## ON *f*-DERIVATIONS OF LATTICES

YILMAZ ÇEVEN AND MEHMET ALI ÖZTÜRK

ABSTRACT. In this paper, as a generalization of derivation on a lattice, the notion of f-derivation for a lattice is introduced and some related properties are investigated.

### 1. Introduction and preliminaries

Lattices play an important role in many fields such as information theory, information retrieval, information access controls and cryptanalysis [2, 6, 13, 8]. Recently the properties of lattices were widely researched [1, 2, 5, 7, 8, 9, 13, 14].

In the theory of rings and near rings, the properties of derivations are an important topic to study [3, 4, 11, 12]. Y. B. Jun and X. L. Xin [10] applied the notion of derivation in ring and near ring theory to BCI-algebras. In [15], J. Zhan and Y. L. Liu introduced the notion of left-right (or right-left) *f*-derivation of a BCI algebra and investigated some properties. In [14], X. L. Xin, T. Y. Li, and J. H. Lu introduced the notion of derivation on a lattice and discussed some related properties.

In this paper, as a generalization of derivation on a lattice, the notion of fderivation of a lattice is introduced and some related properties, which are discussed in [14] for a derivation on a lattice, are investigated for the f-derivation on a lattice. As important results of f-derivations on lattices, distributive and modular lattices are characterized by f-derivations under some conditions.

**Definition 1** ([5]). Let *L* be a nonempty set endowed with operations  $\land$  and  $\lor$ . Then  $(L, \land, \lor)$  is called a lattice if it satisfies the following conditions for all  $x, y, z \in L$ :

(1)  $x \wedge x = x, x \vee x = x$ ,

(2)  $x \wedge y = y \wedge x, x \vee y = y \vee x,$ 

(3)  $(x \wedge y) \wedge z = x \wedge (y \wedge z), (x \vee y) \vee z = x \vee (y \vee z),$ 

(4)  $(x \wedge y) \lor x = x, (x \lor y) \land x = x.$ 

**Definition 2** ([5]). A lattice L is called a distributive lattice if it satisfies the identity (5) or (6) for all  $x, y, z \in L$ :

(5)  $x \wedge (y \vee z) = (x \wedge y) \vee (x \wedge z)$ ,

Received January 30, 2008.

2000 Mathematics Subject Classification. 06B35, 06B99. Key words and phrases. lattice, derivation, f-derivation.

words and phrases. Interee, derivation, j derivation.

©2008 The Korean Mathematical Society



(6)  $x \lor (y \land z) = (x \lor y) \land (x \lor z)$ .

**Definition 3** ([1]). A lattice L is called a modular lattice if it satisfies the following condition for all  $x, y, z \in L$ :

(7) If  $x \le z$ , then  $x \lor (y \land z) = (x \lor y) \land z$ .

**Definition 4** ([5]). Let  $(L, \wedge, \vee)$  be a lattice. A binary relation  $\leq$  is defined by  $x \leq y$  if and only if  $x \wedge y = x$  and  $x \vee y = y$ .

**Definition 5** ([5]). Let L and M be two lattices. The function  $g: L \longrightarrow M$  is called a lattice homomorphism if it satisfies the following conditions for all  $x, y \in L$ :

(8) 
$$g(x \wedge y) = g(x) \wedge g(y),$$

(9) 
$$g(x \lor y) = g(x) \lor g(y).$$

It is known that a homomorphism is called an epimorphism if it is onto.

**Lemma 1** ([14]). Let  $(L, \wedge, \vee)$  be a lattice. Define the binary relation  $\leq$  as the Definition 4. Then  $(L, \leq)$  is a poset and for any  $x, y \in L$ ,  $x \wedge y$  is the g.l.b of  $\{x, y\}$  and  $x \vee y$  is the l.u.b. of  $\{x, y\}$ .

**Definition 6** ([14]). A function  $d: L \longrightarrow L$  on a lattice L is called a derivation on L if it satisfies the following condition

$$d(x \wedge y) = (dx \wedge y) \vee (x \wedge dy).$$

The abbreviation dx is used for d(x) in the above definition.

