Generation of Finite Difference Formulas on Arbitrarily Spaced Grids

By Bengt Fornberg

Abstract. Simple recursions are derived for calculating the weights in compact finite difference formulas for any order of derivative and to any order of accuracy on onedimensional grids with arbitrary spacing. Tables are included for some special cases (of equispaced grids).

1. Introduction. Previously published methods to generate finite difference weights (e.g., references [1]-[5]) have been of considerable complexity and often been limited to derivatives of low order on equidistantly spaced grids. The most ambitious attempt to tabulate weights for many orders of derivatives and to high orders of accuracy appears to be the work by Keller and Pereyra [4]. However, their algorithms (limited to equispaced grids) were very involved, and the resulting tables contain both isolated and systematic errors.

In the present study we describe two simple recursion relations which give the weights for any order of derivative (including the 0th derivative, corresponding to interpolation), approximated to any order of accuracy on an arbitrary grid in one dimension. Since, in general, only four arithmetic operations are needed to determine each weight, the main anticipated application of the present method is to dynamically changing grids. However, the method is also well suited to generate tables of weights. Such tables (in the special case of equispaced grids, up to the 4th derivative and up to 9 weights) are included in the cases of one-sided and centered approximations at a grid point and at a 'half-way point' between grid points.

2. Notation, Algorithm. Given $M \ge 0$, the order of the highest derivative we wish to approximate, and a set of N + 1 grid points (at *x*-coordinates $\alpha_0, \ldots, \alpha_N$; $N \ge 0$), the problem is to find all the weights such that the approximations

$$\left.\frac{d^m f}{dx^m}\right|_{x=x_0} \approx \sum_{\nu=0}^n \delta_{n,\nu}^m f(\alpha_{\nu}), \qquad m=0,1,\ldots,M; \ n=m,m+1,\ldots,N,$$

become of optimal formal order of accuracy (in general of order n-m+1, although it can be higher in special cases). The following algorithm achieves this:

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Enter
$$M$$
, N , x_0 , α_0 , α_1 , α_2 , ..., α_N
 $\delta_{0,0}^0 := 1$
 $c1 := 1$
for $n := 1$ to N do
 $c2 := 1$
for $\nu := 0$ to $n - 1$ do
 $c3 := \alpha_n - \alpha_\nu$
 $c2 := c2 \cdot c3$
if $n \le M$ then $\delta_{n-1,\nu}^n := 0$
for $m := 0$ to min (n, M) do
 $\delta_{n,\nu}^m := ((\alpha_n - x_0)\delta_{n-1,\nu}^m - m\delta_{n-1,\nu}^{m-1})/c3$
next m
next ν
for $m := 0$ to min (n, M) do
 $\delta_{n,n}^m := \frac{c1}{c2}(m\delta_{n-1,n-1}^{m-1} - (\alpha_{n-1} - x_0)\delta_{n-1,n-1}^m)$
next m
 $c1 := c2$
next n

Notes. 1. If the array $\delta_{n,\nu}^m$ initially is zero, the statement "if $n \leq M$ then $\delta_{n-1,\nu}^n := 0$ " can be omitted.

2. In the case of m = 0 (corresponding to interpolation formulas), expressions of the form 'zero*(undefined number)' occur. The result is assumed to be zero.

3. The order in which the α_{ν} (all distinct) are given is significant (since the weights corresponding to all leading subsets of the α_{ν} 's are calculated). There is no restriction on x_0 coinciding with any α_{ν} .

3. Derivation of the Algorithm. For simplicity, assume we seek to approximate the derivatives at the point $x_0 = 0$. Let $\{\alpha_0, \alpha_1, \ldots, \alpha_N\}$ be distinct real numbers and denote

(3.1)
$$\omega_n(x) := \prod_{k=0}^n (x - \alpha_k).$$

The polynomial

(3.2)
$$F_{n,\nu}(x) := \frac{\omega_n(x)}{\omega'_n(\alpha_\nu)(x - \alpha_\nu)}$$

is the one of minimal degree which takes the value 1 at $x = \alpha_{\nu}$ and 0 at $x = \alpha_{k}$, $0 \le k \le n, k \ne \nu$. For an arbitrary function f(x) and nodes $x = \alpha_{\nu}$, Lagrange's interpolation polynomial becomes

