

# Multipath Load Balancing in Multi-Hop Wireless Networks



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# Multi-hop Wireless Networks

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- Nodes with radios
- Self configure to form a network
  
- Cheap and easy to deploy
- Robust
  
- Alternative to traditional wired infrastructure
- “Last mile” Internet access

# Motivation for Load Balancing

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- Multi-hop wireless has low bandwidth
  - Chain with ideal MAC: one quarter channel capacity
- Avoid congestion by distributing load

Can load balancing improve throughput?

# Previous Work

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- Improve reliability with backup paths
- Can decrease delay
- Theoretical analysis: improves aggregate throughput
- Improves performance when used with directional antenna, packet caching, new routing metrics

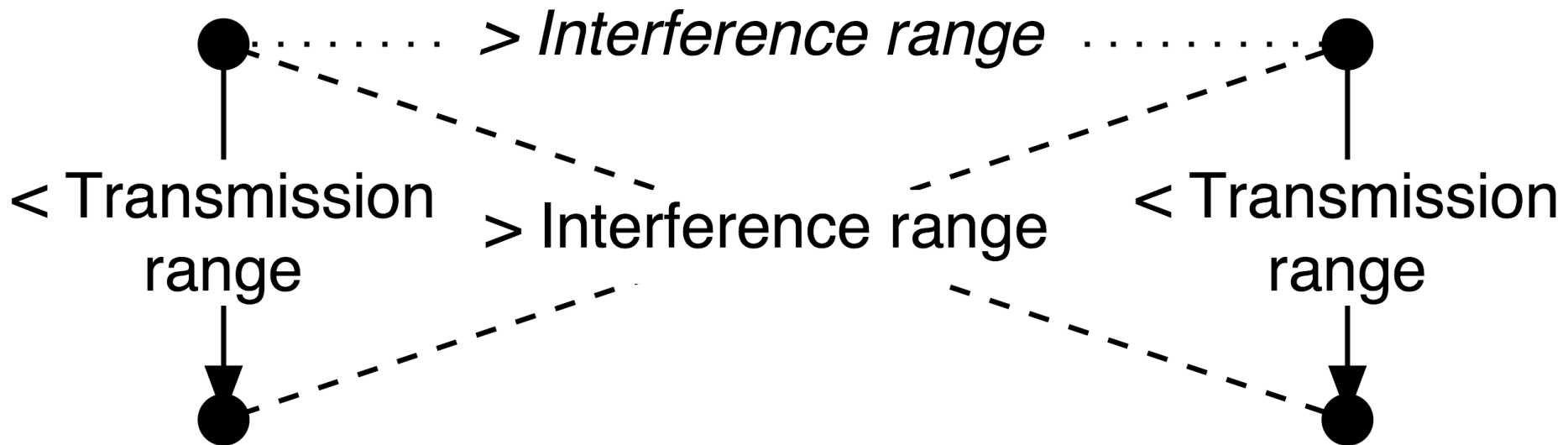
# Understanding Load Balancing

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- No mobility
- Fixed power transmissions
- Single channel
- Omnidirectional antennas

# Protocol Model of Interference

- Nodes must be within transmission range
- No other transmitters within interference range
- Carrier sensing: senders must be outside interference range

