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GAGE: A critical evaluation of genome assemblies and assembly algorithms

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New sequencing technology has dramatically altered the landscape of whole-genome sequencing, allowing scientists to initiate numerous projects to decode the genomes of previously unsequenced organisms. The lowest-cost technology can generate deep coverage of most species, including mammals, in just a few days. The sequence data generated by one of these projects consist of millions or billions of short DNA sequences (reads) that range from 50 to 150 nt in length. These sequences must then be assembled de novo before most genome analyses can begin. Unfortunately, genome assembly remains a very difficult problem, made more difficult by shorter reads and unreliable long-range linking information. In this study, we evaluated several of the leading de novo assembly algorithms on four different short-read data sets, all generated by Illumina sequencers. Our results describe the relative performance of the different assemblers as well as other significant differences in assembly difficulty that appear to be inherent in the genomes themselves. Three overarching conclusions are apparent: first, that data quality, rather than the assembler itself, has a dramatic effect on the quality of an assembled genome; second, that the degree of contiguity of an assembly varies enormously among different assemblers and different genomes; and third, that the correctness of an assembly also varies widely and is not well correlated with statistics on contiguity. To enable others to replicate our results, all of our data and methods are freely available, as are all assemblers used in this study.

[Supplemental material is available for this article.]

The rapidly falling cost of sequencing means that scientists can now attempt whole-genome shotgun (WGS) sequencing of almost any organism, including those whose genomes span billions of base pairs. Interest in genome sequencing of new species has increased rapidly, inspired by high-profile successes such as the panda genome (Li et al. 2010a), the turkey (Dalloul et al. 2010), and several human resequencing efforts (Li et al. 2010b; Schuster et al. 2010; Ju et al. 2011), most of which used reads primarily or exclusively from Illumina sequencers. The read lengths in these projects ranged from 35 to 100 bp, and depth of coverage ranged from 50-fold to 100-fold. In contrast, earlier WGS projects using Sanger sequencing, such as the mouse (Waterston et al. 2002) and dog (Lindblad-Toh et al. 2005) genomes, used read lengths of 750–800 bp and required only sevenfold to 10-fold coverage.

The much deeper coverage of short-read sequencing projects does not entirely compensate for the shorter read length. A side-by-side comparison of the best assemblies produced with short-read data shows that assemblies with longer reads have far better contiguity than the latest short-read assemblies (Gnerre et al. 2011). This illustrates that assembling large genomes from short reads remains a very challenging problem, albeit one that has seen

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considerable progress in just the past two years. Indeed, except for a limited number of specialists in genome assembly, very few scientists know how to optimally design a sequencing strategy and then construct an assembly, and even these experts might not agree. The GAGE (Genome Assembly Gold-standard Evaluations) study was designed to provide a snapshot of how the latest genome assemblers compare on a sample of large-scale next-generation sequencing projects. The study, which was conceived in 2010 in response to the growing use of NGS for de novo assembly and the growing number of genome assembly packages, was designed to help answer questions such as:

- What will an assembly based on short reads look like?
- Which assembly software will produce the best results?
- What parameters should be used when running the software?

As we show below, the answers to these questions depend critically on features of the genome, the design of the sequencing experiments, and on the software used for assembly.

Our results include the full "recipe" that we used for assembling each genome with each assembler. It is important to note in this context that similarly complete instructions are not available for any of the major landmark genomes including human (Lander et al. 2001; Venter et al. 2001) and mouse (Waterston et al. 2002), nor for recently published genomes such as panda (Li et al. 2010a). Whatever the cause, this lack of complete assembly information has made it impossible for others to replicate the assemblies of

major published species. In contrast, we describe all procedures and parameters and provide the complete data sets used for each assembly in our study (see the Supplemental Material). This, coupled with the use of open-source assemblers, should permit replication of any of our results, in contrast with other recent assembly evaluations such as the Assemblathon (Earl et al. 2011) in which the assembly parameters were not described.

We also note that all of the data used in our evaluations were real sequence data from high-throughput sequencing machines, unlike the Assemblathon, which used data from a simulated genome (Earl et al. 2011). Simulated data may not capture the actual patterns of errors in real data or the variability present in naturally occurring genomes.

Results

The data

We chose whole-genome shotgun data from four deep-coverage sequencing projects covering two bacteria, a bee, and the human genome (Table 1). Three of the species were previously sequenced and finished to a very high standard using conventional Sanger technology, and later resequenced using Illumina technology. Having a finished genome allowed us to evaluate the correctness of each assembler on these species. We also included one species for which the "true" assembly is unknown: the bumble bee, Bombus impatiens. This genome is typical of many de novo assembly projects today, where the goal is primarily to create a draft-quality assembly that is the first representative of that species. Correct or not, these assemblies will likely remain for many years as the only reference sequence available.

As Table 1 shows, the four genomes also represent a wide range of genome sizes, from 3 million base pairs (Mb) to 250 Mb (bee) to 3000 Mb (human). For human, however, we used only a single chromosome (chromosome 14) as a representative for the complete genome. We chose to use this smaller sample, just 1/30 of the genome, because some of the assemblers in our comparison would take many weeks to assemble the complete genome, and others would fail entirely. The NGS reads for human derived from a whole-genome sequencing project; we created our data set by first mapping all reads to the genome and then extracting those mapped to chromosome 14 (see Methods).

The Staphylococcus genome has one main chromosome and a small plasmid, while the Rhodobacter genome has two chromosomes and five plasmids. Thus even the bacteria had multiple chromosomes. The read lengths (all Illumina) ranged from 37 to 124 bp.

