

SWIN
BUR
* NE *

SWINBURNE
UNIVERSITY OF
TECHNOLOGY

Multimedia-unfriendly TCP Congestion Control and Home Gateway Queue Management

Lawrence Stewart^α, David Hayes^α,
Grenville Armitage^α, Michael Welzl^β, Andreas
Petlund^β

^α Centre for Advanced Internet Architectures (CAIA),
Swinburne University of Technology

^β Department of Informatics, University of Oslo



Outline



1. Introduction
2. Results
3. Conclusion



Detailed outline (section 1 of 3)

1. Introduction

2. Results

3. Conclusion

1. Introduction

- Motivation
- Not All Queues Are Created Equal
- TCP CC (Congestion Control) Recap
- Methodology



Motivation

- Gateway devices often have excessive drop-tail queue buffering^{1,2}
- TCP CC algorithms each interact differently with network buffers
- Real-time multimedia traffic sharing the gateway bottleneck experiences collateral damage: latency/jitter and loss
- How does choice of CC algorithm and queue implementation impact the inflicted collateral damage?

¹M. Claypool, R. Kinicki, M. Li, J. Nichols and H. Wu, “Inferring Queue Sizes in Access Networks by Active Measurement”, PAM Workshop '04

²M. Dischinger, A. Haeberlen, K. P. Gummadi and S. Saroiu, “Characterizing residential broadband networks”, SIGCOMM IMC '07



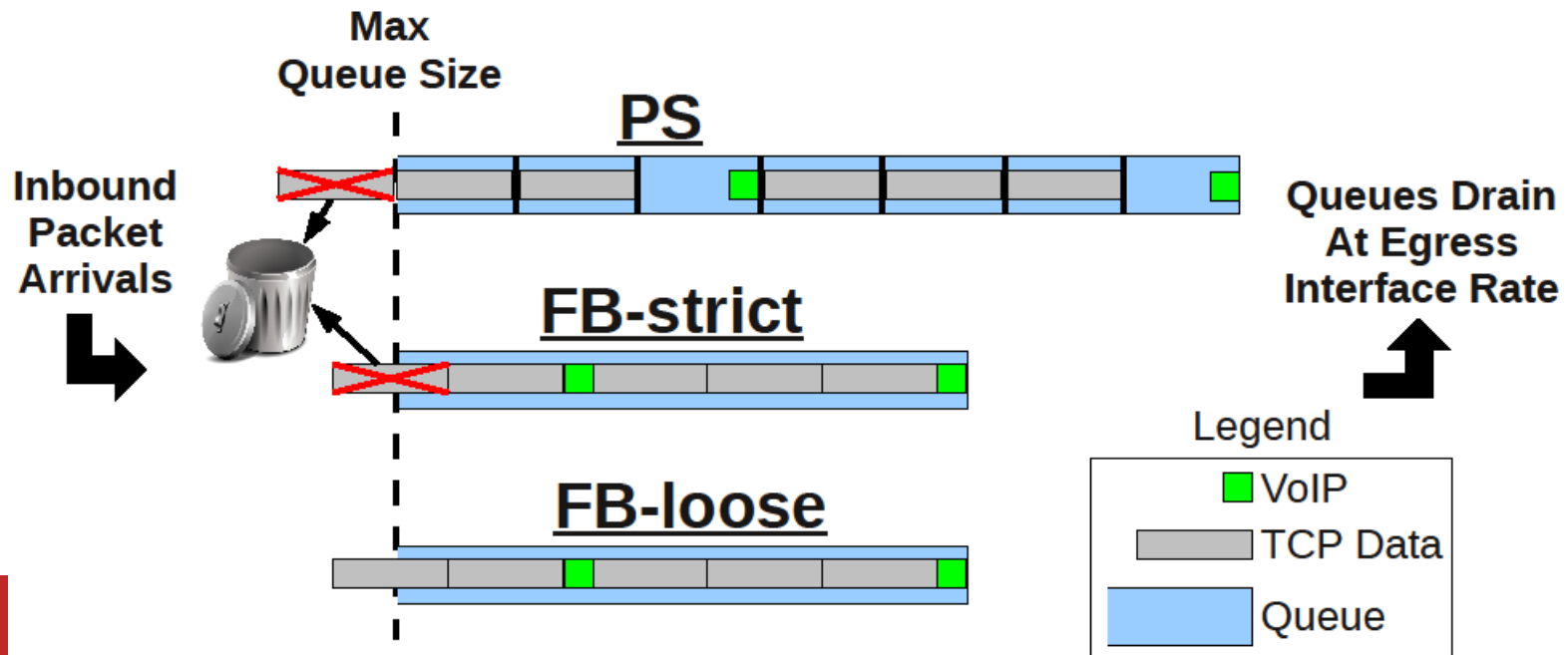
Not All Queues Are Created Equal

■ When is a drop-tail queue full?

