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Image Encryption and Authentication With Elliptic Curve Cryptography and Multidimensional Chaotic Maps

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ABSTRACT Many researchers have used the properties of the popular Elliptic Curve Cryptography(ECC) to devise a stronger and faster image encryption algorithm to assure the secrecy of images during online transmission. In this paper, a robust Elliptic curve based image encryption and authentication model for both grayscale and color images has been proposed. The model uses the secure Elliptic Curve Diffie-Hellman(ECDH) key exchange to compute a shared session key along with the improved ElGamal encoding scheme. 3D and 4D Arnold Cat maps are used to effectively scramble and transform the values of plain image pixels. A well-structured digital signature is used to verify the authenticity of the encrypted image prior to decryption. The model produces good-quality cipher images with an average entropy of 7.9993 for grayscale and 7.99925 for the individual components of color images. The model has high average NPCR of 99.6%, average UACI of 33.3% and low correlation for both grayscale and color images. The model has low computational costs with minimized point multiplication operations. The proposed model is robust with high resilience against statistical, differential, chosen-plaintext(CPA), known-plaintext(KPA) and occlusion attacks.

INDEX TERMS Arnold cat map, chaotic map, digital signature, elliptic curve cryptography, image encryption.

I. INTRODUCTION

Extensive amount of data which consists of a variety of images is transmitted back and forth between users online everyday. Ensuring the availability of correct data to the intended receiver with guaranteed secrecy of data from other users in the network is a challenge. Many image encryption algorithms have been presented for fast and powerful real-time encryption of images. The encryption algorithms include both symmetric and asymmetric encryption algorithms. Symmetric encryption is fast with low computational complexity making it suitable for large data sets. The key distribution and management poses a huge overhead on the symmetric encryption models where each user has his own

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secret key which needs to be shared before communication. Asymmetric encryption eliminates this issue by using two different keys for encryption and decryption.

Elliptic Curve Cryptography(ECC) is a popular asymmetric encryption algorithm which provides higher level security with smaller key sizes with reduced resource consumption making it an ideal choice for resource-constraint devices. The unpredictability introduced by chaotic maps in encryption algorithms, in addition makes the model easy to implement, swift and strong against the attacker. Logistic map, Henon map, Baker's map, Arnold Cat map, Sine map and Lorenz system are some of the well-known chaotic systems used in encryption algorithms. Chaos based image encryption schemes show higher efficacy than the simple symmetric or asymmetric image encryption schemes without it. The encryption schemes that use a combination of chaotic maps

perform better than the ones employing single chaotic map. The use of higher dimensional chaotic maps is known to enhance the security and quality of encryption.

Chen et al. [1] proposed a symmetric image encryption algorithm with 3D cat map and logistic map for shuffling of image data with an added layer of confusion between plain-cipher image after every two rounds of shuffling. The encryption also uses Chen’s chaotic system for key generation. Pareek et al. [2] used the outcome of two logistic maps to randomly opt for one of the eight modes designed to encrypt every pixel in the image. Liu and Wang [3] designed a one-time key stream cipher for color image encryption based on Piecewise Linear chaotic map(PWLCM) to compute the key stream for color image encryption. The generation of key stream depends on another sequence obtained from the Chebyshev map.

Liu et al. [4] devised an encryption algorithm with DNA encoding where the XOR and complement operation is carried out using the random sequences generated by the logistic map to encrypt individual components of the color image. However, Liu et al. [5] proved that the algorithm can be broken by the known-plaintext attack with only one known pair of plain-cipher images. Luo et al. [9] used two chaotic sequences generated by Tent map for the permutation and diffusion of pixels in grayscale images. The initial seed value of Tent map depends on the plain image. Singh and Singh [10] proposed an Elliptic curve(EC) block based image encryption scheme with digital signature for cipher image authentication. The scheme faces the cipher data expansion issue due to additional amount of pixels added during encryption. Singh and Singh [11] presented another refined Elliptic curve based algorithm which uses 2D Arnold Cat map for pixel shuffling and improved ElGamal encryption encoding techniques. The scheme eliminated the data expansion issue.

Xu et al. [12] designed a block based image encryption that uses a quantized chaos matrix obtained from the logistic map for creating X coordinate, Y coordinate and swapping control table to swap pixels together with dynamic indexes for diffusion. Saljoughi and Mirvaziri [22] designed an encryption model which uses 3D logistic maps and XOR operation for row-column pixel permutation. Broumandnia [24] presented a color image encryption that uses a 3D Modular Chaotic Map(3DMCM) for pixel shuffling with exclusive OR and circular shift operation for pixel substitution. Kumar and Girdhar [26] used the Lorenz-Rosler chaotic system for pixel diffusion and 2D logistic maps for confusion as well as DNA cryptography for encryption of pixels of color images.

The paper is organized as follows. Section II delves into the theoretical view of our model. In Section III, the proposed EC based image encryption algorithm and the digital signature is described in detail. In section IV, we review the experimental results acquired from the proposed model with the two scrambling methods. In Section V, we analyse the strength of the proposed model against various statistical and intruder attacks. In Section VI, we discuss the performance

of the proposed model in comparison to the other schemes prevailing in the field of image encryption. In Section VII, the conclusions follow.

II. PRELIMINARIES

A. ELLIPTIC CURVE CRYPTOGRAPHY

The elliptic curve, E in Fig.1 is a non-singular algebraic plane curve defined over a finite field F_p that consists of a set of points satisfying the well known Weierstrass equation of form [27],

$$y^2 \text{ mod } p = x^3 + ax + b \text{ mod } p \tag{1}$$

where $x, y, a, b \in F_p$, and $4a^3 + 27b^2 \text{ mod } p \neq 0$ with a point of infinity, O .

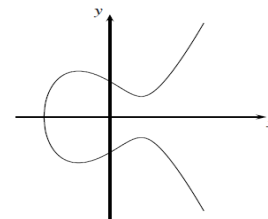


FIGURE 1. An Elliptic Curve E defined over a finite field F_p .

1) GROUP OPERATIONS ON ELLIPTIC CURVE

We assume P and Q as two unique points on the elliptic curve E . The group operations on the curve E are as follows.

- Point addition: The point addition of P and Q is given by

$$P + Q = T \tag{2}$$

The line that joins the given points P and Q intersects the curve E at the point $(-T)$. The mirror image of the point $(-T)$ with respect to the abscissa is thus obtained to be T .

- Point subtraction: The point subtraction of P and Q is given by

$$P - Q = P + (-Q) \tag{3}$$

where $Q = (x, y)$ then $(-Q) = (x, -y)$.

- Point doubling: The addition of the curve point P to itself results in a point S on the curve E . This is known as point doubling and is given by

$$P + P = 2P = S \tag{4}$$

The tangent drawn at the point P intersects the curve at $(-S)$. The mirror image of the point, $(-S)$ with respect to x -axis is the point S .

- Point multiplication: The point multiplication of the point P with a scalar, m is computed by performing m repetitive point additions of P with itself.

$$mP = P + P + P + \dots + P (m \text{ times}) \tag{5}$$

where m is a scalar, $m \in \mathbb{Z}_p$.

