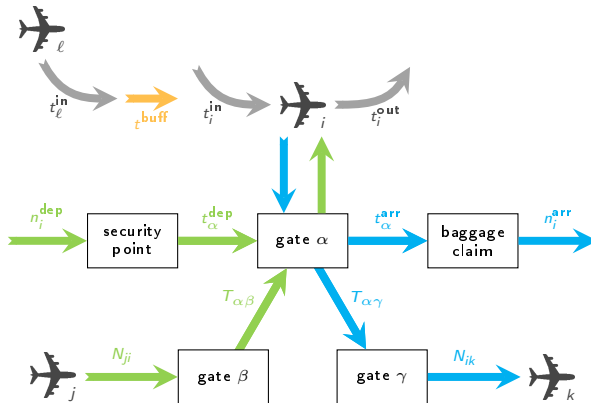
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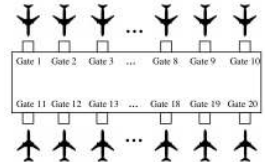


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- t^{buff} buffer time between two flights at the same gate



Question

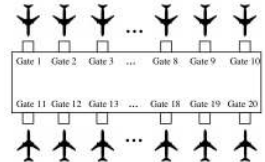
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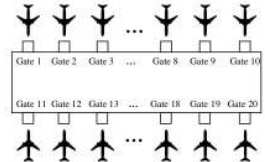


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⇒ Quadratic Assignment Problem



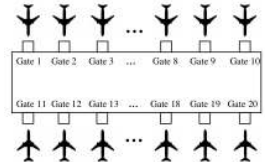
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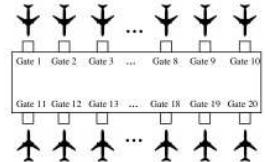
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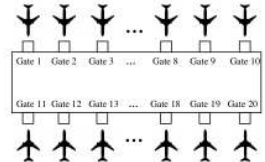
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- fundamental problem in combinatorial optimization
- standard formulation is NP-hard
- seems to exploit possible advantages of the D-Wave machine



FGA Binary Program



Variables $x \in \{0, 1\}^{F \times G}$ with

$$x_{i\alpha} = \begin{cases} 1, & \text{if flight } i \text{ takes gate } \alpha, \\ 0, & \text{otherwise} \end{cases}$$



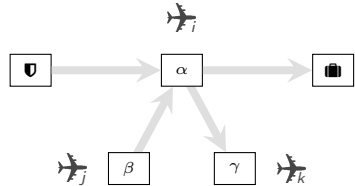
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Minimizing the total transfer time with objective function

$$O(x)$$

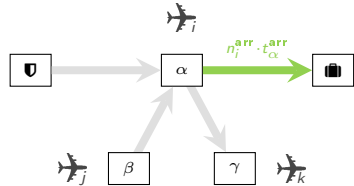


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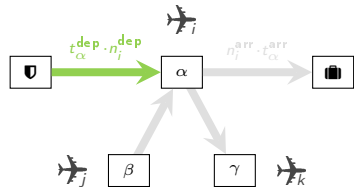


$$O(x) = O_{\text{arr}}(x)$$

$$= \sum_{i\alpha} n_i^{\text{arr}} t_{\alpha}^{\text{arr}} x_{i\alpha}$$



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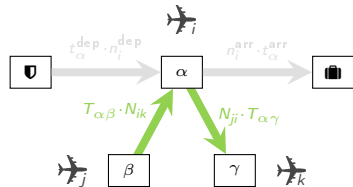
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Minimizing the total transfer time with objective function

$$\begin{aligned} O(x) &= O_{\text{arr}}(x) + O_{\text{dep}}(x) \\ &= \sum_{i\alpha} n_i^{\text{arr}} t_{\alpha}^{\text{arr}} x_{i\alpha} + \sum_{i\alpha} n_i^{\text{dep}} t_{\alpha}^{\text{dep}} x_{i\alpha} \end{aligned}$$



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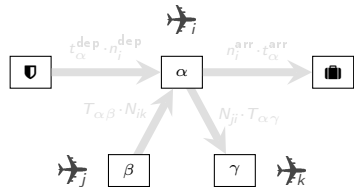
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$$x_{i\alpha} \cdot x_{j\alpha} = 0$$



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$$Q(x) = O(x) + \lambda_{\text{one}} C_{\text{one}}(x) + \lambda_{\text{not}} C_{\text{not}}(x)$$



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Contribution of objective function should not exceed penalty!



