

# On the strength of connectedness of a random hypergraph

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## Abstract

Bollobás and Thomason (1985) proved that for each  $k = k(n) \in [1, n - 1]$ , with high probability, the random graph process, where edges are added to vertex set  $V = [n]$  uniformly at random one after another, is such that the stopping time of having minimal degree  $k$  is equal to the stopping time of becoming  $k$ -(vertex-)connected. We extend this result to the  $d$ -uniform random hypergraph process, where  $k$  and  $d$  are fixed. Consequently, for  $m = \frac{n}{d}(\ln n + (k - 1) \ln \ln n + c)$  and  $p = (d - 1)! \frac{\ln n + (k - 1) \ln \ln n + c}{n^{d-1}}$ , the probability that the random hypergraph models  $H_d(n, m)$  and  $H_d(n, p)$  are  $k$ -connected tends to  $e^{-e^{-c}/(k-1)!}$ .

**Keywords:** random hypergraph; vertex connectivity

## 1 Introduction

Let  $H_d(n, p)$  denote the random  $d$ -uniform hypergraph with vertex set  $[n] := \{1, 2, \dots, n\}$ , where each of the  $\binom{n}{d}$  potential (hyper)edges of cardinality  $d$  is present with probability  $p$ , independently of all other potential edges. Likewise, let  $H_d(n, m)$  be the random  $d$ -uniform hypergraph on  $[n]$ , where  $m$  edges are chosen uniformly at random among all sets of  $m$  potential edges. The model  $H_d(n, m)$  can be gainfully viewed as a snapshot of the random hypergraph process  $\{H_d(n, \mu)\}_{\mu=0}^{\binom{n}{d}}$ , where  $H_d(n, \mu + 1)$  is obtained from  $H_d(n, \mu)$  by inserting an extra edge chosen uniformly at random among all  $\binom{n}{d} - \mu$  remaining potential edges. For  $d = 2$ , these models are the typical random graph models,  $G(n, p)$ ,  $G(n, m)$  and  $\{G(n, \mu)\}_{\mu=0}^{\binom{n}{2}}$ .

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As customary, we say that for a given  $m = m(n)$  ( $p$  resp.) some graph property  $\mathcal{Q}$  holds *with high probability*, denoted *w.h.p.*, if the probability that  $H_d(n, m)$  ( $(H_d(n, p)$  resp.) has property  $\mathcal{Q}$  tends to 1 as  $n \rightarrow \infty$ . Further,  $m(n)$  is the *sharp threshold* for  $\mathcal{Q}$  if for each  $\epsilon > 0$  (fixed), w.h.p.  $H_d(n, (1 - \epsilon)m)$  does not have  $\mathcal{Q}$  and w.h.p.  $H_d(n, (1 + \epsilon)m)$  does have  $\mathcal{Q}$ . For the random hypergraph process, the *stopping time* of  $\mathcal{Q}$ , denoted  $\tau(\mathcal{Q})$  is the first moment that the process has this property  $\mathcal{Q}$ ; we denote the hypergraph process stopped at this time by  $H_d(n, \tau(\mathcal{Q}))$ .

In one of the first papers on random graphs, Erdős and Rényi [4] showed that  $m = \frac{1}{2}n \ln n$  is a sharp threshold for connectivity in  $G(n, m)$ . Later, Stepanov [7] established the sharp threshold of connectivity for  $G(n, p)$  among other results. More recently, Bollobás and Thomason [3] proved the stronger result for the random graph process that w.h.p. the moment the graph process loses its last isolated vertex is also the moment that the process becomes connected; in other words, w.h.p.  $\tau(\text{no isolated vertices}) = \tau(\text{connected})$ . We prove the analogous result for the the random  $d$ -uniform hypergraph process; a consequence of this result is that  $m = \frac{n}{d} \ln n$  is a sharp threshold of connectivity for  $H_d(n, m)$ .

There are various measures for the *strength* of connectedness of a connected graph, but here we will focus on  $k$ -(vertex-)connectivity. For  $k \in \mathbb{N}$ , a hypergraph with more than  $k$  vertices is  $k$ -connected if whenever  $k - 1$  vertices are deleted, along with their incident edges, the remaining hypergraph is connected. Note that the definition of 1-connectedness coincides with connectedness. Necessarily, for a hypergraph to be  $k$ -connected, each vertex must have degree at least  $k$ , because if a vertex  $v$  has degree less than  $k$ , then we can delete a neighbor from each incident edge to isolate  $v$ . However, as commonly seen in these types of results, the main barrier to  $k$ -connectivity in these random graph models arises from such vertices that can be separated from the rest of the graph by the deletion of their neighbors (see for instance Erdős-Rényi [5], Ivchenko [6], Bollobás [1],[2]). Here, we extend this idea to random  $d$ -uniform hypergraphs; in particular, we find that if  $m_0 = \frac{n}{d}(\ln n + (k - 1) \ln \ln n - \omega)$  and  $m_1 = \frac{n}{d}(\ln n + (k - 1) \ln \ln n + \omega)$ ,  $\omega \rightarrow \infty$ , then w.h.p.  $H_d(n, m_0)$  is not  $k$ -connected and w.h.p.  $H_d(n, m_1)$  is  $k$ -connected; also we find an analogous threshold value for  $H_d(n, p)$ .

A stronger result concerns the random graph process where edges are added one after another. Let  $\tau_k := \tau(\text{min degree at least } k)$  and  $T_k := \tau(k\text{-connected})$ ; note that  $\tau_k \leq T_k$ . In [3], Bollobás and Thomason showed that for  $d = 2$  (the graph case) and any  $k = k(n) \in [1, n - 1]$ ,  $P(\tau_k = T_k) \rightarrow 1$ . We extend this result for  $d$ -uniform random hypergraphs albeit for *fixed*  $d$  and  $k$ .

**Theorem 1.1.** *W.h.p., at the moment the  $d$ -uniform hypergraph process loses its last vertex with degree less than  $k$ , this process becomes  $k$ -connected. Formally, for  $d \geq 3$  and  $k \geq 1$  (both fixed),  $P(\tau_k = T_k) \rightarrow 1$ .*

To prove this result, we begin by determining the likely range of  $\tau_k$ , and further that just prior to this window, at some  $m_0$  edges, w.h.p. there are not many vertices of degree less than  $k$ . Then, we prove that w.h.p.  $H_d(n, m_0)$  is almost  $k$ -connected in the sense that whenever  $k - 1$  vertices are deleted, there is a massive component using almost all leftover vertices. Third, we show that w.h.p. to isolate a vertex of  $H_d(n, \tau_k)$ , you would have to

delete at least  $k$  of its neighbors (this is trivially true for graphs, but not so for  $d \geq 3$ ). In particular, we show that w.h.p. edges incident to degree  $k - 1$  vertices have trivial intersection (just the vertex itself). Finally, we show that the probability that  $\tau_k < T_k$ , but these three previous likely events also hold tends to zero, which completes the proof of the theorem. The following corollary is nearly immediate in light of the theorem.