2) ELLIPTIC CURVE DIFFIE–HELLMAN (ECDH) KEY EXCHANGE

The ECDH key exchange is used for secure exchange of a shared session key between two users in a network. The key exchange is carried out as follows.

- User A generates its random secret key, n_A . A then computes the public key P_A using his private key, n_A and generator point, G on the curve as,

$$P_A = n_A * G \tag{6}$$

and sends it to user B.

- User B generates its random secret key, n_B . B then computes the public key P_B using his private key, n_B and generator point, G on the curve as,

$$P_B = n_B * G \tag{7}$$

and sends it to user A.

- Both the users A and B compute the shared ECDH key, Z_{AB} as in (8) and (9). A computes Z_{AB} as follows.

$$Z_{AB} = n_A * P_B = n_A * n_B * G \tag{8}$$

B computes Z_{AB} as follows.

$$Z_{AB} = n_B * P_A = n_B * n_A * G \tag{9}$$

B. ARNOLD'S CAT MAP

Arnold's Cat map [28] is a popular discrete chaotic map that performs scrambling of image pixels by P repetitive shear mapping on the input image to obtain the originally taken image again. The parameter P is defined as the periodicity of the transformation map. The cat map is known for its remarkable mixing properties and high sensitivity to the initial state and control parameters of the map.

1) 2D ARNOLD'S CAT MAP

The two-dimensional Arnold's cat map can be expressed as follows.

$$\begin{pmatrix} x' \\ y' \end{pmatrix} = \begin{pmatrix} 1 & a \\ b & ab + 1 \end{pmatrix} \cdot \begin{pmatrix} x \\ y \end{pmatrix} \text{mod } N \tag{10}$$

where the positive integers, a and b are the control parameters and N is the dimensional value of a $N \times N$ image.

2) EXTENSION TO HIGHER DIMENSIONS

The 2D Arnold's Cat map can be extended to higher dimensional mappings to perform more secure and efficient transforms. The three-dimensional Arnold's mapping is performed with the introduction of two new control parameters, c and d [1]. The general 3D extended form of existing 2D cat map is as follows.

$$\begin{pmatrix} x' \\ y' \\ z' \end{pmatrix} = \begin{pmatrix} 1 & a & 0 \\ b & ab + 1 & 0 \\ c & d & 1 \end{pmatrix} \cdot \begin{pmatrix} x \\ y \\ z \end{pmatrix} \text{mod } N \tag{11}$$

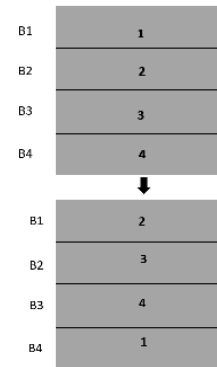


FIGURE 2. 4D Horizontal Block Reordering of Grayscale image.

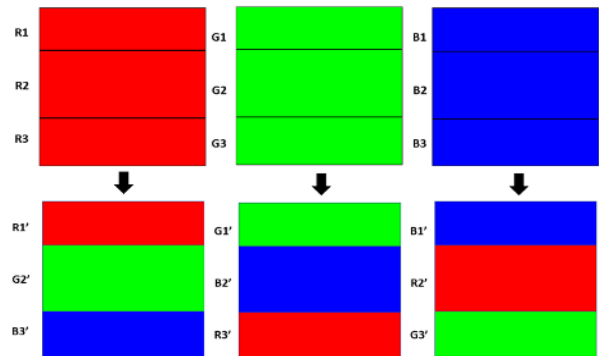


FIGURE 3. 4D Horizontal Block Reordering of a RGB image.

Liu et al. [29] modified the general form of 3D Arnold's map to carry out pixel position scrambling with subsequent value transform as given in (12).

$$\begin{cases} \begin{pmatrix} x' \\ y' \end{pmatrix} = \begin{pmatrix} 1 & a \\ b & ab + 1 \end{pmatrix} \cdot \begin{pmatrix} x \\ y \end{pmatrix} \text{mod } N \\ p' = (cx + dy + p) \text{mod } M \end{cases} \tag{12}$$

where p and p' in equation (12) are the pixel values of the initial and remodeled images, and M is maximum pixel intensity. The values of pixel intensity in an image varies from 0 to 255. The modified 3D cat map can also be extended to 4D mapping by reordering the blocks of pixels between the image components. The block reordering can be carried out either vertically [31] or as horizontally [32]. Fig.2 and Fig.3 show the horizontal block reordering for a grayscale image and shuffling among the individual components of a color(RGB) image.

III. PROPOSED MODEL

The proposed model is divided into three phases viz. ECDH Key Exchange, Encryption/Decryption, and Signing/Verification of encrypted data. The encryption is carried out on blocks of pixels instead of individual pixels to increase efficiency and speed of encryption. The model uses the improved ElGamal encryption encoding [11] for encryption of scrambled pixels. We assume A and B are two users in the network where Sender A wants to transmit data(i.e.images) to

Receiver B in a secure manner. The private and public key pair of Sender, A and Receiver, B are (n_A, P_A) and (n_B, P_B) respectively. A and B compute the ECDH session key, Z_{AB} through their private-public key pairs using equations (8) and (9).

A digital image is characterised by a matrix of smaller image elements called pixels. The number of pixels that can be grouped together is limited by the size of the prime elliptic curve parameter (p) . A 512-bit elliptic curve limits the maximum number of pixels in a group to 64.

The scrambling of pixels using Arnold's Cat map can be performed in two ways.

- **Scheme 1:** 3D Cat map transform
- **Scheme 2:** 4D Cat map transform

The proposed design of the image encryption model, as shown in the Fig. 4 and 5, is explained as follows.

A. ENCRYPTION

- 1) Record the pixel intensities, dimensional size and channel information from the original image. Adjust the values of pixels by random subtraction or addition of 1 or 2 to each pixels.
- 2) Choose a random integer k between 2 to $(n - 1)$, where n is the cyclic order of the 512-bit elliptic curve.
- 3) Scrambling of pixels can be carried out in following ways.

- **Scheme 1: 3D Cat Map Scrambling**

Shuffle the pixels for j rounds using modified 3D Cat map equation in (12) with control parameters $a = 1$ and $b = 1$.

$$j = kG_x \text{ mod } P \quad (13)$$

where $kG_x = x$ -coordinate of kG and P is the period of the 3D Arnold scrambling for a given image of size $N \times N$.

- **Scheme 2: 4D Cat Map Scrambling**

After shuffling the pixels for j rounds using 3D Cat map, perform horizontal block reordering of the scrambled image as described in Fig.2 and Fig.3.

- 4) Partition the scrambled image pixels into groups of 64 each. Convert each group into a single large integer value with 256 as base.
- 5) Pair up every successive two integer values to form the plain text input list, P_M . Compute cipher text list, P_C as follows.

$$P_C = P_M + (kG + Z_{AB}) \quad (14)$$

- 6) Convert cipher text list, P_C to byte values with base 256. Ensure each list has 64 pixel values with required left zeroes padding.
- 7) Convert the cipher pixel values into the cipher image, C .

