# On Directional Convexity* 

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#### Abstract

Motivated by problems from calculus of variations and partial differential equations, we investigate geometric properties of $D$-convexity. A function $f: \mathbf{R}^{d} \rightarrow \mathbf{R}$ is called $D$-convex, where $D$ is a set of vectors in $\mathbf{R}^{d}$, if its restriction to each line parallel to a nonzero $v \in D$ is convex. The $D$-convex hull of a compact set $A \subset \mathbf{R}^{d}$, denoted by $\operatorname{co}^{D}(A)$, is the intersection of the zero sets of all nonnegative $D$-convex functions that are zero on $A$. It also equals the zero set of the $D$-convex envelope of the distance function of $A$. We give an example of an $n$-point set $A \subset \mathbf{R}^{2}$ where the $D$-convex envelope of the distance function is exponentially close to zero at points lying relatively far from $\mathrm{co}^{D}(A)$, showing that the definition of the $D$-convex hull can be very nonrobust. For separate convexity in $\mathbf{R}^{3}$ (where $D$ is the orthonormal basis of $\mathbf{R}^{3}$ ), we construct arbitrarily large finite sets $A$ with $\operatorname{co}^{D}(A) \neq A$ whose proper subsets are all equal to their $D$-convex hull. This implies the existence of analogous sets for rank-one convexity and for quasiconvexity on $3 \times 3$ (or larger) matrices.


## 1. Introduction

Let $X$ be a finite-dimensional real vector space (which can be identified with some $\mathbf{R}^{d}$ ), and let $D \subseteq X$ be a set of vectors, which are thought of as directions. A function $f: X \rightarrow \mathbf{R}$ is called $D$-convex if the restriction of $f$ to each line parallel to a nonzero vector in $D$ is a convex function. The $D$-convex hull of a compact set $A \subset X$, denoted by $\operatorname{co}^{D}(A)$, is defined as the intersection of the zero sets of all nonnegative $D$-convex functions $f: X \rightarrow[0, \infty)$ that are zero on $A$. (Later, in Section 3, we give a more direct characterization of the $D$-convex hull. Also, we remark that this $D$-convex hull is called the functional D-convex hull in [MP], in order to distinguish it from the set-theoretical $D$-convex hull. The latter is not considered in the present paper.)

[^0]The "usual" notion of convexity is obtained for $D=X$. Our investigation is mainly motivated by rank-one convexity, which is a special case of $D$-convexity, where $X$ is the space of real $n \times n$ matrices and $D$ is the set of $n \times n$ matrices of rank one. In what follows, this $D$ will be denoted by $r c$. The rank-one convex hull, as an inner approximation to the so-called quasiconvex hull, is important in the theory of partial differential equations and in the calculus of variations and it was studied in a number of papers, among which we mention only a few: [Mo], [Šv], [BFJK], [DKMŠ], [MŠ], and [MŠ]. The lecture notes [Mü1] can serve as a nice and up-to-date introduction to this area.

Another significant special case of $D$-convexity is that with $D$ being the standard orthonormal basis of $\mathbf{R}^{d}$ : the separate convexity (this $D$ will be denoted by $s c$ ). This arises by restricting the rank-one convexity on the subspace of diagonal matrices, and has been considered in this connection [ $T$ ], but it seems natural and interesting in its own right and was independently studied, e.g., in probability theory [AH].

New Results. In the first part of this paper we concentrate on separate convexity in $\mathbf{R}^{d}$. For the usual convexity, the well-known Carathéodory's theorem holds: if $A \subseteq \mathbf{R}^{d}$ and $x$ lies in the convex hull of $A$, then $x$ is in the convex hull of some at most $(d+1)$-point subset of $A$; we say that the Carathéodory number for convexity in $\mathbf{R}^{d}$ is $d+1$. In [MP], it was proved that the Carathéodory number for separate convexity in $\mathbf{R}^{2}$ is 5 . Here we show that the Carathéodory number for separate convexity in dimensions 3 and higher is infinite. As a consequence, the Carathéodory number for rank-one convexity and for quasiconvexity on $3 \times 3$ matrices is infinite as well. We also report other results concerning minimal nontrivial configurations for separate convexity, and mention outcomes of computer experiments with separate convexity performed by Letocha in his M.Sc. thesis [L].

In Section 3 we give a somewhat more direct description of the $D$-convex hull of a set $A$. While the usual definition takes into account all nonnegative $D$-convex functions vanishing on $A$, we show that $\operatorname{co}^{D}(A)$ is actually the zero set of the $D$-convex envelope of the distance function of $A$. (The $D$-convex envelope of a function $f: X \rightarrow \mathbf{R}$, denoted by $C_{D} f$, is defined as the pointwise supremum of all $D$-convex functions $g$ satisfying $g \leq f$ on $X$.) As was pointed out by one of the referees, such a result (for rank-one convexity) has been known to people working on rank-one convexity and quasiconvexity; supposedly it was proved by Yan. The author was unable to find an explicit reference earlier than [MS2], and so although the result is probably not new, it may be useful to include a proof.

