# MULTI-DIMENSIONAL <br> PIECEWISE POLYNOMIAL <br> CURVE FITTING 

by

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## ABSTRACT

Any set of hyperplanes partitions $R^{n}$ into polygonal pieces. Linearly partitioned piecewise polynomials (LPPP) result when polynomials on each of these pieces are put together in a smooth manner. There is presented an LPPP canonical form which leads to computationally feasible multi-dimensional piecewise polynomial curve fitting.
I. INTRODUCTION

The problem of explicitly describing an $n$-dimensional surface $z=F(x), x$ an $n$-vector, which approximately fits some data points $\left(x_{k}, z_{k}\right)$ arises in mathematical modeling and data analysis with computers. The data points may be the results of a computer simulation (or experimental results, surveys, economic data, etc.) of a function whose values must play a role (where smoothness is required) in a more involved simulation. One way of attempting to describe such a surface is to piece together polynomial functions. This paper is concerned with the process of selecting a piecewise polynomial function which approximates the data.

Section II first describes what the pieces of $R^{n}$ are to be like. Any set of hyperplanes naturally partitions $\mathrm{R}^{\mathrm{n}}$ into polygons. Then the concept of a linearly partitioned piecewise polynomial of degree $s$ is defined. This corresponds to having polynomials of degree $\leq$ s defined on each of the polygons in such a way that the resulting function in $R^{n}$ possesses continuous derivatives of order s-l. Non-degeneracy of a partition is also defined. Section III shows that, for non-degenerate partitions, the piecewise polynomial function must take on an especially simple canonical form. Section IV contains a few comments on how the canonical form is useful in finding the desired surface $z=F(x)$.
II. DEFINITIONS:

A linear partition of $R^{n}$ is a set of vectors $\left(a_{1} ; b_{1}\right) \ldots \ldots$, $\left(a_{m}, b_{m}\right)$, where the $a_{i}^{\prime} s$ are $n$-component vectors and the $b_{i}$ 's are real numbers. Associated with a linear partition are the closed polygonal regions, $R_{I}, I$ a subset of $1, \ldots, m_{\text {, }}$ where $R_{I}$ is defined by

$$
R_{I}=\left\{x \mid a_{i} \cdot x \geq b_{i}, i \varepsilon I, \quad a_{i} \cdot x \leq b_{i}, i \notin I\right\}
$$

and is non-empty.
A linear partition is non-degenerate if for every linearly dependent $a_{i_{1}}, \ldots, a_{i_{k}}$ we have

$$
\operatorname{rank}\binom{a_{i_{1}}}{a_{i_{k}}}<\operatorname{rank}\binom{a_{i_{1}} \cdot b_{i_{1}}}{a_{i_{k}} \cdot b_{i_{k}}}
$$

Nondegeneracy says the intersection of any $k$ of the hyperplanes $a_{i} \cdot x=b_{i}$, which partition $R^{n}$, must have dimension $n-k$. In particular, the intersection of any $n+1$ of these hyperplanes is empty.

Two regions $R_{I}$ and $R_{J}$ are called k-neighbors if either $I=J U\{k\}$ or $J=I U\{k\}$ 。

Linear partitions possess some properties which will be of use in the following section. We present them here.

Proposition I: For any $R_{I}$ and $R_{J}$, there exist a sequence $R_{I_{1}} \ldots \ldots R_{I_{\ell}}$ such that $R_{I_{j}}$ and $R_{I_{j+1}}$ are $k_{j}$-neighbors, $, j=1, \ldots, \ell-1$ and $R_{I_{1}}=R_{I}, R_{I_{\ell}}=R_{J}$.

Proof: Let $U$ be the union of all $R_{K}{ }^{\prime} s$ which can be connected to $R_{I}$ by such chains. $U$ is the union of a finite number of closed sets, so it is closed. We need only show that U is an open set.

