# Generation of Concurrency Control Code using Discrete-Event Systems Theory

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### Motivation

- Optimal usage of today's Multi-core systems requires parallel executable Software
- Efficient and correct concurrent source code is hard to write and debug

→ facilitate development of concurrent source code

### Introduction

- Automatic generation of concurrency control code
  - Input
    - Source code without concurrency control
    - Informal specification of the desired concurrent behaviour
  - Output
    - Source code with concurrency controls

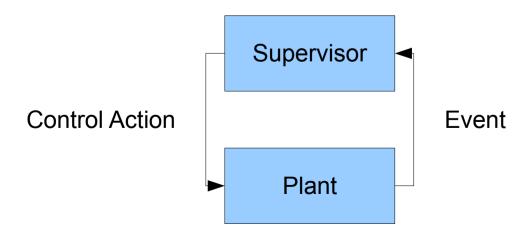
### Introduction

- Aims
  - No Deadlocks
  - No Starvation
  - Minimally restrictive

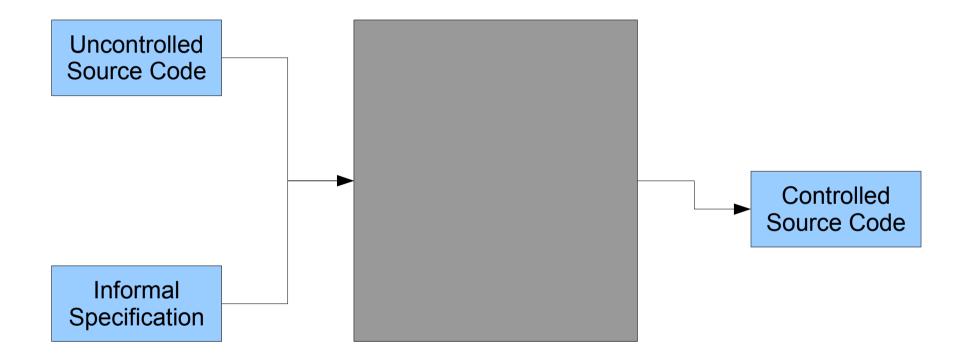
→ using control theory (Discrete-event System)

# DES - Supervisor-Plant

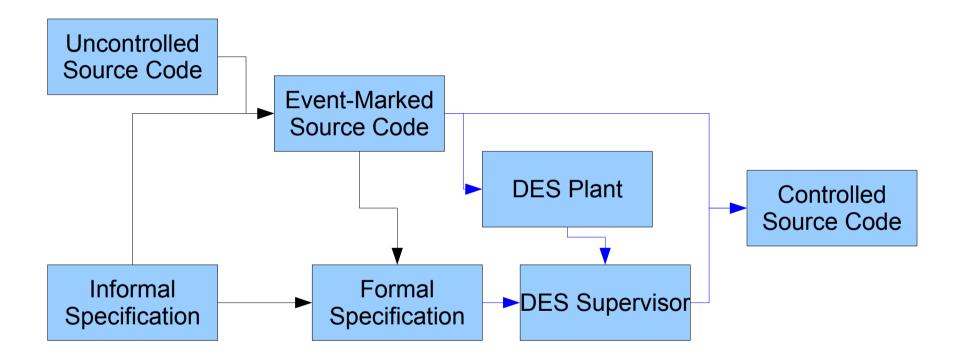
- System(Plant) modelled as a finite-state automaton (FSA)
  - Transitions in FSA are called events
    - Events can be controllable or uncontrollable
- Supervisor modelled as a FSA
  - Enables or disables controllable events



### **Process Overview**



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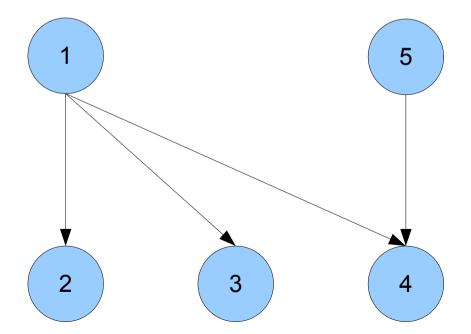


### Relevant Events

- Find relevant events for concurrency control
  - Example
    - Enter/Exit Critical Section → relevant
    - Accessing shared variable → non-relevant
- Mark events in the source code
- Differentiate controllable and uncontrollable events

# Example

- 5 Threads with dependencies
  - $1 \to 2,3,4$
  - $5 \rightarrow 4$



```
Code example for Thread 3:

public void run() {

    // relevant event: T3-start

    System.out.println(id);

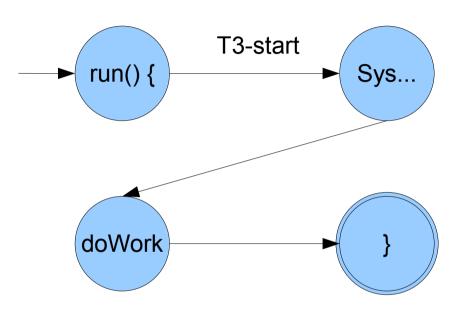
    doWork();

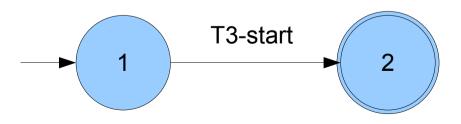
1
```