**Corollary 1.2.** (i) Let  $m = \frac{n}{d} (\ln n + (k - 1) \ln \ln n + c_n)$ , where  $c_n \rightarrow c \in \mathbb{R}$ . W.h.p.  $H_d(n, m)$  is  $(k - 1)$ -connected, but not  $(k + 1)$ -connected. Further, the probability that  $H_d(n, m)$  is  $k$ -connected tends to  $e^{-e^{-c}/(k-1)!}$ .

(ii) Let  $p = (d - 1)! \frac{\ln n + (k-1) \ln \ln n + c_n}{n^{d-1}}$ , where  $c_n \rightarrow c \in \mathbb{R}$ . W.h.p.  $H_d(n, p)$  is  $(k - 1)$ -connected, but not  $(k + 1)$ -connected. Further, the probability that  $H_d(n, p)$  is  $k$ -connected tends to  $e^{-e^{-c}/(k-1)!}$ .

For the remainder of this paper, let  $d \geq 3$  and  $k \geq 1$  be fixed numbers.

## 2 Likely range of $\tau_k$

**Lemma 2.1.** Let  $\omega = \omega(n) \rightarrow \infty$ , but  $\omega = o(\ln \ln n)$ ,  $m_0 = \frac{n}{d} (\ln n + (k - 1) \ln \ln n - \omega)$  and  $m_1 = \frac{n}{d} (\ln n + (k - 1) \ln \ln n + \omega)$ . Then w.h.p.,

(i) the minimum degree of  $H_d(n, m_0)$  is  $k - 1$  and the number of vertices with degree  $k - 1$  is in the interval

$$\left[ \frac{1}{2} \frac{e^\omega}{(k-1)!}, \frac{3}{2} \frac{e^\omega}{(k-1)!} \right]. \quad (1)$$

(ii) there are no vertices of degree  $k - 1$  in  $H_d(n, m_1)$ .

Consequently, w.h.p.  $\tau_k \in [m_0, m_1]$ .

*Proof.* We prove that the number of vertices with degree  $k - 1$ , denoted by  $X$ , is in the interval (1) by Chebyshev's Inequality. Note that a given vertex can be in  $\binom{n-1}{d-1}$  possible edges, so

$$E[X] = nP[\deg(1) = k - 1] = n \frac{\binom{n-1}{k-1} \binom{\binom{n}{d} - \binom{n-1}{d-1}}{m_0 - k + 1}}{\binom{\binom{n}{d}}{m_0}}.$$

Here and elsewhere in this paper, we use the identity  $\binom{N}{m-\ell} = \binom{N}{m} \frac{(m)_\ell}{(N-m+\ell)_\ell}$ , where  $(j)_\ell = j(j-1)\cdots(j-\ell+1)$ , and later, we use the inequality  $\binom{N}{m-\ell} \leq \binom{N}{m} \left(\frac{m}{N-m}\right)^\ell$ . Now

$$E[X] = (1 + O(1/n)) \frac{n \cdot n^{(d-1)(k-1)}}{(k-1)!((d-1)!)^{k-1}} \left(\frac{d! m_0}{n^d}\right)^{k-1} \frac{\binom{\binom{n}{d} - \binom{n-1}{d-1}}{m_0}}{\binom{\binom{n}{d}}{m_0}}.$$

This latter fraction can be sharply approximated.

$$\frac{\binom{\binom{n}{d} - \binom{n-1}{d-1}}{m_0}}{\binom{\binom{n}{d}}{m_0}} = \prod_{i=0}^{m_0-1} \left(1 - \frac{\binom{n-1}{d-1}}{\binom{\binom{n}{d}}{m_0} - i}\right) = \prod_{i=0}^{m_0-1} \left(1 - \frac{d}{n} + O\left(\frac{i}{n^{d+1}}\right)\right)$$

$$\begin{aligned}
&= \exp \left( \sum_{i=0}^{m_0-1} \left[ \frac{-d}{n} + O\left(\frac{1}{n^2}\right) + O\left(\frac{i}{n^{d+1}}\right) \right] \right) \\
&= \left( 1 + O\left(\frac{\ln n}{n}\right) \right) \exp\left(-\frac{dm_0}{n}\right) = \left( 1 + O\left(\frac{\ln n}{n}\right) \right) \frac{e^\omega}{n(\ln n)^{k-1}}. \quad (2)
\end{aligned}$$

Hence

$$\begin{aligned}
E[X] &= \left( 1 + O\left(\frac{\ln n}{n}\right) \right) \frac{e^\omega}{(k-1)!} \left( \frac{dm_0/n}{\ln n} \right)^{k-1} \\
&= \left( 1 + O\left(\frac{\ln \ln n}{\ln n}\right) \right) \frac{e^\omega}{(k-1)!}.
\end{aligned}$$

For the second factorial moment, we have that

$$E[X(X-1)] = n(n-1)P(\deg(1) = \deg(2) = k-1).$$

We break this latter probability over  $i$ , the number of edges that include both vertices 1 and 2. In particular, vertex 1 is in  $k-1-i$  edges that do not contain vertex 2 and vice versa; further there are  $m_0 - 2(k-1) + i$  edges that include neither vertex 1 or 2. Since there are  $\binom{n-2}{d-2}$  potential hyperedges containing both vertices and  $\binom{n-1}{d-1} - \binom{n-2}{d-2}$  potential hyperedges containing one vertex but not the other, we have that

$$P(\deg(1) = \deg(2) = k-1) = \sum_{i=0}^{k-1} \binom{\binom{n-2}{d-2}}{i} \binom{\binom{n-1}{d-1} - \binom{n-2}{d-2}}{k-1-i}^2 \frac{\binom{\binom{n}{d} - 2\binom{n-1}{d-1} + \binom{n-2}{d-2}}{m_0 - 2(k-1) + i}}{\binom{\binom{n}{d}}{m_0}}. \quad (3)$$

Just as in (2), we can estimate this latter fraction

$$\begin{aligned}
\frac{\binom{\binom{n}{d} - 2\binom{n-1}{d-1} + \binom{n-2}{d-2}}{m_0 - 2(k-1) + i}}{\binom{\binom{n}{d}}{m_0}} &= \frac{(m_0)_{2(k-1)-i}}{\left(\binom{\binom{n}{d} - 2\binom{n-1}{d-1} + \binom{n-2}{d-2}}{m_0 - 2(k-1) + i}\right)_{2(k-1)-i}} \times \frac{\binom{\binom{n}{d} - 2\binom{n-1}{d-1} + \binom{n-2}{d-2}}{m_0}}{\binom{\binom{n}{d}}{m_0}} \\
&= \left( 1 + O\left(\frac{\ln \ln n}{\ln n}\right) \right) \left( \frac{m_0}{\binom{n}{d}} \right)^{2(k-1)-i} \frac{e^{2\omega}}{n^2 (\ln n)^{2(k-1)}}.
\end{aligned}$$