The characterization using the distance function suggests an algorithmic approach to computing the $D$-convex hull, via $D$-convexification of the distance function. In Section 4 we show that this approach may be quite problematic in some cases: we exhibit an $n$-point set $A \subset \mathbf{R}^{2}$ and a point $x$ lying relatively far from $\cos ^{s c}(A)$ but such that the separately convex envelope of the distance function has value exponentially close to zero at $x$. Computational experiments indicate bad behavior in this respect (although not as drastic as in the example just mentioned) even for random $n$-point subsets $A$ of the $n \times n \times n$ grid in $\mathbf{R}^{3}$.

In the remaining sections we establish some general properties of $D$-convexity. In Section 5 we show that the $D$-convex envelope of a 1 -Lipschitz function is again 1Lipschitz; we present an argument due to Kirchheim, which is similar to our original proof but simpler. We note that Kirchheim et al. [KKB] recently proved strong results,
somewhat related to the Lipschitz condition, concerning the differentiability of $D$-convex envelopes (as well as quasiconvex envelopes) of differentiable functions. Essentially, they show that if a differentiable function satisfies suitable growth conditions at the infinity, then the envelopes are differentiable as well, while the differentiability may fail without the growth condition.

In [MP], a result on the local behavior of the $D$-convex hull was stated (Corollary 2.9), but, as was pointed out by Kirchheim, it was not sufficiently substantiated. We prove it in Section 6; an independent proof of a somewhat stronger result was recently given by Kirchheim [K].

## 2. New Configurations for Separate Convexity

We call a (finite) set $A \subset X$ nontrivial (for some fixed $D$ ) if $\operatorname{co}^{D}(A) \neq A$, and trivial otherwise. One simple reason for the nontriviality of $A$ is that $A$ contains two points $x, y$ such that the vector $x-y$ is parallel to a direction in $D$ (we say that $A$ has a $D$-connection); for separate convexity in $\mathbf{R}^{d}$, this means that $x$ and $y$ share $d-1$ coordinates.

As was independently discovered by several authors [ Sc ], [AH], [T], [C], a set can be nontrivial without possessing a $D$-connection. In the plane, the four-point configuration $T_{4}=\left\{a_{1}, a_{2}, a_{3}, a_{4}\right\}$ below has a separately convex hull as indicated in the picture (the shaded square and four segments):


Note that this nontrivial configuration is generic nontrivial, meaning that any sufficiently small perturbation of its points again gives a nontrivial set. For separate convexity in $\mathbf{R}^{d}$, a nontrivial configuration in which no two points share the value of any coordinate is necessarily generic (because the combinatorial structure of the separately convex hull is determined by the orderings of the points along the coordinate axes; see Section 2.1 below).

For separate convexity in $\mathbf{R}^{2}$, the situation is relatively simple: any nontrivial set without an $s c$-connection contains a copy of the configuration $T_{4}$ or of its mirror reflection [MP]. Moreover, as was mentioned in the Introduction, the Carathéodory number is 5, meaning that any point in the separately convex hull of $A$ is in the hull of some at most five points of $A$. As we will see in this section, there is no such simple description of nontrivial configurations for separate convexity in higher dimensions. Inclusion-minimal generic nontrivial configuration can be arbitrarily large (and, consequently, the Carathéodory number is infinite), and the number of small configurations is astronomic.

In Sections 2.1-2.4 we always consider separate convexity unless explicitly stated otherwise.

### 2.1. Preliminaries

Let $B \subseteq \mathbf{R}^{d}$ be a set. A point $x \in B$ is called sc-extremal in $B$ if $x$ is contained in no open segment $s \subseteq B$ parallel to one of the coordinate axes.

For the reader's convenience, we briefly review an algorithmic description of the separately convex hull of a finite set $A \subset \mathbf{R}^{d}$ derived in [MP]. For $i=1,2, \ldots, d$, let $x_{i}(A)=\left\{x_{i}(a): a \in A\right\}$, where $x_{i}(a)$ denotes the $i$ th coordinate of $a$, and let $\operatorname{grid}(A)=$ $x_{1}(A) \times x_{2}(A) \times \cdots \times x_{d}(A)$. For a point $x \in G=\operatorname{grid}(A)$, let $a^{i+}$ (resp. $a^{i-}$ ) denote the point of $G$, all of whose coordinates except the $i$ th coincide with those of $a$, and whose $i$ th coordinate is the successor (resp. predecessor) of $x_{i}(a)$ in $x_{i}(G)$ (thus, $a^{i+}$ or $a^{i-}$ need not exist for "border" points of $G$ ). Let $B \subseteq G$ and $a \in B$; we call a point $a \in \operatorname{grid}(A) \operatorname{grid}$-extremal in $B$ if, for each $i=1,2, \ldots, d$, at least one of $a^{i+}, a^{i-}$ either does not exist or does not belong to $B$; intuitively, $a$ is a "local corner" of $B$.

Given a finite $A \subset \mathbf{R}^{d}$, put $B_{0}=\operatorname{grid}(A)$, and for $j=0,1,2, \ldots$, if $B_{j}$ contains a grid-extremal point $b \notin A$, set $B_{j+1}=B_{j} \backslash\{b\}$ and continue with the next $j$. This procedure terminates with a set $B_{j_{0}}$ with all grid-extremal points lying in $A$, and this set describes the separately convex hull of $A$. Namely, let an elementary box for grid $(A)$ be a Cartesian product of the form $I_{1} \times I_{2} \times \cdots \times I_{d}$, where each $I_{i}$ is either $\left\{x_{i}\right\}$ for some $x_{i} \in x_{i}(A)$ or $\left[x_{i}(a), x_{i}\left(a^{i+}\right)\right]$ for an $a \in \operatorname{grid}(A)$. The box complex of $B \subseteq \operatorname{grid}(A)$ consists of the elementary boxes whose corners all lie in $B$. Then, as shown in [MP], $\operatorname{co}^{s c}(A)$ is the union of the box complex of the set $B_{j_{0}}$ obtained by the above algorithm.