**Definition 7** ([14]). Let L be a lattice and d be a derivation on L.

- (1) If  $x \leq y$  implies  $dx \leq dy, d$  is called an isotone derivation,
- (2) If d is one-to-one, d is called a monomorphic derivation,
- (3) If d is onto, d is called an epimorphic derivation.

# 2. The *f*-derivations in lattices

The following definition introduces the notion of f-derivations on lattices.

**Definition 8.** Let L be a lattice. A function  $d : L \longrightarrow L$  is called an fderivation on L if there exists a function  $f : L \longrightarrow L$  such that

(2.1) 
$$d(x \wedge y) = (d(x) \wedge f(y)) \vee (f(x) \wedge d(y))$$

for all  $x, y \in L$ .

It is obvious in the Definition 8 that if f is an identity function then d is a derivation on L. Furthermore, according to Definition 8, a function d on Lcan be an f-derivation only when a function f satisfying equation (2.1) exists. But, to obtain some results, d or f must satisfy some additional conditions as in the following propositions and theorems.

In this paper, we'll abbreviate d(x) as dx and f(x) as fx.

**Example 1.** Let *L* be the lattice of Figure 1 and define a function *d* by d0 = 0, da = a, db = a, dc = c, d1 = c.

702



#### Figure 1

Then d is not a derivation on L since  $a = d(b \wedge 1) \neq (db \wedge 1) \vee (b \wedge d1) = (a \wedge 1) \vee (b \wedge c) = a \vee b = b$ . If we define a function f by f0 = 0, fa = a, fb = a, fc = 1, f1 = 1, then d satisfies the equation (2.1) for all  $x, y \in L$  and so d is an f-derivation on L.

**Example 2.** Let *L* be a lattice and  $a \in L$ . Define a function  $d: L \longrightarrow L$  by  $dx = fx \wedge a$  for all  $x \in L$  where  $f: L \longrightarrow L$  satisfies  $f(x \wedge y) = fx \wedge fy$  for all  $x, y \in L$ . Then *d* is an *f*-derivation. In addition, if *f* is an increasing function then *d* is an isotone derivation.

**Proposition 1.** Let L be a lattice and d be an f-derivation on L. Then the following identities hold for all  $x, y \in L$ .

- a)  $dx \leq fx$ ,
- b)  $dx \wedge dy \leq d(x \wedge y) \leq dx \vee dy$ ,
- c)  $d(x \wedge y) \leq fx \vee fy$ ,

d) If L has a least element 0, then f0 = 0 implies d0 = 0.

*Proof.* a) Since  $dx = d(x \wedge x) = dx \wedge fx$ , we have  $dx \leq fx$ .

b) We have  $dx \wedge fy \leq d(x \wedge y)$  and  $fx \wedge dy \leq d(x \wedge y)$  from the equation (2.1). Since  $dx \leq fx$ , we obtain  $dx \wedge dy \leq fx \wedge dy$  and hence we have  $dx \wedge dy \leq d(x \wedge y)$ . We know that  $dx \wedge fy \leq dx$  and  $fx \wedge dy \leq dy$ . Then we obtain  $d(x \wedge y) = (dx \wedge fy) \vee (fx \wedge dy) \leq dx \vee dy$ .

c) Since  $dx \wedge fy \leq fy$  and  $fx \wedge dy \leq fx$ , we obtain  $d(x \wedge y) \leq fx \vee fy$ .

d) Since  $dx \leq fx$  for all  $x \in L$ , f0 = 0 and 0 is the least element of L, we have  $0 \leq d0 \leq f0 = 0$ .

**Proposition 2.** Let L be a lattice and d be an f-derivation and 1 be the greatest element of L and f1 = 1. Then the following identities hold;

a) If  $fx \leq d1$ , then dx = fx,

b) If  $fx \ge d1$ , then  $dx \ge d1$ .

*Proof.* a) Since  $dx = d(x \wedge 1) = (dx \wedge f1) \vee (fx \wedge d1) = dx \vee fx$ , we have  $fx \leq dx$ . From Proposition 1 a), we obtain dx = fx.

b) Since  $dx = (dx \wedge f1) \lor (fx \wedge d1) = dx \lor d1$ , we have  $dx \ge d1$ .