Packet/Slot based drop if occupancy \geq qlen

Fixed-byte strict drop if occupancy + pkt \geq qlen

Fixed-byte loose drop if occupancy $>$ qlen





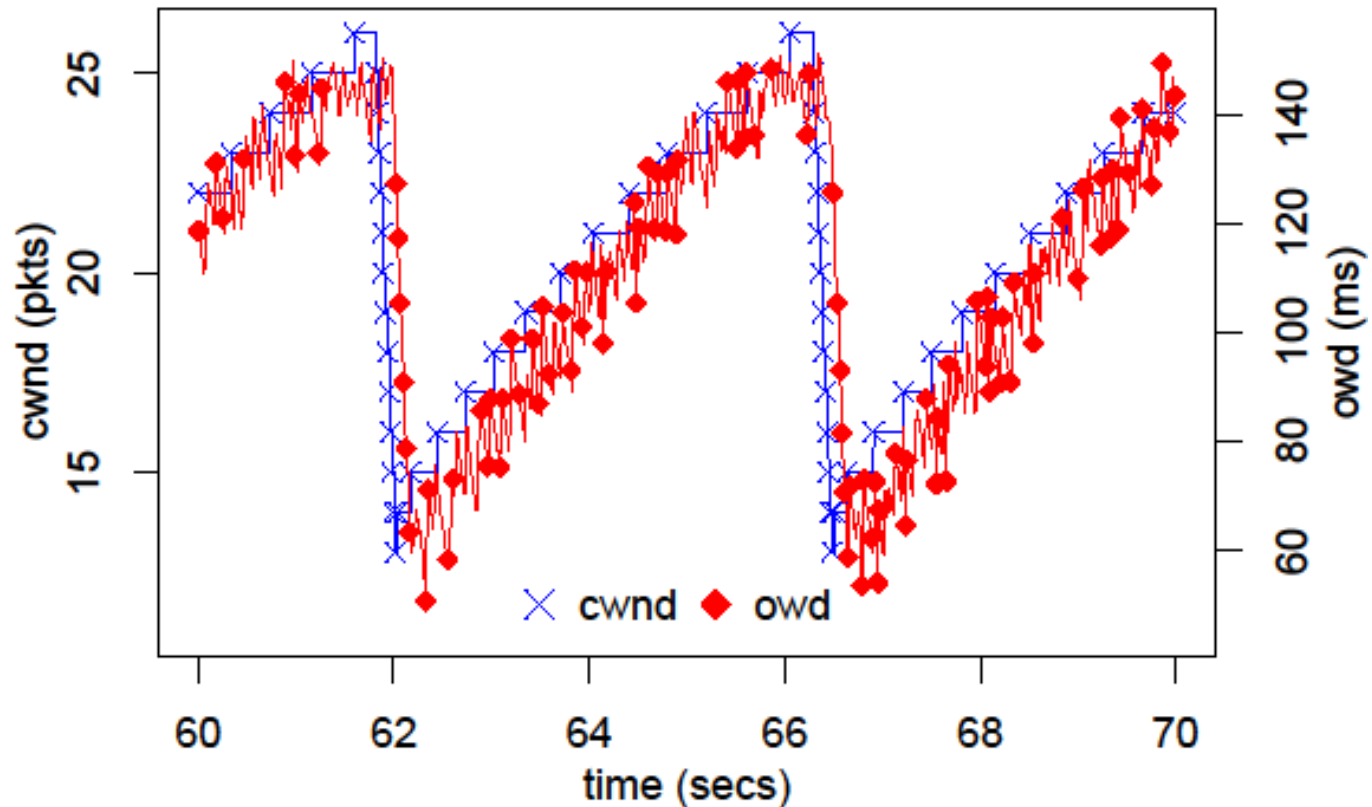
TCP Congestion Control Recap

- NewReno is ailing defacto standard (high latency/speed, wireless)
- Many new algorithm proposals: Compound, CUBIC, HTCP, ...
- BSD still uses NewReno
- Linux uses CUBIC
- Windows Vista / 7 have Compound TCP available for testing

TCP Congestion Control Recap



Downstream NewReno Behaviour



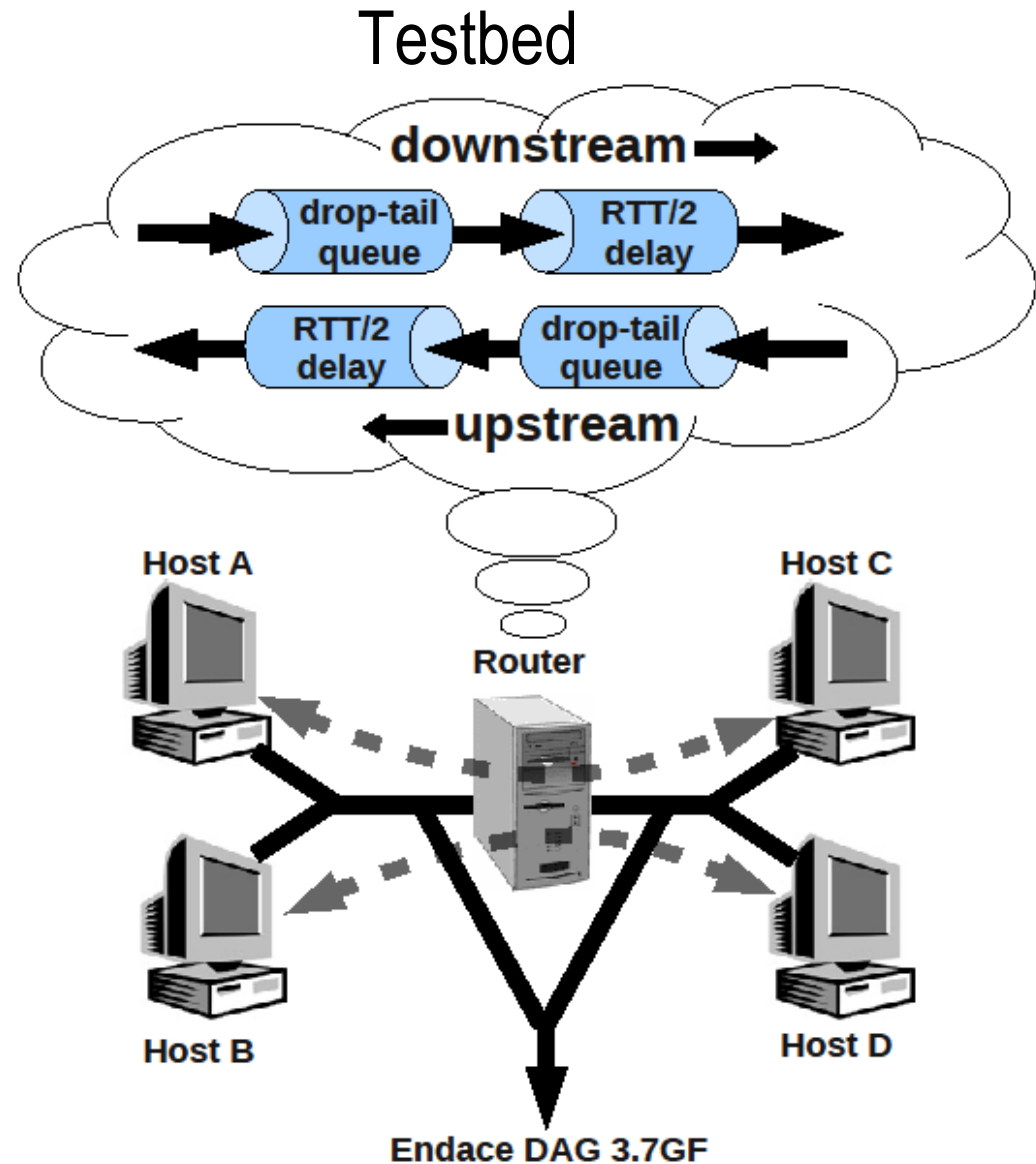
Congestion window oscillation coupled with induced queuing delay.

ADSL 1 type link: 1500/256 kbps, min 100ms RTT, 20 000 byte queue.