Pick $x \in U$. Set $U_{x}$ equal to the union of all $R_{I}$ 's which contain $x$. $x$ is not contained in the union of the $R_{I}$ 's not containing $x$, which is a closed set containing $R^{n}-U_{x}$. So $x$ is an interior point of $U_{x}$. We will show that $U_{x} \subset U$. Suppose $x \in R_{K} \subset U$ : Pick $R_{L} \subset U_{X}$ : Since $x \in R_{K} \cap R_{L}$, we must have $a_{i} \cdot x=b_{i}$, $\forall i \varepsilon(K-L) \cup(L-K)$. Suppose $K-L=\left\{i_{1}, \ldots, i_{r}\right\}, L-K=\left\{j_{1}, \ldots, j_{s}\right\}$. Then set $I_{j}=K-\left\{i_{1} \ldots \ldots, i_{j}\right\}, j=0, \ldots, r$, and $I_{r+j}=\left(K-\left\{i_{1}, \ldots, i_{r}\right\}\right)$
$\left\{j_{1}, \ldots . j_{j}\right\}, j=1, \ldots, s, \quad$ The $R_{I_{j}}, j=0, \ldots, r+s$ are nonempty (because each contains $x$ ) and $R_{I_{0}} \ldots \mathcal{O}_{I_{r+s}}$ is a chain of neighbors connecting $R_{K}$ to $R_{L}$. Thus; $R_{L} C U$, ${ }^{\prime} R_{L} \subset U_{X}$. And $x$ being an interior point of $U_{x} \subset U$, is an interior point of $U$. $U$ being open and closed must equal $R^{n}$. So $R_{J} \subset U$.

Proposition 2: If the partition is non-degenerate, then $R_{I} \neq \phi$ implies $R_{I}$ contains interior points i.e. $R_{I}$ has dimension $n$. Proof: Suppose $R_{I} \neq \phi$. Pick $\hat{x} \in R_{I}$. Let $J=\left\{i \mid \hat{a}_{i} \cdot \hat{x}=\dot{b}_{i}\right\}^{\prime}$. Set $x_{J}=\hat{x}$. If, for some $k \varepsilon J$, there exists $x \in R_{I}$ such that $a_{i} \cdot x=b_{i}$ iff i\&J-\{k\}, then replace $J$ by $J-\{k\}$ and $x_{J}$ by this new $x$. If this procedure eventually exhausts $J$, then the last $x$ is an interior point of $R_{I}$.

Suppose this procedure does not exhaust J. Then it eventually occurs that
a) $a_{j} \cdot x_{J}=b_{j} \cdot j \varepsilon J$ and
b) for $k \varepsilon J / \cap_{I} \min a_{k} \cdot x$
subject to $a_{j} \cdot x=b_{j}, j \varepsilon J-\{k\}$
is equal to $a_{k} \cdot x_{J}=b_{k}$ '
c) for $k \in J-I$
$\max a_{k} \cdot x$
subject to $a_{j} \cdot x=b_{j}, j \varepsilon J-\{k\}$
is equal to $a_{k} \cdot x_{J}=b_{j}$.
Since at least one of $J \cap I$ and $J-I$ is nonempty, we must have that for some $k \varepsilon J$, there exist multipliers of $\alpha_{j}, j \varepsilon J$ such that $\sum_{\substack{j \in J \\ j \neq k}} \alpha_{j} a_{j}=a_{k}$. Thus, the set of $a_{j}{ }^{\prime} s, j \varepsilon J$ is linearly dependent and $a_{j}{ }^{\bullet} x_{j}=b_{j}{ }^{\circ} j \varepsilon J$. This contradicts the non-degeneracy of the partition. Thus, J must be exhausted and the proposition is proved.

Proposition 3: $a_{j} \cdot x=b_{j}$
for some $x \in R_{I}$ implies that
$\operatorname{dim}\left\{x \mid a_{j} \cdot x=b_{j}\right\} \cap R_{I}=n-1$

Proof: Use an argument like that used in the proof of proposition 2 with the additional requirement that $j$ never be removed from $J$.

A function $F$ is said to be alinearly partitioned piecewise polynomial (LPPP) of degrees iff there is a linear
partition of $R^{n}$ such that

1) Fis continuously differentiable of order s-1.
2) $\boldsymbol{F}_{I}\left(F_{\text {restricted }}\right.$ to $\left.R_{I}\right)$ is a polynomial of degree $\leq s, \forall R_{I}$.
3) $E_{I}$ is a polynomial of degree $s$ for some $R_{I}$.
4) If $R_{I}$ and $R_{J}$ are $k$-neighbors then

$$
F_{I}=F_{J} \text { in } R_{I} \cap_{\mathbb{R}_{J}}
$$

Fis a basic LPPP of degree $s$ iff

$$
F(x)=\operatorname{sgn}\left(a_{1} \cdot x-b_{1}\right) \alpha_{1}\left(a_{1} \cdot x-b_{1}\right)^{s}
$$

for some nonzero $\alpha_{1}$, some $n$-vector $a_{1}$, and some real number $b_{1}$. Notice that the above function is an LPPP because with the linear partition consisting of the single vector ( $a_{1}, b_{1}$ )