Table 1. Details of the four next-generation sequence data sets used for the GAGE assembly comparison

Species	S. aureus	R. sphaeroides	Human Chr14	B. impatiens
Size (Mb) Read length Fragment size, Library 1 Number of reads, Library 1 Fragment size, Library 2 Number of reads, Library 2 Fragment size, Library 3 Number of reads, Library 3	2.90 101, 37 180 1,294,104 3500 3,494,070	4.60 101 180 2,050,868 3500 2,050,868	88.29 101 155 36,504,800 2280–2800 22,669,408 35 kb 2,405,064	250 (est.) 124 400 303,118,594 3000–4000 129,118,270 8 kb 65,081,280

The assemblers

We chose eight of the leading genome assemblers, each of which is able to run large, whole-genome assemblies using Illumina-only short read data:

- ABySS (Simpson et al. 2009)
- ALLPATHS-LG (Gnerre et al. 2011)
- Bambus2 (Koren et al. 2011) (http://www.cbcb.umd.edu/software/ bambus).
- CABOG (Miller et al. 2008)
- MSR-CA (http://www.genome.umd.edu/MSR_CA_MANUAL.htm)
- SGA (Simpson and Durbin 2012)
- SOAPdenovo (Li et al. 2010b)
- Velvet (Zerbino and Birney 2008)

All of these are open source assemblers. For each genome and each assembler, we ran multiple assemblies using different parameters until we obtained what appeared to be an optimal or nearoptimal result from that assembler. We used contig and scaffold N50 sizes as the primary metric to determine the best assembly for each program, without consideration of assembly errors. This strategy mimics what is commonly practiced among groups assembling a new genome: the assembly with the largest contigs and scaffolds is usually preferred. Software versions and details of the parameters used for each assembly are given in the Supplemental Material.

Some of these assemblers use a modular design, making it possible to mix and match different modules in different programs. For example, MSR-CA has its own "super-read" module to error-correct high-coverage Illumina reads and extend them into longer reads, which it then processes with modules from CABOG. Bambus2 uses CABOG modules to build contigs and then builds scaffolds from those.

Error correction and data cleaning

One of the most important steps in any assembly, often taking much longer than the assembly itself, is the data cleaning process. WGS data are never perfect, and the various types of errors can cause different problems for different assemblers. High-quality data can produce dramatic differences in the results: for example, one assembly of the Rhodobacter sphaeroides data (using an earlier release of SOAPdenovo) had a contig N50 size of just 233 bp, but after error correction the same assembler achieved a contig N50 of 7793 bp, more than 30 times larger.

Some of the assemblers we ran have their own built-in errorcorrection routines, but we wanted to tease apart the effectiveness of error correction and the assembly algorithms themselves. Therefore, the first step we ran with each of the data sets was an independent error correction method. We allowed assemblers

> that incorporate their own error correction routines to do further corrections in addition to this pre-processing. Abyss, SOAPdenovo, Velvet, and CABOG all produced improved results using error correction provided by a separate program, while the other assemblers were most effective when using their own error correction routines.

> For all data sets, we ran the Quake error corrector (Kelley et al. 2010) to detect and correct sequencing errors. Quake bases its error detection on k-mers that

occur only once or twice in a data set, indicative of a base-calling error. It then tries to replace the lowest-quality base with another base in order to create a *k*-mer that appears to belong to the genome. For most of the data sets, we also ran the ALLPATHS-LG error corrector (Gnerre et al. 2011). Although ALLPATHS-LG is primarily an assembler, we found that use of its corrected reads in some cases led to better assemblies than those based on Quake. Therefore, we extracted the corrected reads from ALLPATHS-LG and used them as another input to all of the assembly algorithms. We ran assemblers using both sets of error-corrected reads and chose the better assembly to report.

For some data sets, additional customized pre-processing was required. For *B. impatiens*, the large insert libraries (3 kb and 8 kb) used an adaptor sequence as part of the library construction protocol. Both libraries had significant numbers of reads that contained adaptor sequences. These adaptors were carefully trimmed out from all reads.

The assemblies

In Tables 2–5, we present snapshots of each assembly using a few metrics: the number, N50 size, and error-corrected sizes of contigs and scaffolds. The N50 value is the size of the smallest contig (or scaffold) such that 50% of the genome is contained in contigs of size N50 or larger. Precise recipes describing how to run each of the assemblers on each of our data sets can be found in the Supplemental Material and at http://gage.cbcb.umd.edu/recipes. These include the parameters used for each assembler as well as the series of steps required to run them, for those assemblers that require multiple steps. If an assembler could not be run on a given data set, then results for that assembler are not included.

Corrected assembly contiguity analysis

It is critical to note here that the statistics in Tables 2–5 can be very misleading if an assembly contains errors; e.g., when two contigs are erroneously concatenated, the resulting assembly has larger contigs, but the assembly is worse, not better. Using the alignments to the reference genomes, we reevaluated the contig sizes for the three finished genomes. For this corrected analysis, we broke contigs at every misjoin and at every indel longer than 5 bases. This produced a revised picture of what the assembly's contiguity statistics would be if every error could be identified and the assembly could be split at that point. Note that errors can be very difficult to find, and assemblies with large numbers of errors present other problems for analysis. To present a more complete picture, Tables 2–4

include the numbers of errors and corrected N50 statistics for each assembler.

Evaluation of assembly accuracy

We assessed the correctness of the assemblies by aligning them to a completed reference genome. Tables 6 and 7 summarize the validation results for the three genomes for which a completed reference is available: *Staphylococcus aureus*, *R. sphaeroides*, and Hs14. A few common assembly problems are readily apparent: many small "chaff" contigs, missing reference sequence, unnecessarily duplicated contigs, repeat compressions, and widespread contig "misjoin" errors. Some of these errors are specific to certain assemblers (e.g., unaligned reference bases), while others are endemic across all of them (e.g., contig misjoins).

For the analysis in Table 6, a "chaff" contig is defined as a single contig <200 bp in length. In many cases, these contigs can be as small as the *k*-mer size used to build the de Bruijn graph (e.g., 36 bp) and are too short to support any further genomic analysis.

One of the more difficult aspects of genome assembly is the estimation of repeat copy numbers. The statistics in Table 6 summarizing duplicated and compressed reference bases illustrate performance of the various assemblers on this task. A duplicated repeat is one that appears in more copies than necessary in the assembly, and a compressed repeat is one that occurs in fewer copies. Interestingly, the duplicated repeats appear to be a preventable problem, one that many of the assemblers handle better than others.