Methodology



- Testbed & NS-2
- 1500/256kbps & 24/1Mbps (ADSL 1 & 2 speeds)
- Byte-based [10k, 100k] & slot-based [7, 67] queues
- RTTs of 24, 50, 100 and 200ms
- Dummynet for bandwidth, queue size & RTT emulation
- Single bulk TCP flow A->C
- 186 byte, $N(\mu=20\text{ms}, \sigma=1\text{ms})$ CBR UDP flow B<->D



Detailed outline (section 2 of 3)



1. Introduction

2. Results

- Goodput
- Collateral Damage to Voice Stream – Latency
- Collateral Damage to Voice Stream - Loss

2. Results

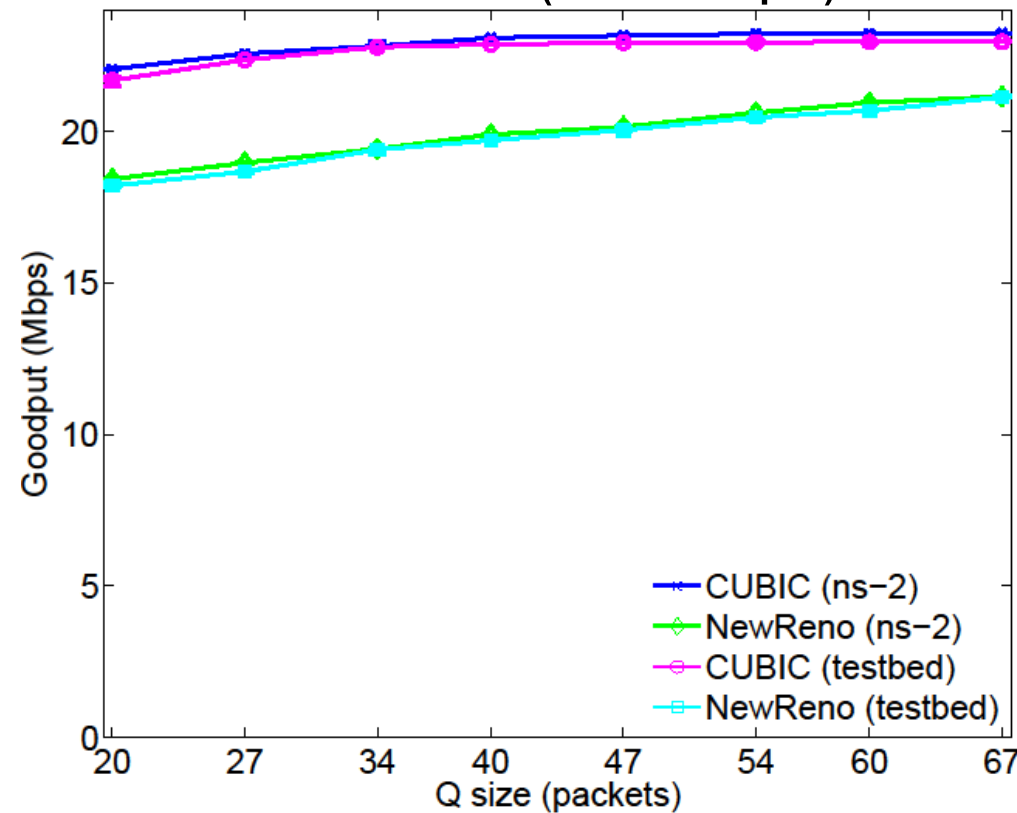
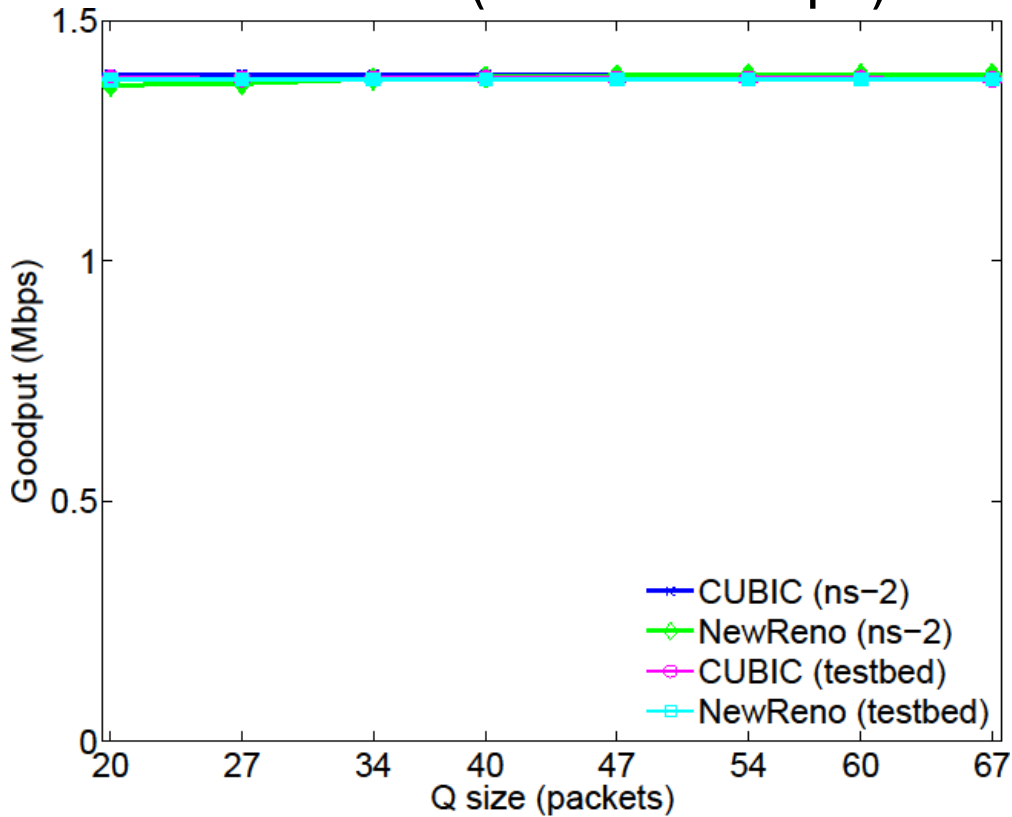
3. Conclusion

Goodput



ADSL 1 (1500/256kbps)

ADSL 2 (24/1Mbps)

