

# Multi-Antenna Communication in Ad Hoc Networks: Achieving MIMO Gains with SIMO Transmission

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**Abstract**—The benefit of multi-antenna receivers is investigated in wireless ad hoc networks, and the main finding is that network throughput can be made to scale linearly with the number of receive antennas  $n_R$  even if each transmitting node uses only a single antenna. This is in contrast to a large body of prior work in single-user, multiuser, and ad hoc wireless networks that have shown linear scaling is achievable when multiple receive and transmit antennas (i.e., MIMO transmission) are employed, but that throughput increases logarithmically or sublinearly with  $n_R$  when only a single transmit antenna (i.e., SIMO transmission) is used. The linear gain is achieved by using the receive degrees of freedom to simultaneously suppress interference and increase the power of the desired signal, and exploiting the subsequent performance benefit to increase the density of simultaneous transmissions instead of the transmission rate. This result is proven in the transmission capacity framework, which presumes single-hop transmissions in the presence of randomly located interferers, but it is also illustrated that the result holds under several relaxations of the model, including imperfect channel knowledge, multihop transmission, and regular networks (i.e., interferers are deterministically located on a grid).

**Index Terms**—Multiantenna communication, MIMO, ad-hoc networks, fading, transmission capacity.

## I. INTRODUCTION

MULTIPLE antenna communication has become a key component of virtually every contemporary high-rate wireless standard (LTE, 802.11n, WiMAX). The theoretical result that sparked the intense academic and industrial investigation of MIMO (multiple-input/multiple-output) communication was the finding that the achievable throughput of a point-to-point MIMO channel scales *linearly* with the minimum of the number of transmit and receive antennas [2], [3]. Linear scaling in the number of transmit antennas can also be achieved in point-to-multipoint (broadcast/downlink) channels [4] even if each receiver has only a single antenna,

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or in the number of receive antennas in multipoint-to-point (multiple access/uplink) channels even if each transmitter has a single antenna [3]. In these two cases, the linear gains are enabled by the simultaneous transmission and reception of multiple data streams.

### A. Overview of Main Results

In this paper we are interested in the throughput gains that multiple antennas can provide in *ad hoc networks*, rather than in channels with a common transmitter and/or receiver. If multiple antennas are added at every node in the network and point-to-point MIMO techniques are used to increase the rate of every individual link (i.e., every hop in a multi-hop route) in the network, then network-wide throughput (i.e., the sum of data rates across all links in the network) naturally increases linearly with the number of antennas per node. Similarly, based on the quoted point-to-multipoint and multipoint-to-point results, linear scaling is also expected to occur if nodes are capable of sending or receiving multiple streams.

The main finding of this paper is that network-wide throughput can be increased linearly with the number of receive antennas, even if only a single transmit antenna is used by each node (i.e., single-input, multiple-output, or SIMO, communication), and each node sends/receives only a single data stream. Furthermore, this gain is achievable using only linear receive processing and does not require any transmit-side channel state information (CSIT). This throughput gain is achieved by linearly increasing (with the number of receive antennas) the density of simultaneous transmissions while keeping the per-link rate fixed, and thus requires additional nodes to communicate as the number of receive antennas is increased.

The main result is obtained by considering an ad hoc network in which transmitters are randomly located on the plane according to a 2-D homogeneous Poisson point process with a particular spatial density. We consider a desired transmit-receive pair separated by a fixed distance, and experiencing interference (assumed to be treated as noise) from all other active transmitters. The received signal and interference are functions of path-loss attenuation and fading, and we assume that for a transmission to be successful, it must be detected with an SINR larger than a defined threshold  $\beta$ . The primary performance metric is the maximum spatial density of transmitters/interferers that can be supported such that the outage probability  $\mathbb{P}[\text{SINR} < \beta]$  is no larger than an outage constraint

$\epsilon$ , and our particular interest is in quantifying how quickly this maximum density increases with the number of receive antennas  $n_R$  (with  $\beta$  and  $\epsilon$  fixed).

If receive beamforming is performed, the receive degrees of freedom can be used to either increase the power of the desired signal (i.e., for array gain) or suppress interference, and these competing objectives are optimally balanced by the SINR-maximizing MMSE (minimum mean-square error) filter. Using the MMSE outage probability lower bound in [5], we develop an upper bound on the maximum density allowable with an MMSE receiver. In conjunction, we develop a density lower bound by analyzing the performance of a novel suboptimal *partial zero forcing* (PZF) receiver, which uses an explicit fraction of the degrees of freedom for array gain and the remainder for interference cancellation. By showing that both the lower and upper bound are linear in  $n_R$ , we can conclude that the maximum transmit density is  $\Theta(n_R)$  – this result applies only to linear techniques, and thus non-linear techniques such as successive interference cancellation could conceivably improve upon this – and we demonstrate that this allows well known metrics like the transmission capacity [6], [7], transport capacity [8], [9], and the expected forward progress [10] to all increase as  $\Theta(n_R)$  as well.<sup>1</sup>

## B. Related Work

In addition to the large body of work on point-to-point and multiuser MIMO systems, this paper is also related to several prior works that have studied the use of multiple receive antennas in ad hoc networks with Poisson distributed transmitters. References [11] and [12] considered precisely the same model as this paper, but studied the performance of slightly different receiver designs that turn out to yield very different performance. In [11] the receive filter is chosen according to the maximal ratio criterion (MRC) and thus only provides array gain,<sup>2</sup> while [12] considers the other extreme where the  $n_R$  antennas are used to cancel the strongest  $n_R - 1$  interferers but no array gain is obtained. Both receiver designs achieve only a *sublinear* density increase with  $n_R$ , with MRC scaling as  $n_R^{2/\alpha}$  and full interference cancellation scaling as  $n_R^{1-2/\alpha}$ , where  $\alpha$  is the path-loss exponent. On the other hand, we show that linear density scaling is achieved if a fraction of the receive degrees of freedom are used for array gain and the other fraction for interference cancellation, or an MMSE receiver is used.

In [13] the performance of the MMSE receiver is investigated in the same setting as in this paper, but for a fixed density network in which additional receive antennas are used to increase the per-link SINR/rate. It is shown that the average per-link SINR increases with  $n_R$  as  $n_R^{\alpha/2}$  which translates into only a logarithmic increase in per-link rate and overall system throughput. In contrast, we study the rather different setting in which the per-link rate is fixed and the density increases with  $n_R$ , and we show that exploiting the antennas

<sup>1</sup>We use the standard convention that a function  $f(n)$  is  $\Theta(n_R)$  if and only if there exist strictly positive constants  $k_1$  and  $k_2$  such that  $n \cdot k_1 \leq f(n) \leq n \cdot k_2$  for sufficiently large  $n$ .

<sup>2</sup>Although reference [11] also considers strategies involving multiple transmit antennas, here we have mentioned only the directly relevant single transmit/multiple-receive scenario results

to increase density instead of rate provides a significantly larger (i.e., linear versus logarithmic in  $n_R$ ) end-to-end benefit (c.f. Section IV-B). The MMSE receiver was also studied in [14], although under the additional assumption of transmit beamforming in the direction of the maximum eigenmode, and in particular the mean and variance of the receiver SINR was computed under an accurate Gamma approximation. The numerical results in [14] hint at the very large density increases that additional receive antennas provide, and our results rigorously show that the density of simultaneous transmissions can be increased linearly with the number of antennas. (Transmit beamforming does not change how density scales with the number of antennas.)

Early work on characterizing the throughput gains from MIMO in ad hoc networks includes [15]–[18] although these generally primarily employed simulations, while more recently [19]–[21] used tools similar to those used in the paper and developed by the present authors. Amongst these [17] is most relevant, as it considers mutually interfering multi-antenna transmit-receive pairs, albeit without an explicit spatial model. In that setting the number of receive antennas is also seen to be a performance bottleneck that can only be overcome if the number of antennas is increased in proportion to the number of interferers. This has some connection to our result showing that the number of antennas should be linearly proportional to the interferer density, although the results differ fundamentally in that there are always an infinite number of interferers in our spatial model.

The remainder of the paper is organized as follows. The system model and key metrics are described in Section II. The main results are derived in Section III, and then various extensions and relaxations of the model are considered in Section IV: in all of these diverse permutations we observe that the linear scaling result still holds. We conclude in Section V.

## II. SYSTEM MODEL AND METRICS

We consider a network in which the set of active transmitters are located according to a 2-D homogeneous Poisson point process (PPP) of density  $\lambda$  (transmitters/m<sup>2</sup>). Each transmitter communicates with a receiver a distance  $d$  meters away from it, where it is assumed that each receiver is randomly located on a circle of radius  $d$  centered around its associated transmitter. (Although meters are used here, this could be substituted with any other distance unit). Note that the receivers are not a part of the transmitter PPP. Each transmitter uses only one antenna, while each receiver has  $n_R$  antennas. The Poisson model is reasonable for uncoordinated networks, such as those using ALOHA. In Section IV-C we briefly examine a regular interferer geometry that is a more appropriate model for networks employing more sophisticated random access techniques.

