Libckpt: Transparent Checkpointing under Unix

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OUTLINE

- Introduction
- Transparent Checkpointing
- User-Directed Checkpointing
- The Mechanics of Checkpointing and Recovery
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Introduction

- Checkpointing is a simple technique for rollback recovery.
- User level
- Transparent Sequential and incremental checkpointing
- User-Directed checkpointing
 little information yield large improvements.

Transparent Checkpointing

- Sequential Checkpointing
- Not completely transparent : change main() to ckpt_target()
- Generates a timer interrupt, and take a sequential checkpoint at each interrupt
- .ckptrc file :
 - checkpointing<on/off>
 - dir <directory>
 - Maxtime <sec>

Transparent Checkpointing

- Incremental Checkpointing
- incremental <on/off>
- old sequential checkpoint file can be deleted.
 but incremental one can't.
- Maxfiles <n>
- use page protection to identify which pages should be save.

Transparent Checkpointing

- Forked Checkpointing
- Main-memory checkpointing
- Saving of the checkpoint to disk is overlapped with application execution.
- fork <on/off>
- main-memory or copy-on-write checkpointing

User-Directed Checkpointing

- Memory exclusion
 - Locations is dead
 - Locations is clean
- Heap variables and variables which reside in statically allocated data segment don't work.
- exclude_bytes(char *addr, int size, int usage)
 include_bytes(char *addr, int size)
- Usage : CKPT_READONLY \ CKPT_DEAD

User-Directed Checkpointing

- Synchronous Checkpointing
- specify points in the program where it is most advantageous for checkpoint
- checkpoint_here()
- mintime <sec>
- asynchronous maxtime synchronous – mintime

User-Directed Checkpointing

```
ckpt target()
main()
                                       struct data *D;
  struct data *D:
                                       FILE *fi, *fo;
 FILE *fi, *fo;
                                       D = allocate data set();
 D = allocate data set();
                                       fi = fopen("input", "r");
  fi = fopen("input", "r");
                                       fo = fopen("output", "w");
  fo = fopen("output", "w");
                                       while(read data(fi, D) != -1) {
  while(read_data(fi, D) != -1) {
                                         perform calculation(D);
    perform calculation(D);
                                         output results(fo, D);
    output_results(fo, D);
                                         exclude_bytes(D, sizeof(struct data),
                                              CKPT DEAD);
                                         checkpoint_here();
                                         include_bytes(D, sizeof(struct data))
```

The Mechanics of Checkpointing and Recovery

- Process creation :
 - auto restore text portion of the process's state, and begins execution.
- Recovery routine :
 - process's stack + data segments
- System state restoration :
 - Save the state of open files.
- Processor state restoration :
 - Program counter stack pointer

The Mechanics of Checkpointing and Recovery

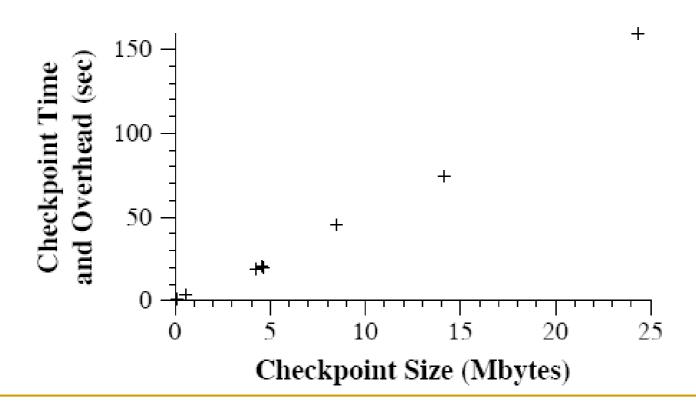
- Save processor state setjmp()
- Record the state of the open file table
- Data state (program's stack, data segment)