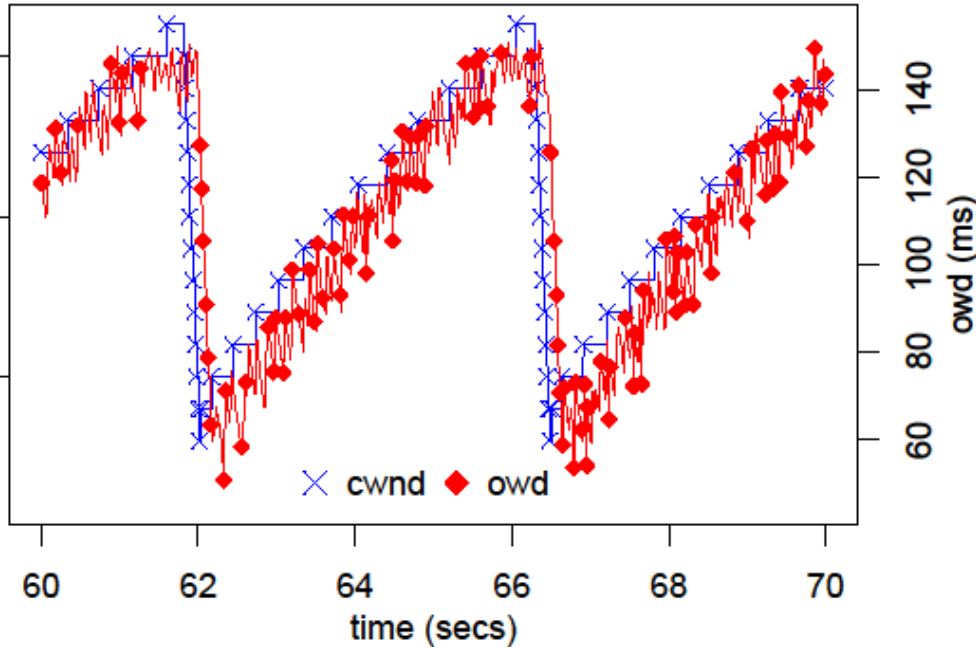


Downstream goodput vs PS queue length - measured (testbed) and simulated (NS-2), 1500/256 kbps & 24/1 Mbps links, 100ms min RTT.

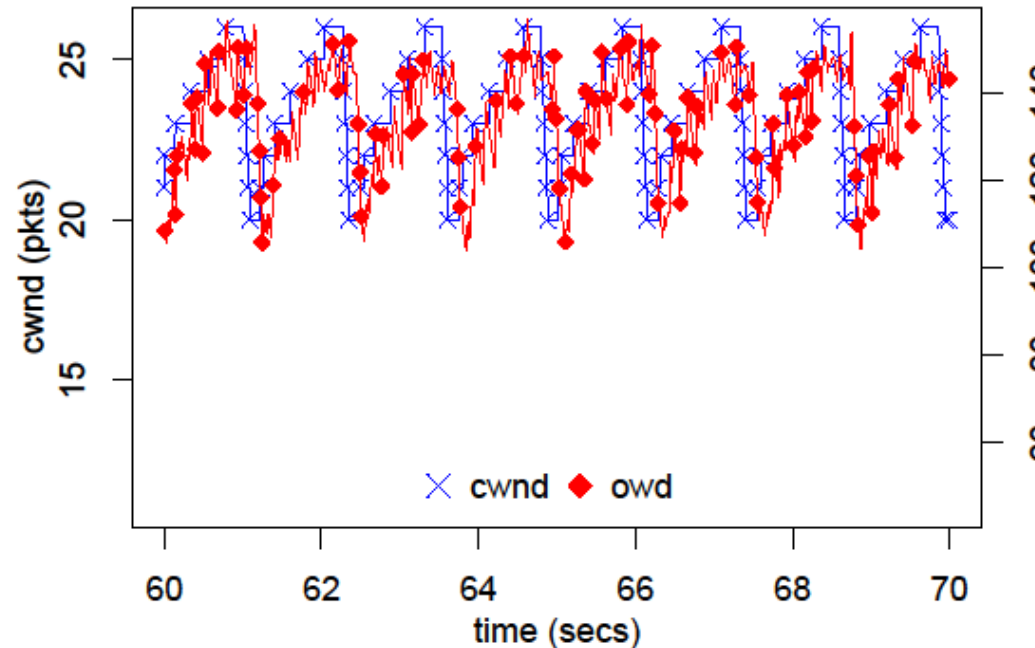
CUBIC Induces More OWD Than NewReno



Downstream NewReno



Downstream CUBIC



Congestion window oscillation coupled with induced queuing delay.