(3.3)
$$p(x) := \sum_{\nu=0}^{n} F_{n,\nu}(x) f(\alpha_{\nu}).$$

The desired weights express how the values of $[d^m p(x)/dx^m]_{x=0}$ vary with changes in $f(\alpha_{\nu})$. Since only one term in p(x) is influenced by changes in each $f(\alpha_{\nu})$, we find

(3.4)
$$\delta_{n,\nu}^m = \left[\frac{d^m}{dx^m}F_{n,\nu}(x)\right]_{x=0}$$

Therefore, the *n*th degree polynomial $F_{n,\nu}(x)$ can also be expressed as

(3.5)
$$F_{n,\nu}(x) = \sum_{m=0}^{n} \frac{\delta_{n,\nu}^{m}}{m!} x^{m}$$

From (3.2) follow (noting that $\omega(x) = (x - \alpha_n)\omega_{n-1}(x)$ implies $\omega'_n(x) = (x - \alpha_n)\omega'_{n-1}(x) + \omega_{n-1}(x)$)

(3.6)
$$F_{n,\nu}(x) = \frac{x - \alpha_n}{\alpha_\nu - \alpha_n} F_{n-1,\nu}(x)$$

and

(3.7)
$$F_{n,n}(x) = \frac{\omega_{n-1}(x)}{\omega_{n-1}(\alpha_n)} = \frac{\omega_{n-2}(\alpha_{n-1})}{\omega_{n-1}(\alpha_n)} (x - \alpha_{n-1}) F_{n-1,n-1}(x) \qquad (n > 1).$$

By substituting the expansion (3.5) into (3.6) and (3.7), and by equating powers of x, the desired recursion relations between the weights are obtained:

(3.8)
$$\delta_{n,\nu}^{m} = \frac{1}{\alpha_{n} - \alpha_{\nu}} (\alpha_{n} \delta_{n-1,\nu}^{m} - m \delta_{n-1,\nu}^{m-1})$$

and

(3.9)
$$\delta_{n,n}^{m} = \frac{\omega_{n-2}(\alpha_{n-1})}{\omega_{n-1}(\alpha_{n})} (m\delta_{n-1,n-1}^{m-1} - \alpha_{n-1}\delta_{n-1,n-1}^{m}).$$

The relation

(3.10)
$$\sum_{\nu=0}^{n} \delta_{n,\nu}^{m} = \begin{cases} 1, & m = 0, \\ 0, & m > 0, \end{cases}$$

could be used instead of (3.9) to obtain $\delta_{n,n}^m$. However, this would increase the operation count and might also cause a growth of errors in the case of floating-point arithmetic.

4. Description of the Tables. Special cases which commonly occur are centered and one-sided approximations on equidistant grids. The particular choices of α_{ν} used for Tables 1-4 correspond to grid spacings $\Delta x = 1$. For other values of Δx , these coefficients should be divided by $(\Delta x)^m$ (where m, as before, is the order of the derivative).

Some weights for centered approximations at a grid point (generated by setting $M =$
4, $N = 8$, $x_0 = 0$ and $\alpha_{\nu} = \{0, 1, -1, 2, -2, 3, -3, 4, -4\}$.