For example, in the *S. aureus* assemblies, ALLPATHS-LG, Bambus2, and SGA all produce only on the order of hundreds of bases in duplications. This may be explained by the tendency of assemblers to output the fewest copies of a repeat that can be explained by the data. In contrast, compressed repeats appear to be a systematic problem with the short-reads assemblers, with all assemblers compressing a significant number of base pairs. Suppression of segmental duplications is a well-known deficiency of modern sequencing and assembly strategies (Kelley and Salzberg 2010).

Single nucleotide polymorphisms (SNPs) and short insertions and deletions (indels), shown in Table 7, also vary by assembler. The number of SNPs and indels varied by an order of magnitude, possibly as a function of the "aggressiveness" of the assembler. An important caveat regarding the human SNPs is that we did not have a true reference for the human sample, NA12878, and this individual genome contains many true SNPs when compared with the human reference genome. However, because we are using a common reference genome and read set, the relative number of

Table 2. Assemblies of S. aureus (genome size 2,872,915)

Assembler		(Contigs		Scaffolds				
	Num	N50 (kb)	Errors	N50 corr. (kb)	Num	N50 (kb)	Errors	N50 corr. (kb)	
ABySS	302	29.2	14	24.8	246	34	1	28	
ALLPATHS-LG	60	96.7	16	66.2	12	1,092	Ö	1,092	
Bambus2	109	50.2	177	16.7	17	1,084	0	1,084	
CABOG			Cou	ld not run: incompatibl	e read length	s in one library		•	
MSR-CA	94	59.2	22	48.2	17	2,412	3	1,022	
SGA	252	4.0	6	4.0	456	208	1	208	
SOAPdenovo	107	288.2	48	62.7	99	332	8	284	
Velvet	162	48.4	28	41.5	45	762	17	126	

The best value for each column is shown in bold. For all assemblies, N50 values are based on the same genome size. The Errors column contains the number of misjoins plus indel errors >5 bp for contigs, and the total number of misjoins for scaffolds. Corrected N50 values were computed after correcting contigs and scaffolds by breaking them at each error. See the evaluation section in the text for details on how errors were identified.

Table 3. Assemblies of *R. sphaeroides* (genome size 4,603,060)

Assembler		Contigs				Scaffolds					
	Num	N50 (kb)	Errors	N50 corr. (kb)	Num	N50 (kb)	Errors	N50 corr. (kb)			
ABySS	1915	5.9	55	4.2	1701	9	3	5			
ALLPATHS-LG	204	42.5	43	34.4	34	3192	0	3192			
Bambus2	177	93.2	368	12.8	92	2439	2	2419			
CABOG	322	20.2	34	17.9	130	66	5	55			
MSR-CA	395	22.1	42	19.1	43	2.976	5	2966			
SGA	3067	4.5	8	2.9	2096	51	0	51			
SOAPdenovo	204	131.7	414	14.3	166	660	3	658			
Velvet	583	15.7	35	14.5	178	353	6	270			

Columns are the same as in Table 2.

SNPs between assemblers should be a valid proxy measure of their single nucleotide errors.

A more aggressive assembler (e.g., SOAPdenovo) is prone to creating more segmental indels as it strives to maximize the lengths of its contigs, while a conservative assembler (e.g., SGA) minimizes errors at the expense of contig size. Interestingly, each assembler has a unique profile of indel error types. Figure 1 shows indel profiles for indels < 100 bp in the ALLPATHS-LG, CABOG, and SOAPdenovo assemblies of human chr14. These plots demonstrate that ALLPATHS-LG and CABOG share a similar error pattern, with the majority of indel errors attributed to misestimation of tandem repeat copy numbers, and a relative balance between compressions and expansions. In contrast, SOAPdenovo shows tandem copy errors with a slight bias toward compressions, in addition to an unusual number of segmental deletions (characterized in Fig. 1 by indels plotted at x > 0 and $y \approx 0$; for more details, see the Supplemental Material). With short reads, tandem repeat length estimation is a notoriously difficult problem—however, many segmental deletions can be avoided with careful use of mate-pair libraries or read threading algorithms.

"Misjoin" errors are perhaps the most harmful type, in that they represent a significant structural error. A misjoin occurs when an assembler incorrectly joins two distant loci of the genome, which most often occurs within a repeat sequence. We have tallied three types of misjoins: (1) inversions, where part of a contig or scaffold is reversed with respect to the true genome; (2) relocations, or rearrangements that move a contig or scaffold within a chromosome; and (3) translocations, or rearrangements between chromosomes. For scaffolds, relocations and indels are grouped together as Reloc/Indel, where an indel error in a scaffold means that a contig (>200 bp in length) has been deleted or inserted

incorrectly. These larger-scale indels are essentially relocations where a contig has been moved. (Note that interchromosomal rearrangements were not possible for our human assembly because only one chromosome was used. Table 7 reports both types of errors under the "Reloc" category, but they are broken out separately in the Supplemental Material.)

One conclusion from our analysis is that no assembler is immune from this type of serious error, and certain assemblers seem to be repeat offenders, while others are consistently more correct. Figure 2 shows a dot plot of the *Rhodobacter* genome as assembled into scaffolds by SOAPdenovo and Velvet. In this example, SOAPdenovo has clearly captured the correct structure of the chromosome and plasmids, and no misjoins are visible at this resolution. However, the Velvet assembly exhibits multiple inversion and relocation errors in the main chromosome. This relative performance is captured in Table 7, where ALLPATHS-LG and SOAPdenovo have the fewest scaffold misjoins (12) and Velvet has the largest (38).