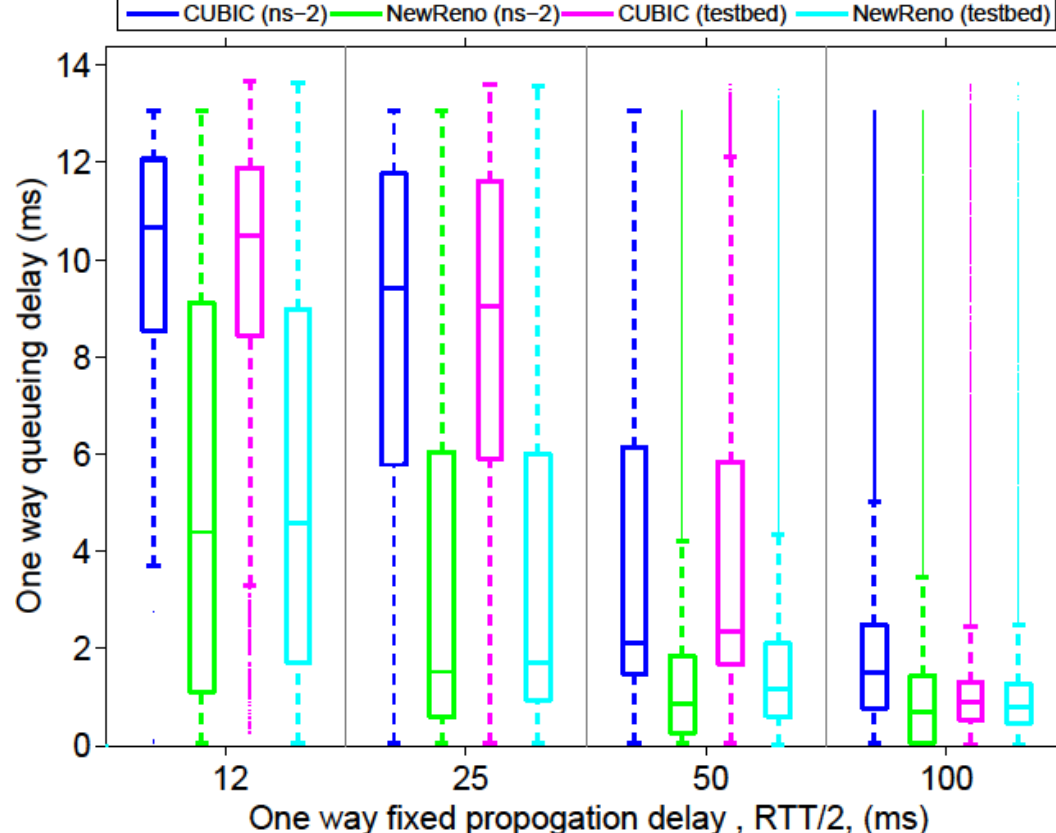
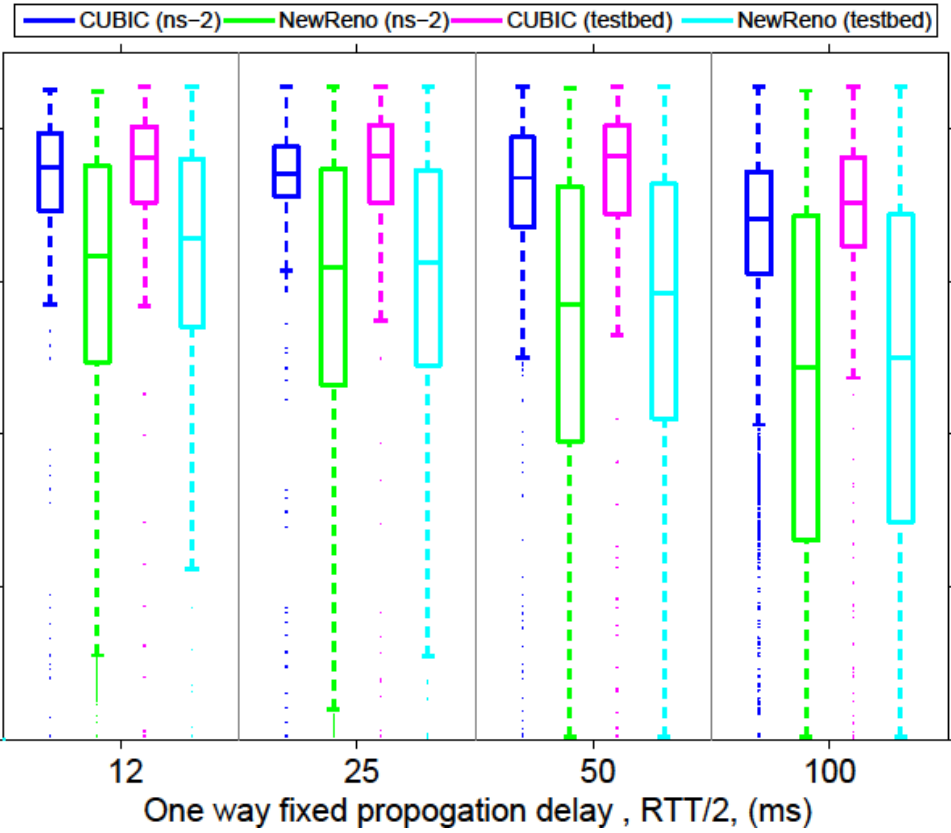
ADSL1 type link: 1500/256kbps, min 100ms RTT, 20 000byte queue.

Collateral Damage to Voice Stream - Latency



ADSL 1 (1500/256kbps)

ADSL 2 (24/1Mbps)



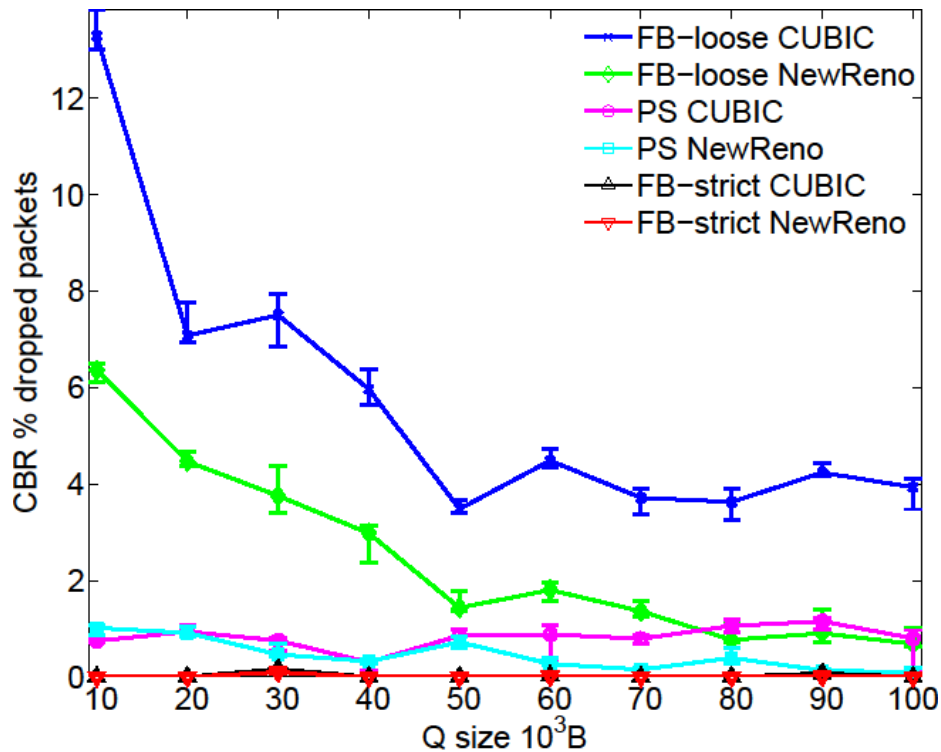
Induced downstream queuing delay - measured (testbed) and simulated (NS-2), 40 000byte FB-loose queue.



