

# FPGA Implementation of Single Precision Floating Point Multiplier Using High Speed Compressors

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## ABSTRACT

Floating point multiplier is one of the vital concerns in every digital system. In this paper, the concepts of High speed compressors are used for the implementation of a High speed single precision binary Floating point multiplier by using IEEE 754 standard. Since compressors are special kind of adder which is capable to add more number of bits at a time, the use of these compressors makes the multiplier faster as compared to the conventional multiplier. For Mantissa calculation, a 24x24 bit multiplier has been developed by using these compressors. Owing to these high speed compressors, the proposed multiplier is implemented using Verilog HDL and it is simulated and synthesized for Xilinx 14.3.

**Keywords:** Floating Point, Multiplication, Single Precision, Verilog HDL.

## I. INTRODUCTION

Multipliers are key components of many high performance systems such as FIR filters, microprocessors, digital signal processors, etc. Multiplication based operations such as multiply and accumulate (MAC) and inner product are among some of the frequently used computation-intensive arithmetic functions currently implemented in many digital signal processing (DSP) applications such as convolution, fast Fourier transform (FFT), filtering and in microprocessors in its arithmetic and logic unit. Since multiplication dominates the execution time of most DSP algorithms, so there is a need of high speed multiplier.

Floating point is a way to represent numbers and do arithmetic in computing machines, ranging from simple calculators to computers. The term floating point is derived from the fact that there is no fixed number of digits before and after the decimal point; that is, the decimal point can float. In general, floating-point representations are slower than fixed-

point Representations, but they can handle a larger Range of numbers. The proposed work deals with implementing an architecture for a fast floating point multiplier compliant with the single precision IEEE 754-standard.

The most common representation is defined by the IEEE Standard for Floating-Point Arithmetic (IEEE 754). It is a technical standard established by the Institute of Electrical and Electronics Engineers (IEEE) and the most widely used standard for floating-point computation. Floating Point numbers represented in IEEE 754 format are used in most of the DSP Processors. It also specifies standards for arithmetic operations and rounding algorithms. Floating point arithmetic is useful in applications where a large dynamic range is required or in rapid prototyping applications where the required number range has not been thoroughly investigated.

The Binary Floating point numbers are represented in Single and Double formats. The Single consist of 32

bits and the Double consist of 64 bits. Floating point number consists of three fields:

**1. Sign (S):** It used to denote the sign of the number i.e. 0 represent positive number and 1 represent negativenumber.

**2. Significand or Mantissa (M):** Mantissa is part of a floating point number which represents the magnitude of the number.

**3. Exponent (E):** Exponent is part of the floating point number that represents the number of places that the decimal point (binary point) is to be moved. Number system is completely specified by specifying a suitable base  $\beta$ , significand (mantissa) M, and exponent E. A floating point number F has the value

$$F = (-1)^S M \beta^E$$

The way floating point operations are executed depends on the data format of the operands. IEEE standards specify a set of floating point data formats, single precision and double precision. The Single precision consists of 32 bits and the Double precision consists of 64 bits. Figure 1 shows the IEEE single and double precision data formats.

Sign S	8 bits - Biased Exponent	23 bits - Unsigned Fraction (f)
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(a) IEEE single precision data format

Sign S	11 bits - Biased Exponent	52 bits - Unsigned Fraction (f)
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(b) IEEE double precision data format

**Figure 1.** IEEE Single and Double Precision data Format

The value of the floating point number represented in single precision format is

$$F = (-1)^S 1.f 2^{E-127}$$

where 127 is the value of bias in single precision data format and exponent E ranges between 1 to 254, and E = 0 and E = 255 are reserved for special values. The value of the floating point number represented in double precision data format is

$$F = (-1)^S 1.f 2^{E-1023}$$

Where 1023 is the value of bias in double precision data format. Exponent E ranges between 1 to 2046, the values of E = 0 and E = 2047 are reserved for special values. The performance of multiplier were analyzed using Xilinx ISE simulation tool.