Effect of multiple libraries on assembly

An important question in the design of any whole-genome sequencing experiment is that of the number and sizes of paired-end libraries to use. Creating long-range paired-end libraries can be very helpful for assembly, but the sequencing protocols are much more costly and technically more difficult. With today's technology, paired-end libraries in the 100–300-bp range are the most economical. To evaluate the effect of library variety and size on assembly, we reassembled the *Rhodobacter* genome using the two original libraries plus one additional library, which consisted of 100-bp reads from 210-bp fragments, downloaded from the Sequence Read Archive. The 210-bp library had approximately the same number of reads as

Table 4. Assemblies of human chromosome 14 (ungapped size 88,289,540)

Assembler		C	ontigs		Scaffolds					
	Num	N50 (kb)	Errors	N50 corr. (kb)	Num	N50 (kb)	Errors	N50 corr. (kb)		
ABySS	51,924	2.0	691	2.0	51,301	2.1	9	2		
ALLPATHS-LG	4529	36.5	2659	21.0	225	81,647	45	4702		
Bambus2	13,592	5.9	8677	4.3	1792	324	143	161		
CABOG	3361	45.3	3032	23.7	479	393	597	26		
MSR-CA	30,103	4.9	4316	4.3	1425	893	1068	94		
SGA	56,939	2.7	831	2.7	30,975	83	19	79		
SOAPdenovo	22,689	14.7	5163	7.4	13,502	455	268	214		
Velvet	45,564	2.3	4541	2.1	3,565	1190	9156	27		

Columns are the same as in Table 2.

Table 5. Assemblies of the bumble bee, *B. impatiens* (estimated size 250 Mb)

		Contigs		Scaffolds				
Assembler	Num	N50 (kb)	E-size (kb)	Num	N50 (kb)	E-size (kb)		
ALLPATHS-LG	Co	uld not r	un: incom	patible lil	orary type	es		
CABOG	22,107	23.5	34.2	1191	1125	1367		
MSR-CA	21,885	32.4	46.9	2551	1246	1528		
SGA		Program crashed: cause unclear						
SOAPdenovo Velvet	15,957 Progra	57.1 im crashe	78.2 ed: insuffic	5800 cient men	1374 nory (256	1608 (GB)		

Column headers have the same meanings as in Table 2.

the 180-bp library. We assembled the genome 32 times, using all combinations of two libraries and the short library along with each assembler. The results are shown in Figure 3 and Supplemental Table 2. For ease of comparison, only two statistics are reported: the number of contigs and the (uncorrected) N50 contig size.

For five of the assemblers, the best N50 statistic was obtained with the 180-bp and 3-kb library combination; however, ABySS, SGA, and MSR-CA obtained better results using the 180-bp and 210-bp combination. The MSR-CA result was almost twice as large, suggesting that it was able to extract more contiguity information from the additional coverage provided by the second short fragment library. This result may also suggest that the 3-kb library contained artifacts that reduced its usefulness for some assemblers. We also note that the use of more than two libraries might produce superior results for some assemblers: The

SOAP denovo assembly of the giant panda genome (Li et al. 2010a) used five libraries with fragment sizes ranging from 150 bp to $10~{\rm kb}$.

Discussion

Comparison of assembly size and contiguity

The tables show very large differences in performance among assemblers, as well as variation in the performance of each individual assembler when applied to different genomes. Note that larger contigs are not always correct, and below we take note of some cases where misassembled contigs produced artificially large N50 values. As Table 6 shows, certain assemblers generate chaff contigs in large amounts. For Hs14, for example, SGA outputs more base pairs in chaff contigs than it does for the rest of the assembly. ABySS also has an unusually high quantity of chaff. This can be indicative of the assembler being unable to integrate short repeat structures into larger contigs, or not properly correcting erroneous bases. These problems might create numerous very short, unambiguous paths through the graph. Alternatively, the other assemblers might simply be eliminating short contigs from their output. In either case, though, this problem can easily be addressed by ignoring the chaff contigs.

Coverage of the reference genome can be measured by the percentage of reference bases aligned to any assembled contig. The best assemblers have both a low incidence of chaff and a high coverage of the reference genome. By this metric, ALLPATHS-LG and CABOG perform admirably well on Hs14 with only 0.03% of the assembly in chaff contigs, and only 2.8% and 1.7% of the chromosome (respectively) missing from the assembly. It would

Table 6. Statistics showing bases that failed to align or were present in different copy numbers in the reference genomes and the assemblies of *S. aureus, R. sphaeroides,* and Hs14

Assembler	Assembly size (%)	Chaff size (%)	Unaligned ref bases (%)	Unaligned asm bases (%)	Duplicated ref bases (%)	Compressed ref bases (%)
S. aureus (2.87 M	b)					
ABySS	127.0	66.00	1.37	< 0.01	23.30	0.99
ALĹPATHS-LG	99.9	0.03	0.62	< 0.01	0.03	1.27
Bambus2	98.5	0	1.32	< 0.01	< 0.01	1.29
MSR-CA	99.6	0.02	1.30	< 0.01	0.83	1.01
SGA	98.5	21.38	1.91	< 0.01	0.03	1.30
SOAPdenovo	101.3	0.35	0.38	0.01	1.44	1.41
Velvet	99.2	0.45	0.79	0.03	0.10	1.28
R. sphaeroides (4.6	60 Mb)					
ABySS	108.0	1.65	3.01	0.15	10.04	0.04
ALLPATHS-LG	99.7	0.01	0.47	0.01	0.38	0.30
Bambus2	94.9	0	4.93	< 0.01	< 0.01	0.24
CABOG	92.1	< 0.01	7.51	0.01	0.12	0.70
MSR-CA	96.9	0.02	3.52	0.04	1.04	0.49
SGA	97.8	4.95	2.31	0.02	0.06	0.92
SOAPdenovo	99.9	0.45	0.88	0.02	1.07	0.51
Velvet	97.8	0.54	1.60	0.01	0.29	0.92
Human chromoso	me 14 (88.29 Mb)					
ABySS	83.1	41.37	17.78	0.03	0.59	0.52
ALĹPATHS-LG	95.6	0.03	2.76	0.03	0.27	2.57
Bambus2	77.3	< 0.01	20.55	0.07	0.14	4.04
CABOG	97.7	0.03	1.68	0.06	0.16	1.71
MSR-CA	92.5	0.18	8.10	0.57	1.69	2.27
SGA	93.3	107.82	6.97	0.06	0.14	2.14
SOAPdenovo	104.9	3.77	1.83	0.60	6.76	3.76
Velvet	84.7	6.25	15.12	0.31	0.09	0.64

The true size of each genome is shown next to the species name. All table values are expressed as a percentage of the true genome size. Column headers are defined in the main text. Additional statistics are provided in the Supplemental Material.