Collateral Damage to Voice Stream - Loss

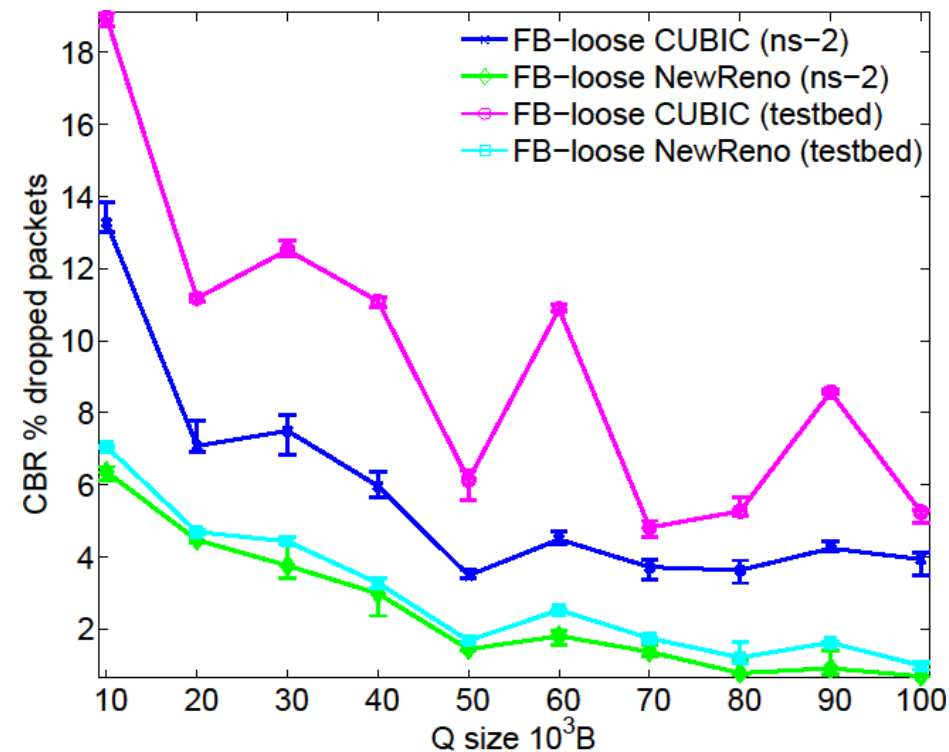
FB-loose, FB-strict & PS

NS-2



FB-loose

Testbed & NS-2

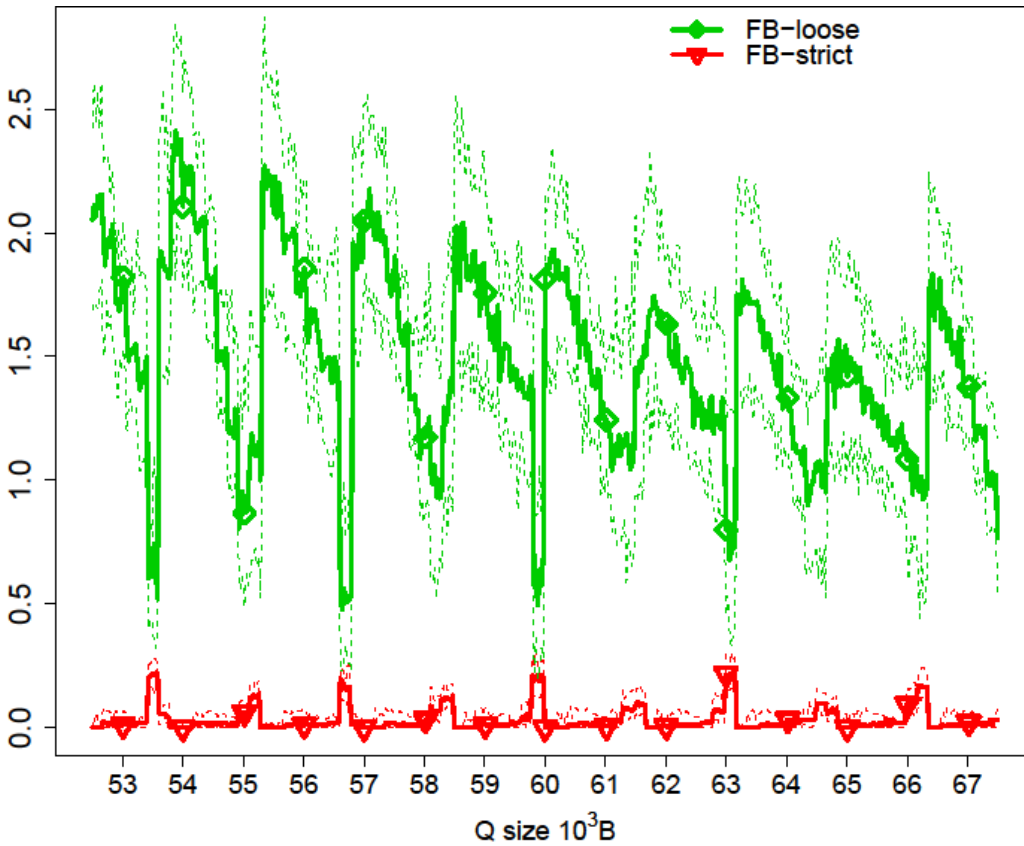


Proportion of CBR packets lost on the downstream for different packet drop mechanisms - 1500/256kbps link, 100ms minimum RTT.