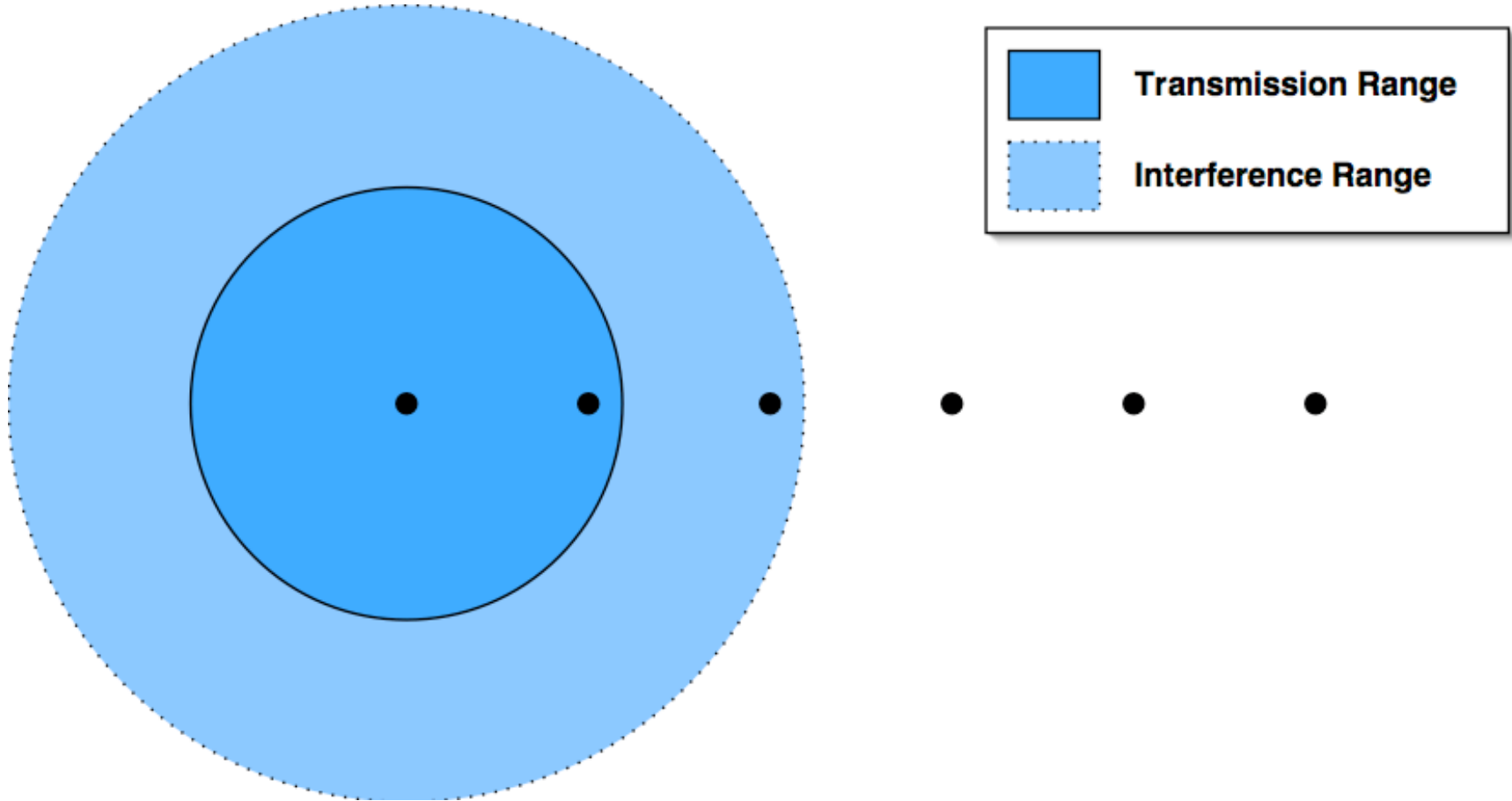


# Simplifying Assumptions

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- No MAC overhead
- Rate limited sender
- Nodes spaced at maximum range
- Fixed sized packets
- Interference range =  $2 \times$  (transmission range)

