Non-Intrusive Load Disaggregation using Graph Signal Processing

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Abstract—With the large-scale roll-out of smart metering worldwide, there is a growing need to account for the individual contribution of appliances to the load demand. In this paper, we design a Graph signal processing (GSP)-based approach for nonintrusive appliance load monitoring (NILM), i.e., disaggregation of total energy consumption down to individual appliances used. Leveraging piecewise smoothness of the power load signal, two GSP-based NILM approaches are proposed. The first approach, based on total graph variation minimization, searches for a smooth graph signal under known label constraints. The second approach uses the total graph variation minimizer as a starting point for further refinement via simulated annealing. The proposed GSP-based NILM approach aims to address the large training overhead and associated complexity of conventional graph-based methods through a novel event-based graph approach. Simulation results using two datasets of real house measurements demonstrate the competitive performance of the GSP-based approaches with respect to traditionally used Hidden Markov Model-based and Decision Tree-based approaches.

Index Terms—energy disaggregation, graph signal processing, energy feedback, smart metering

I. Introduction

Appliance-level load demand information significantly enriches customer energy feedback and improves demand management measures via, for example, appliance load shifting [2]. While appliance-level energy monitors can yield accurate measurements, an alternative approach is preferred that is nonintrusive and accommodates the ever increasing number of appliances in a household, minimizing maintenance and installation costs due to sensor lifetime and networking/security issues. Hence, non-intrusive appliance load monitoring (NILM), i.e., disaggregation of the household aggregate load down to individual appliances, based purely on analytical tools operating only on aggregate load data, has been gaining popularity, especially with ongoing smart meter roll-outs worldwide. The business case for NILM is presented in [3], showing that resulting energy savings significantly surpass the costs of NILM technologies. NILM can support retrofit appliance advice, demand response measures, smart home automation [4], and is an enabler in decision making for home-owners, utilities, appliance manufacturers and policy makers.

Though NILM appeared in the 1980's [5], there has been a recent explosion in the NILM literature to tackle its practical

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challenges. NILM methods can be divided into two groups: steady-state and transient-state methods. Steady-state NILM methods rely on features extracted under steady-state operation of appliances, e.g., changes in steady-state active power [6], [7], [8], [9], reactive power [5], voltage and current waveform [10], [11], steady-state current harmonics and total harmonic distortion [12], [13], or voltage-current trajectory [14]. Transient-state NILM methods identify appliances based on their transient signatures, including transient power [15], high frequency voltage noise [16], [17], harmonics of the transients [18], [19], [20], duration and shape of power/voltage/current transient waveform [21], [22]. Transient-state approaches provide more distinguishable features than steady-state approaches and hence, in general, lead to higher disaggregation accuracy. Transient methods require sampling rates in the order of kHz or MHz [23], unlike steady-state methods, which are more sensitive to power level fluctuations and at low sampling rates, require longer monitoring time to capture all operation cycles [24]. For a more detailed review of NILM, see review papers [21], [23].

The proposed low-rate NILM approach in this paper is motivated by the increasing availability of low-rate data from electricity smart meters that are being deployed at large scale, with an increasing penetration rate, in Europe, Australia and the USA. For example, in the UK [25] and The Netherlands [26] every household will have access to 10-second active power readings. In the USA and Australia, smart meters providing readings at rates in the order of seconds and minutes, are massively deployed. Thus, NILM outputs can be accessible to the average household, without additional metering or monitoring hardware. This has prompted a recent trend in NILM literature tackling the NILM problem at low sampling rates, for example, [6], [7], [8], [27], [28], [29]. However, at low sampling rates, only steady-state features can be extracted reliably. Due to the similarity of steady-state load signatures among many domestic appliances, the NILM problem is particularly challenging.

Most recent low-rate NILM methods are state-based methods that represent each appliance operation using a state machine with distinct state transitions, based on appliance usage patterns. Such probabilistic approaches are usually based on a Hidden Markov Model (HMM) and its variants (see [7], [23], [27], [28], [29] and references therein). Four state-based methods for low-rate NILM, using conditional factorial HMM and Hidden semi-Markov graph models, are proposed in [27], but these methods have high computational complexity and are prone to converge to a local minimum. Another factorial HMM-based method is proposed by [7]

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for disaggregation of active power loads sampled at 1min, using expert knowledge to build initial models for states of known appliances, requiring correctly setting a priori-values for each state for each appliance, which is in turn limited by or strongly dependent on the particular aggregate dataset on which NILM is being performed. The Hierarchical Dirichlet Process Hidden Semi-Markov Model factorial structure is used in [30], removing some limitations of the approach of [27], but at increased complexity. A sparse coding algorithm, that discriminately trains sparse coding dictionaries, is proposed in [31], to learn a probabilistic model for each appliance's load demand over a typical week. The HMM-based method of lower complexity, proposed in [32], reduces the execution time by 72.7 times, but still requires 94 minutes for disaggregating 11 appliances. The main drawback of state-based approaches is the need for expert knowledge to set a-priori values for each appliance state via long periods of training and their high computational complexity, which makes them unsuitable for real-time applications [33].

