# Low Rank Co-Diagonal Matrices and Ramsey Graphs

Vince Grolmusz

Department of Computer Science Eötvös University, H-1053 Budapest HUNGARY E-mail: grolmusz@cs.elte.hu Submitted: March 7, 2000. Accepted: March 12, 2000

#### Abstract

We examine  $n \times n$  matrices over  $Z_m$ , with 0's in the diagonal and nonzeros elsewhere. If m is a prime, then such matrices have large rank (i.e.,  $n^{1/(p-1)} - O(1)$ ). If m is a non-prime-power integer, then we show that their rank can be much smaller. For m = 6 we construct a matrix of rank  $\exp(c\sqrt{\log n \log \log n})$ . We also show, that explicit constructions of such low rank matrices imply explicit constructions of Ramsey graphs.

Keywords: composite modulus, explicit Ramsey-graph constructions, matrices over rings, co-diagonal matrices

## 1 Introduction

In this work we examine matrices over a ring R, such that the diagonal elements of the matrix are all 0's, but the elements off the diagonal are not zero (we shall call these matrices co-diagonal over R). We define the rank of a matrix over a ring, and show that low rank co-diagonal matrices over  $Z_6$  naturally correspond to graphs with small homogenous vertex sets (i.e., cliques and anti-cliques). Consequently, explicitly constructible low rank co-diagonal matrices over  $Z_6$  imply explicit Ramsey graph constructions. Our best construction reproduces the logarithmic order of magnitude of the Ramsey-graph of *Frankl* and *Wilson* [5], continuing the sequence of results on new explicit Ramsey graph constructions of Alon [1] and Grolmusz [6]. Our present result, analogously to the constructions of [6] and [1], can be generalized to more than one color.

Our results give a recipe for constructing explicit Ramsey graphs from explicit low rank co-diagonal matrices over  $Z_6$ , analogously to the way that our results gave a method for constructing explicit Ramsey graphs from certain low degree polynomials over  $Z_6$  in [6]. In this sense, our results may lead to improved Ramsey graph constructions, if lower rank co-diagonal matrix constructions exist.

**Definition 1** Let R be a ring and let n be a positive integer. We say, that  $n \times n$  matrix  $A = \{a_{ij}\}$  is a co-diagonal matrix over R, if  $a_{ij} \in R$ , i, j = 1, 2, ..., n and  $a_{ii} = 0, a_{ij} \neq 0$ , for all i, j = 1, 2, ..., n,  $i \neq j$ .

We say, that A is an upper co-triangle matrix over R, if  $a_{ij} \in R$ , i, j = 1, 2, ..., nand  $a_{ii} = 0, a_{ij} \neq 0$ , for all  $1 \leq i < j \leq n$ . A is a lower co-triangle matrix over R, if  $a_{ij} \in R$ , i, j = 1, 2, ..., n and  $a_{ii} = 0, a_{ij} \neq 0$ , for all  $1 \leq j < i \leq n$ . A matrix is co-triangle, if it is either lower- or upper co-triangle.

We will also need the definition of the rank of a matrix with elements in a ring. The following definition is a generalization of the matrix rank over fields to matrices over rings:

**Definition 2** Let R be a ring and let n be a positive integer. We say, that  $n \times n$  matrix A over R has rank 0 if all of the elements of A are 0. Otherwise, the rank over the ring R of matrix A is the smallest r, such that A can be written as

$$A = BC$$

over R, where B is an  $n \times r$  and C is an  $r \times n$  matrix. The rank of A over R is denoted by rank<sub>R</sub>(A).

It is easy to see, that this definition of the matrix rank coincides with the usual matrixrank over R, when R is a field. The following property of the usual matrix rank also holds:

**Lemma 3** Let R be a ring and let A and A' be two  $n \times n$  matrices. Then  $\operatorname{rank}_R(A + A') \leq \operatorname{rank}_R(A) + \operatorname{rank}_R(A')$ .