der or i der t der t i ver t i ver t	Order of	Approximations at $x = 0$; x-coordinates at nodes:								
$ \begin{array}{c} 0 & 1 \\ \mathbf{f} & \mathbf{e} \end{array} $	o ầ fy	-4	-3	-2	-1	0	1	2	3	4
0	∞					1				
	2				$\frac{-1}{2}$	0	$\frac{1}{2}$			
	4			$\frac{1}{12}$	$\frac{-1}{2}$ $\frac{-2}{3}$ $\frac{-3}{4}$	0	$\frac{2}{3}$	$\frac{-1}{12}$		
1	6		$\frac{-1}{60}$	$\frac{3}{20}$	$\frac{-3}{4}$	0	$\frac{3}{4}$	$\frac{-1}{12}$ $\frac{-3}{20}$	$\frac{1}{60}$	
	8	$\frac{1}{280}$	$\frac{-4}{105}$	$\frac{1}{5}$	$\frac{-4}{5}$	0	$\frac{4}{5}$	$\frac{-1}{5}$	$\frac{4}{105}$	$\frac{-1}{280}$
	2				1	-2	1			
	4			$\frac{-1}{12}$	$\frac{4}{3}$	$\frac{-5}{2}$	$\frac{4}{3}$	$\frac{-1}{12}$		
2	6		$\frac{1}{90}$	$\frac{-1}{12}$ $\frac{-3}{20}$	43 32 85	$\frac{-49}{18}$	4 3 3 2	$\frac{-1}{12}$ $\frac{-3}{20}$	$\frac{1}{90}$	
	8	$\frac{-1}{560}$	$\frac{8}{315}$	$\frac{-1}{5}$	$\frac{8}{5}$	$\frac{-205}{72}$	$\frac{8}{5}$	$\frac{-1}{5}$	$\frac{8}{315}$	$\frac{-1}{560}$
	2			$\frac{-1}{2}$	1	0	-1	$\frac{1}{2}$		
3	4		$\frac{1}{8}$	-1	$\frac{13}{8}$	0	$\frac{-13}{8}$	1	$\frac{-1}{8}$	
	6	$\frac{-7}{240}$	$\frac{\frac{1}{8}}{\frac{3}{10}}$	$\frac{-169}{120}$	$\frac{61}{30}$	0	$\frac{-61}{30}$	$\frac{169}{120}$	$\frac{-3}{10}$	$\frac{7}{240}$
	2			1	-4	6	-4	1		
4	4		$\frac{-1}{6}$	2	$\frac{-13}{2}$	$\frac{28}{3}$	$\frac{-4}{2}$	2	$\frac{-1}{6}$	
	6	$\frac{7}{240}$	$\frac{-2}{5}$	<u>169</u> 60	$\frac{-122}{15}$	<u>91</u> 8	$\frac{-122}{15}$	<u>169</u> 60	$\frac{-2}{5}$	$\frac{7}{240}$

Some weights for centered	l approximations at a	'half-way' point (generated by setting
$M = 4, N = 7, x_0 = 0$ and	and $\alpha_{\nu} = \{1/2, -1/2, 3\}$	3/2, -3/2, 5/2, -3	$5/2, 7/2, -7/2\}).$