B. SIGNING THE ENCRYPTED IMAGE

- 1) Compute the hash value h_{xy} from kG as follows.

$$k_{xy} = kG_x \oplus kG_y \quad (15)$$

$$h_{xy} = SHA_{256}(k_{xy}) \quad (16)$$

- 2) Perform hash of the encrypted image C as h_C and concatenate h_C with h_{xy} to obtain the combined hash value H .

$$h_C = SHA_{256}(C) \quad (17)$$

$$H = (h_C || h_{xy}) \quad (18)$$

- 3) Compute the parameter, k' as follows.

$$Z = Z_{AB_x} \oplus Z_{AB_y} \quad (19)$$

$$k' = Z \oplus k \quad (20)$$

- 4) Calculate the signature (R, U) as follows.

$$R = SHA_{256}(H) \quad (21)$$

$$U = ((k') - R) \text{ mod } n \quad (22)$$

- 5) The encrypted image C is sent to the receiver along with digital signature, (R, U) .

C. VERIFICATION OF ENCRYPTED IMAGE

- 1) Compute the value of kG from (R, U) and Z_{AB} .

$$Z = Z_{AB_x} \oplus Z_{AB_y} \quad (23)$$

$$k = ((R + U) \oplus Z) \quad (24)$$

$$kG = (k * G) \quad (25)$$

- 2) Calculate the hash value h_{xy} from kG .

$$k_{xy} = kG_x \oplus kG_y \quad (26)$$

$$h_{xy} = SHA_{256}(k_{xy}) \quad (27)$$

- 3) Perform hash of the encrypted image C as h_C . Combine h_C with h_{xy} to obtain the hash value H .

$$h_C = SHA_{256}(C) \quad (28)$$

$$H = (h_C || h_{xy}) \quad (29)$$

- 4) Calculate R' from computed hash values. If R' is equal to R received digital signature (R, U) from Sender A, then the signature is verified.

$$R' = SHA_{256}(H) \quad (30)$$

$$R' == R \quad (31)$$

D. DECRYPTION

- 1) Record the pixel intensities, size and related information from the encrypted image, C .
- 2) Partition the encrypted image pixels into groups of 64 each. Convert each group into a single large integer value with 256 as base.
- 3) Pair up the integer values as cipher text list, P_C process and compute plain text list, P_M as follows.

$$P_M = P_C - (kG + Z_{AB}) \quad (32)$$

- 4) Convert calculated plain text list, P_M back to byte values with base as 256 and ensure each list has 64 values with necessary left zero padding.

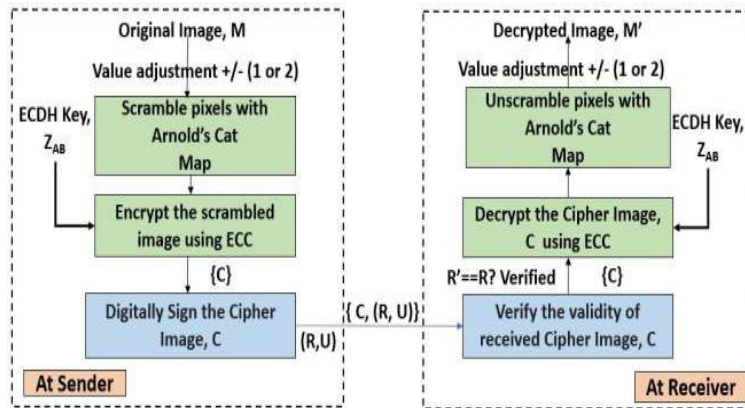


FIGURE 4. Overview of the Proposed Model.

- 5) Unscramble the scrambled pixels in two ways.
 - **Scheme 2:4D Cat map Unscrambling**
Reorder the blocks of scrambled image back to generate the 3D scrambled image.
 - **Scheme 1:3D Cat map Unscrambling**
Unscramble the scrambled pixel values for P-j rounds to obtain the value adjusted image.
- 6) Adjust the values of pixels back by random subtraction or addition of 1 or 2 to each pixel to attain the original image again.

IV. EXPERIMENTAL RESULTS

The proposed model is implemented using Python 3.7 with Spyder 4.1.5 IDE on a HP laptop with system configuration as Intel(R) Core(TM) i5-8250U CPU@1.80GHz with 8GB RAM. The Elliptic curve used in the proposed algorithm is the secure 512-bit curve from ECC Brainpool [33]. The elliptic curve parameters used for implementation are depicted in Table 1.

V. SECURITY ANALYSIS

Various security and statistical tests are carried out to explore the strength of the proposed algorithm with the two implementations of the algorithm using the explained scrambling methods, Scheme 1 and Scheme 2. The proposed model is tested on grayscale and color(RGB) images sourced from the SIPI Image Database [34]. The analyses carried out include histogram analysis, variance analysis, chi-square test, key space analysis, Shannon’s entropy of cipher images, measuring resistance to differential attacks with NPCR(Number of Pixel Change Rate) and UACI(Unified Average Changed Intensity) parameters, similarity measurement (PSNR and SSIM), resistance to intruder attacks, correlation coefficient analysis and computation cost comparison.

A. HISTOGRAM ANALYSIS

Histogram of an image plots the frequency of pixel intensities in the ordinate against the values of the pixel intensities in the abscissa. An effective encryption algorithm produces a

TABLE 1. Elliptic curve parameters used in the implementation of proposed scheme.

Parameter	Values
p	8948962207650232551656602815159153422162609644098354511344597187200057010413552439917934304191956942765446530386427345937963894309923928536070534607816947
a	6294860557973063227666421306476379324074715770622746227136910445450301914281276098027990968407983962691151853678563877834221834027439718238065725844264138
b	3245789008328967059274849584342077916531909009637501918328323668736179176583263496463525128488282611559800773506973771797764811498834995234341530862286627
G _x	6792059140424575174435640431269195087843153390102521881468023012732047482579853077545647446272866794936371522410774532686582484617946013928874296844351522
G _y	6592244555240112873324748381429610341312712940326266331327445066687010545415256461097707483288650216992613090185042957716318301180159234788504307628509330

cipher image with even distribution of frequencies across the histogram graph plot. Grayscale images taken in consideration are the Lena, Barbara, and Mandrill. RGB images taken as input images include Lena, House, and Mandrill. Fig. 6 and Fig.7 illustrate the plots for input plain grayscale and RGB images with their cipher images from Scheme 1 and 2 scrambling methods respectively. The cipher images show the uniform frequency distribution of pixels across the histograms indicating good quality of cipher images.

B. HISTOGRAM VARIANCE ANALYSIS

Variance is a quantitative measure of resistance against statistical attacks by analysing the distribution of pixel frequencies in the histogram of a grayscale image. Variance of a histogram can be computed as follows.

$$VAR(Z) = \frac{1}{P^2} \sum_{i=1}^P \sum_{j=1}^P \frac{(p_i - p_j)^2}{2} \tag{33}$$

P is the grayscale pixel intensity value and the parameters, p_i and p_j are the number of pixels with intensity value i and j.