Event-based NILM approaches have thus emerged [34], which are based on detecting events, usually via edge detection, when the load signal undergoes a statistically significant change indicating appliance use. After event detection, features (e.g., active power signature, increasing/falling edge [8], duration [35], uncorrelated power spectral components [6]) are extracted to classify the events into pre-defined categories, each corresponding to a known appliance. Different classification tools have been used, including support vector machines (SVM), e.g., in [36], neural networks, e.g., in [37], nonnegative tensor factorization [9], k-means [35] and decision trees [8]. Challenges encountered by event detection tools include large measurement noise, including large variance of active power readings for common household appliances, and similarity among active power steady-state signatures of different appliances.

This paper presents a different approach, developing a graph signal processing (GSP) method for steady-state NILM to address the large training overhead and associated complexity of conventional graph-based methods through a novel eventbased graph approach. GSP [38] is an emerging field that relies on expressing the piecewise smoothness of a signal through a graph. In [39], a GSP-based data classifier is proposed that searches for a smooth graph signal under known label constraints, and is applied to image and document datasets. The approach is based on the regularization of graph signals, using the fact that if a signal is piecewise smooth, then the total graph variation is generally small. Inspired by [38] and [39], in this paper, we propose a GSP-based approach for NILM by posing the load disaggregation problem as a single-channel blind source separation problem [31] to perform low-complexity classification of active power measurements. We treat active power measurements as a signal, indexed by the nodes of an undirected graph where each vertex corresponds to the signal sample, and the weights of the edges connecting the vertices reflect the degree of similarity between the nodes, i.e., the weights of the edges enable 'grouping' on/off events from the same appliance. Then, we define an optimization problem that contains the regularization term of the total graph variation, that is, we apply regularization on the constructed graph signal to find the signal with minimum variation. However, unlike the approach in [39], which solves for a smooth graph signal using initially known labels as prior, to avoid over-smoothing, we use the total graph variation minimizer as a starting point to minimise the difference between the total measured power and the sum of the disaggregated loads, deviating from traditional NILM approaches (see [23], [9] and references therein).

The rest of the paper is organized as follows. Sec. II provides a short background on GSP. Sec. III describes the proposed GSP-based NILM algorithms, followed by results in Sec. IV. The last section concludes the paper and highlights future work.

II. GRAPH SIGNAL PROCESSING (GSP)

In this section, we describe some basic concepts of GSP. All matrices are denoted by upper-case bold letters, such as \mathbf{X} . \mathbf{X}^T and $\mathbf{X}^\#$ are the transpose and pseudo-inverse matrices of \mathbf{X} , respectively. An element in the i-th row and j-column of matrix \mathbf{X} is denoted by $X_{i,j}$. Vectors are denoted by lower-case bold letters, such as \mathbf{x} , with the i-th element x_i , and $\mathbf{x}_{i:j}$ denotes a sub-vector $[x_i, x_{i+1}, \ldots, x_j]^T$, for i < j. A set is denoted using calligraphic bold-letters, such as \mathcal{M} , where $|\mathcal{M}|$ denotes its cardinality.

GSP is a novel signal processing concept [38], [40] that effectively captures correlation among data samples in time and space by embedding the structure of signals onto a graph [40] leading to a powerful, scalable and flexible approach suitable for many data mining and signal processing problems, ranging from image denoising and data compression, to classification, biomedical, and environmental data processing (see [38], [39], [40], [41] and references therein). GSP is particularly suitable for data classification when training periods are short and insufficient to build appropriate class models [39].

In GSP, a dataset \mathbf{x} is represented by a discrete signal s indexed by the nodes of a graph $\mathcal{G} = (\mathcal{V}, \mathbf{A})$, where \mathcal{V} is the set of nodes and \mathbf{A} is a weighted adjacency matrix of the graph. Each element $x_i \in \mathbf{x}$ corresponds to a node $v_i \in \mathcal{V}$. The weight of the edge between nodes v_i and v_j reflects the similarity between x_i and x_j and is usually defined using a Gaussian kernel weighting function, which is one of the most used kernels in machine learning for expressing similarity between dataset elements:

$$A_{i,j} = \exp\left\{-\frac{(x_i - x_j)^2}{\sigma^2}\right\},$$
 (1)

where σ is a scaling factor. Then, s, often referred to as *graph signal*, is defined as a mapping from V to a set of complex numbers. For example, s can be a set of classification labels, where s_i is set to the label of the class that x_i belongs to.