By the stationarity of the Poisson process we can consider the performance of an arbitrary TX-RX pair, which we refer to as TX0 and RX0. From the perspective of RX0, the set of interferers (which is the entire transmit process with the exception of TX0) also form a homogeneous PPP due

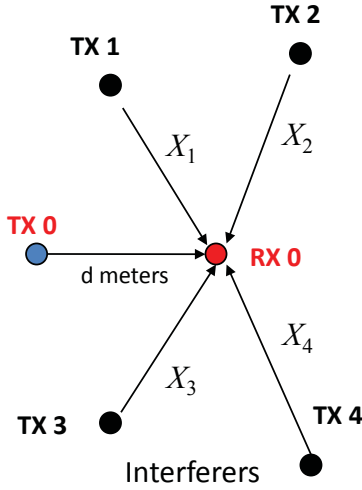


Fig. 1. Example transmit-receive pair with four nearby interferers shown. In addition to the distances  $d$  (desired) and  $X_i$  (interferer  $i$ ) shown, there is a channel vector  $\mathbf{h}_i$ .

to Slivnyak's Theorem; see [7] for additional discussion of this point and further explanation of the basic model. As a result, network-wide performance is characterized by the performance of a single TX-RX pair, separated by  $d$  meters and surrounded by an infinite number of interferers located on the infinite 2-D plane according to a homogeneous PPP with density  $\lambda$  interferers/m<sup>2</sup>. This setup is depicted in Fig. 1.

Assuming a path-loss exponent of  $\alpha$  ( $\alpha > 2$ ) and a frequency-flat channel, the  $n_R$ -dimensional received signal  $\mathbf{y}_0$  at RX0 is given by:

$$\mathbf{y}_0 = d^{-\alpha/2} \mathbf{h}_0 u_0 + \sum_{i \in \Pi(\lambda)} |X_i|^{-\alpha/2} \mathbf{h}_i u_i + \mathbf{z} \quad (1)$$

where  $|X_i|$  is the distance to the  $i$ -th transmitter/interferer (as determined by the realization of the PPP),  $\mathbf{h}_i \in \mathcal{C}^{n_R \times 1}$  is the vector channel from the  $i$ -th transmitter to RX0,  $\mathbf{z} \in \mathcal{C}^{n_R \times 1}$  is complex Gaussian noise with covariance  $\eta \mathbf{I}$ , and  $u_i$  is the data symbol of transmitter  $i$  with  $\mathbb{E}[|u_i|^2] = \rho$ . Consistent with a rich scattering environment, we assume that each of the vector channels  $\mathbf{h}_i$  have iid unit-variance complex Gaussian components, independent across transmitters.

### A. Performance Metrics

If a unit norm receive filter  $\mathbf{v}_0$  is used, the resulting signal-to-interference-and-noise ratio is:

$$\text{SINR} = \frac{\rho d^{-\alpha} |\mathbf{v}_0^\dagger \mathbf{h}_0|^2}{\eta + \sum_{i \in \Pi(\lambda)} \rho |X_i|^{-\alpha} |\mathbf{v}_0^\dagger \mathbf{h}_i|^2}. \quad (2)$$

Without loss of generality, we index the distances  $|X_i|$  in increasing order in order to take advantage of the property that the ordered squared-distances  $|X_1|^2, |X_2|^2, \dots$  follow a 1-D Poisson point process with intensity  $\pi \lambda$  [22]. To simplify notation we define the constant  $\text{SNR} \triangleq \frac{\rho d^{-\alpha}}{\eta}$  as the interference-free signal-to-noise ratio, which allows us to write:

$$\text{SINR} = \frac{|\mathbf{v}_0^\dagger \mathbf{h}_0|^2}{\frac{1}{\text{SNR}} + d^\alpha \sum_{i \in \Pi(\lambda)} |X_i|^{-\alpha} |\mathbf{v}_0^\dagger \mathbf{h}_i|^2}. \quad (3)$$

The received SINR depends on the interferer locations and the vector channels, both of which are random. The outage probability with respect to an SINR threshold  $\beta$  is:

$$P_{\text{out}}(\lambda) = \mathbb{P}[\text{SINR} \leq \beta], \quad (4)$$

which clearly is increasing in  $\lambda$ . The implicit assumption is that the channels and the set of active interferers are constant for the duration of a packet transmission but generally vary across transmissions. As a result, the outage probability accurately approximates the packet error probability experienced by each node; by the stationarity of the process, it also approximates the network-wide packet error probability.

It is often desirable from a system perspective to maintain a constant outage level  $\epsilon$  (e.g., to ensure that higher-layer reliability mechanisms are appropriately utilized), and thus the performance metric of interest is  $\lambda_\epsilon$ , the maximum interferer density such that the outage does not exceed  $\epsilon$ :

$$\lambda_\epsilon \triangleq \max_{\lambda} \{\lambda : P_{\text{out}}(\lambda) \leq \epsilon\}. \quad (5)$$

The *transmission capacity* is the number of successful transmissions per unit area, and can be defined as  $\lambda_\epsilon (1 - \epsilon) b$ , where  $b = \log_2(1 + \beta)$  is the data rate assuming a good channel code; this definition is akin to area spectral efficiency (ASE). Discussion of how this metric translates to end-to-end metrics such as *transport capacity* is provided in Section IV-B.

Although  $\lambda_\epsilon$  depends on the design parameters  $\epsilon$ ,  $\beta$ , and  $n_R$ , in this work we are interested in the behavior of  $\lambda_\epsilon$  with respect to  $n_R$  while the other parameters are kept fixed. (A justification for keeping  $\epsilon$  fixed has already been put forth, and justification for not increasing  $\beta$  is provided in Section IV-B.) Thus, we henceforth denote  $\lambda_\epsilon$  as a function of  $n_R$ .

### B. Receive Filters

The SINR and maximum density depend critically on the receive filter that is used. The receive filter can be used to either boost the power of the desired signal (by choosing  $\mathbf{v}_0$  in the direction of  $\mathbf{h}_0$ ) or to cancel interference, or some combination of the two. In this paper we consider the MMSE receiver, which optimally balances signal boosting and interference cancellation and maximizes the SINR, as well as a sub-optimal partial zero-forcing receiver, which uses a specified number of degrees of freedom for signal boosting and the remainder for cancellation. We assume that the receive filter is chosen based upon knowledge of the signal channel  $\mathbf{h}_0$  and the interfering channels  $\{\mathbf{h}_i\}_{i=1}^\infty$ ; this optimistic assumption of perfect receiver channel state information (CSI) is scrutinized in Section IV-A.

**MMSE Receiver.** From basic results in estimation theory, the MMSE receive filter is given by:

$$\mathbf{v}_0 = \frac{\boldsymbol{\Sigma}^{-1} \mathbf{h}_0}{\|\boldsymbol{\Sigma}^{-1} \mathbf{h}_0\|}. \quad (6)$$

where  $\boldsymbol{\Sigma}$  is the spatial covariance of the interference plus noise

$$\boldsymbol{\Sigma} \triangleq \frac{1}{\text{SNR}} \mathbf{I} + d^\alpha \sum_{i \in \Pi(\lambda)} |X_i|^{-\alpha} \mathbf{h}_i \mathbf{h}_i^\dagger. \quad (7)$$

Note that  $\boldsymbol{\Sigma}$  is the covariance matrix conditioned on the interferer channels and distances  $\{\mathbf{h}_i\}_{i=1}^\infty$  and  $\{|X_i|\}_{i=1}^\infty$ .

Amongst the set of all possible receive filters  $\mathbf{v}_0$ , the MMSE filter maximizes the received SINR. Its corresponding value is:

$$\text{SINR}^{\text{mmse}} = \mathbf{h}_0^\dagger \boldsymbol{\Sigma}^{-1} \mathbf{h}_0. \quad (8)$$

The corresponding outage probability and maximum density are denoted as  $P_{\text{out}}^{\text{mmse}}(\lambda)$  and  $\lambda_\epsilon^{\text{mmse}}(n_R)$ .

**Partial Zero Forcing Receiver.** We also study a suboptimal receiver that explicitly cancels interference from nearby transmitters while using the remaining degrees of freedom to boost the power of the desired signal. More specifically, the filter  $\mathbf{v}_0$  is chosen orthogonal to the channel vectors of the  $k$  nearest interferers  $\mathbf{h}_1, \dots, \mathbf{h}_k$ . The parameter  $k$  must be an integer and satisfy  $k \leq n_R - 1$ . Although performance could conceivably be improved by choosing  $k$  on the basis of the channel realizations, it is sufficient for our purposes to choose  $k$  in an offline fashion. The value of  $k$  is left unspecified for the time being. Amongst the filters satisfying the orthogonality requirement  $|\mathbf{v}_0^\dagger \mathbf{h}_i|^2 = 0$  for  $i = 1, \dots, k$ , we are interested in the one that maximizes the desired signal power  $|\mathbf{v}_0^\dagger \mathbf{h}_0|^2$ . By simple geometry, this corresponds to choosing  $\mathbf{v}_0$  in the direction of the projection of vector  $\mathbf{h}_0$  on the nullspace of vectors  $(\mathbf{h}_1, \dots, \mathbf{h}_k)$ . More precisely, if the columns of the  $n_R \times (n_R - k)$  matrix  $\mathbf{Q}$  form an orthonormal basis for the nullspace of  $(\mathbf{h}_1, \dots, \mathbf{h}_k)$ , then the receive filter is chosen as:

$$\mathbf{v}_0 = \frac{\mathbf{Q}^\dagger \mathbf{h}_0}{\|\mathbf{Q}^\dagger \mathbf{h}_0\|}. \quad (9)$$

The corresponding outage probability and maximum density are denoted  $P_{\text{out}}^{\text{pzf}-k}(\lambda)$  and  $\lambda_\epsilon^{\text{pzf}-k}(n_R)$ , respectively. Notice that  $k = 0$  and  $k = n_R - 1$  correspond to the extremes of MRC ( $\mathbf{v}_0 = \mathbf{h}_0/\|\mathbf{h}_0\|$ ) and interference cancellation of the maximum number of interferers (full zero forcing), respectively. Because the MMSE receiver is SINR-maximizing, we clearly have

$$P_{\text{out}}^{\text{mmse}}(\lambda) \leq P_{\text{out}}^{\text{pzf}-k}(\lambda) \quad \text{and} \quad \lambda_\epsilon^{\text{mmse}}(n_R) \geq \lambda_\epsilon^{\text{pzf}-k}(n_R) \quad (10)$$

for all  $k$  and any set of system parameters.