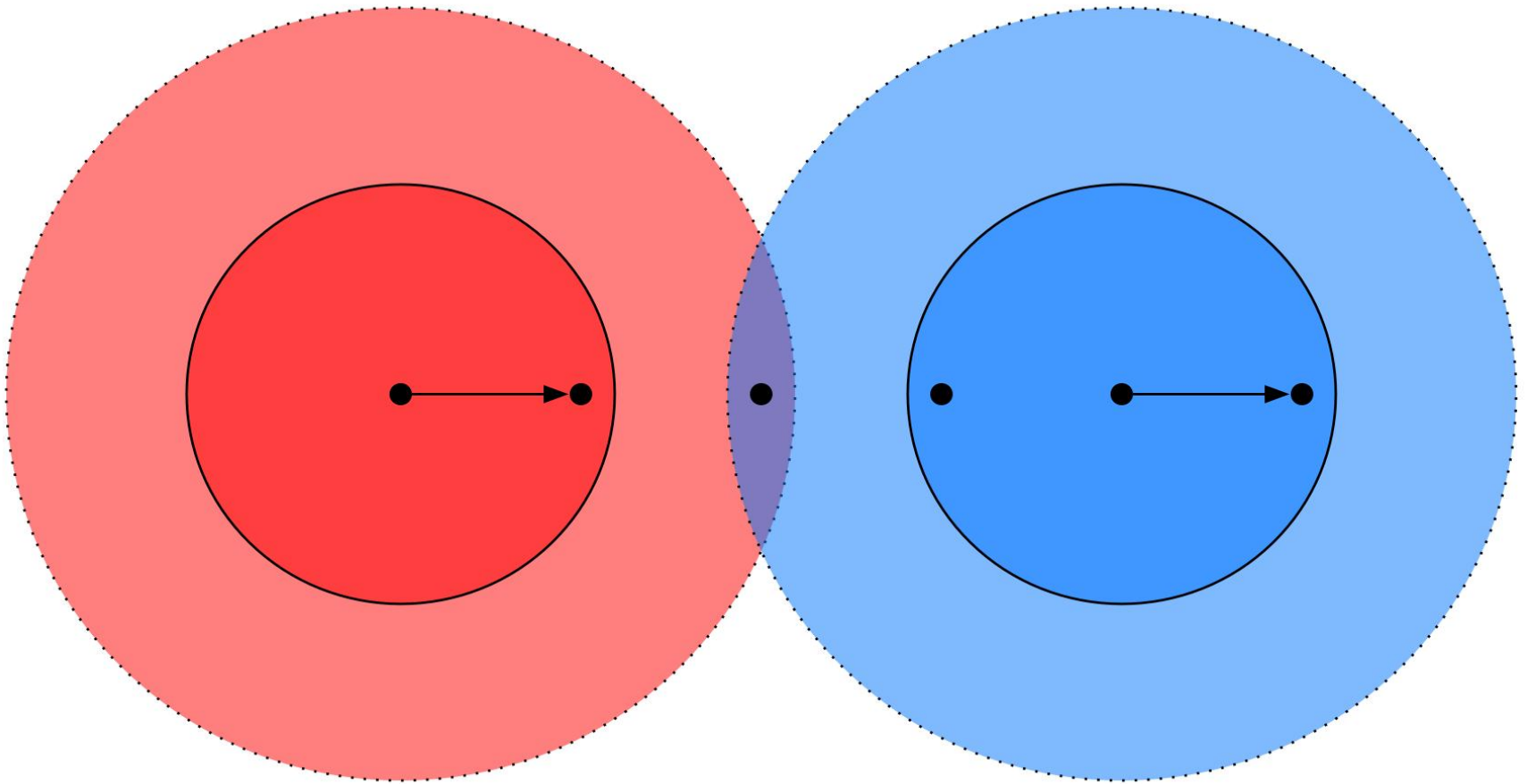
# Chain Topology





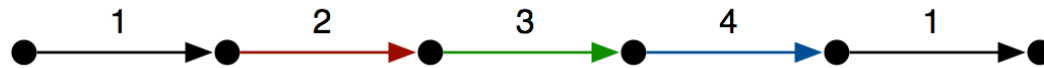
# Chain Topology

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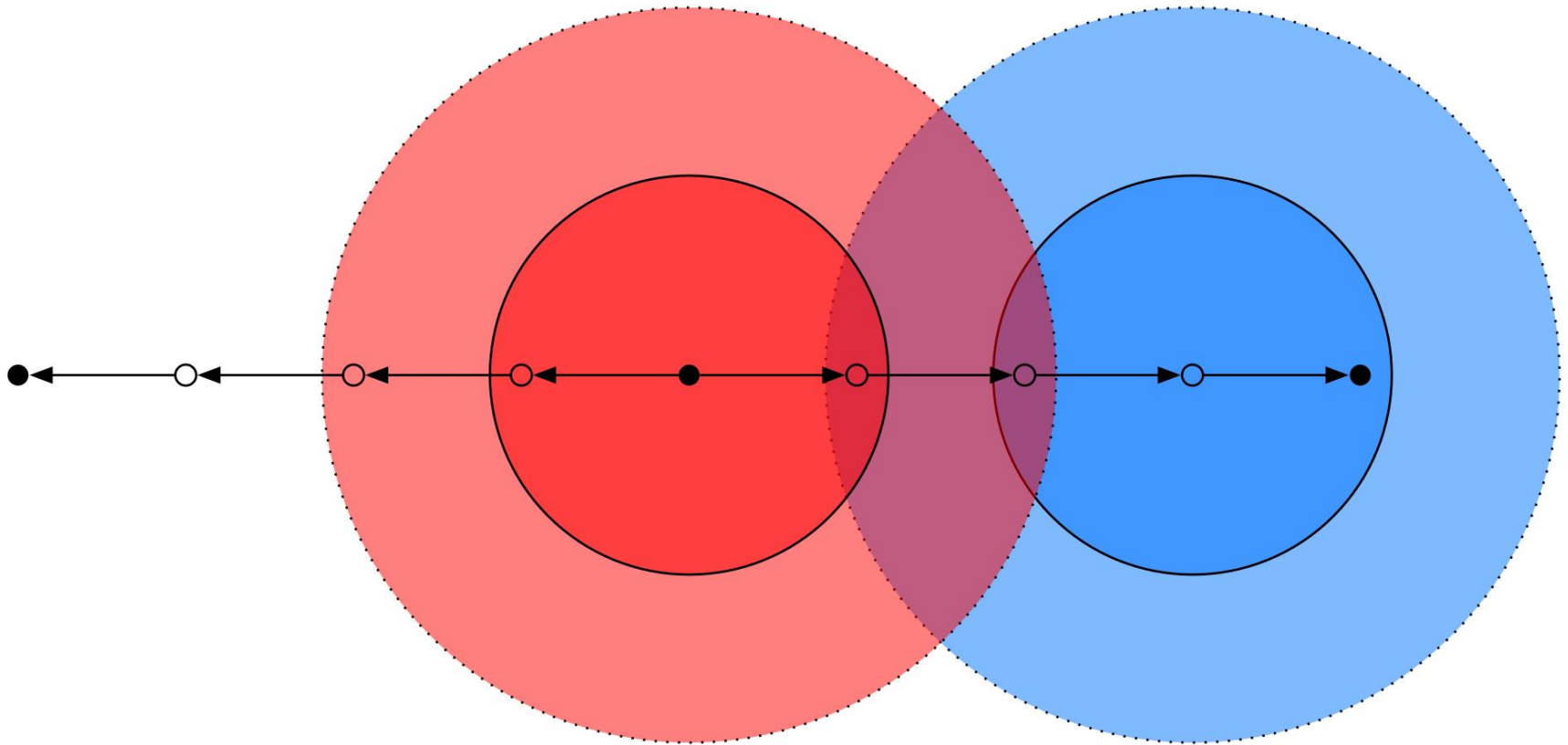
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$$\text{Rate} = \frac{1}{4}$$

# Two Directions: Out

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# Two Directions: Out

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$$\text{Rate} = \frac{1}{3}$$