As a consequence of this algorithmic description, we get that if $B=\operatorname{co}^{s c}(A)$ for $A$ finite, then all $s c$-extremal points of $B$ belong to $A$ and $B$ is the separately convex hull of its extremal points.

### 2.2. Generic Nontrivial Configurations in All Dimensions

The existence of generic nontrivial configurations for separate convexity in $\mathbf{R}^{3}$ was established in [MP]; a 20-point configuration was exhibited. Its nontriviality was verified by applying the above algorithm. Here we generalize the idea of that construction and we present a systematic inductive construction in any dimension.

Theorem 2.1. For any $d \geq 2$, there exists a finite generic nontrivial configuration $A_{d}$ for separate convexity in $\mathbf{R}^{d}$.

Proof. We proceed by induction on the dimension $d$. We need a slightly stronger statement; to state the additional condition, we use the following definition. Let $a \in A$ be an $s c$-extremal point of $\cos ^{s c}(A)$, and let $u \in\left\{e_{1},-e_{1}, e_{2},-e_{2}, \ldots, e_{d},-e_{d}\right\}$ be a direction of some coordinate semiaxis. We call $u$ an inward direction at $a$ if an open neighborhood $U$ of $a$ exists such that, for any $x \in U$, the ray $\{x+t u: t \geq 0\}$ intersects $\cos ^{s c}(A)$.


Fig. 1. Illustration of the construction of $A_{d+1}$.

Our inductive hypothesis is the claim of the theorem with the additional conditions that the origin 0 lie in the interior of $\cos ^{s c}\left(A_{d}\right)$ and that every point $a \in A_{d}$ has an inward direction. The basis of the induction is provided by the configuration $T_{4}$ in the plane.

Suppose that the claim has been proved for $d$, and we want to construct the configuration $A_{d+1} \subset \mathbf{R}^{d+1}$. We refer to the direction of the $x_{d+1}$-axis as "vertical," and to hyperplanes perpendicular to the $x_{d+1}$-axis as "horizontal."

As a first step, we place copies of $A_{d}$ into the horizontal hyperplanes $x_{d+1}=1$ and $x_{d+1}=-1$, and we perturb the points of each copy vertically. Formally, for $a=$ $\left(a_{1}, a_{2}, \ldots, a_{d}\right) \in A_{d}$, we put $a_{+}=\left(a_{1}, a_{2}, \ldots, a_{d}, 1+z_{a}\right)$ and $a_{-}=\left(a_{1}, a_{2}, \ldots, a_{d}\right.$, $-1-z_{a}$ ), where the $z_{a}$ 's are pairwise distinct positive real numbers, and we set $A_{+}=$ $\left\{a_{+}: a \in A_{d}\right\}$ and similarly for $A_{-}$.

Next, we choose pairwise distinct numbers $t_{a} \in(-1,1)$ for $a \in A_{d}$. Let $\varepsilon>0$ be a sufficiently small parameter. Let $P_{a}=\left([0, \varepsilon]^{d}+a\right) \times\left\{t_{a}\right\}$ be a small horizontal "plate" lying at height $t_{a}$ with a corner on the vertical segment connecting $a^{+}$and $a^{-}$(see Fig. 1). Moreover, let $A_{+}^{\prime}=A_{+}+(\varepsilon, \varepsilon, \ldots, \varepsilon, 0)$ be a horizontal translate of $A_{+}$, and put

$$
B=A_{+}^{\prime} \cup A_{-} \cup \bigcup_{a \in A_{d}} P_{a} .
$$

We observe that for any horizontal hyperplane $h=\left\{x_{d+1}=z\right\}$ with $-1 \leq z \leq 1$, $h \cap \cos ^{s c}(B)$ contains an approximate copy $A(h)$ of $A_{d}$ with each point perturbed by at most $\varepsilon$. By the stability of the combinatorial structure of $\operatorname{co}^{s c}\left(A_{d}\right)$ under sufficiently small perturbations, we see that the combinatorial structure of $\operatorname{co}^{s c}(A(h))$ is the same for all $h$ (if $\varepsilon$ is sufficiently small), and an inward direction $u_{a}$ at a point $a \in A_{d}$ remains an inward direction for the corresponding point in each $A(h)$ ( $u_{a}$ is a direction in $\mathbf{R}^{d}$, but from now on we interpret it in $\mathbf{R}^{d+1}$ by appending a zero $x_{d+1}$-coordinate).

For each $a \in A_{d}$, let $c_{a}=\left(a_{1}, a_{2}, \ldots, a_{d}, t_{a}\right)-u_{a}$ be the point reached from the corner of the horizontal plate $P_{a}$ by going one unit against the direction $u_{a}$ (in the horizontal hyperplane $x_{d+1}=t_{a}$ ). Previous considerations show that there exists an open ball $U_{a}$ centered at $c_{a}$, whose radius is independent of $\varepsilon$ (provided that $\varepsilon>0$ is sufficiently small), such that all rays $\left\{x+t u_{a}: t \geq 0\right\}$ for $x \in U_{a}$ intersect $\operatorname{co}^{s c}(B)$.