**Proof:** Let A = BC and A' = B'C', where B is an  $n \times r$  and C is an  $r \times n$  matrix, while B' is an  $n \times r'$  and C' is an  $r' \times n$  matrix. Then A + A' can be given as B''C'', where B'' is an  $n \times (r + r')$  matrix, formed from the union of the columns of B and B', and C'' is an  $(r + r') \times n$  matrix, formed from the union of rows of C and C'.  $\Box$ 

The following theorem shows, that for any prime p, the co-triangle (and, consequently, the co-diagonal) matrices over the p-element field have large rank:

**Theorem 4** Let p be a prime, and let A be an  $n \times n$  co-triangle matrix over  $GF_p$ . Then

$$\operatorname{rank}_{\operatorname{GF}_p}(A) \ge n^{1/(p-1)} - p.$$

**Proof:** We may assume that A is a lower co-triangle matrix. Let  $r = \operatorname{rank}_{\operatorname{GF}_p}(A)$ , and let  $B = \{b_{ij}\}$  be an  $n \times r$ ,  $C = \{c_{ij}\}$  be an  $r \times n$  matrix over  $\operatorname{GF}_p$ , such that:

$$A = BC. (1)$$

For i = 1, 2, ..., n let us consider the following polynomials:

$$P_i(x_1, x_2, \dots, x_r) = \sum_{k=1}^r b_{ik} x_j.$$
 (2)

From (1),

$$P_i(c_{1j}, c_{2j}, \dots, c_{rj}) = \begin{cases} 0, \text{ if } i = j, \\ \neq 0, \text{ if } i > j. \end{cases}$$

Consequently, by the triangle criterion [2], polynomials

$$Q_i(x_1, x_2, \dots, x_r) = 1 - P_i^{p-1}(x_1, x_2, \dots, x_r),$$

for i = 1, 2, ..., n, form a linearly independent set in the vector space of dimension

$$\binom{r+p-2}{p-1} + 1$$

of polynomials of form  $Q + \alpha$ , where Q is an r-variable homogeneous polynomial of degree p-1 and  $\alpha \in GF_p$ . (To prove this without the triangle criterion of [2], one should observe that  $Q_k$  is zero on column *i* of matrix C for i < k, and it is 1 for column *k* of C; so  $Q_i$  cannot be given as a linear combination of some  $Q_{k_i}$ 's, each  $k_j > i$ .) Consequently,

$$n \le \binom{r+p-2}{p-1} + 1 \le (r+p)^{p-1}.$$
(3).

We are interested in the following question:

Question. Let  $R = Z_m$ , what is the minimum rank of an  $n \times n$  co-triangle (or co-diagonal) matrix over R?

If m = p a prime, then by Theorem 4 we have that the rank should be at least  $n^{1/p-1} - p$ . What can we say for non-prime m's?

The main motivation of this question is the following theorem:

**Theorem 5** Let  $A = \{a_{ij}\}$  be an  $n \times n$  co-triangle matrix over  $R = Z_6$ , with  $r = \operatorname{rank}_{Z_6}(A)$ . Then there exists an n-vertex graph G, containing neither a clique of size r + 2 nor an anti-clique of size

$$\binom{r+1}{2} + 2.$$

**Proof:** Suppose, that A is a lower co-triangle matrix. If the  $Z_6$  rank of A is r, then both the GF<sub>2</sub> and GF<sub>3</sub> ranks of A are at most r. Let  $V = \{v_1, v_2, \ldots, v_n\}$ . For any i > j, let us connect  $v_i$  and  $v_j$  with an edge, if  $a_{ij}$  is odd. Then any clique of size t will correspond to a  $t \times t$  lower co-triangle minor over GF<sub>2</sub>, so from (3),

$$t \le r+1.$$

Any anti-clique of size t will correspond to a  $t \times t$  lower co-triangle minor over  $GF_3$ , so from (3),

$$t \le \binom{r+1}{2} + 1. \tag{4}$$

From Theorem 5 one can get a lower bound for the rank, using estimations for the Ramsey numbers. Our original bound was significantly improved by *Noga Alon*, who allowed us to include his proof here.