We define a graph signal's total Lipschitz smoothness [38], [42] with respect to the intrinsic structure of the underlying graph \mathcal{G} as:

$$\frac{1}{2} \sum_{i=1}^{N} \sum_{j=1}^{N} A_{i,j} \left(s_i - s_j \right)^2. \tag{2}$$

It can be shown that (2) is equal to $\mathbf{s}^T \mathbf{L} \mathbf{s}$ [38], where an $N \times N$ Laplacian matrix \mathbf{L} is defined as

$$\mathbf{L} = \mathbf{D} - \mathbf{A},\tag{3}$$

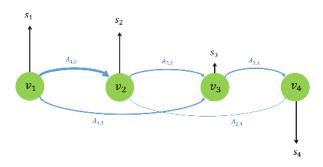


Fig. 1. A GSP example with four nodes.

which is a real symmetric matrix and can be seen as a difference operator for the graph signal s [38]. In (3), \mathbf{D} is an $N \times N$ diagonal matrix where for $k = 1, \ldots, N$, $D_{k,k} = \sum_{j=1}^{N} A_{j,k}$. Eigenvalues of \mathbf{L} carry the notion of frequency spectrum of \mathbf{s} , where values of eigenvectors associated with low eigenvalues (low frequency) change less rapidly [38] across the nodes, which can be used to design graph signal filters.

Fig. 1 shows an example with four nodes. The thickness of the graph edges reflect the correction between the nodes and vertical lines correspond to the graph signal values. If the graph signal s is piecewise smooth with respect to the graph \mathcal{G} , then (2) will be generally small, which can be used as a prior for regularization. Indeed, in classification [39], the elements that are strongly correlated will be connected via high-weight edges and associated with the same classification labels s_i making s varying smoothly across the connected nodes in the graph.

III. METHODOLOGY

A. Problem Formulation

Let \mathcal{M} be the set of all known appliances in the house and $p(t_i)$ the measured aggregate active power of the entire house measured at time t_i . Without loss of generality, in the following, we denote $p(t_i)$ as $p(t_i) = p_i \geq 0$. Let $p_j^m \geq 0$ be the power load of appliance $m \in \mathcal{M}$ at time instance t_j .

Let
$$\Delta p_i = p_{i+1} - p_i$$
 and $\Delta p_i^m = p_{i+1}^m - p_i^m$. Then,

$$p_i = \sum_{m=1}^{|\mathcal{M}|} p_i^m + n_i, \tag{4}$$

where, n_i is the measurement noise that also includes baseload and all unknown appliances running. The disaggregation task is, for i = 1, ..., N and $m \in \mathcal{M}$, given p_i , to estimate p_i^m .

B. GSP-based Disaggregation

To tackle the disaggregation problem using the GSP framework, we construct a graph $\mathcal{G}=(\mathcal{V},\mathbf{A})$, where each vertex $v_i\in\mathcal{V}$ is associated to one sample $\Delta p_i,\ i=1,\ldots,N$. For training, we assume availability of p_i and p_i^m , for $i=1,\ldots,n< N$, for all $m\in\mathcal{M}$. Then, the task is to estimate p_i^m , for $n< i\leq N$.

Let $Thr_m \ge 0$ be a power threshold for appliance m which is set during training (see Sec. IV) in such a way that if

the magnitude of the appliance active power change is larger than the power threshold, then we assume that the appliance changed its operation state (e.g., switched on/off). Then, we define an N-length graph signal \mathbf{s}^m for Appliance m as:

$$s_i^m = \begin{cases} +1, & \text{for } |\Delta p_i^m| \ge Thr_m \text{ and } i \le n \\ -1, & \text{for } |\Delta p_i^m| < Thr_m \text{ and } i \le n \\ 0, & \text{for } i > n. \end{cases}$$
 (5)

Note that s_i^m can be seen as a set of classification labels, where during training $(i \leq n)$ s_i^m is set to +1 (element i belongs to Appliance m class) or -1 (element i does not belong to the class), depending whether the appliance changed state or not. Since for the testing dataset (i > n) we do not know if the appliance was running, we set corresponding values of s_i^m to 0.

We can now calculate adjacency matrix **A** according to (1), where $x_i = \Delta p_i$. The graph smoothness can be calculated using (2).

Let $\mathbf{r}_i = \left[\Delta p_i^1, \Delta p_i^2, \cdots, \Delta p_i^{|\mathcal{M}|}\right]$ and let the difference between the actual aggregate power and the sum of disaggregated appliance powers be:

$$f(\mathbf{r}_i) = \left\| \Delta p_i - \sum_{m=1}^{|\mathcal{M}|} \Delta p_i^m \right\|_2^2. \tag{6}$$

We pose the disaggregation optimization problem as minimization of $\sum_{i=n+1}^N f(\mathbf{r}_i)$ using piecewise smoothness as a prior by introducing (2) as a regularization term, i.e.,

$$\min_{[\mathbf{r}_{n+1},\cdots,\mathbf{r}_N]} \sum_{i=n+1}^{N} f(\mathbf{r}_i) + \omega \sum_{m \in M} \left\| \mathbf{s}^{m^T} \mathbf{L} \mathbf{s}^m \right\|_2^2.$$
 (7)

Note that (7) defines an optimal solution as the smoothest solution that minimizes (6), where ω is a parameter that trades off smoothness and the minimization (6).