SunOS

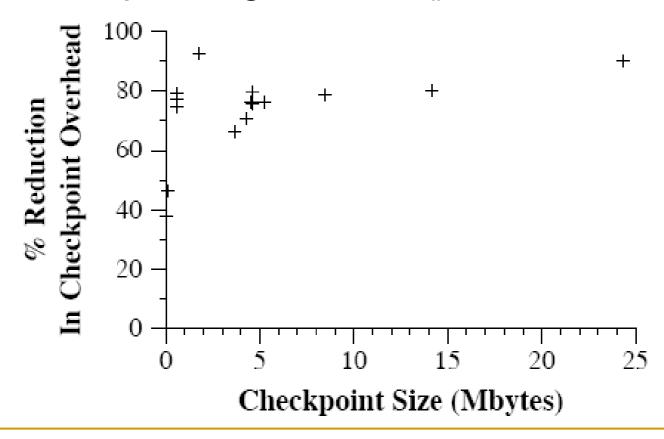
Application	Abbreviation	Language	Running	Maximum	Checkpoint
			Time	Checkpoint	Interval
			(mm:ss)	Size (Mbytes)	(min)
Matrix Multiplication	MAT	C	15:20	4.6	2
Linear Equation Solver	SOLVE	FORTRAN	13:42	4.6	2
Cellular Automata	CELL	C	17:39	8.4	2
Shallow Water Model	WATER	FORTRAN	25:54	13.1	3
Multicommodity Flow	MCNF	FORTRAN	18:38	24.3	6

Table 1: Description of application instances

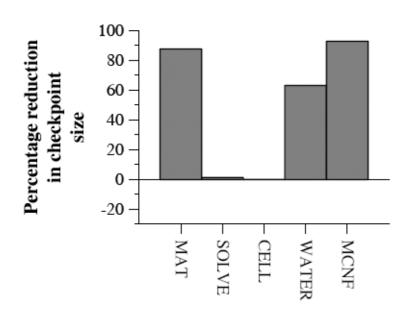
Sequential Checkpointing

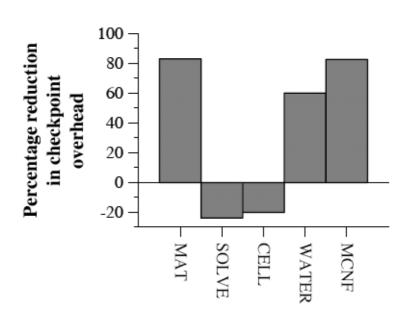


Checkpointing with fork()



Incremental Checkpointing





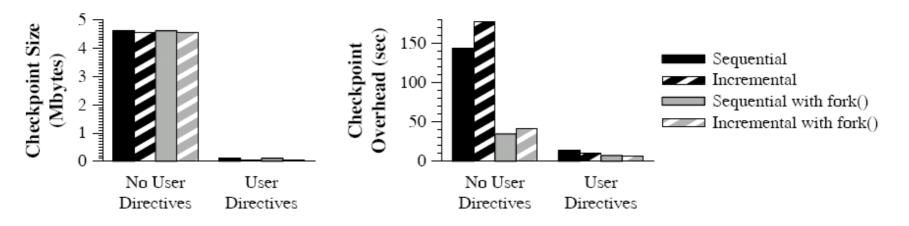


Figure 6: Results of User-Directed Checkpointing on the SOLVE Application

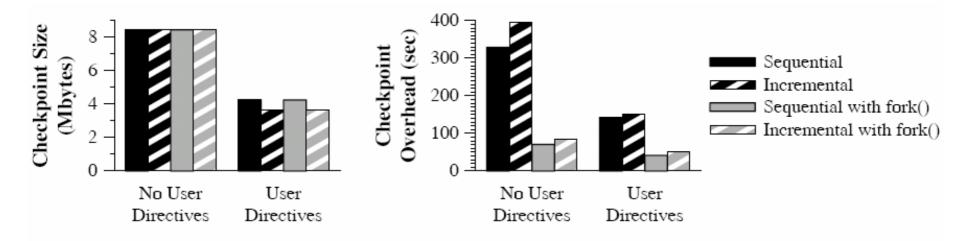


Figure 7: Results of User-Directed Checkpointing on the CELL Application

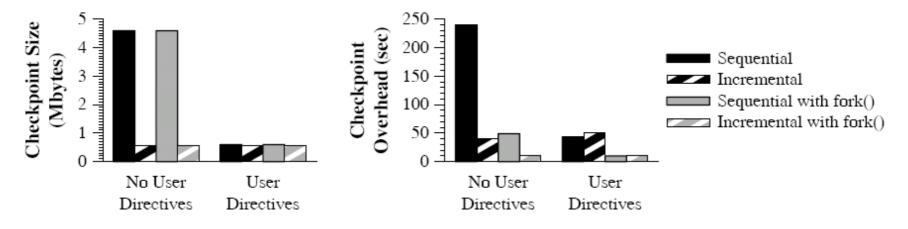


Figure 8: Results of User-Directed Checkpointing on the MAT Application

Conclusion

Ease of use and low overhead

 Future research : employ compiler analysis to assist user-directed checkpointing