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Index Theory of One Dimensional Quantum Walks and Cellular Automata

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Abstract: If a one-dimensional quantum lattice system is subject to one step of a reversible discrete-time dynamics, it is intuitive that as much “quantum information” as moves into any given block of cells from the left, has to exit that block to the right. For two types of such systems — namely quantum walks and cellular automata — we make this intuition precise by defining an *index*, a quantity that measures the “net flow of quantum information” through the system. The index supplies a complete characterization of two properties of the discrete dynamics. First, two systems S_1, S_2 can be “pieced together”, in the sense that there is a system S which acts like S_1 in one region and like S_2 in some other region, if and only if S_1 and S_2 have the same index. Second, the index labels connected components of such systems: equality of the index is necessary and sufficient for the existence of a continuous deformation of S_1 into S_2 . In the case of quantum walks, the index is integer-valued, whereas for cellular automata, it takes values in the group of positive rationals. In both cases, the map $S \mapsto \text{ind } S$ is a group homomorphism if composition of the discrete dynamics is taken as the group law of the quantum systems. Systems with trivial index are precisely those which can be realized by partitioned unitaries, and the prototypes of systems with non-trivial index are shifts.

1. Introduction

Quantum walks and quantum cellular automata are quantum lattice systems with a discrete step dynamics, which is reversible, and satisfies a causality constraint: In each step only finitely many neighboring cells contribute to the state change of a given cell. This leads to an interesting interplay between the conditions of reversibility (unitarity) and causality, which is the subject of this article.

Starting point of the analysis is a simple intuition: for any connected group of cells in a one dimensional system as much “quantum information” as moves into the subsystem from the left has to move out at the other end. Moreover, this “flow” is a conserved

quantity, in the sense that it remains constant over the spatial extent of the system. It can thus be determined locally at any point.

Making this intuition precise, we associate with every such lattice system an *index*, a quantity measuring the net flow of information. The index theory developed in this work completely resolves three, a priori very different, classification problems:

(1) *Find all locally computable invariants.* It is shown that there exists a “crossover” between two systems S_1, S_2 if and only if their indices coincide. More precisely, a crossover S between S_1 and S_2 is a system which acts like S_1 on a negative half line $\{x|x \leq a\}$ and like S_2 on a positive half line $\{x|x \geq b\}$. Clearly, a locally computable invariant must assign the same value to two systems if there exists a crossover between them. It follows that any invariant is a function of the index.

(2) *Classify dynamics up to composition with local unitaries.* A natural way of constructing dynamics which respect both reversibility and causality is by concatenating layers of block unitaries. In every step, one would decompose the lattice into non-overlapping finite blocks and implement a unitary operation within every block. Such *local unitary implementations* are conceptually related to the gate model of quantum information. Not every time evolution may be realized this way: a uniform right-shift of cells serves as the paradigmatic counter-example. We show that the systems with local implementations are precisely those with trivial index. Consequently, equivalence classes of dynamics modulo composition with block unitaries are labeled by their indices.

(3) *Determine the homotopy classes.* It is proved that two systems may be continuously deformed into each other (with a uniform bound on the causality properties along the connecting path) if and only if they have the same index.

We will consider the above questions, and define indices, for two kinds of systems. *Quantum walks* are, on the one hand, the quantum analogs of classical random walks. On the other hand, they are discrete time analogs of a standard quantum particle “hopping” on a lattice according to a Hamiltonian which is a lattice version of the momentum operator $i\partial_x$. The index defined for these systems is the same as a quantity called “flow” by Kitaev [12]. Intuitively, this measures the mean speed of a quantum walk, expressed in units of “state space dimensions shifted to the right per time step”. The mathematical background has been explored, in a more abstract setting, by Avron, Seiler and Simon [3]. Kitaev’s work treated the first classification problem above. We will re-prove his results with an eye on generalizations to cellular automata, and will supply solutions to questions (2) and (3). Although the quantum walks seemed to be comparatively straightforward initially, the intuition gained from this case served us well in setting up the theory for the much more involved case of cellular automata. This allowed us to build an abstract index theory covering both cases with almost identical arguments (Sects. 4 and 5).

Cellular automata are characterized by the property that whatever state is possible in one cell (e.g., a superposition of empty/occupied) can be chosen independently for each cell. Expressed in terms of particles this means that we are necessarily looking at a “gas” system of possibly infinitely many particles. The basic definition of quantum cellular automata was given in [16]. On the one hand, the setting considered here is more restrictive than [16], covering only one-dimensional systems. On the other hand, we are allowing for non-translationally invariant dynamics — a strong generalization over the earlier paper. In fact, having completed the present work we feel that the translation invariance assumed in [16] was obscuring the fundamental interplay between reversibility and causality. Accordingly, we obtain here a stronger structure result, even though it is built on the same key ideas. Throughout, there is a strong interplay between local and

global properties. For example, the following statement is an immediate consequence of our main Theorem 9: If a nearest neighbor cellular automaton has somewhere a cell of dimension n , and somewhere else a cell of dimension m , coprime to n , then it can be globally implemented as a product of two partitioned unitary operations.

Our paper is organized as follows: After giving two examples in Sect. 2, we begin by a mathematical description of what we mean by quantum walks (Sect. 3.1) and quantum cellular automata (Sect. 3.2). We then describe the notion of locally computable invariants, and why they should form an abelian group (Sect. 5). A similar general explanation of the notion of local implementation is given in Sect.4. The detailed theory for quantum walks is in Sect. 6, and in Sect. 7 for cellular automata. This includes the proof that an index previously defined in the classical translationally invariant case [10], coincides with our index for this special case (Sect. 7.4). We close with an outlook on variants of index theory for either higher dimensional systems or automata with only approximate causality properties (Sect. 8).

2. Examples

Before introducing the mathematical setting, we will illustrate the problems treated in this paper by giving two concrete examples.

2.1. Particle hopping on a ring. The simplest example is given by a single particle on a ring of N sites arranged in a circle. More precisely, the Hilbert space we are considering is \mathcal{C}^N with basis vectors $\{|e_0\rangle, \dots, |e_{N-1}\rangle\}$. The vector $|e_i\rangle$ is taken to represent a “particle localized at position i ”. One step of a reversible discrete-time dynamics is simply given by an arbitrary unitary $U \in U(\mathcal{C}^N)$. We will consider two such time evolutions U_0, U_1 defined with respect to the standard basis by

$$U_0 : |e_i\rangle \mapsto |e_i\rangle, \quad U_1 : |e_i\rangle \mapsto |e_{(i+1) \bmod N}\rangle \tag{1}$$

respectively. The first unitary is the trivial evolution and the second one models a uniform movement of the particle with velocity one site per time step to the right.

The physical interpretation of this simple model seems clear: we can think of a lattice version of a particle with one spatial degree of freedom, where we have introduced cyclic boundary conditions to get a simple, finite description. The causality property defining a quantum walk then expresses the physically reasonable assumptions that couplings are local and dynamics preserve locality.

In this setting, it is natural to think of the time evolution as being generated by a Hamiltonian: $U_t = e^{itH}$. Such a Hamiltonian formula would allow us to extend the dynamics to arbitrary real times $t \in \mathcal{R}$. To recover H , we need to take a logarithm of U_1 . This operation is of course not uniquely defined, but the ansatz

$$H = \sum_{k=0}^{N-1} \frac{k}{2\pi} |f_k\rangle\langle f_k| \tag{2}$$

in terms of the Fourier basis

$$|f_k\rangle = \frac{1}{\sqrt{N}} \sum_{j=0}^{N-1} e^{\frac{2\pi i}{N}kj} \tag{3}$$

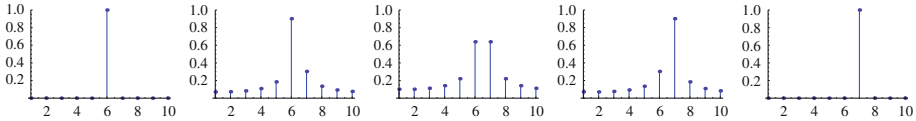


Fig. 1. Absolute value squared of a vector initially localized at site 6 (of 10) under the action of the time evolution e^{itH} for $t \in \{0, .25, .5, .75, 1\}$. The Hamiltonian is given by Eq. (2). The dynamics is not causal for non-integer values of t

seems particularly appealing. It can easily be checked to be compatible with our previous definitions of U_0 and U_1 .

Does this definition yield physically satisfactory dynamics for all times $t \in [0, 1]$? Hardly. As demonstrated in Fig. 1, a particle initially localized at site 6 ($N = 10$) will spread out over the entire ring during the interval $t \in [0, \frac{1}{2}]$, and refocus to site 7 during $t \in [\frac{1}{2}, 1]$. Any sensible notion of “causality” is violated for non-integer values of t . Note that this contrasts with the time evolution generated by the momentum operator $H = i\partial_x$ of a continuous-variable system. The latter fulfills $(e^{itH}\psi)(x) = \psi(x - t)$, thus preserving the localization properties of vectors $\psi \in L^2(\mathcal{R})$ for any $t \in \mathcal{R}$.

The discussion immediately raises several questions:

Is there a way to continuously interpolate between U_0 and U_1 while preserving causality? Our particular choice (2) for H could have been unfortunate. Conceivably, there is a better-suited, potentially time-dependent Hamiltonian which does the job. More physically phrased: is it possible to discretize the spacial degrees of freedom, but not the temporal ones, of a free single particle, while respecting causality? (The answer is: no, this is not possible).

Is there a simple way of deciding that the answer to the first question is negative? (Yes: the index of U_0 is 0, the index of U_1 is 1. The index is constant on connected components with uniformly bounded interaction length).

Further questions we will answer include: Is there a “gate-model” implementation of U_1 ? More precisely, can I write U_1 in terms of a constant-depth sequence of unitary operations, each of which acts non-trivially only on a constant number of sites? (No). Can I engineer a large system on $N' \gg N$ sites, endowed with a global time evolution U , such that the restriction of U to some region of contiguous sites looks like U_0 and the restriction of U to some other region looks like U_1 ? (No).

All these statements are made precise in Theorem 3.

2.2. Cluster state preparation. In this section, we consider N spin-1/2 particles arranged on a circle. The Hilbert space of the system is $\mathcal{H} = \bigotimes_{i=1}^N \mathbb{C}^2$. In the previous example, it was clear what it meant for a state vector ψ to be “localized” in some region $\Lambda \subset [1, N]$: namely this was the case if $\langle \psi | e_i \rangle = 0$ for all $i \notin \Lambda$. For state vectors on tensor product spaces, on the other hand, there seems to be no satisfactory notion of “locality”. To circumvent this problem, we focus on observables instead. An observable $A \in \mathcal{B}(\mathcal{H})$ is localized in a region Λ if A acts like the identity on all tensor factors outside of Λ .

Reversible dynamics on such a system is again represented by a unitary $U \in U(\mathcal{H})$. We use the symbol α to denote the action by conjugation of U on observables: $A \mapsto \alpha(A) = UAU^*$.

Let $\{\sigma_x^{(i)}, \sigma_y^{(i)}, \sigma_z^{(i)}\}$ be the Pauli matrices acting on the i^{th} spin. Since one can form a basis of $\mathcal{B}(\mathcal{H})$ from products of the Pauli matrices $\{\sigma_x^{(i)}, \sigma_z^{(i)}\}_{i=1}^N$ acting on single spins alone, it suffices to specify the effect of α on these $2N$ matrices in order to completely determine the dynamics. For example, we can set

$$\alpha(\sigma_x^{(i)}) = \sigma_z^{(i-1)} \otimes \sigma_x^{(i)} \otimes \sigma_z^{(i+1)}, \tag{4}$$

$$\alpha(\sigma_z^{(i)}) = \sigma_z^{(i)}. \tag{5}$$

It is a simple exercise to verify that the operators on the right-hand side fulfill the same commutation relations as the $\{\sigma_x^{(i)}, \sigma_z^{(i)}\}$. This is sufficient to ensure that a unitary U implementing the time evolution α actually exists.

As in the previous example, we can ask more refined questions about U . For example: can we implement U by a fixed-depth circuit of nearest-neighbor unitaries? Can one interpolate between U and the trivial time evolution while keeping observables localized along the path? Is there a simple numerical invariant which would allow us to easily decide these questions?

In this particular case an educated guess gives rise to affirmative answers to all these questions. Indeed, set

$$V_t^{(i,i+1)} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & e^{it\pi} \end{pmatrix} \tag{6}$$

with respect to the standard basis of the i^{th} and $(i + 1)^{\text{th}}$ spin. Define

$$U_t = \prod_{i=1}^N V_t^{(i,i+1)} \tag{7}$$

(the product is well-defined because V_t commutes with its translates, so the order in which the product is taken is immaterial). Then one checks that α corresponds to U_1 , whereas U_0 is the trivial evolution. Clearly U_t continuously interpolates between these two cases, and, by construction, does not increase the localization region of an observable by more than two sites. We note that U_1 is known in the quantum information literature as the interaction used to generate *graph states* [9, 15].

In Sect. 7, we will set up a general theory for answering the questions posed above—including in cases where one is not so lucky to have an explicit parametrization in terms of gates at hand.

3. Systems

3.1. Quantum walks. We consider a quantum system with a spatial degree of freedom $x \in \mathbb{Z}$, and at every site or “cell” x a finite dimensional “one-cell Hilbert space” \mathcal{H}_x . These spaces can be thought of as describing the internal states of the system as opposed to the external, spatial variables x . The Hilbert space of the system is

$$\mathcal{H} = \bigoplus_{x=-\infty}^{\infty} \mathcal{H}_x, \tag{8}$$

The specification of the \mathcal{H}_x will be called a *cell structure*.

We call a unitary operator U on \mathcal{H} *causal*,¹ or a *quantum walk*, if, for any $x \in \mathbb{Z}$, there are some $x_- < x_+$ such that $\phi \in \mathcal{H}_x \subset \mathcal{H}$ implies $U\phi \in \bigoplus_{y=x_-}^{x_+} \mathcal{H}_y$. We assume that x_{\pm} are both non-decreasing as functions of x , and go to $\pm\infty$ when x does.

By $U_{yx} : \mathcal{H}_x \rightarrow \mathcal{H}_y$ we denote the block matrix corresponding to the direct sum (8), i.e.

$$U \bigoplus_x \phi_x = \bigoplus_y \sum_x U_{yx} \phi_x. \tag{9}$$

The causality of U implies that, for any x , only finitely many y give non-zero summands.

Grouping. The spatial variable x of a walk and the internal degrees of freedom, described in \mathcal{H}_x , are largely interchangeable. In one direction we can choose a basis $|x, 1\rangle, \dots, |x, d\rangle$ in some \mathcal{H}_x , and replace the point x by the sequence of points $(x, 1), \dots, (x, d)$, each with a one-dimensional space $\mathcal{H}_{x,i}$ of internal states. Because in the above definition, we assumed that the spatial variable ranges over \mathbb{Z} , groupings have to be followed up by a relabeling of sites in the obvious way. In the other direction, we can “fuse together” several cells x_1, \dots, x_k , getting a new cell X with internal state space $\mathcal{H}_X = \bigoplus_{i=1}^k \mathcal{H}_i$. In either case it is clear how to adjust the neighborhood parameters x_{\pm} .

Hence we can either regard our system as one without internal degrees of freedom, and Hilbert space $\mathcal{H} = \ell^2(\mathbb{Z})$. Typically this may involve some large neighborhoods $[x_-, x_+]$. Or else, we can group cells until we get a nearest neighborhood system, i.e., $x_{\pm} = x \pm 1$, at the expense of having to deal with high-dimensional \mathcal{H}_x .