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Comparing coefficients in worst cases for

- not assigning a flight to any gate

$$\lambda_{\text{one}} > \max_{i, \alpha} \left(n_i^{\text{dep}} t_{\alpha}^{\text{dep}} + n_i^{\text{arr}} t_{\alpha}^{\text{arr}} + \max_{\beta} T_{\alpha\beta} \sum_j N_{ij} \right)$$



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- assigning a pair of forbidden flights to the same gate

$$\begin{aligned} \lambda_{\text{not}} > \max_{i,\alpha,\gamma} & \left((n_i^{\text{dep}} t_{\alpha}^{\text{dep}} - n_i^{\text{dep}} t_{\gamma}^{\text{dep}}) + (n_i^{\text{arr}} t_{\alpha}^{\text{arr}} - n_i^{\text{arr}} t_{\gamma}^{\text{arr}}) \right. \\ & \left. + \max_{\beta} (T_{\alpha\beta} - T_{\gamma\beta}) \sum_j N_{ij} \right) \end{aligned}$$



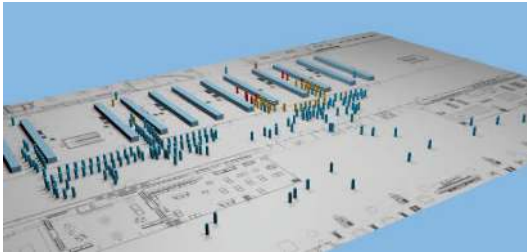
Airport Data

- Flight schedule for one day from a mid-sized European airport



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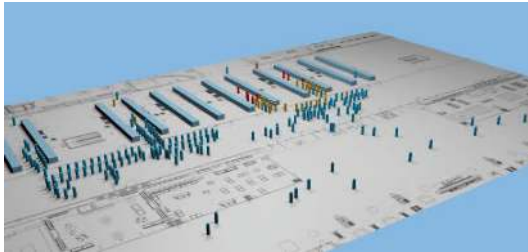
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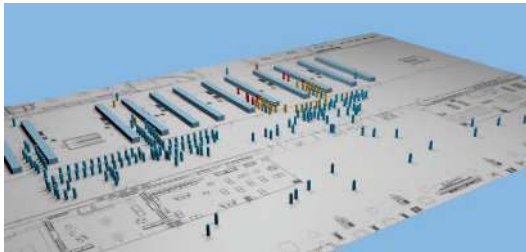
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- Extracted total instance: 293 flights and 97 gates



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Airport Data

- Flight schedule for one day from a mid-sized European airport
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- ⇒ Over 28000 binary variables with about 400 Mio. couplings

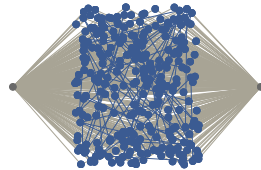


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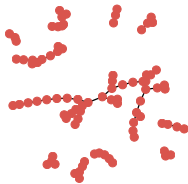
Instance Preprocessing