(7) is a hard optimization problem especially since $|\mathcal{M}|$ and N-n can be large. Thus, we propose two approximate solutions, one minimizing only the second term in (7), and the other minimizing both terms iteratively.

C. Solution 1: Total Graph Variation Classifier

If we assume, as in [39], that the true classification labels form a low-frequency graph signal \mathbf{s}^m , then for each Appliance m, an individual classifier can be defined that minimizes $\left\|\mathbf{s}^{m^T}\mathbf{L}\mathbf{s}^m\right\|_2^2$, i.e., one that finds the smoothest signal.

We call this classifier, the total graph variation classifier. The intuition behind it is that the labeled training samples for $i \leq n$ that are close in value to the unknown samples, j > n, will have large edge weights $A_{i,j}$, and so a smooth graph-signal prior will ensure that the testing samples have similar classification as these training samples.

Since $\mathbf{s}_{1:n}^m$ is known, determined during training, we can simplify the smoothness term as [41]:

$$\mathbf{s}^{m^{T}} \mathbf{L} \mathbf{s}^{m} = \mathbf{s}_{1:n}^{m^{T}} \mathbf{L}_{1:n,1:n} \mathbf{s}_{1:n}^{m} + \mathbf{s}_{1:n}^{m^{T}} \mathbf{L}_{1:n,n+1:N} \mathbf{s}_{n+1:N}^{m} + \mathbf{s}_{n+1:N}^{m^{T}} \mathbf{L}_{n+1:N,1:n} \mathbf{s}_{1:n}^{m} + \mathbf{s}_{n+1:N}^{m^{T}} \mathbf{L}_{n+1:N,n+1:N} \mathbf{s}_{n+1:N}^{m}.$$
(8)

Since matrix **A** is symmetric, **D** is diagonal and **L** is diagonally symmetric, it follows that:

$$\mathbf{s}_{1:n}^{m^{T}} \mathbf{L}_{1:n,n+1:N} \mathbf{s}_{n+1:N}^{m} = \mathbf{s}_{n+1:N}^{m^{T}} \mathbf{L}_{n+1:N,1:n} \mathbf{s}_{1:n}^{m}.$$
(9)

Using (9), and since $\mathbf{s}_{1:n}^m$ is constant, the first term does not affect minimization, we minimize (8) as:

$$\arg\min \left\| \mathbf{s}^{m^T} \mathbf{L} \mathbf{s}^m \right\|_2^2 =$$

$$\arg\min \left\{ 2\mathbf{s}_{n+1:N}^{m^T} \mathbf{L}_{n+1:N,1:n} \mathbf{s}_{1:n}^m + \right.$$

$$\mathbf{s}_{n+1:N}^{m^T} \mathbf{L}_{n+1:N,n+1:N} \mathbf{s}_{n+1:N}^m \right\}.$$
(10)

This is an unconstrained quadratic programming problem with a closed form solution [41], [43]:

$$\mathbf{s}^{m^*} = \mathbf{L}_{n+1:N,n+1:N}^{\#} \left(-\mathbf{s}_{1:n}^{m^T} \right) \mathbf{L}_{1:n,n+1:N}^{T}. \tag{11}$$

Once \mathbf{s}^{m^*} is determined, for i>n, if $s_i^{m^*}>T_s$, then, Appliance m changed state, $\Delta p_i^{m^*}$ is set to the mean of p_i^m which is calculated through the training process; otherwise, Appliance m did not change its state, and $\Delta p_i^{m^*}=0$.

In contrast to [39], where the threshold T_s is set to zero, we set our classification threshold $T_s = 0.5$, which imposes that only samples that are highly correlated with the training samples will be assigned to the same class. The value of 0.5 was found heuristically to yield the fewest false positives without increasing the number of false negatives.

We repeat minimisation of the smoothness term for all appliances $m \in \mathcal{M}$, where after each appliance has been disaggregated, its contribution to the total load is removed by subtracting its mean from the aggregate. Note that the same nodes are used for all appliances, but matrix **A** changes with updated \mathbf{r}_i^* , $i = n + 1, \dots, N$.

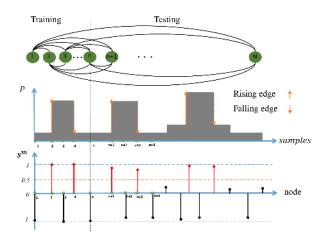


Fig. 2. An example of GSP-based disaggregaoin.

An example is given in Fig. 2. The top figure shows the generated graph nodes and connections between the nodes. Note that each node corresponds to one active power reading shown in the middle graph. The graph signal (shown on the bottom graph) contains classification labels for each power edge. During testing (i>n) all calculated values of \mathbf{s}^m above threshold $T_s=0.5$ will be classified to the Appliance m class, i.e., there was appliance state change or event.

The flow chart of the algorithm is shown in Fig. 3. The complexity of the approach depends on N-n since it is necessary to find the pseudo-inverse of an $(N-n)\times(N-n)$ real-valued matrix, which can be done in $O((N-n)^3)$ time.