Table 7. Statistics on insertions, deletions, and misassembly errors in the various assemblies of S. aureus, R. sphaeroides, and Hs14

		Ind	els		Contigs		Scaffolds			
Assembler	SNPs	≤5 bp	>5 bp	Misjoins	Inv	Reloc	Misjoins	Inv	Reloc/indel	
S. aureus (2.87 Mb)										
ABySS	258	20	9	5	3	2	1	1	0	
ALĹPATHS-LG	79	4	12	4	0	4	0	0	0	
Bambus2	28	56	164	13	2	11	0	0	0	
MSR-CA	191	23	10	12	6	6	3	3	0	
SGA	32	2	2	4	1	3	_	_	_	
SOAPdenovo	246	25	31	17	1	16	8	1	7	
Velvet	217	6	14	14	5	9	17	5	12	
R. sphaeroides (4.60	Mb)									
ABySS	692	288	34	21	2	19	3	0	3	
ALĹPATHS-LG	218	150	37	6	0	6	0	0	0	
Bambus2	189	149	363	5	0	5	2	0	2	
CABOG	536	145	24	10	1	9	5	4	1	
MSR-CA	807	179	32	10	1	9	5	2	3	
SGA	336	116	4	4	0	4	-	-	-	
SOAPdenovo	527	155	406	8	0	8	3	1	2	
Velvet	413	148	27	8	0	8	6	6	7	
Human chromosom	e 14 (88.29 Mb)									
ABySS	60,408	9987	678	13	6	7	9	9	0	
ALĹPATHS-LG	55,317	27,559	2558	101	44	57	45	0	45	
Bambus2	64,869	17,141	5411	3266	1722	1544	143	37	106	
CABOG	81,125	28,420	2883	149	46	103	597	389	208	
MSR-CA	153,104	21,933	3082	1234	653	581	1068	210	858	
SGA	70,976	15,483	681	150	90	60	_	_	_	
SOAPdenovo	98,185	21,347	3902	1261	520	741	268	17	251	
Velvet	79,399	17,505	4172	369	199	170	9156	3824	5332	

Column headers are defined in the main text.

appear that these assemblers are able to resolve the complex repeat structure of the human genome by a combination of accurate error correction and good use of mate-pair information. Despite its performance on Hs14, however, CABOG leaves more of R. sphaeroides uncovered (7.5%) than any other assembler.

To provide a context, it is also worth considering whether some genomes are intrinsically more difficult to assemble than the others. Assembly difficulty is partly a function of repetitiveness, which also interacts with read length: In general, a repeat sequence creates a gap unless the reads fully contain (and are longer than) the repeat. Assemblers can fill in many of these gaps using pairedend information, as long as the paired-end distances are longer than the repeats. One measure of repetitiveness is K-mer uniqueness (Schatz et al. 2010), defined as the percentage of a genome that is covered by unique sequences of length K. We computed this ratio for the three known genomes in our study and compared it with the full human genome and the nematode Caenorhabditis elegans (Fig. 4). As the figure shows, the two bacteria are less repetitive than Hs14, and Hs14 is noticeably less repetitive than the full human genome.

Importance of error correction

For all four genomes and for all eight assemblers used in GAGE, the best assemblies were created from reads that had been processed through extensive error correction routines. As noted above, contig sizes after error correction often increased dramatically, as much as 30-fold. This highlights the critical importance of data quality to a good assembly. For most of the assemblers, the best results came from using reads that had been corrected either by Quake or by ALLPATHS-LG (for details, see the Supplemental Material). MSR-CA and SGA produced better results using their own built-in error correction routines, but in all cases, error correction was a key part of the assembly process.

S. aureus

Table 2 shows that SOAPdenovo produced much larger contigs for S. aureus than any of the other systems, with an N50 size of 288 kb.

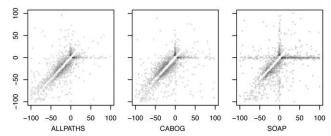


Figure 1. Comparison of the indel profiles for three assemblies of human Chr14. Every indel in the assembly is defined by the two aligned segments on either side. For each indel, the x-axis displays the distance between the two adjacent segments in the reference, and the y-axis displays the distance in the query. Thus, the point x = 100, y = 0 indicates a 100-bp deletion in the assembly, relative to the reference. Deletions from the assembly lie *below* the line y = x, and insertions in the assembly lie *above*. The indels can be roughly categorized by quadrant: (*top right*) divergent sequence; (bottom right) segmental assembly deletion; (bottom left) tandem repeat collapse/expansion; (top left) segmental assembly insertion. No points lie on the line y = x because only indels >5 bp are displayed. For details, see the Supplemental Methods.

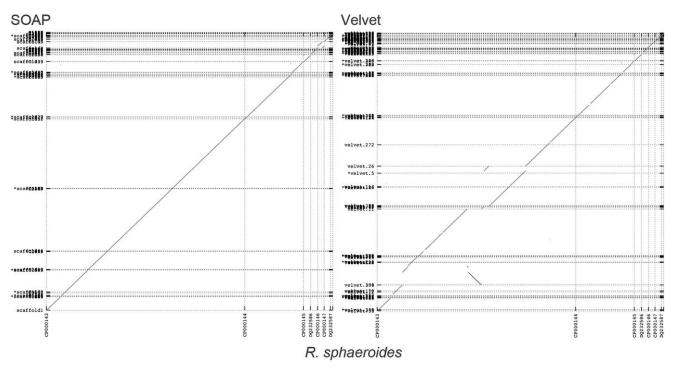


Figure 2. A dot-plot comparison of the SOAPdenovo and Velvet scaffolds of *R. sphaeroides*. The finished reference chromosomes are plotted on the *x*-axis and the assembly scaffolds on the *y*-axis. Dotted lines indicate scaffold or chromosome boundaries. The apparent rearrangement at the *top right* of the SOAPdenovo plot is an artifact of the circular reference plasmid.