Using these asymptotics, one can show that the  $i$ 'th term in (3) is on the order of  $\frac{e^{2\omega}}{n^{2+i}(\ln n)^i}$ . In particular, the sum of the terms over  $i \in [1, k-1]$ , is  $O(n^{-3})$ . Therefore

$$\begin{aligned}
P(\deg(1) = \deg(2) = k-1) &= \binom{\binom{n-1}{d-1} - \binom{n-2}{d-2}}{k-1}^2 \frac{\binom{\binom{n}{d} - 2\binom{n-1}{d-1} + \binom{n-2}{d-2}}{m_0 - 2(k-1)}}{\binom{\binom{n}{d}}{m_0}} + O(n^{-3}) \\
&= \left( 1 + O\left(\frac{\ln \ln n}{\ln n}\right) \right) \frac{e^{2\omega}}{((k-1)!)^2 n^2};
\end{aligned}$$

whence

$$E[X(X-1)] = \left( 1 + O\left(\frac{\ln \ln n}{\ln n}\right) \right) E[X]^2.$$

Consequently,

$$\text{var}[X] = E[X] + O\left((E[X])^2 \frac{\ln \ln n}{\ln n}\right).$$

By Chebyshev's Inequality,  $X$  is concentrated around its mean and in particular, w.h.p.  $X$  is in the interval (1). To finish the proof of part (i), it remains to show that w.h.p. there are no vertices of degree less than  $k - 1$ , which can be done by a first moment argument using similar techniques to the asymptotics of  $E[X]$ . Similarly, for part (ii), one can easily show that the expected number of vertices of degree  $k - 1$  in  $H_d(n, m_1)$  tends to zero as well.  $\square$

### 3 $H_d(n, m_0)$ is almost $k$ -connected

Now we will establish that w.h.p.  $H_d(n, m_0)$  is almost  $k$ -connected in the sense that if  $k - 1$  vertices are deleted, then there remains a massive component containing almost all left-over vertices. To this end, we prove an analogous statement for the random Bernoulli hypergraph  $H_d(n, p)$  and use a standard conversion lemma to obtain the desired result for  $H_d(n, m_0)$ . In this next lemma, we pick a specific version of  $m_0$ , one where  $\omega = \ln \ln \ln n$ .

**Lemma 3.1.** *Let  $m'_0 = \frac{n}{d} (\ln n + (k - 1) \ln \ln n - \ln \ln \ln n)$  and  $p = m'_0 / \binom{n}{d}$ . With high probability,*

(i)  $H_d(n, p)$  has the property “whichever  $k - 1$  vertices are deleted, there remains a giant component which includes all but up to  $\ln n$  leftover vertices.”

(ii)  $H_d(n, m'_0)$  has the property “whichever  $k - 1$  vertices are deleted, there remains a giant component which includes all but up to  $\ln n$  leftover vertices.”

*Proof.* (i) Given a set of  $k - 1$  vertices,  $\mathbf{v} = \{v_1, \dots, v_{k-1}\}$ , let  $\mathcal{F}(\mathbf{v})$  be the event that if the vertices  $\mathbf{v}$  are deleted from  $H_d(n, p)$  along with their incident edges, then there remains no components of size at least  $n - (k - 1) - \ln n$ . In particular, we wish to show that w.h.p.  $H_d(n, p)$  is not in  $\mathcal{F}(\mathbf{v})$  for any  $\mathbf{v}$ . Using the union bound over all  $k - 1$  element sets of  $[n]$  as well as symmetry, we find that

$$P\left(\bigcup_{\mathbf{v}} \mathcal{F}(\mathbf{v})\right) \leq \binom{n}{k-1} P(\mathcal{F}(\mathbf{v}^*)), \quad (4)$$

where  $\mathbf{v}^* = \{n - (k - 1) + 1, \dots, n - 1, n\}$ . Note that the remaining hypergraph left after deleting  $\mathbf{v}^*$  from  $H_d(n, p)$  is distributed as  $H_d(n', p)$ ,  $n' := n - (k - 1)$  (this is the primary reason that we consider the Bernoulli hypergraph  $H_d(n, p)$  rather than  $H_d(n, m)$ ). Therefore  $P(\mathcal{F}(\mathbf{v}^*))$  is precisely the probability that  $H_d(n', p)$  does not have a component of size at least  $n' - \ln n$ .

To bound  $P(\mathcal{F}(\mathbf{v}^*))$ , we note that any hypergraph on  $n'$  vertices without a component of size at least  $n' - \ln n$  has a set of vertices  $S$  such that there are no edges between  $S$  and  $[n'] \setminus S$  where  $|S| \in [\ln n, n' - \ln n]$ . To see this fact, consider a hypergraph  $H$  on  $n'$  vertices without such a large component and let  $L_1, \dots, L_\ell$  be the vertex sets of the

components of  $H$  in increasing order by their cardinalities. Then there is some minimal  $j$  so that

$$\ln n \leq \left| \cup_{i=1}^j L_i \right| < \ln n + n' - \ln n = n'.$$

Further,  $L_{j+1}$  is not empty and since  $|L_j| \leq |L_{j+1}|$ , we have that  $|L_j| \leq n'/2$  and

$$\ln n \leq \left| \cup_{i=1}^j L_i \right| < \ln n + n'/2 < n' - \ln n.$$

Clearly, there are no edges including a vertex of  $S := \cup_{i=1}^j L_i$  and  $[n'] \setminus S$ .