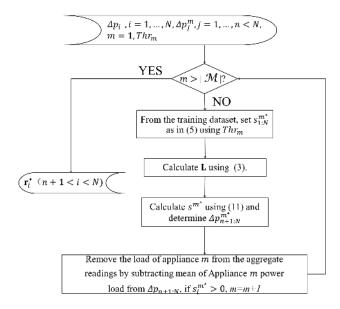


Fig. 3. Flow chart for Solution1: Total Graph Variation Classifier.

D. Solution 2: GSP + Simulated Annealing Refinement

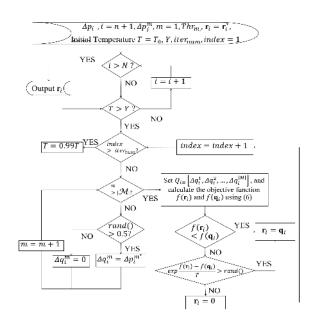


Fig. 4. Flow chart for Solution 2. $\mathbf{r}_i^*, i=n+1,\cdots,N$ is the solution found by Solution 1.

While Solution 1 finds the smoothest graph signal under given constraints, it may over-smooth the result. To avoid this, we introduce a sub-optimal solution for solving (7) based on minimising both terms.

First, we apply Solution 1 based on (11), which is the starting point for the simulated annealing method [44] that will attempt to refine the result by minimising the first term of (7). We demonstrate in the next section that simulated annealing provides a solution identical or close to a greedy full-search method, which has exponential complexity in $|\mathcal{M}|$.

The flow chart of Solution 2 is shown in Fig. 4. Note that the input is, \mathbf{r}_i^* , $i=n+1,\cdots,N$, which is a solution found by Solution 1. This solution is refined via iterative simulated annealing, where $iter_{num}$ denotes the number of iterations. In each iteration, a candidate solution \mathbf{q}_i is formed by randomly setting appliances on/off. Step $\exp\left\{\frac{f(\mathbf{r}_i)-f(\mathbf{q}_i)}{T}\right\} > rand()$ ensures that when the "temperature" T is high, the algorithm does not accept a worse solution, where rand() is a function that returns a random number in the interval (0,1). We heuristically demonstrate the convergence of the algorithm in the next section.

IV. EXPERIMENTAL RESULTS AND DISCUSSION

In this section, we present the following results: (1) Relative performance of using Simulated Annealing (SA) vs. Full Search, (2) Relative performance of Solution 1, SA only and Solution 2, (3) Comparison with state-of-the-art methods of [8] and [7], (4) Computational complexity.

We use the publicly available REDD dataset that contains load data from US houses [45], downsampled to 1min resolution, and the REFIT dataset [4], one of the largest UK datasets that contains active power measurements, sampled at 8 sec resolution, and collected over a continuous period of 2 years from 20 UK homes. The REDD dataset contains clean data and a small number of unknown appliances. On the other hand, the REFIT dataset contains many unknown appliances and high variations in the baseload. For validation purposes, household appliances, for which timestamped individual power consumption is available via submetering, are treated as known appliances. They are: Dishwasher (DW), Refrigerator (REFR), Microwave (MW), Washer dryer (WD), Kitchen outlet (KO), Stove (ST), air-conditioning high (AC_H) and low (AC_L) state, Electronics (EL), Washing machine (WM), Kettle (KE), Electric shower (ES), Electric heater (EH), Freezer (FRZ), Fridge-freezer (FFRZ). Baseload and Unknown appliances are abbreviated as BL and UN, respectively.

For training with the proposed methods, for each appliance, we used a period in the aggregate dataset when only that appliance is running (together with the BL). If a new appliance is introduced in the household, the training dataset is updated with that particular appliance's signature, comprising samples representing a full run from on to off. No retraining needs to be performed for other appliances. At least 14 days worth of data is used for testing.

 Thr_m is always set to one half of the difference between mean values of Appliance m's consecutive states. For example, if a two-state appliance, on and off, Thr_m would be half of the power value in the on-state. The scaling factor σ is picked during training in the area from the first non-zero value of the smoothness term to the inflection point, which was shown to provide the highest performance. For SA, temperature threshold Y=0.01, $T_0=100|\mathcal{M}|$ and $iter_{num}=1000$ which

trades off performance and complexity. For both proposed algorithms, we use windows of size 1000 samples, which ensures low complexity (see Subsection IV-D).

The evaluation metrics used are precision (PR), recall (RE) and F-Measure (F_M) [46] defined as:

$$PR = TP/(TP + FP) \tag{12}$$

$$RE = TP/(TP + FN) \tag{13}$$

$$F_M = 2 * (PR * RE)/(PR + RE),$$
 (14)

where true positive (TP) is recorded when the detected appliance was actually used, false positive (FP) is when the the detected appliance was not running, and false negative (FN) indicates that the appliance operation was not detected. Precision captures the correctness of detection - the higher the PR, the fewer the FPs. On the other hand, high RE means a low number of FNs, which implies that a higher percentage of appliance state changes are detected correctly. F_M balances PR and RE.

