# Q-Learning-Based Teaching-Learning Optimization for Distributed Two-Stage Hybrid Flow Shop Scheduling with Fuzzy Processing Time 

Bingjie Xi and Deming Lei*


#### Abstract

Two-stage hybrid flow shop scheduling has been extensively considered in single-factory settings. However, the distributed two-stage hybrid flow shop scheduling problem (DTHFSP) with fuzzy processing time is seldom investigated in multiple factories. Furthermore, the integration of reinforcement learning and metaheuristic is seldom applied to solve DTHFSP. In the current study, DTHFSP with fuzzy processing time was investigated, and a novel Q-learning-based teaching-learning based optimization (QTLBO) was constructed to minimize makespan. Several teachers were recruited for this study. The teacher phase, learner phase, teacher's self-learning phase, and learner's self-learning phase were designed. The Q-learning algorithm was implemented by 9 states, 4 actions defined as combinations of the above phases, a reward, and an adaptive action selection, which were applied to dynamically adjust the algorithm structure. A number of experiments were conducted. The computational results demonstrate that the new strategies of QTLBO are effective; furthermore, it presents promising results on the considered DTHFSP.


Key words: teaching-learning based optimization; Q-learning algorithm; two-stage hybrid flow shop scheduling; fuzzy processing time

## 1 Introduction

A two-stage hybrid flow shop scheduling problem (THFSP) consists of two stages, at least one of which consists of parallel machines. This NP-hard problem is a special case of the hybrid flow shop scheduling problem (HFSP). This topic has attracted much attention in the past decade, and a number of results have been obtained in single-factory ${ }^{[1-14]}$ and multifactory settings ${ }^{[15,16]}$.

Various methods, such as the exact, heuristic, and metaheuristic methods, have been applied to solve THFSP in a single-factory setting. For instance, Allaoui and Artiba ${ }^{[1]}$ studied the problem with availability constraint by using an exact method and several heuristics to minimize makespan. Yang ${ }^{[2]}$

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developed some heuristics for particular cases of the problem with dedicated machines. Wang and Liu ${ }^{[3]}$ presented a branch-and-bound based heuristic for solving the problem with dedicated machines. Wang et al. ${ }^{[4]}$ proposed a problem-tailored constructive heuristic with local search for bi-criteria energyefficient THFSP. Feng et al. ${ }^{[5]}$ developed an exact method and two heuristics for THFSP with interval processing time, while Wang and Liu ${ }^{[6]}$ proposed a genetic algorithm for no-wait THFSP. Tan et al. ${ }^{[7]}$ presented a hybrid decomposition approach with variable neighborhood search for the problem with batch processing machines. Fan et al. ${ }^{[8]}$ presented a mutant firefly algorithm to optimize the simultaneous rate for the arrival and on-time delivery rate. Past studies also solved the THFSP by using other metaheuristics, such as artificial immune system ${ }^{[9]}$, tabu search ${ }^{[4,10]}$, estimation of distribution algorithm ${ }^{[11]}$, artificial bee colony ${ }^{[12]}$, and imperialist competitive algorithm ${ }^{[13]}$. Lang et al. ${ }^{[14]}$ investigated the neuro-evolution of augmenting topology to generate and parameterize artificial neural networks on
determining allocation and sequencing decisions in THFSP with family setup times.
In recent years, production has transferred from single factories to multiple factories with the further development of globalization. As a result, distributed scheduling problems in multiple factories have become the main topic of production scheduling in recent years. The distributed two-stage hybrid flow shop scheduling problem (DTHFSP) is also considered in this work. Lei and Wang ${ }^{[15]}$ and Cai et al. ${ }^{[16]}$ proposed a shuffled frogleaping algorithm with memeplex grouping and a collaborative variable search (CVS), respectively for DTHFSP with makespan minimization and DTHFSP with sequence-dependent setup time (SDST) and fuzzy processing time.

As stated above, previous works have mainly investigated the THFSP in single-factory settings, while the DTHFSP has been seldom studied. The THFSP exists in many real-life manufacturing processes, with extensive applications in multifactory production in many manufacturing industries, such as semiconductors. To quickly respond to rapid changes in customer demands and global market requirements, it is necessary to extend the THFSP by considering the DTHFSP. However, uncertainty extensively exists in real-life production processes, and these cannot be neglected intentionally because the obtained schedule would not be able to meet the requirements of real-life production once uncertainty is neglected. Moreover, neglecting uncertainty in multiple factories can lead to more losses than neglecting it in a single factory, so it is important to deal with the DTHFSP with uncertainty.
Uncertainty is often depicted by fuzzy theory. Fuzzy scheduling problems have been considered in single ${ }^{[17-25]}$ and multiple factories ${ }^{[16,26,27]}$. For example, Shao et al. ${ }^{[26]}$ studied a distributed fuzzy blocking flow shop scheduling problem in multiple homogeneous factories, after which they presented two constructive heuristics and two iterated greedy to minimize fuzzy makespan. Zheng et al. ${ }^{[27]}$ dealt with multi-objective fuzzy distributed hybrid flow shop scheduling and developed a cooperative coevolution algorithm with an iterated greedy estimation of distribution algorithm. The DTHFSP with fuzzy processing time is solved by using CVS ${ }^{[16]}$. Obviously, fuzzy distributed scheduling does not attract sufficient attention, and the DTHFSP with fuzzy processing time is seldom considered.

Meanwhile, teaching-learning based optimization (TLBO) is a population-based algorithm inspired by the act of passing on knowledge from a teacher to a student within a classroom. The TLBO possesses a simple structure and fewer parameters and is easily understood and implemented ${ }^{[28]}$. In recent years, the TLBO has been successfully used to solve various production scheduling problems in single and multiple factories ${ }^{[29-36]}$. The competitive performances of the TLBO have been verified and tested in these works.

At the same time, the integration of reinforcement learning (RL) into metaheuristics has also become a hot topic, with some studies investigating production scheduling in recent years ${ }^{[37-41]}$. The integration of RL and metaheuristic can lead to the dynamic selection of search operators or adaptively adjusted parameter settings, among others ${ }^{[42,43]}$. As a result, integrating RL with metaheuristic can improve the performance of the latter, making it an effective approach to obtaining high quality solutions.

The extensive applications of the TLBO to scheduling and the positive impact of the integration of RL and metaheuristic motivate us to construct a Q-learning-based teaching-learning based optimization (QTLBO) to solve the DTHFSP with fuzzy processing time. Moreover, RL is used to dynamically adjust the TLBO's algorithm structure, thereby resulting in new applications of RL with metaheuristics.

In this study, the DTHFSP with fuzzy processing time is studied, and a novel algorithm called the QTLBO is constructed through the integration of the Q-learning algorithm and the TLBO to minimize makespan. We recruited a group of teachers and adopted four phases: teacher, learner, teacher's selflearning, and learner's self-learning. The Q-learning algorithm is constructed by 9 states, 4 actions, a reward, and an adaptive action selection to dynamically adjust the algorithm structure of the QTLBO. After conducting extensive experiments, the computational results demonstrate that the new strategies are effective and efficient. Furthermore, the QTLBO is a competitive algorithm for the considered problem.

The remainder of the paper is organized as follows. The problem description is given in Section 2, followed by an introduction to the TLBO and Q-learning in Section 3. Section 4 presents the QTLBO for the problem. Numerical experiments on the QTLBO are reported in Section 5, and the conclusions are
summarized in Section 6. Some topics for future research are also proposed in Section 6.

## 2 Problem Description

The DTHFSP with fuzzy processing time consists of $n$ jobs $J_{1}, J_{2}, \ldots, J_{n}$ processed by $F$ factories located in different sites. Each factory $f$ possesses a two-stage hybrid flow shop with $m_{f}$ identical parallel machines $M_{1, S_{f}+1}, M_{1, S_{f}+2}, \ldots, M_{1, S_{f}+m_{f}}$ at the first stage and a single machine $M_{2, f}$ at the second stage, $S_{1}=0$. Here $S_{f}=\sum_{l=1}^{f-1} m_{l}, f>1, W=S_{F+1}$ indicates the total number of parallel machines. All jobs are available at time zero.

The triangular fuzzy number (TFN) is used to depict the processing time and setup time. Specifically, $p_{i l f}=\left(p_{i l f}^{1}, p_{i l f}^{2}, p_{i l f}^{3}\right)$ is the processing time of job $J_{i}$ at stage $l$ of factory $f$. SDST is considered. Meanwhile, $u_{j i l f}=\left(u_{j i s f}^{1}, u_{j i s f}^{2}, u_{\text {jilf }}^{3}\right)$ denotes the setup time for job $J_{i}$ on each machine of stage $l$ in factory $f$ when job $J_{j}$ is processed directly before it on the same machine. At the same time, $u_{0 i l f}=\left(u_{0 i s f}^{1}, u_{0 i s f}^{2}, u_{0 i l f}^{3}\right)$ is the initial fuzzy setup time of the first job $J_{i}$ processed on the machine at stage $l$ of factory $f$.
The DTHFSP has some constraints on jobs and machines: (1) each machine can process at most one operation at a time, (2) no jobs may be processed on more than one machine at a time, (3) operations cannot be interrupted, and (4) all machines are available at all times, that is, maintenance, failure, and breakdown are not considered. Figure 1 shows the layout of the DTHFSP.
The DTHFSP possesses three subproblems: the


Fig. 1 Layout of the distributed two-stage hybrid flow shop scheduling problem (DTHFSP).
factory assignment used to decide which jobs are assigned to each factory, machine assignment, and scheduling. Cai et al. ${ }^{[16]}$ analyzed the relationships among these subproblems and combined two assignment problems into machine assignments.
The goal of the problem is to minimize makespan until all constraints are met:

$$
\begin{equation*}
C_{\max }=\max _{i=1,2, \ldots, n}\left\{C_{i}\right\} \tag{1}
\end{equation*}
$$

where $C_{\max }=\left(C_{\max }^{1}, C_{\max }^{2}, C_{\max }^{3}\right)$ is the fuzzy makespan, and $C_{i}$ indicates the fuzzy completion time of job $J_{i}$.