However, after comparing it with the reference genome, we found that SOAPdenovo contained multiple assembly errors (Table 2). Breaking the assembly at these errors produced a much smaller N50 value of 63 kb. The N50 size for ALLPATHS-LG was initially 97 kb, and with many fewer assembly errors, breaking the contigs reduced the N50 value less dramatically, to 66 kb, making it the best of the assemblers on this genome. MSR-CA's corrected N50 of 48 kb placed it below SOAPdenovo, but with about half as many assembly errors (34 vs. 65), MSR-CA would appear preferable to SOAPdenovo.

ALLPATHS-LG, MSR-CA, and Bambus2 all produced very large scaffolds, with MSR-CA producing a single scaffold containing the entire main chromosome. However, this scaffold contained several inversions, and only ALLPATHS-LG and Bambus2 produced scaffolds with no major errors.

Note that CABOG was not run on *S. aureus* because one of the two paired-end libraries contains reads of just 37 bp, and CABOG has a minimum read length of 64 bp.

R. sphaeroides

For *Rhodobacter* (Table 3), Bambus2 had the smallest number of contigs and scaffolds, with relatively large N50 sizes in both categories. The largest contigs were built by SOAPdenovo (with an N50 size of 132 kb), followed by Bambus2 (93 kb) and ALLPATHS-LG (42 kb).

As with *Staphylococcus*, however, the errors in the assemblies made some, particularly SOAPdenovo, appear to be better than they really were. With 422 errors, SOAPdenovo was the most errorprone of all the assemblers for *Rhodobacter*, and after breaking contigs at these errors, its N50 size was just 14.3 kb, dropping it to fifth place for contiguity. Bambus2 had almost as many errors and dropped even further after correction, to 12.8 kb. ALLPATHS-LG's contiguity dropped the least, and after correction its contig N50 of 34.4 kb was the best, followed by MSR-CA at 19.1 kb.

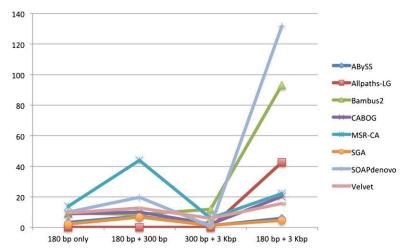


Figure 3. Assemblies of *R. sphaeroides* using four different combinations of paired-end libraries as input to the assemblers. Each run used either one library (180 bp only) or a different combination of two libraries from 180 to 3000 bp. Note that N50 values are uncorrected; see Table 3 for the true N50 sizes for the 180 bp + 3 kb combination, which are much lower in some instances; e.g., SOAPdenovo has a corrected N50 of 14.3 kb (rather than 131.7 kb) for assembly with the 180-bp and 3-kb libraries.

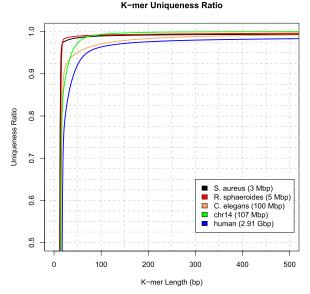


Figure 4. K-mer uniqueness ratio for the three genomes assembled in GAGE: the bacteria *S. aureus* and *R. sphaeroides* and human chromosome 14. The ratio is defined as the percentage of a genome that is covered by unique (i.e., non-repetitive) DNA sequences of length *K.* Shown for comparison are the k-mer uniqueness ratios for the full human genome and for the nematode *C. elegans*.

ALLPATHS-LG also produced the best scaffolding results, with the main chromosome entirely spanned by a single scaffold, closely followed by MSR-CA and Bambus2. SOAPdenovo's scaffolding results were a distant fourth place, approximately five times smaller than ALLPATHS-LG. An important caveat on these results is that the *Rhodobacter* data set was created following the ALLPATHS-LG "recipe" for library construction, which makes it an ideal data set for that assembler.

Although the overall results were similar for the two bacterial data sets, the sizes of the contigs were generally much larger for *S. aureus*, and the contigs for a given assembler varied by as much as sixfold (for ABySS). This variation illustrates how one of the most important variables in predicting assembly contiguity may be the genome itself, which is an element that cannot be controlled.

Human chromosome 14

For the human chromosome data, most of the assemblers produced relatively poor results, and the differences between the best and worst assemblers were dramatic. As with Rhodobacter, the sequencing strategy and mate-pair data were designed specifically for ALLPATHS-LG, and the creators of some of the assemblers might not have anticipated or taken full advantage of this type of data (particularly the library with overlapping mates). Regardless of the reason, ALLPATHS-LG and CABOG clearly outperformed all of the other assemblers in the contiguity statistics shown in Table 4. CABOG's contigs were 30% larger than those from ALLPATHS-LG (45.3 kb vs. 36.5 kb), but both were far larger than those produced by any of the other methods, most of which built contigs in the 2-4-kb range. Even more dramatic was the exceptionally large scaffold produced by ALLPATHS-LG, which contained almost the entire chromosome in one scaffold of 81.6 Mb. The largest scaffold generated by any other assembler was one produced by Velvet, at only 4.6 Mb.

After adjusting for misassemblies (Table 4), CABOG remained slightly ahead of ALLPATHS-LG, with both dropping substantially,

to 23.7 kb and 21.0 kb, respectively. They remained far ahead of the third-best assembler, SOAPdenovo, with an N50 size of just 7.4 kb. It is also important to note that all of the leading performers had thousands of assembly errors on this chromosome, which translates into tens of thousands of errors on a full human genome. Fewer errors were found in the assemblies of ABySS (704 errors) and SGA (981 errors), but their more-cautious approaches produced very small contig N50 sizes of 2.0 and 2.7 kb. Thus, despite all efforts at error correction and repeat identification, assembly of a mammalian genome from NGS data remains an extremely challenging problem.