1) $\frac{\cdots \partial^{s-1}}{\partial x_{1}^{j 1} \ldots \partial x_{n}^{j}} \cdot\left(\operatorname{sgn}\left(a_{1} \cdot x-b_{1}\right) \alpha_{1}\left(a_{1} \cdot x-b\right)^{s}\right)$ $=s!a_{1}^{j_{1}} \ldots a_{n}^{j_{n}}\left(a_{1} \cdot x-b_{1}\right) \alpha_{1}\left(a_{1} \cdot x-b_{1}\right)$
2) +3) $\quad F_{\Phi}(x)=-\alpha_{i}\left(a_{1} \cdot x-b_{1}\right)^{S}$, and

$$
F_{\{1\}}(x)=\alpha_{1}\left(a_{1} \cdot x-b_{1}\right)^{s}
$$

4) $F_{\phi}(x)=F_{\{1\}}(x)=0$, if $x \in R_{\phi} \cap R_{\{1\}}$

## III. CANONICAL FORM:

An LPPP (with partition $\left(a_{1}, b_{1}\right), \ldots,\left(a_{m}, b_{m}\right)$ ) is said to be in canonical form, if it is expressed as

$$
F(x)=F_{0}+\sum_{i=1}^{m} F_{i}
$$

where $F_{0}$ is a polynomial and $F_{i}$ is a basic LPPP with partition ( $\mathrm{a}_{\mathrm{i}}, \mathrm{b}_{\mathrm{i}}$ ).

Theorem: If $F$ is an LPPP of degree $s$, with non-degenerate partition, then $F$ can be expressed in canonical form. Proof: We first prove a few lemmas.

Lemma 1: If $R_{I}$ and $R_{J}$ are $k$-neighbors, then

$$
F_{I}(x)=F_{J}(x) \forall x 3 a_{k} \cdot x=b_{k}
$$

Proof: By proposition 3, $\operatorname{dim} R_{I} \cap R_{J}=n-1$.
Since $F_{I}(x)=F_{J}(x) \quad x \in R_{I} \cap_{R_{J}}$ an $n-1$ dimensional subset of $\left\{x \mid a_{k} \cdot x=b_{k}\right\}$, and $F_{I}$ and $F_{J}$ are polynomials, we must have $F_{I}(x)=F_{J}(x), \quad \forall x 3 a_{k} \cdot x=b_{k}$.

Lemma 2: If $R_{I}$ and $R_{I-\{k\}}$ are non-empty, then there exists a real number $\alpha_{k}^{I}$ such that $F_{I}(x)-F_{I-\{k\}}(x)=2 \alpha_{k}^{I}\left(a_{k} \cdot x-b_{k}\right)^{s}$. Proof: Let $Y_{1}, \ldots, Y_{n}$ be a new orthogonal coordinate system which is a linear transformation of the $x_{1}, \ldots x_{n}$ coordinate system such that $y_{1}=a_{k}{ }^{\bullet} x-b_{k}$.

Now,
i) $F_{I}\left(0, Y_{2}, \ldots, Y_{n}\right)=F_{I-\{k\}}\left(0, Y_{2}, \ldots, Y_{n}\right)$, follows from lemma 1 , and
ii)

for $i_{1}+\ldots+i_{n}=r \leq s-1$, follows from the fact that $F$ is continuously differentiable of order sol.

Expansion of $F_{I}$ and $F_{I-\{k\}}$ in Taylor series about any point in $R_{I} \cap_{R_{I}}-\{k\}$ and use of i) and ii) gives

$$
F_{I}(y)-F_{I-\{k\}}(y)=\alpha y_{1}^{s} .
$$

Taking $\alpha_{k}^{I}=\frac{1}{2} \alpha$, we get

$$
F_{I}(x)-F_{I-\{k\}}(x)=2 \alpha_{k}^{I}\left(a_{k} \cdot x-b_{k}\right)^{s}
$$

Lemma 3: If $k, \ell \in I$ and $R_{I}$ and $R_{I-\{\ell\}}$ satisfy
$R_{I} \cap_{R_{I-\{\ell}} \cap\left\{x \mid a_{k} \cdot x=b_{k}\right\} \neq \phi_{r}$ then $\alpha_{k}^{I}=\alpha_{k}^{I-\{\ell\}}$.