Most of these definitions and constructions are easily generalized to higher-dimensional lattices. It is therefore instructive to identify the feature which restricts our results to the one-dimensional case. Indeed, it lies in the fact that one can choose a partitioning into intervals $[a_i, b_i] \subset \mathbb{Z}$ such that the sites below a_i and above b_i interact only through the interval $[a_i, b_i]$. Formally, that is a consequence of demanding $\lim_{x \rightarrow \pm\infty} x_{\pm} = \pm\infty$. It is easy to see that such a separation need not be possible in a two-dimensional lattice, even if every cell has a finite neighborhood. In this more general setup, neighborhood relations up to regroupings may be described in terms of *coarse geometry* [14], a theme we will not pursue here.

Translationally invariant walks. The simplest way to define a cell structure is to choose a Hilbert space \mathcal{H}_0 , and to set $\mathcal{H}_x \equiv \mathcal{H}_0$ for all $x \in \mathbb{Z}$. We then have the unitary equivalence $\mathcal{H} \cong \ell^2(\mathbb{Z}) \otimes \mathcal{H}_0$. In that case we can define the *shift* operation S and its powers by

$$S^n(|x\rangle \otimes \phi_0) = |x+n\rangle \otimes \phi_0. \tag{10}$$

In this setting one frequently looks at translationally invariant walks, i.e., unitaries U commuting with S . More generally, there might be some period p such that $[U, S^p] = 0$. Clearly, it is natural in this case to group p consecutive cells, so that after grouping one gets a strictly translationally invariant walk.

¹ We are indebted to a referee who drew our attention to the inflationary use of the word “local” in our manuscript, where, among other things, quantum walks were called “local” unitaries. We changed this to “causal” for the crucial finite propagation property of walks and cellular automata. This is reminiscent of relativistic propagation in algebraic quantum field theory and in keeping with usage in signal processing (where x would be time, and $x_- = x$). In quantum information it agrees with [2,4]. The terminology disagrees with what some field theorists would probably say [18], and with [16].

The space \mathcal{H}_0 can then either be considered as an internal degree of freedom of a walking particle, or as a *coin* so that shift steps (possibly depending on the internal state) are alternated with unitary coin tosses $\mathbb{1} \otimes C$. Translationally invariant systems will be treated in more detail in Sect. 6.4.

Periodic boundary conditions. Since we are after a local theory of quantum walks, global aspects — like the distinction between a walk on \mathbb{Z} and a walk on a *large ring* of M sites — are secondary, as long as the interaction length $L = \max |x_+ - x_-|$ remains small in comparison with M . In fact, from any walk on a ring we can construct one on \mathbb{Z} which locally looks the same. More formally, let the sites of the ring be labeled by the classes \mathbb{Z}_M of integers modulo M , and identified with $\{0, \dots, M - 1\}$. Then we extend the cell structure by setting $\mathcal{H}_{x+kM} = \mathcal{H}_x$ for all $k \in \mathbb{Z}$. In order to extend the unitary U on the ring to a walk \widehat{U} on \mathbb{Z} we set

$$\widehat{U}_{xy} = \begin{cases} 0 & \text{if } |x - y| > L \\ U_{x'y'} & \text{if } |x - y| \leq L, x' \equiv x, y' \equiv y \pmod{M}. \end{cases} \tag{11}$$

For the second line to be unambiguous, we require $2L < M$. To verify unitarity we need to compute

$$\sum_y \widehat{U}_{xy}^* \widehat{U}_{yz} = \sum_{y'} U_{x'y'}^* U_{y'z'}, \tag{12}$$

where we have assumed that $|x - z| \leq 2L$, because otherwise the left hand side is zero anyhow. Note that for each summation index y only one class $y' \in \mathbb{Z}_M$ can occur in the sum on the right-hand side. Moreover, every class y' appears, although possibly with a zero contribution. But the sum on the right is $\delta_{x'z'} \mathbb{1}_{x'}$, which together with a similar argument for UU^* proves the unitarity of \widehat{U} .

From the point of view of index theory, the walk U on the ring and \widehat{U} on the line are the same. However, if we iterate U , the interaction length $(x_+ - x_-)$ increases, and eventually non-zero matrix elements can occur anywhere in U^n . In this sense, the set of quantum walks on a ring does not form a group. This is the reason why the theory of walks on \mathbb{Z} is more elegant and more complete. From now on we will therefore consider walks on \mathbb{Z} only.

3.2. Cellular automata. Once again we consider a system in which a finite dimensional Hilbert space \mathcal{H}_x is associated with every site $x \in \mathbb{Z}$. However, rather than combining these in a direct sum, we take their tensor product. In plain English this means that, for any two sites x, y , rather than having a system of type \mathcal{H}_x at position x *or* a system of type \mathcal{H}_y at position y , as in a quantum walk, we now have a system of type \mathcal{H}_x at position x *and* a system of type \mathcal{H}_y at position y . In contrast to the infinite direct sum of Hilbert spaces, the infinite tensor product is not well-defined. Since we want to look at local properties, we could work with a “potentially” infinite product, i.e., some finite product with more factors added as needed in the course of an argument. But it is easier to work instead with the observable algebras \mathcal{A}_x , equal to the operators on the Hilbert space \mathcal{H}_x , or equivalently the algebra $\mathcal{M}_{d(x)}$, where $d(x) = \dim \mathcal{H}_x$ (\mathcal{M}_d denotes the algebra of $d \times d$ -matrices). In analogy to the definition for walks, we will refer to the specification of the algebras \mathcal{A}_x as the *cell structure*.

For the observable algebras associated to sets $\Lambda \subset \mathbb{Z}$, we use the following notations: for finite Λ , $\mathcal{A}(\Lambda)$ is the tensor product of all \mathcal{A}_x with $x \in \Lambda$. For $\Lambda_1 \subset \Lambda_2$ we identify $\mathcal{A}(\Lambda_1)$ with the subalgebra $\mathcal{A}(\Lambda_1) \otimes \mathbb{1}^{\Lambda_2 \setminus \Lambda_1} \subset \mathcal{A}(\Lambda_2)$. For infinite $\Lambda \subset \mathbb{Z}$ we denote by $\mathcal{A}(\Lambda)$ the C^* -closure of the increasing family of finite dimensional algebras $\mathcal{A}(\Lambda_f)$ for finite $\Lambda_f \subset \Lambda$, also called the quasi-local algebra [5]. In particular, the algebra of the whole chain is $\mathcal{A}(\mathbb{Z})$, sometimes abbreviated to \mathcal{A} .

A *cellular automaton* with cell structure $\{\mathcal{A}_x\}_{x \in \mathbb{Z}}$ is an automorphism α of $\mathcal{A} = \mathcal{A}(\mathbb{Z})$ such that, for some functions $x \mapsto x_{\pm}$ as specified in Sect. 3.1, each $\alpha(\mathcal{A}_x) \subset \mathcal{A}([x_-, x_+])$. Note that the restricted homomorphisms $\alpha_x : \mathcal{A}_x \rightarrow \mathcal{A}([x_-, x_+])$ uniquely determine α , because every observable acting on a finite number of cells can be expanded into products of one-site observables. These *local rules* α_x have to satisfy the constraint that the algebras $\alpha_x(\mathcal{A}_x)$ for different x commute element-wise. In that case they uniquely determine an endomorphism α . For examples and various construction methods for cellular automata we refer to [16].

Exactly as in the case of quantum walks we can group cells together for convenience. Whereas the dimensions for subcells add up for quantum walks ($\dim \bigoplus_{x \in \Lambda} \mathcal{H}_x = \sum_{x \in \Lambda} \dim \mathcal{H}_x$) we get $\mathcal{A}(\Lambda) \cong \mathcal{M}_d$ with the product $d = \prod_{x \in \Lambda} d(x)$.

By considering the time evolution of observables, we have implicitly chosen to work in the Heisenberg picture. The expectation value of the physical procedure (i) prepare a state ρ , (ii) run the automaton for k time steps, (iii) measure an observable A would thus be given by the expression $\rho(\alpha^k(A))$. Accordingly, we choose a convention for the *shift* on a chain with isomorphic cells, which at first seems inverted relative to the definition (10) for walks. We define it as the automorphism σ with $\sigma(\mathcal{A}_x) = \mathcal{A}_{x-1}$, acting according to the assumed isomorphism of all the cell algebras. Thus if one prepares a certain state, it will be found shifted to the right after one step of σ , in accordance with (10) although in that case $U\mathcal{H}_x = \mathcal{H}_{x+1}$.

4. Local Implementability

We have defined the causality properties of walks and cellular automata *axiomatically*, i.e., as a condition on the input-output behavior of the maps U and α . Alternatively, one may take a *constructive* approach. Here, one would list a set of operations that should certainly be included in the set of local dynamics, and refer to any given time evolution as being locally implementable if it can be represented as a sequence of these basic building blocks. Both methods are equally valid, and in this section we will completely analyze their relation. From the axiomatic point of view this might be called a “structure theorem”, whereas from the constructive point of view one would call it a “characterization theorem”.

In the case at hand, there is a natural choice of building blocks. Namely, we can partition the system into some subsets (“blocks”) of sites, and apply a unitary operation separately to each subsystem in the partition. (Note that the unitaries would be combined by a direct sum for walks and by a tensor product for cellular automata). For such maps the interplay between unitarity and causality is trivial: causality puts no constraint whatsoever on the choice of unitaries acting in each block. Moreover, it allows the overall operation to be resolved into a sequence of steps, in which one block operation is done after the other. This picture is close to the gate model of quantum computation [13]: here each block unitary would correspond to one “gate” involving some subset of registers, so that these gates do not disturb each other. The fact that they can be executed in parallel is expressed by saying that these infinitely many gates nevertheless represent an operation of *logical depth 1*.

For partitioned unitary operations the various block unitaries obviously commute. Commutation is really the essential feature if we want to resolve the overall time step into a sequence of block unitary steps. Indeed, consider a family of commuting unitaries U_j , each localized in a finite subset Λ_j of some infinite lattice (not necessarily one-dimensional). We only need that the cover by the Λ_j is locally finite, i.e., each point x is contained in at most finitely many Λ_j . Then the product $\prod_j U_j$ implements a well-defined operation on localized elements. In the cellular automaton case (where localization just means $U_j \in \mathcal{A}(\Lambda_j)$), we define the action on a local observable A as

$$\alpha(A) = \left(\prod_j U_j^* \right) A \left(\prod_j U_j \right), \tag{13}$$

with the understanding that both products range over the same index set, namely those j for which Λ_j meets the localization region of A . Here the products can be taken without regard to operator ordering, since we assumed that the U_j commute. Including additional factors U_j on both sides will not change $\alpha(A)$, since such factors can be “commuted past” all other $U_{j'}$, $U_{j'}^*$ and A to meet the corresponding U_j^* and cancel. So the product is over *all* j , in the sense of a product over any sufficiently large finite set. Similar considerations apply for the case of walks on general lattices.

Now if a QCA α is represented in the form (13), we can also represent it as a product of partitioned operations: indeed, we only need to group the U_j into families within which all Λ_j are disjoint. The product of each family is obviously a partitioned unitary and under suitable uniformity conditions on the cover we only need a finite product of such partitioned operations to represent α , typically $s + 1$ factors, where s is the spatial dimension of the lattice. Hence we consider the representation as a product of partitioned unitaries as essentially equivalent to the representation by commuting unitaries as in (13). In either case we will say that the system is *locally implementable*.

We now come to the basic result for implementing general walks or cellular automata by commuting unitaries — provided we are allowed to enlarge the system. The key feature of these extensions is that they work in arbitrary (not necessarily one-dimensional) lattices and that the ancillary system is a copy of the system itself, on which we implement simultaneously the inverse operation. In the following result, we allow the underlying “lattice” X to be any countable set. General neighborhood schemes are described as controlled sets in some coarse structure [14]. For the present paper it suffices to describe causality in terms of a metric d on X , of which we only assume that all balls $\mathcal{N}_L(x) = \{y \mid d(x, y) \leq L\}$ are finite sets. The causality condition for walks on X is then that there is some “interaction radius” L such that in (9) $U_{yx} = 0$ for $d(x, y) > L$. Similarly, for QCAs, the causality condition reads $\alpha(\mathcal{A}_x) \subset \mathcal{A}(\mathcal{N}_L(x))$. For walks U, V on the same lattice we simply write $U \oplus V$ for a walk with one-cell Hilbert spaces $\mathcal{H}_x \oplus \mathcal{K}_x$, where \mathcal{H}_x are the one-cell spaces for U and \mathcal{K}_x those for V . This splits the total Hilbert space into $\mathcal{H} \oplus \mathcal{K}$, and $U \oplus V$ acts according to this direct sum. Similarly, we define the tensor product $\alpha \otimes \beta$ acting on two parallel systems combined in a tensor product.

Proposition 1. (1) *For any quantum walk U , the walk $U \oplus U^*$ is locally implementable.*
 (2) *For any cellular automaton α , the automaton $\alpha \otimes \alpha^{-1}$ is locally implementable.*

Proof. (1) We are considering a doubled system in which the one-cell Hilbert space at x is $\mathcal{H}_x \oplus \mathcal{H}_x$. Let S_x denote the unitary operator on the doubled system which swaps

these two summands, and acts as the identity on the one-cell spaces of all other sites. Now consider the unitaries

$$T_x = (U^* \oplus \mathbb{1})S_x(U \oplus \mathbb{1}).$$

These commute, because they are the images of the commuting transformations S_x under the same unitary conjugation. Moreover, they are localized near x by the causality properties postulated for U . Hence their infinite product defines a walk unitary, as discussed above. This unitary is

$$\prod_x T_x = (U^* \oplus \mathbb{1})S(U \oplus \mathbb{1}) = S(\mathbb{1} \oplus U^*)(U \oplus \mathbb{1}) = S(U \oplus U^*),$$

where we have used that $S = \prod_x S_x$ is just the global swap of the two system copies. Hence $U \oplus U^* = (\prod_x S_x)(\prod_x T_x)$ is locally implemented.

(2) Essentially the same idea works for cellular automata [2]. Again we consider the unitaries $S_x \in \mathcal{A}_x \otimes \mathcal{A}_x$, which swap the two tensor factors, so that $S_x(A_x \otimes B_x) = (B_x \otimes A_x)S_x$. Now consider the unitary elements

$$T_x = (\text{id} \otimes \alpha)[S_x].$$

Here we have written the arguments of an automorphism in brackets, to distinguish it from grouping parentheses, and thus eliminate a possible source of confusion in the coming computations. As images of a family of commuting unitaries under an automorphism, the T_x are themselves a commuting family of unitaries. Moreover, they are localized in $\mathcal{A}_x \otimes \mathcal{A}_{\mathcal{N}(x)}$. Hence they implement a cellular automaton β . We determine it by letting it act first on a localized element of the form $A_x \otimes \mathbb{1}$ with $A_x \in \mathcal{A}_x$,

$$\begin{aligned} \beta[A_x \otimes \mathbb{1}] &= \left(\prod_y T_y\right)^* (A_x \otimes \mathbb{1}) \left(\prod_y T_y\right) \\ &= (\text{id} \otimes \alpha) \left[\prod_y S_y\right] (\text{id} \otimes \alpha) [A_x \otimes \mathbb{1}] (\text{id} \otimes \alpha) \left[\prod_y S_y\right] \\ &= (\text{id} \otimes \alpha) \left[\left(\prod_y S_y\right) (A_x \otimes \mathbb{1}) \left(\prod_y S_y\right)\right] \\ &= (\text{id} \otimes \alpha)[\mathbb{1} \otimes A_x] = \mathbb{1} \otimes \alpha[A_x]. \end{aligned}$$

A similar computation shows that $\beta[\mathbb{1} \otimes \alpha[B_x]] = B_x \otimes \mathbb{1}$. Since α is an automorphism, this is the same as $\beta[\mathbb{1} \otimes B_x] = \alpha^{-1}[B_x] \otimes \mathbb{1}$. Using the homomorphism property of β , we get for general localized elements A, B that $\beta[A \otimes B] = \alpha^{-1}[B] \otimes \alpha[A]$. Hence following β by a global swap (implemented by $\prod_x S_x$) we have implemented $\alpha \otimes \alpha^{-1}$ locally. \square

5. The Group of Locally Computable Invariants

In this section we take up the idea of a *locally computable invariant* and show that, for either walks or automata, these invariants necessarily form an abelian group. The group multiplication reflects both the composition and the parallel application to a double chain. We postpone to later sections the question whether nontrivial invariants exist, i.e., at this stage it might well be that the group described here is trivial. Later on we

will determine this group to be $(\mathbb{Z}, +)$ for quantum walks (see Sect. 6) and (\mathbb{Q}_+, \cdot) , the multiplicative group of positive fractions, for cellular automata (see Sect. 7). In this section, in order not to double each step, we will describe the arguments for the case of walks, and just comment at the end on the necessary changes for the cellular automaton case.