Let $g_{a}$ be the hyperplane perpendicular to $u_{a}$ and containing the point $c_{a}$. Identify $\mathbf{R}^{d}$ with $g_{a}$ so that 0 is placed into $c_{a}$ and the axes directions in $\mathbf{R}^{d}$ remain parallel to the axes directions in $\mathbf{R}^{d+1}$, and let $M_{a}$ be a copy of $A_{d}$ in $g_{a}$ scaled by the factor of $\sqrt{\varepsilon}$.

In this way, we may assume that $M_{a}$ is contained in the ball $U_{a}$ and, moreover, that the union of rays emanating from points of $\operatorname{co}^{s c}\left(M_{a}\right)$ in the direction $u_{a}$ contains the whole plate $P_{a}$ in its interior. As a final step of the construction, shift each point $b \in M_{a}$ by $-y_{b} u_{a}$, where the $y_{b}$ 's are pairwise distinct positive real numbers. This yields a set $\tilde{M}_{a}$.

Set

$$
C=B \cup \bigcup_{a \in A_{d}} \tilde{M}_{a},
$$

and define $A_{d+1}$ as the set of all $s c$-extremal points of $\operatorname{co}^{s c}(C)$. By the remark at the end of Section 2.1, we have $\cos ^{s c}(C)=\operatorname{co}^{s c}\left(A_{d+1}\right)$, and $A_{d+1}$ consists of points of $A_{-} \cup A_{+}^{\prime} \cup \bigcup_{a \in A_{d}} \tilde{M}_{a}$ plus possibly $s c$-extremal points of the $P_{a}$ 's. However, since $\operatorname{co}^{s c}(C)$ contains each $P_{a}$ in its interior, we get $A_{d+1} \subseteq A_{-} \cup A_{+}^{\prime} \cup \bigcup_{a \in A_{d}} \tilde{M}_{a}$. So $A_{d+1}$ is finite, nontrivial with 0 lying in the interior of $\operatorname{co}^{s c}\left(A_{d+1}\right)$, and it is easily checked from the construction that no two of its points share a common coordinate hyperplane, and hence $A_{d+1}$ is also generic. Finally, for the points of $A_{-}$and $A_{+}^{\prime}$, inward directions are $(0,0, \ldots, 0,1)$ and $(0,0, \ldots, 0,-1)$, respectively (because an open neighborhood of each $P_{a}$ is contained in $\operatorname{co}^{s c}\left(A_{d+1}\right)$ ), and, for each $b \in \tilde{M}_{a}, u_{a}$ can be chosen as an inward direction. This finishes the proof of Theorem 2.1.

### 2.3. Computer Experiments in Dimension 3

The algorithm for the three-dimensional separately convex hull reviewed in Section 2.1 was fine-tuned and implemented by Letocha [L]. The correctness of the implementation was checked by the comparison of many results with an earlier, slower implementation by Matoušek. Letocha noticed that the generic nontrivial configuration constructed in [MP] is not inclusion-minimal, and an 18-point generic nontrivial configuration can be obtained from it by removing two suitable points. This is also the smallest generic nontrivial configuration known so far.

Letocha conducted extensive computer search for inclusion-minimal generic nontrivial configurations. At each experiment, independent random permutations $\pi_{1}$ and $\pi_{2}$ of $\{1,2, \ldots, n\}$ were generated, and the (generic) set

$$
\begin{equation*}
A=\left\{\left(i, \pi_{1}(i), \pi_{2}(i)\right): i=1,2, \ldots, n\right\} \tag{1}
\end{equation*}
$$

was considered. (For example, for $n=100$, the computation of the separately convex hull for such a set took about 0.2 s on a Pentium II, 300 MHz machine.) If $A$ turned out to be nontrivial (successful experiment), it was checked for inclusion-minimality, and as soon as a point $a \in A$ with $A \backslash\{a\}$ nontrivial was found, it was removed. This was repeated until an inclusion-minimal nontrivial set was obtained. As noted in [MP], the existence of a single generic nontrivial configuration implies that the probability of success in this experiment tends to 1 as $n \rightarrow \infty$. However, it turned out that the probability of success is quite large even for fairly small $n$; for $n=65$ it is (estimated to be) slightly over 0.5 , and for $n=78$ it exceeds 0.9 .

Minimal configurations with sizes between 18 and 28 were discovered by this method. The 18 -point configurations were most frequent; for $n=66$, with about $55 \%$ of success-
ful experiments, about $36 \%$ of the successful experiments led to 18 -point configurations, $33 \%$ to 19 -point ones, and $19 \%$ to 20 -point ones. From such data, one can estimate from below the number of distinct minimal nontrivial configurations of the "canonical" form (1). For example, for $n=40$, an 18 -point configuration was observed in about $0.12 \%$ of the cases (in $10^{6}$ experiments). Thus, up to the small statistical uncertainty, we get that the probability $P(40,18)$ of a random set (1) containing a minimal nontrivial 18 -point configuration as a subset is at least 0.0012 . A random 40-point set (1) contains $\binom{40}{18} 18$-point subsets, each of which can be regarded as a random 18 -point set of the form (1). These random subsets are not independent, but certainly we have $P(40,18) \leq\binom{ 40}{18} P(18)$, where $P(18)$ is the probability of a random 18-point set (1) being minimal nontrivial. Since the total number of 18 -point sets $(1)$ is $(18!)^{2}$, we can conclude that the number of distinct 18 -point minimal generic nontrivial configurations is at least about $(18!)^{2} P(40,18) /\binom{40}{18} \geq 4 \times 10^{17}$.