When the TFN is used to represent the processing data of the scheduling problem, three operators are often used to build a schedule: addition operator, max operator, and ranking operator.

The addition operator is used to compute the completion time of the job and is defined for TFNs $B=\left(b_{1}, b_{2}, b_{3}\right)$ and $H=\left(h_{1}, h_{2}, h_{3}\right)$ by

$$
\begin{equation*}
B+H=\left(b_{1}+h_{1}, b_{2}+h_{2}, b_{3}+h_{3}\right) \tag{2}
\end{equation*}
$$

The ranking operator is about the comparison regarding fuzzy completion time. If $c_{1}(B)>c_{1}(H)$ or $c_{1}(B)=c_{1}(H)$ and $\quad c_{2}(B)>c_{2}(H)$ or $\quad c_{1}(B)=c_{1}(H)$, $c_{2}(B)=c_{2}(H)$, and $c_{3}(B)>c_{3}(H)$, then $B>H$, where $c_{3}(B)=b_{3}-b_{1}, c_{1}(B)=\left(b_{1}+2 b_{2}+b_{3}\right) / 4$, and $c_{2}(B)=b_{2}$.

The max operator is applied to decide the fuzzy beginning time and is given by

$$
\begin{equation*}
\text { if } B>H \text {, then } B \vee H=B \text {; else } B \vee H=H \tag{3}
\end{equation*}
$$

Cai et al. ${ }^{[16]}$ also considered the above problem and provided an example and a schedule. With respect to makespan, each job $J_{i}$ has a fuzzy completion time $C_{i}$, then $C_{\text {max }}=C_{1} \vee C_{2} \vee \cdots \vee C_{n}$.

## 3 Introduction to the TLBO and the $Q$ Learning Algorithm

The search process of the TLBO acts on population $P$, which is considered as a class of learners and has a teacher phase and a learner phase. In the teacher phase, a teacher $x_{\text {teacher }}$, which is the best solution obtained so far, passes its knowledge to learners, and a solution $x$ is modified by

$$
\begin{equation*}
x_{\text {new }}=x+\operatorname{rand} \times\left(x_{\text {teacher }}-T \times x_{\text {mean }}\right) \tag{4}
\end{equation*}
$$

where $T$ is a teaching factor that can have a value of either 1 or $2, x_{\text {mean }}$ is the current mean value of all solutions, and rand is a random number following
uniform distribution in $[0,1]$.
In the learner phase, a learner $x$ increases its knowledge by interacting with its classmate $x_{j}$ :

$$
x_{\text {new }}= \begin{cases}x+\operatorname{rand} \times\left(x-x_{j}\right), & \text { if } f(x)<f\left(x_{j}\right)  \tag{5}\\ x+\operatorname{rand} \times\left(x_{j}-x\right), & \text { else }\end{cases}
$$

where $f(x)$ indicates the fitness of $x$.
In the above phases, $x_{\text {new }}$ is accepted if $x_{\text {new }}$ gives a better objective than $x$.

Equations (4) and (5) cannot directly generate feasible solutions to scheduling problems because of their combinatorial feature. The TLBO is often discrete when it is used to solve scheduling problems ${ }^{[29-37]}$. In the current study, The QTLBO is also discrete.
In recent years, RL has been applied to solve a wide variety of complex problems ${ }^{[37-40,42-45]}$, and its main components are a learning agent, an environment, states, actions, and rewards. The Q-learning algorithm ${ }^{[46]}$ is the most commonly used RL algorithm. It works by learning an action-value function that expresses the cumulative reward of taking a certain action in a given state. It has been successfully integrated with metaheuristics for production scheduling ${ }^{[37-40]}$. Its simplest form is defined by

$$
\begin{align*}
& Q\left(s_{t}, a_{t}\right) \leftarrow Q\left(s_{t}, a_{t}\right)+ \\
& \alpha\left(r_{t+1}+\gamma \max _{a} Q\left(s_{t+1}, a\right)-Q\left(s_{t}, a_{t}\right)\right) \tag{6}
\end{align*}
$$

where $\alpha$ denotes the learning rate, often set to 0.1 (this setting is adopted in this study), $\gamma$ indicates the discount factor, $r_{t+1}$ is the reward by taking $a_{t}$ on $s_{t}$, and $\max Q\left(s_{t+1}, a\right)$ represents the largest $Q$ value in the Q-table at state $s_{t+1}$.
Action selection is executed based on the Q-table. Initially, all elements of the Q-table are zero, which means that the agent does not have any learning experience. Here, $\varepsilon$-greedy is often used, and this is expressed as follows: if a random number rand $<\varepsilon$, then randomly select an action $a$; otherwise, select an action $a$ that maximizes $Q$ values, that is, $a=$ $\arg \max _{a^{\prime}} Q\left(s_{t}, a^{\prime}\right)$.

## 4 QTLBO for the DTHFSP with Fuzzy Processing Time

The integration of RL and the TLBO is hardly investigated in the literature. Thus, the current paper proposed an effective way to combine an RL algorithm named Q-learning and the TLBO to solve the considered DTHFSP.

### 4.1 Coding and decoding

For the DTHFSP with fuzzy processing time, it is required to optimize scheduling and machine assignment of the first stage because factory assignment can be inferred from machine assignment ${ }^{[16]}$. Thus, a two-string representation is used as a solution to the problem.

For the DTHFSP with $n$ jobs and $W$ parallel machines at stage 1 , a solution is denoted by a machine assignment string $\left[M_{1, \theta_{1}}, M_{1, \theta_{2}}, \ldots, M_{1, \theta_{n}}\right]$ and a scheduling string $\left[q_{1}, q_{2}, \ldots, q_{n}\right]$, where machine $M_{1, \theta_{i}}$ is allocated for job $J_{i}$, and $q_{i}$ is the real number of job $J_{i}$. The second string is a random key; job permutation is obtained by sorting all its genes. Figure 2 shows the coding strings.

For example, a solution has a string $[0.2,0.3,0.45$, $0.55,0.09$ ] and a string [ $M_{1,2}, M_{1,3}, M_{1,1}, M_{1,3}, M_{1,4}$ ]. In the machine assignment string, $J_{1}$ is allocated on $M_{1,1}$, $J_{2}$ is on $M_{1,3}$, and so on, job permutation [5,1,2,3,4] is obtained by sorting all genes in the ascending order of $q_{i}$.

The decoding procedure is shown as follows. In each factory $f$, for each parallel machine $M_{1, k}$, all its allocated jobs are decided on according to the machine assignment string. Then, a permutation of these jobs is obtained by sorting these jobs in the ascending order of their $q_{i}$. Next, the first job of the permutation is initiated, all the jobs of the permutation on $M_{1, k}$ are processed sequentially, and the processing of each job on $M_{2, f}$ is completed after the processing at stage 1 is finished. On $M_{1,3}, J_{2}$ and $J_{4}$ are processed and a permutation of these jobs is presented ${ }^{[2,4]}$.

Meanwhile, in the decoding process, the beginning time of a job on a machine is computed by using the max operator. The completion time is calculated with the addition operator, and the ranking operator compares the makespan of each solution in order to decide the elite solution, among others. The initial population $P$ with $N$ solutions is randomly produced. All solutions are sorted in the ascending order of their makespan.


Fig. 2 Descriptions on coding strings.

### 4.2 Four phases of the QTLBO

A group $G$ of teachers is constructed by selecting $\beta \times N$ best solutions from population $P$, where $\beta$ is a real number. Each member of $G$ acts as a teacher. Aside from the teacher and learner phases, the teacher's and learner's self-learning are also added. Thus, a total of four phases were used.
The teacher phase is shown below. For each learner $x \in P \backslash G$, the following steps are repeated $\eta_{1}$ times. First, a teacher $x_{\text {teacher }} \in G$ is randomly selected. Second, for $x$ and $x_{\text {teacher }}$, a two-point crossover of the machine assignment string between them is executed. If the new solution $x_{\text {new }}$ is better than $x$, then $x_{\text {new }}$ substitutes for $x$; else do a two-point crossover of scheduling string between them and replace $x$ with $x_{\text {new }}$ if $x_{\text {new }}$ is better than $x$.

Next, the learner phase is described as follows. For each learner $x \in P \backslash G$, repeat the following steps $\eta_{2}$ times, randomly select $\delta \in[1, N]$ and a solution $x_{j}$, $j \in[1, \delta]$, and execute a two-point crossover between $x$ and $x_{j}$. Two cases are considered. In the first case, only the two-point crossover of the machine assignment string is done. In the second case, crossovers between $x$ and $x_{j}$ are executed as in the teacher phase, where $\eta_{i}$ is an integer. A teaching-learning ration $\mu$ is defined as $\eta_{1}: \eta_{2}$.
The two other phases are based on neighborhood searches $N S_{1}$ and $N S_{2}$, which are constructed by using six neighborhood structures. Neighborhood structure $\mathcal{N}_{1}$ produces new solutions by removing a randomly selected job $J_{i}$ from the factory with the biggest completion time to the factory with the shortest completion time. $\mathcal{N}_{2}$ is depicted as follows. First, randomly choose jobs $J_{i}, J_{j}$ from the factory $f$ with the biggest completion time, then swap $M_{1, \theta_{i}}$ and $M_{1, \theta_{j}}$ on the machine assignment string. When a job is removed from the factory with the biggest completion time in $\mathcal{N}_{1}$ or two jobs are swapped in $\mathcal{N}_{2}$, there is a high probability that the makespan will diminish.