# Two Directions: In

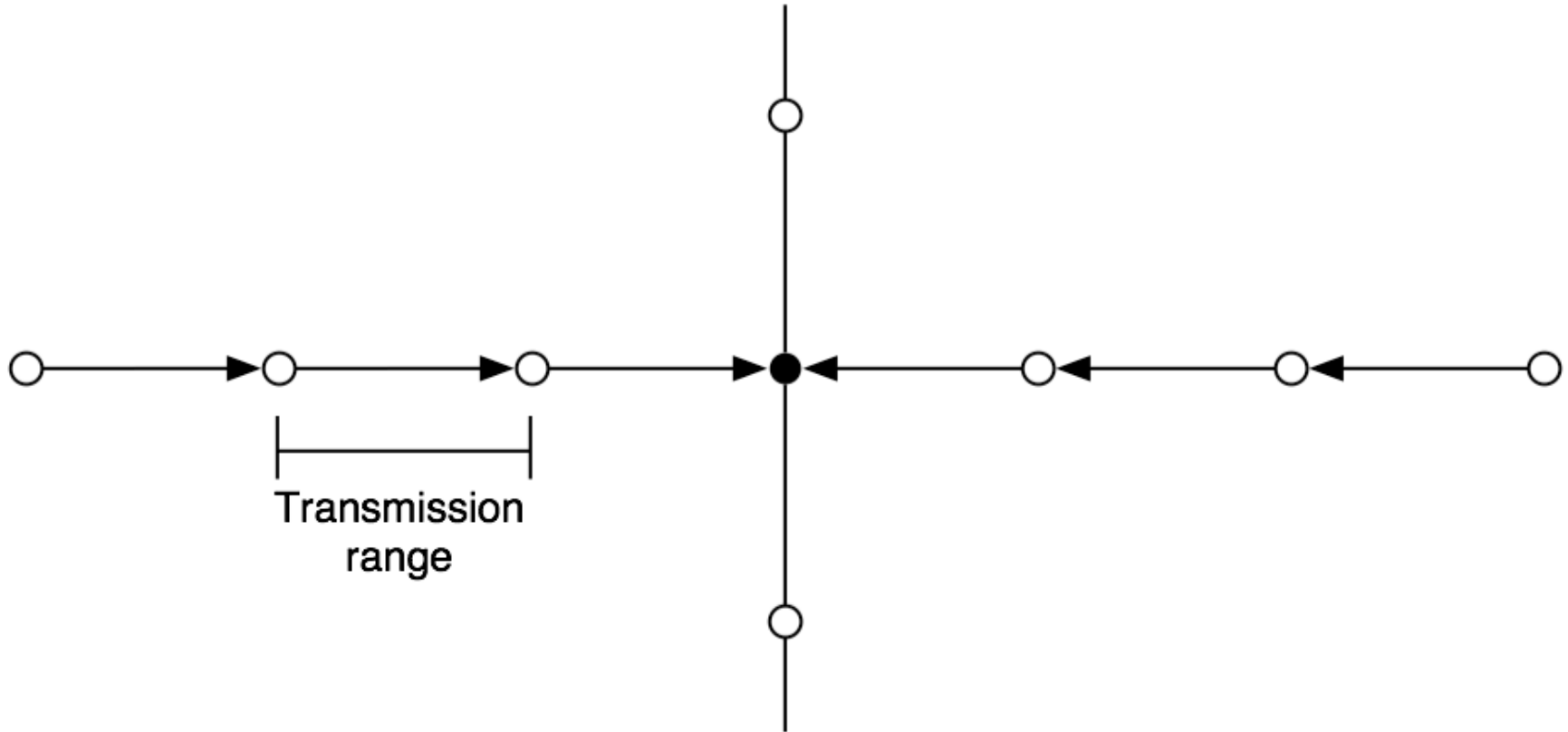
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$$\text{Rate} = \frac{1}{2}$$

# Cross Topology

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# Cross Throughput

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<b>Dir.</b>	<b>Paths (l=2T)</b>			
	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>
<b>Out</b>	$\frac{1}{4}$	$\frac{1}{3}$	$\frac{1}{3}$	$\frac{1}{3}$
<b>In</b>	$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{2}$	$\frac{1}{2}$