Although suboptimal, it is beneficial to study the PZF receiver because it is generally more amenable to analysis than the MMSE filter and because its simple structure allows us to clearly understand why linear scaling is achievable. Note, however, that an MMSE filter should be used in practice because it also is linear and its CSI requirements are less stringent (PZF requires knowledge of individual interferer channels, whereas MMSE only requires knowledge of the spatial covariance of the aggregate interference). Furthermore, note that after the submission of this paper a closed-form expression for the outage probability with an MMSE receiver was derived in [23], [24].

### III. MAIN RESULTS: DENSITY SCALING WITH RECEIVE ANTENNAS

In this section we prove the main result of the paper, which is that  $\lambda_\epsilon^{\text{mmse}}(n_R)$  and  $\lambda_\epsilon^{\text{pzf}-k}(n_R)$  both scale linearly with  $n_R$ . We prove this result in two parts: we first show that a lower bound to  $\lambda_\epsilon^{\text{pzf}-k}(n_R)$ , and thus to  $\lambda_\epsilon^{\text{mmse}}(n_R)$ , increases linearly with  $n_R$ , and then show that an upper bound on  $\lambda_\epsilon^{\text{mmse}}(n_R)$  also increases linearly with  $n_R$ .

#### A. Lower Bound: Achievability of Linear Scaling with Partial Zero Forcing

First, we show that linear scaling is achievable by finding a lower bound on  $\lambda_\epsilon^{\text{pzf}-k}(n_R)$  that is linear in  $n_R$ . In order to develop the bound, we first statistically characterize the signal and interference coefficients when the PZF receiver is used. These characterizations are a consequence of the basic result that the squared-norm of the projection of a  $n_R$ -dimensional vector with iid unit-variance complex Gaussian components onto an independent  $s$ -dimensional subspace is a  $\chi_{2s}^2$  random variable.<sup>3 4</sup>

We denote the signal and interference coefficients as

$$S \triangleq |\mathbf{v}_0^\dagger \mathbf{h}_0|^2 \quad (11)$$

$$H_i \triangleq |\mathbf{v}_0^\dagger \mathbf{h}_i|^2 \quad i = 1, 2, \dots \quad (12)$$

and characterize the statistics of these coefficients in the following lemma.

*Lemma 1:* For PZF- $k$ , the signal coefficient  $S$  is  $\chi_{2(n_R-k)}^2$ , the interference terms  $H_1, \dots, H_k$  are zero, and coefficients  $H_{k+1}, H_{k+2}, \dots$  are iid unit-mean exponential (i.e.,  $\chi_2^2$ ). Furthermore,  $S, H_{k+1}, H_{k+2}, \dots$  are mutually independent.

*Proof:* See Appendix A. ■

Using this statistical characterization and the definitions in (11)-(12) the received SINR is

$$\text{SINR}^{\text{pzf}-k} = \frac{S}{\frac{1}{\text{SNR}} + d^\alpha \sum_{i=k+1}^{\infty} |X_i|^{-\alpha} H_i} \quad (13)$$

where the  $S$  and  $H_i$  terms are characterized in Lemma 1, the quantities  $|X_{k+1}|^2, |X_{k+2}|^2, \dots$  are the  $k+1, k+2, \dots$  ordered points of a 1-D PPP with intensity  $\pi\lambda$ , and the ordered points are independent of the signal and interference terms. The aggregate interference power for PZF- $k$  is denoted as:

$$I_k \triangleq d^\alpha \sum_{i=k+1}^{\infty} |X_i|^{-\alpha} H_i. \quad (14)$$

and the expectation of this interference power is characterized in the following lemma:

*Lemma 2:* For  $k > \frac{\alpha}{2} - 1$ , the expected interference power is characterized as:

$$\mathbb{E}[I_k] = (\pi d^2 \lambda)^{\frac{\alpha}{2}} \sum_{i=k+1}^{\infty} \frac{\Gamma(i - \frac{\alpha}{2})}{\Gamma(i)} \quad (15)$$

$$< (\pi d^2 \lambda)^{\frac{\alpha}{2}} \left(\frac{\alpha}{2} - 1\right)^{-1} \left(k - \left\lceil \frac{\alpha}{2} \right\rceil\right)^{1 - \frac{\alpha}{2}}, \quad (16)$$

where  $\Gamma(\cdot)$  is the gamma function and  $\lceil \cdot \rceil$  is the ceiling function, with the upper bound valid for  $k > \lceil \frac{\alpha}{2} \rceil$ .

*Proof:* See Appendix B. ■

To derive the main result for PZF, we use Lemma 2 to upper bound  $\mathbb{E}[1/\text{SINR}^{\text{pzf}-k}]$  and then combine this with Markov's

<sup>3</sup>This result is shown for real-valued vectors/matrices in [25], and easily extends to the complex setting. Because the subspace is independent of the vector and because the vector is spatially isotropic, without loss of generality one can assume that the projection is performed on the space spanned by the first  $s$  elementary basis vectors. As a result, the projection operation zeroes all but the first  $s$  elements of the vector, from which the result follows.

<sup>4</sup>Throughout the paper we abide by communications literature convention and define a  $\chi_{2s}^2$  random variable to have PDF  $f(x) = \frac{x^{s-1} e^{-x}}{(s-1)!}$ , which may differ slightly from the definition in probability literature.

inequality to reach an outage probability upper bound, and, inversely, a density lower bound:

*Theorem 1:* The outage probability with PZF- $k$  is upper bounded by:

$$P_{\text{out}}^{\text{pzf}-k}(\lambda) \leq \frac{\beta \left( (\pi d^2 \lambda)^{\frac{\alpha}{2}} \left( \frac{\alpha}{2} - 1 \right)^{-1} \left( k - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1 - \frac{\alpha}{2}} + \frac{1}{\text{SNR}} \right)}{n_{\text{R}} - k - 1} \quad (17)$$

for  $\left\lceil \frac{\alpha}{2} \right\rceil < k < n_{\text{R}} - 1$ . In turn, the maximum density  $\lambda_{\epsilon}^{\text{pzf}-k}(n_{\text{R}})$  is lower bounded by:

$$\lambda_{\epsilon}^{\text{pzf}-k}(n_{\text{R}}) \geq \left( \frac{\epsilon}{\beta} \right)^{\frac{2}{\alpha}} \frac{\left( \frac{\alpha}{2} - 1 \right)^{\frac{2}{\alpha}}}{\pi d^2} \left( n_{\text{R}} - k - 1 - \frac{\beta}{\epsilon \text{SNR}} \right)^{\frac{2}{\alpha}} \times \left( k - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1 - \frac{2}{\alpha}} \quad (18)$$

for any  $k$  satisfying  $\left\lceil \frac{\alpha}{2} \right\rceil < k < n_{\text{R}} - 1 - \frac{\beta}{\epsilon \text{SNR}}$ .

*Proof:* The outage upper bound is derived by rewriting the outage probability as the tail probability of random variable  $1/\text{SINR}$  and then applying Markov's inequality as follows:

$$\begin{aligned} P_{\text{out}}^{\text{PZF}-k}(\lambda) &= \mathbb{P} \left[ \frac{1}{\text{SINR}^{\text{pzf}-k}} \geq \frac{1}{\beta} \right] \\ &\stackrel{(a)}{\leq} \beta \cdot \mathbb{E} \left[ \frac{1}{\text{SINR}^{\text{pzf}-k}} \right] \\ &\stackrel{(b)}{=} \beta \cdot \mathbb{E} \left[ I_k + \frac{1}{\text{SNR}} \right] \mathbb{E} \left[ \frac{1}{S} \right] \\ &\stackrel{(c)}{<} \frac{\beta \left( (\pi d^2 \lambda)^{\frac{\alpha}{2}} \left( \frac{\alpha}{2} - 1 \right)^{-1} \left( k - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1 - \frac{\alpha}{2}} + \frac{1}{\text{SNR}} \right)}{n_{\text{R}} - k - 1}, \end{aligned}$$

where (a) is due to Markov's inequality, (b) is due to (13) and the independence of  $I_k$  and  $S$ , and (c) follows from Lemma 2 and because  $S$  is  $\chi_{2(n_{\text{R}}-k)}^2$  and  $\mathbb{E}[1/\chi_{2l}^2] = 1/(l-1)$  for  $l > 1$ . Setting this bound equal to  $\epsilon$  and then solving for  $\lambda$  yields the associated lower bound to  $\lambda_{\epsilon}$ . ■

It is worthwhile to note that the  $\left( n_{\text{R}} - k - 1 - \frac{\beta}{\epsilon \text{SNR}} \right)^{2/\alpha}$  term in the  $\lambda_{\epsilon}$  lower bound is the density increase due to array gain (i.e., increased signal power), while the  $\left( k - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1-2/\alpha}$  term is the density increase due to interference cancellation. Thus, the bound succinctly illustrates the tradeoff between array gain and interference cancellation.