When $J_{j}$ is selected at random from a randomly chosen factory (not factory $f$ in $\mathcal{N}_{2}$ ), $\mathcal{N}_{3}$ is thus obtained. Meanwhile, $\mathcal{N}_{4}$ is used to generate solutions by exchanging two randomly selected genes in the scheduling string. $\mathcal{N}_{5}$ is shown as follows: randomly decide $q_{i}$ and $q_{j}$, and insert $q_{i}$ into position $j-1$ in scheduling string if $j>1$. If $j=1$, insert $q_{i}$ into position $j$ in scheduling string. In addition, $\boldsymbol{N}_{6}$ is depicted as follows: randomly select $q_{i}, q_{j}$, and $q_{k}$ in the scheduling
string and set a new sequence for them. For example, $q_{1}=0.8, q_{3}=0.75$, and $q_{5}=0.78$; hence, the new sequence may be $q_{1}=0.78, q_{3}=0.8$, and $q_{5}=0.75$. Thus, $\mathcal{N}_{4}, \mathcal{N}_{5}$, and $\mathcal{N}_{6}$ act on the scheduling string. $N S_{1}$ is shown as follows. For solution $x$, let $g=1$, and produce a solution $x_{\text {new }} \in \mathcal{N}_{g}(x)$. If $x_{\text {new }}$ is better than $x$, then $x=x_{\text {new }}$; else, $g=g+1$. Next, obtain $x_{\text {new }} \in \mathcal{N}_{g}(x)$. If $x_{\text {new }}$ is better than $x$, then $x=x_{\text {new }}$; else $g=g+1$. Finally, generate $x_{\text {new }} \in \mathcal{N}_{g}(x)$ and $x=x_{\text {new }}$ if $x_{\text {new }}$ is better than $x . N S_{2}$ differs from $N S_{1}$ in that $g$ is initially set to be 4 and not 1 initially.

The teacher's self-learning phase is executed as follows. For each $x_{\text {teacher }} \in G$, execute $N S_{1} \eta_{2}$ times. The learner's self-learning phase is depicted as follows: for each learner $x$, execute one of the following two cases $\eta_{2}$ times: Case 1 -only $N S_{1}$ is used, and Case $2-N S_{1}$ and $N S_{2}$ are done sequentially.

### 4.3 Q-learning algorithm

In this study, the Q-learning algorithm is integrated with the TLBO to dynamically select some phases. To realize this goal, the population evaluation results are used to depict state $s_{t}$, and a combination of some phases is applied to describe action $a_{t}$. As a result, the action selection will lead to the dynamical selection of the algorithm structure.

Three indices are used to evaluate population $P$. First, trial* is used to describe the global best solution $x_{g b}^{t} \in P$. Initially, trial $=0$. If $x_{g b}^{t}$ is not updated on generation $t$, trial $^{*}=$ trial $^{*}+1 ;$ otherwise, trial ${ }^{*}=0$. Discrete degree $D^{t}$ on generation $t$ is defined by

$$
\begin{equation*}
D^{t}=\frac{2 \sum_{i=1}^{N-1} \sum_{j=i+1}^{N}\left|C_{\max }^{i, t}-C_{\max }^{j, t}\right|}{N(N-1)} \tag{7}
\end{equation*}
$$

where $C_{\max }^{i, t}$ indicates the makespan of solution $x_{i} \in P$ on generation $t$.

Genotype difference $G D^{t}$ of $\beta \times N$ best solutions in $P$ is shown below. The set $\Theta$ is the set of $\beta \times N$ solutions of $P$ with the smallest makespan. For each $x \in \Theta$, job permutation is first obtained by sorting all genes of scheduling string, then let $\eta_{x}^{t}=0$, for each $y \in \Theta, y \neq x$, starting with the first position of the machine assignment string of $x$ and $y$. For each position, $x$ and $y$ have different genes, then $\eta_{x}^{t}=\eta_{x}^{t}+1$; the same process is done on the job permutation of $x$ and $y$.

$$
\begin{equation*}
G D^{t}=\sum_{x \in \Theta} \eta_{x}^{t} /|\Theta| \tag{8}
\end{equation*}
$$

where $\eta_{x}^{t}$ is an integer related to $x$ on generation $t$.

As shown in Table 1, nine states are defined based on the above three indices, where $\Delta D^{t}=D^{t}-D^{t-1}$, $\Delta G D^{t}=G D^{t}-G D^{t-1}$.
Four actions are given by using some phases of the QTLBO. Action $a(1)$ is composed of the teacher and learner phases. Action $a(2)$ consists of the teacher phase, the first case of the learner's self-learning phase, and the learner phase. In the teacher phase of $a(2)$, only a two-point crossover of the machine assignment string is executed. Action $a(3)$ is made up of the teacher phase, the first case of the learner's self-learning phase, and the second case of the learner's phase. When $a(4)$ is executed, the teacher's self-learning phase, the teacher phase, and the second case of the other two phases are applied sequentially. When other actions like $a(1)$ are completed, their own phases are also executed sequentially.
Here, the reward is defined by

$$
\begin{equation*}
r_{t+1}=\frac{w_{1} \times \rho_{1}+w_{2} \times \rho_{2}+w_{3} \times \rho_{3}}{\rho_{1}+\rho_{2}+\rho_{3}} \tag{9}
\end{equation*}
$$

where $\quad \rho_{1}=\left(f\left(x_{g b}^{t}\right)-f\left(x_{g b}^{t+1}\right)\right) / f\left(x_{g b}^{t}\right), \quad \rho_{2}=\Delta D^{t+1} / D^{t}$, $\rho_{3}=\Delta G D^{t+1} / G D^{t}$, and $w_{i}$ is the real number, $w_{1}+w_{2}+w_{3}=1$.
We believe that the elite solution is more important than the other two indices used for the state, so $w_{1}>w_{2}, w_{1}>w_{3}$. We set $w_{1}=0.5, w_{2}=0.3, w_{3}=0.2$ by experiments.

An adaptive $\varepsilon$-greedy selection is proposed to decide action on generation $t$. Here $\varepsilon$ is defined as $\varepsilon=\varepsilon_{0}-\left(\varepsilon_{0}-0.01\right) \times R_{\text {cur }} / R_{\max }$, where $R_{\max }$ is the running time as a stopping condition, and $R_{\text {cur }}$ indicates the current time. In addition, 0.01 is used to guarantee $\varepsilon \geqslant 0.01$.

### 4.4 Algorithm description

The above Q-learning algorithm is the main part of the QTLBO and is used to decide its algorithm structure. This is a new application of Q-learning to

Table 1 Station set.

| State | Description |
| :---: | :---: |
| 1 | trial $^{*}=0, \Delta D^{t}>0$ |
| 2 | trial $^{*}=0, \Delta D^{t}<0$ |
| 3 | trial $^{*}>0, \Delta D^{t}=0$ |
| 4 | trial $^{*}>0, \Delta D^{t}>0, \Delta G D^{t}<0$ |
| 5 | trial $^{*}>0, \Delta D^{t}>0, \Delta G D^{t}=0$ |
| 6 | trial $^{*}>0, \Delta D^{t}>0, \Delta G D^{t}>0$ |
| 7 | trial $^{*}>0, \Delta D^{t}<0, \Delta G D^{t}<0$ |
| 8 | trial $^{*}>0, \Delta D^{t}<0, \Delta G D^{t}=0$ |
| 9 | trial $^{*}>0, \Delta D^{t}<0, \Delta G D^{t}>0$ |

metaheuristics. Yet, this feature is hardly considered in the previous TLBOs.

The QTLBO consists of the initialization, sorting, Qlearning step, and execution steps of the chosen action. The latter three steps are repeated until the stopping condition is met. The flow chart of the QTLBO is shown in Fig. 3. Meanwhile, the sorting step is used to sort all solutions of the population in the ascending order of their makespan, which decides the group of teachers.

Unlike the previous TLBO, sorting and Q-learning are applied to choose an action, and the chosen action includes some phases of the QTLBO.

## 5 Computational Experiment

Extensive experiments were conducted to evaluate the performance of the QTLBO for the DTHFSP with fuzzy processing time and SDST. All experiments were implemented by using Microsoft Visual C++ 2019 and run on an 8.0 GB RAM 2.4 GHz CPU PC.

### 5.1 Instance and comparative algorithms

One hundred instances were generated according to Cai et al. ${ }^{[16]}$ Tables 2 and 3 show the information on these instances.

Ying and Lin ${ }^{[47]}$ studied distributed HFSP with


Fig. 3 Algorithm flow chart.