Generic nontrivial configurations of 17 or fewer points were never encountered in the experiments, and they must be much more rare than those with 18 points, if they exist at all.

### 2.4. Arbitrarily Large Minimal Configurations

First we exhibit an arbitrarily large inclusion-minimal nontrivial set in $\mathbf{R}^{3}$ which is not generic. For $n=1,2, \ldots$, the set $A_{n}$ has $3 n+3$ points $d, e, f$ and $a_{i}, b_{i}, c_{i}, i=$ $1,2, \ldots, n$. The construction for $n=3$ is drawn in Fig. 2. The cube drawn by thin line is included solely for better visualization; the points are drawn by dots and the set $B_{3}=\operatorname{co}^{s c}\left(A_{3}\right)$ by thick lines. For other $n$, the construction is analogous, but the "stairs" produced by $a_{i}, b_{i}, c_{i}$ are made smaller and $n$ of them are put in. Using the algorithm in Section 2.1, it is not difficult to check that $B_{n}=\cos ^{s c}\left(A_{n}\right)$. (The inclusion $B_{n} \subseteq \cos ^{s c}\left(A_{n}\right)$ is especially easy, since all $s c$-extreme points of $B_{n}$ lie in $A_{n}$.) If we remove, for example, the point $c_{3}$ from $A_{3}$, the algorithm allows us to remove the segment of $B_{3}$ ending in $c_{3}$, and then, successively, the segments ending in $b_{3}, a_{3}, c_{2}, b_{2}, \ldots, e$,


Fig. 2. The minimal nontrivial set $A_{3}$ and its $s c$-hull.
and $d$. The set $A_{3} \backslash\left\{c_{3}\right\}$ is trivial, and the situation with removing any other point is entirely analogous.

Proposition 2.2. The Carathéodory numberfor the separate convexity in $\mathbf{R}^{d}$, ford $\geq 3$, is infinite. In fact, arbitrarily large (finite) inclusion-minimal nontrivial configurations exist.

Recently, Müller proved that if $\mathbf{R}^{d}$ is identified with the space Diag $_{d}$ of $d \times d$ diagonal matrices in the obvious manner $\left(\left(x_{1}, x_{2}, \ldots, x_{d}\right)\right.$ becomes the matrix with $x_{1}, x_{2}, \ldots, x_{d}$ on the diagonal and zeros elsewhere), and if $f:$ Diag $_{d} \rightarrow \mathbf{R}$ is any separately convex function, then for any $\varepsilon>0$ and a compact set $K \subset \operatorname{Diag}_{d}$, a quasiconvex function $g$ on the space $M^{d \times d}$ of all $d \times d$ matrices exists with $|f(x)-g(x)|<\varepsilon$ for all $x \in K$ ([Mü2] deals with the case $d=2$ and announces the result for an arbitrary $d$ ). Consequently, for compact $A \subseteq$ Diag $_{d}$, the separately convex hull of $A$ within Diag ${ }_{d}$ equals the quasiconvex hull, and also the rank-one convex hull, of $A$ in $M^{d \times d}$. (We remark that the result just mentioned is not obvious even for rank-one convexity.) Therefore, we get

Corollary 2.3. The Carathéodory number for quasiconvexity, as well as for rank-one convexity, on $d \times d$ matrices, $d \geq 3$, is infinite, and arbitrarily large (finite) inclusionminimal nontrivial configurations exist.

The configuration $A_{n}$ constructed above is not generic, but an arbitrarily large minimal generic nontrivial configuration for separate convexity can be obtained as well (for rank-one convexity or quasiconvexity, the existence of such configurations is open at present). The idea is to replace each point of $A_{n}$ by a small (perturbed) copy of the planar configuration $T_{4}$. We first note that if the horizontal rectangle $R$ as in Fig. 3 lies in the hull of some set, and the four points $a_{1}, \ldots, a_{4}$ are in the set, then the rectangle $R_{1}$ also lies in the hull. Then such rectangles (and the corresponding 4-tuples) can be arranged cyclically, similar to the segments in Fig. 2, as depicted in Fig. 4 (the position of the 4 -tuples is indicated schematically by thick segments). The resulting configuration is generic nontrivial. It is not minimal (it turns out that two points suffice at each turn of the "stairs"; Fig. 5 shows a minimal subconfiguration obtained for $n=3$ ), but if we delete any of the 4 -tuples, we get a trivial configuration, and hence any minimal nontrivial subset has at least $3(n+1)$ points.


Fig. 3. Four suitable points and $R$ generate the rectangle $R_{1}$.


Fig. 4. Constructing arbitrarily large generic nontrivial configuration.

We note that by applying this construction with $n=1$ and selecting a minimal nontrivial subset, we arrive at the nicely symmetric 18 -point configuration (the smallest known size) shown in Fig. 6. The picture displays the separately convex hull as the appropriate box complex; the two-dimensional elementary boxes are shown semitransparent. It is also remarkable that, unlike the usual convex hulls, the separately convex hulls in dimension 3 need not be contractible.


