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- In C and C++ perennial source of security vulnerabilities
- In Java, C# program crashes due to exceptions



# Preventing Memory Safety Errors

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  - Can be hard to understand and fix detected errors
  - Does not help programmer write safe-by-construction code
- **Dynamic approaches** to memory safety
  - + Widely used in managed languages
  - But only transforms vulnerability into program crash
  - Run-time overhead

## Key Idea: Use Program Synthesis

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- 1 Programmer specifies **which** parts of the program should be guarded

Example:

```
if(???) {R} else { /* handle error */ }
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# Key Idea: Use Program Synthesis

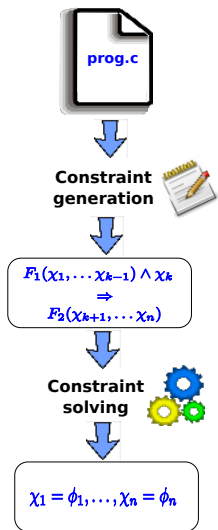
## Key Idea: Program synthesis to guarantee memory safety

- 1 Programmer specifies **which** parts of the program should be guarded
- 2 Our technique **synthesizes** correct and optimal guards that guarantee memory safety
  - Optimal means as weak and as simple as possible

Example:

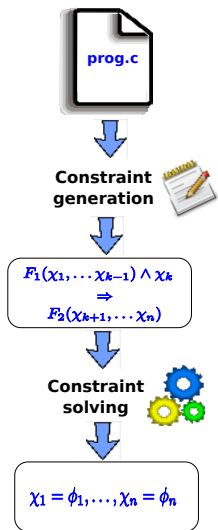
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# Solution Overview



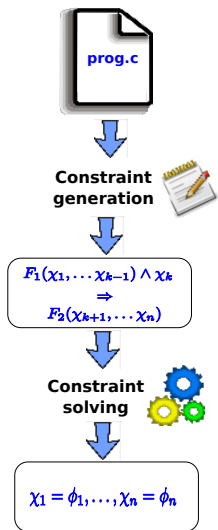
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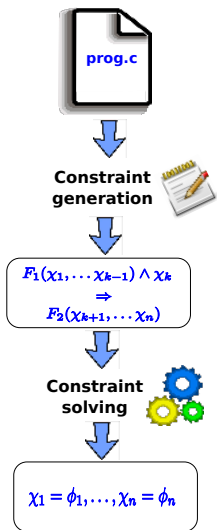
- Represent unknown guards using placeholders



## 1 Constraint Generation:

- Represent unknown guards using placeholders
- Perform dual forward and backward analysis to generate constraint for each unknown

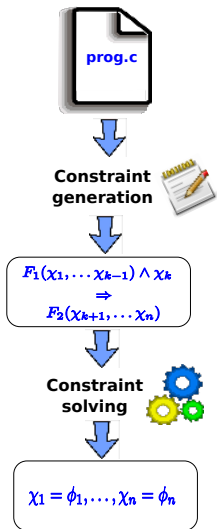
# Solution Overview



## 1 Constraint Generation:

- Represent unknown guards using placeholders
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## 2 Constraint Solving:



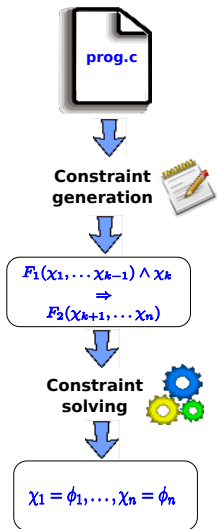
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- An extended abduction algorithm for solving constraint system with multiple unknowns

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## 1 Constraint Generation:

- Represent unknown guards using placeholders
- Perform dual forward and backward analysis to generate constraint for each unknown

## 2 Constraint Solving:

- An extended abduction algorithm for solving constraint system with multiple unknowns
- Guarantees Pareto-optimality

# Constraint Generation Overview

$\phi$  

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if(??)
{
    ...
}
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- At synthesis point, compute postcondition  $\phi$  of code above ??

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- Compute precondition  $\psi$  that ensures **memory safety** of code guarded by ??

$\psi$

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if(??)  
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  ...  
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# Constraint Generation Overview

$\phi$

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$\psi$

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if(??)  
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- At synthesis point, compute postcondition  $\phi$  of code above ??
- Compute precondition  $\psi$  that ensures **memory safety** of code guarded by ??
- Condition to guarantee memory safety:

$$\phi \wedge ?? \models \psi$$



# Key Insight: Abduction

**Solution:** Abductive inference





## **Solution:** Abductive inference

- Given facts  $F$  and desired outcome  $O$ , find simple explanatory hypothesis  $E$  such that

$$F \wedge E \models O \text{ and } \text{SAT}(F \wedge E)$$



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- $O \equiv$  memory safety precondition  $\psi$
- $E \equiv$  Solution for ??



$$F_1(\vec{\chi}_1) \wedge \chi_1 \Rightarrow F'_1(\vec{\chi}_1)$$

$$F_2(\vec{\chi}_2) \wedge \chi_2 \Rightarrow F'_2(\vec{\chi}_2)$$

⋮

$$F_n(\vec{\chi}_n) \wedge \chi_n \Rightarrow F'_n(\vec{\chi}_n)$$

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- Cannot directly use abduction because constraints have multiple unknowns
- New iterative, stratification-based algorithm for solving constraint system
- Uses abduction as a helper procedure
- Resulting solution is Pareto-optimal
  - Cannot improve solution for one unknown without making others worse

## Example

- Code snippet from Unix Coreutils with protected memory access