Therefore,

$$P(\mathcal{F}(\mathbf{v}^*)) \leq \sum_{s=\ln n}^{n'-\ln n} P(\exists S \subset [n'], |S| = s, \text{ no edge between } S \text{ and } [n'] \setminus S),$$

and by symmetry over all such vertex sets  $S$ ,

$$P(\mathcal{F}(\mathbf{v}^*)) \leq \sum_{s=\ln n}^{n'-\ln n} \binom{n'}{s} P(\text{no edge between } [s] \text{ and } [n'] \setminus [s]).$$

Further, this latter probability is symmetric about  $s = n'/2$  (i.e. the probabilities corresponding to  $s$  and  $n' - s$  are equal). Hence

$$P(\mathcal{F}(\mathbf{v}^*)) \leq 2 \sum_{s=\ln n}^{\lfloor n'/2 \rfloor} \binom{n'}{s} P(\text{no edge between } [s] \text{ and } [n'] \setminus [s]).$$

The number of potential edges that contain at least one vertex from  $[s]$  and at least one vertex from  $[n'] \setminus [s]$  is  $\binom{n'}{d} - \binom{s}{d} - \binom{n'-s}{d}$ . Hence

$$\begin{aligned} P(\mathcal{F}(\mathbf{v}^*)) &\leq 2 \sum_{s=\ln n}^{\lfloor n'/2 \rfloor} \binom{n'}{s} (1-p)^{\binom{n'}{d} - \binom{s}{d} - \binom{n'-s}{d}} \leq 2 \sum_{s=\ln n}^{\lfloor n'/2 \rfloor} \binom{n'}{s} e^{-p(\binom{n'}{d} - \binom{s}{d} - \binom{n'-s}{d})} \\ &=: 2E_1 + 2E_2, \end{aligned}$$

where  $E_1$  and  $E_2$  are the sums over  $S_1 := [\ln n, n/(\ln n)]$  and  $S_2 := (n/(\ln n), \lfloor n'/2 \rfloor]$  respectively. We begin with analyzing  $E_2$  since these bounds will be cruder and simpler.

Trivially,

$$E_2 \leq \sum_{s \in S_2} \binom{n'}{s} e^{-p \binom{n'}{d} + p \max_{t \in S_2} (\binom{t}{d} + \binom{n'-t}{d})} \leq 2n' e^{-p \binom{n'}{d} + p \max_{t \in S_2} (\binom{t}{d} + \binom{n'-t}{d})}. \quad (5)$$

Now let's take on these binomial coefficient terms. Trivially  $\binom{\nu}{d} \leq \frac{\nu^d}{d!}$ , and so

$$\binom{t}{d} + \binom{n'-t}{d} \leq \frac{t^d + (n'-t)^d}{d!}.$$

Further, the function  $f(t) = t^d + (n' - t)^d$  is decreasing for  $t \in S_2$ , so we have that

$$\max_{t \in S_2} \left( \binom{t}{d} + \binom{n' - t}{d} \right) \leq \frac{1}{d!} \left( \left( \frac{n}{\ln n} \right)^d + \left( n' - \frac{n}{\ln n} \right)^d \right).$$

Now our bound in (5) becomes

$$E_2 \leq 2^{n'} \exp \left( -p \binom{n'}{d} + \frac{p}{d!} \left( \frac{n}{\ln n} \right)^d + \frac{p}{d!} \left( n' - \frac{n}{\ln n} \right)^d \right).$$

In the previous expression, the leading order terms in the first and third terms will cancel and the middle term is absorbed in the error. Namely, we have that

$$E_2 \leq 2^{n'} \exp \left( -\frac{p(n')^d}{d!} + O(\ln n) + O \left( \frac{n \ln n}{(\ln n)^d} \right) + \frac{p(n')^d}{d!} - \frac{pd}{d!} \frac{n^d}{\ln n} + O \left( \frac{n}{\ln n} \right) \right),$$

or equivalently

$$\begin{aligned} E_2 &\leq 2^{n'} \exp \left( -\frac{pn^d}{(d-1)! \ln n} + O \left( \frac{n}{\ln n} \right) \right) \\ &= \exp((\ln 2 - 1)n + o(n)) \leq e^{-n/4}. \end{aligned} \tag{6}$$

Now let's take on the sum  $E_1$ . We begin with taking the leading order terms in the exponent.

$$E_1 \leq \sum_{s \in S_1} \binom{n'}{s} \exp \left( -p \binom{n'}{d} \left( 1 - \frac{\binom{n'-s}{d}}{\binom{n'}{d}} - O \left( \frac{s^d}{n^d} \right) \right) \right). \tag{7}$$

Uniformly over  $s \in S_1$ , we have that

$$\frac{\binom{n'-s}{d}}{\binom{n'}{d}} = \left( 1 - \frac{s}{n'} + O \left( \frac{s}{n^2} \right) \right)^d = 1 - \frac{ds}{n'} + O \left( \frac{s^2}{n^2} \right).$$

Consequently, the exponent in (7) is

$$p \binom{n'}{d} \left( \frac{ds}{n} + O \left( \frac{s^2}{n^2} \right) \right) = (\ln n + (k-1) \ln \ln n - \ln \ln \ln n) (s + O(s^2/n)) + O(1);$$

whence there is some (fixed)  $\gamma > 0$  such that for all  $s \in S_1$ ,

$$p \binom{n'}{d} \left( \frac{ds}{n} + O \left( \frac{s^2}{n^2} \right) \right) \geq (\ln n - \ln \ln \ln n) (s - \gamma s^2/n) - \gamma. \tag{8}$$

Using the bound  $\binom{n'}{s} \leq (en/s)^s$  as well as (8), (7) becomes

$$E_1 \leq \sum_{s \in S_1} \left( \frac{en}{s} \right)^s \exp \left( -(\ln n - \ln \ln \ln n) (s - \gamma s^2/n) + \gamma \right)$$

$$= e^\gamma \sum_{s \in S_1} \exp(s(1 - \ln s + \ln \ln \ln n + \gamma s(\ln n - \ln \ln \ln n)/n)).$$

However, for  $s \in S_1$ , we have that

$$s \leq \frac{n}{\ln n} \leq \frac{2n}{\ln n - \ln \ln \ln n} \implies s(\ln n - \ln \ln \ln n)/n \leq 2.$$

Therefore

$$E_1 \leq e^\gamma \sum_{s \in S_1} \exp(s(1 - \ln s + \ln \ln \ln n + 2\gamma)).$$

This sum is dominated by the first term ( $s = \ln n$ ) because ratios of consecutive terms uniformly tend to zero. Consequently, we have that

$$E_1 = O[\exp(-\ln n \ln \ln n + o(\ln n \ln \ln n))]. \quad (9)$$

Summing our bounds for  $E_1$  and  $E_2$  ((9) and (6), respectively), we have that

$$P(\mathcal{F}(\mathbf{v}^*)) \leq \exp(-\ln n \ln \ln n + o(\ln n \ln \ln n)), \quad (10)$$

which most definitely is  $o(n^{-(k-1)})$  and so by the bound (4), part (i) of the lemma is proved.