Table 2 Factory information on 100 instances.

| Instance | $F$ | $m_{f}$ |
| :---: | :---: | :---: |
| $1-10$ | 2 | 3,3 |
| $11-20$ | 2 | 2,4 |
| $21-31$ | 3 | $3,3,3$ |
| $32-42$ | 3 | $2,3,4$ |
| $43-53$ | 4 | $2,2,2,2$ |
| $54-64$ | 4 | $2,2,3,3$ |
| $65-73$ | 4 | $2,3,4,5$ |
| $74-82$ | 5 | $3,3,3,3,3$ |
| $83-91$ | 5 | $2,3,3,3,4$ |
| $92-100$ | 5 | $2,3,4,5,6$ |

Table 3 Information of number of jobs on 100 instances.

| Instance | $n$ |
| :---: | :---: |
| $1,11,21,32,43,54$ | 30 |
| $2,12,22,33,44,55$ | 40 |
| $3,13,23,34,45,56,65,74,83,92$ | 50 |
| $4,14,24,35,46,57,66,75,84,93$ | 60 |
| $5,15,25,36,47,58,67,76,85,94$ | 70 |
| $6,16,26,37,48,59,68,77,86,95$ | 80 |
| $7,17,27,38,49,60,69,78,87,96$ | 90 |
| $8,18,28,39,50,61,70,79,88,97$ | 100 |
| $9,19,29,40,51,62,71,80,89,98$ | 120 |
| $10,20,30,41,52,63,72,81,90,99$ | 150 |
| $31,42,53,64,73,82,91,100$ | 200 |

multiprocessor tasks and presented a self-tuning iterated greedy (SIG) to minimize makespan. Li et al. ${ }^{[48]}$ developed an improved artificial bee colony (IABC) for distributed HFSP with SDST and unrelated parallel machines. The proposed IABC adopts the factory assignment rule, the greedy iterative strategy, and a hybrid search strategy. Li et al.[49] developed an improved brainstorm optimization algorithm for type-2 fuzzy distributed hybrid flow shop scheduling. These


Fig. 4 Results on average value and $\mathrm{S} / \mathrm{N}$ ratio on $\boldsymbol{c}_{1}(\mathrm{Min})$.
and $\mu$ are 2, the QTLBO produces smaller average $c_{1}$ (Min) and bigger $\mathrm{S} / \mathrm{N}$ ratio than the QTLBO with other levels. The suggested settings are $N=80, \gamma=0.9$, $\varepsilon_{0}=0.5, \beta=0.15$, and $\mu=5: 2$.
The above settings were obtained by $c_{1}$ (Min). We tested these settings on all instances and found that the above settings can produce better results than the other settings; thus, the above settings were adopted.
Except the stopping condition, the TLBO only has $N=80$. For three comparative algorithms, all parameters except the stopping condition were directly selected from previous works ${ }^{[45-47]}$.
We ran the QTLBO, TLBO, and three comparative algorithms 10 times on each instance. Tables 5-10 show the computational results of all algorithms, in which Max indicates the worst one of 10 elite solutions obtained in 10 runs, and $A v g$ is the average value of the 10 elite solutions. Table 11 presents the results of the paired sample t-test, where the significance level is 0.05 . Figure 5 describes the convergence curves of all algorithms on four instances, while Fig. 6 describes the best schedule of instance 1 obtained by the QTLBO.

As shown in Tables 5-10, the QTLBO produces better Min than the TLBO on all instances; moreover, there are significant differences between the Min of the QTLBO and the Min of the TLBO. Notably, the former performs better than the latter in terms of convergence. The same conclusion also can be drawn from Table 11 and Fig. 5. The QTLBO also obtains smaller Avg and Max than the TLBO on all instances. Furthermore, the QTLBO possesses significant performance advantages in terms of average results and stability. Obviously, the new strategies of the QTLBO have had a positive impact on its effective and efficient performance.

When the QTLBO is compared with the IABC, it can be found from Tables 7 and 8 that the QTLBO obtains smaller Min than the IABC in 92 of 100 instances. This means that the IABC generates better Min than the QTLBO on only 8 instances. Meanwhile, the QTLBO outperforms the IABC in terms of convergence. As shown in Tables $5,6,9$, and 10, the QTLBO produces smaller Max than the IABC on 89 instances and better Avg than the IABC on 85 instances. The QTLBO also performs better than the IABC in terms of average results and stability. It can be seen from Table 5-10 that the QTLBO has better performance than the SIG and IBSO. The SIG and IBSO cannot produce better solutions than the QTLBO on all instances. Moreover,
the Min of QTLBO is significantly smaller than that of the SIG and IBSO. Finally, the QTLBO performs better than the SIG and IBSO in relation to the other two metrics.

Table 11 and Fig. 5 present the convergence difference between the QTLBO and its comparative algorithm. The superior performance of the QTLBO mainly results from its integration with the Q -learning algorithm and four phases. These four phases are a combination of global search and neighborhood search. The Q-learning algorithm is used to dynamically select the algorithm structure with some phases. The exploration is intensified because of the dynamical selection of algorithm structure, and the possibility of falling local optima also can be diminished notably. Thus, it can be concluded that the QTLBO is an effective method for the DTHFSP with fuzzy processing time.

## 6 Conclusion

Many studies have investigated the integration of RL and metaheuristic for production scheduling. This study provides a new path to integrate RL with the TLBO. Unlike existing works, this study applies an RL algorithm named Q-learning to dynamically adjust the algorithm structure of the QTLBO. In the QTLBO, a group of teachers was used. The teacher phase, learner phase, teacher's self-learning phase, and learner's selflearning phase were designed. Then, the Q-learning algorithm was implemented by 9 states, 4 actions defined as combinations of the above phases, a reward, and an adaptive action selection, after which they are applied to dynamically adjust algorithm structure. A number of experiments were conducted. The computational results demonstrate that the new strategies of the QTLBO are effective and that the QTLBO provides promising results on the considered DTHFSP.

Distributed scheduling has also been extensively considered; however, distributed scheduling with uncertainty is not studied fully. Thus, in future works, we will attempt to solve the distributed scheduling problem with uncertainty by using various metaheuristics. Previous works have also mainly used many kinds of RL algorithms, particularly Q-learning. Related to this, the integration of other RL algorithms with metaheuristics for production scheduling can also be the subject of our future research.

Table 5 Computational results of all algorithms on Avg for instances 1-50.