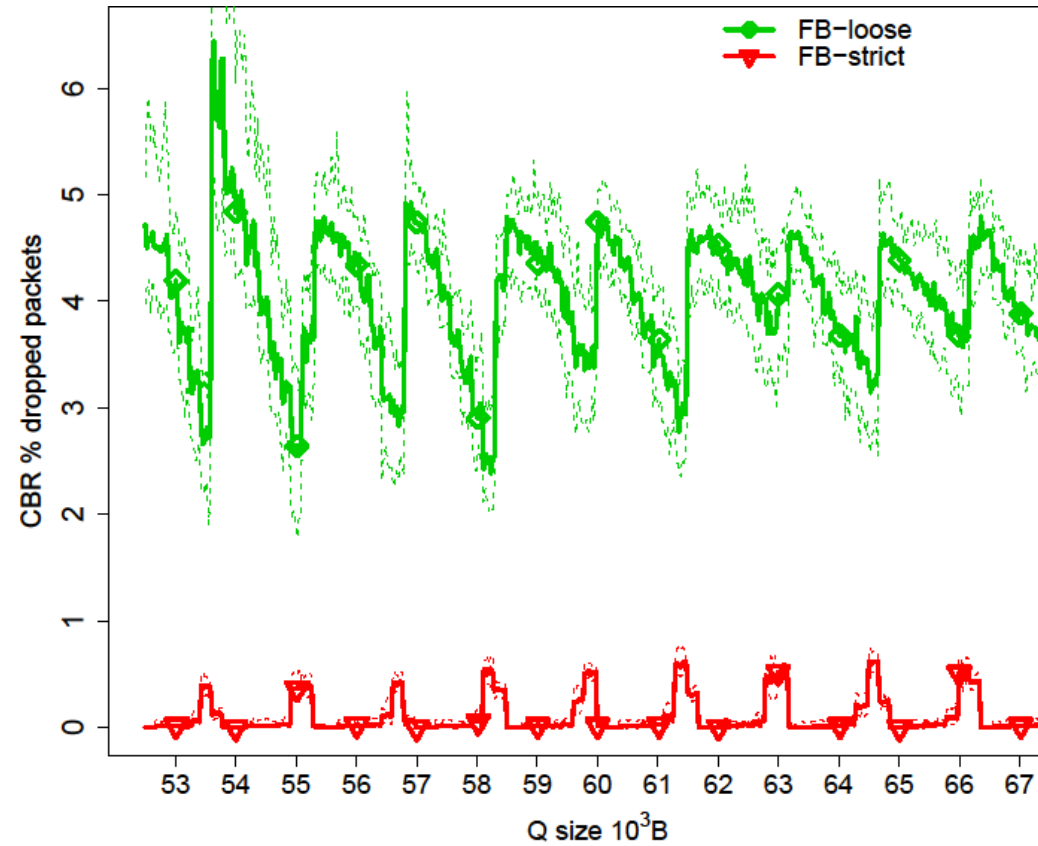


Detailed Simulation: Impact of Queue Size on Loss

NewReno



CUBIC

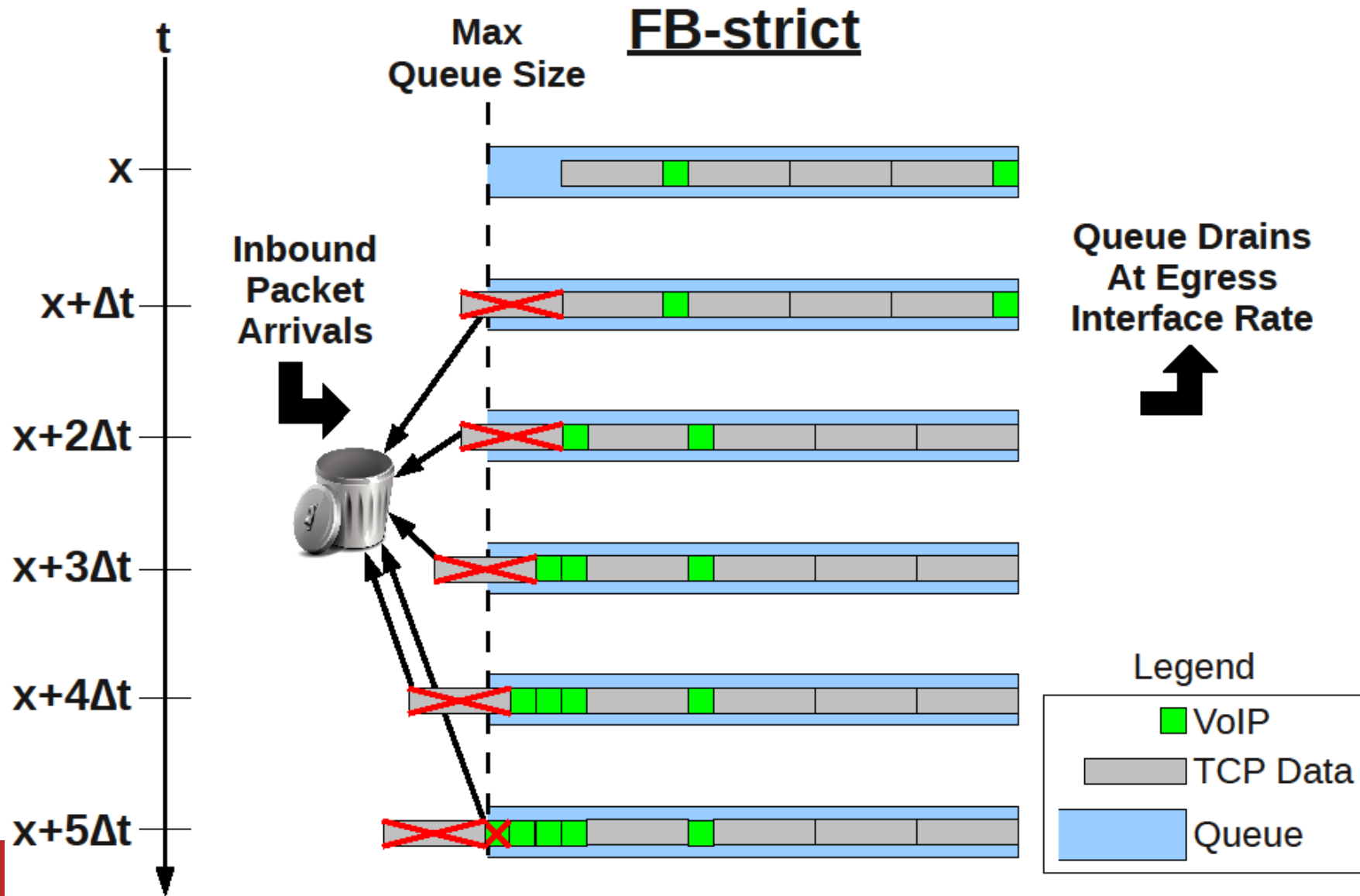


[52500, 67500] in 30 byte increments, 1500/256kbps, 100ms min RTT.