In order to show the achievability of linear scaling, we need only appropriately increase  $k$  with  $n_{\text{R}}$ . If we choose the number of cancelled interferers  $k = \theta n_{\text{R}}$  for some constant  $0 < \theta < 1$ , the density lower bound becomes:

$$\begin{aligned} \lambda_{\epsilon}^{\text{pzf}-\theta n_{\text{R}}}(n_{\text{R}}) &\geq \left( \frac{\epsilon}{\beta} \right)^{\frac{2}{\alpha}} \frac{\left( \frac{\alpha}{2} - 1 \right)^{\frac{2}{\alpha}}}{\pi d^2} (1 - \theta)^{\frac{2}{\alpha}} \theta^{1 - \frac{2}{\alpha}} \\ &\quad \times \left( n_{\text{R}} - \frac{1 + \frac{\beta}{\epsilon \text{SNR}}}{1 - \theta} \right)^{\frac{2}{\alpha}} \left( n_{\text{R}} - \theta^{-1} \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1 - \frac{2}{\alpha}}. \end{aligned} \quad (19)$$

Because the conditions for Theorem 1 are satisfied for sufficiently large  $n_{\text{R}}$  if  $k = \theta n_{\text{R}}$  with  $0 < \theta < 1$  (for any  $\epsilon > 0$  and  $0 < \beta < \text{SNR}$ ), the lower bound scales linearly with  $n_{\text{R}}$ . This result is formally stated as follows:

*Lemma 3:* For any  $\theta$  satisfying  $0 < \theta < 1$ ,

$$\frac{\lambda_{\epsilon}^{\text{pzf}-\theta n_{\text{R}}}(n_{\text{R}})}{n_{\text{R}}} \geq \left( \frac{\epsilon}{\beta} \right)^{\frac{2}{\alpha}} \frac{\left( \frac{\alpha}{2} - 1 \right)^{\frac{2}{\alpha}}}{\pi d^2} (1 - \theta)^{\frac{2}{\alpha}} \theta^{1 - \frac{2}{\alpha}}, \quad (20)$$

for sufficiently large  $n_{\text{R}}$ .

This perhaps surprising scaling result can be intuitively understood by examining how the expected signal and interference power increase with  $n_{\text{R}}$ . Choosing  $\theta < 1$  ensures that the signal power, which is  $\chi_{2(1-\theta)n_{\text{R}}}^2$ , increases linearly with  $n_{\text{R}}$ . Based on the upper bound in Lemma 2 we can see that the condition  $\theta > 0$  ensures that the interference power increases only linearly with  $n_{\text{R}}$  if  $\lambda$  is linear in  $n_{\text{R}}$ . These linear terms are offsetting, and thus allow an approximately constant SINR to be maintained as  $\lambda$  is increased linearly with  $n_{\text{R}}$ .

On the other hand, linear scaling does not occur if  $k$  is kept constant. Specifically, if  $k = \kappa$  for any constant  $\kappa$ , then the signal power increases linearly with  $n_{\text{R}}$  as desired but the interference power increases too quickly with the density (as  $\lambda^{\alpha/2}$ ), thereby limiting the density growth to  $n_{\text{R}}^{2/\alpha}$ . Linear scaling also does not hold at the other extreme where all but a fixed number of degrees of freedom are used for cancellation: if  $k = n_{\text{R}} - \kappa$  for any constant  $\kappa$ , then the interference power scales appropriately with the density but the signal power is  $\chi_{2\kappa}^2$  and thus does not increase with  $n_{\text{R}}$ , thereby limiting the density increase to  $n_{\text{R}}^{1-2/\alpha}$ . These results are consistent with the findings of [11] and [12].

### B. Upper Bound: The MMSE Receiver

While the earlier result showed that a lower bound to  $\lambda_{\epsilon}^{\text{mmse}}(n_{\text{R}})$  scales linearly with  $n_{\text{R}}$ , we make our scaling characterization more precise by finding an upper bound to  $\lambda_{\epsilon}^{\text{mmse}}(n_{\text{R}})$  that also scales linearly with  $n_{\text{R}}$ . In order to obtain such a bound, we utilize the MMSE performance outage probability lower bound from [5]. Extended to the model in this paper, the bound states:

$$P_{\text{out}}^{\text{mmse}}(\lambda) \geq \mathbb{P} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{\sum_{i=n_{\text{R}}}^{\infty} |X_i|^{-\alpha} H_i} \leq \beta \right] \quad (21)$$

where  $\mathbf{h}_0$  and  $|X_1|, |X_2|, \dots$  are defined as before, and the random variables  $H_1, H_2, \dots$  are iid, unit-mean exponential random variables (independent of all other random variables). The bound in [5] is for fixed interferer distances and no thermal noise. However, by averaging over the interferer locations (according to the PPP) and using the fact that outage probability is increasing in the noise power  $\eta$ , we obtain (21) and see that it also holds in the presence of noise.

The SIR expression in (21) is closely related to the SINR characterization for PZF- $k$  in (13). The denominator of the SIR in (21) is precisely as if a PZF receiver with  $k = n_{\text{R}} - 1$  is used (i.e., the nearest  $n_{\text{R}} - 1$  interferers are cancelled, and the effective fading coefficients from the uncanceled interferers are iid exponential), while the numerator corresponds to PZF with  $k = 0$ . Thus, the bound in (21) corresponds to an idealized setting where the receive filter cancels the nearest  $n_{\text{R}} - 1$  interferers but still is in the direction of  $\mathbf{h}_0$ .

We translate this outage lower bound into a density upper bound in a manner that is complementary to Theorem 1: we upper bound the expected SIR and then apply Markov's

inequality to the success probability to obtain the following result:

*Theorem 2:* The outage probability with an MMSE receiver is lower bounded by:

$$\lambda_{\text{out}}^{\text{mmse}}(\lambda) \geq 1 - \frac{d^{-\alpha}}{\beta} \left( \frac{2n_{\text{R}} + 1 + \frac{\alpha}{2}}{\pi\lambda} \right)^{\alpha/2} \quad (22)$$

and, in turn,  $\lambda_{\epsilon}^{\text{mmse}}(n_{\text{R}})$  is upper bounded by:

$$\lambda_{\epsilon}^{\text{mmse}}(n_{\text{R}}) \leq \frac{2n_{\text{R}} + 1 + \frac{\alpha}{2}}{\pi d^2 \beta^{2/\alpha} (1 - \epsilon)^{2/\alpha}}. \quad (23)$$

*Proof:* See Appendix C. ■

This upper bound, which by (10) also applies to PZF, scales linearly with  $n_{\text{R}}$ . Combining Theorems 1 and 2, it is clear that  $\lambda_{\epsilon}^{\text{mmse}}(n_{\text{R}})$  and  $\lambda_{\epsilon}^{\text{pzf}}(n_{\text{R}})$  both scale linearly with  $n_{\text{R}}$ , although numerical results in III-E will confirm that MMSE is better by a non-negligible constant factor.

Because the SINR with a PZF receiver differs only slightly from the expression in (21), we can use precisely the argument of Appendix C to derive an upper bound on  $\lambda_{\epsilon}^{\text{pzf}-k}(n_{\text{R}})$ . While in the proof of Theorem 2 we retained the contribution of the nearest  $n_{\text{R}} + 1$  uncanceled interferers in (21), for PZF we retain the contribution of the nearest  $l$  uncanceled interferers. This results in the following bound:

$$\lambda_{\epsilon}^{\text{pzf}-k}(n_{\text{R}}) \leq \frac{k + l + \alpha/2}{\pi d^2 \beta^{2/\alpha} (1 - \epsilon)^{2/\alpha}} \left( \frac{n_{\text{R}} - k}{l - 1} \right)^{2/\alpha} \quad (24)$$

which, unlike the PZF density lower bound in Theorem 1, applies for any  $0 \leq k \leq n_{\text{R}} - 1$ , and furthermore holds for any integer  $l > 1$ .

For  $k = 0$  (MRC), if we choose  $l = 2$  the upper bound becomes

$$\lambda_{\epsilon}^{\text{pzf}}(0) \leq \frac{2 + \alpha/2}{\pi d^2 \beta^{2/\alpha} (1 - \epsilon)^{2/\alpha}} n_{\text{R}}^{2/\alpha}, \quad (25)$$

while for  $k = n_{\text{R}} - 1$  (full ZF) we choose  $l = n_{\text{R}} + 1$  to get

$$\lambda_{\epsilon}^{\text{pzf}}(n_{\text{R}} - 1) \leq \frac{2 + \alpha/(2n_{\text{R}})}{\pi d^2 \beta^{2/\alpha} (1 - \epsilon)^{2/\alpha}} n_{\text{R}}^{1-2/\alpha}. \quad (26)$$

For MRC the upper bound is  $O(n_{\text{R}}^{2/\alpha})$  while for full zero-forcing it is  $O(n_{\text{R}}^{1-2/\alpha})$ . These upper bounds complement the matching lower bounds in [11] and [12], respectively.

### C. Array Gain v. Interference Cancellation

Because the MMSE receiver implicitly balances, through (6), array gain and interference cancellation, it is not evident how the MMSE utilizes the receive degrees of freedom. For the sake of simplicity we limit the discussion in this section to the case where thermal noise is negligible, while keeping in mind that the MMSE filter considers the balance between noise power and interference power through (6). If the eigenvalues of the interference covariance  $\Sigma$  are roughly equal then the MMSE filter is nearly in the direction of  $\mathbf{h}_0$ ; on the other hand, if the eigenvalues are very disparate then the MMSE filter is (approximately) in the direction of the projection of  $\mathbf{h}_0$  on the subspace orthogonal to the directions of the strong interfering eigenmodes. Thus, the fraction of degrees of freedom used for array gain instead of interference

cancellation depends critically on the spread of the eigenvalues of  $\Sigma$ , which turns out to depend on the path loss exponent  $\alpha$ .

The MMSE receiver can be understood by studying the PZF receiver, and in particular by finding the value of  $\theta$  that maximizes the PZF density. Based on Lemma 3 it is clear that the PZF density lower bound depends on  $\theta$  only through the term  $(1 - \theta)^{\frac{2}{\alpha}} \theta^{1 - \frac{2}{\alpha}}$  for large  $n_{\text{R}}$ .