| ance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | (1021.2,1131.6,1246.4) | (1062.7,1174.5,1347.9) | .6) | (1100,1202.8,1354.1) | ) |
| 2 | 7) | (1445.5,1592.3,1854.1) | (1382.8,1534.4,1719.9) | (1459.3,1603.6 | 4) |
| 3 | (1653.5,1844.4,2084.6) | (1748. | (1653.3,1845.9,2086.6) | (1774.6,1961.5,223 | .8) |
| 4 | (1995.9, | (2100.2,2313.9,2692.9) | (1993.5,2193.8,2486.2) | (2124.2,2330.3,264 | 2077.9,2293.3,2574.1) |
| 5 | (2325.4,2556,2871.1) | (243 | (23 | (2 | 1) |
| 6 | (2645.5,2939.6,3289.3) | (2774.8,3081,3544.1) | (26 | (2786.6,3090.8,3510.9) | 9) |
| 7 | (2) | (3133.7,3476.6 | (31 | (31 | 4) |
| 8 | (3308.8, | 4) | (3509.1,3876.4,4393.6) | (35 | 4) |
| 9 | (403 | (4184.9,4609.2,5298.9) | (4212.2,4632.6,5240.7) | (4226.1,4649.6,5304.3) | 7) |
| 10 | (5004.6, | (5209.7,5757.3,6583.6) | (5249.7,5770.7,6557.6) | (524 | (5215,5749.4,6522.1) |
| 11 | (1044.7,1142.7, | (1092.3,1198.1,1412.4) | (1045.2,1146.6,1284.1) | $(1102,1199,1347)$ | (1068.3,1172.6,1316.6) |
| 12 | (13 | (14 | (1345.6,1503.8,1699.8) | (1445.6,1588,1804.2) | (1394.4,1547.3,1755.4) |
| 13 | (1676.1,1846.6,20 | (17 | (1652.7,1842.3,2042.4) | (1771.8,1955.2,2192.6) | (1718.7,1910.4,2138.4) |
| 14 | (1 | (21 | (2008.2,2228.4,2509.4) | (2 | (2097.9,2320.6,2619.2) |
| 15 | (2330.4,2585.4,2911.4) | (2490.2,2749.8,3146.4) | (2359.9,2594,2896.4) | (2479.7,2736,3080.6) | (2451.5,2698.4,3008.8) |
| 16 | (2623.9,2895.1,3243. | (2831.9,3110,3546.2) | (2667.3,2959.3,3281.4) | (2809.7,3097.3,3494.7) | (2762.2,3045.1,3411.4) |
| 17 | (2985.5,3290.5,3719. | (3170.5,3502,4036) | (31 | (3161.5,3493.4,3959.9) | 1) |
| 18 | (3300.4,3665.5,4083. | (3527.3,3882.1,4437.7) | (35 | (3534. | $(3489.3,3838.5,4340.7)$ |
| 19 | (3951,4365.5,4889.9) | (4229.5,4653.7,5304.4) | $(4193.2,4618.3,5198.7)$ | (4197.4,4639.7,5224 | (4179.7,4597.2,5148) |
| 20 | (4976.6,5494.2,6103. | (5266.5,5811.4,6650 | (52 | (5225.2,5773.7,6559 | (5197.7,5740.3,6456) |
| 21 | 96.5,769.4,86 | (749.6,824.9,984.3) | (710.3,784.1,851.8) | 6.6,951.9) | .7) |
| 22 | (924,1004.7,1 | (978.1,1076.5,1271.8) | (913.1,1015.7,1113.8) | 32.2) | (956.3,1059.4,1168.2) |
| 23 | (1118.1,12 | (1207.6,1342.1,1599) | (1119.6,1262.4,1436.4) | (1223.6,1360.9,1548) | (1186.1,1320.3,1502.5) |
| 24 | (1336.3 | (1450.9,1 | (1355.8,1507,1648.8) | (14 | (1434.6,1575.9,1756.6) |
| 25 | (155 | (1658.8, | (553,1728.9,1931.4) | (1678,1857,2115.1) | (1637.1,1818.8,2067.6) |
| 26 | (1757.3, | (1922.9, | 6.6) | (1932.2,2136.7,2433.6) | (1883.6,2074,2361.9) |
| 27 | (1967.1 | (2147 | (2141.3,2357.5,2670) | (2160.9,2382,2693.5) | (2112.5,2336.5,2620.6) |
| 28 | (2202.2, | (2353 | (2357.3,2605.5,2972.5) | (2396 | (2327.9,2581,2925.4) |
| 29 | (260 | (2817 | (2828.2,3116.1,3498.9) | (2842.3,3138.5,3534.2) | 7.7) |
| 30 | (3344.7 | (3501.4,3861.8,4418.4) | (3525.6,3889.7,4379.2) | (3542.6,3906.9,4421.9) | $(3497.5,3845.9,4335.3)$ |
| 31 | (4421.4,4925.6,5556. | $(4651.5,5161,5912.7)$ | (4 | (4735.9,5239.4,5930.1) | $(4665.8,5153.8,5820.9)$ |
| 32 | 03.9, | (757.6,837.1,1026) | (696.4,775.8,876.2) | 67.8,842.6,961.2) | ) |
| 33 | (914.4,1024,1136 | (966.7 | (920.5,1033.1,1162.6) | (1004.5,1113.4,1252.5) | (965,1068.9,1205.8) |
| 34 | (1118.9,1251.5,141 | (1201.1,1335.9,1579.3) | .6,1348) | (1227.4,1350.1,1513.8) | 175.4,1302.7,1442.5) |
| 35 | (1326.1,1463.2,1638 | .6, | 7) | (1427,1579.7,1792.7) | (1398.2,1558.8,1747.9) |
| 36 | (1560.7,1739,195 | 3) | 5.6,1909.2) | (1691.1,1862 | (1) |
| 37 | (1805.2,1981.8,222 | (1918.4,2112,2443.4) | 786.2,1995.7,2236) | (1922.9,2128.9,2410.3) | (859.2,2067.9,2333.6) |
| 38 | (1992.5,2209.4,2463.3) | (2160.8,2384.4,2750.1) | (2121.4,2332.5,2621.3) | (2160.6,2367.9,2668.1) | (2089.6,2306.8,2596.4) |
| 39 | (2259.4,2497.9,2763.8) | (2385.7,2626.4,3038.8) | (2364.2,2614.3,2940.4) | (2376.8,2623.8,2943 | 2328.9,2575.4,2904) |
| 40 | (2665.8,2943.1,3291.4) | (2863.2,3160.6,3628.6) | (2843.5,3133.7,3492.7) | (2871.9,3153.7,3554 | (2794.7,3068.3,3482.7) |
| 41 | (3364.6,3728.9,4112.2) | (3558.4,3898.1,4480.2) | (3495.6,3865.8,4375.4) | (3524.9,3912.8,4394.) | $(3468.8,3840.8,4315.1)$ |
| 42 | (4449.2,4924.2,5502.) | (4711.5,5188.1,5970.5) | 700,5181.8,5856.2) | (4723.1,5199.1,5860.9) | (4673.7,5148.5,5781.9) |
| 43 | (540,606.5,6 | 88.4,652.4,7 | (540.5,608.5,687.5) | (603.6,671.3,764) | (569.2,629.2,715.3) |
| 44 | (696.8,761.6,860.3) | (749.2,830.9,1016.3) | (696.4,770.4,865.2) | (788.1,858.3,979.2) | (744,816.7,923.8) |
| 45 | (855.9,953.3,1072.3) | (937.7,1034.7,1220.8) | (872,959.5,1058.7) | (972.2,1065.5,1204.9) | (919.4,1014.6,1133.5) |
| 46 | (1021.2,1132.7,1276.4) | $(1101,1213.2,1442)$ | (1003,1111.7,1242.6) | (1111.4,1223.9,1384.8) | (1080.3,1192.4,1342.1) |
| 47 | (1180.9,1309.1,1449.2) | $(1278.7,1422.5,1648.6)$ | (1192.9,1332.5,1499.2) | (1304,1438.6,1618.6) | (1253.5,1379.7,1560.5) |
| 48 | (1345.7,1488.5,1664.4) | $(1449.2,1589.5,1837.5)$ | (1333.9,1485.2,1660) | (1476,1623.5,1834.5) | (1417.6,1566.6,1768.3) |
| 49 | (1495.6,1664.4,1866.7) | $(1625,1793.7,2094)$ | (1637.5,1813.6,2034.8) | (1661.6,1839.1,2079.1) | (1600.7,1774.8,2006.6) |
| 50 | (1670.3,1852.3,206 | 790.4,1980.2,2291.7) | (1796.4,1988.8,2247.4) | (1818.5,2018.2,22 | ) |

Table 6 Computational results of all algorithms on Avg for instances 51-100.

