# On Double and Multiple Interval Graphs 

William T. Trotter, Jr. UNIVERSITY OF SOUTH CAROLINA

Frank Harary<br>UNIVERSITY OF MICHIGAN


#### Abstract

In this paper we discuss a generalization of the familiar concept of an interval graph that arises naturally in scheduling and allocation problems. We define the interval number of a graph $G$ to be the smallest positive integer $t$ for which there exists a function $f$ which assigns to each vertex $u$ of $G$ a subset $f(u)$ of the real line so that $f(u)$ is the union of $t$ closed intervals of the real line, and distinct vertices $u$ and $v$ in $G$ are adjacent if and only if $f(u)$ and $f(v)$ meet. We show that (1) the interval number of a tree is at most two, and (2) the complete bipartite graph $K_{m, n}$ has interval number $\lceil(m n+1) /(m+n)\rceil$.


## 1. INTRODUCTION

A graph $G$ is called an interval graph if there is a function $f$ that assigns to each vertex $u$ of $G$ a closed interval of the real line $R$ so that distinct vertices $u, v$ of $G$ are adjacent if and only if $f(u) \cap f(v) \neq \varnothing$. Structural characterizations of interval graphs have been provided by Lekkerkerker and Boland [7] who specified the forbidden subgraphs, Gilmore and Hoffman [2] in terms of cycles, and Fulkerson and Gross [1] in terms of matrices. Definitions not given here can be found in Ref. 5.
In this paper, we consider a generalization of the concept of an interval graph; we are motivated by scheduling and allocation problems that arise when a graph is used to model constraints on interactions between
components of a large scale system. For a graph $G$, we define* the interval number of $G$, denoted $i(G)$, as the smallest positive interger $t$ for which there exists a function $f$ which assigns to each vertex $u$ of $G$ a subset $f(u)$ of $R$ which is the union of $t$ (not necessarily disjoint) closed intervals of $R$ and distinct vertices $u, v$ of $G$ are adjacent if and only if $f(u) \cap f(v) \neq \varnothing$. The function $f$ is called a $t$-representation of $G$. Thus $G$ is an interval graph if and only if its interval number is one. Obviously every graph $G$ with $p$ vertices has an interval number $i(G) \leq p-1$, and thus $i(G)$ is well defined.

A number $m$ is called an upper bound for a representation $f$ of a graph $G$ when $m>r$ for every number $r$ in $f(u)$ and every vertex $u$ of $G$.

We will frequently find it convenient to impose an additional restriction on a representation of a graph. A $t$-representation $f$ of a graph $G$ is said to be displayed if for every vertex $u$ of $G$, there exists an open interval $I_{u}$ contained in $f(u)$ so that $I_{u} \cap f(v)=\varnothing$ for every vertex $v$ in $G$ with $u \neq v$.

Recall that for any tree $T$, the tree $T^{\prime}$ is obtained by removing all the endvertices of $T$. A caterpillar is a tree $T$ for which $T^{\prime}$ is a path. It was noted in Harary and Schwenk [6] that $T$ is a caterpillar if and only if $T$ does not contain the subdivision graph of $K_{1,3}$ as a subtree.

Theorem 1. If $T$ is a tree, then $i(T)=1$ if $T$ is a caterpillar and $i(T)=2$ if it is not.

Proof. If $T$ is a tree and does not contain the subdivision graph of $K_{1,3}$ as a subtree, then it follows from the forbidden subgraph characterization of Ref. 7 that $T$ is an interval graph. On the other hand, if $T$ contains this subdivision graph, then $T$ is not an interval graph and $i(T) \geq 2$.

Now we proceed by induction on the number of vertices to show that every tree has a displayed 2-representation. If $T$ is the one point tree, the result is trivial. Next assume that for some $k \geq 1$, every tree on $k$ vertices has a displayed 2-representation and let $T$ be a tree with $k+1$ vertices.

Choose an endvertex $u$ of $T$ and let $f$ be a displayed 2-representation of the tree $T-u$. Let $v$ be the unique vertex adjacent to $u$ in $T$ and let $I_{v}$ be an open interval contained in $f(v)$ so that $I_{v} \cap f(w)=\varnothing$ for every vertex $w$ in $T-u$ with $w \neq v$. Choose a closed interval A contained in $I_{v}$.

[^0]Now choose an upper bound $m$ for $f$ and define $g(w)=f(w)$ for every vertex $w$ in $T-u$ and $g(u)=A \cup[m, m+1]$. It is clear that $g$ is a displayed 2-representation of $T$ and our proof is complete.