B. impatiens

Unlike the other three genomes, the bumble bee (*B. impatiens*) does not have a finished reference. Based on the results above, contiguity and size statistics should be interpreted very cautiously; it is possible that assembly errors, if known, would dramatically change these values, as they did in our experiments on *S. aureus* above. Nonetheless, we found that SOAPdenovo generated contigs with nearly double the N50 size of CABOG, 57 kb versus 24 kb. The scaffold N50 sizes were all similar, although SOAPdenovo's were slightly larger than the others. Worth noting here is that in experiments using an earlier (2010) release of SOAPdenovo, it could only produce contigs with an N50 of 6.4 kb, indicating a substantial improvement in that assembler in its more recent version.

Most of the other assemblers could not assemble these data at all, for various reasons. ALLPATHS-LG could not be used because it requires at least one library with overlapping mate pairs, which this project did not have. The other assemblers appeared to be unable to handle the large number of reads (\sim 500 million), and most of them crashed, often after several days running on a 256-GB multi-core computer. This illustrates an underappreciated fact of genome assembly with current technology: For larger genomes, the choice of assemblers is often limited to those that will run without crashing.

Shared assembly errors

To address the question of whether assembly errors were common or different among all of the algorithms, we looked at the intersections of errors on the assembly of Hs14. Insofar as the errors are unique, then it might be beneficial to merge the results of multiple assemblers to produce a consensus assembly. We focused on errors >5 bp, which include the collapse or expansion of small tandem repeats as well as larger errors. As shown in Figure 5, Bambus2, Velvet, and SOAPdenovo had significantly more unique errors than the other assemblers, ranging from just over 2000 (SOAPdenovo) to 4000 (Bambus2). SGA had by far the fewest unique errors. Among the shared errors, ALLPATHS-LG and CABOG had the largest numbers, suggesting that these two assemblers might agree with one another and possibly that some of their errors might represent true haplotype differences. Finally, there were about 200 errors shared by all eight assemblers, indicating that these are likely true variations in the target genome rather than errors.

Conclusions

Figure 6 summarizes the results across the three genomes for which the true assembly is available. ALLPATHS-LG demonstrated consistently strong performance based on contig and scaffold size, with the best trade-off between size and error rate, as shown in the figure. MSR-CA also performed relatively well, although with more

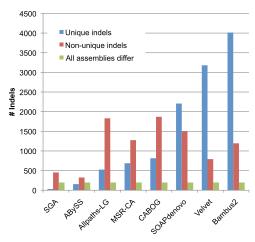


Figure 5. Comparison of insertion and deletion errors among all eight assemblers for human chromosome 14. (Blue) The indel errors >5 bp in length that are unique to each assembler. (Red bars) Indel errors made by at least one other assembler. (Green bars) Indels shared by all assemblers, which might represent true differences between the target genome and the reference.

errors than ALLPATHS-LG. Bambus2 seems to be a very capable scaffolder, as shown in Figure 6, but its contigs contain numerous small errors. (An explanation for this result is that contig merging is a very recent addition to Bambus2, one that is still under development.) The latter two assemblers use parts of the CABOG assembler for many of their core functions, and in this respect their performance is not independent. SOAPdenovo produced results that initially seemed superior to most assemblers, but on closer inspection it generated many misassemblies that would be impossible to detect without access to a reference genome. Despite its poor performance on human, SOAPdenovo performed very well on the bacteria, creating contigs that were eight times larger than it built on the human data. Finally, Table 7 and Figure 6 show that Velvet had a particularly high error rate for its scaffolds, creating many more inversions and translocations than any other algorithm.

As illustrated by the differences between the original and corrected N50 values in Tables 2-4, an assembler can produce a large N50 value by using an overly aggressive assembly strategy, which, in turn, will yield a higher number of errors. In contrast, more conservative assemblers might produce smaller contigs but fewer errors. For the genomes examined here, ALLPATHS-LG and CABOG stood out as assemblers capable of producing both high contiguity and high accuracy. SOAPdenovo often produced similar or larger N50 values, but it appears to achieve this by sacrificing correctness. For all three of the previously sequenced genomes, SOAPdenovo showed a higher rate of chaff, duplications, compressions, SNPs, indels, and misjoins than CABOG and ALLPATHS-LG. Considering all metrics, and with the caveat that it requires a precise recipe of input libraries, ALLPATHS-LG appears to be the most consistently performing assembler, both in terms of contiguity and correctness.

For all of the assemblers, contig sizes for the human chromosome assembly were smaller than contigs for either of the bacterial genomes. The problem would only be more difficult if we had used the entire genome rather than a single chromosome. We conclude that, despite very significant improvements in assembly technology, the problem of assembling a large genome from short reads remains very difficult. The remarkable gains in sequencing throughput of recent years will require further improvements, especially in read length and in paired-end protocols, before we are likely to see accurate, highly contiguous mammalian assemblies. Thanks to algorithmic improvements, the assemblers used in this study can handle very large data volumes, but they will need longer-range linking information if they are to match or exceed the quality of assemblies based on Sanger sequencing technology.

Finally, we should note that all of the assemblers considered here are under constant development, and many will be improved by the time this analysis appears. Evaluations of assemblers such as GAGE are useful snapshots of performance, but ongoing reevaluation will be necessary as algorithms and sequencing technology change. Assembly evaluations should also be reproducible, which requires that the complete recipes for running these complex programs should be provided, as we have done here for the first time.

Methods

Data for *S. aureus* were downloaded from the Sequence Read Archive (SRA) at NCBI, accession numbers SRX007714 and SRX016063. The *R. sphaeroides* data have SRA accessions SRX033397 and SRX016063. The SRA libraries downloaded had higher coverage than was needed for most experiments. Each library was therefore randomly sampled to create a data set with $45 \times$ genome coverage, giving a total of $90 \times$ coverage for each genome.