## II. FLOATING POINT MULTIPLIER ALGORITHM

According to IEEE754 standard, the representation of a 32 bit binary floating point number is consists of sign, exponent and mantissa component. During calculation of floating point multiplication different operations are performed on each component. The detail algorithm is described as below

1. Calculation of the sign bit; i.e. SA XOR SB.
2. Exponent is calculated by adding the exponent of EA and EB. After that, bias the addition by 127 to get the final exponent. i.e. EA + EB - 127.
3. Add 1 before the mantissa bit to make the 23 bit into 24 bit after that multiplies together the 24 bit to get 48 bit result.
4. Normalizing the result, to get the required 23 bit mantissa for the final result.
5. Combine the calculated sign, exponent and mantissa components to get the desired multiplication result.

### Example of Floating Point Multiplier

Consider two floating point numbers a = -18.0 and b = +9.5

Expected floating point product = (-18.0) x (+9.5)  
= - 171.0

a = - 10010.0 = - 00010010.0

$= - 1.001000000000000000000000 \times 2^4$   
 $b = +1001.1 = + 00001001.1$   
 $= + 1.001100000000000000000000 \times 2^3$   
 sign of a = 1 = s\_a  
 sign of b = 0 = s\_b  
 biased exponent of a =  $127 + 4 = 131 = 10000011 = e_a$   
 biased exponent of b =  $127 + 3 = 134 = 10000010 = e_b$   
 mantissa of a =  $001000000000000000000000 = \text{mant}_a$   
 mantissa of b =  $001100000000000000000000 = \text{mant}_b$   
 $\text{fp}_a = 1\ 10000011\ 001000000000000000000000$   
 $= \text{C1900000h}$   
 $\text{fp}_b = 0\ 10000010\ 001100000000000000000000$   
 $= 41180000h$

Calculation of sign of the product 's\_out':  
 $s_{\text{out}} = s_a \text{ xor } s_b = 1 \text{ xor } 0 = 1$   
 Calculation of exponent of the product 'e\_out':

Step1: Add e\_a and e\_b to get the sum  
 $10000011 + 10000010 = 1\ 00000101$

Step 2: Bias of 127 is subtracted from the sum to  
 exponent of the output  
 $1\ 00000101 - 01111111 = 10000110 = e_{\text{out}}$

Calculation of mantissa of the product 'm\_out':  
 Step 1: Extract both the mantissas by adding 1 as MSB  
 for normalization to form 24-bit mantissas  
 24-bit mant\_a =  $100100000000000000000000$   
 24-bit mant\_b =  $100110000000000000000000$

Step 2: multiply 24-bit mant\_a and mant\_b to get 48-  
 bit product  
 $(100100000000000000000000) \times$   
 $(100110000000000000000000)$   
 $= 010101011000$   
 $0000$

Step 3 : Leading 1 of the 48-bit is found and the  
 remaining bits are truncated to 23-bit output mantissa  
 value to get the mantissa of the output.  
 $m_{\text{out}} = 010101100000000000000000,$   
 $e_{\text{out}} = 100000110$   
 Floating Point Product(in binary) =  $1\ 10000110$   
 $010101100000000000000000 = \text{C32B0000h}$

biased exponent =  $10000110 = 134$   
 unbiased exponent =  $134 - 127 = 7$   
 Floating Point Product(in decimal)  
 $= - 1.010101100000000000000000 \times 2^7$   
 $= - 10101011.000000000000000000$   
 $= -171.0$

### III. DIFFERENT TYPES OF COMPRESSOR ARCHITECTURE

In binary multiplication a large number of partial  
 product additions are carried out. In conventional  
 multiplier full adders are used for partial product  
 addition. With full adder maximum of 3 inputs can be  
 added at a time. Accordingly to add all the partial  
 products large numbers of full adders are required.  
 Hence, to reduce the number of adders in this  
 implementation High speed compressors are  
 introduced. Different compressors are developed  
 based upon the concept of binary counter property.