```
int main(int argc,
        char** argv)
{
    if(argc<=1) return -1;
    argv++; argc--;

    optind=0;
    while(...) {
        optind++;
        if(*) {argv++;
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    if(??) {
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```

## Example

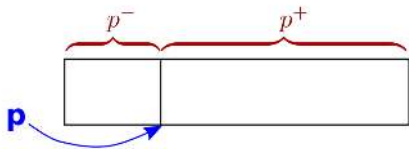
- Code snippet from Unix Coreutils with protected memory access
- **Convention:** For pointer  $p$ :

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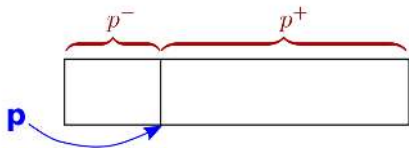


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# Example

- Code snippet from Unix Coreutils with protected memory access
- **Convention:** For pointer  $p$ :
  - $p^+$  represents distance to end of memory block
  - $p^-$  represents distance from beginning of memory block



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int main(int argc,
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## Example Cont.

- **First Step:** Compute what is known at ??  $\Rightarrow$  **postcondition**  $\phi$

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$$argv^+ = argc \wedge argv^- = 0$$

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- From language semantics:

$$argv^+ = argc \wedge argv^- = 0$$

- From computing the strongest postcondition:

$$argv^+ = argc \wedge \\ argv^- \geq 1 \wedge optind \geq 0$$

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- **Second Step:** Compute what needs to hold at ?? to ensure memory safety  
⇒ precondition  $\psi$

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int main(int argc,
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## Example Cont.

- **Second Step:** Compute what needs to hold at ?? to ensure memory safety  
⇒ precondition  $\psi$
- Buffer access:

$$\begin{aligned}optind + 1 &< argv^+ \wedge \\optind + 1 &\geq -argv^-\end{aligned}$$

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- Solve abduction problem

$\phi \wedge ?? \models \psi$  where

$$\phi : \quad \begin{array}{l} argv^+ = argc \wedge \\ argv^- \geq 1 \wedge optind \geq 0 \end{array}$$

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- **Solution:**  $argc - optind > 1$

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- Removed conditionals used to prevent memory safety errors



- Evaluated technique on the Unix Coreutils and parts of OpenSSH
- Removed conditionals used to prevent memory safety errors
- Used our new technique to synthesize the missing guards





## Experiments Cont.

| Program                 | Lines | # holes | Time (s) | Memory | Synthesis successful? | Bug? |
|-------------------------|-------|---------|----------|--------|-----------------------|------|
| Coreutils hostname      | 160   | 1       | 0.15     | 10 MB  | Yes                   | No   |
| Coreutils tee           | 223   | 1       | 0.84     | 10 MB  | Yes                   | Yes  |
| Coreutils runcon        | 265   | 2       | 0.81     | 12 MB  | Yes                   | No   |
| Coreutils chroot        | 279   | 2       | 0.53     | 23 MB  | Yes                   | No   |
| Coreutils remove        | 710   | 2       | 1.38     | 66MB   | Yes                   | No   |
| Coreutils nl            | 758   | 3       | 2.07     | 80 MB  | Yes                   | No   |
| SSH - sshconnect        | 810   | 3       | 1.43     | 81 MB  | Yes                   | No   |
| Coreutils mv            | 929   | 4       | 2.03     | 42 MB  | Yes                   | No   |
| SSH - do_authentication | 1,904 | 4       | 3.92     | 86 MB  | Yes                   | Yes  |
| SSH - ssh_session       | 2,260 | 5       | 4.35     | 81 MB  | Yes                   | No   |

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# Summary

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- Two key ingredients:

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- Two key ingredients:
  - **Constraint generation:** Generates VCs with placeholders using dual forward and backward reasoning
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- Experimental validation of our approach

# Questions?