Remark 1. Note that if d1 = 1, since  $d1 \le f1$ , we have f1 = 1 where 1 the greatest element of L. In this case, from Proposition 2 a), we get d = f.

Let L be a lattice and d be an f-derivation of L. Define a set  $F = \{x \in L : dx = fx\}$ .

**Proposition 3.** Let L be a lattice and d be an f-derivation. If f is an increasing function, then  $y \leq x$  and  $x \in F$  implies  $y \in F$ .

*Proof.* Note that, since  $x \in F$  and  $dy \leq fy \leq fx = dx$ , we get  $dy = d(x \wedge y) = (dx \wedge fy) \lor (fx \wedge dy) = (fx \wedge fy) \lor dy = fy \lor dy = fy$ .  $\Box$ 

**Proposition 4.** Let L be a lattice and d be an isotone f-derivation on L. Then for any  $x, y \in L$ ,  $dx = dx \lor (fx \land d(x \lor y))$ .

*Proof.* Since d is an isotone f-derivation, we know that for all  $x, y \in L$ ,  $dx \leq d(x \lor y) \leq f(x \lor y)$ . Hence we have  $dx = d((x \lor y) \land x) = (d(x \lor y) \land fx) \lor (f(x \lor y) \land dx) = dx \lor (fx \land d(x \lor y))$ .

**Proposition 5.** Let L be a lattice and d be an isotone f-derivation. If  $x, y \in F$  and f is a decreasing function, then  $x \lor y \in F$ .

*Proof.* Since  $x \leq x \lor y$  and  $y \leq x \lor y$ , we have  $f(x \lor y) \leq fx$  and  $f(x \lor y) \leq fy$  respectively. Then we obtain  $f(x \lor y) \leq fx \lor fy = dx \lor dy \leq d(x \lor y)$  since d is an isotone f-derivation. It is known that  $d(x \lor y) \leq f(x \lor y)$ , hence we get  $x \lor y \in F$ .  $\Box$ 

**Theorem 1.** Let L be a lattice with greatest element 1 and d be an f-derivation on L. Let f1 = 1 and  $f(x \land y) = fx \land fy$  for all  $x, y \in L$ . Then the following conditions are equivalent:

(1) d is an isotone f-derivation,

(2)  $dx \lor dy \le d(x \lor y)$ ,

(3)  $dx = fx \wedge d1$ ,

 $(4) \ d(x \wedge y) = dx \wedge dy.$ 

*Proof.* (1)  $\implies$  (2) : Suppose that d is an isotone f-derivation. We know that  $x \leq x \lor y$  and  $y \leq x \lor y$ . Since d is isotone,  $dx \leq d(x \lor y)$  and  $dy \leq d(x \lor y)$ . Hence we obtain  $dx \lor dy \leq d(x \lor y)$ .

 $(2) \Longrightarrow (1)$ : Suppose that  $dx \lor dy \le d(x \lor y)$  and  $x \le y$ . Then we have  $dx \le dx \lor dy \le d(x \lor y) = dy$ .

 $(1) \Longrightarrow (3)$ : Suppose that d is an isotone f-derivation. We have  $dx \le d1$ . It is known that  $dx \le fx$  from Proposition 1 a). Then we get  $dx \le fx \land d1$ . From Proposition 4, for y = 1, we have  $dx = dx \lor (fx \land d1) = fx \land d1$ .

(3)  $\implies$  (4) : Assume that (3) holds. Then  $d(x \wedge y) = f(x \wedge y) \wedge d1 = fx \wedge fy \wedge d1 = (fx \wedge d1) \wedge (fy \wedge d1) = dx \wedge dy$ .

 $(4) \Longrightarrow (1)$ : Let  $d(x \wedge y) = dx \wedge dy$  and  $x \leq y$ . Since  $dx = d(x \wedge y) = dx \wedge dy$ , we get  $dx \leq dy$ .