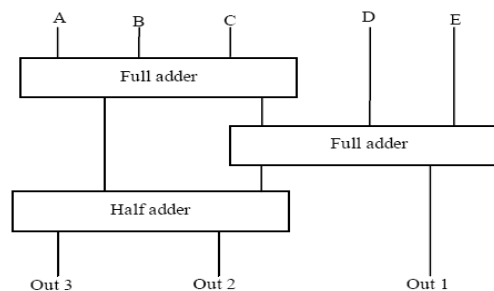
#### A. 5-3 compressor

In 5-3 compressor, maximum of five inputs can be  
 added at a time wherein the output is a three  
 bitnumber. Table-1 shows the counter property of 5-3  
 compressor.

**Table 1.** Counter Property of 5-3 Compressor

Input Condition	Out3	Out2	Out1	Decimal Value
All inputs are zero	0	0	0	0
Any one input is one	0	0	1	1
Any two input are one	0	1	0	2
Any three input are one	0	1	1	3
Any four input are one	1	0	0	4
Any five input are one	1	0	1	5

5-3 compressor is implemented by using two full  
 adders and one half adders as shown in Figure2.



**Figure 2.** Structural Design of 5-3 Compressor

### B. 4-3 compressor

Some basic different compressors architecture are designed .In Figure. 3 shows the block diagram of 4-3 compressor which is designed by using one full adder and two half adder.

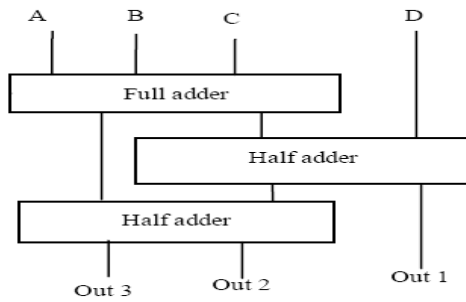


Figure 3. Structural Design of 4-3 Compressor

### C. 6-3 compressor

For the design of 6-3 compressor, two full adders are used. Figure 4 shows the block diagram of 6-3 compressor. For the parallel addition purpose ripple carry adder as seen in Figure 5 is used.

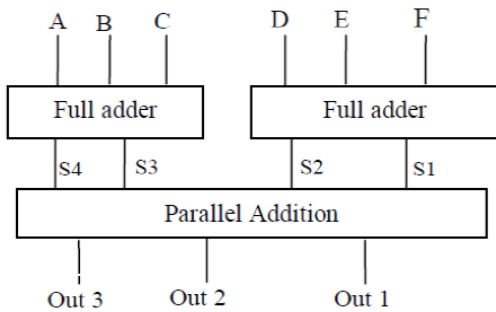


Figure 4. Structural Design of 6-3 Compressor

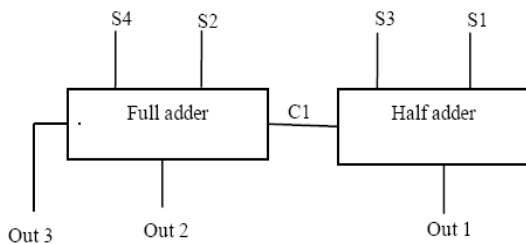


Figure 5. Parallel Addition Unit for 6-3 Compressor

### D. 7-3 compressor

Similarly for the design of 7-3 compressor in Figure 6, one full adder and one 4-3 compressor are used. Its parallel addition unit is given in figure 7.