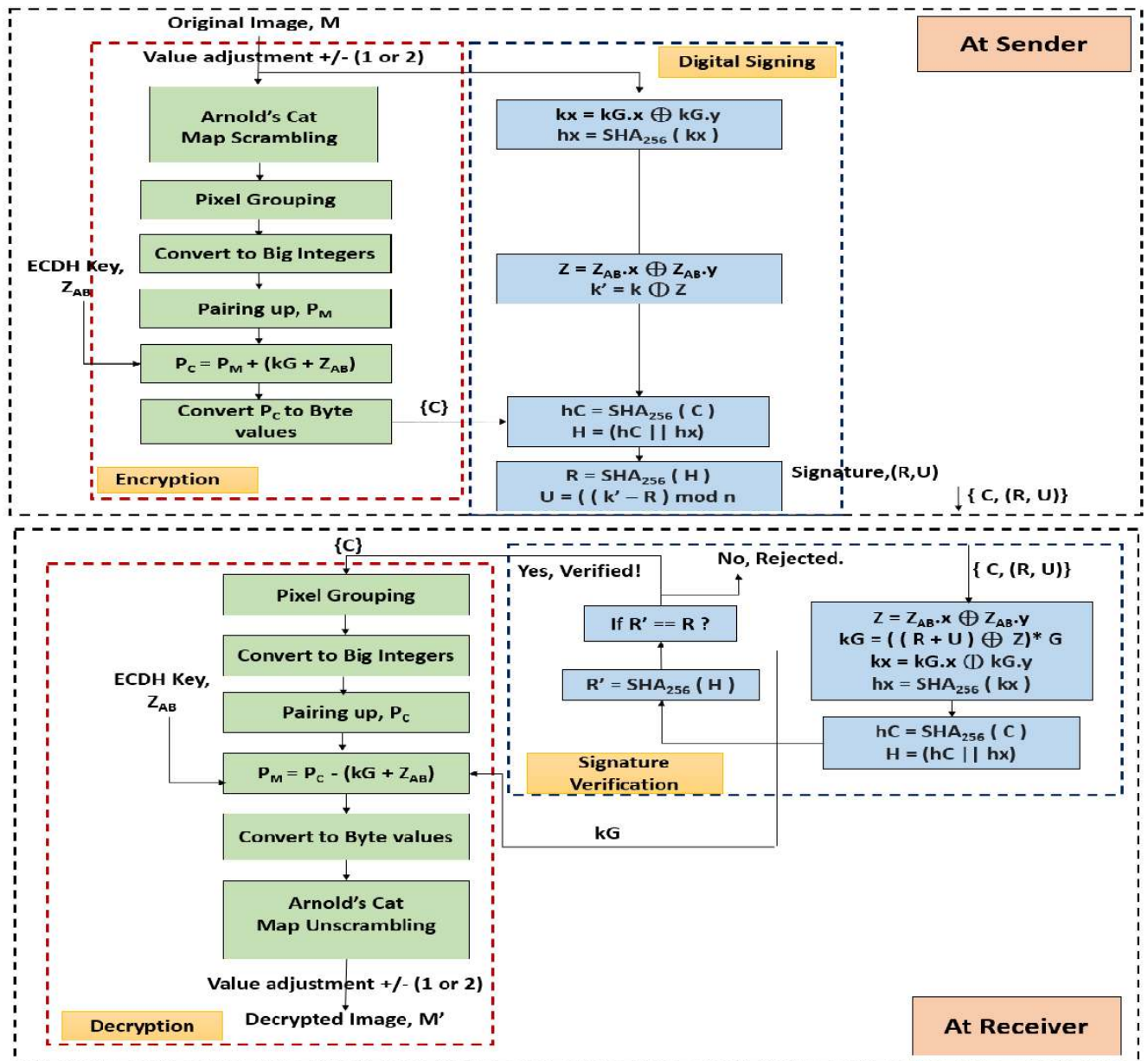


FIGURE 5. Proposed design model.

Lower is the computed variance, higher is the uniformity of pixel distribution of the histogram. From Table 2, we infer that the cipher images obtained from the proposed model have lower values of variance indicating high resistance to statistical attacks.

C. CHI-SQUARE TEST ANALYSIS

Chi-square test measures the uniformity of pixel distribution in the histogram of grayscale images quantitatively. The Chi-square value can be evaluated as follows.

$$\chi_{test}^2 = \sum_{i=0}^{255} \frac{(OF_i - EF_i)^2}{EF_i} \tag{34}$$

where the OF_i and EF_i are the observed and expected values of frequency of pixel intensity value i . EF_i is

$$EF_i = \frac{M \times N}{I_{max}} \tag{35}$$

I_{max} is the total number of pixel intensities for a grayscale image. Table 3 depicts the calculated values of Chi-square for cipher images obtained from the proposed model with significance level of 1% and 5% at 255 degrees of freedom. The ideal values of Chi-square at 1% and 5% significance level are $\chi_{255,0.01}^2 = 310.457$ and $\chi_{255,0.05}^2 = 293.2478$. Lower the value of Chi-square, higher is the uniformity of pixel distribution in the histogram. At 1% and 5%

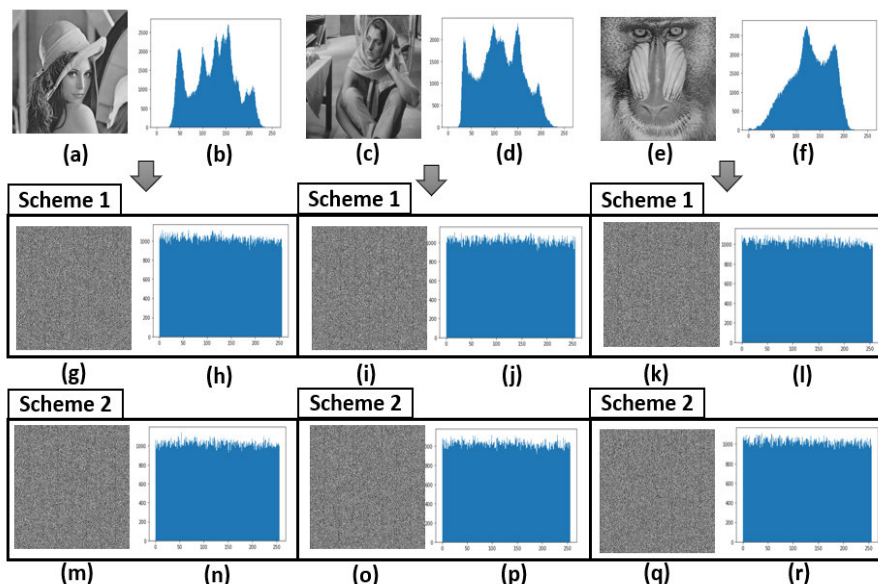


FIGURE 6. (a-f):Plain grayscale images with their histograms, (g-l):Cipher images with Scheme 1 scrambling with their histograms, (m-r):cipher images with Scheme 2 scrambling with respective histogram.