| Instance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 51 | (2043.6,2244.6,2506.5) | (2148.2,2378.4,2756.6) | (2140,2372.5,2677.2) | 8) | (2 |
| 52 | (2505.9,2785.5,3108.2) | (2667.2,2951.6,3404.8) | (2688.9,2960.2,3315.9) | ,2993.3,3365.1) | (2639.8,2921.6,3246.1) |
| 53 | (3352.4,3708,4154.6) | (3546.5,3918.7,4501.9) | (3551.1,3919.3,4442.4) | (3575.3,3945.9,4488.6) | (3497.4,3869,4331.5) |
| 54 | (541.8 | 02.3, | 612.3,662.4) | (604,667.4,739.3) | ) $3,629.2,696.1$ ) |
| 55 | (695.545,764.909,855) | (745.6,822.5,996.4) | (695.9,762.8,859) | 7,981.3) | .3) |
| 56 | (865.1,960,1052.1) | 3) | (859,953,1048.9) | 1056.4,1198.3) | ,1009.1,1127.3) |
| 57 | (1018.1,1121.2,1267.4) | $(1115.3,1222.5,1433.4)$ | 000.1,1120,12 | (121,1229.8,1387) | $(1078,1190.3,1340.6)$ |
| 58 | 7) | (274.3,1414.8,1666.2) | 8) | 2.9) | (1256.8,1394.6,1549.9) |
| 59 | (1350.1,1485.4,1662.4) | (1446.9,1606.2,1861.2) | (1351.3,1490,1656.9) | 634.3,1826.6) | 8) |
| 60 | 4.7,1663.8,1863) | (616.7,1784.6,2093.3) | (1627,1793.9,2033.6) | (650.2,1813.1,2056.7) | $(1591.5,1761.9,1989.7)$ |
| 61 | (1679,1857.7,2070.6) | (1804.9,1992.3,2315.8) | (1794.6,1977.4,2236.4) | (1843.7,2028.8,2306.7) | (1762.3,1942,2194.9) |
| 62 | (2022.5,2248.4,2522.4) | (2138.7,2375.8,2789.2) | (2160.6,2385.2,2675.9) | (2189.9,2425,2734.8) | (2113.6,2339.2,2650.6) |
| 63 | (2523.6,2789.4,3113) | (2674.1,2949.4,3408.8) | (2694.1,2970.6,3328.7) | (2695.8,2979.9,3380.2) | $(2645.7,2920.5,3292.4)$ |
| 64 | (3374.7,3727.6,4166.5) | (3566.3,3926.4,4521.6) | (3573.8,3946.7,4464.4) | (3645.6,4025.7,4561.6) | (3528.8,3893.3,4427.5) |
| 65 | (855.6,954.3,1052.8) | (933.9,1025.9,1233.8) | (867.8,958.4,1037.3) | (963,1063.1,1190.4) | (911.2,1004.4,1102.7) |
| 66 | (1012.8,1118.2,1266.2) | (1110.4,1226.6,1462.3) | (1004.1,1115.8,1237 | (1117.4,1233.8,1386.8) | (1079.2,1190.9,1342) |
| 67 | (1175.4,1301.1 | (1294.4,1435.1,1669.5) | (1169.9,1312.2 | (1300.7,1425.4,1609.2) | (1246.7,1376.8,1563.7) |
| 68 | (1337.9,1482.4, | (1444.8,1592.4,1855.6) | ( $1380.8,1525.8$ | 42.5,1869.9) | $(1432.5,1581.9,1778)$ |
| 69 | (1510.2,1672.8,1 | (1638.6,1805.6,2100.1) | (1638.4,1793.2,2023. | (1659,1828.6,2070.8) | (1601.5,1768.3,1989.3) |
| 70 | (1661.5,1844.7,205 | (1804.4,1996.4,2333.9) | (1792.1,1982.6 | (1835.7,2039.8,2310.2) | $(1767,1952.2,2205.8)$ |
| 71 | (2026.6,2251.8,2540) | (2167.4,2394.4,2777.6) | (2148.5,2377.7, | (2187.2,2418.1,2722.6) | (2111.4,2331.9,2632.5) |
| 72 | (2505.9,2774.9,3092. | 685.1,2963.7,3420) | (2684.2,2962.6, | (2711.3,2996.7,3413.6) | (2634.5,2913.4,3307.9) |
| 73 | (3331.7,3695.2,4132.4) | (3571,3941.8,4565.4) | (3545.5, | (3577.1,3950.7,4488.3) | (3523.2,3888.4,4370.7) |
| 74 | (683.3,765.8,8 | 6.2,832.3,1006.1) | (678.6,758.2,847.6) | 3) | (737.1,810.4,918.8) |
| 75 | (815.2,908.8,1003.3) | (907.7,996.6,1196.3) | (814.3,905.9,1028.4) | (930.7,1023.2,1164) | (876.3,967.6,1100.1) |
| 76 | (951.3,1054.9,1191.7) | (1042.4,1149.1,1364.2) | (947.5,1049.6,1179.7) | (1069.4,1185.4,1333.6) | (1012.8,1124.1,1268.5) |
| 77 | (1075.2, | 83.3,1305,1576.7) | (1132.9,1252.8,1379.2) | (1227.3,1343.6,1527) | (1150.8,1269.9,1430.8) |
| 78 | (1214.5,1348.5, | (1315.4,1458.1,1711) | (1323.8,1457.6,1644.7) | (1349.2,1481.8,1686.5) | $(1289.3,1427.1,1613.4)$ |
| 79 | (1350.6, | $(1470.1,1617.2,1886.9)$ | 2.2) | (1478.4,1633.2,1864.9) | $(1436.6,1588.3,1786.9)$ |
| 80 | (1608.1 | 737.6,1908.1,2 | 753,1933.1,2195) | (1809.9,1998.1,2268.3) | $(1716.1,1890.9,2141.7)$ |
| 81 | (2031.6,2245.6,25 | (2162,2375.1,2745.5) | (2182.1,2405.8,2707.5) | (2240.4,2461.1,2760.1) | $(2141.7,2362.2,2640.8)$ |
| 82 | (2707.4,2995.4,3368.4) | (2845.1,3145.1,3613.5) | (2950.2,3245.5,3664 | 915.5,3214,3643.2) | (2856.1,3134,3526.4) |
| 83 | (687.7,765.8,862.2) | 8.3,850.9 | 8.2,762.4,840.4) | (805.4,890.5,1002) | 742.1,814.1,916.4) |
| 84 | (815.4,904.7,1003.5) | (906.2,995.6,1193.7) | (813.1,907.2,1009.2) | 4.4,1047.8,1193.8) | (884.5,970.7,1096) |
| 85 | (946.7,1049.9,1174.1) | (1039.1,1156.4,1398.7) | (13.6,1043.1,1159.8) | (1066.6,1184.1,1352) | (1010.8,1117.4,1265.1) |
| 86 | (1079.1,1197.4,1327.1) | $(1178.9,1309.5,1549.4)$ | (1131.9,1251.7,1407) | (1200.9,1324.3,1484.8) | (1150.8,1276.8,1436.1) |
| 87 | (1211,1346.8,1501.8) | (1329.6,1465,1709.5) | $(1311.9,1450.7,1642.9)$ | (1346.3,1477.1,1678.6) | (1282.9,1424.3,1591.9) |
| 88 | (1340.3,1503.6,1686.1) | $(1462.8,1614.3,1902.9)$ | (1473.3,1628.5,1840.8) | (1491.2,1649.4,1861.4) | 436.7,1595.1,1797.7) |
| 89 | (1605.5,1788.7,1990.1) | 750.6,1926,2250.8) | (1769.7,1949.8,2204.4) | (1809.9,1991.4,2252.9) | (1718.9,1896.1,2129.9) |
| 90 | (2017,2237.2,2486.3) | (2174,2387.5,2778.2) | (2180.2,2394.1,2707.7) | (2254.9,2472,2803) | (2132.3,2351.7,2645.6) |
| 91 | (2705,3000.1,3371.3) | (2866.5,3167.1,3641.6) | (2914,3205.1,3631.8) | (2947.8,3253.1,3703.9) | (2832.7,3127.1,3548.4) |
| 92 | (690.2,762.9,8 | 766.7,840,1033.5) | (678.7,752.3,842.4) | (799.9,883,996.1) | (733,806.3,915.4) |
| 93 | (815.2,903.3,1005.9) | (905,1001.1,1184.6) | (823.6,910.8,999.4) | (933.9,1028.9,1168.4) | (886.5,973,1092.6) |
| 94 | (950.4,1054,1167.2) | (1052.1,1158.5,1379.7) | (940.6,1035,1159.4) | (1065.3,1165.4,1322.2) | (1007.3,1112.5,1253.2) |
| 95 | (1086.2,1196.7,1336.8) | (1199.8,1323.3,1554.5) | (1151.9,1275.1,1430.8) | (1234.3,1363.7,1551.7) | (1152.6,1274.4,1438.3) |
| 96 | (1229.2,1359.8,1502.8) | (1326.5,1465,1743.5) | (1321.3,1454.8,1641.4) | (1372.1,1511.1,1715.1) | (1295.8,1422.1,1608.8) |
| 97 | (1347.7,1490.4,1642.8) | (1459.8,1620.2,1921.7) | (1458.7,1606.6,1831.5) | (1501.5,1660.1,1866.2) | (1433.6,1584.4,1779.3) |
| 98 | (1611.9,1794,2020.7) | (1760.8,1941.1,2264.9) | (1767.4,1941.6,2228.1) | (1804.3,1991.4,2281.7) | (1724,1898.1,2160.1) |
| 99 | (2028.8,2236.5,2494.1) | (2195.9,2419.3,2824.3) | (2154.8,2390.7,2701.3) | (2174.6,2419.4,2733.8) | (2112.3,2346.1,2645.2) |
| 00 | (2722.2,3012.2,3359.3) | (2952,3243.5,3742.6) | (2935.4,3231.7,3685.9) | (2917.4,3209.6,3650.3) | (2828.7,3119,3512.4) |

Table 7 Computational results of all algorithms on Min for instances 1-50.

