

ON A PROBABILISTIC GRAPH-THEORETICAL METHOD

JAROSLAV NEŠETŘIL AND VOJTĚCH RÖDL

ABSTRACT. We introduce a method by means of which one can simply prove the existence of sparse hypergraphs with large chromatic number. Moreover this method gives the full solution of an Erdős-Ore problem.

Introduction. The existence of sparse (i.e. without short cycles) graphs with a large chromatic number is a classical combinatorial problem. It has been answered affirmatively by probabilistic means by Erdős [1] and for hypergraphs by Erdős and Hajnal [2] and a construction was provided by Lovász [5]. While in [8] the present authors suggested a different construction of these objects, the purpose of this note is to provide a new hopefully simpler probabilistic method.

This method allows to prove existence of sparse hypergraphs which contain a certain ordered subhypergraph for every ordering of its vertices. This has been asked by Erdős [4] in response to a theorem of Ore and Gallai (Corollary 3). The results may be strengthened so as to answer a question of Bollobás as well (Corollary 4).

Preliminaries. A graph G is a couple (V, E) where V is a set (of vertices) and $E \subseteq [V]^2 = \{e \subseteq V; |e| = 2\}$. A k -graph (k -uniform hypergraph) is a couple where V is a set and $E \subseteq [V]^k = \{e \subseteq V; |e| = k\}$. An embedding $f: (V, E) \rightarrow (V', E')$ is a 1-1 mapping which satisfies

(1) f is 1-1;

(2) $e \in E \Leftrightarrow \{f(x); x \in e\} \in E'$. A cycle of length s in a k -graph (V, E) is a sequence

$$x_0, e_1, x_1, e_2, \dots, e_s, x_s, e_0$$

which satisfies $x_i \in e_i$, $x_{i-1} \in e_i$ for $i \in \{1, 2, \dots, s\}$ and $x_s \in e_0$, $x_0 \in e_0$ and there are i, j with $e_i \neq e_j$.

Chromatic number of a k -graph is a minimal number of colours which are sufficient for colouring of vertices in such a way that no edge is monochromatic. It can be proved easily that a k -graph (V, E) does not contain 2-cycles iff $|e \cap e'| \leq 1$ for all $e \neq e'$, $e, e' \in E$ and generally a k -graph does not contain cycles of length $< s$ iff $|\cup E'| \geq (k-1)|E'| + 1$ for every $E' \subset E$, $|E'| < s$ (see e.g. [2], [5]).

The following lemma is a key fact for our method. The lemma may be

Received by the editors May 20, 1977 and, in revised form, January 6, 1978.

AMS (MOS) subject classifications (1970). Primary 05A99, 05C20.

Key words and phrases. Ordering, partition, graph.

© American Mathematical Society 1978

easily proved by probabilistic argument (see e.g. [3]). We sketch a proof for completeness.

LEMMA. *For all positive integers k and s there exists a k -graph (X, E) , $|X| = n$ without cycles of length $< s$ and with $|E| > n^{1+1/s}$ edges for all n sufficiently large.*

PROOF. Let us consider a set \mathfrak{M}_x of all k -graphs $(X, E)_i |X| = n$ with $m = 2\lceil n^{1+1/s} \rceil$ edges. Then

$$|\mathfrak{M}_x| = \binom{n}{m}$$

and the average number of edges contained in cycles of length $< s$ is less than

$$\sum_{s=2}^{s-1} c(s, j) \binom{n}{(k-1)j} \frac{\binom{n}{m-j}}{\binom{n}{m}} = o(n)$$

where $c(s, j) > 0$ is a function of s, j which does not depend on n .

Consequently for all n sufficiently large there exists an example of a k -graph $G = (X, E)$, $|X| = n$, $|E| = 2\lceil n^{1+1/s} \rceil$ such that G contains at most $\lceil n^{1+1/s} \rceil$ edges contained in circuits of length $< s$. After deleting these edges we are left a k -graph with at least $\lceil n^{1+1/s} \rceil$ edges without cycles of length $< s$.

We find it convenient to use this lemma for constructing of special sparse graphs. This will be clear from below.

Results and applications.