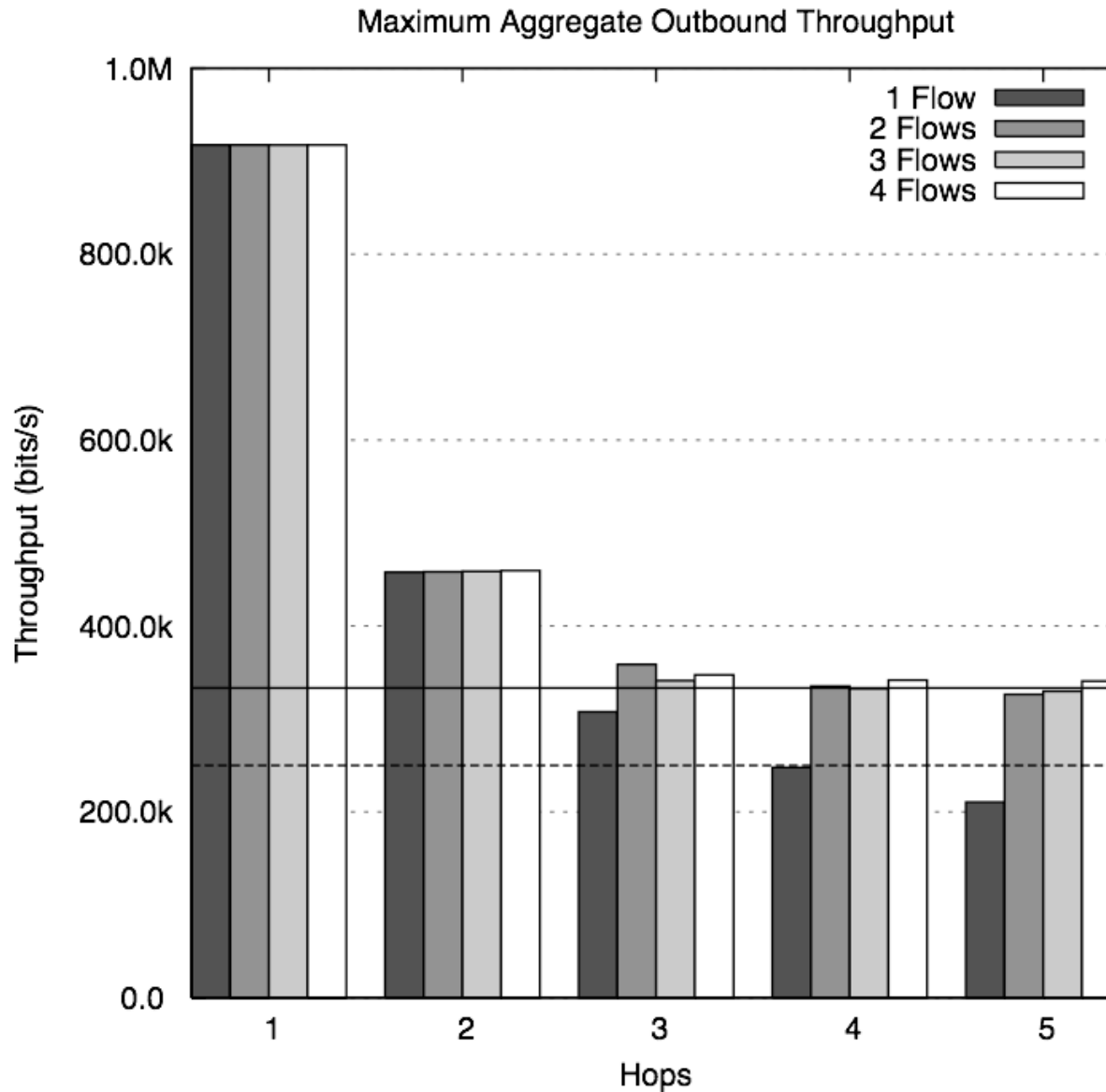
# More Realistic Model

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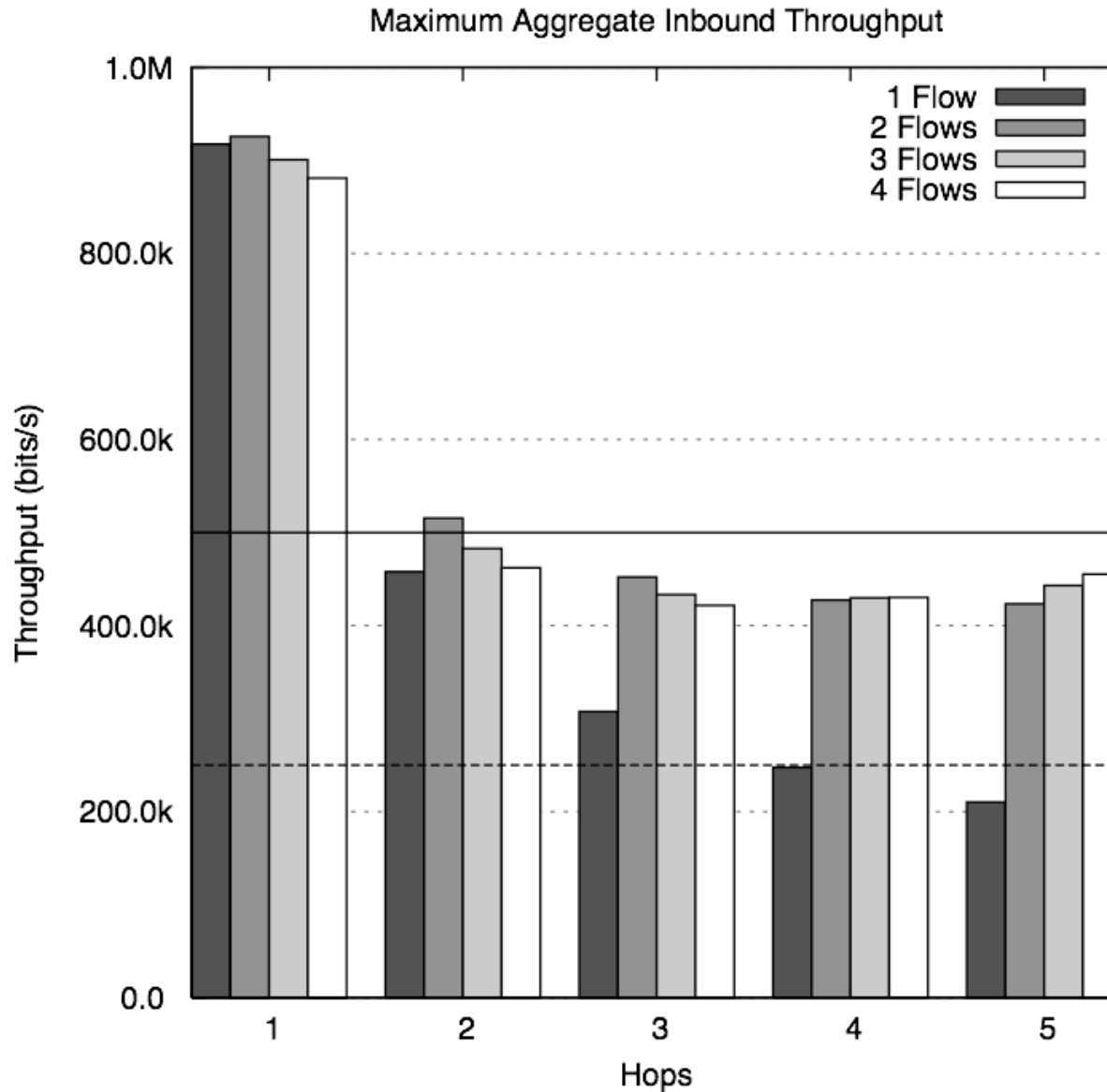
- MAC protocol: 802.11
- Power capture model of interference
  - If  $\text{SNR} > \text{threshold}$ : packet received
  - Two ray ground model
- Simulated with ns2
  - $T = 250\text{m}$ ,  $I = 550\text{m} = 2.2 T$
- 1 Mbps data rate, 1500 byte packets
- CBR sources, rates scaled from low to high load