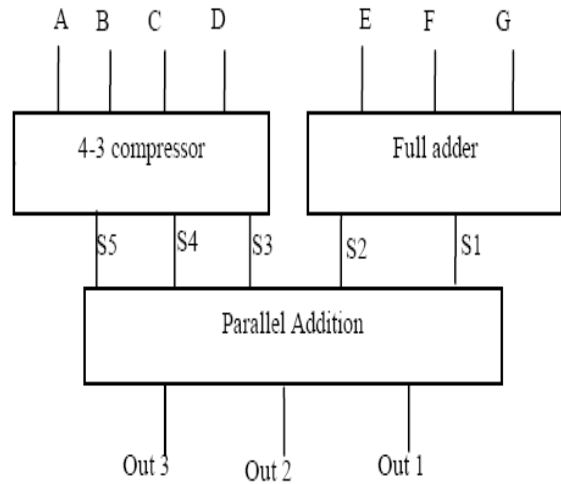


Figure 6. Structural Design of 7-3 Compressor

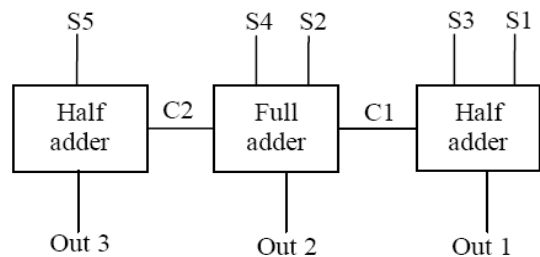


Figure 7. Parallel Addition Unit for 7-3 Compressor

### E. 15-4 compressor

For the implementation of 15-4 compressor five full adders, two compressors, and parallel adder unit are used. Figure.8 shows the block diagram of 15-4 compressor and its optimized parallel adder is shown in figure 9.

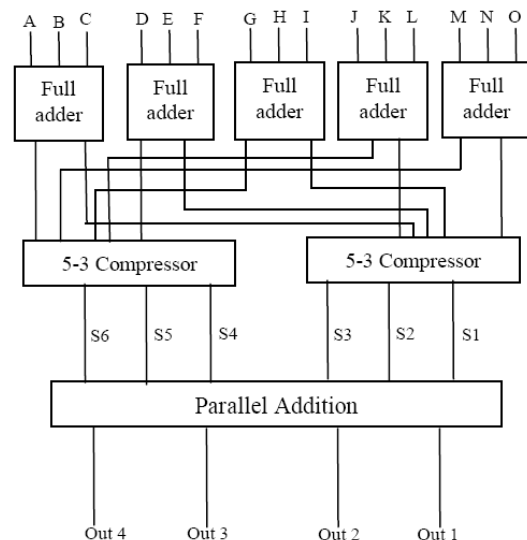
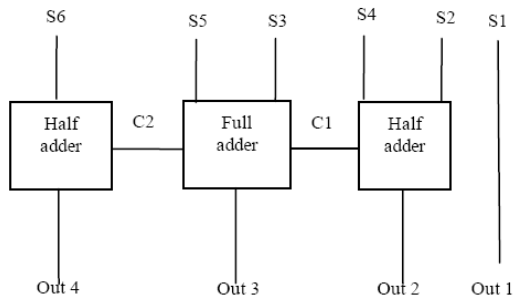