Suppose we have defined a property $\mathbf{P}(U)$, which is defined for any quantum walk U , and which can be determined solely on the basis of a finite collection of the block matrices U_{xy} . More specifically, if we write the walk in nearest-neighbor form by grouping, we call a property $\mathbf{P}(U)$ *locally computable* if we can compute it from the restriction of U on any interval of length ≥ 2 . The crucial part of this definition is, of course, that the result obtained in this way must be the same for any interval we may select for the computation, a property which we stress by calling \mathbf{P} a locally computable invariant.

Suppose now that two walks U_1 and U_2 *share a patch*, in the sense that there is a long interval $[x_1, x_2] \subset \mathbb{Z}$, on which the Hilbert spaces \mathcal{H}_x for $x \in [x_1, x_2]$ have the same dimensions and, after the choice of a suitable isomorphism, the unitaries U_1 and U_2 restricted to these subspaces act in the same way. We assume that the interval is sufficiently long to determine \mathbf{P} . Then local computability just means that we must have $\mathbf{P}(U_1) = \mathbf{P}(U_2)$. In other words, \mathbf{P} must be constant on each equivalence class of the relation of “sharing a patch”. So the theory of locally computable invariants is really equivalent to characterizing the classes of the transitive hull of this relation: We will write $U_1 \sim U_n$, if there is a chain of walks U_1, U_2, \dots, U_n such that, for all i , U_i and U_{i+1} share a patch. In contrast to the relation of sharing a patch, this equivalence relation no longer makes any requirements about the sizes of any one-cell Hilbert spaces in the walks U_1 and U_n . The most comprehensive locally computable property is now just the property of U to belong to some equivalence class: all other locally computable properties are functions of this class. Our aim thus shifts to computing the set \mathcal{J} of equivalence classes for “ \sim ”. The equivalence class of a walk U will be denoted by $\text{ind}(U) \in \mathcal{J}$, and called its (abstract) *index*.

Triviality of locally implementable systems. Let us first make the connection to the questions of the previous section: suppose that a walk or automaton is locally implementable, i.e., the product of a collection of block partitioned unitaries. Compare this with a system in which all unitaries, whose localization intersects the positive half axis, are replaced by the identity. Clearly, this acts like the identity on all cells on the positive axis, and we can further modify the system by making it trivial (0-dimensional \mathcal{H}_x or 1-dimensional algebras \mathcal{A}_x) for $x > 0$. Clearly these systems share a large patch (most of the negative axis), so they are equivalent. In other words, locally implementable systems have the same index as the identity on a trivial chain.

Crossovers. A very useful fact about the relation $U_1 \sim U_2$ is that it implies a *prima facie* much stronger relation: It is equivalent to the property that there is a “crossover” walk U_c , which coincides with U_1 on a negative half line $\{x|x \leq a\}$ and coincides with U_2 on a positive half line $\{x|x \geq b\}$.

Proof. Let us denote the relation just described by $U_1 \approx U_2$. Then $U_1 \approx U_2 \Rightarrow U_1 \sim U_2$, because U_1 and U_2 each share an infinite patch with U_c .

In the converse direction, if U_1 and U_2 share a patch, we can define U_c to be the walk whose one-cell Hilbert spaces \mathcal{H}_x are those of U_1 for x to the left of the shared patch, and are those of U_2 for x to the right of the shared patch. Similarly, we define the unitary U_c to coincide with U_1 to the left and with U_2 on the right. On the shared patch we can

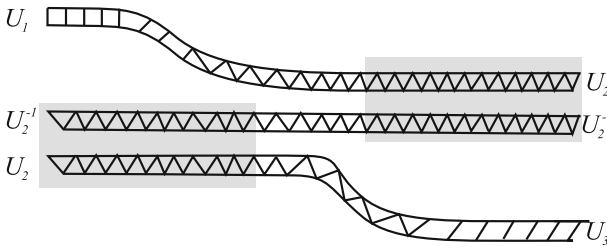


Fig. 2. Combining a crossover from U_1 to U_2 with a crossover from U_2 to U_3 to obtain a crossover from U_1 to U_3 . The shaded double chains can be fused to a single cell by Prop. 1

choose either one, since the two walks coincide. Since the shared patch was assumed to be sufficiently long this does not lead to an ambiguity for either U_c or U_c^{-1} . Hence $U_1 \approx U_2$.

In order to cover the case that U_1 and U_2 are linked by a chain in which any neighbors share a patch, we only need to prove that $U_1 \approx U_2$ is a transitive relation.

In order to prove transitivity, consider the walks U_1, U_2, U_3 with crossovers U_{12} and U_{23} as shown in Fig. 2. There we also included a copy of the inverse of U_2 . We take the overall picture as a representation of the walk $U_{12} \oplus U_2^{-1} \oplus U_{23}$. Now consider the strands of U_{12} and U_2^{-1} to the right of the crossover region of U_{12} . Since $U_2 \oplus U_2^{-1}$ is locally implementable by Prop. 1, we can replace this pair of strands by a trivial system, still retaining a legitimate unitary operation for the rest. Similarly, we can fuse the strands of U_{23} and U_2^{-1} to the left of the crossover region of U_{23} to nothing. This results in a unitary operator, which coincides with U_1 on a left half axis and with U_3 on a positive half axis, i.e., a crossover U_{13} . \square

Grouping. Suppose we regroup some finite collection of the cells. Clearly, this does not affect cells far away, and we immediately get a crossover. Hence the index does not change when regrouping cells, even if this is carried out in parallel. We will implicitly use this in the sequel by regrouping sites in whatever way is most convenient.

Parallel chains. Define the direct sum $U \oplus V$ of walks as in the previous section. Then if U_c (resp. V_c) is a crossover between U_1 and U_2 (resp. V_1 and V_2), $U_c \oplus V_c$ is obviously a crossover between $U_1 \oplus V_1$ and $U_2 \oplus V_2$. Hence the class $\text{ind}(U \oplus V) \in \mathcal{J}$ depends only on the equivalence classes of U and V , and we can define an “addition” of indices by $\text{ind}(U) + \text{ind}(V) = \text{ind}(U \oplus V)$. This addition is abelian, because there is a trivial crossover between $U \oplus V$ and $V \oplus U$, just exchanging the summands on a half chain. Moreover, since an inverse is defined as $-\text{ind}(U) = \text{ind}(U^*)$ via Prop. 1, we conclude that \mathcal{J} becomes an abelian group.

Products. Now suppose U and V are walks on the same cell structure, so that UV makes sense. We claim that there is a crossover between UV and $U \oplus V$. Hence we also get

$$\text{ind}(UV) = \text{ind}(U \oplus V) = \text{ind}(U) + \text{ind}(V).$$

Indeed, consider the cell structure on which $U \oplus V$ is defined, which has one-cell Hilbert spaces $\mathcal{H}_x \oplus \mathcal{H}_x$ at each site. Now let S_+ denote the unitary which acts as the swap on all $\mathcal{H}_x \otimes \mathcal{H}_x$ with $x > 0$ and leaves the subspaces with $x \leq 0$ unchanged. Then $(U \oplus \mathbb{1})S_+(\mathbb{1} \oplus V)S_+$ is a crossover between $U \oplus V$ and $UV \oplus \mathbb{1} \sim UV$.

Modifications for cellular automata. The concept of crossovers and the arguments for the group structure can be taken over verbatim, with the replacements $U \mapsto \alpha$, $\mathcal{H}_x \mapsto \mathcal{A}_x$, $\oplus \mapsto \otimes$. Of course, the index group will also be different, and we will adopt the convention to write it multiplicatively. The product formula thus reads $\text{ind}(\alpha\beta) = \text{ind}(\alpha \otimes \beta) = \text{ind}(\alpha)\text{ind}(\beta)$.

Numerical representation and shift subgroups. To make the abstract theory of this section useful, one needs to establish an isomorphism of the index group with some explicitly known group. The natural way to do that is to identify *generators*, i.e., some particular walks which cannot be implemented locally, but are sufficient to generate arbitrary walks up to locally implementable factors. Although it is far from obvious at this point, it will turn out later that, for walks as well as for cellular automata, the role of generators is played by the shifts. Since there is only one kind of generators, it suffices to assign numbers as “index values” to the shifts to establish an isomorphism of the abstract index groups with groups of numbers.

For walks, the shift was introduced in Sect. 3.1. We denote by S_d the shift on a system with d -dimensional internal degree of freedom ($\dim \mathcal{H}_x = d$ for all x). Similarly, let σ_d denote the shift automorphism on a chain with cell algebra \mathcal{M}_d . We tentatively demand

$$\text{ind}(S_d) = d \tag{14}$$

and, similarly,

$$\text{ind}(\sigma_d) = d. \tag{15}$$

This has to be consistent for shifts on parallel chains. Since $S_d \oplus S_e = S_{d+e}$ and $\sigma_d \otimes \sigma_e = \sigma_{d \cdot e}$, this requires that we take the indices of walks as a group of numbers under addition, and for the cellular automata as a group of numbers under multiplication. Indeed, we will show that the above formulas fix an isomorphism of the abstract index group \mathcal{J} to $(\mathbb{Z}, +)$ for quantum walks, and to the group (\mathbb{Q}_+, \cdot) of positive fractions for cellular automata.

6. Index for Quantum Walks

6.1. Pedestrian definition. The following is the basic definition of this chapter. To the best knowledge of the authors, it is due to Kitaev [12], who calls this quantity the *flow* of a walk U .

Definition 2. For any walk U , we define the **index** as

$$\text{ind } U = \sum_{x \geq 0 > y} (\text{Tr}(U_{xy})^* U_{xy} - \text{Tr}(U_{yx})^* U_{yx}). \tag{16}$$

Note that the sum is finite by virtue of the definition of causal unitaries. Clearly, for the simple shift we get $\text{ind } S_1 = 1$, confirming Eq. (14).

Of course, we will show presently that this quantity has all the properties of the abstract index discussed in the previous sections. However, from the definition given here it seems miraculous that such a quantity should be always an integer, and independent of the positioning of the cut. To see this it is better to rewrite this quantity in the following way.

6.2. Operator theoretic definition. We introduce the projection P for the half axis $\{x \geq 0\}$, i.e., the projection onto the subspace $\bigoplus_{x \geq 0} \mathcal{H}_x$. Then, for $\phi_x \in \mathcal{H}_x$ and $\phi_y \in \mathcal{H}_y$ we get

$$\langle \phi_x | PU - UP | \phi_y \rangle = \begin{cases} 0 & x \geq 0 \text{ and } y \geq 0 \\ -\langle \phi_x | U_{xy} | \phi_y \rangle & x < 0 \text{ and } y \geq 0 \\ \langle \phi_x | U_{xy} | \phi_y \rangle & x \geq 0 \text{ and } y < 0 \\ 0 & x < 0 \text{ and } y < 0 \end{cases}. \quad (17)$$

Hence, for any pair (x, y) the commutator $[U, P]$ has just the signs used in the definition of the index, and we get

$$\text{ind } U = \text{Tr } U^*[P, U] = \text{Tr}(U^*PU - P). \quad (18)$$

Note that for the the trace on the right hand side we cannot use linearity of the trace to write it as the difference of two (equal!) terms, because this would result in an indeterminate expression $\infty - \infty$.

6.3. Fundamental properties of the index for walks.

Theorem 3. 1. $\text{ind } U$ is an integer for any walk U .

2. $\text{ind } U$ is locally computable, and uniquely parameterizes the equivalence classes for the relation \sim from Sect. 5, hence can be identified with the abstract index defined there.
3. $\text{ind}(U_1 \oplus U_2) = \text{ind}(U_1) + \text{ind}(U_2)$, and, when U_1 and U_2 are defined on the same cell structure, $\text{ind}(U_1 U_2) = \text{ind}(U_1) + \text{ind}(U_2)$. Moreover, for the shift of d -dimensional cells: $\text{ind } S_d = d$.
4. $\text{ind } U = 0$ if and only if U admits a “local decoupling”, i.e. there is a unitary V , which acts like the identity on all but finitely many \mathcal{H}_x , such that UV is block diagonal with respect to the decomposition $\mathcal{H} = (\bigoplus_{x \leq 0} \mathcal{H}_x) \oplus (\bigoplus_{x \geq 1} \mathcal{H}_x)$.
5. $\text{ind } U = 0$ if and only if it is locally implementable (see Sect. 4). In this case it can be written as a product of just two partitioned unitaries. When U is regrouped in nearest neighbor form, then the partitioned unitaries can be chosen to couple only pairs of nearest neighbors.
6. $\text{ind } U_1 = \text{ind } U_0$ if and only if U_0 and U_1 lie in the same connected component, i.e., there is a norm continuous path $[0, 1] \ni t \mapsto U_t$ of causal unitaries of uniformly bounded interaction length L with the specified boundary values.

The rest of this subsection is devoted to the proof of this result. According to Eq. (18), the index is closely related to a difference of projections. If these were finite dimensional, we could just use linearity to get the difference of two integers. The following lemma shows that the result is still an integer when the difference of the two projections has finite rank. Actually it is even sufficient for the difference to be trace class, and with a careful discussion of the trace, it is sufficient for the ± 1 eigenspaces of the difference to be finite dimensional [3]. Here we include the simple case sufficient for our purposes.

Lemma 4. Let Q, P be orthogonal projections in a Hilbert space \mathcal{H} , such that $Q - P$ has finite rank. Then

1. The range $\mathcal{R} = (Q - P)\mathcal{H}$ is an invariant subspace for both Q and P .
2. $\text{Tr}(Q - P)$ is an integer.

3. *There is a unitary operator V such that $V\phi = \phi$ for all $\phi \perp (Q - P)\mathcal{H}$, and such that $Q \geq V^*PV$ or $Q \leq V^*PV$.*
4. *If $\text{Tr}(Q - P) = 0$, the V from the previous item satisfies $Q = V^*PV$.*

Proof. 1. Follows from the identity

$$Q(Q - P) = Q(\mathbb{1} - P) = (Q - P)(\mathbb{1} - P),$$

and its analogue for P .

2. Clearly, we can evaluate the trace in a basis of \mathcal{R} since the basis elements from \mathcal{R}^\perp contribute only zeros. Since the restrictions of Q and P to \mathcal{R} are projections on a finite dimensional space,

$$\text{Tr}(Q - P) = \text{Tr}_{\mathcal{R}}(Q - P) = \text{Tr}_{\mathcal{R}}(Q) - \text{Tr}_{\mathcal{R}}(P)$$

is the difference of two natural numbers.

3&4. Obviously, we can find such a unitary on \mathcal{R} with the corresponding property for the restrictions of Q and P to \mathcal{R} . We then extend V to be the identity on \mathcal{R}^\perp . When $\text{Tr}(Q - P) = 0$, this V is a unitary mapping from $Q\mathcal{R}$ to $P\mathcal{R}$. \square

Proof (of Theorem 3). None of the statements, or values of the index will change under grouping, except part of Item 5, which requires nearest neighbor form. Therefore we will assume without loss that all walks are nearest neighbor. We will use Lemma 4 with $Q = U^*PU$.