In addition, we use the average normalized error metric to measure the total energy difference between power estimated by the NILM algorithm and the actual power consumed, across all known appliances, which is defined as:

$$ANE = \frac{\left| \sum_{i=1}^{N} \bar{p}_i - \sum_{i=1}^{N} \hat{p}_i \right|}{\sum_{i=1}^{N} \bar{p}_i},$$
 (15)

where \hat{p}_i is the power estimated by the NILM algorithm from all disaggregated appliances $m \in \mathcal{M}$ at time i and \bar{p}_i is the actual power consumed by all known appliances at time i. This measure is useful, e.g., for appliance-itemized billing, when quantifying, across a fixed period of time, the error incurred by the NILM algorithm in estimating the total power consumed by individual appliances.

A. Full-Search vs Simulated Annealing

Fig. 5 shows an example of the convergence of the SA algorithm. It can be seen that the method converges after less than 300 iterations. Similar results are obtained for different datasets.

A full-search method can be used to minimize (6) off-line when $|\mathcal{M}|$ is small. Assuming only two-state appliances (i.e., ON/OFF), for each sample i, each appliance can either be switched on or off, or does not change state. Thus, with full-search, there are $3^{|\mathcal{M}|}$ possible combinations that should be inspected for each sample. Table I shows the F_M value comparison between the full-search method and SA for two houses from the REDD dataset. One can see that the proposed suboptimal SA approach finds a solution that is either identical to or very close to the full-search method.

B. SA, Solution 1 and Solution 2 Comparison

In this subsection we compare three approaches: (1) Total Graph Variation (Solution 1); (2) minimizing the first term of (7) only, i.e., (6), using SA; (3) Solution 2, incorporating the latter two approaches. As shown in Tables II and III, SA and

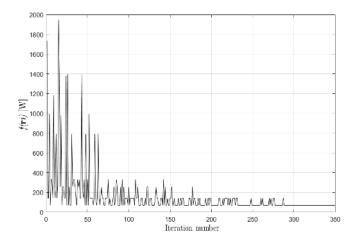


Fig. 5. Convergence of the simulated annealing method for House 2 from the REDD dataset.

	REFR	ST	MW	KO	EL	AC_H	AC_L	DW
FS H2	0.84	0.31	0.91	0.83	-	-	-	0.62
SA H2	0.84	0.30	0.91	0.83	-	-	-	0.62
FS H6								
SA H6	0.77	0.75	0.77	0.55	0.48	0.92	0.53	-

Solution 1, lead to significantly worse F_M performance for some appliances than treating them jointly (Solution 2).

When minimizing Equation (6), SA only uses known mean values of p_i^m and thus does not account for small fluctuations in actual instantaneous p_i^m . SA results in Table II are poor for Stove, since Stove (mean $p_i^m = 408\text{W}$) is often confused with the low-power state of Dishwasher (mean $p_i^m = 349\text{W}$). Similarly, in Table III, Kitchen Outlet and Electronics have similar operating power states, hence they are often incorrectly labelled.

SA only does not outperform Solution 1, however the advantage of Solution 2 (integrating Solution 1 and SA) is consistent across all appliances and especially significant for appliances such as Electronics and AC.

Appliance	REFR	ST	MW	KO	DW
Avg.Power [W]	171	408	1840	1056	1201 (349)
Solution1	0.81	0.81	0.91	0.85	0.83
SA	0.84	0.30	0.91	0.83	0.62
Solution2	0.84	0.86	0.93	0.88	0.84

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Appliance	REFR	ST	MW	KO	EL	AC_H	AC_L
Avg.Power [W]	149	1724	1352	946	815	2376	357
Solution1	0.77	0.92	0.91	0.86	0.28	0.94	0.53
SA						0.92	0.56
Solution2	0.78	0.92	0.91	0.88	0.66	1.00	0.79

C. Comparison with state of the art

A comparison of performance of Solution 2 with the state-of-the-art NILM approaches, namely Decision Tree (DT) approach [8] and HMM-based approach [7] is shown in Tables IV, V, and VI for REDD Houses 1, 2, 6, respectively.

TABLE IV

COMPARISON BETWEEN THE PROPOSED SOLUTION 2 (P), HMM AND

DT-BASED METHODS FOR REDD HOUSE 1.

Appliance	REFR	MW	DW	KO	WD
F_{M_p}	0.88	0.70	0.57	0.39	0.89
$F_{M_{HMM}}$	0.97	0.50	0.13	0	0
$F_{M_{DT}}$	0.88	0.85	0.39	0.19	0.88

TABLE V
COMPARISON BETWEEN THE PROPOSED SOLUTION 2 (P) AND HMM AND DT-BASED METHODS FOR REDD HOUSE 2.

Appliance	ST	REFR	KO	MW	DW
F_{M_p}	0.86	0.84	0.88	0.93	0.84
$F_{M_{HMM}}$	0.21	0.90	0.68	0.47	0.04
$F_{M_{DT}}$	0.33	0.97	0.92	0.97	0.32

TABLE VI
COMPARISON BETWEEN THE PROPOSED SOLUTION 2 (P), HMM AND
DT-BASED METHODS FOR REDD HOUSE 6.