THEOREM 1 ([1], [2]). *For all positive integers k, n, s there exists a k -graph (X, F) without cycles of length $< s$ with chromatic number $> n$.*

PROOF. Put $p = n(k - 1) + 1$. Let (X, E) be a p -graph without cycles of length $< s$, $|X| = N$ and with $\lceil N^{1+1/s} \rceil$ edges. Let \mathfrak{G}_x be the family of all k -graphs (X, F) which contain in each edge of E exactly one edge of F . The following holds:

(1) $|\mathfrak{G}_x| = \binom{p}{k}^{\lceil N^{1+1/s} \rceil}$ as (X, E) does not contain 2-cycles the choice of edges of F in edges of E is mutually independent.

(2) G does not contain cycles of length $< s$ for each $G \in \mathfrak{G}_x$.

However the number of k -graphs in the class \mathfrak{G}_x which admit a given n -coloration is less then $(\binom{p}{k} - 1)^{\lceil N^{1+1/s} \rceil} + 1$. The theorem follows as

$$n^N \left(\binom{p}{k} - 1 \right)^{\lceil N^{1+1/s} \rceil} < |\mathfrak{G}_x| = \binom{p}{k}^{\lceil N^{1+1/s} \rceil}$$

for N sufficiently large.

THEOREM 2. *Let $G = ((V, \leq), E)$ be an ordered k -graph (i.e. (V, E) is a*

k-graph; (V, \leq) is a totally ordered set) without cycles of length $< s$. Then there exists a *k*-graph (V', E') without cycles of length $< s$ such that for every ordering (V', \leq') there exists a monotone mapping $f: (V, \leq) \rightarrow (V', \leq')$ which is an embedding $(V, E) \rightarrow (V', E')$.

PROOF. Let (V, E) be a *k*-graph. Denote by m the number of *k*-graphs with the vertex set V which are isomorphic to (V, E) . If $m = 1$ one may put $(V', E') = (V, E)$. (One can show easily that in this case either $E = \emptyset$ or $E = [V]^k$.)

Let us suppose that $m > 1$. Put $|V| = p$. Let (X, U) be a p -graph without cycles of length $< s$ with $[N^{1+1/s}]$ edges. Let \mathcal{G}_x be the class of all *k*-graphs $G = (X, F)$ which satisfy:

- (1) $G|e = (e, F \cap [e]^k) \simeq (V, E)$ for each edge $e \in U$;
- (2) $F = \cup_{e \in U} (F \cap [e]^k)$.

Then (i) G does not contain a cycle of length $< s$ for each $G \in \mathcal{G}_x$;

(ii) $|\mathcal{G}_x| = m^{[N^{1+1/s}]}$.

On the other hand for each ordering \leq of the set X the number of those $G = (X, F) \in \mathcal{G}_x$ for which there exists no monotone embedding $((V, \leq), E) \rightarrow ((X, \leq), F)$ is less than $(m - 1)^{[N^{1+1/s}] + 1}$.

Consequently the set of those $G \in \mathcal{G}_x$ which admit an ordering (X, \leq) for which there is no monotone embedding

$$((V, \leq), E) \rightarrow ((X, \leq), F)$$

is less than $N!(m - 1)^{[N^{1+1/s}]}$. This proves the theorem as the last quantity is $o(m^{[N^{1+1/s}]})$.

COROLLARY 3. For every s there exists a graph $G = (V, E)$ without cycles of length $< s$ which contains for every ordering \leq of its vertices a cycle of length s with the ordering given in the Figure 1.

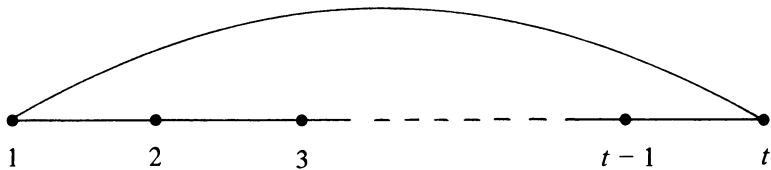


FIGURE 1

(For $s = 3$ the statement is evident, for $s = 4$ it was proved by Ore. Gallai showed that the Grötsch graph is an example for $s = 4$. For $s > 4$ this was asked by Erdős [4].)