Figure 8. Architecture of 15-4 Compressor

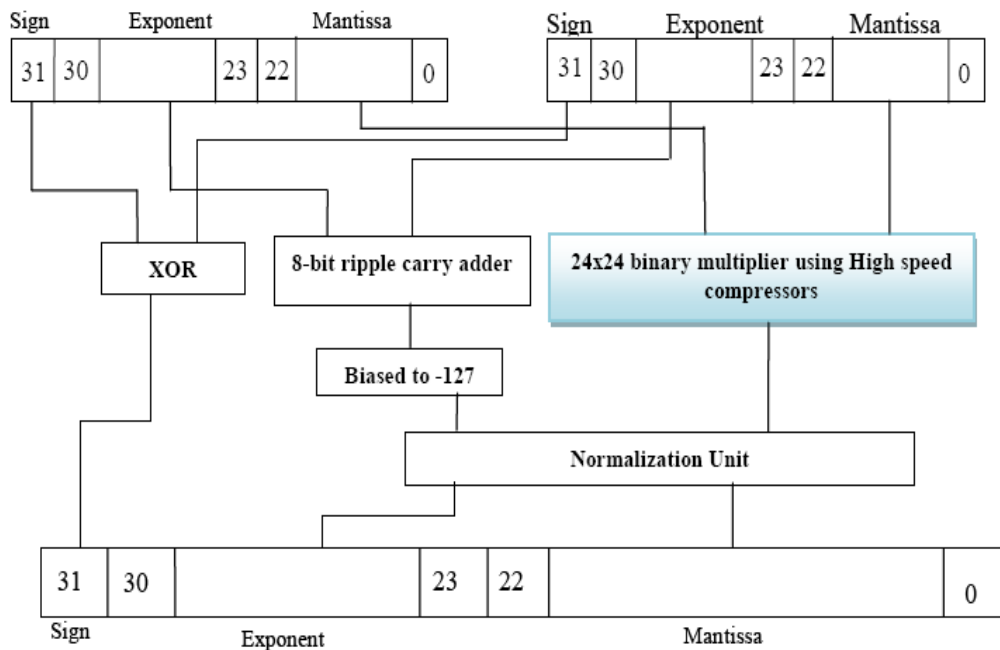


**Figure 9.** Parallel Adder Unit for 15-4 Compressor

By using the concept of adders all the compressors are implemented. These compressors are used in the calculation of partial product addition in different stages of binary multiplication.

#### IV. SINGLE PRECISION FLOATING POINT MULTIPLIER

The proposed architecture for Single Precision Floating Point Multiplier using High Speed Compressors is given in Figure 10.



**Figure 10.** Proposed Architecture of Single Precision Floating Point Multiplier

From the calculation perspective whole floating point multiplication is divided into four sections.

- A. Sign section
- B. Exponent section
- C. Mantissa section
- D. Normalization section

##### A. Sign Section

In sign section, the sign bit for the final result is calculated. The sign bit is calculated by performing XOR operation on the sign bits of the two operands.

##### B. Exponent Section

In this section the two operands are added directly by using 8-bit ripple carry adder. After addition, biasing is required to get the final exponent. For biasing decimal value 127 is subtracted. Here subtraction is done by using 2's complement method.

##### C. Mantissa Section

Mantissa calculation is the most important part of the floating point multiplier. The whole performance of the floating point multiplier is depending upon the mantissa calculation section. In mantissa calculation 24x24 bit binary multiplier is required to do the multiplication of mantissas of two operands. In this

paper binary multiplier is implemented by using high speed compressors, which reduces the delay. In single precision floating point number, mantissa consists of 23 bits. So initially normalization is done in the mantissa by adding 1 at the MSB. After normalization the mantissa of each operand became 24 bits. After binary multiplication of the two normalized 24-bit operands a 48-bit number will generate as a multiplication output. Again this 48-bit number is normalized to 23 bits to get the final mantissa.

In this paper 24-bit multiplier is proposed and implemented by using the concept of high speed compressors. After multiplication of 24x24-bit binary multiplier large numbers of partial products are obtained. Further those partial products are added by using different compressors at different stages. At first stage of partial product addition, no need of any addition because only one partial product is there. Which directly goes as it is as LSB. At the second stage one half adder is required because only two partial products are here. The sum bit taken as output for that stage and carry will go to the next stage. At third stage, partial products are increased to three and one carry from previous stage and hence a total four partial products are present. Addition of these four partial products is done by using 4-3 compressor. In the next stages the number of partial products goes on increasing. So, higher order compressors are required for addition. Similarly in the 4th stage 7-3 compressor, 5th stage 8-4 compressor, 6th stage 9-4 compressor, and 7th stage 10-4 compressor are used. As the stage increases higher order compressors are used. At 24th stage, partial products are increased to 28 including the carry from previous stages. For the partial addition of 24th stage 28-5 compressor is used. After 24th stage number of partial products are decreasing. So for the addition of those partial products previously used compressors are used.