Fig. 5. A minimal 30-point generic nontrivial configuration and its sc-hull.


Fig. 6. An 18 -point generic configuration with its $s c$-hull.

## 3. Envelope of the Distance Function Defines the Hull

Theorem 3.1. Let $D \subseteq \mathbf{R}^{d}$ be a set of directions containing a basis of $\mathbf{R}^{d}$. Let $A \subset \mathbf{R}^{d}$ be a compact set. Let $\delta_{A}$ be the function giving distance from $A$; that is, $\delta_{A}(x)=$ $\inf _{y \in A}\|x-y\|$. Let $Z$ be the zero set of $C_{D} \delta_{A}$. Then $Z=\operatorname{co}^{D}(A)$.

Proof. The inclusion $\operatorname{co}^{D}(A) \subseteq Z$ is clear. To prove the opposite inclusion, consider a point $x_{0} \notin \operatorname{co}^{D}(A)$. This means that a nonnegative $D$-convex function $f$ exists with $f\left(x_{0}\right)>0$ and $f(A)=0$. Our goal is to produce a $D$-convex function $g$ with $g\left(x_{0}\right)>0$ and satisfying $g \leq \delta_{A}$ everywhere. Then we will also have that the $D$-convex envelope of $\delta_{A}$ majorizes $g$, and in particular it cannot be zero at $x_{0}$.

Choose the coordinate system so that $0 \in A$. Let $B(0, R), R \geq 1$, be a (closed) ball containing both $x_{0}$ and $A$. Set

$$
\eta=\frac{1}{4} \inf \left\{\frac{\delta_{A}(x)}{\max \left(\delta_{A}(x), f(x)\right)}: x \in B(0,2 R) \backslash A\right\}
$$

We want to prove that $\eta>0$. We recall that the function $f$, being $D$-convex, is locally Lipschitz [MP, Observation 2.3], and hence Lipschitz on any compact set. Let $C$ be the Lipschitz constant of $f$ on $B(0,2 R)$. If $x \in B(0,2 R)$ lies at distance $t>0$ from $A$, then there is a point $a \in A$ at distance $t$ from $x$, and since $f(a)=0$ we have $f(x) \leq C t$. From this we get $\eta \geq 1 / 4 C$.

Now we know that $\eta f$ is a $D$-convex function nonzero at $x_{0}$ satisfying $\eta f \leq \delta_{A}$ everywhere on the ball $B(0, R)$; we still need to extend it on the whole $\mathbf{R}^{d}$ so that it remains below $\delta_{A}$ everywhere. To this end, we use a standard trick from convex analysis for extending a convex function defined on a ball.

Define a function $g$ by setting

$$
g(x)=\left\{\begin{array}{lll}
\max (\eta f(x),\|x\|-R) & \text { for } & \|x\| \leq 2 R \\
\|x\|-R & \text { for } & \|x\|>2 R
\end{array}\right.
$$

First we note that on $B(0, R), g$ coincides with $\eta f$, and hence $g\left(x_{0}\right)>0$. We also have $g \leq \delta_{A}$ everywhere (because $\eta f \leq \delta_{A}$ on $B(0,2 R)$ and $\|x\|-R \leq \delta_{A}(x)$ ). It remains to show that $g$ is $D$-convex. Clearly it is $D$-convex on $B(0,2 R)$, being a maximum of two $D$-convex functions there. We note that for all $x$ with $\frac{3}{2} R \leq\|x\| \leq 2 R$ we have

$$
\eta f(x) \leq \frac{1}{4} \delta_{A}(x) \leq \frac{1}{4}\|x\| \leq \frac{1}{2} R \leq\|x\|-R .
$$

This means that in the annulus $\frac{3}{2} R \leq\|x\| \leq 2 R, g(x)$ coincides with $\|x\|-R$, and from this it is routine to check the $D$-convexity of $g$ on the whole $\mathbf{R}^{d}$.

## 4. Nonrobustness of the $\boldsymbol{D}$-Convex Hull

For separate convexity, there is a simple exact algorithm for computing the $s c$-hulls of finite sets, but for rank-one convexity (or for any $D$ with more than $d$ directions in $\mathbf{R}^{d}$ ), no such algorithm is known so far. In view of Theorem 3.1, it seems natural to try to approximate the $D$-convex hull by approximately computing the $D$-convex envelope of the distance function and taking the "near-zero" set. However, even for separate convexity in the plane, this is generally unrealistic, because enormous accuracy would be required.

Proposition 4.1. For each $n \geq 1$, there exist: an ( $n+1$ )-point set $A \subset \mathbf{R}^{2}$ contained in the $m \times m$ integer grid with $m=O(n)$ and with $\operatorname{co}^{s c}(A)=A$, and two points $b_{0}$ and $b_{n}$, both at distance 1 from $A$, such that any separately convex function $f$ that is zero on A satisfies $f\left(b_{n}\right) \leq n^{-n} f\left(b_{0}\right)$.