- Removing corrupted data

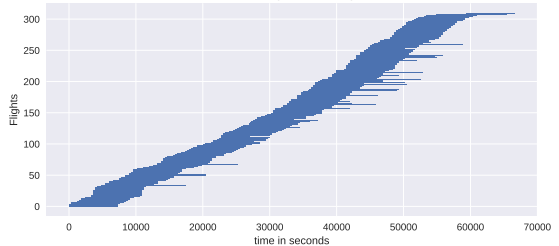
Total problem



Transfer Passengers



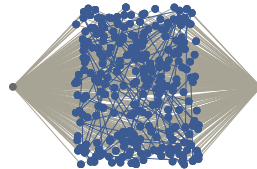
Time at the gates for all flights



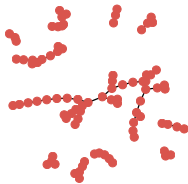
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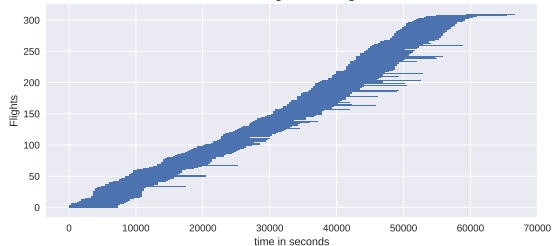
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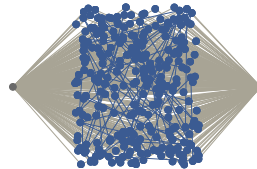
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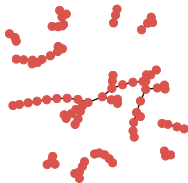
Instance Preprocessing

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- Reducing to flights with transfers only

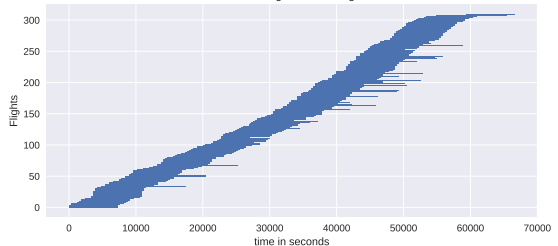
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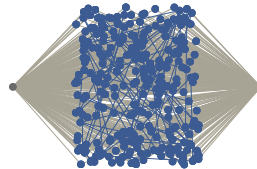
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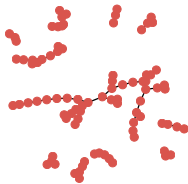
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- ⇒ 89 remaining flights with 80 transfers

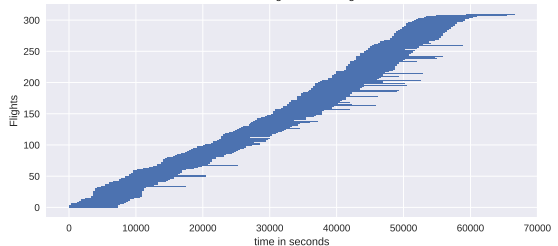
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Transfer Passengers



Time at the gates for all flights



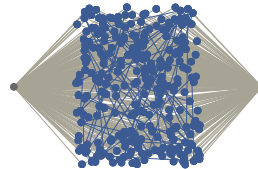
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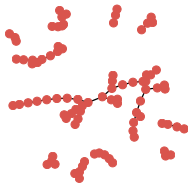
⇒ 89 remaining flights with 80 transfers

⇒ more than 35 gates

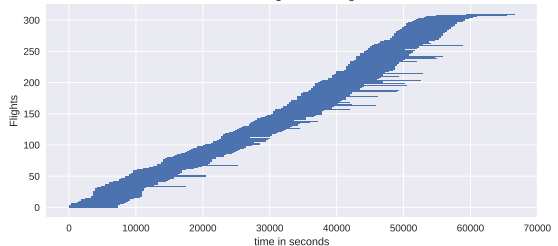
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Transfer Passengers