1. This follows directly from Eq. (18) and Lemma 4, Item 2.

2. Let P' be the projection onto another half axis, say $x \geq x_0$. Then $P - P'$ is finite rank and hence $(U^*PU - P) - (U^*P'U - P') = P' - P - U^*(P - P')U$ is the difference of two finite rank operators with equal trace. Hence the index does not depend on the cut position, and since formula (2) clearly involves only matrix elements at most 1 site from the cut, it is a locally computable invariant. It remains to be shown that it is a complete invariant, i.e., that $\text{ind } U_1 = \text{ind } U_2$ implies $U_1 \sim U_2$ in the sense of Sect. 5. This will be done in connection with Item 4 below.

3. This follows from Sect. 5. But a direct proof (for the product) is also instructive:

$$(U_1U_2)^*P(U_1U_2) - P = (U_2^*PU_2 - P) + U_2^*(U_1^*PU_1 - P)U_2$$

is the sum of two finite rank operators, of which we can take the trace separately.

4. Apply Lemma 4, Item 3, to get V with $P = V^*(U^*PU)V$, and hence $PUV = UVV^*P$. The fact that $V - \mathbb{1}$ vanishes on all but finitely many \mathcal{H}_x follows from its construction: $V - \mathbb{1}$ vanishes on the complement of $(P - Q)\mathcal{H} \subset \mathcal{H}_{-1} \oplus \mathcal{H}_0$, for the cut “ $-1|0$ ” used in Def. 2. Note that this implies $V \sim \mathbb{1}$, and also $UV \sim \mathbb{1}$, since a unitary which has no matrix elements connecting $x \geq 0$ and $x < 0$ clearly allows a crossover with the identity. From the product formula for locally computable invariants we therefore get that $\text{ind } U = 0$ implies $U \sim \mathbb{1}$. Obviously, this extends to other values of the index: if $\text{ind } U_1 = \text{ind } U_2$ we have $\text{ind } U_1^*U_2 = 0$, hence $U_1^*U_2 \sim \mathbb{1}$ and hence $U_1 \sim U_2$. This completes the proof of Item 2.

5&6. These items each contain a trivial direction: We have already shown in Sect. 5 that locally implementable walks have trivial index. Moreover, it is clear from Definition 2 that the index is a continuous function, and must hence be constant on each connected component. The non-trivial statement in Item 5. is that walks with trivial index are indeed implementable, and in Item 6. that walks with vanishing index can

be connected to the identity (the rest then follows by multiplication). In either case, an explicit construction is required, and it will actually be the same one.

So let $U = U_1$ be a walk and $U = 0$. Let V_0 denote the decoupling unitary for the cut $-1|0$, obtained in the proof of Item 4, and define similar unitaries V_k for the cuts at $2k - 1|2k$ such that UV_k has no non-zero matrix elements $(UV_k)_{xy}$ with $y < 2k \leq x$. Let H_k denote a hermitian operator located on the same subspaces as $V_k - \mathbb{1}$, such that $V_k = \exp(iH_k)$. We will take all H_k bounded in norm by the same constant (π will do). Then since they live on orthogonal subspaces, their sum $H = \sum_k H_k$ is well-defined and also bounded. Now let $V(t) = \exp(itH)$, which is a norm continuous function of t , because $\|H\| \leq \pi$. The endpoint $V(1)$ can also be defined by this product formula $V(1) = \prod_k V_k$, because on each subspace $\mathcal{H}_{2k-1} \oplus \mathcal{H}_{2k}$ only one of these unitaries is different from $\mathbb{1}$. Moreover, $UV(1)$ has no matrix elements $y < 2k \leq x$ for any k , i.e., it is block diagonal for a decomposition of \mathbb{Z} into pairs $\{2k, 2k + 1\}$. Now take a similar Hamiltonian path deforming each block in this matrix decomposition separately to the identity. Specifically, we take $W(0) = \mathbb{1}$ and $W(1) = UV(1)$. Then $t \mapsto W(t)V(t)^*$ is a norm continuous path (although no longer a unitary group), with the endpoints $\mathbb{1}$ and U . Moreover, each unitary $W(t)$ or $V(t)^*$ is based on a partition into neighboring pairs so that, for no t , $W(t)V(t)^*$ has any non-zero matrix element between sites with $|x - y| > 2$. This proves the remaining statement in Item 5 (for $t = 1$), and also the statement about uniformly bounded neighborhoods in Item 6. \square

6.4. The translation invariant case. Suppose that U commutes with some power of the shift. It is then useful to group spaces \mathcal{H}_x into larger blocks to get commutation with the shift itself. That is, in this section we assume all $\mathcal{H}_x \equiv \mathcal{K}$ to be equal, and $U_{xy} = U_{x-y}$, where by a slight abuse of notation the single-index quantity U_x is defined as U_{x0} . The width L is the largest x such that $U_x \neq 0$ or $U_{-x} \neq 0$. It is natural to diagonalize U using the Fourier transform. We define $\mathcal{F} : \ell^2(\mathbb{Z}) \otimes \mathcal{K} \rightarrow \mathcal{L}^2([-\pi, \pi]) \otimes \mathcal{K}$ by $\mathcal{F}(\Psi)(p) = \frac{1}{\sqrt{2\pi}} \sum_x e^{ipx} \Psi(x)$. This is to be read as a \mathcal{K} -valued equation, where we use the natural identification of $\mathcal{L}^2([-\pi, \pi]) \otimes \mathcal{K}$ with the set of \mathcal{K} -valued square integrable functions on $[-\pi, \pi]$. Similarly, we identify $\ell^2(\mathbb{Z}) \otimes \mathcal{K}$ with the \mathcal{K} -valued square summable sequences. Then $\mathcal{F}U\mathcal{F}^*$ becomes the multiplication operator by the p -dependent matrix

$$\widehat{U}(p) = \sum_{x=-L}^L U_x e^{ipx}. \tag{19}$$

Note that this is a Laurent polynomial in e^{ip} . The largest degree of e^{ip} in the polynomial is $x_+ - x_-$, which is constant by translation invariance. The lowest degree is $x - x_-$. Further, $\widehat{U}(p)$ must be a unitary operator on \mathcal{K} for every $p \in \mathbb{R}$. Taking these facts together, we conclude that both the determinant $\det \widehat{U}(p) = f(p)$ and its inverse $1/f(p) = \det \widehat{U}(p)^*$ are Laurent polynomials as well. But this is only possible if f is actually a monomial, say proportional to $\exp(inp)$ for some integer n . We claim that this n is the index:

Proposition 5. *For a translation invariant walk,*

$$\det \widehat{U}(p) = C e^{ip \operatorname{ind}(U)}, \tag{20}$$

for some phase constant C .

As a simple example consider the shift on a chain with $\dim \mathcal{H}_0 = 1$. We already noted after Definition 2 that this has index 1. The corresponding p -dependent unitary is the number $\widehat{U}(p) = e^{ip}$, so this also gives index 1. For unitaries acting on each site separately in the same way, we get agreement because $\text{ind } U = 0$, and $\widehat{U}(p)$ is independent of p . Note also that both sides of 20 have the same behavior under composition and direct sums. This proves the formula for all walks which can be composed of shifts and sitewise rotations. Actually, *all* translationally invariant walks can be represented in this way [7], but we prefer to give a direct proof of the proposition without invoking this decomposition.

Proof. From Definition 2 we get

$$\begin{aligned} \text{ind } U &= \sum_{x=0}^{\infty} \sum_{y=-\infty}^{-1} \text{Tr} |U_{x-y}|^2 - \text{Tr} |U_{y-x}|^2 \\ &= \sum_{n=-\infty}^{\infty} n \text{Tr} |U_n|^2 = \sum_{nm} \delta_{nm} n \text{Tr}(U_m^* U_n) \\ &= \frac{1}{2\pi i} \int_{-\pi}^{\pi} dp \text{Tr} \left(\widehat{U}(p)^* \frac{d\widehat{U}(p)}{dp} \right). \end{aligned}$$

On the other hand, for any invertible matrix function \widehat{U} ,

$$\frac{d}{dp} \det \widehat{U}(p) = \det \widehat{U}(p) \text{Tr} \left(\widehat{U}(p)^{-1} \frac{d\widehat{U}(p)}{dp} \right).$$

Hence with $\det \widehat{U}(p) = \exp(ipn)$ the above integrand is actually constant equal to in , and $\text{ind}(U) = n$. \square

The properties of a walk depend crucially on the properties of the eigenvalues $e^{i\omega_1(p)}, \dots, e^{i\omega_d(p)}$ of $\widehat{U}(p)$ (see Fig. 3). Clearly, $p \mapsto \widehat{U}(p)$ is an analytic family of operators, so we can follow the eigenvalues as analytic curves [11]. The derivatives of the eigenvalues determine the *group velocity*: Let $P_\alpha(p)$ be the eigenprojection of $\widehat{U}(p)$ belonging to the eigenvalue $\exp i\omega_\alpha(p)$. Then the group velocity operator can be written as

$$G(p) = \lim_{t \rightarrow \infty} \frac{1}{t} X(t) = \sum_{\alpha} P_{\alpha}(p) \frac{d\omega_{\alpha}(p)}{dp}, \tag{21}$$

where X denotes the position operator, and at degenerate eigenvalues the projections $P_{\alpha}(p)$ have to be chosen appropriately, as dictated by analytic perturbation theory. The limit is in the strong resolvent sense. Hence the probability distribution for the selfadjoint operator G in a state ρ is equal to the asymptotic position distribution starting from ρ in “ballistic” scaling[1]. In particular, when the internal state is unpolarised, i.e., when the initial state is of the form $\rho = \sigma \otimes \mathbb{1}/d$, we find [12]

$$\langle X(t) \rangle = \langle X(0) \rangle + \frac{t}{d} \text{ind}(U).$$

In this sense the index has direct relevance as a mean speed of the walk.

It is, of course, suggestive to connect the topological properties of the eigenvalue curves in Fig. 3 with the index. In principle, these curves are allowed to cross each other.

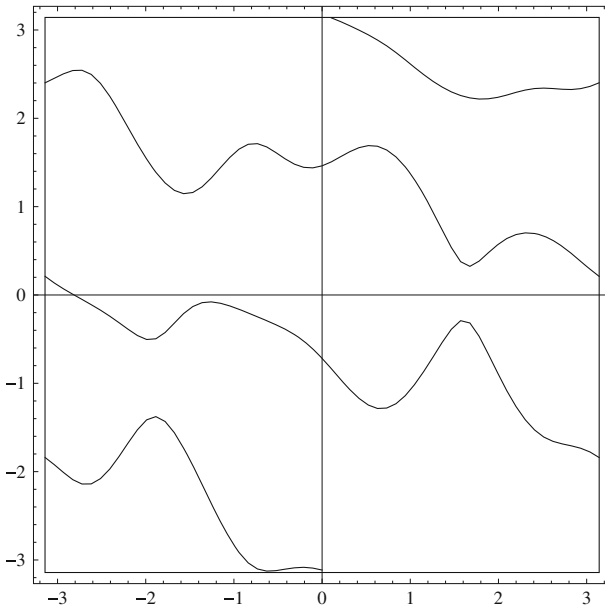


Fig. 3. Example for eigenvalues of $\widehat{U}(p)$, with $\dim \mathcal{K} = 3$, $L = 5$, $\text{ind } U = -1$. In this case the spectrum is a single curve on the torus. The index can be computed by the signed number of crossings of any horizontal line, or as the sum of the derivatives of all branches

So in general, we have several branches of curves, which wind several times around the torus before closing. The number of curves and their winding numbers would appear to be a topological invariant, but this is not true because of the “avoided crossing” phenomenon, in which a small perturbation of the walk turns an eigenvalue crossing into a close encounter of two separate curves (suggested also at $p = 1.6$ in Fig. 3). Hence only the *sum of the winding numbers* is stable with respect to small perturbations, and this is indeed equal to $\text{ind } U$. To see this, note that at every p the sum of the derivatives of all branches is equal to the index. Therefore the sum of the winding angles of the branches taken from $p = -\pi$ to $p = \pi$ is $2\pi \text{ind } U$. The winding angles of the closed curve components are just suitable sub-sums of this, and this partition is changed at avoided crossing points.

We close this section by establishing a variant of the main theorem for the translationally invariant case. Of course, most of this follows by simply specializing. The only question which requires a new argument is whether the path connecting two translationally invariant walks can be chosen to go entirely through translationally invariant walks. This is established in the following proposition.

Proposition 6. *Let U be a translationally invariant causal unitary with $\text{ind } U = 0$. Then we can find a norm continuous path $t \mapsto U_t$ of translationally invariant causal unitaries of bounded width such that $U_0 = \mathbb{1}$ and $U_1 = U$.*

Proof. We use the factorization theorem for “paraunitary” operators [7] (see also [17]), which states that $\widehat{U}(p)$ can be written as a finite product

$$\widehat{U}(p) = V_0 \prod_{k=1}^r \widehat{W}_{m_k}(p) V_k$$

of constant unitaries V_k and the elementary causal unitaries

$$\widehat{W}_m(p) = \begin{pmatrix} e^{imp} & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix}.$$

As a quantum walk these correspond to a shift of only the first internal state by m positions, leaving all other internal states fixed. This walk has index m , and the product formula yields $\text{ind } U = \sum_k m_k$. The maximal polynomial degree of any matrix element in $e^{\pm ip}$ depends on the localization region $[x_+, -x_-]$, and is clearly bounded by $L_{\max} = \sum_k |m_k|$. Hence we can contract the walk to $\mathbb{1}$ by contracting each of the V_k to $\mathbb{1}$, never exceeding width L_{\max} on the way. \square

7. The Index for Cellular Automata

For cellular automata we proceed in analogy to the case of walks, i.e., by defining directly a locally computable quantity as the index $\text{ind } \alpha$ of a walk automorphism α . We then establish that it is actually a complete locally computable invariant and, at the same time that it characterizes the connected components of QCAs.

As a technical preparation we need some background on commutation properties of algebras spanning several cells. It is basically taken from [16], and included here to make the presentation here self-contained.

7.1. Support algebras. For defining the index we need to find a quantitative characterization of “how much” of one cell ends up in another. To this end we introduce the notion of support algebras. Consider a subalgebra $\mathcal{A} \subset \mathcal{B}_1 \otimes \mathcal{B}_2$ of a tensor product. What is the position of \mathcal{A} relative to the tensor structure? Here we answer a relatively simple part of this question: which elements of $\mathcal{B}_1, \mathcal{B}_2$ are actually needed to build \mathcal{A} ?

For the following definition with lemma, recall that \mathcal{A}' denotes the commutant $\{a | \forall a_1 \in \mathcal{A} : [a, a_1] = 0\}$.