Proof: $\quad a_{\ell}{ }^{\circ} \mathrm{x}=\mathrm{b}_{\ell}$ implies that

$$
F_{I}(x)=F_{I-\{\ell\}}(x) \text { and } F_{I-\{k\}}(x)=F_{I-\{k ; \ell\}}(x)
$$

by lemma 1. Thus,

$$
F_{I}(x)-F_{I-\{k\}}(x)=F_{I-\{\ell\}}(x)-F_{I-\{k, \ell\}}(x) \forall x \cdot 3 a_{\ell} \cdot x=b_{\ell}
$$

That is, $2 \alpha_{k}^{I}\left(a_{k} \cdot x-b_{k}\right)=2 \alpha_{k}^{I-\{\ell\}}\left(a_{k} \cdot x-b_{k}\right), \forall x \geqslant a_{\ell} \cdot x=b_{\ell}$. Since $a_{k} \neq a_{\ell}$ (because $a_{k} \cdot x=b_{k}$ and $a_{\ell} \cdot x=b_{\ell}$ intersect and the partition is non-degenerate), we must have $\alpha_{k}^{I}=\alpha_{k}^{I-\{\ell\}}$. Lemma 4: For fixed $k$, all $\alpha_{k}^{I \prime s}$ are the same.

Proof: $\alpha_{k}^{I_{k}} s$ are only defined for $I \varepsilon d$ the set of those I's for which $k \in I$ and $R_{I} \cap\left\{x \mid a_{k} \cdot x=b_{k}\right\} \neq \phi$. Lemma 3 shows that if any two of these $R_{I}{ }^{\prime} s$ are $\ell$-neighbors then the $\alpha_{k}^{I}{ }^{I} s$ are equal. It follows that if two of these, say $R_{I_{1}}$ and $R_{I_{2}}$, are connected by a chain $R_{I_{1}}, \ldots, R_{I_{h}}, I_{j} \varepsilon \& j=1, \ldots, h$ where $R_{I}$. and $R_{I_{j+1}}$ are $\ell_{j}$-neighbors, $j=1, \ldots, h-1$, then $\alpha_{k}^{I_{1}}=\alpha_{k}^{I_{n}}$. An argument like that used in the proof of proposition 1 (with $\left\{x \mid a_{k} \cdot x=b_{k}\right\}$ replacing $R^{n}$ ) shows that such a chain exists for any $I_{0}$ Jed. Thus, $\alpha_{k}^{I}=\alpha_{k}^{J}$ for any I, Jed\& Henceforth, the superscript will be dropped and we will use $\alpha_{k}$.

We now turn to the proof of our theorem. Pick a nonempty $R_{I} \cdot$ Set $F_{0}(x)=F_{I}(x)-\sum_{i \varepsilon I} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}+\sum_{i \notin I} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}$ and $F_{i}(x)=\operatorname{sgn}\left(a_{i} \cdot x-b_{i}\right) \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}, i=1, \ldots, m$. We claim that $F=F_{0}+\sum_{i=1}^{m} F_{i} . \quad$ clearly, $F_{I}(x)=F_{0}(x)+\sum_{i=1}^{m} F_{i}(x) \quad \forall x \in R_{I}$. Pick any other $R_{J}$. By proposition 1 there exist $R_{I_{1}} \ldots, R_{I_{\ell}}$ such that $R_{I_{1}}=R_{I^{\prime}} R_{I_{\ell}}=R_{J}$ and $R_{I_{j}}$ and $R_{I_{j+1}}$ are $k$-neighbors , $j=1, \ldots . . \quad$ Using lemmas 2 and 4 we get

$$
F_{J}(x)=F_{I}(x)+\sum_{k_{i} \varepsilon L^{+}} 2 \alpha_{k_{j}}\left(a_{k_{j}} \cdot x-b_{k_{j}}\right)^{s}-\sum_{k_{i} \varepsilon L^{-}} 2 \alpha_{k_{j}}\left(a_{k_{j}} \cdot x-b_{k_{j}}\right)^{s}
$$

where $L^{+}=\left\{k_{j} \mid I_{j+1}=I_{j}+\left\{k_{j}\right\}\right\}$ and $L^{-}=\left\{k_{j} \mid I_{j+1}=I_{j}-\left\{k_{j}\right\}\right\}$. Notice that
i) if i $\varepsilon I \cap J$, then $i$ appears in $L^{+}$the same number of times as it appears in $\mathrm{L}^{-}$,
ii) if iعI-J, then $i$ appears in $L^{+}$one time less often than it appears in $\mathrm{L}^{-}$, and
iii) if iعJ-I, then $i$ appears in $L^{+}$one more time than it appears in $L^{-}$.
i), ii) and iii) lead to
(1) $F_{J}(x)=F_{I}(x)+\sum_{i \varepsilon J-I} 2 \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}-\sum_{i \varepsilon I-J} 2 \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}$.