Time at the gates for all flights

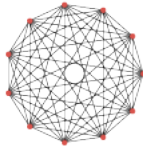


Instance Reduction

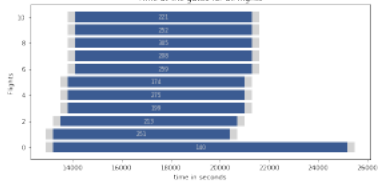
- Slicing by time intervals

Transfer Passengers

Forbidden Flight Pairs

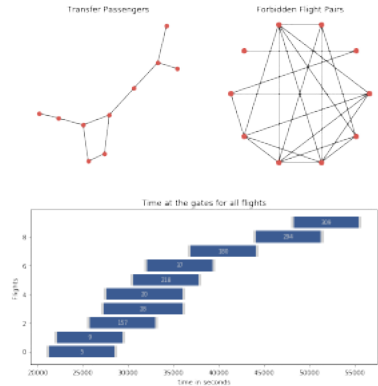
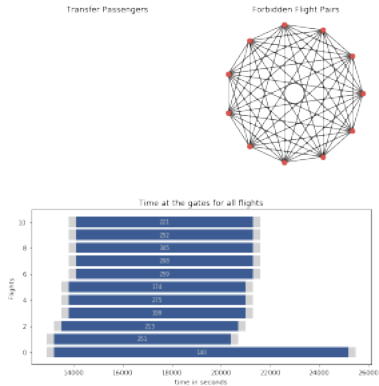


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Instance Reduction

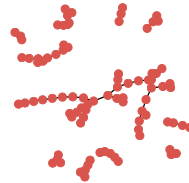
- Slicing by time intervals
- Connected components of transfer passenger graph



Instance Reduction

- Use connected components of transfer passenger graph to find small hard instances

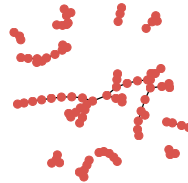
Transfer Passengers



Instance Reduction

- Use connected components of transfer passenger graph to find small hard instances
- Choose number of gates close number of maximal clique

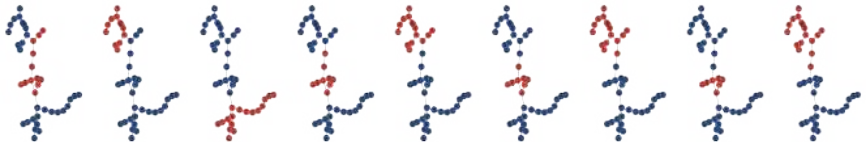
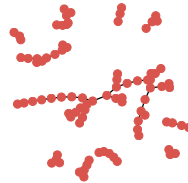
Transfer Passengers



Instance Reduction

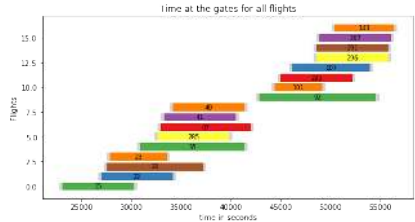
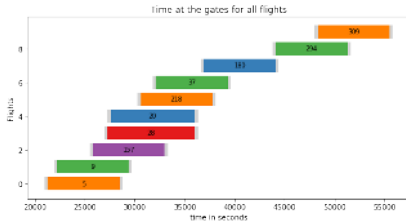
- Use connected components of transfer passenger graph to find small hard instances
- Choose number of gates close number of maximal clique
- Split largest connected component of the transfer passenger graph

Transfer Passengers



Classical Solution

- Getting exact solutions using classical solver SCIP



Embedding

- Standard D-Wave API



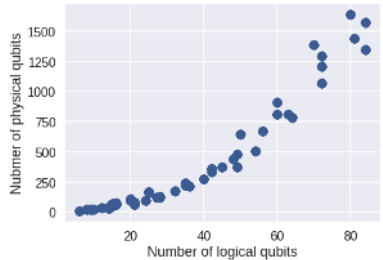
Embedding

- Standard D-Wave API
- 5 embeddings for each instance



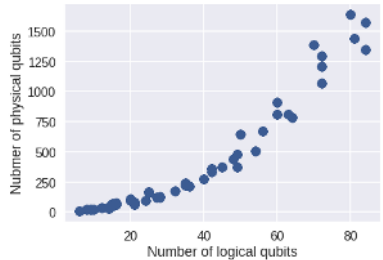
Embedding

- Standard D-Wave API
- 5 embeddings for each instance
- Quadratic overhead



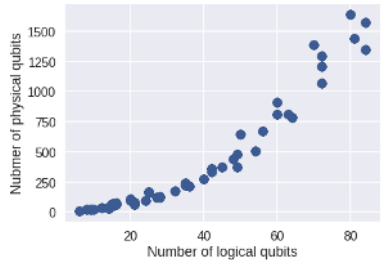
Embedding

- Standard D-Wave API
- 5 embeddings for each instance
- Quadratic overhead
- Maximum 90 logical qubits