## 2. COMPLETE BIPARTITE GRAPHS

We now derive our main result. We use the notation $\lceil x\rceil$ to represent the smallest integer among those which are at least as large as $x$.

Theorem 2. The interval number of the complete bipartite graph $K_{m, n}$ is given by

$$
i\left(K_{m, n}\right)=\lceil(m n+1) /(m+n)\rceil .
$$

Proof. We first show that $i\left(K_{m, n}\right) \geq\lceil(m n+1) /(m+n)\rceil$. Suppose that $f$ is a $t$-representation of $K_{m, n}$. Without loss of generality, we may assume that for each vertex $u$ in $K_{m, n}, f(u)$ is the union $A_{1}(u) \cup A_{2}(u) \cup \cdots \cup$ $A_{t}(u)$ of $t$ pairwise disjoint closed intervals.

We now use $f$ to determine a graph $G$. The vertices of $G$ are the ordered pairs of the form ( $u, i$ ) where $u$ is a vertex in $K_{m, n}$ and $1 \leq i \leq t$ with distinct vertices $(u, i)$ and $(v, j)$ adjacent in $G$ when $A_{i}(u) \cap$ $A_{i}(v) \neq \varnothing$. The function $g$ defined by $g(u, i)=A_{i}(u)$ is a 1-representation of $G$ so $G$ is an interval graph. Since $G$ is bipartite, it is triangle-free. Since $G$ is an interval graph, it does not contain a cycle of four or more vertices as an induced subgraph. Therefore, $G$ is a forest. Note that $G$ has $(m+n) t$ vertices and at most $(m+n) t-1$ edges.

Now suppose that $e=\{u, v$,$\} is an edge of K_{m, n}$. Then there exist integers $i, j$ with $A_{i}(u) \cap A_{j}(v) \neq \varnothing$, and we may therefore define a function $h$ from the edge set of $K_{m, n}$ to the edge set of $G$ by setting $h(e)=h(\{u, v\})=\{(u, i),(v, j)\}$. Clearly, $h$ is a one-to-one function and since $K_{m, n}$ has $m n$ edges, we see that $m n \leq(m+n) t-1$, i.e., $t \geq$ $\lceil(m n+1) /(m+n)\rceil$.

We will now show that $i\left(K_{m, n}\right) \leq\lceil(m n+1)(m+n)\rceil$. Let $t=$ $[(m n+1) /(m+n)\rceil$. We will construct an interval graph $G$ with a 1representation $g$. We will then construct a t-representation $f$ of $K_{m, n}$ by appropriately choosing, for each vertex $u$ of $K_{m, n}, t$ intervals from the range of $g$ as the intervals whose union is $f(u)$.

We begin by labeling the vertices of $K_{m, n}$ with the symbols $a_{1}, a_{2}, \ldots, a_{m}, b_{1}, b_{2}, \ldots, b_{n}$ so that $a_{i}$ is adjacent to $b_{j}$ for all $i$ and $j$. Without loss of generality, we may assume $m \geq n$. Let $A=$ $\{1,2,3, \ldots, m\}$ and $B=\{1,2,3, \ldots, n\}$.

We next construct a graph $T$ whose vertex set is

$$
\left\{u_{k}: 1 \leq k \leq n t\right\} \cup\left\{v_{k}: 1 \leq k \leq n t-1\right\} \cup\left\{w_{i j}: 1 \leq i \leq m, 1 \leq j \leq n\right\}
$$

where $T$ has the following adjacencies: $v_{k}$ is adjacent to $u_{k}$ and $u_{k+1}$ for $k=1,2, \ldots, n t-1$ and $w_{i j}$ is adjacent to $u_{\mathrm{i}}$ for $i=1,2, \ldots, m$ and $j=1,2, \ldots, n$. The graph $T$ is a caterpillar and, by Theorem 1 , is also an interval graph. Consequently any induced subgraph of $T$ is also an interval graph.