Lemma 7. *Let \mathcal{B}_1 and \mathcal{B}_2 be finite dimensional C^* -algebras, and $\mathcal{A} \subset \mathcal{B}_1 \otimes \mathcal{B}_2$ a subalgebra. Then*

1. *There is a smallest C^* -subalgebra $\mathcal{C}_1 \subset \mathcal{B}_1$ such that $\mathcal{A} \subset \mathcal{C}_1 \otimes \mathcal{B}_2$. We call this the **support algebra** of \mathcal{A} on \mathcal{B}_1 , and denote it by $\mathcal{C}_1 = \mathbf{S}(\mathcal{A}, \mathcal{B}_1)$.*
2. *Consider a basis $\{e_\mu\} \subset \mathcal{B}_2$, so that every $a \in \mathcal{A}$ has a unique expansion $a = \sum_\mu a_\mu \otimes e_\mu$ with $a_\mu \in \mathcal{B}_1$. Then $\mathbf{S}(\mathcal{A}, \mathcal{B}_1)$ is generated by all the elements a_μ arising in this way.*
3. *The commutant of $\mathbf{S}(\mathcal{A}, \mathcal{B}_1)$ in \mathcal{B}_1 is characterized as $\{b \in \mathcal{B}_1 | b \otimes \mathbb{1} \in \mathcal{A}'\}$.*

Proof. We can pick out the terms a_μ by applying a suitable functional ω_μ from the dual basis to the second factor, i.e., by applying the map $\text{id} \otimes \omega_\mu : \mathcal{B}_1 \otimes \mathcal{B}_2 \rightarrow \mathcal{B}_1$, which takes $b_1 \otimes b_2$ to $\omega_\mu(b_2)b_1$. Clearly, if $a \in \mathcal{C}_1 \otimes \mathcal{B}_2$, so that a can be expanded into simple tensors using only elements from \mathcal{C}_1 in the first factor, we find $a_\mu = (\text{id} \otimes \omega_\mu)(a) \in \mathcal{C}_1$. Hence the algebra described in Item 2 must be contained in any \mathcal{C}_1 satisfying Item 1. Since it also satisfies Condition 1, we have identified the unique smallest \mathcal{C}_1 . The Characterization 3 follows by looking at commutators of the form $[b \otimes \mathbb{1}, a]$, and expanding a as above. \square

This construction was introduced to the QI community by Zanardi [19], who applied it to the algebra generated by an interaction Hamiltonian, and consequently called it an “interaction algebra”. Of course, we can apply the construction also to the second factor, so that

$$\mathcal{A} \subset \mathbf{S}(\mathcal{A}, \mathcal{B}_1) \otimes \mathbf{S}(\mathcal{A}, \mathcal{B}_2) \subset \mathcal{B}_1 \otimes \mathcal{B}_2 . \tag{22}$$

The crucial fact we need about support algebras is that “commutation of algebras with overlapping localization happens on the intersection”. More precisely, we have the following

Lemma 8. *Let $\mathcal{A}_1 \subset \mathcal{B}_1 \otimes \mathcal{B}_2$ and $\mathcal{A}_2 \subset \mathcal{B}_2 \otimes \mathcal{B}_3$ be subalgebras such that $\mathcal{A}_1 \otimes \mathbb{1}_3$ and $\mathbb{1}_1 \otimes \mathcal{A}_2$ commute in $\mathcal{B}_1 \otimes \mathcal{B}_2 \otimes \mathcal{B}_3$. Then $\mathbf{S}(\mathcal{A}_1, \mathcal{B}_2)$ and $\mathbf{S}(\mathcal{A}_2, \mathcal{B}_2)$ commute in \mathcal{B}_2 .*

Proof. Pick bases $\{e_\mu\} \subset \mathcal{B}_1$ and $\{e'_\nu\} \subset \mathcal{B}_2$, and let $a \in \mathcal{A}_1$ and $a' \in \mathcal{A}_2$. Then we may expand uniquely: $a = \sum_\mu e_\mu \otimes a_\mu$ and $a' = \sum_\nu a'_\nu \otimes e'_\nu$. Then by assumption

$$0 = [a \otimes \mathbb{1}_3, \mathbb{1}_1 \otimes a'] = \sum_{\mu\nu} e_\mu \otimes [a_\mu, a'_\nu] \otimes e'_\nu .$$

Now since the elements $e_\mu \otimes e'_\nu$ are a basis of $\mathcal{B}_1 \otimes \mathcal{B}_3$, this expansion is unique, so we must have $[a_\mu, a'_\nu] = 0$ for all μ, ν . Clearly, this property also transfers to the algebras generated by the a_μ and a'_ν , i.e., to the support algebras introduced in the lemma. \square

7.2. Defining the index. Now consider a cellular automaton α on a cell structure \mathcal{A}_x . By regrouping, if necessary, we may assume that it has only nearest neighbor interactions. Now consider any two neighboring cells $\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}$, and their image under α , i.e.,

$$\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}) \subset (\mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x}) \otimes (\mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2}) .$$

We apply the support algebra construction to this inclusion, obtaining two algebras

$$\mathcal{R}_{2x} = \mathbf{S}(\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), (\mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x})) , \tag{23}$$

$$\mathcal{R}_{2x+1} = \mathbf{S}(\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), (\mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2})) . \tag{24}$$

Intuitively, the algebras \mathcal{R}_y with even index become larger when information flows to the right, whereas the ones with odd index describe a flow to the left. Of course, this intuition will be made precise below, but at this stage one can at least check these statements for the shift: When $\alpha(\mathcal{A}_y) = \mathcal{A}_{y-1}$ (recall the convention made at the end of Sect. 3.2) we have $\mathcal{R}_{2x} = \mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x}$ and $\mathcal{R}_{2x+1} = \mathbb{C}\mathbb{1}$. This is to be contrasted with $\mathcal{R}_y = \mathcal{A}_y$ for the identity.

Continuing with our construction, observe that by Lemma 8, the subalgebras \mathcal{R}_{2x+1} and \mathcal{R}_{2x+2} commute in the algebra $(\mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2})$. Algebras \mathcal{R}_y which are further away commute anyhow, since they are contained in disjoint cells. We conclude that *all* \mathcal{R}_y commute.

By definition of support algebras, $\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}) \subset \mathcal{R}_{2x} \otimes \mathcal{R}_{2x+1}$, so that the algebras \mathcal{R}_x together generate an algebra containing $\alpha\mathcal{A}(\mathbb{Z})$. Because α is an automorphism, this is the same as $\mathcal{A}(\mathbb{Z})$. Now if any \mathcal{R}_x had a non-trivial center, i.e., if there were an element $X \in \mathcal{R}_x$ commuting with all of \mathcal{R}_x but not a multiple of $\mathbb{1}$, this X would also be

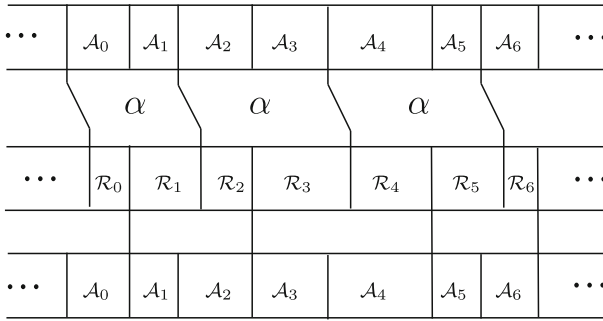


Fig. 4. Cell structure with support algebras. (Read top to bottom) If the width of cells is taken as log dimension, the index can be read off the slope in the boxes representing mapping by α

in the center of the entire quasi-local algebra $\mathcal{A}(\mathbb{Z})$. However, this center is known to be trivial [5]. We conclude that each \mathcal{R}_x must have trivial center, and hence be isomorphic to $\mathcal{M}_{r(x)}$ for some integer $r(x)$. Moreover, the inclusion noted at the beginning of this paragraph cannot be strict, since otherwise we would find an element in the relative commutant, which would once again be in the center $\alpha\mathcal{A}(\mathbb{Z})$. To summarize, we must have

$$\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}) = \mathcal{R}_{2x} \otimes \mathcal{R}_{2x+1}, \tag{25}$$

$$\text{hence } d(2x)d(2x + 1) = r(2x)r(2x + 1). \tag{26}$$

On the other hand, the commuting full matrix algebras \mathcal{R}_{2x+1} and \mathcal{R}_{2x+2} together span the tensor product isomorphic to $\mathcal{M}_{r(2x+1)r(2x+2)}$ inside $(\mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2})$. Again the inclusion cannot be strict, because otherwise the automorphism would not be onto. From this we get the second relation and dimension equation

$$\mathcal{R}_{2x+1} \otimes \mathcal{R}_{2x+2} = \mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2}, \tag{27}$$

$$\text{hence } r(2x + 1)r(2x + 2) = d(2x + 1)d(2x + 2).$$

These relations are summarized pictorially in Fig. 4. They give us the the first two equalities in

$$\frac{r(2x)}{d(2x)} = \frac{d(2x + 1)}{r(2x + 1)} = \frac{r(2x + 2)}{d(2x + 2)} \equiv \text{ind } \alpha. \tag{28}$$

In other words, this quantity is constant along the chain, and hence we are entitled to define it as the index $\text{ind } \alpha$. The even/odd asymmetry only comes from the construction, by which the even \mathcal{R}_{2x} describe a flow to the right (increasing x), and the odd \mathcal{R}_{2x+1} are associated with a flow to the left. By shifting the entire construction, we could switch the even/odd distinction, and define, for any \mathcal{A}_x , both the ascending and the descending \mathcal{R}_x . In any case, for the shift σ_d of a d -dimensional chain we get $d(y) \equiv d$, $r(2x) = d^2$, and $r(2x - 1) = 1$, and hence $\text{ind } \sigma_d = d$, as announced in Eq. (15).

It is part of the local computability property that we have a lot of freedom in choosing the cell structure for which we want to evaluate the index. Since one typically wants to use this freedom to simplify the computation, we will now summarize the constraints. It is clear that there are three subalgebras involved in the computation, playing the rôle of, say the above $\mathcal{A}_0, \mathcal{A}_1, \mathcal{A}_2$ for determining \mathcal{R}_1 . Let us call these $\mathcal{A}_L, \mathcal{A}_M, \mathcal{A}_R$ to

emphasize that these algebras need not be part of the original cell structure, and we are free to choose them within certain limits. Let us start by fixing some algebra $\mathcal{A}_M \cong \mathcal{M}_d$, a full matrix algebra contained in some local algebra, whose crucial property is to split the system: we must have an isomorphism of the total algebra with $\mathcal{A}_{<M} \otimes \mathcal{A}_M \otimes \mathcal{A}_{>M}$, where the outer factors contain suitable infinite half chain algebras, and such that

$$[\alpha(\mathcal{A}_{<M}), \mathcal{A}_{>M}] = [\alpha(\mathcal{A}_{>M}), \mathcal{A}_{<M}] = \{0\}.$$

Clearly, this imposes a lower bound on the size of \mathcal{A}_M in terms of the interaction length of the automorphism. Now we choose finite dimensional matrix subalgebras $\mathcal{A}_R \subset \mathcal{A}_{>M}$ and $\mathcal{A}_L \subset \mathcal{A}_{<M}$ such that $\alpha(\mathcal{A}_M) \subset \mathcal{A}_L \otimes \mathcal{A}_M \otimes \mathcal{A}_R$. These three can be taken as part of a nearest neighbor cell structure, so that the above arguments give

$$\left. \begin{aligned} \mathcal{A}_M &\cong \mathcal{M}_d \\ \mathcal{S}(\alpha(\mathcal{A}_L \otimes \mathcal{A}_M), \mathcal{A}_M \otimes \mathcal{A}_R) &\cong \mathcal{M}_r \end{aligned} \right\} \Rightarrow \text{ind } \alpha = \frac{d}{r}. \tag{29}$$

Note that there is no harm in choosing \mathcal{A}_R larger than necessary: the support algebra, being the minimal algebra needed to build the tensor product, will simply not change. By a similar argument, we can choose \mathcal{A}_L too large without changing this support algebra.

7.3. Fundamental properties of the index for cellular automata.

Theorem 9. 1. *ind α is a positive rational for every α . When the automaton is regrouped in nearest neighbor form, both the numerator and the denominator of ind α in canceled form divide every cell dimension.*

2. *ind α is locally computable, and uniquely characterizes the equivalence classes for the relation \sim from Sect. 5. It can hence be identified with the abstract index defined there.*
3. *ind $(\alpha \otimes \alpha') = \text{ind } \alpha \text{ind } \alpha'$ and, when α and α' are defined on the same cell structure, $\text{ind } (\alpha\alpha') = \text{ind } \alpha \text{ind } \alpha'$.
If, for some y , $\alpha(\mathcal{A}((-\infty, y])) \subset \mathcal{A}((-\infty, y))$, then $\text{ind } \alpha \in \mathbb{N}$.
Moreover, for the shift of d -dimensional cells: $\text{ind } \sigma_d = d$.*
4. *An automaton α has index 1 if and only if it can be implemented locally. In this case it can be written as a product of just two partitioned unitary automorphisms. If α is partitioned in nearest neighbor form, the partitioned automorphisms can be taken to couple pairs of nearest neighbors only.*
5. *Two automata α_0, α_1 on the same cell structure have the same index if and only if they can be deformed to each other; i.e., there is a strongly continuous path $[0, 1] \ni t \mapsto \alpha_t$ of automorphisms, all with the same neighborhoods, and with the specified boundary values.*

Proof. Items (1,3) $\text{ind } \alpha \in \mathbb{Q}_+$ follows immediately from the construction in Sect. 7.2, particularly Eq. (28). Let $\text{ind } \alpha = \frac{p}{q}$ be the fraction in canceled form. Then, from this equation $p = nr(2x)$ and $q = nd(2x)$, where n is the canceled factor. Hence q divides $d(2x)$ and from the second fraction in (28) we find that p divides $d(2x + 1)$. By shifting the construction by one cell, we find the remaining divisibility statements. From Eq. (28) we also get the product formulas Item 3. Suppose that α maps some left half chain into itself. Then we choose a cell partition so that $y = 2x + 1$ in the setting of (23). Then $\mathcal{R}_{2x+1} \subset \mathcal{A}_{2x+1}$, and since these are full matrix algebras with the same unit, the quotient $\text{ind } \alpha = d_{2x+1}/r_{2x+1}$ in (28) is integer. The value for the shift was verified as an example after (28).

Items (2,4,5) From Eq. (28) local computability is obvious. From the general discussion in Sect. 5 we also get that locally implementable α have $\text{ind } \alpha = 1$. The hard part, which is needed to identify the abstract index with the concrete formula is the converse. For this the crucial step is the following claim:

Let α and α' be nearest neighbor cellular automata on the same cell structure and with the same index. Then there are unitaries $V_x \in \mathcal{A}_x \otimes \mathcal{A}_{x+1}$ such that the two locally implemented automorphisms

$$\beta(A) = \left(\prod_x V_{2x-1}\right)^* A \left(\prod_x V_{2x-1}\right),$$

$$\gamma(A) = \left(\prod_x V_{2x}\right)^* A \left(\prod_x V_{2x}\right)$$

satisfy: $\alpha'\gamma = \beta\alpha$.

Before proving this claim, let us see how it implies the statements in the theorem. By the general theory of Sect. 5 locally implementable automorphisms β, γ are \sim -equivalent to the identity, and $\alpha \sim \beta\alpha = \alpha'\gamma \sim \alpha'$. Hence equality of the indices as defined by (29) implies crossover equivalence, and hence the equality of all locally computable invariants. This proves Item 2. The converse in Item 4 follows by taking $\alpha' = \text{id}$, giving the local implementation $\alpha = \beta^{-1}\gamma$ of any index 1 automorphism α .

Finally, it is clear for Item 5 that we can connect α and α' with the same index by the required continuous path: we just need to contract each unitary V_x in β, γ to the identity, to obtain a path $\alpha_t = \beta_t \alpha \gamma_t^{-1}$ with $\alpha_0 = \alpha$ and $\alpha_1 = \alpha'$. This path will not be continuous in the norm on automorphisms, i.e., we cannot make $\|\alpha_t - \alpha_s\|$ small, since this would already fail for one-site operations $\alpha_t(A) = \left(U_t^{\otimes \infty}\right)^* A U_t^{\otimes \infty}$ with $t \mapsto U_t$ norm continuous. However, for any finitely localized observable A , $t \mapsto \alpha_t(A)$ will be continuous in norm, which is the claim of strong continuity made in the theorem. We remark that an important part of the proof of Item 5 is missing at this point: We did not exclude the possibility that there are continuous paths linking automorphisms of different index. This will be achieved by Prop. 13, an expression for the index which is manifestly continuous with respect to strongly continuous deformations.