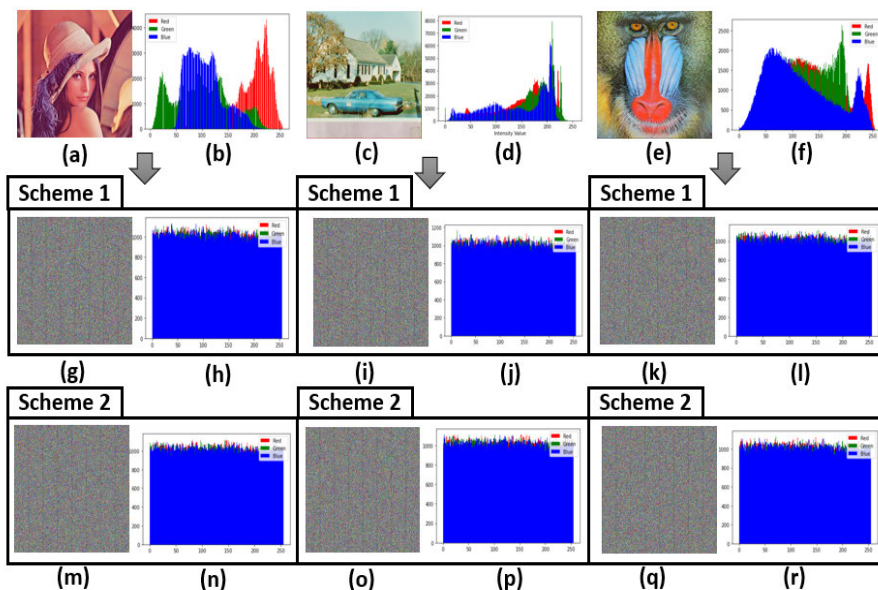


FIGURE 7. (a-f):Plain RGB images with their histograms, (g-l):Cipher images with Scheme 1 scrambling with their histograms, (m-r):Cipher images with Scheme 2 scrambling with respective histogram.

significance levels, the cipher images obtained from Scheme 1 and 2 show Chi-square values lower than the ideal values. The calculated Chi-square values satisfy the hypothesis and implies grayscale pixel uniformity in the cipher images derived from the proposed model.

D. KEY SPACE

The robustness of an algorithm depends upon the key size of the cryptosystem. The proposed algorithm relies on the inability of solving the exponentially difficult Elliptic Curve

Discrete Logarithm Problem(ECDLP), in feasible amount of time, to provide higher level of security with shorter key sizes. The implementation of the proposed model uses the ECC Brainpool [33] standard 512-bit Elliptic curve which can successfully withstand brute-force attacks.

E. SHANNON’S ENTROPY ANALYSIS

The Shannon’s theory [35] states that entropy is directly proportional to the degree of uncertainty present in the data. An encryption algorithm should ideally produce a cipher

TABLE 2. Variance of various grayscale images.

Scheme	Image	Variance	
		Plain	Cipher
Scheme 1	Lena	632097.48	958.04
	Barbara	576404.48	991.95
	Mandrill	750395.63	1026.14
	Peppers	555344.7	1003.05
Scheme 2	Lena	632097.48	876.32
	Barbara	576404.48	977.73
	Mandrill	750395.63	942.47
	Peppers	555344.7	934.98
Ref. [23]	Lena	634734	980.8
	Barbara	433040	1013.2
	Mandrill	627520	1008.3
	Peppers	448850	950.9

TABLE 3. Chi-square analysis of proposed model.

Scheme	Images	χ^2_{test}	Results
Scheme 1	Lena	275.89	Pass
	Barbara	263.87	Pass
	Mandrill	286.44	Pass
	Peppers	282.01	Pass
Scheme 2	Lena	251.52	Pass
	Barbara	247.37	Pass
	Mandrill	266.72	Pass
	Peppers	241.97	Pass

image with entropy value of 8. The entropy values are computed for the cipher images obtained from both implementations of the model in Table 4 and 5. The cipher images of the tested grayscale as well as color images display better entropy values than existing image encryption schemes.

F. RESISTANCE TO DIFFERENTIAL ATTACKS

Differential attacks are a form of cryptanalysis to find the secret key by tracing the differences in the cipher data due to minimal changes in the plain data. The Number of Changing Pixel Rate(NPCR) and Unified Average Changed Intensity(UACI) values, computed from the cipher images acquired from a pair of nearly similar plain images, are used to measure the resilience of the algorithm against such attacks. The ideal value of NPCR for an encryption is 100% with UACI value being around 33%. From Table 6 and 7, it can be ascertained that the cipher images derived from the scheme approximate near-ideal values for the NPCR and UACI parameters.

$$NPCR = \frac{\sum_{i,j} D(i,j)}{T} \times 100\% \tag{36}$$

$$UACI = \frac{\sum_{i,j} |C(i,j) - C'(i,j)|}{I_{max} \times T} \times 100\% \tag{37}$$

where

$$D(i,j) = \begin{cases} 0 & \text{if } C(i,j) = C'(i,j) \\ 1 & \text{if } C(i,j) \neq C'(i,j) \end{cases} \tag{38}$$

T = Total number of pixels, C and C' are the two cipher images, and I_{max} = Total number of Pixel Intensities.

G. SIMILARITY MEASUREMENT

The Peak Signal-to-Noise ratio(PSNR) and Structural Similarity(SSIM) index values measure the similarity between the plain and cipher images. The PSNR value edges nearer to infinity as Mean Squared Error(MSE) approaches zero. SSIM index estimates the perceived quality of images from its structural information. Lower the PSNR and SSIM values are between plain and cipher data, better is the quality of encryption. The measured values of PSNR and SSIM for the proposed schemes are given in Table 8. The SSIM values of cipher images from the tested grayscale and RGB images approach zero with reduced values of PSNR indicating good quality of encryption. The Scheme 2 has lower similarity between original and cipher images than Scheme 1 making it a better mode of encryption.

$$PSNR = 20 \times \log_{10} \frac{P_{max}}{\sqrt{MSE}} \tag{39}$$

where P_{max} is the maximum pixel intensity.

H. KNOWN-PLAINTEXT ATTACK(KPA) AND CHOSEN-PLAINTEXT ATTACK(CPA)

Known-Plaintext and Chosen-Plaintext attacks are two attack models for cryptanalysis where intruder can exploit the knowledge of one or more plaintext-ciphertext pairs from an encryption method to reveal secret information. In case of Known-Plaintext attacks, we assume the intruder has access to one or more plain images and their corresponding cipher images. The proposed algorithm generates different and unique cipher images for the same plain image due to use of the random parameter k, value adjustment, and efficient scrambling methods in each session of encryption.