der ridvartiov fe	Order of	Approximations at $x = 0$; x-coordinates at nodes:									
fě	fÿ	-7/2	-5/2	-3/2	-1/2	1/2	3/2	5/2	7/2		
	2				$\frac{1}{2}$	$\frac{1}{2}$					
	4			$\frac{-1}{16}$	$\frac{9}{16}$	$\frac{9}{16}$	$\frac{-1}{16}$				
0	6		$\frac{3}{256}$	$\frac{-25}{256}$	$\frac{75}{128}$	$\frac{75}{128}$	$\frac{-25}{256}$	$\frac{3}{256}$			
	8	$\frac{-5}{2048}$	$\frac{49}{2048}$	$\frac{-245}{2048}$	$\frac{1225}{2048}$	$\frac{1225}{2048}$	$\tfrac{-245}{2048}$	$\frac{49}{2048}$	$\frac{-5}{2048}$		
	2				-1	1					
	4			$\frac{1}{24}$	$\frac{-9}{8}$	<u>9</u> 8	$\frac{-1}{24}$				
1	6		$\frac{-3}{640}$	$\frac{25}{384}$	$\frac{-75}{64}$	$\frac{75}{64}$	$\frac{-25}{384}$	$\frac{3}{640}$			
	8	$\frac{5}{7168}$	$\frac{-49}{5120}$	$\tfrac{245}{3072}$	$\tfrac{-1225}{1024}$	$\frac{1225}{1024}$	$\tfrac{-245}{3072}$	$\frac{49}{5120}$	$\frac{-5}{7168}$		
	2			$\frac{1}{2}$	$\frac{-1}{2}$	$\frac{-1}{2}$	$\frac{1}{2}$				
2	4		$\frac{-5}{48}$	$\frac{13}{16}$	$\frac{-17}{24}$	$\frac{-17}{24}$	$\frac{13}{16}$	$\frac{-5}{48}$			
	6	$\tfrac{259}{11520}$	$\tfrac{-499}{2304}$	$\frac{1299}{1280}$	$\tfrac{-1891}{2304}$	$\frac{-1891}{2304}$	$\frac{1299}{1280}$	$\tfrac{-499}{2304}$	$\frac{259}{11520}$		
	2			-1	3	-3	1				
3	4		$\frac{1}{8}$	$\frac{-13}{8}$	$\frac{17}{4}$	$\frac{-17}{4}$	<u>13</u> 8	$\frac{-1}{8}$			
	6	$\frac{-37}{1920}$	$\frac{499}{1920}$	$\frac{-1299}{640}$	$\tfrac{1891}{384}$	$\frac{-1891}{384}$	<u>1299</u> 640	$\frac{-499}{1920}$	$\frac{37}{1920}$		
	2		$\frac{1}{2}$	$\frac{-3}{2}$	1	1	$\frac{-3}{2}$	$\frac{1}{2}$			
4	4	$\frac{-7}{48}$	<u>59</u> 48	$\frac{-45}{16}$	<u>83</u> 48	<u>83</u> 48	$\frac{-45}{16}$	<u>59</u> 48	$\frac{-7}{48}$		