The next step in the construction is to color some, but not all, of the vertices of $T$ using the elements of $A$ as colors. We begin by assigning to $u_{1}, u_{2}, \ldots, u_{n t}$ the colors

$$
1,2,3, \ldots, n, 1,2,3, \ldots, n, \ldots, 1,2,3, \ldots, n
$$

in order. Note that each color from $B$ is used exactly $t$ times.
Now let $s=n-t$; then $2 s \leq n-1$. Suppose that $S$ is a set of either $2 s$ or $2 s-1$ consecutive vertices from the sequence $v_{1}, v_{2}, \ldots, v_{n t-1}$. Consider a subset $S^{\prime}$ of $S$ that contains $s$ vertices, no two of which are consecutive. Then let $B^{\prime}$ be the subset of $B$ consisting of those integers $j$ for which there is a vertex $v$ from $S^{\prime}$ and a vertex $u$ adjacent to $v$ with $u$ having color $j$. It is easy to verify that $B^{\prime}$ must contain $2 s$ elements, i.e., the $s$ vertices of $S^{\prime}$ are adjacent to $2 s$ distinctly colored vertices.

The next step is to assign colors to the first $m s$ vertices in the sequence $v_{1}, v_{2}, \ldots, v_{n t-1}$. Note that $t=\lceil(m n+1) /(m+n)\rceil$ and $s=n-t$ imply that $m s \leq n t-1$. At this point, we must consider two cases depending on the parity of $m$. If $m$ is even, then assign the vertices $v_{1}, v_{2}, \ldots, v_{m s}$ the colors

$$
1,2,1,2, \ldots, 1,2,3,4,3,4, \ldots, 3,4, \ldots, m-1,
$$

in order. Note that each color in $A$ is to be used exactly $s$ times. If $m$ is odd, we modify this scheme as follows. We first assign color $m$ to $v_{1}, v_{n+3}, v_{2 n+5}, \ldots, v_{(s-1)(n+2)+1}$. Note that for each $j=1,2,3, \ldots, 2 s$, there are integers $k, l$ for which $u_{k}$ is adjacent to $v_{l}$, where $v_{l}$ has color $m$ and $u_{k}$ has color $j$. Next assign to the ( $m-1$ )s vertices in the sequence $v_{1}, v_{2}, \ldots, v_{m s}$, which were not assigned color $m$, the colors

$$
1,2,1,2, \ldots, 1,2,3,4,3,4, \ldots, 3,4, \ldots, m-2, m-1, \ldots, m-2, m-1
$$

in order. Again we note that each color in $A$ is to be used exactly $s$ times.
When $m$ is even, observe that each color $i$ from $A$ is assigned to $s$ nonconsecutive vertices in a block of $2 s-1$ consecutive vertices from the sequence $v_{1}, v_{2}, \ldots, v_{n t-1}$. When $m$ is odd, we observe that distinct vertices that have been assigned color $m$ are at least $n+2$ apart in the
sequence $v_{1}, v_{2}, \ldots, v_{n t-1}$. Therefore, we observe that each color $i$ from A with $i \neq m$ is assigned to $s$ nonconsecutive vertices in a block of $2 s$ or $2 s-1$ consecutive vertices in the sequence $v_{1}, v_{2}, \ldots, v_{n t-1}$. For each color $i \in A$, define the set
$B(i)=\left\{j \in B\right.$ : There exist integers $k, l$ with $u_{k}$ adjacent to $v_{1}$ for which $u_{k}$ has been assigned color $\dot{j}$ and $v_{l}$ has been assigned color $\left.i\right\}$.
We conclude that for all values of $m$ and for every color $i$ from $A$, the set $B(i)$ contains exactly $2 s$ elements.

The next step in the construction is to assign colors to some, but not all, of the vertices in $\left\{w_{i j}: 1 \leq i \leq m, 1 \leq j \leq n\right\}$. The construction is the same for all values of $m$. Let $i$ be an element of $A$; assign color $i$ to vertex $w_{i j}$ if and only if $j$ is an element of $\boldsymbol{B}-\boldsymbol{B}(i)$. Now let

$$
U_{1}=\left\{v_{k}: 1 \leq k \leq n t-1\right\} \cup\left\{w_{i j}: 1 \leq i \leq m, 1 \leq j \leq n\right\}
$$

and let

$$
U_{2}=\left\{u_{k}: 1 \leq k \leq n t\right\} .
$$

Observe that for each color $i$ from $A$, exactly $t$ vertices of $U_{1}$ have been assigned color $i$, and for each color $j$ from $B$, exactly $t$ vertices from $U_{2}$ have been assigned color $j$; furthermore, there exist adjacent vertices $u^{\prime}, u^{\prime \prime}$ with $u^{\prime}$ from $U_{1}, u^{\prime \prime}$ from $U_{2}, u^{\prime}$ having color $i$, and $u^{\prime \prime}$ having color $j$.