**Theorem 6** Let  $A = \{a_{ij}\}$  be an  $n \times n$  co-triangle matrix over  $R = Z_6$ . Then

$$\operatorname{rank}_{Z_6}(A) \ge \frac{\log n}{2\log\log n} - 2.$$

**Proof:** By the result of *Ramsey* [7] and *Erdős* and *Szekeres* [4], every *n*-vertex graph has either a clique on k, or an anti-clique on  $\ell$  vertices, if

$$n \ge \binom{k+\ell-2}{k-1}$$

If we set  $k = \lfloor \frac{1}{2} \frac{\log n}{\log \log n} \rfloor$ , and  $\ell = \lfloor \log^2 n \rfloor$ , then we get from Theorem 5, that both  $r+2 \leq k$  and  $\binom{r+1}{2} + 2 \leq \ell$  cannot be satisfied, and this completes the proof.  $\Box$ 

The proof of Theorem 5 also proves

**Theorem 7** Suppose, that there exists an explicitly constructible  $n \times n$  co-triangle matrix  $A = \{a_{ij}\}$  over  $R = Z_6$ , with  $r = \operatorname{rank}_{Z_6}(A)$ . Then one can explicitly construct an n-vertex Ramsey-graph, without homogenous vertex-sets of size

$$\binom{r+1}{2} + 2.$$

Our main result is that there do exist explicitly constructible low-rank co-diagonal matrices over  $Z_6$ , implying explicit Ramsey-graph constructions.

**Theorem 8** There exists a c > 0 such that for all positive integer n, there exists an explicitly constructible  $n \times n$  co-diagonal matrix  $A = \{a_{ij}\}$  over  $R = Z_6$ , with

$$\operatorname{rank}_{Z_6}(A) \le 2^{c\sqrt{\log n \log \log n}}$$

Theorem 8 together with Theorem 5, gives an explicit Ramsey-graph construction on n vertices, without a homogeneous vertex-set of size  $2^{c'\sqrt{\log n \log \log n}}$ , for some c' > 0, or in other words, an explicit Ramsey-graph construction on

$$2^{\frac{c''\log^2 t}{\log\log t}}$$

vertices, without homogeneous vertex-set of size t, for some c'' > 0. This bound was first proven by *Frankl* and *Wilson* [5] with a larger (better) constant than our c'', using the famous Frankl-Wilson theorem [5]. We also gave a construction, using the BBR polynomial [3] and also the Frankl-Wilson theorem in [6].

A generalization of our main result for ring  $Z_m$ , where *m* has more than two prime divisors:

**Theorem 9** For any  $m = p_1^{\alpha_1} p_2^{\alpha_2} \dots p_{\ell}^{\alpha_{\ell}}$ , where the  $p_i$ 's are distinct primes, there exists a  $c = c_m > 0$  such that for all positive integer n, there exists an explicitly constructible  $n \times n$  co-diagonal matrix  $A = \{a_{ij}\}$  over  $R = Z_m$ , with

$$\operatorname{rank}_{Z_m}(A) \le 2^{c\sqrt{\log n(\log \log n)^{\ell-1}}}.$$

# 2 Constructing Low Rank mod 6 Co-Diagonal Matrices

In this section we prove Theorems 8 and 9.

Our main tool is the following theorem (choosing m = 6 and  $\ell = 2$ ):

**Theorem 10 (Barrington, Beigel, Rudich[3])** Given  $m = p_1^{\alpha_1} p_2^{\alpha_2} \dots p_{\ell}^{\alpha_{\ell}}$  where the  $p_i$  are distinct primes, then there exists an explicitly constructible multi-linear polynomial P with integer coefficients, with k variables, and of degree  $O(k^{1/\ell})$  which satisfies for  $x \in \{0,1\}^k$ , that P(x) = 0 over  $Z_m$  iff  $x = (0, 0, \dots, 0)$ .