Now to prove the claim, let α and α' have the same cell structure and the same index. Then we carry out the construction of Sect. 7.2 for both automorphisms, resulting in some intermediate algebras \mathcal{R}_x and \mathcal{R}'_x .

Since the indices coincide, formula (28) demands that these are full matrix algebras of the same dimensions. For example,

$$\mathcal{R}_{2x-1} \otimes \mathcal{R}_{2x} = \mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x} = \mathcal{R}'_{2x-1} \otimes \mathcal{R}'_{2x}.$$

Clearly, there is a unitary operator $V_{2x-1} \in \mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x}$ so that $V_{2x-1}^* \mathcal{R}_y V_{2x-1} = \mathcal{R}'_y$ for $y = 2x$ and for $y = 2x - 1$. We can take all these unitaries together as implementing one partitioned automorphism $\beta(A) = \left(\prod_x V_{2x-1}\right)^* A \left(\prod_x V_{2x-1}\right)$. By definition, it satisfies $\beta(\mathcal{R}_y) = \mathcal{R}'_y$ for all $y \in \mathbb{Z}$.

Now consider the action of α and α' on $\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}$. We now get two isomorphisms

$$\alpha' : \mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1} \rightarrow \mathcal{R}'_{2x} \otimes \mathcal{R}'_{2x+1} \quad \text{and}$$

$$\beta\alpha : \mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1} \rightarrow \mathcal{R}_{2x} \otimes \mathcal{R}_{2x+1} \rightarrow \mathcal{R}'_{2x} \otimes \mathcal{R}'_{2x+1}.$$

Hence $(\alpha')^{-1}\beta\alpha$ restricts to an automorphism of $\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}$, and can therefore be implemented by a unitary $V_{2x} \in \mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}$. These unitaries together implement γ , and we get the desired equation $\alpha'\gamma = \beta\alpha$. \square

7.4. Index for classical reversible automata. In this section, we will review a common notion of index for reversible classical cellular automata, and show that it coincides with our definition. In the context of this paper, a reversible classical cellular automaton can be defined as a particular case of a quantum cellular automaton. In each cell \mathcal{A}_x we single out a maximal abelian subalgebra \mathcal{D}_x . With respect to a suitable choice of basis, \mathcal{D}_x is then the set of diagonal matrices. As a finite dimensional abelian algebra, we can regard \mathcal{D}_x as the set of complex valued functions on a finite set \mathbb{A}_x , called the *alphabet* of the cell, which at the same time serves as the set of basis labels for the orthonormal basis in which \mathcal{D}_x is diagonal. The global C^* -algebra of the classical system is then the infinite tensor product $\mathcal{D}(\mathbb{Z}) = \bigotimes_{x=-\infty}^{\infty} \mathcal{D}_x$. It is canonically isomorphic to the algebra of continuous functions on the compact cartesian product space $\mathbb{A}_{\mathbb{Z}} = \prod_{x=-\infty}^{\infty} \mathbb{A}_x$, also known as the space of infinite configurations. We use lower case letters such as c for such configurations, and denote by $c(x) \in \mathbb{A}_x$ the configuration of the cell at x .

Now let α be a QCA with the property that $\alpha(\mathcal{D}) \subset \mathcal{D}$. Then the restriction of α to \mathcal{D} is an automorphism of \mathcal{D} , which must be of the form $(\alpha f)(c) = f(\Phi(c))$, where Φ is a homeomorphism on configurations. The causality conditions on α are readily expressed in terms of Φ , and, together with analogous arguments for the inverses show that Φ is a reversible classical cellular automaton in the usual sense, apart from the requirement of translation invariance. There are some subtle points to note about the correspondence $\alpha \rightarrow \Phi$:

- Suppose that in the above argument we start from a general, not necessarily strictly causal automorphism α of the quasi-local algebra $\mathcal{A}(\mathbb{Z})$. We still get a continuous Φ on the compact space $\mathbb{A}_{\mathbb{Z}}$. By the definition of the product topology this means that the local configurations $\Phi(c)_x$ after the time step depend on only finitely many $c(y)$. In the translation invariant case this means that Φ is a cellular automaton with finite neighborhood. In fact, this argument is used to establish that the inverse of a reversibly cellular automaton also has a finite neighborhood. This is a rather surprising sharpening of the causality condition. However, we are appealing here to a highly non-constructive compactness argument, which gives no control on the size of the neighborhoods, or (barring translation invariance) on the uniformity of the neighborhoods. For example, we can apply it to a cellular automaton α on a 2D lattice, whose quasi-local C^* -algebra is isomorphic to that of a 1D automaton. Hence the condition of α being an automorphism is not strong enough to give a 1D automaton in the sense defined above.
- The mapping is onto, i.e., every classical reversible cellular automaton can be “quantized”. The argument is very simple for finite lattices, e.g., a regular lattice with periodic boundary conditions: one labels the basis of a Hilbert space by the classical configurations. Then the classical automaton Φ is a permutation of the basis vectors, which can be interpreted as a unitary operator via $U_{\Phi}|c\rangle = |\Phi(c)\rangle$. Then for all observables we set $\alpha(A) = U_{\Phi}^* A U_{\Phi}$. One needs to check that this transformation is causal in the quantum sense [16], in particular that off-diagonal local operators that are finitely localized (i.e. localized on a finite number of cells) keep this property under the action of α . Indeed one gets a bound on the quantum neighborhoods, which involves both the neighborhoods of Φ and the neighborhoods of Φ^{-1} . The same computation provides a formula for $\alpha(A)$, for A finitely localized, in terms of the classical rule, and this can be used to define α also for infinite lattices. We will not, however, make this explicit here.
- The mapping $\alpha \rightarrow \Phi$ is not injective. Indeed an ambiguity is inherent in the construction just described: we can choose different bases with the same diagonal operators

$|a\rangle\langle a|$ by choosing a phase for each basis vector. This amounts to changing α by an on-site unitary, which is certainly irrelevant for index purposes. But we can consider this more generally: suppose that two cellular automata α, β restrict to the same automorphism on the diagonal algebra \mathcal{D} . Then $\alpha\beta^{-1}$ leaves \mathcal{D} point-wise fixed. In a finite lattice, so that α is unitarily implemented, the implementing unitary hence commutes with all elements of the maximally abelian algebra \mathcal{D} , hence must itself be diagonal. It is suggestive that this also holds in a localized form on the infinite lattice, so $\alpha\beta^{-1}$ would be a product of commuting unitaries and hence have trivial index. The following result builds on this intuition.

Proposition 10. *Suppose that α and β be quantum cellular automata taking the diagonal algebra \mathcal{D} into itself, with the same restriction to \mathcal{D} . Then $\text{ind } \alpha = \text{ind } \beta$.*

Before coming to the proof of this proposition we single out two arguments of independent interest, each of which can be used to draw the main conclusion, without discussing in detail the structure of local phase factors. The first criterion uses the absence of propagation. The second uses the global transpose map $\Theta : \mathcal{A}(\mathbb{Z}) \rightarrow \mathbb{Z}$. It is defined as the matrix transpose on each local algebra, in a basis in which \mathcal{D} is diagonal. Since transposition is isometric, and consistent with the embeddings $A \mapsto A \otimes \mathbb{1}$ it extends to the whole algebra. Θ is a linear anti-homomorphism (meaning $\Theta(AB) = \Theta(B)\Theta(A)$), and, for every automorphism α , $\Theta\alpha\Theta$ is again an automorphism.

Lemma 11. *1. Let α be a nearest neighbor cellular automaton such that, for some finite interval $[z_-, z_+]$ we have $\alpha^n(\mathcal{A}([0, 1])) \subset \mathcal{A}([z_-, z_+])$ for all $n \in \mathbb{Z}$. Then $\text{ind } \alpha = 1$.
 2. For any cellular automaton α : $\text{ind } (\Theta\alpha\Theta) = \text{ind } \alpha$.*

Proof. 1. The index $\text{ind } \alpha^n$ can be expressed as a ratio of subcell dimensions of $\mathcal{A}([z_-, z_+])$, hence is uniformly bounded in n . But since $\text{ind } \alpha^n = (\text{ind } \alpha)^n$ this implies $\text{ind } \alpha \leq 1$. With the same argument for the inverse we get $\text{ind } \alpha \geq 1$.

2. By assumption, the global transposition is made with respect to product basis, so that for a tensor product $\mathcal{B} = \mathcal{B}_1 \otimes \mathcal{B}_2$ of cells we get $\Theta_{\mathcal{B}} = \Theta_{\mathcal{B}_1} \otimes \Theta_{\mathcal{B}_2}$. We will drop the indices on Θ in the sequel. Then it is clear from the definition (Lemma 7) that the support algebra construction behaves naturally under global transposition, i.e., we have $\mathbf{S}(\Theta\mathcal{A}, \Theta\mathcal{B}_1) = \Theta\mathbf{S}(\mathcal{A}, \mathcal{B}_1)$. Moreover, when \mathcal{B}_1 is a tensor product of cells we get $\Theta\mathcal{B}_1 = \mathcal{B}_1$. Hence in (23) we find for the automorphism $\tilde{\alpha} = \Theta\alpha\Theta$:

$$\begin{aligned} \tilde{\mathcal{R}}_{2x} &= \mathbf{S}\left(\Theta\alpha\Theta(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), (\mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x})\right) \\ &= \mathbf{S}\left(\Theta\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), \Theta(\mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x})\right) \\ &= \Theta\mathbf{S}\left(\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), (\mathcal{A}_{2x-1} \otimes \mathcal{A}_{2x})\right) = \Theta\mathcal{R}_{2x}. \end{aligned}$$

Since this has the same dimension as \mathcal{R}_{2x} , we find from (28) that $\text{ind } (\Theta\alpha\Theta) = \text{ind } \alpha$. □

Proof (of Prop. 10). Due to the multiplication formula, we only need to consider the case that α is equal to the identity ($= \beta$) on \mathcal{D} . Consider any operator $A_x \in \mathcal{A}_x$ in a single cell, and let $D_y \in \mathcal{D}_y$ be a diagonal element in another cell. Then

$$\alpha(A_x)D_y = \alpha(A_x)\alpha(D_y) = \alpha(A_x D_y) = \alpha(D_y A_x D_y) = D_y \alpha(A_x).$$

Hence the finitely localized element $\alpha(A_x)$ commutes with all diagonal operators on the neighboring sites, and $\alpha(A_x) \in \mathcal{A}_x \otimes \bigotimes_{y \neq x} \mathcal{D}_y$. This algebra is best seen as a direct sum of copies of \mathcal{A}_x , labeled by configurations $c = \{c(y)\}$ with $c(y) \in \mathbb{A}_y$ of all cells $y \neq x$ in the localization region of $\alpha(\mathcal{A}_x)$. A homomorphism of \mathcal{A}_x into this algebra splits into one homomorphism into each summand, which in turn is given by a unitary $U_x(c)$. Hence we can summarize the action of α on \mathcal{A}_x as

$$\alpha(A_x) = \sum_c U_x(c)^* A_x U_x(c) \otimes P(c), \tag{30}$$

where P_c is the minimal projection of the diagonal algebra corresponding to c . We also know that diagonal elements of \mathcal{A}_x are fixed, so $U_x(c)$ is itself diagonal, say $U_x(c)|a\rangle = u(a, c)|a\rangle$. This leads to

$$\alpha(|a\rangle\langle b|_x) = \sum_c \frac{u(b, c)}{u(a, c)} |a\rangle\langle b|_x \otimes P(c). \tag{31}$$

The commutation of $\alpha(\mathcal{A}_x)$ and $\alpha(\mathcal{A}_y)$ for $x \neq y$ introduces further conditions on the phase functions u . But rather than analyzing these in detail, we use the lemma to conclude directly from Eq. (31) that $\text{ind } \alpha = 1$. To this end, note that by applying the homomorphism α to (31) and using $\alpha(P(c)) = P(c)$, we get a corresponding formula for the iterate of α :

$$\alpha^n(|a\rangle\langle b|_x) = \sum_c \left(\frac{u(b, c)}{u(a, c)} \right)^n |a\rangle\langle b|_x \otimes P(c), \tag{32}$$

for any $n \in \mathbb{Z}$. Clearly, the localization region of this operator does not increase with n , so by the first part of the lemma we get $\text{ind } \alpha = 1$. Alternatively, we can apply Θ to the equation, using $\Theta(P(c)) = P(c)$. This reverses each of the ketbra operators, so $\Theta \alpha \Theta = \alpha^{-1}$. Hence $\text{ind } \alpha = 1$ also follows with the second part of the lemma. \square

An index for classical reversible cellular automata has been defined albeit only in the translationally invariant case. According to G. A. Hedlund ([8], Sect. 14), the definition is due to L. R. Welch, so we will call it the Welch index i_W here. For the definition itself we will follow Kari [10], where it is introduced in Sect. 3. We will show that this coincides with the quantum index. Hence the quantum index is a possible extension to non-translationally invariant systems. It is very likely that the theory in [10] can also be extended directly, but we have not gone to the trouble of checking all the details.

In the translation invariant case, all cell alphabets $\mathbb{A}_x \equiv \mathbb{A}$ are the same, and a cellular automaton is a map $\phi : \mathbb{A}^{\mathbb{Z}} \rightarrow \mathbb{A}^{\mathbb{Z}}$. Let r be a ‘‘large enough’’ integer, and R_ϕ^r the set of $4r$ -tuples of the form

$$\left(c(0), \dots, c(2r - 1), (\phi c)(-r), \dots, (\phi c)(r - 1) \right),$$

where c runs over all infinite configurations. Then the Welch index of ϕ is defined as

$$i_W(\phi) = \frac{|R_\phi^r|}{|\mathbb{A}|^{3r}}. \tag{33}$$

Clearly, for the identity only $|\mathbb{A}|^3$ different letters occur here, so $i_W(\text{id}) = 1$. Similarly, for a shift we get $i_W(S) = |\mathbb{A}|$, and parallel application of ϕ and ϕ' to parallel chains yields $i_W(\phi \times \phi') = i_W(\phi)i_W(\phi')$.

The non-trivial results about the index, and the structure theory of reversible, translationally invariant classical cellular automata is developed in [10], with key results analogous to our paper: the expression (33) does not depend on r (provided it is large enough). The product formula holds, and an automaton ϕ is locally implementable iff $i_W(\phi) = 1$. Moreover, every such automaton can be decomposed into shifts and locally implementable ones. Since a classical “local implementation” implies a partitioned representation of the quantum automaton, we can put these facts together to conclude that

$$i_W(\phi) = \text{ind } \alpha, \tag{34}$$

for any quantum cellular automaton which restricts on the diagonal subalgebra to a classical CA given by ϕ . In this sense our theory is a direct generalization of Kari’s work, extended by the aspects of deformation classes (which make no sense in the classical discrete setting) and local computability (which makes no sense in the translation invariant setting).

7.5. Interlude: More analogies between walks and cellular automata. The two definitions of the index, (16) for walks and (29) for cellular automata are not directly analogous. Here we would like to point out the differences, and discuss how to make the analogy between these two cases even tighter by supplying the missing analogous definitions.

The definition (29) considers a part of the system split into three parts L-M-R. Based on suitable inclusions, it gives a formula for the index, which immediately makes obvious that it is always a positive rational. In contrast, the walk expression Eq. (16) is a difference of numbers which can take arbitrary real positive values, and only one cut of the system is considered. Moreover, (16) made it very easy to prove the continuity of the index under deformations, whereas neither (29) nor the abstract considerations of Sects. 4 and 5 clarify continuity for the index of cellular automata. Since continuity is an important feature of our index theory, we will need an appropriate expression also for the automaton case, and the analogies laid out in this subsection are intended to motivate the form of this formula.