Now let $G$ be the subgraph of $T$ generated by the colored vertices and let $g$ be a 1 -representation of $G$. The final step in the construction is to use $g$ to define a $t$-representation $f$ of $K_{m, n}$. But this is accomplished simply by defining

$$
f\left(a_{i}\right)=\cup\left\{g\left(u^{\prime}\right): u^{\prime} \text { is a vertex from } U_{1} \text { and } u^{\prime} \text { has color } i\right\}
$$

$$
\text { for } i=1,2, \ldots, m
$$

and

$$
\begin{array}{r}
f\left(b_{j}\right)=\cup\left\{g\left(u^{\prime \prime}\right): u^{\prime \prime} \text { is a vertex from } U_{2} \text { and } u^{\prime \prime} \text { has color } j\right\} \\
\text { for } j=1,2, \ldots, n .
\end{array}
$$

It is trivial to verify that $f$ is a $t$-representation of $K_{m, n}$.

## 3. OTHER RESULTS

A preliminary version of this paper included a proof of the following result.

Theorem 3. If $G$ has $p$ vertices, then $i(G) \leq\lceil p / 3\rceil$.
This theorem may be established using a two-part argument in which it is proved inductively that a graph on $3 n$ vertices has an $n$-representation and a triangle-free graph on $3 n$ vertices has a displayed $n$-representation. The proof of the second part makes use of Turán's theorem for the maximum number of edges in a triangle-free graph.

However, the authors did not believe that the upper bound on the interval number of a graph provided by Theorem 3 was best possible. Motivated by the observation that the complete bipartite graph $K_{2 n, 2 n}$ has $4 n$ vertices and interval number $n+1$, the authors conjectured that if $G$ is a graph with $p$ vertices, then $i(G) \leq\lceil(p+1) / 4\rceil$.

The concept of interval number has been independently investigated by Griggs and West [4]. They obtained the formula given in Theorem 1 for the interval number of a tree as well as the upper bound given in Theorem 3. They also made the same conjecture concerning the maximum interval number of a graph with $p$ vertices. And they also provided an upper bound on the interval number of a graph in terms of the maximum degree of a vertex in the graph. Specifically, they showed that if the maximum degree of a vertex in a graph $G$ is $d$, then $i(G) \leq$ $\lceil(d+1) / 2\rceil$. This last result allowed them to determine that the interval number of the $n$-cube $Q_{n}$ is $\lceil(n+1) / 2\rceil$, which answered a problem posed in the preliminary version of this paper.

The authors have recently learned that Griggs [3] has established the conjecture by proving that if $G$ has $4 n-1$ vertices, then $i(G) \leq n$.

## 4. AN OPEN PROBLEM

Lekkerkerker and Boland [7] gave a forbidden subgraph characterization of interval graphs by listing the collection $\mathscr{\Phi}_{2}$ of graphs defined by

$$
\mathscr{I}_{2}=\{G: i(G)=2 \text { but } i(H)=1
$$

for every proper induced subgraph $H$ of $G\}$.
We propose the general problem of finding for $t \geq 3$, the collection

$$
\mathscr{I}_{t}=\{G: i(G)=t \text { but } i(H) \leq t-1 \text { for every proper subgraph } H \text { of } G\}
$$

The problem for $t=3$ seems to both manageable and interesting since from applied viewpoint, graphs that are the intersection graphs of a family of sets each of which is the union of two intervals of the real line have practical significance, e.g., two work periods separated by a lunch break. By double interval graphs, we mean graphs with interval number
two. Theorem 2 shows that $K_{2 n, 2 n}$ is in $\mathscr{\Phi}_{n+1}$ for every $n \geq 1$ and that $K_{2 n-1,2 n+2}$ is in $\mathscr{I}_{n+1}$ for every $n \geq 2$. In particular, we note then that a forbidden subgraph characterization of double interval graphs will include $K_{4,4}$ and $K_{3,6}$.

## References

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[^0]:    * Roberts [8] has studied another generalization of interval graphs. He defines the boxicity of a graph $G$ as the smallest positive integer $t$ for which there exists a function $f$ which assigns to each vertex $u$ of $G$ a sequence $f(u)(1), f(u)$ 2), $\ldots, f(u)(t)$ of closed intervals of $R$ so that distinct vertices $u, v$ of $G$ are adjacent if and only if $f(u)(i) \cap f(v)(i) \neq \varnothing$ for $i=1,2,3, \ldots, t$.