**Theorem 2.** Let L be a modular lattice and d be an f-derivation on L.

(a) d is an isotone f-derivation if and only if  $d(x \wedge y) = dx \wedge dy$ ,

(b) If d is an isotone f-derivation where  $f(x \lor y) = fx \lor fy$ , then dx = fx implies  $d(x \lor y) = dx \lor dy$ .

*Proof.* (a) Suppose that d is an isotone f-derivation. Since  $x \wedge y \leq x$  and  $x \wedge y \leq y$ , we get  $d(x \wedge y) \leq dx$  and  $d(x \wedge y) \leq dy$ . Hence  $d(x \wedge y) \leq dx \wedge dy$ . Also using Proposition 1 a) and the fact  $dx \wedge fy \leq dx \leq fx$ , we get

$$dx \wedge dy = (dx \wedge dy) \wedge (fx \wedge fy)$$
  

$$\leq (dx \vee dy) \wedge (fx \wedge fy)$$
  

$$= ((dy \vee dx) \wedge fy) \wedge fx$$
  

$$= (dy \vee (dx \wedge fy)) \wedge fx$$
  

$$= (dx \wedge fy) \vee dy) \wedge fx$$
  

$$= (dx \wedge fy) \vee (fx \wedge dy)$$
  

$$= d(x \wedge y).$$

Conversely, let  $d(x \wedge y) = dx \wedge dy$  and  $x \leq y$ . Since  $dx = d(x \wedge y) = dx \wedge dy$ , we have  $dx \leq dy$ .

(b) Suppose that d is an isotone f-derivation and dx=fx . Using Proposition 4 and since L is modular lattice, we have

$$dy = dy \lor (fy \land d(x \lor y))$$
  
=  $(dy \lor fy) \land d(x \lor y)$   
=  $fy \land d(x \lor y).$ 

Hence using the hypothesis, we obtain

$$dx \lor dy = dx \lor (fy \land d(x \lor y))$$
  
=  $(dx \lor fy) \land d(x \lor y)$   
=  $(fx \lor fy) \land d(x \lor y)$   
=  $f(x \lor y) \land d(x \lor y)$   
=  $d(x \lor y).$ 

**Theorem 3.** Let L be a distributive lattice and d be an f-derivation on L where  $f(x \lor y) = fx \lor fy$ . Then the following hold:

(1) d is an isotone f-derivation implies  $d(x \wedge y) = dx \wedge dy$ ,

(2) d is an isotone f-derivation if and only if  $d(x \lor y) = dx \lor dy$ .

*Proof.* (1) Since d is isotone f-derivation, we know that  $d(x \wedge y) \leq dx \wedge dy$ . From Proposition 1 a), we get

$$dx \wedge dy = (dx \wedge fx) \wedge (dy \wedge fy)$$
$$= (dx \wedge fy) \wedge (fx \wedge dy)$$

$$\leq (dx \wedge fy) \lor (fx \wedge dy) = d(x \wedge y).$$

Hence we have  $d(x \wedge y) = dx \wedge dy$ .

(2) Let d is an isotone f-derivation. From (1), we have  $d(x \wedge y) = dx \wedge dy$ . Then, from Proposition 1 a) and Proposition 4, we get  $dy = dy \vee (fy \wedge d(x \vee y)) = (dy \vee fy) \wedge (dy \vee d(x \vee y)) = fy \wedge d(x \vee y)$  and similarly  $dx = fx \wedge d(x \vee y)$ . Then we obtain

$$dx \lor dy = (fx \land d(x \lor y)) \lor (fy \land d(x \lor y))$$
  
=  $(fx \lor fy) \land d(x \lor y)$   
=  $f(x \lor y) \land d(x \lor y)$   
=  $d(x \lor y).$ 

Conversely, suppose that  $d(x \lor y) = dx \lor dy$  and  $x \le y$ . Then since  $dy = d(x \lor y) = dx \lor dy$ , we have  $dx \le dy$ .

**Theorem 4.** Let L be a lattice. If there exists an f-derivation d on L such that  $d(x \lor y) = dx \lor dy$  for all  $x, y \in L$  and f is an epimorphism, then L is a distributive lattice.