d	<u> </u>						, ,,,,	<u> </u>		
d Ori dva rt	O ra		Approxi	mations a	at $x = 0;$					
e a r t	e c r u r r			nates at						
ov fe	accuracy	0	1	2	3	4	5	6	7	8
0	∞	1								
	1	-1	1							
	2	$\frac{-3}{2}$	2	$\frac{-1}{2}$						
	3	$\frac{-11}{6}$	3	$\frac{-3}{2}$	$\frac{1}{3}$					
	4	$\frac{-25}{12}$	4	-3	$\frac{4}{3}$	$\frac{-1}{4}$				
1	5	$\frac{-137}{60}$	5	-5	$\frac{10}{3}$	$\frac{-5}{4}$	$\frac{1}{5}$			
	6	$\frac{-49}{20}$	6	$\frac{-15}{2}$	$\frac{20}{3}$	$\frac{-15}{4}$	$\frac{6}{5}$	$\frac{-1}{6}$		
	7	$\frac{-363}{140}$	7	$\frac{-21}{2}$	$\frac{35}{3}$	$\frac{-35}{4}$	$\frac{21}{5}$	$\frac{-7}{6}$	$\frac{1}{7}$	
	8	$\frac{-761}{280}$	8	-14	$\frac{56}{3}$	$\frac{-35}{2}$	$\frac{56}{5}$	$\frac{-14}{3}$	$\frac{8}{7}$	$\frac{-1}{8}$
	1	1	-2	1						
	2	2	$^{-5}$	4	-1					
	3	$\frac{35}{12}$	$\frac{-26}{3}$	$\frac{19}{2}$	$\frac{-14}{3}$	$\frac{11}{12}$				
2	4	$\frac{15}{4}$	$\frac{-77}{6}$	$\frac{107}{6}$	-13	$\frac{61}{12}$	$\frac{-5}{6}$			
	5	$\frac{203}{45}$	$\frac{-87}{5}$	$\frac{117}{4}$	$\frac{-254}{9}$	$\frac{33}{2}$	$\frac{-27}{5}$	$\frac{137}{180}$		
	6	<u>469</u> 90	$\frac{-223}{10}$	$\tfrac{879}{20}$	$\frac{-949}{18}$	41	$\frac{-201}{10}$	$\tfrac{1019}{180}$	$\frac{-7}{10}$	
	7	$\tfrac{29531}{5040}$	$\frac{-962}{35}$	$\frac{621}{10}$	$\frac{-4006}{45}$	<u>691</u> 8	$\frac{-282}{5}$	<u>2143</u> 90	$\frac{-206}{35}$	<u>363</u> 560
	1	-1	3	-3	1					
	2	$\frac{-5}{2}$	9	-12	7	$\frac{-3}{2}$				
	3	$\frac{-17}{4}$	$\frac{71}{4}$	$\frac{-59}{2}$	$\frac{49}{2}$	$\frac{-41}{4}$	$\frac{7}{4}$			
3	4	$\frac{-49}{8}$	29	$\frac{-461}{8}$	62	$\frac{-307}{8}$	13	$\frac{-15}{8}$		
	5	$\frac{-967}{120}$	$\frac{638}{15}$	$\frac{-3929}{40}$	$\frac{389}{3}$	$\frac{-2545}{24}$	$\frac{268}{5}$	$\frac{-1849}{120}$	$\frac{29}{15}$	
	6	$\frac{-801}{80}$	$\frac{349}{6}$	$\frac{-18353}{120}$	$\tfrac{2391}{10}$	$\frac{-1457}{6}$	$\frac{4891}{30}$	$\frac{-561}{8}$	$\frac{527}{30}$	$\frac{-469}{240}$
	1	1	-4	6	-4	1				
	2	3	-14	26	-24	11	-2			
4	3	<u>35</u> 6	-31		$\frac{-242}{3}$		-19	$\frac{17}{6}$		
	4	$\frac{28}{3}$	$\frac{-111}{2}$		$\frac{-1219}{6}$		$\frac{-185}{2}$	$\frac{82}{3}$	$\frac{-7}{2}$	
	5	<u>1069</u> 80	$\frac{-1316}{15}$	<u>15289</u> 60	$\frac{-2144}{5}$	<u>10993</u> 24	$\frac{-4772}{15}$	2803 20	$\frac{-536}{15}$	$\frac{967}{240}$