Proof. The construction is shown in Fig. 7. There is an auxiliary gray square with side $n$ in the middle, and the points $a_{0}, a_{1}, \ldots, a_{n}$ are placed in a spiral-like configuration around the square; the scaling is such that the distance of $a_{0}$ and $b_{0}$ is 1 , as well as each of the distances $a_{i} b_{i}$. Since $f\left(a_{1}\right)=0$ and the distance $b_{0} a_{1}$ is at least $n$, the convexity of $f$ on the line $b_{0} a_{1}$ implies $f\left(b_{1}\right) \leq(1 / n) f\left(b_{0}\right)$, and induction (along the indicated lines) yields $f\left(b_{i}\right) \leq n^{-i} f\left(b_{0}\right)$. Finally, the triviality of $A$ is easy (one can check that there is no $T_{4}$ configuration, or apply the algorithm).


Fig. 7. A configuration with nonrobust hull.

Of course, one can hope that configurations with this bad behavior are exceptional and that we suffice with much smaller precision for computing the $D$-convex envelope for "usual" examples. Computational experiments, described next, indicate that one has to be careful even with not too large random configurations.

For various values of $n$, random $n$-point sets $A \subset \mathbf{R}^{3}$ of the form (1) were generated. Recall that they are subsets of the grid $G=\{1,2, \ldots, n\}^{3}$. For such an $A$, the function $f_{A}: G \rightarrow\{0,1\}$, with $f_{A}(x)=0$ for $x \in A$ and $f_{A}(x)=1$ otherwise, was considered. As shown in [MP], the points of $G \cap \cos ^{s c}(A)$ are exactly the zero set of the separately convex envelope of $f_{A}$ on $G$ (where the separately convex envelope on $G$ is the largest function $G \rightarrow \mathbf{R}$ that is below $f_{A}$ and satisfies the convexity condition for any triple of points of $G$ lying on a line parallel to a coordinate axis). The function $f_{A}$ was separately convexified on the grid $G$ by a straightforward iterative algorithm: convexify along the lines parallel to the $x$-axis, then along the $y$-axis, then along the $z$-axis, and repeat until the maximum change of the function's value in a single iteration drops below a small threshold (chosen as $10^{-13}$ for double-precision arithmetic). Up to small rounding errors, this algorithm provides an upper bound on the values of $C_{s c} f_{A}$ on $G$ (and, hopefully, should provide good approximation to the actual values of the separately convex envelope, but no error bound seems to be available).

A measure of the accuracy required for correct computation of $\cos ^{s c}(A)$ by this method is the smallest value of $C_{s c} f_{A}(x)$ for $x \in G \backslash \cos ^{s c}(A)$ (where $\operatorname{co}^{s c}(A)$ was determined by the exact combinatorial algorithm). In the experiments, this value was typically quite small even for moderate values of $n$. For example, while it was typically between $10^{-3}$ and $10^{-4}$ for $n=20$, already for $n=40$ it was usually below $10^{-7}$ and values as small as $2 \times 10^{-11}$ appeared in a few cases. For $n=50$, the values were
often below $10^{-13}$ and cannot be considered reliable anymore with the double-precision computations.

Still, knowing that the value of the $D$-convex envelope of the distance function of $A$ (or of some other suitable function) is reasonably large at some point $x$, we can conclude that $x \notin \operatorname{co}^{D}(A)$. Approximate computation of $D$-convex envelopes might thus yield at least a reasonable outer approximation of the $D$-convex hull. Nonetheless even here there is a problem with controlling the error of the approximation of the envelope. The most natural method of computing the $D$-convex envelope (employed above for separate convexification on a grid), namely iterative one-dimensional convexifications along the directions in $D$, is likely to provide an upper bound on the values of the envelope. However, controlling the error, and getting a reliable bound from below, appears challenging. As was remarked above, no error bounds seem to be available even for separate convexification on a grid.

## 5. Lipschitz Constant is Preserved by the D-Convex Envelope

Recall that for a function $f: \mathbf{R}^{d} \rightarrow \mathbf{R}$ and a direction vector $v, C_{\{v\}} f$ denotes the convex envelope of $f$ taken in direction $v$; that is, $f$ is convexified (independently) along each line parallel to $v$. For simplicity, we write only $C_{v} f$ for $C_{\{v\}} f$.

Lemma 5.1. Let $v \in \mathbf{R}^{d}$ be a nonzero direction vector. Let $f$ be a real 1-Lipschitz function defined on $\mathbf{R}^{d}$. Then $C_{v} f$ is 1-Lipschitz as well. Here "1-Lipschitz" may be taken with respect to an arbitrary translation-invariant metric on $\mathbf{R}^{d}$.

The following simple proof was communicated to me by Kirchheim; my previous formulation was more complicated.

Proof. For brevity, denote the function $C_{v} f$ by $g$. Let $x, y \in \mathbf{R}^{d}$ be two arbitrary points; it suffices to prove

$$
\begin{equation*}
g(x) \leq g(y)+\rho(x, y), \tag{2}
\end{equation*}
$$

where $\rho$ is the considered translation-invariant metric. Let $\varepsilon>0$ be arbitrary, and let $\ell_{y}$ be the line parallel to $v$ containing $y$. Since the point $(y, g(y)+\varepsilon)$ is above the convex envelope of $f$ restricted to $\ell_{y}$, there are two points $y_{1}, y_{2} \in \ell_{y}$ such that $y$ lies between $y_{1}$ and $y_{2}$ and $g(y)+\varepsilon>t f\left(y_{1}\right)+(1-t) f\left(y_{2}\right), t \in[0,1]$. Then we have

$$
\begin{aligned}
g(x) & =g(y+(x-y)) & & \\
& \leq t f\left(y_{1}+(x-y)\right)+(1-t) f\left(y_{2}+(x-y)\right) & & \text { (convexity of } \left.g \text { on } \ell_{x}\right) \\
& \leq t f\left(y_{1}\right)+(1-t) f\left(y_{2}\right)+\rho(x, y) & & \text { (as } f \text { is 1-Lipschitz) } \\
& <g(y)+\varepsilon+\rho(x, y) . & &
\end{aligned}
$$

Since $\varepsilon>0$ was arbitrary, (2) is proved.
Corollary 5.2. Let $f$ be a 1-Lipschitz real function defined on $\mathbf{R}^{d}$. Then $C_{D} f$ is 1Lipschitz as well.