To determine the dependence of the PZF upper bound in (24), we first minimize the bound with respect to  $l$ , for large  $n_{\text{R}}$  and  $k = \theta n_{\text{R}}$ . A simple calculation finds that  $l = \left( \frac{2/\alpha}{1 - 2/\alpha} \right) \theta n_{\text{R}}$  is the minimizer, and with this choice of  $l$  the dependence of the density upper bound also occurs only through the term  $(1 - \theta)^{\frac{2}{\alpha}} \theta^{1 - \frac{2}{\alpha}}$ .

Thus, the upper and lower bounds depend on  $\theta$  only through the term  $(1 - \theta)^{\frac{2}{\alpha}} \theta^{1 - \frac{2}{\alpha}}$ . By taking the derivative (w.r.t.  $\theta$ ) and solving, we find that the maximizing value of  $\theta$  is:

$$\theta^* = 1 - \frac{2}{\alpha}. \quad (27)$$

As  $\alpha \rightarrow 2$  the degrees of freedom should be used to boost signal power rather than to cancel interference (i.e.,  $\theta^* \rightarrow 0$ ), because far-away interference is significant and so cancelling a few nearby interferers provides a smaller benefit than using the antennas for array gain. At the other extreme,  $\theta^* \rightarrow 1$  as the path loss exponent increases because the power from nearby interferers begins to dominate and thus the antennas are more profitably used for interference cancellation than array gain.

As it turns out, the optimizing value  $\theta^*$  and the above intuition are also consistent with the MMSE receiver. Fig. 2 contains a plot of  $\mathbb{E} \left[ |\mathbf{v}_0^\dagger \mathbf{h}_0|^2 / (|\mathbf{h}_0|^2 |\mathbf{v}_0|^2) \right]$ , the expectation of the squared correlation between the normalized MMSE filter and channel vector, versus  $\alpha$  for  $n_{\text{R}} = 8$  (no thermal noise). This quantity effectively measures the fraction of degrees of freedom used for array gain. For the PZF receiver this metric is precisely equal to  $n_{\text{R}} - k$ , and thus would be equal to  $1 - \theta^* = 2/\alpha$  if  $k$  was chosen as  $k = \theta^* n_{\text{R}}$ . Although the MMSE receiver implicitly balances array gain and cancellation (as opposed to the explicit balance for the PZF receiver), the plot shows that the MMSE receiver also utilizes (approximately) a fraction  $2/\alpha$  of its receive degrees of freedom for array gain. Similar to the intuition stated for the behavior of  $\theta^*$  for PZF, the eigenvalues of  $\Sigma$  become more disparate as  $\alpha$  is increased, and the MMSE receiver takes advantage of this by performing more interference cancellation (and thus providing less array gain) when  $\alpha$  is larger.

### D. Improved Lower and Upper Bounds

Although the bounds developed in Sec. III-A and III-B are sufficient to show linear scaling, both of them are quite loose. This looseness is primarily due to the use of Markov's inequality, and the following more accurate bounds are instead derived through application of Chebychev's inequality:

*Theorem 3:* The outage probability for the PZF filter is

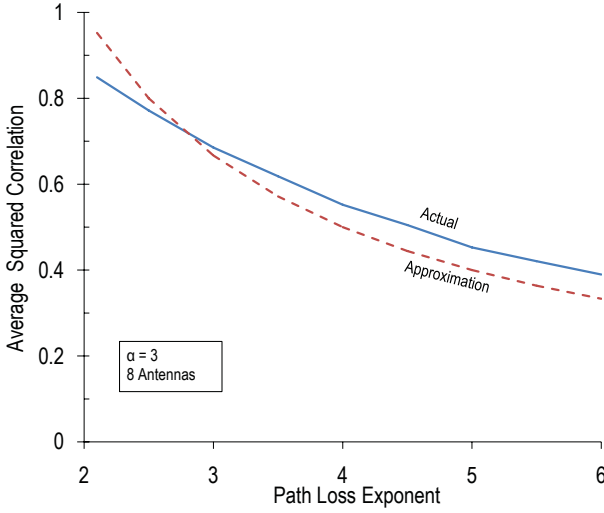


Fig. 2. Average squared correlation between signal channel and MMSE filter as a function of path loss exponent  $\alpha$ . Also shown is the approximation  $2/\alpha$ .

upper bounded by

$$P_{\text{out}}^{\text{pzf-k}}(\lambda) \leq \mathbb{P}(S \leq \sigma^*) + \beta^2 \text{Var}^{\text{ub}}(I_k) \int_{\sigma^*}^{\infty} \frac{1}{\left(s - \frac{\beta}{\text{SNR}} - \beta \mathbb{E}[I_k]\right)^2} f_S(s) ds. \quad (28)$$

where  $S \sim \chi_{2(n_R-k)}^2$ ,  $I_k$  is defined in (14),  $\sigma^* = \beta \left( \mathbb{E}[I_k] + \frac{1}{\text{SNR}} + \sqrt{\text{Var}^{\text{ub}}(I_k)} \right)$ , and

$$\text{Var}^{\text{ub}}(I_k) = (\pi d^2 \lambda)^\alpha \left( \sum_{i=k+1}^{\infty} \left( \mathbb{E}[T_i^{-\alpha}] + \text{Var}(T_i^{-\frac{\alpha}{2}}) \right) + 2 \sum_{i=k+1}^{\infty} \sqrt{\text{Var}(T_i^{-\frac{\alpha}{2}})} \sum_{j=i+1}^{\infty} \sqrt{\text{Var}(T_j^{-\frac{\alpha}{2}})} \right).$$

*Theorem 4:* The outage probability for the MMSE filter is lower bounded by

$$P_{\text{out}}^{\text{mmse}}(\lambda) \geq 1 - \frac{\text{Var}^{\text{ub}}(S/I_{n_R})}{(\beta - \mathbb{E}^{\text{lb}}[S/I_{n_R}])^2} \quad (29)$$

where  $S \sim \chi_{2n_R}^2$ ,  $I_{n_R}$  is defined in (14), and

$$\mathbb{E}^{\text{lb}}[S/I_{n_R}] = n_R / (\pi d^2 \lambda)^{\frac{\alpha}{2}} \sum_{i=n_R}^{\infty} \mathbb{E}[T_i^{-\frac{\alpha}{2}}] \quad (30)$$

$$\text{Var}^{\text{ub}}(S/I_{n_R}) = \frac{n_R(n_R + 1)}{(\pi d^2 \lambda)^\alpha \left( \sum_{i=n_R}^{\infty} e^{-(\gamma + \frac{\alpha}{2} \psi_0(i))} \right)^2} - \frac{n_R^2}{(\pi d^2 \lambda)^\alpha \left( \sum_{i=n_R}^{\infty} \frac{\Gamma(i - \frac{\alpha}{2})}{\Gamma(i)} \right)^2}, \quad (31)$$

and  $\gamma$  is the Euler-Mascheroni constant and  $\psi_0(i)$  is the poly-Gamma function.

In the two theorems the random variables  $T_i$  are each chi-square with  $2i$  degrees of freedom, and have moments characterized by:

$$\mathbb{E}[T_i^b] = \frac{\Gamma(i+b)}{\Gamma(i)} \quad (32)$$

$$\text{Var}(T_i^b) = \frac{\Gamma(i+2b)}{\Gamma(i)} - \left( \frac{\Gamma(i+b)}{\Gamma(i)} \right)^2, \quad \text{if } i+b > 0. \quad (33)$$

Although space constraints preclude inclusion of the proofs, they can be found in the online version of this paper [26]. By equating these bounds to  $\epsilon$  and solving (numerically) for  $\lambda$ , the PZF and MMSE densities can be lower and upper bounded, respectively.

#### E. Numerical Results

In Figures 3 and 4 the numerically computed maximum densities for the MMSE receiver  $\lambda_\epsilon^{\text{mmse}}(n_R)$  and the PZF receiver  $\lambda_\epsilon^{\text{pzf-k}}(n_R)$  with  $k = \theta^* n_R$  are plotted on a log-log scale versus  $n_R$  for  $\alpha = 3$  and  $\alpha = 4$ , along with the PZF lower bounds (from Theorems 1 and 3), the MMSE upper bounds (from Theorems 2 and 4), and the densities for MRC and full zero-forcing ( $\lambda_\epsilon^{\text{pzf-k}}(n_R)$  with  $k = 0$  and  $k = n_R - 1$ , respectively). In each plot, the tighter of the upper and the tighter of the lower bounds correspond to the Chebychev-based bounds in the previous section. The bounds and numerically computed densities are representative of the linearly increasing density for PZF and MMSE, whereas MRC and full zero forcing both exhibit much poorer scaling. Figures 5 and 6 provide linear plots of the maximum density versus  $n_R$  for more realistic numbers of antennas. Even a few antennas allow for very large density gains, and thus the asymptotic scaling results also are indicative of performance for small values of  $n_R$ . The plots also make it clear that MMSE and PZF are strongly preferred to MRC or full zero forcing, and also that a non-negligible benefit is afforded by using the optimal MMSE filter rather than PZF.

#### IV. GENERALIZATIONS AND EXTENSIONS OF THE MODEL

In this section, we explore three of the potentially controversial aspects of our model to show that the linear scaling result is not an artifact of our model and assumptions: we remove the assumption of perfect CSI at the receiver, we evaluate the benefit of antennas from an end-to-end perspective, and we consider the importance of the interferer geometry.