Appliance	ST	REFR	KO	MW	AC	EL
F_{M_p}	0.92	0.77	0.88	0.91	0.88	0.66
$F_{M_{HMM}}$	0	0.88	0	0	0.12	0.03
$F_{M_{DT}}$	0.67	0.99	0	0	0.89	0

The proposed method was also tested using the noisy REFIT dataset [4]. The REFIT households were monitored remotely and uninterruptedly, while they were going about their usual domestic routines. Each house contains numerous appliances that were not monitored, including oven, lights, chargeable devices, small electronics, etc., which are considered unknown and contribute significantly towards noise. Tables VII and VIII show results for REFIT Houses 2 and 17, respectively. These two houses were selected as two houses which had relatively fewer unknown appliances compared to other houses in the dataset.

TABLE VII

COMPARISON BETWEEN THE PROPOSED SOLUTION 2 (P), HMM AND DT-BASED METHODS REFIT HOUSE 2.

Appliance	FRZ	WM	DW	TV	MW	KE
F_{M_n}	0.77	0.55	0.62	0.49	0.95	0.88
$F_{M_{HMM}}$	0.49	0.26	0	0.06	0.01	0.01
$F_{M_{DT}}$	0.33	0.73	0.36	0	0.95	0.58

Tables IX, X, and XI show the relative contribution of known appliances to the total aggregate load, for REDD Houses 1 and 2 over a period of two weeks and REFIT Houses 2 and 17 over a period of one month (October 2015). The proposed Solution 2 can disaggregate over 60% of the household's total load, demonstrating its effectiveness in accounting for individual appliance demand.

TABLE VIII

COMPARISON BETWEEN THE PROPOSED SOLUTION 2 (P), HMM AND DT-BASED METHODS FOR REFIT HOUSE 17.

Appliance	FRZ	FFRZ	KE	MW	WM
F_{M_p}	0.64	0.76	0.96	0.81	0.76
$F_{M_{HMM}}$	0.32	0.19	0.01	0.22	0.01
$F_{M_{DT}}$	0.81	0.79	0.97	0.71	0.50

TABLE IX
REDD HOUSE 1 AND 2. THE TOTAL DEMAND FOR TWO WEEKS WAS 158
KWH AND 77 KWH FOR HOUSES 1 AND 2, RESPECTIVELY.

	REFR	MW	ST	KO	DW	WD	BL	UN
	14%	8%	-	1%	7%	10%	22%	39%
H2	31%	8%	2%	1%	5%	-	16%	37%

TABLE X
REFIT HOUSE 2. TOTAL MONTHLY DEMAND WAS 372 KWH.

FFRZ	WM	DW	TV	MW	KE	ES	BL	UN
7%	4%	12%	< 1%	< 1%	6%	15%	18%	35%

TABLE XI REFIT HOUSE 17. TOTAL MONTHLY DEMAND WAS 341 kWh.

FRZ	FFRZ	KE	MW	WM	EH	ES	BL	UN
15%	7%	8%	2%	2%	11%	8%	21%	27%

With respect to disaggregation accuracy, the ANE measure given by (15), which measures the discrepancy between the true consumption and the disaggregated values, is, for REDD Houses 1 and 2, 7.33% and 6.91%, respectively, and for REFIT Houses 2 and 17, the ANE is 8.97% and 9.24%, respectively. These results demonstrate that a very small percentage of the total load is wrongly disaggregated.

1) Discussion: As can be seen from Tables IV, V, and VI, the proposed method significantly outperforms the HMM-based approach for all appliances in the REDD dataset except the refrigerator. HMM usually performs well disaggregating the refrigerator due to continuous and sole operation (i.e., without any other appliances running) during the night and hence large data availability for learning and improving initial models [7], [8]. The poor performance for other appliances can be attributed to the short training period. The proposed method shows better or similar performance to the method of [8] expect for microwave in Houses 1 and 2, and refrigerator in most houses. These appliances have very small power fluctuations during operation, and hence the decision tree classifier based on the rising and falling power edge works especially well.

Tables VII and VIII show poor HMM results for the REFIT dataset due to the noise and many unknown appliances. The proposed solution is again better than or equal to the benchmark methods for most of the appliances except washing machine in House 2, where DT performs the best due to very distinctive high-state washing machine power edges.

The results for both REDD and REFIT dataset demonstrate that the proposed method provides more accurate disaggregation than the benchmark methods for most appliances. The difference is especially pronounced for the kitchen appliances, namely Kettle, Microwave, and Stove. This is due to the fact that Stove and other kitchen appliances normally have a short operation time and relatively high power, thus machine-learning based approaches cannot generate probabilistic models that accurately capture appliance operation. The HMM-based approach is sensitive to noise and suffers from a short training period. The DT-based method works well for appliances that have small fluctuations in load during their steady-state operation.