# Cross: Throughput Out



# Cross: Throughput In



# End Points: Observations

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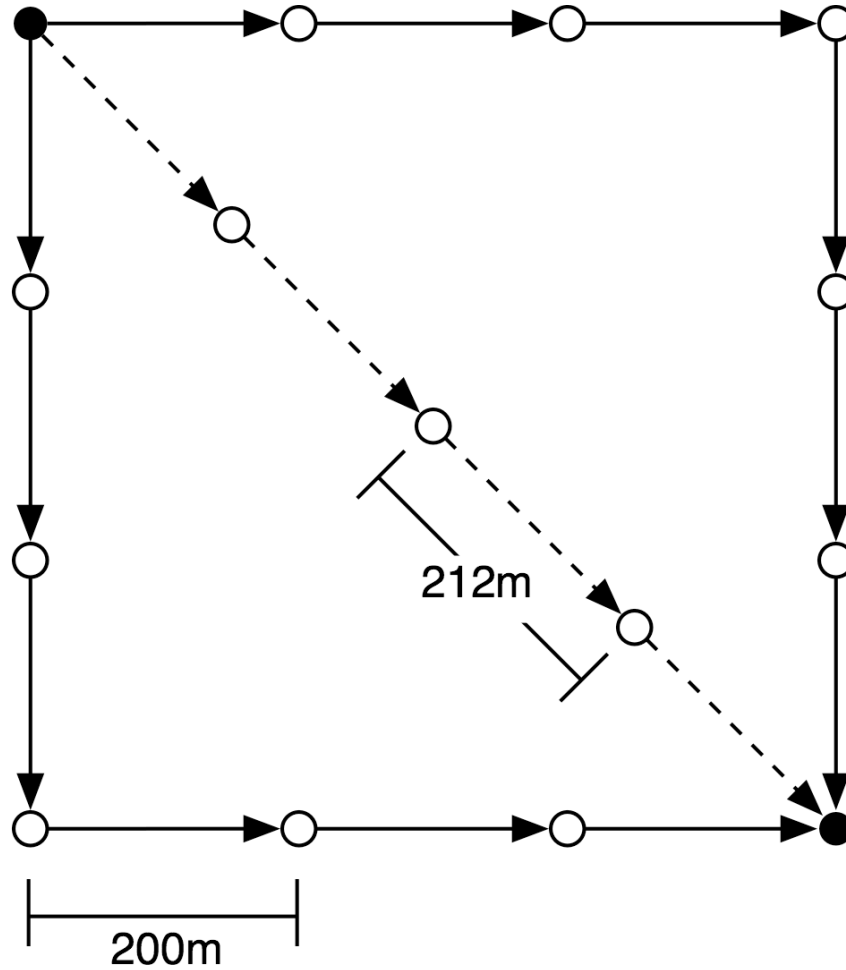
- Protocol model results match ns2 results
- Load balancing can improve throughput
  - Up to 101% increase in throughput
- 2 hops or less: no benefit
- Diminishing returns after adding second flow
- No delay improvement

# Simple Multipath Topology

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- Two flows
- At least three hops in the shortest path
- Concurrent transmissions must be outside interference range
  - ns2: Physical separation  $> 550\text{m}$
- Simple case:  $4 \times 4$  grid

# Simple Multipath: 4x4 Grid



# 4x4 Grid Performance

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<b>Metric</b>	<b>Single Path</b>	<b>Edge Path</b>	<b>Multipath</b>
Path Length (hops)	4	6	6
Throughput (bps)	252 720	196 440	267 840
Avg. Delay at 120 kbps	54.4 ms	80.8 ms	78.9 ms

# Grid Routing

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- Routing using node location
  - Half of the paths have  $> 35\%$  throughput improvement
- Heuristic using network topology
  - Half of the paths have  $> 20\%$  throughput improvement
- Some paths have  $80\%$  throughput improvement

# Load Balancing Conclusions

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- Can improve throughput
- Increases delay
  - Longer paths
  - Higher probability of collision
- Need at least three hops
- Longer paths are better
- Diminishing returns with more than two flows
- Very sensitive to interference



# Future Work

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- Multiple gateways
- Using TCP
- Multiple flows
- Multi-channel networks
- Random topologies

# Questions?

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