#### D. Normalization Section

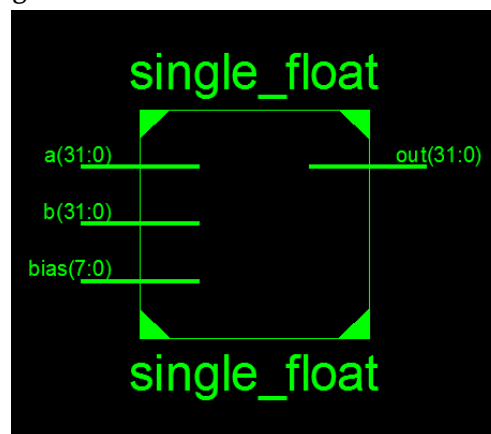
In the Normalization section, normalization of exponent and mantissa are performed. According to the 47th bit (result of the 24x24-bit binary multiplier) normalization is done.

- i. When 47th bit of 24x24-bit binary multiplier is binary one, mantissa is normalized to 23 bits by taking 46th to 24th bit position number and exponent is increased by decimal value one.
- ii. When 47th bit of 24x24-bit binary multiplier is binary zero, mantissa is normalized to 23 bits by taking 45th to 23rd bit position number and there is no increment in the exponent value.

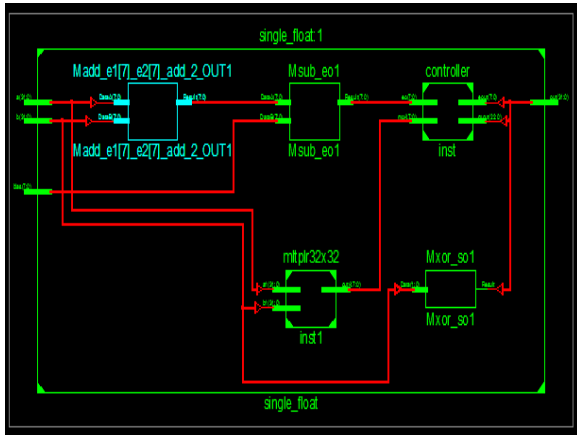
#### V. RESULTS

The program is written in Verilog code for the implemented single precision floating point multiplier and the results are verified. The simulation results and synthesized results are shown in Xilinx 14.3.

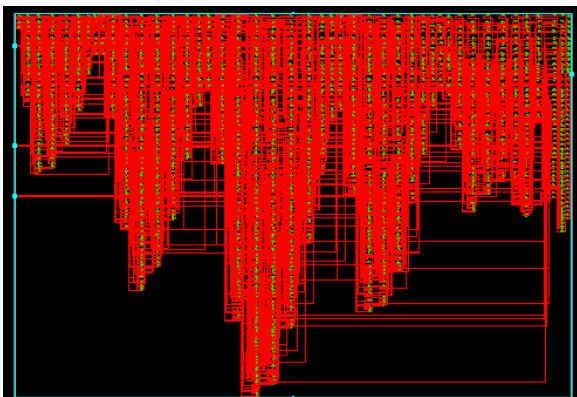
Block diagram



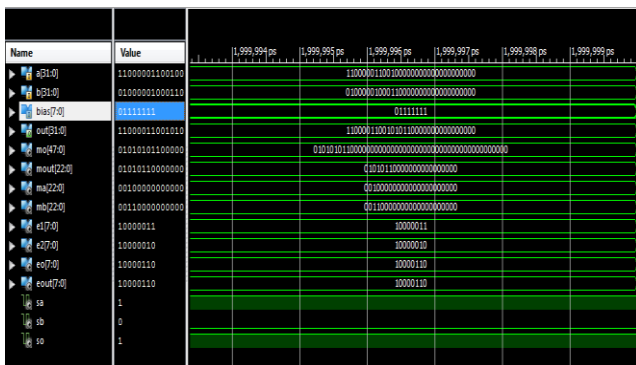
## RTL schematic



## Technology schematic



## Simulation results



## VI. CONCLUSION

In this paper, single precision floating point multiplier based on the IEEE-754 format is successfully implemented on FPGA. The modules are written in Verilog HDL to optimize implementation on FPGA.

## VII. REFERENCES

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