Part (ii) is established by using a standard conversion technique between  $H_d(n, m)$  and  $H_d(n, p)$ . For any hypergraph property  $\mathcal{A}$ , we have that

$$P(H_d(n, p) \in \mathcal{A}) = \sum_{m=0}^{\binom{n}{d}} P(H_d(n, m) \in \mathcal{A})P(e(H_d(n, p)) = m), \quad (11)$$

where  $e(H)$  is the number of edges of  $H$ . Therefore, for any (possible)  $m$ ,

$$P(H_d(n, p) \in \mathcal{A}) \geq P(H_d(n, m) \in \mathcal{A})P(e(H_d(n, p)) = m),$$

whence

$$P(H_d(n, m) \in \mathcal{A}) \leq \frac{P(H_d(n, p) \in \mathcal{A})}{P(e(H_d(n, p)) = m)}.$$

For  $m = \Theta(n \ln n)$  and  $p = m/\binom{n}{d}$ , one can show that

$$P(e(H_d(n, p)) = m) = \binom{\binom{n}{d}}{m} p^m (1-p)^{\binom{n}{d}-m} = \Theta(m^{-1/2}).$$

Hence in our case,

$$P(H_d(n, m'_0) \in \cup_{\mathbf{v}} \mathcal{F}(\mathbf{v})) = O\left(\sqrt{n \ln n} P(H_d(n, p) \in \cup_{\mathbf{v}} \mathcal{F}(\mathbf{v}))\right).$$

In the proof of part (i), we found that this latter probability tends to zero superpolynomially fast.  $\square$



## 4 Quasi-disjoint Edges

For the random graph process ( $d = 2$ ), it was found that the main barrier to  $k$ -connectivity is the presence of vertices of degree less than  $k$ , which could be isolated with the deletion of their neighbors (see Erdős-Rényi [5], Ivchenko [6], Bollobás [1],[2]). We will find a similar situation for the random hypergraph process.

However, we run into an additional issue here for hypergraphs. Even if the degree of a vertex  $v$  is  $k$ , we could isolate  $v$  with the deletion of less than  $k$  vertices. For instance, if all of  $v$ 's edges also include vertex  $w$ , then the deletion of just  $w$  from the hypergraph (along with its incident edges) will isolate  $v$  from the rest of the hypergraph. Our ultimate goal in this section is to show that w.h.p. each vertex of  $H_d(n, \tau_k)$  has at least  $k$  edges whose pairwise intersections are precisely  $\{v\}$ ; in this case, for any vertex, you would need to delete at least  $k$  of its neighbors to isolate it. To this end, we first prove that w.h.p.  $H_d(n, m_0)$  has this property for vertices with degree at least  $k$  and as nearly as could be expected for vertices with degree  $k - 1$ .

A set of edges  $E$  incident to vertex  $v$  is *quasi-disjoint* if all pairwise intersections of these edges are  $\{v\}$ ; formally, if  $e, f \in E$ ,  $e \neq f$ , then  $e \cap f = \{v\}$ .

**Lemma 4.1.** *Let  $m_0 = \frac{n}{d}(\ln n + (k - 1) \ln \ln n - \omega)$ , where  $\omega \rightarrow \infty$ , but  $\omega = o(\ln \ln n)$ . W.h.p.,  $H_d(n, m_0)$  is such that*

(i) *the incident edges of a degree  $k - 1$  vertex form a quasi-disjoint set,*

(ii) *vertices with degree at least  $k$  have a quasi-disjoint set of incident edges with size at least  $k$ .*

*Proof.* Note that both parts of this lemma are trivially true for  $k = 1$  and part (i) is also trivially true for  $k = 2$ . Let  $X(j, \ell)$  be the number of vertices whose maximum quasi-disjoint set has size  $j$  and whose degree is  $j + \ell$ . To prove this lemma, it suffices to show that w.h.p. for  $j \leq k - 1$  and  $\ell \geq 1$ , we have that  $X(j, \ell) = 0$ , which is shown by a first moment argument. Now

$$E[X(j, \ell)] = nP(j, \ell),$$

where  $P(j, \ell)$  is the probability that a generic vertex  $v$  has a maximum quasi-disjoint set of size  $j$  and whose degree is  $j + \ell$ . To bound this probability, note that  $v$  has a set of  $j$  quasi-disjoint edges and each of the remaining  $\ell$  edges must have at least one vertex from the  $j(d - 1)$  neighbors from the quasi-disjoint edges; further the remaining  $m_0 - j - \ell$  edges do not include  $v$ . Hence

$$\begin{aligned} P(j, \ell) &\leq \binom{\binom{n-1}{d-1}}{j} \binom{\binom{j(d-1)}{1} \binom{n-2}{d-2}}{\ell} \frac{\binom{\binom{n}{d} - \binom{n-1}{d-1}}{m_0 - j - \ell}}{\binom{\binom{n}{d}}{m_0}} \\ &\leq n^{(d-1)j} \left( \frac{e j (d-1) n^{d-2}}{\ell (d-2)!} \right)^\ell \left( \frac{m_0}{\binom{\binom{n}{d}}{d} - \binom{\binom{n-1}{d-1}}{d-1} - m_0} \right)^{j+\ell} \frac{\binom{\binom{n}{d} - \binom{n-1}{d-1}}{m_0}}{\binom{\binom{n}{d}}{m_0}}. \end{aligned}$$

We gave sharp asymptotics for the last fraction in (2). Here and throughout the rest of the paper, we will use  $f \leq_b g$  for  $f = O(g)$  when the formula for  $g$  becomes too bulky.

Therefore

$$P(j, \ell) \leq_b (\ln n)^j \frac{e^\omega}{n(\ln n)^{k-1}} \left( \frac{e j (d-1) n^{d-2} m_0}{\ell (d-2)! \left( \binom{n}{d} - \binom{n-1}{d-1} - m_0 \right)} \right)^\ell;$$

whence

$$P(j, \ell) \leq_b \frac{e^\omega}{n} \left( C \frac{\ln n}{n} \right)^\ell,$$

for  $C = 2(k-1)(d-1)$  (independent of  $j \leq k-1$  and  $\ell \geq 1$ ). Thus

$$\sum_{j=0}^{k-1} \sum_{\ell \geq 1} E[X(j, \ell)] \leq_b e^\omega \sum_{\ell \geq 1} \left( C \frac{\ln n}{n} \right)^\ell \leq_b e^\omega \frac{\ln n}{n} \rightarrow 0,$$

which completes the proof of the lemma.  $\square$

**Lemma 4.2.** *W.h.p. each vertex of  $H_d(n, \tau_k)$  has a quasi-disjoint set of incident edges with size at least  $k$ .*

*Proof.* This lemma is trivially true for  $k = 1$ . Suppose that  $k \geq 2$ . Let  $A_n$  be the event that  $H_d(n, \tau_k)$  has a vertex that does *not* have a quasi-disjoint set of edges with size at least  $k$ ; we wish to show that  $P(A_n) \rightarrow 0$ . Let  $m_0, m_1$  be as defined in Lemma 2.1. We have proved that w.h.p.  $\tau_k \in [m_0, m_1]$  and that  $H_d(n, m_0)$  does not have vertices of degree less than  $k-1$ . Further, w.h.p. the number of degree  $k-1$  vertices in  $H_d(n, m_0)$  is less than  $\frac{3e^\omega}{2(k-1)!}$  (Lemma 2.1). In addition, w.h.p.  $H_d(n, m_0)$  has the two properties of the previous lemma (Lemma 4.1). Let  $B_n$  be the intersection of these four likely events. To prove the lemma, it suffices to show that  $P(A_n \cap B_n) \rightarrow 0$ .