COROLLARY 4. For every s there exists a graph (V, E) without cycles of length $< s$ which is not a subgraph of a Hasse-diagram of a partially ordered set.

This follows immediately from Corollary 3 and answers a question of B. Bollobás.

Theorem 2 is related to the concept of ordering property of a class of graphs and hypergraphs defined in [6]. Let \mathfrak{R} be a class of hypergraphs. We say that \mathfrak{R} has ordering property if for every hypergraph $G = (V, E) \in \mathfrak{R}$ there exists a hypergraph $G' = (V', E') \in \mathfrak{R}$ such that for every ordering \prec of V and every ordering \preceq of V' there exists a monotone mapping $f: (V, \prec) \rightarrow (V', \preceq)$ which is an embedding $G \rightarrow G'$.

The ordering property plays an important role in the study of partition properties of classes of hypergraphs. Using this concept one can reformulate the above Theorem 2.

THEOREM 5. *Let \mathfrak{A} be a finite set of 2-connected graphs. Let $\mathfrak{R} = \text{Forb}(\mathfrak{A})$ be the class of all finite graphs which do not contain any graph $\in \mathfrak{A}$ as an induced subgraph. (Thus $G \in \text{Forb}(\mathfrak{A})$ iff there exists no embedding $A \rightarrow G$ for any $A \in \mathfrak{A}$.) Then \mathfrak{R} has the ordering property.*

It is not a simple matter to prove the ordering property of a class \mathfrak{R} constructively. This was done for the class of all finite graphs in [11], for the class of all finite graphs without K_k in [9], for the class of all finite graphs without cycles of length 3, 5, . . . $2k + 1$ in [10]. (The type representation of finite graphs was mainly used.)

Concluding remarks. 1. Theorem 5 has a hypergraph analogue.

2. The ordering property plays an important role in the study of partition properties of graphs. In fact this was the original motivation of this paper.

REFERENCES

- [1] P. Erdős, *Graph theory and probability*, Canad. J. Math. **11** (1959), 34–38.
- [2] P. Erdős and A. Hajnal, *On chromatic number of graphs and set systems*, Acta Math. Acad. Sci. Hungar. **17** (1966), 61–69.
- [3] P. Erdős and J. Spencer, *Probabilistic methods in combinatorics*, Akadémiai Kiadó, Budapest; North-Holland, Amsterdam; Academic Press, New York, 1974.
- [4] P. Erdős, *Some unsolved problems in graph theory and combinatorial analysis*, Combinatorial Mathematics and its Applications (Proc. Conf., Oxford, 1969), Academic Press, London and New York, 1971, pp. 97–99.
- [5] L. Lovász, *On chromatic number of finite set-systems*, Acta Math. Acad. Sci. Hungar. **19** (1968), 59–67.
- [6] J. Nešetřil and V. Rödl, *Type theory of partition properties of graphs*, Recent advances in graph theory, (Proc. Second Czechoslovak Sympos., Prague, 1974), Academia, Prague, 1975, pp. 405–512.
- [7] ———, *Partition of subgraphs*, Recent advances in graph theory (Proc. Second Czechoslovak Sympos., Prague, 1974), Academia, Prague, 1975, pp. 413–423.
- [8] ———, *A short proof of the existence of highly chromatic graphs without short cycles*, J. Combinatorial Theory Ser. B (to appear).

[9] ———, *Ramsey property of classes of graphs without forbidden complete subgraphs*, *J. Combinatorial Theory* **20** (1976), 243–249.

[10] ———, *Ramsey property of graphs without short odd cycles*, *Math. Slovaca* (to appear).

[11] V. Růdl, *Generalization of Ramsey theorem and dimension of graphs*, Thesis, Charles University, Prague, 1973.

KZAA MFF KU, CHARLES UNIVERSITY, SOKOLOVSKÁ 83, 18600, PRAHA 8, CZECHOSLOVAKIA

KM FJFI CVUT, CZECH TECHNICAL UNIVERSITY, HUSOVA 5, 11000 PRAHA 1, CZECHOSLOVAKIA