Let k be the smallest integer such that  $n \leq k^k$ . Let  $B = \{0, 1, 2, \dots, k-1\}$ . Let us define  $\delta : B \times B \to \{0, 1\}$  as follows:

$$\delta(u, v) = \begin{cases} 1, & \text{if } u = v, \\ 0 & \text{otherwise.} \end{cases}$$

Then matrix  $\overline{A}$  is defined as follows: both the rows and the columns of  $\overline{A}$  correspond to the elements of the set  $B^k$ . The entry of matrix  $\overline{A}$  in the intersection of a row, corresponding to  $u = (u_1, u_2, \ldots, u_k) \in B^k$  and of a column, corresponding to  $v = (v_1, v_2, \ldots, v_k) \in B^k$  is the number:

$$P(1 - \delta(u_1, v_1), 1 - \delta(u_2, v_2), \dots, 1 - \delta(u_k, v_k)).$$
(5)

If u = v, then all of the  $\delta(u_i, v_i)$ 's are 1, so the value of P is 0. So the diagonal of  $\overline{A}$  is all-0, but no other elements of the matrix are 0 over  $Z_6$ , consequently,  $\overline{A}$  is co-diagonal over  $Z_6$ .

Multi-linear polynomial P has degree  $O(\sqrt{k})$ , so (5) can be written as the sum of

$$\binom{k}{\leq c \lfloor \sqrt{k} \rfloor} = \sum_{i=0}^{c \lfloor \sqrt{k} \rfloor} \binom{k}{i} < k^{c\sqrt{k}}$$
(6)

monomials of the form:

$$a_{i1,i2,\ldots,is}\delta(u_{i1},v_{i1})\delta(u_{i2},v_{i2}),\ldots\delta(u_{is},v_{is}),$$
(7)

where c is positive, (in fact, c < 3 is also satisfied),  $a_{i1,i2,...,is}$  is an integer between 0 and 5, and  $s \leq c\sqrt{k}$ .

Since the (u, v) entry of A is the value (5), and (5) can be written as the sum of monomials in (7), matrix  $\overline{A}$  can be written as the sum of matrices  $D_{i1,i2,\ldots,is}$ , where the entry of matrix  $D_{i1,i2,\ldots,is}$  in the intersection of a row, corresponding to  $u = (u_1, u_2, \ldots, u_k) \in B^k$  and of a column, corresponding to  $v = (v_1, v_2, \ldots, v_k) \in B^k$  is equal to the value of (7).

It is easy to verify that  $D_{i1,i2,...,is}$  can be written into the following form (applying the same, suitable permutation to the rows and columns):

Let us observe, that the number of all-1 square minors, covering the diagonal is  $k^s$ . Then, from Lemma 3 the rank of  $D_{i1,i2,...,is}$  is  $k^s$ ,  $s \leq c\sqrt{k}$ . It follows from this and from (6), that the rank of  $\overline{A}$  is at most  $k^{2c\sqrt{k}}$ .

Let matrix A be defined as the  $n \times n$  upper left minor of matrix  $\overline{A}$ . Obviously, A is also a co-diagonal matrix, and its rank is at most  $k^{2c\sqrt{k}}$ . Due to the choice of k the statement follows.  $\Box$ 

The proof of Theorem 9 follows the same steps as the proof of Theorem 8. If m has  $\ell$  prime divisors, then polynomial P has degree  $O(k^{1/\ell})$ , so matrix  $D_{i1,i2,...,is}$  has rank at most  $k^s$ ,  $s \leq ck^{1/\ell}$ , and co-diagonal matrix A has rank at most  $k^{2ck^{1/\ell}}$ , and this proves the theorem.

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