In Chosen-Plaintext attacks, the intruder can choose certain plaintext images and use the respective cipher images to unveil the secret key. We assume the intruder chooses two special images i.e. all black and all white images as the plain images for the CPA attack. From Fig. 8 and Tables 9 and 10, we infer that the proposed method produces good quality cipher images for the special input images with low correlation, uniform cipher pixel distribution and high entropy

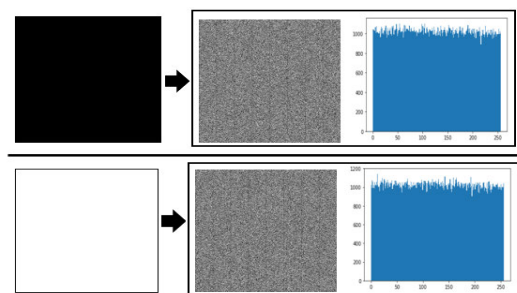


FIGURE 8. Plain All Black and All White images with their respective encrypted images obtained from the proposed encryption model, and the histograms of the encrypted images.

TABLE 4. Shannon’s entropy analysis for grayscale images.

Image	Entropy Value							
	Scheme 1	Scheme 2	Ref. [23]	Ref. [9]	Ref. [14]	Ref. [19]	Ref. [18]	Ref. [20]
Lena	7.9993	7.9993	7.9993	7.9993	7.9991	7.9992	7.9992	7.9991
Barbara	7.9993	7.9993	7.9993	-	-	-	-	-
Peppers	7.9993	7.9994	7.9994	7.9992	-	-	-	-
Mandrill	7.9993	7.9993	7.9993	7.9992	-	-	-	-

TABLE 5. Shannon’s entropy analysis for RGB images.

Image	Component	Entropy Value					
		Scheme 1	Scheme 2	Ref. [39]	Ref. [40]	Ref. [41]	Ref. [51]
Lena	Red	7.9992	7.9993	7.9992	7.9972	7.9993	7.9966
	Green	7.9992	7.9992	7.9993	7.9973	7.9993	7.9972
	Blue	7.9992	7.9993	7.9992	7.9972	7.9993	7.9967
Mandrill	Red	7.9993	7.9992	7.9991	7.9972	7.9923	-
	Green	7.9992	7.9992	7.9991	7.9972	7.9803	-
	Blue	7.9992	7.9992	7.9993	7.9972	7.9986	-
Peppers	Red	7.9992	7.9993	7.9989	7.9971	7.9962	7.9910
	Green	7.9992	7.9993	7.9991	7.9975	7.9929	7.9918
	Blue	7.9993	7.9993	7.9989	7.9974	7.9925	7.9905

TABLE 6. NPCR and UACI analysis for grayscale tested images.

Type	Grayscale	Scheme 1	Scheme 2	Ref. [7]	Ref. [14]	Ref. [23]	Ref. [38]
NPCR (%)	Lena	99.61	99.632	99.60	99.61	99.6113	99.59
	Barbara	99.63	99.62	-	-	99.5796	-
	Peppers	99.62	99.63	-	-	99.6109	-
	Mandrill	99.62	99.61	-	-	99.6112	-
UACI (%)	Lena	33.34	33.34	33.44	33.32	33.4682	33.41
	Barbara	33.35	33.34	-	-	33.4296	-
	Peppers	33.31	33.33	-	-	33.4836	-
	Mandrill	33.34	33.34	-	-	33.4919	-

TABLE 7. Average values for NPCR and UACI analysis for RGB tested images.

Type	RGB		Scheme 1	Scheme 2	Ref. [49]	Ref. [50]	Ref. [51]
Average NPCR (%)	Lena	Red	99.6200	99.6223	99.6052	99.6100	99.6001
		Green	99.6105	99.6125	99.6060	99.6100	99.5998
		Blue	99.6097	99.6185	99.6113	99.6100	99.5997
	Mandrill	Red	99.5937	99.6243	99.6024	99.6200	99.6099
		Green	99.6223	99.6132	99.6252	99.6200	99.6058
		Blue	99.6105	99.6059	99.6004	99.6200	99.5956
	Peppers	Red	99.6269	99.6093	99.6060	99.6100	-
		Green	99.6028	99.6292	99.6286	99.6100	-
		Blue	99.6013	99.6128	99.5874	99.6100	-
Average UACI (%)	Lena	Red	33.2830	33.3438	33.4280	32.2300	33.3575
		Green	33.3327	33.2787	33.4966	33.9600	33.4287
		Blue	33.3277	33.3651	33.3779	34.3500	33.3683
	Mandrill	Red	33.3216	33.3588	33.4311	33.4600	33.3743
		Green	33.3115	33.3359	33.4500	32.2200	33.3829
		Blue	33.2757	33.3345	33.4935	34.2300	33.5604
	Peppers	Red	33.3182	33.3151	33.4959	31.0300	-
		Green	33.3691	33.3561	33.4874	34.0000	-
		Blue	33.3692	33.3614	33.4302	34.7800	-

making it difficult for the intruder to elicit secret information from the plain-cipher image pairs.

I. OCCLUSION ATTACK ANALYSIS

The encrypted images can suffer from data loss due to varying reasons while transmitting data between sender and receiver. The ability to recognise the original images from the distorted images signifies the strength of the proposed model against such occlusion attacks. Figures 9 and 10 depict the

decrypted grayscale and color images acquired from the distorted cipher images through Scheme 1 and 2. The decrypted images show that the proposed model can resist the occlusion attacks successfully.

J. CORRELATION COEFFICIENT ANALYSIS

Correlation coefficient is a statistical measure of strength of relationship between the original and encrypted image.

TABLE 8. Similarity analysis between two proposed schemes.

Image Type	Image	PSNR		SSIM	
		Scheme 1	Scheme 2	Scheme 1	Scheme 2
Grayscale	Lena	27.897	27.894	0.0097	0.0091
	Barbara	27.895	27.894	0.0094	0.0087
	Mandrill	27.902	27.898	0.0092	0.0078
	Peppers	27.903	27.898	0.0086	0.0084
RGB	Lena	27.899	27.895	0.0099	0.0098
	Mandrill	27.896	27.894	0.0097	0.0087
	Peppers	27.9	27.896	0.0086	0.0084

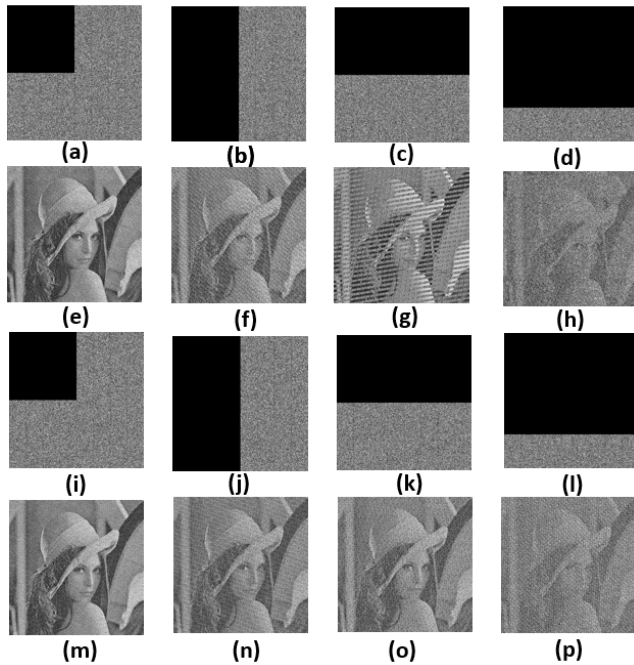


FIGURE 9. Decrypted grayscale images from Scheme 1 and 2 under Occlusion attack analysis. (a-d) are cipher images from Scheme 1 having 25%, 50% vertical, 50% horizontal and 75% data loss. (e-h) are the corresponding decrypted images for Scheme 1 Occlusion analysis. (i-l) are cipher images from Scheme 2 having 25%, 50% vertical, 50% horizontal and 75% data loss. (m-p) are the corresponding decrypted images for Scheme 2 Occlusion analysis.