Let  $\tilde{V}_0$  be the vertex set of vertices of degree  $k-1$  in  $H_d(n, m_0)$ . Note that

$$P(A_n \cap B_n) = \sum_{V_0 \subset [n], |V_0| \leq 3e^\omega / (2(k-1)!)} P(A_n \cap B_n \cap \{\tilde{V}_0 = V_0\}).$$

On the event that  $A_n$  and  $B_n$  occur and  $\tilde{V}_0 = V_0$ , necessarily some edge  $e_m$  is added in the hypergraph process at some step  $m \in [m_0, m_1]$  such that  $e_m$  includes both a vertex  $v \in V_0$  and one of  $v$ 's  $(k-1)(d-1)$  neighbors in  $H_d(n, m_0)$ . Thus

$$\begin{aligned} P(A_n \cap B_n \cap \{\tilde{V}_0 = V_0\}) &\leq \sum_{m=m_0}^{m_1} \frac{\binom{|V_0|}{1} \binom{(k-1)(d-1)}{1} \binom{n-2}{d-2}}{\binom{n}{d} - m} P(\tilde{V}_0 = V_0) \\ &\leq_b (m_1 - m_0) e^\omega \frac{\binom{n-2}{d-2}}{\binom{n}{d} - m_1} P(\tilde{V}_0 = V_0) \leq_b \frac{\omega e^\omega}{n} P(\tilde{V}_0 = V_0). \end{aligned}$$

Therefore

$$P(A_n \cap B_n) \leq_b \frac{\omega e^\omega}{n} \sum_{V_0 \subset [n], |V_0| \leq 3e^\omega / (2(k-1)!)} P(\tilde{V}_0 = V_0) \leq \frac{\omega e^\omega}{n} \rightarrow 0.$$

$\square$

## 5 W.h.p. $H_d(n, \tau_k)$ is $k$ -connected

Now that we have sufficient knowledge about the structure of  $H_d(n, m_0)$  and low-degree vertices in  $H_d(n, \tau_k)$ , we can prove our main Theorem.

**Theorem 1.1.** *W.h.p.  $H_d(n, \tau_k)$  is  $k$ -connected. In short, w.h.p.  $\tau_k = T_k$ .*

*Proof.* Let  $m'_i = \frac{n}{d} (\ln n + (k-1) \ln \ln n + (-1)^{i+1} \ln \ln \ln n)$ , for  $i = 0, 1$ . By Lemma 2.1, we have shown that w.h.p.  $\tau_k \in [m'_0, m'_1]$  and by Lemma 4.2, each vertex of  $H_d(n, \tau_k)$  has a quasi-disjoint edge set of size at least  $k$ . Further, the property “whichever  $k-1$  vertices are deleted, there remains a giant component which includes all but up to  $\ln n$  leftover vertices,” denoted  $\mathcal{Q}$ , is an increasing property (closed under the addition of edges). Therefore, by Lemma 3.1,  $H_d(n, \tau_k)$  has property  $\mathcal{Q}$  as well. To prove this theorem, it suffices to show that the probability that these three likely events hold yet  $\tau_k < T_k$  tends to zero.

To this end, for  $m \in [m'_0, m'_1]$ , let  $C_m$  be the event that  $H_d(n, m)$  is not  $k$ -connected, but each vertex has a quasi-disjoint edge set of size at least  $k$  and  $H_d(n, m)$  has property  $\mathcal{Q}$ . To prove this theorem, it suffices to prove that  $P(\cup C_m) \rightarrow 0$ . We will in fact show that

$$P(C_m) \leq_b \frac{(\ln n)^{dk+k+1}}{n^{d-1}}, \quad (12)$$

uniformly over  $m \in [m'_0, m'_1]$ . In this case,  $P(\cup C_m) \leq_b \frac{(\ln n)^{dk+k+2}}{n^{d-2}} \rightarrow 0$ , as desired. All that remains is to prove the bound (12).

On the event  $C_m$ , there are  $k-1$  vertices,  $w_1, \dots, w_{k-1}$  such that upon their deletion, there is a component of size  $n' - s$  for some  $s \in [1, \ln n)$ . In fact, since each remaining vertex must have at least one incident edge, we must have that  $s \geq d$ . Let  $S$  be the set of vertices not in this large component. By the union bound over all  $k-1$  element sets of  $[n]$  and sets  $S$ ,  $|S| = s$ , as well as symmetry, we have that

$$P(C_m) \leq \binom{n}{k-1} \sum_{s=d}^{\ln n} \binom{n-(k-1)}{s} P_s,$$

where  $P_s$  is the probability that each vertex of  $H_d(n, m)$  has a quasi-disjoint edge set of size at least  $k$  and that after the deletion of  $\{n'+1, \dots, n'+(k-1) = n\}$  from  $H_d(n, m)$ , the vertices  $[n'-s]$  form a component; in this case,  $S = \{n'-s+1, \dots, n'\}$ . We now turn to showing that  $P_s$  tends to zero sufficiently fast.

Suppose that  $H$  is some hypergraph in the event corresponding to  $P_s$ . After the deletion of the  $k-1$  vertices, we know that vertex  $w := \{n'\}$  from  $S$  has at least one incident edge, which necessarily must reside completely within  $S$ . Further, before deletion, any incident edge to  $w$  must be completely contained within  $S$  or this edge must contain one of the  $k-1$  to-be-deleted vertices. Moreover, there are at least  $k$  edges incident to  $w$  before the deletion. Therefore

$$P_s \leq \sum_{i=1}^k P_s(i),$$

where  $P_s(i)$  is the corresponding probability to when there are (at least)  $i$  incident edges to  $w$  contained within  $S$  and (at least)  $k - i$  incident edges to  $w$  that contain at least one of the to-be-deleted vertices. To bound the number of hypergraphs contributing to  $P_s(i)$ , we choose  $i$  potential edges within  $S$  containing  $w$ ,  $k - i$  potential edges that include  $w$  and at least one to-be-deleted vertex; then we choose the remaining  $m - k$  edges among all potential edges except those that include a vertex of  $S$  and  $d - 1$  vertices of  $[n' - s]$  (which necessarily can not be present). Note that these last chosen edges can include  $w$  as well. Therefore

$$P_s(i) \leq \binom{\binom{s-1}{d-1}}{i} \binom{\binom{1}{1} \binom{k-1}{1} \binom{n-2}{d-2}}{k-i} \binom{\binom{n}{d} - \binom{s}{1} \binom{n'-s}{d-1}}{m-k} \frac{1}{\binom{\binom{n}{d}}{m}}.$$