Dimension based formula for walks. Let us first set up an index formula for walks in analogy with (29). The analog of the support algebra is the “support subspace of \mathcal{K}_{12} in \mathcal{K}_2 ”, denoted $\mathbf{S}(\mathcal{K}_{12}, \mathcal{K}_2)$, which is defined for any subspace $\mathcal{K}_{12} \subset \mathcal{K}_1 \oplus \mathcal{K}_2$ of the orthogonal direct sum of Hilbert spaces. Namely, it is the smallest subspace $\mathcal{L} \subset \mathcal{K}_2$ such that $\mathcal{K}_{12} \subset \mathcal{K}_1 \oplus \mathcal{L}$. Then the analog of Lemma 8 holds in the sense that subspaces $\mathcal{K}_{12} \subset \mathcal{K}_1 \oplus \mathcal{K}_2$ and $\mathcal{K}_{23} \subset \mathcal{K}_2 \oplus \mathcal{K}_3$ are orthogonal iff $\mathbf{S}(\mathcal{K}_{12}, \mathcal{K}_2) \perp \mathbf{S}(\mathcal{K}_{23}, \mathcal{K}_2)$ are orthogonal. Now consider subspaces $\mathcal{H}_L \oplus \mathcal{H}_M \oplus \mathcal{H}_R \subset \mathcal{H}$ chosen with the localization constraints as in Sect. 7.2. In particular, we require $U\mathcal{H}_L \perp \mathcal{H}_R$, and $U\mathcal{H}_R \perp \mathcal{H}_L$. Then the direct analog of (29) reads

$$\begin{aligned} \text{ind } U &= \dim \mathbf{S}(U(\mathcal{H}_L \oplus \mathcal{H}_M), \mathcal{H}_M \oplus \mathcal{H}_R) - \dim \mathcal{H}_M \\ &= \text{rank}(P_{MR}U P_{LM}) - \dim \mathcal{H}_M. \end{aligned} \tag{35}$$

Here the second equality, in which P_{LM} is the projection onto $\mathcal{H}_L \oplus \mathcal{H}_M$ etc., follows with $\mathbf{S}(U(\mathcal{H}_L \oplus \mathcal{H}_M), \mathcal{H}_M \oplus \mathcal{H}_R) = P_{LM}U P_{MR}\mathcal{H}$.

Proof (of Eq. (35)). Consider the block matrix for U with respect to the decomposition

$$\mathcal{H} = \mathcal{H}_{-\infty} \oplus \mathcal{H}_L \oplus \mathcal{H}_M \oplus \mathcal{H}_R \oplus \mathcal{H}_{+\infty},$$

where the pieces at the ends contain the appropriate infinite half chains. Using the causality assumptions, we find that

$$U = \left(\begin{array}{ccc|cc} * & * & 0 & 0 & 0 \\ * & * & * & 0 & 0 \\ \hline 0 & \mathbf{*} & \mathbf{*} & * & 0 \\ 0 & \mathbf{0} & \mathbf{*} & * & * \\ \hline 0 & 0 & 0 & * & * \end{array} \right), \tag{36}$$

where the asterisks stand for any possibly non-zero block. We have highlighted in bold-face the block $P_{MR}U P_{LM}$ appearing in (35), and introduced two separating lines, namely the separation $-\infty, L, M | R, +\infty$ on the domain side (i.e., a vertical line) and the horizontal separation $-\infty, L | M, R, +\infty$ on the range side of U . These separators do *not* cross on the diagonal, which is why we cannot simply compute the index from the highlighted block via Eq. (16). However, this is easily amended by multiplying with a suitable shift: we can introduce a basis in each \mathcal{H}_x , and hence in \mathcal{H} , effectively making all underlying cell dimensions one-dimensional. In this representation we can introduce a shift operation S , and clearly $U' = S^{|M|}U$ will be a unitary with the same matrix elements as U shifted vertically by $|M| = \dim \mathcal{H}_M$. Obviously, $\text{ind } U' = \text{ind } U + |M|$, which explains the second term in (35), and leaves us with proving that $\text{ind } U' = \text{rank } (P_{MR}U P_{LM})$. Clearly, for this task the further subdivision of the blocks is irrelevant, and we can consider a general block decomposed unitary operator

$$U = \left(\begin{array}{c|c} U_{11} & 0 \\ \hline U_{21} & U_{22} \end{array} \right)$$

with a finite rank upper right corner. We have to show that $\text{rank } U_{21} = \text{Tr}(U_{21}^* U_{21})$. But from the unitarity equation it follows that U_{21} is a partial isometry, so $U_{21}^* U_{21}$ is its domain projection, whose dimension is indeed the rank of U_{21} . \square

Half neighborhoods and one-cut dimension formulas. The proof of formula (35) was essentially by reduction to the case of “half neighborhoods”, i.e., the case that $[x_-, x_+] = [x, x + 1]$, in which no influence ever spreads to the left. Then one of the off-diagonal blocks of the unitary entering (16) vanishes, and we saw directly that the other block gives an integer contribution, which can be interpreted as a dimension.

Similarly, for a half-neighborhood cellular automaton we demand $\alpha(\mathcal{A}_x) \subset \mathcal{A}_{x-1} \otimes \mathcal{A}_x$, where once again we have chosen the convention to match information traveling to the right, observing the Heisenberg picture. In particular, this condition is satisfied by the shift. For such automata we can simplify the index formula in a way quite analogous to the case of half-neighborhood walks. Indeed, setting

$$\mathcal{T}_x = \mathbf{S}(\alpha(\mathcal{A}_x), \mathcal{A}_x) \quad \text{and} \quad \mathcal{N}_x = \mathbf{S}(\alpha(\mathcal{A}_x), \mathcal{A}_{x-1}), \tag{37}$$

we can employ the same arguments as in Sect. 7 to conclude that these commute, and must be full matrix algebras $\mathcal{T}_x \cong \mathcal{M}_{t(x)}$ and $\mathcal{N}_x \cong \mathcal{M}_{n(x)}$. Then, further following the previous reasoning, $\alpha(\mathcal{A}_x) = \mathcal{N}_x \otimes \mathcal{T}_x$ and $\mathcal{A}_x = \mathcal{T}_x \mathcal{N}_{x+1}$. This yields the dimension equation

$$d(x) = n(x)t(x) = t(x)n(x - 1). \tag{38}$$

Hence the integers $n(x)$ do not depend on x . In fact, $\text{ind } \alpha = n(x)$. This follows readily from the observation that

$$\mathcal{R}_{2x+1} = \mathbf{S}\left(\alpha(\mathcal{A}_{2x} \otimes \mathcal{A}_{2x+1}), (\mathcal{A}_{2x+1} \otimes \mathcal{A}_{2x+2})\right) = \mathcal{T}_{2x+1} \otimes \mathbb{1}_{2x+2}.$$

Therefore, from Eq. (28) we get

$$\text{ind } \alpha = d(2x + 1)/r(2x + 1) = d(2x + 1)/t(2x + 1) = n(2x + 1).$$

In fact, for any cell structure on which a shift is available, we could have used this slightly simplified setup to define the index of any α by first shifting and regrouping to a half-neighborhood automaton, and correcting by a factor depending on the size of the necessary shift. However, since such a shift is not available in general, there was no gain in this approach.

For nearest neighbor automata the structure of support algebras using just a single cut, i.e., algebras of the form $\mathbf{S}(\alpha(\mathcal{A}_L), \mathcal{A}_R)$ are not sufficient to define the index. As a counterexample, consider a unitary evolution $A \mapsto U^*AU$ for $U \in \mathcal{A}_L \otimes \mathcal{A}_R$ with $\mathcal{A}_L = \mathcal{A}_R = \mathcal{M}_d$. Clearly, as a locally implementable operation, this has always trivial index. But the support algebra written above can be just about anything. For example, taking a ‘‘controlled unitary’’ $U = \sum_i |i\rangle\langle i| \otimes U_i$, we have $\mathbf{S}(\alpha(\mathcal{A}_R), \mathcal{A}_L)$ as the algebra of diagonal matrices, and $\mathbf{S}(\alpha(\mathcal{A}_L), \mathcal{A}_R)$ generated by the $U_i^*U_j$, which for $\mathcal{A}_R \cong \mathcal{M}_r$ can easily be \mathcal{A}_R . On the other hand, one can easily construct an automorphism with the same support algebras, but index d .

The algebraic structure of support algebras across a single cut is also insufficient for us in another way: it is a discrete structure, hence does not go to any trivial value as we deform an α to the identity. Therefore, if only to get the continuity of the index we need to look at some continuously varying quantities. We found a formula with just these properties by looking at cellular automaton analogs of the simple walk formula (16). The following section is devoted to the proof of this formula.

7.6. One-cut quotient formula. Notation for normalized traces. A continuously varying quantity measuring the difference between subalgebras and going to a trivial value as they come to coincide may be some kind of angle, or overlap, between the linear subspaces. It is natural to measure such angles with respect to the only scalar product between algebra elements, which is canonically defined in our context. Indeed, let τ denote the *normalized trace* on the algebra \mathcal{A} of the entire chain. That is, on any matrix subalgebra $\mathcal{B} \subset \mathcal{A}$, $\mathcal{B} \cong \mathcal{M}_d$, we define $\tau(A) = \frac{1}{d} \text{Tr}(A)$, where Tr is the usual matrix trace on \mathcal{M}_d , which is 1 on minimal projections. The reason for this normalization is that in contrast to the matrix trace, the value of τ does not change if we consider A to be embedded as $A \otimes \mathbb{1}$ in some larger subalgebra $\mathcal{B} \otimes \mathcal{B}_1$, and is hence a well defined state on chain algebra \mathcal{A} . In fact, it is the unique state (normalized positive functional) on \mathcal{A} with the property that $\tau(AB) = \tau(BA)$. This characterization is purely algebraic, which implies that, for every automorphism α of \mathcal{A} and any $A \in \mathcal{A}$, we have $\tau(\alpha(A)) = \tau(A)$.

We now define the scalar product between algebra elements by

$$\langle x|y \rangle = \tau(x^*y). \tag{39}$$

The completion of the algebra \mathcal{A} as a Hilbert space with this scalar product is called the GNS-Hilbert space \mathcal{H}_τ of the tracial state τ . We write $|x\rangle \in \mathcal{H}_\tau$ for the vector obtained by embedding $x \in \mathcal{A}$ in the completion. The trace of operators on \mathcal{H}_τ will be denoted by \mathbf{Tr} , to avoid confusion with the matrix trace Tr of some elements of \mathcal{A} , which is also

used later. Since τ is preserved by any automorphism α , we can define a unitary operator V_α on \mathcal{H}_τ with

$$V_\alpha|x\rangle = |\alpha(x)\rangle. \tag{40}$$

Consider now a finite or infinite dimensional subalgebra $\mathcal{B} \subset \mathcal{A}$. By P we denote the orthogonal projection onto the closure of \mathcal{B} in \mathcal{H}_τ . For finite matrix algebras $\mathcal{B} \cong \mathcal{M}_d$ the matrix units $e_{ij} \cong |i\rangle\langle j|$ clearly form a basis \mathcal{B} , and one readily verifies that $\{\sqrt{d}|e_{ij}\rangle\}_{i,j=1}^d$ is an orthonormal basis of this d^2 -dimensional subspace. It is sometimes also helpful to represent the projection P as an integral over unitaries. That is, for a finite dimensional subalgebra $\mathcal{B} \subset \mathcal{A}$ we have

$$P = d^2 \int dU |U\rangle\langle U|, \tag{41}$$

where the integral is over the unitary group of \mathcal{B} , and dU denotes the normalized Haar measure. For infinite dimensional subalgebras such formulas are not available. We will only need infinite dimensional projections of this type for half chain algebras, which are the closure of an increasing net of finite dimensional subalgebras \mathcal{B}_n . In this case the family of projections P_n associated with the approximating algebras is also increasing, and converges strongly to P .

When $\mathcal{B}_1 \subset \mathcal{A}$ and $\mathcal{B}_2 \subset \mathcal{A}$ are commuting matrix subalgebras, the corresponding matrix units, say $e_{ij}^{(1)}$ and $e_{ab}^{(2)}$, can be taken together as a set of matrix units for $\mathcal{B}_1 \otimes \mathcal{B}_2$, and we get

$$\langle e_{ij}^{(1)} | e_{ab}^{(2)} \rangle = \tau(e_{ji}^{(1)} e_{ab}^{(2)}) = \frac{1}{d_1 d_2} \delta_{ij} \delta_{ab}.$$

Therefore we get

$$P_1 P_2 = |\mathbb{1}\rangle\langle \mathbb{1}|. \tag{42}$$

This equation also holds for infinite pieces of the chain, e.g., a right and a left half chain localized on disjoint subsets of \mathbb{Z} . This is readily seen by approximating each half chain by finite matrix algebras and using the strong convergence of projections. Note that (42) also implies that P_1 and P_2 commute, and the projections $P_i - |\mathbb{1}\rangle\langle \mathbb{1}|$ with the intersection removed are orthogonal.

Overlap of algebras. Of course, if two algebras do not commute, which requires that the localization regions of \mathcal{B}_1 and \mathcal{B}_2 overlap, the geometric position of the subspaces $P_i \mathcal{H}_\tau$ is not so simple. Even if the algebras have trivial intersection $\mathcal{B}_1 \cap \mathcal{B}_2 = \mathbb{C}\mathbb{1}$, so that $(P_1 \mathcal{H}_\tau) \cap (P_2 \mathcal{H}_\tau) = \mathbb{C}|\mathbb{1}\rangle$ the vectors in the remainder can now have angles different from $\pi/2$, and may even approximate each other. This leads to the following definition of a quantitative notion of the overlap of two algebras, which will be the basis of the index formula we develop in this section. We state it together with a few of its basic properties. A matrix algebra here always means a subalgebra, which is isomorphic to \mathcal{M}_d for some $d < \infty$ and contains the identity of \mathcal{A} .

Lemma 12. *For any two subalgebras $\mathcal{B}_1, \mathcal{B}_2 \subset \mathcal{A}$, with corresponding orthogonal projections P_1, P_2 on \mathcal{H}_τ , we define their **overlap** as*

$$\eta(\mathcal{B}_1, \mathcal{B}_2) = \sqrt{\text{Tr}(P_1 P_2)} \in [0, \infty]. \tag{43}$$

Then

1. $\eta(\mathcal{B}_1, \mathcal{B}_2) = 1$, for commuting matrix algebras.
2. $\eta(\mathcal{B}_1, \mathcal{B}_2) \geq 1$, for any two subalgebras (with unit).
3. Let $\mathcal{B}_1, \mathcal{B}_2 \subset \mathcal{B}_{12}$ and $\mathcal{B}_3, \mathcal{B}_4 \subset \mathcal{B}_{34}$ be matrix subalgebras such that \mathcal{B}_{12} and \mathcal{B}_{34} commute. Then

$$\eta(\mathcal{B}_1\mathcal{B}_3, \mathcal{B}_2\mathcal{B}_4) = \eta(\mathcal{B}_1, \mathcal{B}_2)\eta(\mathcal{B}_3, \mathcal{B}_4). \tag{44}$$

4. Let $\mathcal{B}_1, \mathcal{B}_2, \mathcal{B}_3, \mathcal{B}_4$ be a chain of matrix algebras such that $[\mathcal{B}_i, \mathcal{B}_j] = \{0\}$, except for $i, j = 2, 3$. Then $\eta(\mathcal{B}_1\mathcal{B}_2, \mathcal{B}_3\mathcal{B}_4) = \eta(\mathcal{B}_2, \mathcal{B}_3)$.
5. Let $\mathcal{B}_1, \mathcal{B}_2, \mathcal{B}_3$ be commuting matrix algebras with $\mathcal{B}_2 \cong \mathcal{M}_d$. Then $\eta(\mathcal{B}_1\mathcal{B}_2, \mathcal{B}_2\mathcal{B}_3) = d$.
6. Let $\mathcal{B}_1 \cong \mathcal{M}_d$ and \mathcal{B}_2 be finite dimensional matrix algebras, and α an automorphism of the ambient algebra such that $\|\alpha(x) - x\| \leq \varepsilon\|x\|$ for $x \in \mathcal{B}_1$. Then

$$|\eta(\alpha(\mathcal{B}_1), \mathcal{B}_2) - \eta(\mathcal{B}_1, \mathcal{B}_2)| \leq \varepsilon d^2.$$

Proof. Item 1 follows directly from Eq. (42). For Item 2, note that $P_i \geq |\mathbb{1}\rangle\langle\mathbb{1}|$. Since we are only considering subalgebras containing the identity of the ambient algebra, the parenthetical remark is only added for emphasis. Item 3 follows by observing that the normalized trace on $\mathcal{B}_{12}\mathcal{B}_{34} \cong \mathcal{B}_{12} \otimes \mathcal{B}_{34}$ is fixed to be the product of normalized traces on the subalgebras. In this tensor product representation one readily verifies that the projection belonging to $\mathcal{B}_1\mathcal{B}_3$ is the tensor product $P_1 \otimes P_3$, and similarly for P_2, P_4 . Then the formula follows because the trace \mathbf{Tr} also factorizes.