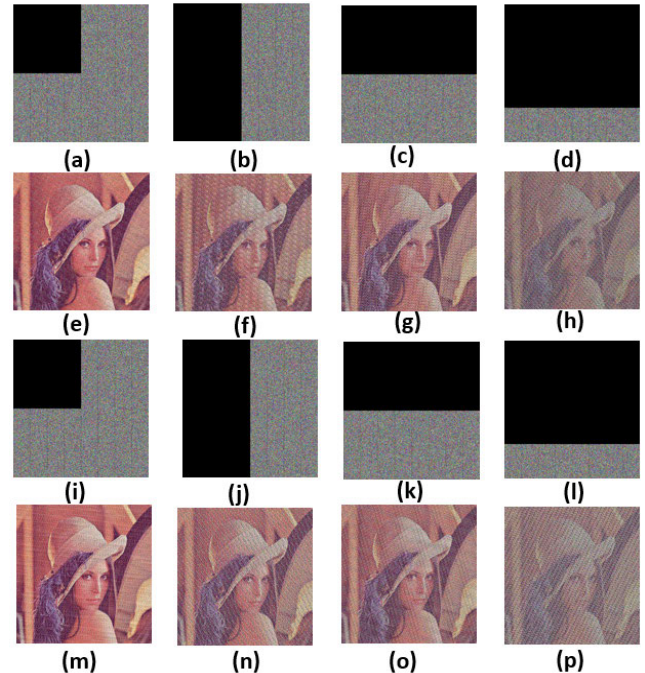


FIGURE 10. Decrypted color images from Scheme 1 and 2 under Occlusion attack analysis. (a-d) are cipher images from Scheme 1 having 25%, 50% vertical, 50% horizontal and 75% data loss. (e-h) are the corresponding decrypted images for Scheme 1 Occlusion analysis. (i-l) are cipher images from Scheme 2 having 25%, 50% vertical, 50% horizontal and 75% data loss. (m-p) are the corresponding decrypted images for Scheme 2 Occlusion analysis.

TABLE 9. Histogram uniformity analysis for special input images.

Image	Proposed		Variance	χ^2_{test}	Results
All black	Scheme 1	Plain	267386880	-	Pass
		Cipher	1000.22	286.54	
All white	Scheme 1	Plain	267386880	-	Pass
		Cipher	1007.16	260.99	
All black	Scheme 2	Plain	267386880	-	Pass
		Cipher	999.35	260.89	
All white	Scheme 2	Plain	267386880	-	Pass
		Cipher	883.95	254.62	

A plain image has high correlation between the adjoining pixels. The correlation graph for the plain image as a consequence is dense. A strong and desirable cipher image should have an evenly distributed graph with low correlation values between the adjoining pixels. We compute correlation coefficient between the plain and cipher images horizontally,

vertically and diagonally. The correlations graph plots for the grayscale Lena image and its cipher image from the proposed model are illustrated in Fig. 11. Fig. 12 shows the correlation graph plots for the component blue of color Lena image and its cipher image. Table 11 and 12 display the comparison of correlation coefficient values obtained from Scheme 1 and Scheme 2 implementation of the tested grayscale and RGB images with those from recent image encryption algorithms. The correlation coefficient values for tested grayscale and RGB images are closer to zero in comparison to the existing encryption models.

K. COMPUTATIONAL COST COMPARISON

A robust image encryption algorithm should perform encryption on images with minimal consumption of computing resources. The time consumption of an EC based encryption

TABLE 10. Performance of proposed model on special input images.

Image	Proposed		Entropy	Correlation Coefficient		
				Horizontal	Vertical	Diagonal
All black	Scheme 1	Plain	0	-	-	-
		Cipher	7.9992	0.00171	0.00717	-0.00133
All white	Scheme 1	Plain	0	-	-	-
		Cipher	7.9993	0.00159	0.0043	0.00082
All black	Scheme 2	Plain	0	-	-	-
		Cipher	7.99928	-0.0023	0.0069	-0.0025
All white	Scheme 2	Plain	0	-	-	-
		Cipher	7.99933	0.0035	0.00046	0.00057

TABLE 11. Comparison of correlation coefficient values for grayscale images.

Algorithm	Image	Horizontal		Vertical		Diagonal	
		Plain	Cipher	Plain	Cipher	Plain	Cipher
Scheme 1	Barbara	0.85973	-0.0016	0.95908	0.0031	0.84181	-0.0012
	Lena	0.97189	0.0019	0.98498	0.0017	0.95928	0.0011
	Mandrill	0.86652	-0.00077	0.75864	0.0019	0.72613	0.00068
	Peppers	0.97918	0.0011	0.98264	0.0047	0.96797	-0.00024
Scheme 2	Barbara	0.85973	-0.0015	0.95908	0.0016	0.84181	0.00073
	Lena	0.97189	0.0012	0.98498	0.0016	0.95928	-0.0007
	Mandrill	0.86652	-0.0002	0.75864	0.0036	0.72613	0.0002
	Peppers	0.97918	-0.0012	0.98264	0.0024	0.96797	0.00021
Ref. [12]	Lena	0.9503	-0.0226	0.9775	0.0041	0.9275	0.0368
Ref. [18]	Lena	0.9325	0.0074	0.9139	-0.0094	0.9469	-0.0054
Ref. [23]	Barbara	0.9689	0.0024	0.8956	0.0031	0.8536	-0.0013
	Lena	0.9858	0.0019	0.9801	-0.0024	0.9669	-0.0011
	Mandrill	0.7251	0.0024	0.8558	0.0011	0.6920	-0.0008
	Peppers	0.9807	-0.0028	0.9752	0.0039	0.9636	-0.00024
Ref. [36]	Lena	0.9771	0.0925	0.9631	0.0430	0.9469	-0.0054
Ref. [37]	Peppers	0.9295	0.0048	0.9294	0.0062	0.8771	0.0030
Ref. [38]	Mandrill	0.7508	-0.0061	0.8562	0.0130	0.7153	0.0017

algorithm is determined by the number of time intensive point multiplications it has. Lesser the number of point multiplications, faster is the image encryption algorithm. In Table 13, we carried out a comparative analysis of time consumption among recent EC based image encryption algorithms. The proposed model has lower computational costs with lesser number of point multiplications.