##### A. Effect of Imperfect CSI

The objective of this section is to illustrate that reasonable performance is achieved even if the receiver has to estimate the CSI, instead of assuming this information is *a priori* provided to the receiver. In order to design the optimal MMSE filter, the receiver requires an estimate of  $\mathbf{h}_0$ , the signal channel, and  $\Sigma$ , the interference (plus noise) covariance. The desired channel can be estimated at the receiver via pilot symbols and the effects of such training error are well understood [27]. On the other hand, it is not as clear how the receiver can estimate  $\Sigma$  and what effect estimation error has on performance.

Recall that the covariance  $\Sigma$  depends on the interferer locations, and thus on the active interferers, as well as

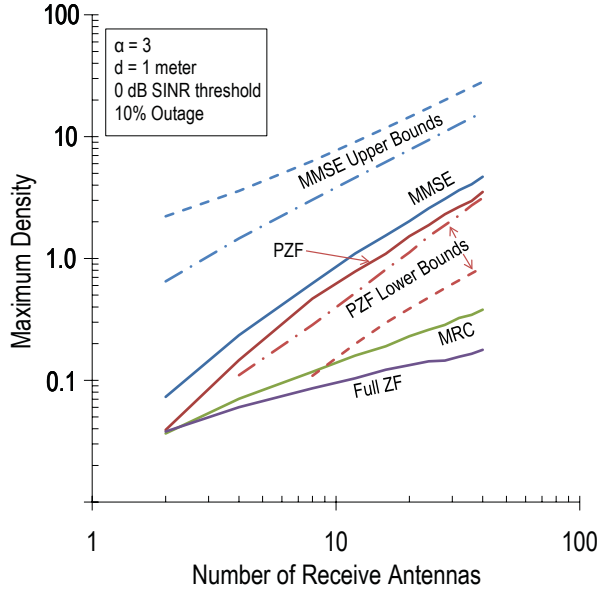


Fig. 3. Density versus  $n_R$  for  $\epsilon = .1$ ,  $\beta = 1$ ,  $\alpha = 3$ ,  $\theta = \frac{1}{3}$ ,  $d = 1$ .

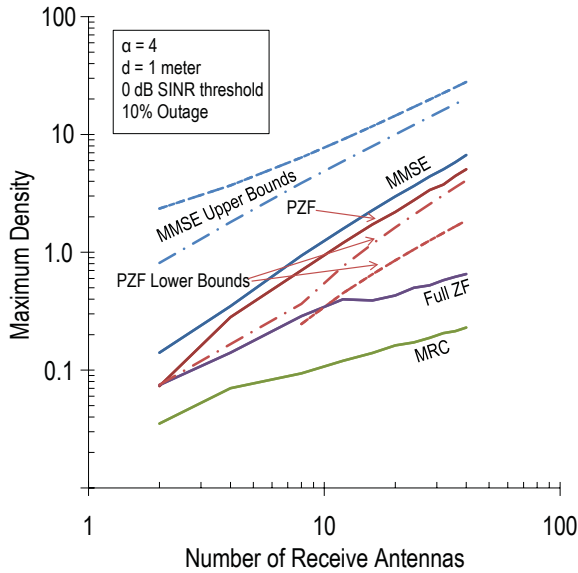


Fig. 4. Density versus  $n_R$  for  $\epsilon = .1$ ,  $\beta = 1$ ,  $\alpha = 4$ ,  $\theta = \frac{1}{2}$ ,  $d = 1$ .

the instantaneous channel realizations. Although coordinated transmission of pilots seems infeasible in a decentralized network, the receiver can estimate the covariance by listening to interferer transmissions, in the absence of desired signal. If the desired transmitter remains quiet for  $K$  symbols, the receiver can use the  $K$  observations of noise plus interference to form the sample covariance

$$\hat{\Sigma} \triangleq \frac{1}{K} \sum_{i=1}^K \mathbf{r}_i \mathbf{r}_i^\dagger \quad (34)$$

where  $\mathbf{r}_i$  represents the  $i$ -th observation of the noise plus

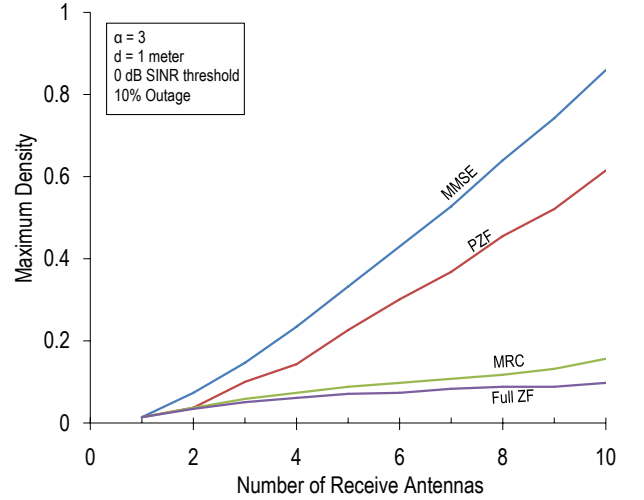


Fig. 5. Density versus  $n_R$  for  $\epsilon = .1$ ,  $\beta = 1$ ,  $\alpha = 3$ ,  $\theta = \frac{1}{3}$ ,  $d = 1$ .

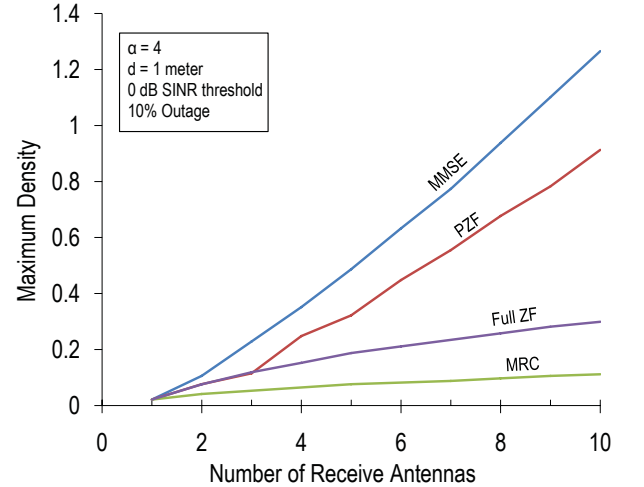


Fig. 6. Density versus  $n_R$  for  $\epsilon = .1$ ,  $\beta = 1$ ,  $\alpha = 4$ ,  $\theta = \frac{1}{2}$ ,  $d = 1$ .

interference. Assuming, for simplicity, knowledge of  $\mathbf{h}_0$ , the receiver can then use the filter  $\hat{\Sigma}^{-1} \mathbf{h}_0$  and the corresponding SINR is

$$\text{SINR} = \frac{(\mathbf{h}_0^\dagger \hat{\Sigma}^{-1} \mathbf{h}_0)^2}{\mathbf{h}_0^\dagger \hat{\Sigma}^{-\dagger} \hat{\Sigma} \hat{\Sigma}^{-1} \mathbf{h}_0}. \quad (35)$$

This SINR was analyzed in [28] assuming that all interferers transmit independent Gaussian symbols, and it was shown that for every  $\Sigma$ , the expected SINR using filter  $\hat{\Sigma}^{-1} \mathbf{h}_0$  (expectation w.r.t the distribution of  $\hat{\Sigma}$ ), and the SINR using the correct filter  $\Sigma^{-1} \mathbf{h}_0$  are related according to:

$$\mathbb{E}_{\hat{\Sigma}} \left[ \frac{(\mathbf{h}_0^\dagger \hat{\Sigma}^{-1} \mathbf{h}_0)^2}{\mathbf{h}_0^\dagger \hat{\Sigma}^{-\dagger} \hat{\Sigma} \hat{\Sigma}^{-1} \mathbf{h}_0} \right] = \left( 1 - \frac{n_R - 1}{K + 1} \right) \mathbf{h}_0^\dagger \Sigma^{-1} \mathbf{h}_0, \quad (36)$$

where from (8),  $\mathbf{h}_0^\dagger \Sigma^{-1} \mathbf{h}_0$  is the SINR when the proper MMSE filter is used. By taking an additional expectation with



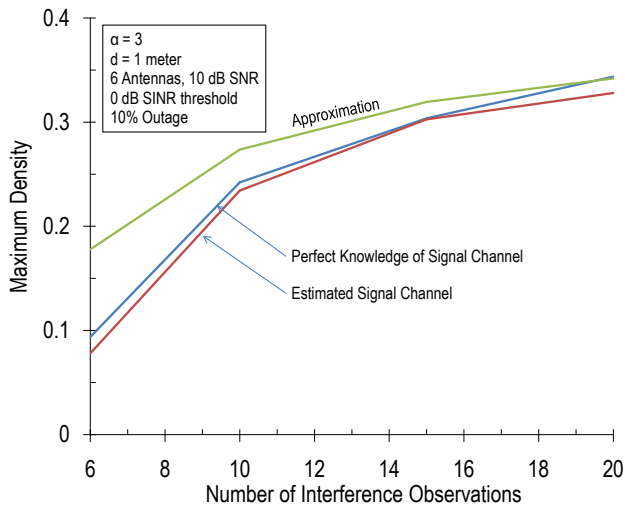


Fig. 7. Maximum density versus  $K$ , number of interference observations, for  $\epsilon = .1$ ,  $\beta = 1$ ,  $\alpha = 3$ ,  $d = 1$ , and SNR = 10 dB.

respect to  $\Sigma$  (which is determined by the interferer locations and channels), we see that the expected SINR using an MMSE filter based upon the sample covariance  $\hat{\Sigma}$  is precisely a factor of  $1 - \frac{n_R - 1}{K + 1}$  smaller than the expected SINR with perfect knowledge of  $\Sigma$ . As expected, this factor is increasing in  $K$  and converges to one as  $K \rightarrow \infty$ , because  $\hat{\Sigma} \rightarrow \Sigma$  as  $K \rightarrow \infty$ . If  $K = 2n_R - 3$ , the expected SINR is decreased by 3 dB.