First, we use trivial bounds on the first two binomial terms. Then we use the inequality  $\binom{N-j}{j} \leq \binom{N}{j} e^{-j\ell/N}$ . Namely, note that

$$\begin{aligned} P_s(i) &\leq s^{di} k^k n^{(d-2)(k-i)} \left( \frac{m}{\binom{n}{d} - s \binom{n'-s}{d-1} - m} \right)^k \frac{\binom{\binom{n}{d} - s \binom{n'-s}{d-1}}{m}}{\binom{\binom{n}{d}}{m}} \\ &\leq_b s^{dk} n^{(d-2)(k-1)} \left( \frac{\ln n}{n^{d-1}} \right)^k \exp \left( -\frac{s \binom{n'-s}{d-1} m}{\binom{n}{d}} \right). \end{aligned}$$

Further, for  $s \leq \ln n$ , we have that

$$\frac{s \binom{n'-s}{d-1} m}{\binom{n}{d}} = \frac{s d m}{n} + O \left( \frac{(\ln n)^2}{n} \right) \geq \frac{s d m'_0}{n} + o(1).$$

Hence

$$\begin{aligned} P_s(i) &\leq_b (\ln n)^{dk} n^{(d-2)(k-1)} \left( \frac{\ln n}{n^{d-1}} \right)^k \left( \frac{\ln \ln n}{n (\ln n)^{k-1}} \right)^s \\ &\leq (\ln n)^{dk+k} n^{2-d-k} \left( \frac{\ln \ln n}{n (\ln n)^{k-1}} \right)^s, \end{aligned}$$

which no longer depends on  $i$ . Therefore

$$P_s \leq_b (\ln n)^{dk+k} n^{2-d-k} \left( \frac{\ln \ln n}{n (\ln n)^{k-1}} \right)^s,$$

and

$$P(C_m) \leq_b n^{k-1} \sum_{s=d}^{\ln n} \frac{n^s}{s!} (\ln n)^{dk+k} n^{2-d-k} \left( \frac{\ln \ln n}{n (\ln n)^{k-1}} \right)^s.$$

Now taking on this sum, we find that

$$P(C_m) \leq_b \frac{(\ln n)^{dk+k}}{n^{d-1}} \sum_{s=d}^{\ln n} \frac{1}{s!} \left( \frac{\ln \ln n}{(\ln n)^{k-1}} \right)^s \leq \frac{(\ln n)^{dk+k}}{n^{d-1}} \exp \left( \frac{\ln \ln n}{(\ln n)^{k-1}} \right),$$

and we find that  $P(C_m) \leq_b \frac{(\ln n)^{dk+k+1}}{n^{d-1}}$ , as desired.

□

## 6 Sharp Threshold of $k$ -connectivity

As a consequence of Theorem 1.1, for any  $m$ , we have that

$$\begin{aligned} P(H_d(n, m) \text{ is } k\text{-connected}) &= P(T_k \leq m) = P(\tau_k \leq m) + o(1) \\ &= P(\min\text{-deg } H_d(n, m) \geq k) + o(1). \end{aligned} \quad (13)$$

We use this fact to determine the probability that  $H_d(n, m)$  and  $H_d(n, p)$  is  $k$ -connected in the critical window.

**Corollary 1.2.** *(i) Let  $m = \frac{n}{d}(\ln n + (k-1)\ln \ln n + c_n)$ , where  $c_n \rightarrow c \in \mathbb{R}$ . W.h.p.  $H_d(n, m)$  is  $(k-1)$ -connected, but not  $(k+1)$ -connected. Further the probability that  $H_d(n, m)$  is  $k$ -connected tends to  $e^{-e^{-c}/(k-1)!}$ .*

*(ii) Let  $p = (d-1)! \frac{\ln n + (k-1)\ln \ln n + c_n}{n^{d-1}}$ , where  $c_n \rightarrow c \in \mathbb{R}$ . W.h.p.  $H_d(n, p)$  is  $(k-1)$ -connected, but not  $(k+1)$ -connected. Further the probability that  $H_d(n, p)$  is  $k$ -connected tends to  $e^{-e^{-c}/(k-1)!}$ .*

*Proof.* **(i)** First, note that w.h.p.  $\tau_{k-1} < m$  and  $\tau_{k+1} > m$  by Lemma 2.1. Therefore, by Theorem 1.1, w.h.p.  $H_d(n, m)$  is  $(k-1)$ -connected, but not  $(k+1)$ -connected. In the lemma following this proof, we show that  $X$ , the number of vertices of degree  $k-1$  in  $H_d(n, m)$  is asymptotically Poisson with parameter  $e^{-c}/(k-1)!$ . Thus

$$P(\min\text{-deg } H_d(n, m) \geq k) = P(\text{Poi}(e^{-c}/(k-1)!) = 0) + o(1) = e^{-e^{-c}/(k-1)!} + o(1).$$

Using the equation (13) finishes off the proof.

**(ii)** This part will be proved from **(i)** using a standard conversion technique similar to the one used in Lemma 3.1. Since the number of edges in  $H_d(n, p)$ , denoted  $e(H_d(n, p))$ , is binomially distributed on  $N := \binom{n}{d}$  trials with success probability  $p$ , we have that

$$e(H_d(n, p)) = Np + O_p\left(\sqrt{Np(1-p)}\right) = \frac{n}{d}(\ln n + (k-1)\ln \ln n + c_n) + O_p\left(\sqrt{n \ln n}\right).$$

Therefore, if  $m_{+,-} := \frac{n}{d}(\ln n + (k-1)\ln \ln n + c_n^{\pm,-})$ , where  $c_n^{\pm,-} = c_n \pm \ln n/\sqrt{n}$ , then w.h.p.  $m_- \leq e(H_d(n, p)) \leq m_+$ . Using this fact, (11) becomes

$$P(H_d(n, p) \in \mathcal{A}) = \sum_{m=m_-}^{m_+} P(H_d(n, m) \in \mathcal{A})P(e(H_d(n, p)) = m) + o(1).$$

Notice that  $c_n^{\pm,-} \rightarrow c$ . By part **(i)**, if  $\mathcal{A}$  is  $\{(k-1)\text{-connected}\}$  or  $\{\text{not } (k+1)\text{-connected}\}$ , then  $P(H_d(n, m_{+,-}) \in \mathcal{A}) \rightarrow 1$ ; also, if  $\mathcal{A} = \{k\text{-connected}\}$ , then  $P(H_d(n, m_{+,-}) \in \mathcal{A}) \rightarrow e^{-e^{-c}/(k-1)!}$ . To finish off the proof that  $P(H_d(n, p) \in \mathcal{A})$  has the same limits, we will use the fact that these properties are monotone.