Some weights for one-sided approximations at a grid point (generated by setting M = 4, N = 8, $x_0 = 0$ and $\alpha_{\nu} = \{0, 1, 2, 3, 4, 5, 6, 7, 8\}$).

der Orivatiover of	Order of	-1/2		imations inates at 3/2	at $x = 0$; nodes: 5/2	7/2	9/2	11/2	13/2	15/9
1.6	1	1/2	1/2	0/2				11/2	15/2	15/2
	2	$\frac{1}{2}$	$\frac{1}{2}$							
	3	2 3 8	2 3 4	$\frac{-1}{8}$						
	4	_5	<u>15</u>	-5	_1					
0	5	16 _35	16 <u>35</u>	16 -35	16 7	5				
Ů		128 63	$\frac{32}{315}$	-105	32 63	$\overline{128} - 45$	7			
	6	$\frac{\overline{256}}{231}$	256 693	<u>-1155</u>	$\frac{128}{128}$ 231	256 -495	256 77	-21		
	7	1024 _429	512 <u>3003</u>	$\frac{1000}{1024}$	256 3003	1024	512	$\frac{\overline{1024}}{-273}$	22	
	8	2048	2048	2048	$\overline{2048}$	$\frac{-2145}{2048}$	$\frac{1001}{2048}$	2048	$\frac{33}{2048}$	
	9	$\tfrac{6435}{32768}$	<u>6435</u> 4096	$\frac{-15015}{8192}$	<u>9009</u> 4096	$\frac{-32175}{16384}$	<u>5005</u> 4096	$\tfrac{-4095}{8192}$	$\frac{495}{4096}$	$\frac{-429}{32768}$
	2	-1	1							
	3	$\frac{-23}{24}$	$\frac{7}{8}$	$\frac{1}{8}$	$\frac{-1}{24}$					
	4	$\frac{-11}{12}$	$\frac{17}{24}$	$\frac{3}{8}$	$\frac{-5}{24}$	$\frac{1}{24}$				
1	5	$\frac{-563}{640}$	$\frac{67}{128}$	$\frac{143}{192}$	$\frac{-37}{64}$	$\frac{29}{128}$	$\frac{-71}{1920}$			
	6	$\frac{-1627}{1920}$	$\frac{211}{640}$	<u>59</u> 48	$\frac{-235}{192}$	$\frac{91}{128}$	$\frac{-443}{1920}$	$\frac{31}{960}$		
	7	$\frac{-88069}{107520}$	$\frac{2021}{15360}$	$\frac{28009}{15360}$	$\frac{-6803}{3072}$	$\frac{5227}{3072}$	$\frac{-12673}{15360}$	<u>3539</u> 15360	$\frac{-3043}{107520}$	
	8	$\frac{-1423}{1792}$	$\frac{-491}{7168}$	$\frac{7753}{3072}$	$\frac{-18509}{5120}$	$\frac{3535}{1024}$	$\frac{-2279}{1024}$	$\frac{953}{1024}$	$\frac{-1637}{7168}$	$\frac{2689}{107520}$
	1	1	-2	1						
	2	$\frac{3}{2}$	$\frac{-7}{2}$	$\frac{5}{2}$	$\frac{-1}{2}$					
	3	$\frac{43}{24}$	$\frac{-14}{3}$	$\frac{17}{4}$	$\frac{-5}{3}$	$\frac{7}{24}$				
2	4	<u>95</u> 48	$\frac{-269}{48}$	<u>49</u> 8	$\frac{-85}{24}$	59 48	$\frac{-3}{16}$			
	5	<u>12139</u> 5760	<u>-6119</u> 960	<u>3091</u> 384	$\frac{-1759}{288}$	$\frac{1211}{384}$	$\frac{-919}{960}$	739		
	6	25333	-80813	2553	$\frac{-21457}{2304}$	14651	-3687	5760 <u>8863</u>	$\frac{-211}{2004}$	
	7	11520 <u>81227</u>	11520 -67681	256 34151	-16747	2304 <u>5669</u>	1280 -76621	11520 1699	2304 -5647	21719
		35840	8960	2880	1280	512	11520	640	8960	322560
	1 2	-1 -2	3 7	-3 -9	1 5	-1				
	3	$\frac{-23}{8}$	91 8	$\frac{-71}{4}$	<u>55</u>	-43	$\frac{7}{8}$			
3	4	-29	127	4 -29	4 <u>115</u>	8 <u>-133</u>	<u>43</u>	<u>-3</u>		
-	5	<u>-8197</u>	8 <u>39139</u>	-27219	4 <u>19699</u>	8 <u>-15043</u>	8 <u>12099</u>	4 	1237	
	6	1920 -2317	1920 <u>47707</u>	640 <u>-7443</u>	384 <u>158471</u>	384 <u>-30037</u>	<u>32091</u>	<u>-40087</u>	1920 <u>1961</u>	<u>-357</u>
		480	1920	128	1920	384	640	1920	384	640
	1	1 5	-4 -23	6	-4	1 17	-3			
	2	$\frac{5}{2}$ <u>101</u>	$\frac{-23}{2}$ -87	21 373	-19 -319	$\frac{17}{2}$ 273	$\frac{-3}{2}$	41		
4	3	24	4	<u>373</u> 8	$\frac{-319}{6}$	8	$\frac{-47}{4}$	$\frac{41}{24}$. 05	
	4	287 48	$\frac{-1639}{48}$	$\frac{1341}{16}$	$\frac{-5527}{48}$	4613 48	$\frac{-783}{16}$	677 48	$\frac{-85}{48}$	1107
	5	$\frac{14861}{1920}$	$\frac{-1447}{30}$	21299 160	$\frac{-25651}{120}$	$\tfrac{42119}{192}$	$\frac{-2951}{20}$	$\frac{30437}{480}$	$\frac{-1903}{120}$	$\frac{1127}{640}$

Some weights for one-sided approximations at a 'half-way' point (generated by setting M=4, N=8, $x_0=0$ and $\alpha_{\nu} = \{-1/2, 1/2, 3/2, 5/2, 7/2, 9/2, 11/2, 13/2, 15/2\})$.

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Exxon Research and Engineering Company Annandale, New Jersey 08801 *E-mail*: bfornbe@erenj.bitnet

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