The results for multi-state appliances (dishwasher, washing machine) are generally worse for all tested algorithms. This is due to first, the fact that low-power operating states are often difficult to detect since they are 'hidden' in the baseload and noise. Second, since these appliances are on for a very longer period, many appliances are likely to run in parallel, adding to noise. Finally, multi-state appliances are not used frequently, thus it is more difficult to isolate good training periods.

On the other hand, cold appliances, refrigerator, freezer and fridge-freezer, are always on and have regular periodic signatures; thus the algorithms show good accuracy. However, due to higher level of noise from many unknown appliances, slightly worse results are obtained for the REFIT dataset with a higher false positive rate.

The TV in REFIT House 2 is hard to identify since it has relative low operating power and thus it is often hidden by noise and baseload. In addition, TV runs usually for a long period of time, and thus many appliances will run in parallel. Still, the proposed method is more successful than the benchmark methods, since it is less sensitive than HMM to the training dataset that does not have any instances of TV running alone and it is more successful in resolving the cases when multiple appliances run in parallel than DT. Electronics in REDD House 6 include different electronics equipment which produce complex power signatures that lead to worse results. The proposed algorithm again shows more robustness than the benchmark methods in this situation, since it is less sensitive to fluctuations in steady-state power signature during training and testing.

Tables IX, X, and XI show that unknown appliances, including lighting, whose consumption could not be validated, make up only under 40% of the total load. The ANE results indicate that the discrepancy between the actual load due to known appliances and NILM-estimated load is very low. However, the ANE results of REFIT houses are slightly worse than in the case of REDD dataset, mostly due to additional noise from unknown appliances and multi-state appliances, i.e., washing machine and dishwasher. Note that, each REFIT house contained over 40 operational appliances (see [4] for monitoring details), many of which could not be validated via a time-diary or appliance-level load measurements. Moreover, considering that lighting, which contributes towards the 'Unknown' category in Table X and XI, accounts for about 20% of a household's domestic consumption in October in UK (see [4] and references therein), the results demonstrate the potential of the proposed technology in effectively disaggregating smart meter aggregate loads.

D. Computational complexity

The proposed algorithm was implemented in Matlab2015b and executed on an Intel(R) Core(TM) i5-3470 CPU @ 3.20GHz machine running Windows 7 64-bit. Table XII shows the computational time needed to disaggregate the refrigerator, which is periodically on, in the REDD House 1 for various testing set sizes N-n. The shaded row shows the accuracy when the testing dataset of size N-n is split into windows of 1000 samples each, and each window is independently processed, this way reducing the dimension of matrix L and hence lowering the complexity of calculation in (11). The bottom unshaded row shows the accuracy when the entire testing dataset is processed in one go - it also shows indirectly the effect of using only one 'window' of size 1000, 2000, 3000 etc. Similar F_M values in all cases confirm that splitting the testing dataset in smaller manageable 'windows' significantly reduces execution time without a loss in performance. The table also shows that the proposed method performs well with a small training overhead, i.e., increasing the number of samples in the testing dataset does not improve performance. Note that a window size of less than 1000 would not capture a full run of appliances with long operation period, such as washing machines or AC.

TABLE XII

COMPUTATION TIME OF THE PROPOSED SOLUTION 2 FOR REFR IN REDD HOUSE 1. THE SHADED ROWS SHOW RESULTS OBTAINED USING SMALL WINDOWS OF 1000-SAMPLES EACH. THE BOTTOM ROWS SHOW RESULTS WHERE THE TRAINING DATASET IS NOT SPLIT.

N-n	1000	2000	3000	4000	5000	6000	7000
Time [s]	1.33	3.22	3.86	4.41	4.99	5.52	6.01
F_M	0.89	0.89	0.90	0.90	0.90	0.89	0.89
Time [s]	1.33	16.85	44.81	90.96	153.99	243.88	406.69
F_M	0.89	0.89	0.90	0.90	0.90	0.90	0.90

When two weeks of data are used, comprising just over N=20,000 samples, the computation time on average for REDD House 1, 2 and 6, is between 10 to 12 sec for disaggregating one appliance. This is faster than the HMM-based method which disaggregates the same amount of data in 40-50 sec, as reported in [35].

V. CONCLUSION

This paper presented two NILM algorithms based on the emerging concept of GSP. The first approach minimizes the total graph variation. The second approach further refines the total graph variation solution using simulated annealing. Experimental results show the competitiveness of the methods with respect to two NILM methods, and were demonstrated over two datasets with a range of appliances. We also discuss the relative performance of the proposed methods for different appliances and how robust the methods are to short training periods, and how fast this can be implemented through effective windowing without performance loss. The proposed methods could work with conventional smart meters, e.g., accessing 10 second data via the Consumer Access Device, and do not require any additional hardware installation. Future work will consist of using confusion matrix results of other similar

measures to attempt to identify in more detail weaknesses and efficient real-time implementation of the proposed algorithms and integration into smart home decision support systems for demand response as well as designing advanced energy feedback mechanisms.

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