To create the human chromosome 14 data set, reads sequenced from cell line GM12878 were downloaded from the SRA under the following accession numbers: SRR067780, SRR067784, SRR067785, SRR067787, SRR067789, SRR067791–SRR067793, SRR067771, SRR067773, SRR067776–SRR067779, SRR067781, SRR067786, SRR068214, SRR068211, SRR068335. Reads came from one short fragment library (mean read length 101

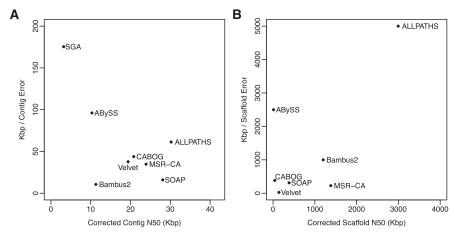


Figure 6. Average contig (A) and scaffold (B) sizes, measured by N50 values, versus error rates, averaged over all three genomes for which the true assembly is known: *S. aureus, R. sphaeroides,* and human chromosome 14. Errors (vertical axis) are measured as the average distance between errors, in kilobases. N50 values represent the size N at which 50% of the genome is contained in contigs/scaffolds of length N or larger. In both plots, the best assemblers appear in the *upper right*.

bp, fragment size 155 bp), two short jump libraries (101-bp mean read length, 2536-bp mean insert size), and two fosmid libraries (76-bp mean read length, 35,295-bp mean insert length). The original set of >1 billion reads was mapped against the entire human genome (GRCh37/hg19) using Bowtie (Langmead et al. 2009); reads mapping to multiple locations were randomly distributed across those locations (parameters: -128 -n 3 -e 300 -3 20 -M 1-best). Only reads mapping to Hs14 were retained. Each read in a pair was mapped separately to allow for inclusion of real distribution of insert sizes (including chimeric reads) and to avoid excessively filtering the data so as to better reflect the distribution in the original data set. The overall coverage of Hs14 was $60\times$, as shown in Supplemental Figure 1, and the number of gaps in coverage was 108, with gap sizes ranging from 1 to 2412 bp.

The *B. impatiens* data were sequenced at the Keck Center for Comparative and Functional Genomics, University of Illinois and released for public use by Gene Robinson.

Reads were error-corrected using both Quake and the ALLPATHS-LG error corrector (for details, see the Supplemental Methods). All assemblers were run using multiple parameters and with corrected and uncorrected reads as input; the best assembly for each genome was chosen.

For the three previously finished genomes, N50 sizes were computed based on the known size of the genome. For the bumble bee, N50 sizes used the estimated genome size of 250 Mb. Contigs and scaffolds of 200 bp or longer were used for all computations.

Because N50 size might sometimes be a misleading statistic, we also computed another statistic, which we call E-size. The E-size is designed to answer the question: If you choose a location (a base) in the reference genome at random, what is the expected size of the contig or scaffold containing that location? This statistic is one way to answer the related question: How many genes will be completely contained within assembled contigs or scaffolds, rather than split into multiple pieces? E-size is computed as:

$$E = \sum_{c} \frac{(L_c)^2}{G},$$

where $L_{\rm C}$ is the length of contig C, and G is the genome length estimated by the sum of all contig lengths. E-size is computed similarly for scaffolds. To be consistent across all assemblies, we only considered contigs and scaffolds of 200 bp or longer in computing the E-size, and we used a constant value of G for all assemblies of a given genome. After computing E-sizes for all assemblies and all genomes, we found that they correlated very closely with N50 sizes in every case, validating our choice of N50 size as a representative assembly size metric. E-sizes for all assemblies can be found in Supplemental Table 1.

For evaluating correctness, alignment statistics and misassemblies were tallied using the program *dnadiff* (Phillippy et al. 2008) from MUMmer v3.23 (Kurtz et al. 2004). *dnadiff* operates by constructing local pairwise alignments between a reference and query genome using the Nucmer aligner. The aligned segments are then filtered to obtain a globally optimal mapping between the reference and query segments, while allowing for rearrangements, duplications, and inversions. This technique was later described in detail by Dubchak et al. (2009) as the SuperMap algorithm. Conveniently, this method identifies both a one-to-one mapping of segments as well as any duplicated sequences. When applied to assembly mapping, it can be used to measure the quantity and types of common misassemblies.

To create the alignments, contigs <200 bp were excluded, and the remainder were aligned using *nucmer* (Kurtz et al. 2004) with the options "-maxmatch -l 30 -banded -D 5." Combined with its default options, this invocation requires a minimum exact-match

anchor size of 30 bp and a minimum combined anchor length of 65 bp per cluster. Clusters are further required to have no more than 90 bp separation or more than five inserted bases between any two adjacent anchors. Acceptable clusters are then used to seed banded Smith-Waterman alignments (Smith and Waterman 1981). After running *nucmer*, alignments with <95% identity or >95% overlap with another alignment were discarded using *delta-filter. dnadiff* was then executed on the remaining alignments with default parameters, and correctness statistics were tabulated from its output (see the Supplemental Material).

For the scaffolds, we calculated three types of errors: indels, where there is an incorrect interleaving of multiple scaffolds; inversions, where a scaffold switches strands within a chromosome; and translocations, where a scaffold maps to multiple chromosomes in the reference. We also counted the number of gaps where the scaffold gap-size estimate is at least 1 kb off and the average absolute difference between the scaffold gap estimate and true gap size in each assembly. Details of how the scaffolds were aligned are in the Supplemental Material.

Any alignment-based metric is subject to the accuracy of the underlying alignments. Because complex repeat structures made the correct determination of alignment boundaries difficult in some cases, the figures presented here are to be taken only as estimates of the various features of each assembly. This is especially true of the misjoin features, which penalize small contig misassemblies just as severely as more major rearrangements. However, even allowing for some alignment-based error, the relative performance of each assembler would likely remain the same, and we should emphasize that all assemblies were analyzed with identical methods and against the same reference genomes.

Data access

All data sets, including error-corrected reads for each genome, are freely available from http://gage.cbcb.umd.edu/data.

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