*Proof.* We know from Example 2 that the function d defined by  $dx = fx \wedge c$  for  $c \in L$  where f is a homomorphism is an f-derivation on L. Also suppose that f is onto and  $d(x \vee y) = dx \vee dy$  for all  $x, y \in L$ . Then, for all  $a, b \in L$  there exist  $u, v \in L$  such that fu = a and fv = b. Hence

$$(a \lor b) \land c = (fu \lor fv) \land c$$
  
=  $f(u \lor v) \land c$   
=  $d(u \lor v)$   
=  $du \lor dv$   
=  $(fu \land c) \lor (fv \land c)$   
=  $(a \land c) \lor (b \land c).$ 

Since every distributive lattice is a modular lattice, we have the following corollary.

**Corollary 1.** Let L be a lattice. If there exists an f-derivation d on L such that  $d(x \lor y) = dx \lor dy$  for all  $x, y \in L$  and f is an epimorphism, then L is a modular lattice.

Acknowledgements. The authors are highly grateful to the referees for their valuable comments and suggestions for the paper.

706

### References

- R. Balbes and P. Dwinger, *Distributive Lattices*, University of Missouri Press, Columbia, Mo., 1974.
- [2] A. J. Bell, The co-information lattice, 4<sup>th</sup> Int. Symposium on Independent Component Analysis and Blind Signal Separation (ICA2003), Nara, Japan, 2003, 921–926.
- [3] H. E. Bell and L. C. Kappe, Rings in which derivations satisfy certain algebraic conditions, Acta Math. Hungar. 53 (1989), no. 3-4, 339–346.
- [4] H. E. Bell and G. Mason, On derivations in near-rings, Near-rings and near-fields (Tubingen, 1985), 31–35, North-Holland Math. Stud., 137, North-Holland, Amsterdam, 1987.
- [5] G. Birkhoof, *Lattice Theory*, American Mathematical Society, New York, 1940.
- [6] C. Carpineto and G. Romano, Information retrieval through hybrid navigation of lattice representations, International Journal of Human-Computers Studies 45 (1996), 553–578.
- [7] C. Degang, Z. Wenxiu, D. Yeung, and E. C. C. Tsang, Rough approximations on a complete completely distributive lattice with applications to generalized rough sets, Inform. Sci. 176 (2006), no. 13, 1829–1848.
- [8] G. Durfee, Cryptanalysis of RSA using algebraic and lattice methods, A dissertation submitted to the department of computer sciences and the committe on graduate studies of Stanford University (2002), 1–114.
- [9] A. Honda and M. Grabisch, Entropy of capacities on lattices and set systems, Inform. Sci. 176 (2006), no. 23, 3472–3489.
- [10] Y. B. Jun and X. L. Xin, On derivations of BCI-algebras, Inform. Sci. 159 (2004), no. 3-4, 167–176.
- [11] K. Kaya, Prime rings with  $\alpha$ -derivations, Hacettepe Bull. Natural. Sci. and Eng. 16-17 (1987/1988), 63–71.
- [12] E. Posner, Derivations in prime rings, Proc. Amer. Math. Soc. 8 (1957), 1093–1100.
- [13] R. S. Sandhu, Role hierarchies and constraints for lattice-based access controls, Proceedings of the 4<sup>th</sup> Europan Symposium on Research in Computer Security, Rome, Italy, 1996, 65–79.
- [14] X. L. Xin, T. Y. Li, and J. H. Lu, On derivations of lattices, Inform. Sci. 178 (2008), no. 2, 307–316.
- [15] J. Zhan and Y. L. Liu, On f-derivations of BCI-algebras, Int. J. Math. Math. Sci. (2005), no. 11, 1675–1684.

YILMAZ ÇEVEN DEPARTMENT OF MATHEMATICS FACULTY OF ARTS AND SCIENCES SÜLEYMAN DEMIREL UNIVERSITY 32260 ISPARTA, TURKEY *E-mail address*: yceven@fef.sdu.edu.tr

MEHMET ALI ÖZTÜRK DEPARTMENT OF MATHEMATICS FACULTY OF ARTS AND SCIENCES ADIYAMAN UNIVERSITY 02040 ADIYAMAN, TURKEY *E-mail address*: mehaliozturk@gmail.com