Although the result of [28] applies to the expected SINR, numerical results confirm that the results also apply to the outage scenario considered here. Therefore, a system using the sample covariance from  $K$  observations and SINR threshold  $\beta$  has, approximately, the same maximum density as a system with perfect CSI and SINR threshold  $\beta / \left(1 - \frac{n_R - 1}{K + 1}\right)$ . From the various bounds, we see that the maximum density depends on the SINR threshold as  $\beta^{-2/\alpha}$ . Thus, using the  $K$ -observation sample covariances instead of the true covariance reduces the density, approximately, by a factor of  $\left(1 - \frac{n_R - 1}{K + 1}\right)^{2/\alpha}$ . If we choose  $K = 2n_R - 3$  then the loss factor is  $2^{2/\alpha}$  for all  $n_R$ ; therefore, by appropriately scaling  $K$  linearly with  $n_R$  performance within a constant factor of the perfect CSI benchmark is achieved.

To solidify these conclusions, in Fig. 7 the maximum density is plotted versus  $K$  for a 6 antenna system with SNR = 10 dB. The curves correspond to perfect knowledge of  $\mathbf{h}_0$ , estimation of  $\mathbf{h}_0$  on the basis of two interference-free pilots (each at 10 dB), and the approximation of the perfect CSI density (0.41 in this case) multiplied by  $\left(1 - \frac{n_R - 1}{K + 1}\right)^{2/\alpha}$ . From the figure we see that estimation of  $\mathbf{h}_0$  does not significantly reduce density, the approximate density expression is reasonably accurate, and that choosing  $K$  on the order of 10 or 20 leads to a density reasonable close to the perfect CSI benchmark. Indeed, even if such estimation must be performed for every transmission, the overhead is reasonable in light of the fact that packets are typically on the order of hundreds of symbols.

## B. End-to-End Throughput

The transmission capacity quantifies per-hop performance, whereas end-to-end throughput depends on the range and rate of each transmission and the spatial intensity of such transmissions. Thus, a legitimate question is whether the linear scaling of transmission capacity translates into linear scaling of end-to-end throughput? In the context of the transport capacity [8], which is a widely accepted end-to-end metric, this question can be answered in the affirmative. Transport capacity is the product of rate and distance summed over all transmissions and thus is proportional to  $\lambda_\epsilon (1 - \epsilon) d \log_2(1 + \beta)$ , i.e., the product of the successful transmission density, per-hop distance, and per-hop rate. Since the transport capacity is linearly proportional to  $\lambda_\epsilon$ , the linear density scaling established in Theorems 1 and 2 also translates into linear scaling of the transport capacity.

Although increasing the density with  $n_R$  leads to linear scaling of the transport capacity, it is not *a priori* clear if the antennas should instead be used to increase the transmission rate and/or range. Based on the scaling results in Theorems 1 and 2, we see that  $\lambda_\epsilon d^2 \beta^{2/\alpha} \propto n_R$ . Thus, if the density and rate (i.e., SINR threshold  $\beta$ ) are kept constant, then the range can be increased at order  $d \propto \sqrt{n_R}$ . Alternatively the SINR threshold can be increased at order  $\beta \propto n_R^{2/\alpha}$ , which translates to increasing per-hop rate approximately as  $\frac{\alpha}{2} \log_2(1 + n_R)$ . Because transport capacity is proportional to  $\lambda_\epsilon d \log_2(1 + \beta)$ , using the receive antennas to increase per-hop range only increases transport capacity at order  $\sqrt{n_R}$  while increasing per-hop rate leads to an even poorer logarithmic increase (consistent with [13]). Therefore, the most efficient use of the receive array, from an end-to-end perspective, is to increase the density of simultaneous transmissions rather than the per-transmission rate or distance.

These points can be argued concretely within the framework of *expected forward progress* (EFP), a metric introduced in [10] that is defined as

$$\text{EFP} = \nu p \cdot \mathbb{E}[X_0] \cdot \log_2(1 + \beta), \quad (37)$$

where  $\lambda = \nu p$  is the density of transmitters subject to an ALOHA protocol where the entirety of nodes in the network (that have density  $\nu$ ) transmit with probability  $p$  and act as receivers (and hence relays) with probability  $1 - p$ , where  $p$  is a design parameter that is optimized offline. An opportunistic routing protocol is employed through which the successful receiving node (i.e., relay) offering the most geographic progress towards a defined destination direction is selected to forward each transmitter's packet, and the quantity  $\mathbb{E}[X_0]$  is the expected progress offered by such relay. The expected distance is inversely proportional to the transmitter density  $p$ , and by understanding how the optimum  $p$  changes with respect to  $n_R$  we can determine if receive antennas are more effectively used for increasing transmit density ( $p$  increasing rapidly with  $n_R$ ) or for increasing per-hop range ( $p$  approximately constant).

Fig. 8 contains plots of EFP versus transmission probability  $p$  for  $n_R = 1$  to  $n_R = 8$  and the optimum value of  $p$  is seen to increase approximately linearly with  $n_R$ , confirming that it is more effective to use the antennas to increase

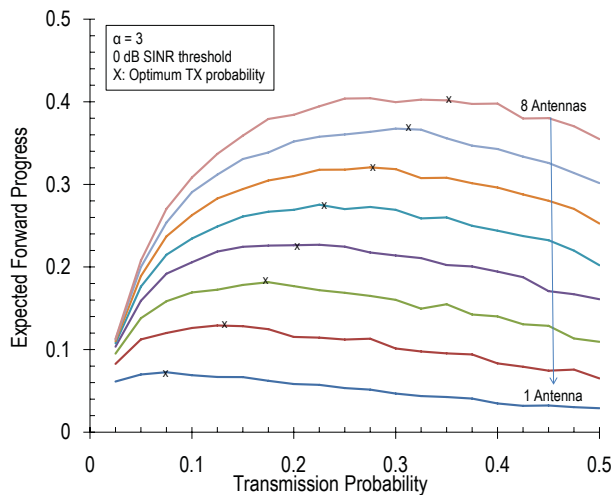


Fig. 8. Expected forward progress vs. transmission probability ( $p$ ) for  $n_R$  ranging from 1 (bottom) to 8 (top), for  $\alpha = 3$ ,  $\beta = 0$  dB, and no noise. The optimizing values of  $p$  are denoted with an X.

density. (Closer examination shows that the expected distance per communication, i.e. the per-hop range, is approximately constant with respect to  $n_R$  for the optimizing value  $p$ .) The plot also illustrates that the EFP itself is linear with  $n_R$ , confirming that linear scaling holds for multihop wireless networks. To see that increasing the rate rather than density also is suboptimal, if  $p$  is kept fixed to the small value of 0.075 when  $n_R = 8$  and the spectral efficiency  $\log_2(1 + \beta)$  is set to 4 (this value leads to the same expected per-hop range as the optimizing  $p$ ), then the resulting EFP is 0.3 instead of the 0.4 achievable if rate is kept fixed and  $p$  (i.e. the density) is increased.

### C. Effect of Interference Geometry

In this paper we have assumed a homogeneous Poisson distribution for the node locations, which is a realistic model if the users take up random locations and do not coordinate their transmissions. One might reasonably wonder, however, if the linear scaling result is an artifact of this model, since in a Poisson field the nearest interferers dominate and so interference cancellation might be far more profitable in this setup than in a more regular network. A “good” MAC protocol would seemingly space out the active transmitters at any instance, to avoid dominant interference. As a simple manifestation of such a MAC, we consider a regular network where the interferers take up positions on a square grid with edges of length  $1/\sqrt{\lambda}$  and study the performance of such a network through Monte Carlo simulation (the regular interferer spacing makes analysis very difficult). From Fig. 9 we see that a regular network allows for a larger density of simultaneous transmissions, but only by a constant factor that is independent of  $n_R$ . Based on this, we conjecture that the linear scaling result holds for any network geometry in which nodes are reasonably scattered in space.

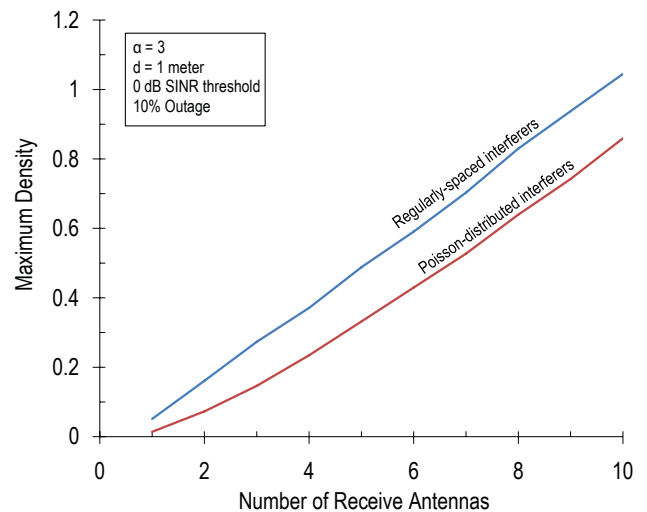


Fig. 9. The optimal density for both Regular and Poisson networks increases linearly with  $n_R$ , which regular networks having a slightly higher achievable density since nearby (dominant) interferers do not exist.