Finally Item 4 follows by taking $\mathcal{B}_1\mathcal{B}_4$ and $\mathcal{B}_2\mathcal{B}_3$ as the pairing of Item 2, and using Item 1 to conclude that $\eta(\mathcal{B}_1, \mathcal{B}_3) = 1$.

To prove the continuity estimate in Item 6, note that overlaps are ≥ 1 , so that by a gradient estimate on the square root function

$$|\eta(\alpha(\mathcal{B}_1), \mathcal{B}_2) - \eta(\mathcal{B}_1, \mathcal{B}_2)| \leq \frac{1}{2} |\mathbf{Tr}((\tilde{P}_1 - P_1)P_2)| \leq \frac{1}{2} \|\tilde{P}_1 - P_1\|_1,$$

where \tilde{P}_1, P_1, P_2 are the projections belonging to $\alpha(\mathcal{B}_1), \mathcal{B}_1, \mathcal{B}_2$, and $\|X\|_1$ denotes the trace norm of X . We use the representation of P_1 in the form (41), so that

$$\tilde{P}_1 - P_1 = d^2 \int dU (|\alpha(U)\rangle\langle\alpha(U)| - |U\rangle\langle U|).$$

We will estimate the trace norm of this expression by estimating the integrand, and using that, for any unit vectors ϕ, ψ in a Hilbert space, $\| |\phi\rangle\langle\phi| - |\psi\rangle\langle\psi| \|_1 \leq 2\|\phi - \psi\|$. Indeed, for unitaries like U and $\alpha(U)$ the vectors $|U\rangle$ and $|\alpha(U)\rangle$ have norm 1 in \mathcal{H}_τ . Moreover, $\|\alpha(U) - U\|^2 = \tau((\alpha(U) - U)^*(\alpha(U) - U)) \leq \varepsilon^2\|U\|\tau(\mathbb{1}) = \varepsilon^2$. Hence

$$\frac{1}{2} \|\tilde{P}_1 - P_1\|_1 \leq \frac{2d^2}{2} \varepsilon \int dU = d^2\varepsilon.$$

□

Index formula. With these preparations we can state the main result of this section:

Proposition 13. *Let $\mathcal{A}_< = \mathcal{A}_{(-\infty, 0]}$ and $\mathcal{A}_> = \mathcal{A}_{[1, \infty)}$, and $\mathcal{A}_L, \mathcal{A}_R$ be the algebras of any two neighboring cell algebras for a nearest neighbor grouping of the chain. Then*

$$\text{ind}(\alpha) = \frac{\eta(\alpha(\mathcal{A}_<), \mathcal{A}_>)}{\eta(\alpha(\mathcal{A}_>), \mathcal{A}_<)} = \frac{\eta(\alpha(\mathcal{A}_L), \mathcal{A}_R)}{\eta(\alpha(\mathcal{A}_R), \mathcal{A}_L)}. \tag{45}$$

Moreover, if $t \mapsto \alpha_t$ is a strongly continuous family of cellular automata with the same cell structure and neighborhood scheme, then $\text{ind}(\alpha_t)$ is constant.

Proof. The square of the numerator of the second index expression is $\text{Tr} V_\alpha P_L V_\alpha^* P_R$, and we first verify that this expression is unchanged if we choose a larger L and R , say $L' = L_1 \cup L$ and $R' = R \cup R_1$. Indeed, the algebras $\alpha(\mathcal{A}_{L_1}), \alpha(\mathcal{A}_L), \mathcal{A}_R, \mathcal{A}_{R_1}$ satisfy the conditions of Lemma 12, Item 4. Moreover, arguing as for formula (42) we see that not only the trace, but even the operator is independent of an enlargement. Hence taking a strong limit we obtain the corresponding expression for the infinite half-chains.

From Lemma 12, Item 3, it is clear that the expression thus defined satisfies the tensor product property, in fact, numerator and denominator do so independently. Moreover, the formula is valid for shift automorphisms by virtue of Item 5 of the lemma.

Now let σ be a tensor product of shifts with $\text{ind} \sigma = \text{ind} \alpha$, and let $\text{ind}' \alpha$ be the value the formula in the proposition gives for any automorphism α . Then we have $\text{ind}'(\alpha \otimes \sigma^{-1}) = \text{ind}'(\alpha)(\text{ind} \sigma)^{-1} = \text{ind}' \alpha / (\text{ind} \alpha)$. So it remains to prove that $\text{ind}' \alpha = 1$ for every α with $\text{ind} \alpha = 1$, which by Thm. 9 means that α is implemented by two layers of block unitaries. Moreover, we can forget all unitaries acting only on one side of the separation, since they do not change the algebras. Only one unitary $U_{LR} \in \mathcal{A}_L \otimes \mathcal{A}_R$ connecting the cells immediately adjacent to the cut remains. Hence it only remains to prove the lemma below.

For the continuity statement observe that we can make a nearest neighbor grouping jointly for all α_t , so we can apply Lemma 12, Item 6, to see that both denominator and numerator of the second fraction in (45) are continuous. Moreover, the denominator is ≥ 1 , so $\text{ind}(\alpha_t)$ is a continuous function of t . On the other hand, numerator and denominator of $\text{ind} \alpha$ have to divide every cell dimension, so there is only a finite choice of possible values, given the neighborhood structure. Hence $\text{ind}(\alpha_t)$ is constant. \square

In order to state the remaining lemma, let us consider any automorphism α of a tensor product $\mathcal{A}_L \otimes \mathcal{A}_R$ of two finite dimensional matrix algebras. We can take $\mathcal{A}_x = \mathcal{B}(\mathcal{H}_x)$, for some d_x -dimensional Hilbert spaces \mathcal{H}_x . In that case, the automorphism α is implemented by conjugation with a unitary: $\alpha(x) = UxU^*$, for some $U \in U(\mathcal{H}_L \otimes \mathcal{H}_R)$. We will express the fraction appearing in (45) in terms of U .

In the following calculation, it turns out to be convenient to introduce orthonormal bases in the Hilbert spaces involved. Denote these bases by $\{|i\rangle\} \subset \mathcal{H}_L$ and $\{|a\rangle\} \subset \mathcal{H}_R$ respectively. These bases allow us to define the notion of a *partial transpose* U^Γ of U . We set

$$\langle ia|(U^\Gamma)|jb\rangle = \langle ja|U|ib\rangle.$$

While the partial transpose depends on the basis used to define it, one may easily convince oneself that the expression appearing in next lemma is independent of that choice.

Lemma 14. *When $\alpha(x) = UxU^*$ is an automorphism of the tensor product $\mathcal{A}_L \otimes \mathcal{A}_R$ of finite dimensional matrix algebras, we have*

$$\eta(\alpha(\mathcal{A}_L), \mathcal{A}_R)^2 = \frac{1}{d_A d_B} \text{Tr} \left((U^\Gamma U^{\Gamma*})^2 \right). \tag{46}$$

Moreover, this expression is invariant under the substitution $\mathcal{A}_L \leftrightarrow \mathcal{A}_R$.

Proof. We introduce matrix units $e_{ij} = |i\rangle\langle j|$ as above and use that the $\sqrt{d} e_{ij}$ form an orthonormal basis of \mathcal{A}_L . Thus the projection P_L on \mathcal{H}_τ corresponding to this algebra is $P_L = d_L \sum_{ij} |e_{ij}\rangle\langle e_{ij}|$. Of course, the transformed projection \tilde{P}_L for $\alpha(\mathcal{A}_L)$ is obtained by substituting $|\alpha(e_{ij})\rangle$ for $|e_{ij}\rangle$. Similarly, we set $f_{ab} = |a\rangle\langle b|$ and, accordingly, $P_R = d_R \sum_{ab} |f_{ab}\rangle\langle f_{ab}|$. Then the left hand side of (46) becomes

$$\eta(\alpha(\mathcal{A}_L), \mathcal{A}_R)^2 = \text{Tr}(\tilde{P}_L P_R) = d_L d_R \sum_{ijab} |\langle f_{ab} | \alpha(e_{ij}) \rangle|^2,$$

with the scalar products

$$\begin{aligned} \langle f_{ab} | \alpha(e_{ij}) \rangle &= \sum_{kc} \tau(e_{kk} f_{ba} U e_{ij} f_{cc} U^*) \\ &= \frac{1}{d_L d_R} \sum_{kc} \langle ka | U | ic \rangle \langle jc | U^* | kb \rangle \\ &= \frac{1}{d_L d_R} \sum_{kc} \langle ia | U^\Gamma | kc \rangle \langle kc | U^{\Gamma*} | jb \rangle \\ &= \frac{1}{d_L d_R} \langle ia | U^\Gamma U^{\Gamma*} | jb \rangle. \end{aligned} \tag{47}$$

Altogether we get Eq. (46).

The second claim becomes a simple corollary: we must verify that the right-hand side is unchanged by the substitution $U \mapsto U^*$. Since the adjoint operation commutes with the partial transpose, this amounts to a cyclic rearrangement under the trace. \square

8. Outlook

Two directions of generalization of the theory presented here are especially suggestive and are, in fact, the subject of ongoing work in our group. We briefly comment on the prospects.

8.1. Approximate causality. In many situations in physics causality is only approximately satisfied, e.g., as a bound $\|U_{xy}\| \leq f(|x - y|)$ for some function f going to zero at infinity. For example, the unitary groups generated by finite range Hamiltonians would satisfy this, but are never strictly causal in the sense required here. For the Key Lemma 4 powerful generalizations exist [3]. From these it is clear that the index of approximately causal unitaries is definable, integer valued, satisfies a product formula, and is zero for the unitaries arising from Hamiltonian subgroups. The part of the theory which is less clear is related to the converses, namely the construction of strictly causal unitary walks with the same index, approximating a given approximately causal unitary.

In the case of cellular automata the right notions of approximate causality are not clear. Ideally, one would only demand that α be an automorphism of the quasi-local algebra [5]. By definition, this means that the image of any localized element can be approximated in norm by localized ones. The idea of support algebras (which largely relies in its present form on picking a finite basis) is certainly too simplistic, and in any case it is unlikely that such algebras will always turn out to be finite dimensional matrix algebras. One may speculate whether the index should take on also irrational values, but this seems unlikely, because of its dependence on the cell structure: for chains of homogeneous cell dimension 2, the index is always a power of 2, and not a dense set of rationals.

8.2. Higher lattice dimension. As we have shown, in lattice dimension 1 three possible classifications of walks and automata coincide: (1) the classification by locally computable invariants, (2) the classification modulo locally implementable, unitaries, and (3) the classification up to homotopy. It is very unlikely that these three coincide in higher dimension as well. In the one-dimensional theory we could allow the local systems to grow, but also the localization regions. In the higher dimensional case we will use a translation invariant metric to bound the neighborhood sizes of a “local” system. Generalizations can be built on coarse geometry [14].

Locally computable invariants are probably trivial. For example, an arbitrarily large patch of the shift automorphism can be connected to the identity outside a finite enlargement of the patch. In this sense the shifts have the same invariants as the identity. To get a more interesting theory, one should take other regions for the definition of “locally” computable, e.g., computability on cones [6], or computability *outside* of any arbitrarily large region.

The classification modulo local implementability is especially interesting from the physical point of view, but it might turn out to be rather wild. For example, consider some self-intersection free path in the lattice, which comes from infinity and goes to infinity. Fix walks/cellular automata, which are equal to the identity off this path and allow any of the one-dimensional systems along this path. As long as we fix the path, we can apply the one dimensional theory. For two paths, which keep a finite distance from each other, it is easy to envisage local swap-type unitaries taking one path to the other, which would bring the corresponding path-related indices under the same roof. However, if the paths move away from each other, there will be no such local operation connecting them, so systems with non-trivial indices along these paths fall into different equivalence classes modulo local implementability. However, the equivalence classes of paths modulo “keeping a finite distance” are a rather unmanageably large set. A useful classification cannot be expected. Incidentally, the same class of examples shows that the “invariants computable outside any finite region” will give a wild set.

For the homotopy classification of walks there is already a theory, based on the K -theory of C^* -algebras and its connection with coarse geometry. Indeed, the K_1 group of a C^* -algebra just classifies the connected components of its automorphism group. This theory will most naturally apply to approximately causal walks, since strict causality cannot even be stated simply in terms of the C^* -algebra of the whole system. The connection with coarse geometry is being explored, for example, by Ralf Meyer and his group in Göttingen. Surprisingly (to us) it turns out that using Bott periodicity one can see that the K_1 group of approximately causal walks alternates between \mathbb{Z} in odd dimension and 0 in even dimension. Unfortunately, this theory does not apply readily to the case of cellular automata.

8.3. Higher dimensional translation invariant systems. In order to tame the wildness indicated in the previous subsection, one can restrict attention to translationally invariant systems. Immediately, the index of walks gets an obvious definition. The Fourier transform \widehat{U} of the walk is now a Laurent polynomial in the variables $\exp(ip_k)$, where (p_1, \dots, p_s) is the momentum vector. Then $\det \widehat{U}$ is an invertible polynomial and we again conclude $\det \widehat{U}(p) = \exp(i \sum_k n_k p_k)$ for some integers n_k . The lattice vector (n_1, \dots, n_s) can be called the index, in direct generalization of the one-dimensional case.

For cellular automata, which are products of partial shifts and local block unitaries, we can just define the index via the shift content contained in such a representation, obtaining some vector with rational components. Since it is not known, whether any QCA is of this form, this is very unsatisfactory. As a step in the right direction, one can define an index by reduction to the 1D case, without using any special decomposition: Suppose we choose vectors a_1, \dots, a_{n-1} in an n -dimensional lattice, and we identify sites differing by integer multiples of these vectors. Then if the a_k are large enough with respect to the interaction length of the automaton, this gives a well-defined cellular automaton evolution $\alpha[a_1, \dots, a_{n-1}]$ on the quotient lattice, which now has only one unbounded direction, so we can assign an index to it. Then we call a vector $q \in \mathbb{R}^n$ the index of α , if

$$\text{ind } \alpha[a_1, \dots, a_{n-1}] = \det[q, a_1, \dots, a_{n-1}], \quad (48)$$

where on the right-hand side the square bracket denotes the matrix with the specified column vectors. It can be verified easily that this gives the result indicated before, for any cellular automaton, which has a decomposition into partial shifts and partitioned unitaries. But in general it is not even clear that the left hand side must depend linearly on each a_k .

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