### **DES Plant**

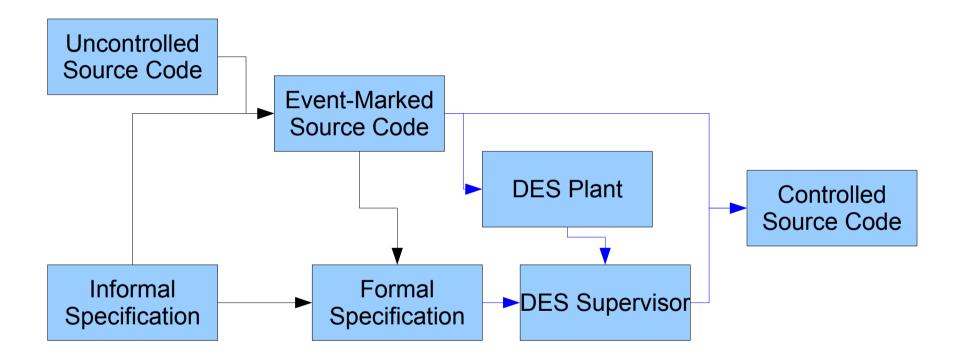
- Build a Finite-state automaton (FSA)
   representing all possible event sequences for
   each Thread
  - Build control-flow graph(CFG) from source code
  - Reduce CFG
    - All relevant events remain
    - All non-relevant events important for CFG structure remain

### Example





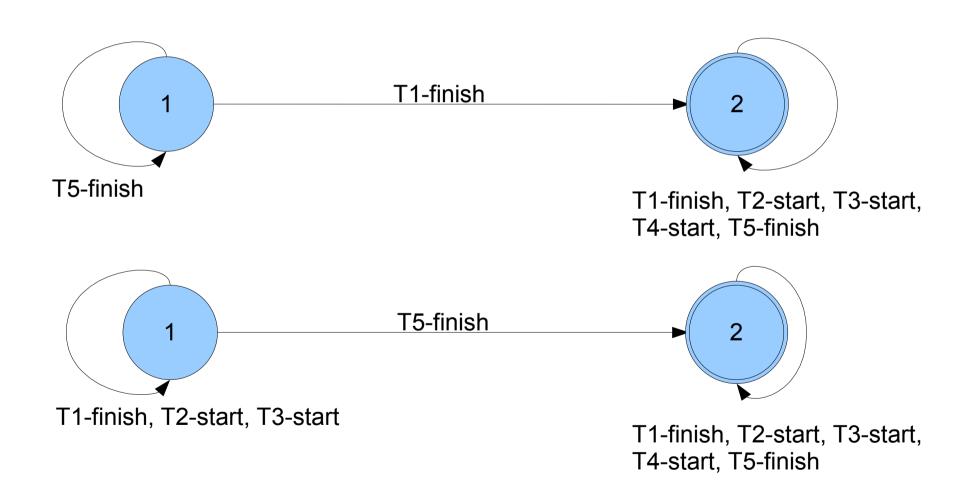
### **Process Overview**



# Formal Specification

- Specifies the allowed subset of event sequences
  - FSA for each restriction
  - Only restrictive

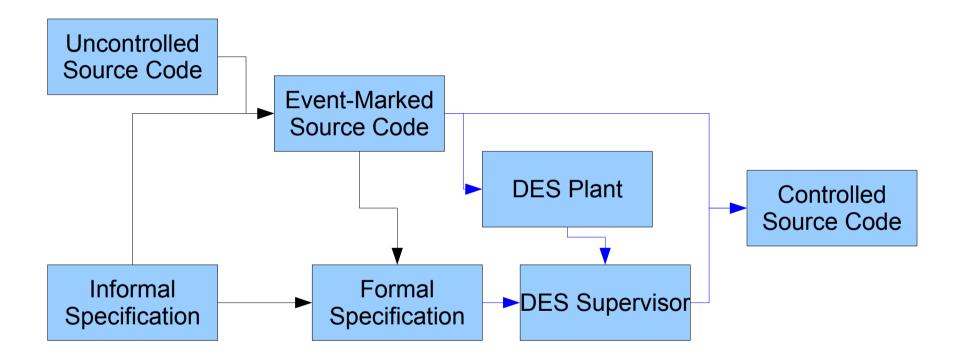
# Example



# **DES Supervisor**

- Only one supervisor (simplification)
- Different FSA's need to be combined into a monolithic specification
  - Scalability issue

### **Process Overview**



### **Code Generation**

- Supervisor needs to block non allowed controllable events
  - Generates a Semaphore for each controllable event set to the initial state
  - Supervisor state changing function enables/disables controllable events

### Example – Thread 3

```
//relevant event: T3start
while (true) {
    if(Synchronizer.stateChangeTest("T3start",
        Synchronizer.T3start))
        break;
    Synchronizer.T3start.acquireUninterruptibly();
    Synchronizer.T3start.release();
}
```

### Example – Supervisor I

```
public static synchronized Boolean stateChangeTest(String event,
Semaphore eventBlocker) {
   if (!(eventBlocker == null)) {
       if (!eventBlocker.tryAcquire()) {
           return false;
       eventBlocker.release();
   changeSupervisorState(event);
   return true;
```

### Example – Supervisor II

```
private static void changeSupervisorState(String event) {
   if (event.equals("T1finish")) {
      switch(Synchronizer.stateTracker) {
      case(0):
         Synchronizer.T3start.release();
         Synchronizer.T2start.release();
         Synchronizer.stateTracker = 1;
      break;
      case(1):
```

### Verification

- Discrete-event theory is proven correct and non blocking
- Formal proof for algorithm is still needed
- Modelchecker (Java Pathfinder)

Input not reliable!

### Limitations

- No dynamic threads generation allowed
- Monolithic supervisor not very efficient
- Starvation not properly addressed

### Conclusion

- DES theory can be applied to concurrency
- Chosen FSA-based version of DES not expressive enough
- Further work needed

# Questions

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