## V. CONCLUSION

The main takeaway of this paper is that very large throughput gains can be achieved in ad hoc networks using only receive antennas in conjunction with linear processing. In a point-to-point link receive antennas only provide array gain, which translates into a linear SNR and thus logarithmic rate increase (in the number of receive antennas). In an ad hoc network, however, receive antennas can also be used to cancel interference and this possibility turns out to yield much more significant benefits. In particular, the main result of the paper showed that using receive antennas to cancel interference and obtain some array gain allows the density of simultaneous transmissions to be increased linearly with the number of receive antennas when nodes transmit using only a single antenna. This result only requires channel state information at the receiver, which can be reasonably estimated, and the conclusion was seen to be robust to the particular interferer geometry.

From an end-to-end perspective, this linear increase in the density of simultaneous transmissions naturally translates into a linear increase in network-wide throughput. In addition, our analysis showed that receive antennas are in fact best utilized by increasing the density of simultaneous transmissions rather than increasing the per-hop rate or range. Finally, although the single-transmit/multiple-receive antenna setting may appear artificial, subsequent work on this model has shown that it can be detrimental to employ multiple transmit antennas when channel state information is not available to the transmitter [24], [29]. As a result, the single transmit/multiple receive antenna setting is indeed very relevant.

## APPENDIX A PROOF OF LEMMA 1

By the definition of  $\mathbf{v}_0$ , the quantity  $|\mathbf{v}_0^\dagger \mathbf{h}_0|^2$  is the squared-norm of the projection of vector  $\mathbf{h}_0$  on  $\text{Null}(\mathbf{h}_1, \dots, \mathbf{h}_k)$ . This nullspace is  $n_R - k$  dimensional (with probability one) by basic

properties of iid Gaussian vectors and is independent of  $\mathbf{h}_0$  by the independence of the channel vectors, and thus  $|\mathbf{v}_0^\dagger \mathbf{h}_0|^2$  is  $\chi_{2(n_{\text{R}}-k)}^2$  [25]. The second property holds by the definition of the PZF- $k$  receiver. To prove the third property, note that  $\mathbf{v}_0$  depends only on  $\mathbf{h}_0, \mathbf{h}_1, \dots, \mathbf{h}_k$  and thus is independent of  $\mathbf{h}_{k+1}, \mathbf{h}_{k+2}, \dots$ . Because the distribution of each channel vector is rotationally invariant (i.e., the distributions of  $\mathbf{W}\mathbf{h}_i$  and  $\mathbf{h}_i$  are the same for any unitary matrix  $\mathbf{W}$ ), we can perform a change of basis such that  $\mathbf{v}_0 = [1 \ 0 \ \dots \ 0]^T$ . After this change of basis, each  $\mathbf{v}_0^\dagger \mathbf{h}_i$  (for  $i \geq k+1$ ) is simply equal to the first component of  $\mathbf{h}_i$ . As a result  $\mathbf{v}_0^\dagger \mathbf{h}_{k+1}, \mathbf{v}_0^\dagger \mathbf{h}_{k+2}, \dots$  are iid complex Gaussians; thus the squared norms are iid exponentials, and furthermore these terms are independent of  $S$ .

#### APPENDIX B PROOF OF LEMMA 2

First we have

$$\mathbb{E} \left[ \sum_{i=k+1}^{\infty} |X_i|^{-\alpha} H_i \right] = \sum_{i=k+1}^{\infty} \mathbb{E} [|X_i|^{-\alpha} H_i] \quad (38)$$

$$= \sum_{i=k+1}^{\infty} \mathbb{E} [|X_i|^{-\alpha}] \quad (39)$$

where  $\mathbb{E} [|X_i|^{-\alpha} H_i] = \mathbb{E} [|X_i|^{-\alpha}] \mathbb{E} [H_i] = \mathbb{E} [|X_i|^{-\alpha}]$  due to the independence of  $|X_i|$  and  $H_i$  and the fact that  $\mathbb{E} [H_i] = 1$ . Because  $|X_1|^2, |X_2|^2, \dots$  are a 1-D PPP with intensity  $\pi\lambda$ , random variable  $\pi\lambda|X_i|^2$  is  $\chi_{2i}^2$  and thus has PDF  $f(x) = \frac{x^{i-1} e^{-x}}{\Gamma(i)}$ . Therefore

$$\mathbb{E} \left[ (|X_i|^2)^{-\alpha/2} \right] = (\pi\lambda)^{\frac{\alpha}{2}} \int_0^{\infty} x^{-\alpha/2} \frac{x^{i-1} e^{-x}}{(i-1)!} dx \quad (40)$$

$$= (\pi\lambda)^{\frac{\alpha}{2}} \frac{\Gamma(i - \frac{\alpha}{2})}{\Gamma(i)}. \quad (41)$$

This quantity is finite only for  $i > \frac{\alpha}{2}$ , and thus the expected power from the nearest uncanceled interferer is finite only if  $k+1 > \frac{\alpha}{2}$ .

To reach the upper bound, we use the following inequality from [12] (which is derived using Kershaw's inequality to the gamma function):

$$\frac{\Gamma(i - \frac{\alpha}{2})}{\Gamma(i)} < \left( i - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{-\frac{\alpha}{2}} \quad (42)$$

where  $\lceil \cdot \rceil$  is the ceiling function and we require  $i > \lceil \frac{\alpha}{2} \rceil$ . Therefore

$$\sum_{i=k+1}^{\infty} \frac{\Gamma(i - \frac{\alpha}{2})}{\Gamma(i)} < \sum_{i=k+1}^{\infty} \left( i - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{-\frac{\alpha}{2}} \quad (43)$$

$$\leq \int_k^{\infty} \left( x - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{-\frac{\alpha}{2}} dx \quad (44)$$

$$= \left( \frac{\alpha}{2} - 1 \right)^{-1} \left( k - \left\lceil \frac{\alpha}{2} \right\rceil \right)^{1 - \frac{\alpha}{2}}, \quad (45)$$

where the inequality in the second line holds because  $x^{-\frac{\alpha}{2}}$  is a decreasing function.

#### APPENDIX C PROOF OF THEOREM 2

The outage upper bound is obtained by keeping the interference contribution of only the nearest  $l$  uncanceled interferers in (21) and applying Markov's inequality to the *success* probability:

$$1 - \mathbf{P}_{\text{out}}^{\text{mmse}}(\lambda) \stackrel{(a)}{\leq} \mathbb{P} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{\sum_{i=n_{\text{R}}}^{\infty} |X_i|^{-\alpha} H_i} \geq \beta \right] \quad (46)$$

$$\stackrel{(b)}{\leq} \mathbb{P} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{\sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} |X_i|^{-\alpha} H_i} \geq \beta \right] \quad (47)$$

$$\stackrel{(c)}{\leq} \mathbb{P} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{|X_{n_{\text{R}}-1+l}|^{-\alpha} \sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} H_i} \geq \beta \right] \quad (48)$$

$$\stackrel{(d)}{\leq} \frac{1}{\beta} \mathbb{E} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{|X_{n_{\text{R}}-1+l}|^{-\alpha} \sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} H_i} \right] \quad (49)$$

where (a) follows from (21), (b) because decreasing the interference increases the SIR and thus the success probability, (c) because  $|X_i|$  are increasing in  $i$  and the function  $(\cdot)^{-\alpha}$  is decreasing, and (d) is due to Markov's inequality. By the independence of the various random variables:

$$\begin{aligned} & \mathbb{E} \left[ \frac{d^{-\alpha} \|\mathbf{h}_0\|^2}{|X_{n_{\text{R}}-1+l}|^{-\alpha} \sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} H_i} \right] \\ &= \mathbb{E} [\|\mathbf{h}_0\|^2] \mathbb{E} \left[ \frac{d^{-\alpha}}{\sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} H_i} \right] \mathbb{E} [|X_{n_{\text{R}}-1+l}|^{\alpha}] \\ &= \frac{n_{\text{R}}}{l-1} (\pi d^2 \lambda)^{-\alpha/2} \frac{\Gamma(n_{\text{R}}-1+l+\frac{\alpha}{2})}{\Gamma(n_{\text{R}}-1+l)}. \end{aligned} \quad (50)$$

where we have used the fact that  $\|\mathbf{h}_0\|^2 \sim \chi_{2n_{\text{R}}}^2$ ,  $\sum_{i=n_{\text{R}}}^{n_{\text{R}}-1+l} H_i$  is the sum of  $l$  iid exponentials and thus is  $\chi_{2l}^2$ , and  $|X_{n_{\text{R}}-1+l}|^2 \sim \frac{1}{\pi\lambda} \chi_{2(n_{\text{R}}-1+l)}^2$ .

By applying Kershaw's inequality, which states  $\Gamma(x+1)/\Gamma(x+s) < \left( x - \frac{1}{2} + \sqrt{s+1/4} \right)^{1-s} \forall x > 0$  and  $0 < s < 1$ , and the property  $\Gamma(x+1) = x\Gamma(x)$ , we have:

$$\frac{\Gamma(n_{\text{R}}-1+l+\frac{\alpha}{2})}{\Gamma(n_{\text{R}}-1+l)} \leq \left( n_{\text{R}}-1+l+\frac{\alpha}{2} \right)^{\alpha/2}. \quad (51)$$

Substituting (50) and (51) into (49) yields:

$$\mathbf{P}_{\text{out}}^{\text{mmse}}(\lambda) \geq 1 - \frac{1}{\beta} \frac{n_{\text{R}}}{l-1} (\pi\lambda)^{-\alpha/2} \left( n_{\text{R}}-1+l+\frac{\alpha}{2} \right)^{\alpha/2} \quad (52)$$

By choosing  $l = n_{\text{R}} + 1$  we obtain the desired outage probability lower bound, and by setting this bound to  $\epsilon$  and solving we get the density upper bound.

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