| Instance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $(1019,1132,1245)$ | $(1053,1164,1324)$ | $(1020,1131,1248)$ | (1079,1187,1344) | $(1045,1153,1289)$ |
| 2 | $(1380,1531,1712)$ | $(1432,1585,1844)$ | $(1379,1536,1713)$ | $(1451,1583,1780)$ | $(1431,1565,1768)$ |
| 3 | $(1651,1836,2091)$ | $(1745,1917,2252)$ | $(1649,1845,2079)$ | (1749,1934,2182) | $(1711,1908,2165)$ |
| 4 | (1997,2187,2497) | $(2099,2313,2639)$ | $(1980,2209,2453)$ | (2094,2307,2626) | $(2077,2295,2538)$ |
| 5 | (2322,2542,2885) | $(2415,2651,3096)$ | $(2328,2545,2876)$ | $(2413,2659,3012)$ | $(2404,2641,2960)$ |
| 6 | $(2639,2926,3299)$ | (2749,3058,3544) | (2667,2967,3290) | (2747,3047,3444) | $(2753,3041,3433)$ |
| 7 | (2990,3313,3663) | (3113,3474,3906) | (3124,3464,3899) | $(3070,3409,3810)$ | $(3090,3445,3845)$ |
| 8 | (3287,3686,4133) | $(3478,3843,4427)$ | $(3514,3887,4329)$ | $(3469,3821,4339)$ | (3460,3842,4282) |
| 9 | $(4025,4451,4940)$ | $(4152,4567,5289)$ | $(4173,4595,5239)$ | $(4162,4599,5155)$ | (4172,4600,5182) |
| 10 | $(4995,5499,6176)$ | (5163,5740,6518) | $(5298,5761,6465)$ | $(5215,5768,6525)$ | $(5162,5730,6548)$ |
| 11 | $(1038,1140,1284)$ | $(1074,1190,1406)$ | $(1039,1146,1281)$ | $(1075,1184,1323)$ | $(1050,1162,1304)$ |
| 12 | $(1353,1483,1684)$ | $(1401,1565,1801)$ | $(1345,1499,1699)$ | $(1407,1561,1740)$ | $(1393,1548,1728)$ |
| 13 | $(1662,1840,2087)$ | $(1715,1900,2198)$ | $(1665,1847,2007)$ | $(1756,1942,2109)$ | $(1727,1902,2087)$ |
| 14 | $(1988,2224,2479)$ | $(2121,2343,2704)$ | (2010,2225,2495) | $(2115,2339,2620)$ | (2082,2314,2602) |
| 15 | $(2336,2586,2894)$ | $(2436,2714,3164)$ | (2363,2587,2892) | $(2462,2699,2985)$ | $(2452,2687,2998)$ |
| 16 | (2620,2887,3247) | (2776,3067,3540) | (2658,2964,3250) | $(2784,3085,3436)$ | (2742,3051,3357) |
| 17 | (2947,3277,3752) | $(3159,3475,4003)$ | $(3132,3469,3914)$ | (3075,3396,3911) | (3084,3424,3940) |
| 18 | $(3300,3642,4069)$ | $(3493,3856,4409)$ | (3493,3840,4396) | $(3513,3865,4335)$ | (3474,3826,4325) |
| 19 | $(3928,4340,4858)$ | $(4192,4588,5293)$ | $(4195,4599,5200)$ | $(4151,4605,5162)$ | $(4157,4590,5109)$ |
| 20 | $(4938,5446,6106)$ | $(5194,5735,6570)$ | $(5196,5729,6507)$ | (5170,5749,6532) | (5162,5724,6444) |
| 21 | (700,764,863) | $(746,810,964)$ | (709,782,848) | $(751,829,926)$ | $(718,805,907)$ |
| 22 | $(918,1002,1119)$ | $(967,1062,1274)$ | $(904,1006,1128)$ | $(985,1080,1200)$ | $(953,1056,1156)$ |
| 23 | $(1116,1248,1402)$ | $(1185,1324,1573)$ | $(1115,1258,1423)$ | $(1175,1331,1516)$ | $(1188,1323,1478)$ |
| 24 | (1340,1494,1627) | $(1445,1567,1848)$ | $(1359,1502,1634)$ | $(1436,1570,1783)$ | $(1416,1577,1737)$ |
| 25 | $(1550,1712,1903)$ | $(1661,1819,2124)$ | $(1532,1712,1961)$ | $(1643,1839,2084)$ | $(1644,1819,2036)$ |
| 26 | $(1756,1936,2189)$ | (1903,2107,2425) | $(1797,1991,2295)$ | $(1907,2091,2398)$ | $(1859,2053,2385)$ |
| 27 | $(1958,2183,2499)$ | $(2110,2325,2723)$ | $(2106,2342,2708)$ | $(2112,2328,2631)$ | (2099,2334,2602) |
| 28 | (2201,2444,2739) | (2309,2577,3003) | $(2366,2589,2916)$ | $(2370,2599,2951)$ | $(2332,2581,2888)$ |
| 29 | $(2663,2951,3266)$ | (2752,3054,3525) | (2835,3114,3472) | $(2825,3125,3459)$ | $(2805,3088,3402)$ |
| 30 | $(3373,3700,4070)$ | $(3492,3829,4365)$ | $(3529,3876,4347)$ | $(3519,3896,4417)$ | $(3480,3849,4315)$ |
| 31 | $(4398,4917,5542)$ | $(4612,5120,5892)$ | $(4699,5163,5823)$ | $(4699,5199,5852)$ | $(4668,5140,5777)$ |
| 32 | $(699,776,873)$ | $(739,831,1026)$ | $(696,771,875)$ | $(743,824,919)$ | $(726,806,904)$ |
| 33 | $(906,1018,1150)$ | $(965,1072,1233)$ | $(915,1030,1154)$ | $(991,1097,1223)$ | $(971,1066,1184)$ |
| 34 | $(1092,1251,1432)$ | $(1159,1321,1588)$ | $(1103,1233,1348)$ | $(1214,1332,1455)$ | $(1178,1296,1424)$ |
| 35 | $(1321,1461,1624)$ | $(1459,1604,1873)$ | $(1313,1461,1684)$ | $(1396,1559,1764)$ | $(1387,1547,1756)$ |
| 36 | $(1555,1732,1956)$ | $(1671,1842,2109)$ | $(1541,1710,1905)$ | $(1657,1848,2081)$ | $(1617,1801,2069)$ |
| 37 | (1818,1980,2202) | $(1896,2098,2419)$ | (1784,1994,2211) | $(1885,2084,2392)$ | (1843,2052,2342) |
| 38 | (1956,2194,2500) | $(2139,2358,2765)$ | $(2123,2318,2590)$ | $(2118,2339,2632)$ | $(2069,2297,2568)$ |
| 39 | $(2246,2506,2736)$ | $(2368,2588,3030)$ | $(2352,2605,2915)$ | $(2335,2579,2901)$ | (2329,2574,2850) |
| 40 | $(2655,2950,3268)$ | (2865,3131,3566) | (2850,3123,3428) | (2807,3108,3461) | (2771,3052,3449) |
| 41 | $(3355,3735,4088)$ | (3546,3892,4420) | $(3493,3843,4343)$ | $(3471,3865,4370)$ | $(3471,3833,4296)$ |
| 42 | $(4449,4910,5493)$ | $(4680,5167,5914)$ | $(4658,5183,5833)$ | $(4650,5144,5836)$ | $(4654,5129,5758)$ |
| 43 | $(545,603,655)$ | $(584,645,769)$ | $(546,605,678)$ | $(583,634,730)$ | $(561,622,707)$ |
| 44 | $(695,757,860)$ | $(755,827,982)$ | $(694,770,854)$ | $(775,845,942)$ | $(734,814,913)$ |
| 45 | $(855,952,1066)$ | $(927,1020,1217)$ | $(846,960,1043)$ | $(930,1029,1133)$ | $(904,1006,1131)$ |
| 46 | $(1018,1131,1273)$ | $(1085,1197,1403)$ | $(996,1103,1245)$ | $(1087,1186,1351)$ | $(1063,1198,1334)$ |
| 47 | $(1183,1290,1465)$ | (1272,1410,1625) | $(1188,1341,1469)$ | (1289,1394,1582) | (1247,1370,1553) |
| 48 | $(1336,1478,1670)$ | $(1428,1588,1798)$ | $(1338,1483,1637)$ | $(1435,1602,1790)$ | $(1400,1558,1763)$ |
| 49 | $(1491,1665,1850)$ | $(1620,1766,2073)$ | $(1628,1793,2057)$ | $(1651,1804,2034)$ | $(1593,1772,1991)$ |
| 50 | $(1685,1849,2009)$ | $(1766,1962,2264)$ | $(1789,1982,2232)$ | $(1786,1987,2201)$ | $(1738,1942,2198)$ |

Table 8 Computational results of all algorithms on Min for instances 51-100.