VI. DISCUSSION

The performance of the proposed model is tested on two implementations of model, Scheme 1 using 3D Arnold Cat map and Scheme 2 using the 4D Arnold Cat map. The extended Arnold Cat maps add additional layers of value substitution after the numerous rounds of pixel scrambling which increases the security as well as quality of encryption. The model generates better-quality cipher images as evident from the uniform pixel frequency distribution in histogram analysis as can be seen in the Fig. 6 and Fig.7. The cipher images obtained are unique and distinct with high entropy, low PSNR and SSIM values and reduced correlation coefficient values between neighbouring pixels. The model has high resistance to differential attacks with an average NPCR of 99.62% and UACI of 33.3%. From the performance comparisons between the proposed model and existing encryption schemes in Tables 4–13, we infer that the proposed scheme has higher efficacy with stronger resistance against the security and statistical attacks.

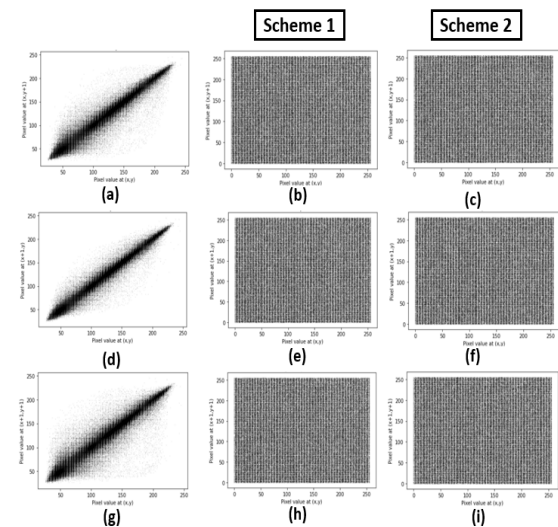


FIGURE 11. Correlation graphs of plain and cipher grayscale LENA images from their respective proposed Scheme 1 and Scheme 2 implementations. Here, (a - c) are horizontal correlation graphs, (c - e) are vertical correlation graphs and (f-i) are diagonal correlation graphs.

Scheme 2(4D Arnold Cat Map) performs better than Scheme 1(3D Arnold Cat Map) scrambling in comparison. The proposed algorithm is time-efficient with fewer time-consuming EC point multiplications. A lightweight digital signature algorithm using the secure SHA-256 hash

TABLE 12. Comparison of correlation coefficient values for RGB images.

Scheme	Image	Type	Horizontal		Vertical		Diagonal	
			Plain	Cipher	Plain	Cipher	Plain	Cipher
Scheme 1	Mandrill	Red	0.92306	0.00046	0.86596	0.006	0.85434	-0.0002
		Green	0.86548	0.00037	0.76501	0.0079	0.73479	-0.002
		Blue	0.90734	-0.00026	0.88089	0.0048	0.83986	-0.0001
	Lena	Red	0.97977	-0.00086	0.98932	0.0062	0.969694	0.0036
		Green	0.96907	-0.00065	0.98249	0.0053	0.955546	-0.0024
		Blue	0.93274	-0.00154	0.957604	0.0029	0.918286	-0.0009
	House	Red	0.95355	0.0016	0.95788	0.0077	0.92243	0.0035
		Green	0.93909	0.00039	0.942326	0.0064	0.890085	-0.0016
		Blue	0.97248	0.0041	0.968604	0.0058	0.944495	0.0032
Scheme 2	Mandrill	Red	0.92306	0.00023	0.86596	0.0027	0.85434	0.00008
		Green	0.86548	0.0034	0.76501	0.00037	0.73479	0.00027
		Blue	0.90734	0.0008	0.88089	0.0018	0.83986	0.0012
	Lena	Red	0.97977	-0.0007	0.98932	0.0035	0.969694	-0.00009
		Green	0.96907	-0.0015	0.98249	0.0047	0.955546	0.0017
		Blue	0.93274	-0.0009	0.957604	0.0068	0.918286	0.00103
	House	Red	0.95355	-0.002	0.95788	0.0041	0.92243	-0.0017
		Green	0.93909	-0.0016	0.942326	0.0008	0.890085	0.0008
		Blue	0.97248	0.0012	0.968604	0.0019	0.944495	0.0031
Ref. [10]	House	Red	0.9467	-0.0067	0.9310	0.0004	0.8987	0.0147
		Green	0.9203	0.0177	0.9120	0.0175	0.8502	-0.0025
		Blue	0.9686	-0.0153	0.9526	-0.0001	0.9271	-0.0207
Ref. [42]	Mandrill	Red	0.9280	0.0186	0.8650	-0.006	0.8538	-0.0013
		Green	0.8625	0.0066	0.7697	0.0164	0.7256	0.0092
		Blue	0.9087	0.0067	0.8859	0.0012	0.8427	0.0172
	Lena	Red	0.9326	0.0035	0.9624	-0.004	0.907	-0.0410
		Green	0.9222	-0.0097	0.9546	0.0053	0.8804	-0.0085
		Blue	0.8938	0.0185	0.9343	0.0106	0.8634	-0.017

TABLE 13. Computational cost comparison analysis.

Scheme	Number of Point Multiplications required				
	Encryption	Decryption	Signature	Verification	Total
Proposed	1	0	0	1	2
Ref. [10]	2	1	1	2	6
Ref. [23]	2	1	×	×	3
Ref. [25]	1	1	0	1	3

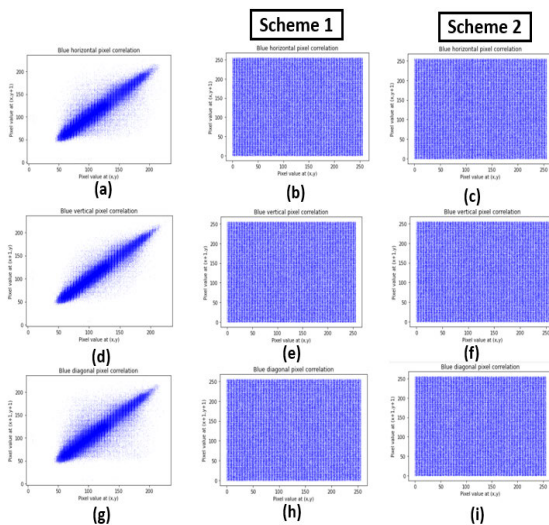


FIGURE 12. Correlation graphs of Blue component of plain and cipher color LENA images from their respective proposed Scheme 1 and Scheme 2 implementations. Here, (a - c) are horizontal correlation graphs, (c - e) are vertical correlation graphs and (f-i) are diagonal correlation graphs.

function is used to verify the validity of cipher image before decryption for faster decryption of correct cipher images.

VII. CONCLUSION

The proposed image encryption and authentication model intends to enhance the quality of encryption of both grayscale and color images through the discretized chaotic 3D and 4D Arnold Cat maps. The model strengthens the quality of cipher images compared to existing schemes with higher entropy, lower correlation, higher average NPCR and UACI, lower PSNR and SSIM values along with the ability to actively thwart the Chosen-Plaintext(CPA) and Known-Plaintext attacks(KPA). The proposed model is proven to be robust, lightweight and competent against the statistical and cryptanalytic attacks. In future work, the encryption model can be amended to support other forms of real-time multimedia encryption such as audio, video and more.

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