Now an *increasing* (*decreasing*) property is a property that is closed under the addition (deletion resp.) of edges. For any increasing property  $\mathcal{A}$ , we have that  $P(H_d(n, m) \in \mathcal{A}) \leq P(H_d(n, m') \in \mathcal{A})$  for any  $m \leq m'$ . This fact is obvious when you consider  $H_d(n, m)$  and

$H_d(n, m')$  to be snapshots of the random hypergraph process  $\{H_d(n, \mu)\}_{\mu=0}^N$ . Moreover, if  $\mathcal{A}$  is an increasing property, then

$$P(H_d(n, m_-) \in \mathcal{A}) + o(1) \leq P(H_d(n, p) \in \mathcal{A}) \leq P(H_d(n, m_+) \in \mathcal{A}) + o(1);$$

further, if  $\mathcal{A}$  is decreasing, then the inequalities above are reversed. Consequently, for a monotone property  $\mathcal{A}$  such that both  $P(H_d(n, m_{+,-}) \in \mathcal{A})$  tend to the same number, then  $P(H_d(n, p) \in \mathcal{A})$  does as well.  $\square$

**Lemma 6.1.** *Let  $m = \frac{n}{d}(\ln n + (k-1) \ln \ln n + c_n)$ , where  $c_n \rightarrow c \in \mathbb{R}$ . W.h.p. the number of vertices of degree  $k-1$ , denoted by  $X$ , converges in distribution to a Poisson random variable with parameter  $e^{-c}/(k-1)!$ .*

*Proof.* We prove this lemma using the method of moments (see [2] for a description of this method). In order to prove the lemma, it suffices to show that for each  $r \in \mathbb{N}$  (fixed),

$$\lim_{n \rightarrow \infty} E[(X)_r] = \left( \frac{e^{-c}}{(k-1)!} \right)^r.$$

To compute the  $r$ 'th factorial moment, note that

$$E[(X)_r] = E_0 + E_1 + \dots + E_{(k-1)r},$$

where  $E_j$  is the expected number of ordered  $r$ -tuples of vertices of degree at most  $k-1$  such that there are exactly  $(k-1)r-j$  edges containing at least one of these vertices. We will see that the terms other than  $E_0$  are negligible.

Let's first consider  $E_0$ . Since the number of edges is  $(k-1)r$ , each of these  $r$  vertices have degree  $k-1$  and must necessarily not be adjacent; so

$$E_0 = (n)_r \binom{n-r}{d-1}^r \frac{\binom{n-r}{d}^{m-r(k-1)}}{\binom{n}{d}^m} = (1 + O(n^{-1})) n^r \left( \frac{n^{(k-1)(d-1)}}{(k-1)!((d-1)!)^{k-1}} \right)^r \frac{\binom{n-r}{d}^{m-r(k-1)}}{\binom{n}{d}^m}.$$

Taking on this last factor, we have that

$$\binom{\binom{n-r}{d}}{m-r(k-1)} = \binom{\binom{n-r}{d}}{m} \frac{(m)_{r(k-1)}}{\left( \binom{n-r}{d} - m + r(k-1) \right)_{r(k-1)}}.$$

By sharply approximating this last fraction, we have that

$$E_0 = \left( 1 + O\left( \frac{\ln \ln n}{\ln n} \right) \right) \left( \frac{n (\ln n)^{k-1}}{(k-1)!} \right)^r \frac{\binom{n-r}{d}^m}{\binom{n}{d}^m}.$$

Note that

$$\binom{\binom{n-r}{d}}{m} = \frac{1}{m!} \left( \binom{n-r}{d} \right)^m \prod_{i=0}^{m-1} \left( 1 - \frac{i}{\binom{n-r}{d}} \right) = \frac{1}{m!} \left( \binom{n-r}{d} \right)^m \left( 1 + O\left( \frac{m^2}{n^d} \right) \right).$$

We also have  $d \geq 3$  so that

$$\binom{\binom{n-r}{d}}{m} = \frac{1}{m!} \binom{n-r}{d}^m \left(1 + O\left(\frac{(\ln n)^2}{n}\right)\right);$$

similarly, we have that

$$\binom{\binom{n}{d}}{m} = \frac{1}{m!} \binom{n}{d}^m \left(1 + O\left(\frac{(\ln n)^2}{n}\right)\right).$$

Further

$$\left(\frac{\binom{n-r}{d}}{\binom{n}{d}}\right)^m = \left(\left(\frac{n-r}{n} + O\left(\frac{1}{n^2}\right)\right)^d\right)^m = \exp\left(dm\left(-\frac{r}{n} + O(n^{-2})\right)\right),$$

and

$$\frac{\binom{\binom{n-r}{d}}{m}}{\binom{\binom{n}{d}}{m}} = \left(1 + O\left(\frac{(\ln n)^2}{n}\right)\right) e^{-rdm/n} = \left(1 + O\left(\frac{(\ln n)^2}{n}\right)\right) \left(\frac{e^{-c_n}}{n(\ln n)^{k-1}}\right)^r. \quad (14)$$

Thus

$$E_0 = \left(1 + O\left(\frac{\ln \ln n}{\ln n}\right)\right) \left(\frac{e^{-c_n}}{(k-1)!}\right)^r.$$

Now we turn to  $E_j$ , for  $j \geq 1$ . Note that  $E_j$  is less than the expected number of such  $r$ -tuples where these  $r$  vertices have exactly  $r(k-1) - j$  adjacent edges (we dropped the degree condition). Then, for  $j \geq 1$ , we have that

$$E_j \leq \binom{n}{r} \frac{\binom{\binom{n-r}{d}}{m - (k-1)r + j}}{\binom{\binom{n}{d}}{m}}.$$

In particular, we have that

$$E_j \leq n^r \left(\binom{n}{d} - \binom{n-r}{d}\right)^{(k-1)r-j} \left(\frac{m}{\binom{n-r}{d} - m}\right)^{(k-1)r-j} \frac{\binom{\binom{n-r}{d}}{m}}{\binom{\binom{n}{d}}{m}}.$$

Using the fact that

$$\binom{n}{d} - \binom{n-r}{d} = r \binom{n-r}{d-1} + O(n^{d-2}) \leq rn^{d-1}$$

along with the bound (14), we have that

$$E_j \leq_b n^{r+(d-1)[(k-1)r-j]} \left(\frac{\ln n}{n^{d-1}}\right)^{(k-1)r-j} \left(\frac{1}{n(\ln n)^{k-1}}\right)^r = \frac{1}{(\ln n)^j},$$

which completes the proof of the lemma. □

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