As noted above

$$
F_{I}(x)=F_{0}(x)+\sum_{i \in I} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}-\sum_{i \not 又 I} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}
$$

Substituting in (1) gives

$$
\begin{aligned}
& F_{J}(x)=F_{0}(x)+\sum_{i \in J} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s}-\sum_{i \nless J} \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)^{s} \\
& =F_{0}(x)+\sum_{i=1}^{m} F_{i}(x), \forall x \in R_{J}
\end{aligned}
$$

IV. CURVE FITTING:

Suppose we have data points $\left(x_{i}, y_{i}\right)$, $i=1, \ldots, r$, where $x_{i}=\left(x_{i j}, \ldots, x_{i n}\right)$, and we wish to determine a function
$F(x)$ such that $F\left(x_{i}\right) \quad \approx_{i}, i=1, \ldots, r$. We are interested in synthesizing $F$ as a piecewise polynomial function of degree s possessing continuous derivatives of order s-1. We could concentrate on functions of the form

$$
F(x)=F_{0}(x)+\sum_{i=1}^{m} F_{i}(x)
$$

where $F_{0}(x)$ is a polynomial of degree $s$ and $F_{i}(x)=\operatorname{sgn}\left(a_{i} \cdot x-b_{i}\right) \alpha_{i}\left(a_{i} \cdot x-b_{i}\right)$. The planes $\dot{a}_{i} \cdot x=b_{i}$ partition $R^{n}$. into regions over which the pieces of $F$ are defined.

According to section III, if the partition is non-degenerate then $F_{0}+\sum_{i=1}^{m} F_{i}$ encompasses all possible piecewise polynomials of degree s possessing continuous derivatives of order s-l.

We wish to determine the coefficients in $F(x)=F_{0}(x)+\sum_{i=1}^{m} F_{i}(x)$ so that $F(x)$ gives the best fit to the data points. In the discussion below we take "best" to mean that the coefficients are chosen so that $\sum_{i=1}^{r}\left(F\left(x_{i}\right)-y_{i}\right)^{2}$ is minimized. (The comments in part 1 are also pertinent in the case when $\max _{i=1, \ldots \ldots, r}\left|F\left(x_{i}\right)-y_{i}\right|$ is minimized). There are two levels of complication.

1. Fixed Partition. We assume that the partition (the ( $a_{i}, b_{i}$ )'s) are predetermined. The parameters to be determined are the coefficients in $F_{0}(x)$ and the $\alpha_{i}{ }^{\prime} s$. These parameters occur linearly in $F(x)$. The problem is merely a linear regression problem.

The number of parameters is equal to $N(s, n)+m$, where $N(s, n)$ is the no. of coeff. in a polynomial of degree s in n-variables.

For example, any piecewise cubic in 3 dimensions which has a non-degenerate linear partition of $m$ planes; involves $14+m$ parameters. For $m=4$. such a piecewise cubic could have 15 pieces.
2. Free Partition. We assume that the number of planes is pre-determined but that their positions are to be chosen. The parameters to be determined are the coefficients in $F_{0}(x)$ and the $\alpha_{i}, a_{i_{j}}, b_{i}, i=1, \ldots, m, j=1, \ldots, n$. This is a problem in non-linear regression. Notice that the $F_{i}(x)$ possess continuous derivatives of order (s-l) with respect to the $\alpha_{i}, a_{i_{j}}$, and $b_{i}$. All higher order derivatives exist except for derivatives of $F_{i}\left(x_{k}\right)$ of order $s$ with respect to the $a_{i_{j}}, b_{i}$ when the data point $x_{k}$ satisfies $a_{i} \cdot x_{k}=b_{i}$. In such cases, it is reasonable to set the derivative equal to zero. (This last comment may be of interest when s = 0 or l.)

Nonlinear regression techniques, such as appear in [2], can be used on this type problem. The fact that second derivatives of $F$ with respect to the parameters are mostly zero may be of some use.

Once the $\left(a_{i}, b_{i}\right)$ 's are fixed, the other coefficients are uniquely determined by linear regression. Thus, the sum of squares of the deviations is a well-defined function of the $\left(a_{i}!b_{i}\right)$ !s. An unconstrained minimization technique,
such as appears in [1], can be used to find $a_{i_{j}}, b_{i}$, $i=1, \ldots, m, j=1, \ldots, n$ which minimize this objective function. If this involves too many parameters it may be useful to fix the $a_{i}{ }_{j}$ 's and only regress on the $b_{i}$ 's.

Suppose $s=3, n=3$, and $m=4$. The nonlinear regression problem would involve 30 parameters. The minimization problem would involve 16 parameters. The minimization problem in which the partition planes are predetermined except for translation, would involve 4 parameters.

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