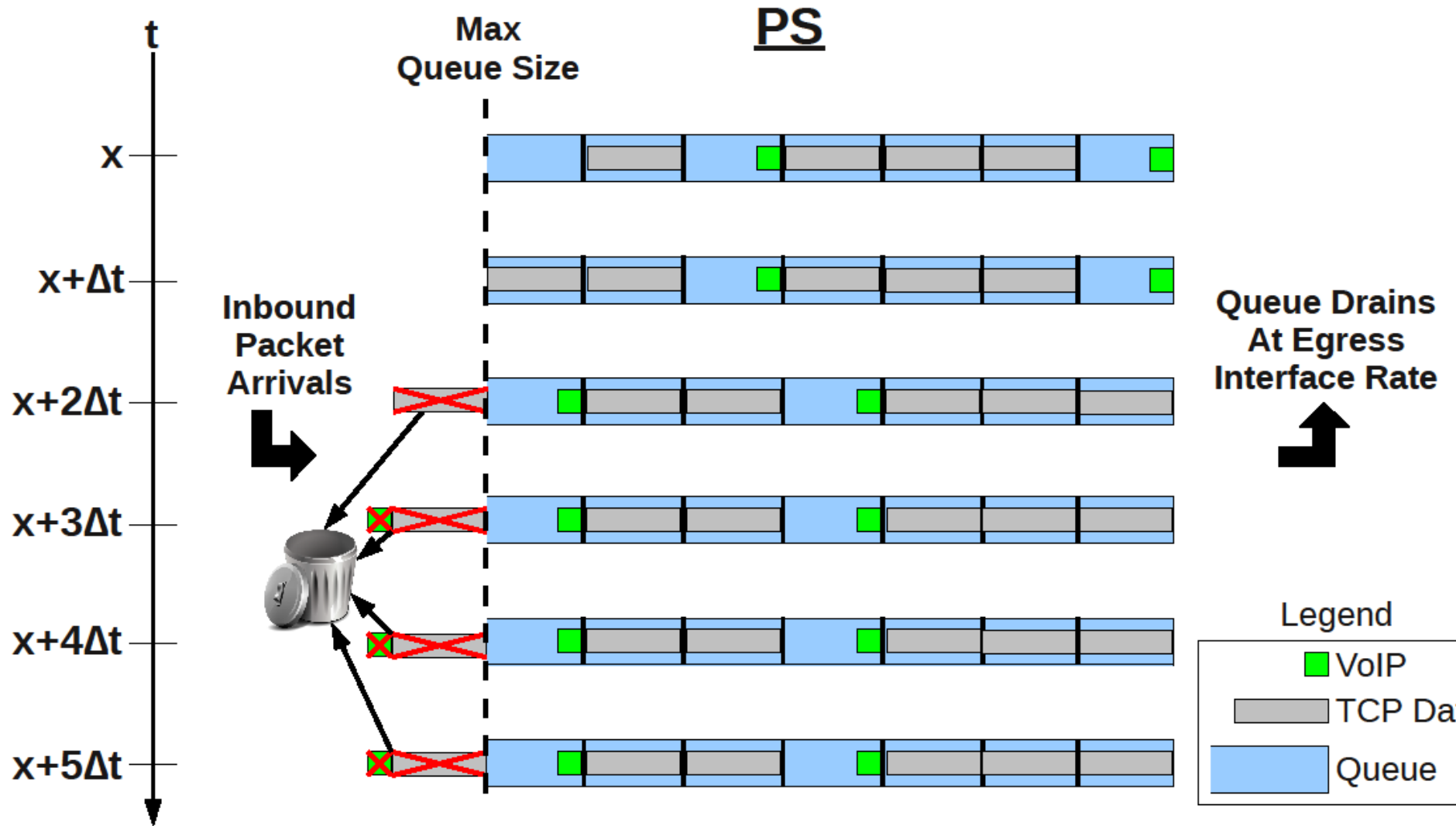
Dashed line envelope - 5% & 95% confidence intervals.



Time Series Analysis of FB-Strict

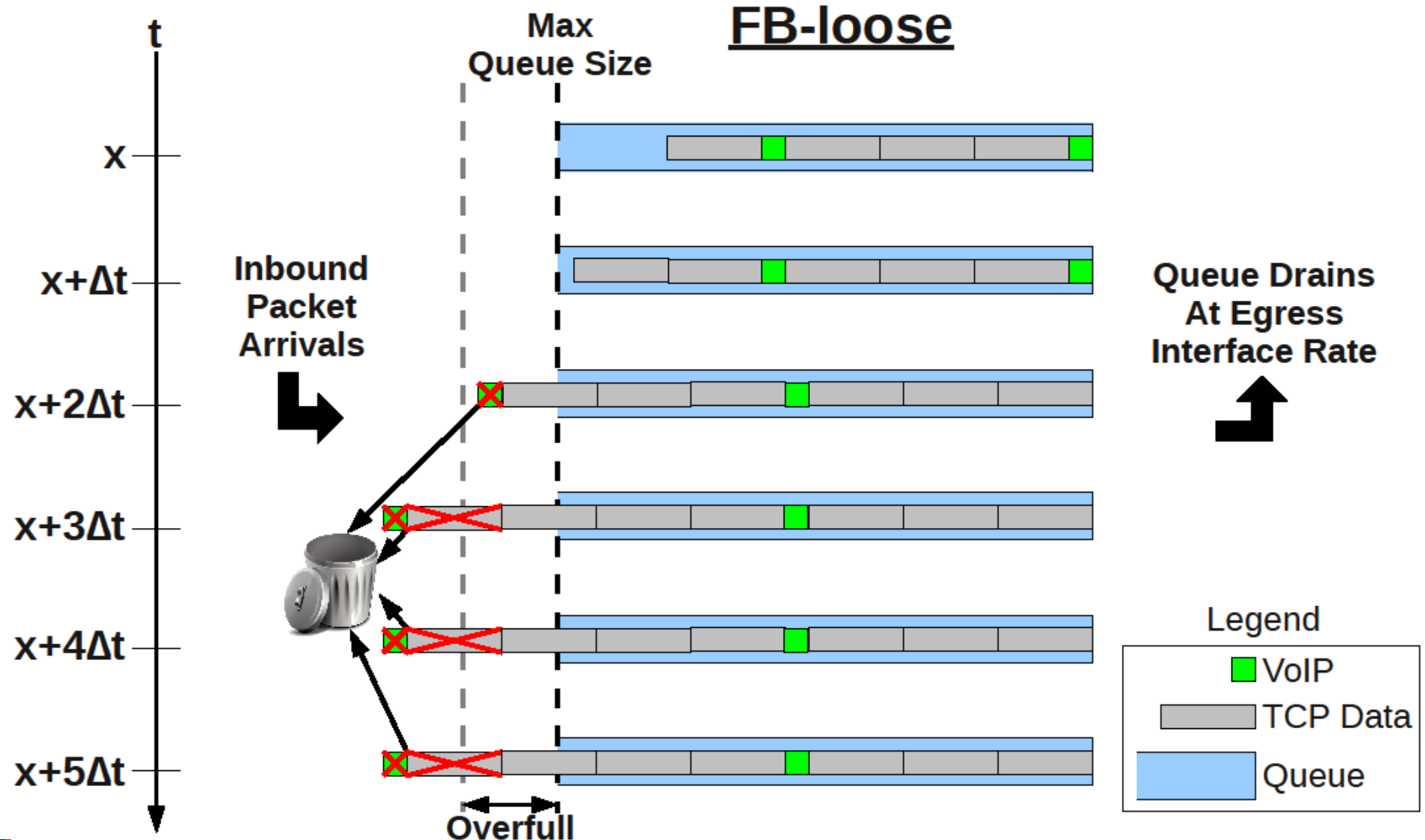


Time Series Analysis of PS





Time Series Analysis of FB-Loose





Detailed outline (section 3 of 3)

1. Introduction
2. Results
3. Conclusion
 - Summary & Open Questions



Summary & Open Questions

■ Summary

- CUBIC keeps buffers fuller: latency & loss increase
- FB-Loose disadvantages small, real-time multimedia packets
- Byte-based queue sizes around integer multiples of TCP packet size correspond with lowest FB-Loose loss & highest FB-Strict loss
- Dummynet implements FB-Loose byte-based queues

■ Open Questions

- Does using TCP for interactive traffic change things?
- Can we develop guidelines for home gateway manufacturers?