## Proof. Set

$$
\bar{f}=\inf \left\{C_{v_{1}} C_{v_{2}} \cdots C_{v_{n}} f: v_{1}, v_{2}, \ldots, v_{n} \in D, n=1,2,3, \ldots\right\}
$$

(a pointwise infimum). Clearly $C_{D} f \leq \bar{f}$, and it is straightforward to verify that $\bar{f}$ is a $D$-convex function, hence $\bar{f}=C_{f}$. At the same time, all the functions $C_{v_{1}} C_{v_{2}} \cdots C_{v_{n}} f$ are 1-Lipschitz by Lemma 5.1, and hence $C_{D} f$ is 1-Lipschitz as well.

There is an interesting consequence for $D$-convex envelopes of distance functions to sets.

Corollary 5.3. Let $A \subseteq \mathbf{R}^{d}$ be a set and let $B=\operatorname{co}^{D}(A)$. Then $C_{D} \delta_{A}=C_{D} \delta_{B}$, where $\delta_{X}$ denotes the distance function of a set $X$.

Proof. Since $A \subseteq B$, we have $\delta_{A} \geq \delta_{B}$ and hence also $C_{D} \delta_{A} \geq C_{D} \delta_{B}$. To see the opposite inequality, we note that $C_{D} \delta_{A}$ is a $D$-convex function and hence it is zero on $B$. Moreover, it is 1 -Lipschitz by Corollary 5.2, and therefore $C_{D} \delta_{A} \leq \delta_{B}$. Taking $D$-convex envelopes on both sides of the last inequality yields $C_{D} \delta_{A} \leq C_{D} \delta_{B}$.

## 6. A Locality Result

The following result was claimed in [MP] as Corollary 2.9:
Proposition 6.1. Let $A \subseteq \mathbf{R}^{d}$ be contained in a (functionally) $D$-convex set $C$, which is a disjoint union of compact sets $C_{1}, \ldots, C_{k}$. Then $\operatorname{co}^{D}(A)=\bigcup_{i=1}^{k} \operatorname{co}^{D}\left(A \cap C_{i}\right)$.

As was pointed out by Kirchheim (private communication, 1998), this is probably not an immediate consequence of the previous theorem in [MP] (which states that if $C_{1}, C_{2} \subseteq \mathbf{R}^{d}$ are disjoint compact sets with $C_{1} \cup C_{2}$ being $D$-convex, then $C_{1}$ and $C_{2}$ are $D$-convex as well). Here we give a full proof. The result was independently proved, together with some other related properties of $D$-convex hulls, by Kirchheim [K].

Proof. It suffices to prove the following statement: Let $B, C \subset \mathbf{R}^{d}$ be disjoint compact sets whose union is $D$-convex, and let $K \subseteq B$; then $B \cap \operatorname{co}^{D}(K \cup C)=\operatorname{co}^{D}(K)$. Indeed, in the situation of Proposition 6.1, we set $B=C_{1}, C=C_{2} \cup \cdots \cup C_{k}, K=A \cap C_{1}$, and we use the monotonicity of the $D$-convex hull.

Put $f=C_{D} \delta_{B \cup C}$. Fix $\beta>0$ such that $B_{2 \beta} \cap C=\emptyset$ (where $B_{\varepsilon}$ denotes the $\varepsilon$ neighborhood of $B$ ), and let $S=\overline{B_{2 \beta} \backslash B_{\beta}}$. By Theorem 3.1, $f$ is positive on $S$, and so it is bounded away from zero there, by the compactness of $S$.

Let $f_{K}=C_{D} \delta_{K}$. This $f_{K}$ is positive on the compact set $S$ (since $c^{D}(K) \subseteq B$ ). Choose $\eta>0$ so that $\eta f_{K} \leq f$ on $S$. Define a function $g$ as follows:

$$
g= \begin{cases}\max \left(\eta f_{K}, f\right) & \text { on } B_{2 \beta} \\ f & \text { elsewhere }\end{cases}
$$

Since $f=0$ on $B \cup C$, the zero set of $g$ contains $C$, and its intersection with $B$ equals the zero set of $f_{K}$, i.e., $\operatorname{co}^{D}(K)$. It remains to check that $g$ is $D$-convex. On $B_{2 \beta}, g$ is $D$-convex as the maximum of two $D$-convex functions. Outside of $B_{\beta}$, we have $g=f$.

If a line $\ell$ intersects both $B_{\beta}$ and the complement of $B_{2 \beta}$, it shares a segment of length at least $\beta$ with $S$; consequently, $g$ is $D$-convex everywhere on $\mathbf{R}^{d}$.

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