Embedding

- Standard D-Wave API
- 5 embeddings for each instance
- Quadratic overhead
- Maximum 90 logical qubits
($\#Variables = \#Flights \cdot \#Gates$)



Annealing Results

- 10000 runs with different parameter settings



Annealing Results

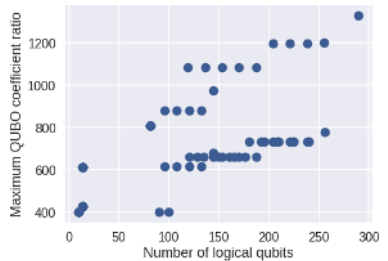
- 10000 runs with different parameter settings
- Even for the small instances:



Annealing Results

- 10000 runs with different parameter settings
- Even for the small instances:

Ratio of maximal to minimal coefficient too large!

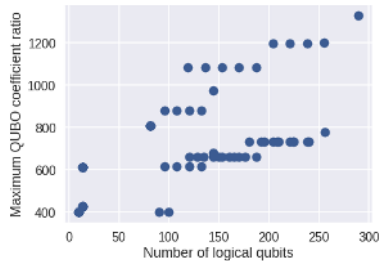


Annealing Results

- 10000 runs with different parameter settings
- Even for the small instances:

Ratio of maximal to minimal coefficient too large!

⇒ Success probability close to 0



Annealing Results

- Running structurally equivalent but random instances with much smaller ratios



Annealing Results

- Running structurally equivalent but random instances with much smaller ratios
- 10000 runs with different parameter settings



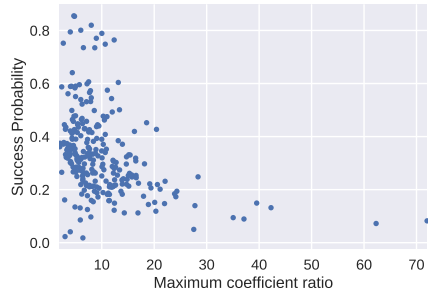
Annealing Results

- Running structurally equivalent but random instances with much smaller ratios
- 10000 runs with different parameter settings
- Differentiated results



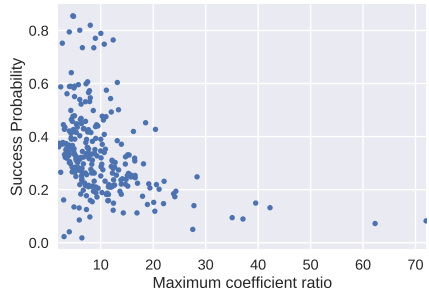
Annealing Results

- Running structurally equivalent but random instances with much smaller ratios
- 10000 runs with different parameter settings
- Differentiated results
 - ratio needs to be < 30



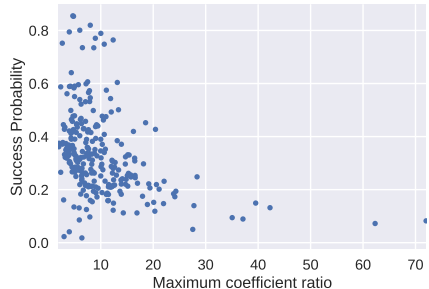
Annealing Results

- Running structurally equivalent but random instances with much smaller ratios
- 10000 runs with different parameter settings
- Differentiated results
 - ratio needs to be < 30
 - wide distribution



Annealing Results

- Running structurally equivalent but random instances with much smaller ratios
- 10000 runs with different parameter settings
- Differentiated results
 - ratio needs to be < 30
 - wide distribution
 - due to other parameters like intra chain coupling of embedding



Outlook

- Adjust embedding/annealing parameters for more clear results



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- Investigate new D-Wave features



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- Using hybrid approach to incorporate classical results



Outlook

- Adjust embedding/annealing parameters for more clear results
- Investigate new D-Wave features
- Using hybrid approach to incorporate classical results
- Using QAOA algorithm (related gate based approach)



Thank You