| Instance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 51 | (2040,2222,2494) | $(2128,2377,2719)$ | (2118,2346,2688) | (2129,2364,2641) | (2115,2305,2645) |
| 52 | $(2468,2768,3147)$ | $(2642,2901,3432)$ | $(2672,2939,3284)$ | $(2636,2908,3294)$ | $(2636,2919,3209)$ |
| 53 | $(3313,3674,4198)$ | $(3539,3905,4419)$ | $(3504,3907,4446)$ | $(3538,3898,4409)$ | (3470,3864,4270) |
| 54 | $(540,600,671)$ | $(584,654,825)$ | $(567,608,646)$ | $(597,649,699)$ | $(554,623,699)$ |
| 55 | $(690,756,863)$ | $(732,805,978)$ | $(694,763,847)$ | $(760,825,920)$ | $(732,796,919)$ |
| 56 | $(857,962,1041)$ | $(909,1009,1231)$ | $(863,941,1054)$ | $(925,1003,1179)$ | $(906,1006,1122)$ |
| 57 | $(1006,1125,1254)$ | $(1111,1214,1368)$ | $(992,1117,1229)$ | $(1097,1220,1366)$ | $(1068,1192,1312)$ |
| 58 | $(1157,1301,1463)$ | $(1279,1388,1657)$ | $(1161,1305,1450)$ | $(1259,1410,1588)$ | $(1260,1380,1553)$ |
| 59 | $(1350,1473,1660)$ | $(1423,1588,1858)$ | $(1347,1488,1636)$ | $(1471,1597,1800)$ | $(1411,1539,1762)$ |
| 60 | $(1497,1660,1856)$ | $(1601,1756,2023)$ | $(1641,1779,2015)$ | $(1606,1796,2045)$ | $(1588,1751,1963)$ |
| 61 | $(1670,1853,2066)$ | (1795,1982,2300) | $(1784,1953,2249)$ | $(1789,1974,2205)$ | $(1748,1929,2180)$ |
| 62 | (2000,2243,2512) | $(2114,2351,2784)$ | $(2143,2373,2666)$ | (2152,2374,2669) | $(2097,2329,2643)$ |
| 63 | (2480,2767,3142) | $(2674,2912,3369)$ | $(2667,2949,3361)$ | $(2655,2951,3292)$ | (2651,2934,3233) |
| 64 | $(3376,3731,4116)$ | $(3506,3897,4505)$ | $(3523,3911,4426)$ | $(3531,3920,4438)$ | $(3493,3881,4419)$ |
| 65 | $(851,938,1081)$ | $(923,1006,1192)$ | $(860,960,1021)$ | $(911,1013,1133)$ | $(914,1001,1083)$ |
| 66 | (1006,1120,1257) | $(1089,1213,1459)$ | (1000,1112,1230) | $(1108,1205,1324)$ | $(1048,1186,1338)$ |
| 67 | $(1179,1294,1454)$ | $(1274,1421,1639)$ | $(1163,1307,1438)$ | $(1260,1395,1569)$ | $(1241,1373,1560)$ |
| 68 | $(1328,1471,1656)$ | $(1420,1576,1843)$ | $(1380,1518,1681)$ | $(1443,1594,1801)$ | $(1442,1570,1743)$ |
| 69 | $(1496,1654,1895)$ | $(1615,1809,2074)$ | $(1614,1778,2007)$ | $(1635,1790,2077)$ | $(1599,1761,1981)$ |
| 70 | (1651,1834,2055) | (1754,1980,2310) | $(1770,1970,2227)$ | $(1806,2001,2270)$ | (1764,1938,2186) |
| 71 | (2017,2224,2571) | $(2105,2360,2780)$ | $(2159,2373,2624)$ | $(2120,2355,2607)$ | $(2079,2325,2614)$ |
| 72 | (2486,2762,3096) | (2632,2934,3414) | (2691,2961,3282) | $(2669,2932,3333)$ | $(2601,2888,3333)$ |
| 73 | (3295,3677,4091) | $(3566,3910,4483)$ | $(3522,3917,4423)$ | (3512,3848,4451) | $(3485,3895,4357)$ |
| 74 | $(680,759,859)$ | $(748,825,1004)$ | $(668,750,861)$ | $(751,822,944)$ | $(730,803,894)$ |
| 75 | $(806,909,999)$ | (877,980,1152) | $(806,903,1026)$ | $(897,989,1117)$ | (860,958,1120) |
| 76 | $(945,1048,1195)$ | $(1042,1141,1329)$ | $(930,1044,1169)$ | $(1038,1139,1262)$ | $(1019,1117,1257)$ |
| 77 | $(1072,1191,1309)$ | $(1155,1274,1539)$ | $(1117,1233,1372)$ | $(1195,1300,1441)$ | $(1126,1266,1420)$ |
| 78 | (1206,1340,1508) | $(1297,1431,1701)$ | $(1295,1448,1634)$ | $(1311,1435,1635)$ | $(1264,1418,1628)$ |
| 79 | $(1355,1494,1638)$ | $(1443,1593,1828)$ | $(1462,1610,1802)$ | $(1431,1578,1815)$ | $(1434,1575,1781)$ |
| 80 | (1612,1774,1996) | $(1726,1886,2197)$ | $(1734,1916,2197)$ | $(1719,1921,2191)$ | $(1716,1882,2125)$ |
| 81 | (2002,2223,2567) | $(2113,2351,2729)$ | $(2171,2405,2661)$ | $(2148,2383,2643)$ | $(2107,2362,2641)$ |
| 82 | (2705,2976,3351) | (2846,3120,3576) | (2914,3192,3548) | (2843,3173,3589) | (2854,3116,3534) |
| 83 | $(690,768,846)$ | $(761,843,998)$ | $(699,760,829)$ | $(748,840,927)$ | $(731,807,926)$ |
| 84 | $(806,902,1009)$ | $(901,983,1205)$ | $(808,908,995)$ | (889,985,1116) | (886,961,1100) |
| 85 | $(954,1043,1167)$ | $(1027,1126,1403)$ | $(956,1045,1113)$ | $(1033,1133,1288)$ | (1002,1117,1239) |
| 86 | $(1091,1196,1301)$ | $(1147,1288,1545)$ | $(1125,1243,1391)$ | $(1148,1278,1462)$ | $(1125,1261,1435)$ |
| 87 | $(1219,1348,1478)$ | $(1318,1448,1691)$ | $(1306,1448,1615)$ | $(1313,1427,1603)$ | $(1262,1409,1610)$ |
| 88 | $(1334,1497,1684)$ | $(1422,1597,1906)$ | $(1459,1614,1814)$ | $(1456,1602,1807)$ | $(1418,1603,1770)$ |
| 89 | (1617,1774,1973) | $(1719,1901,2204)$ | $(1744,1938,2189)$ | (1742,1940,2179) | $(1729,1880,2088)$ |
| 90 | (2022,2227,2468) | $(2143,2373,2713)$ | $(2183,2384,2684)$ | $(2155,2343,2672)$ | $(2120,2345,2610)$ |
| 91 | $(2713,2986,3345)$ | (2859,3142,3636) | (2864,3160,3573) | (2886,3142,3592) | (2806,3104,3563) |
| 92 | $(702,769,806)$ | $(739,809,1008)$ | $(675,751,838)$ | $(760,825,926)$ | $(733,804,902)$ |
| 93 | $(798,902,1010)$ | $(888,992,1169)$ | (820,909,987) | $(888,988,1128)$ | $(891,969,1067)$ |
| 94 | $(970,1055,1135)$ | $(1023,1136,1320)$ | $(958,1030,1140)$ | $(1031,1143,1307)$ | $(998,1099,1260)$ |
| 95 | $(1079,1197,1322)$ | $(1176,1300,1511)$ | $(1154,1270,1410)$ | $(1152,1271,1460)$ | $(1137,1275,1403)$ |
| 96 | $(1213,1355,1496)$ | $(1306,1451,1734)$ | $(1307,1444,1625)$ | $(1293,1435,1650)$ | $(1284,1418,1603)$ |
| 97 | $(1336,1486,1647)$ | $(1444,1592,1908)$ | $(1451,1598,1823)$ | $(1486,1606,1805)$ | $(1415,1580,1783)$ |
| 98 | $(1601,1791,2006)$ | $(1722,1911,2266)$ | $(1755,1916,2219)$ | $(1739,1910,2187)$ | (1714,1900,2126) |
| 99 | $(2014,2238,2476)$ | $(2151,2392,2818)$ | $(2167,2385,2651)$ | $(2153,2380,2704)$ | $(2070,2336,2653)$ |
| 100 | $(2718,3022,3288)$ | (2845,3147,3685) | (2872,3181,3621) | (2825,3130,3558) | $(2845,3105,3487)$ |

Table 9 Computational results of all algorithms on Max for instances 1-50.

| Instance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | $(1022,1133,1247)$ | $(1075,1181,1368)$ | $(1024,1136,1249)$ | $(1132,1222,1399)$ | (1057,1164,1306) |
| 2 | $(1381,1534,1726)$ | $(1449,1606,1884)$ | $(1387,1538,1721)$ | $(1484,1631,1855)$ | $(1446,1581,1779)$ |
| 3 | $(1652,1846,2102)$ | $(1760,1956,2243)$ | (1656,1854,2099) | $(1795,1971,2293)$ | $(1745,1927,2160)$ |
| 4 | (2008,2190,2498) | $(2099,2339,2732)$ | (2006,2195,2505) | $(2147,2344,2673)$ | (2081,2307,2586) |
| 5 | $(2329,2550,2894)$ | $(2443,2697,3096)$ | $(2332,2552,2886)$ | $(2475,2729,3044)$ | $(2452,2676,2971)$ |
| 6 | $(2651,2945,3334)$ | (2810,3104,3570) | (2681,2986,3348) | $(2824,3151,3659)$ | $(2764,3055,3519)$ |
| 7 | $(3008,3341,3694)$ | $(3121,3497,3998)$ | (3163,3490,3879) | $(3161,3497,3944)$ | $(3144,3458,3851)$ |
| 8 | (3307,3704,4189) | $(3514,3869,4475)$ | (3520,3885,4431) | (3579,3970,4400) | (3453,3857,4365) |
| 9 | $(4037,4446,5013)$ | $(4226,4653,5342)$ | $(4226,4635,5296)$ | $(4261,4686,5372)$ | $(4156,4612,5281)$ |
| 10 | (5036,5550,6241) | $(5255,5782,6605)$ | $(5253,5780,6607)$ | $(5269,5812,6670)$ | $(5215,5768,6559)$ |
| 11 | $(1047,1149,1271)$ | $(1109,1209,1416)$ | (1041,1154,1287) | $(1106,1210,1362)$ | $(1068,1173,1337)$ |
| 12 | $(1361,1491,1681)$ | $(1446,1595,1861)$ | $(1359,1516,1684)$ | $(1461,1607,1851)$ | $(1398,1556,1750)$ |
| 13 | $(1674,1848,2084)$ | $(1794,1989,2270)$ | $(1647,1840,2066)$ | $(1783,1985,2281)$ | $(1732,1916,2158)$ |
| 14 | $(1988,2225,2506)$ | $(2180,2393,2768)$ | (2027,2231,2510) | $(2190,2406,2750)$ | $(2087,2328,2645)$ |
| 15 | $(2344,2591,2928)$ | (2496,2784,3196) | $(2342,2603,2909)$ | (2499,2774,3161) | (2437,2710,3045) |
| 16 | (2627,2914,3259) | (2893,3150,3573) | (2675,2956,3313) | $(2852,3113,3524)$ | (2781,3066,3416) |
| 17 | (2958,3291,3782) | $(3210,3536,4045)$ | (3175,3508,3921) | $(3218,3538,3977)$ | (3129,3462,3924) |
| 18 | (3314,3686,4099) | $(3553,3914,4468)$ | (3513,3870,4438) | $(3563,3917,4428)$ | (3537,3860,4297) |
| 19 | $(3973,4376,4918)$ | $(4267,4706,5332)$ | $(4207,4635,5202)$ | $(4224,4685,5306)$ | $(4196,4601,5186)$ |
| 20 | (5000,5497,6158) | $(5320,5843,6761)$ | $(5216,5768,6508)$ | $(5236,5812,6622)$ | $(5215,5752,6469)$ |
| 21 | $(694,769,880)$ | $(734,831,1023)$ | $(713,786,853)$ | $(784,873,967)$ | $(738,810,896)$ |
| 22 | $(922,1007,1128)$ | $(986,1082,1322)$ | $(907,1015,1134)$ | $(1020,1125,1256)$ | $(963,1067,1165)$ |
| 23 | $(1123,1246,1422)$ | $(1218,1354,1616)$ | $(1130,1280,1424)$ | $(1258,1383,1590)$ | $(1184,1321,1524)$ |
| 24 | $(1332,1497,1665)$ | $(1493,1624,1807)$ | $(1347,1505,1682)$ | $(1442,1620,1894)$ | $(1441,1579,1774)$ |
| 25 | $(1557,1726,1917)$ | (1688,1864,2153) | $(1555,1728,1958)$ | $(1734,1893,2149)$ | $(1649,1819,2080)$ |
| 26 | $(1775,1962,2185)$ | (1936,2114,2517) | $(1822,2002,2330)$ | $(1957,2167,2480)$ | (1901,2098,2352) |
| 27 | $(1980,2208,2454)$ | $(2177,2360,2777)$ | (2178,2374,2642) | $(2248,2443,2763)$ | (2108,2352,2655) |
| 28 | $(2230,2461,2737)$ | $(2407,2607,3025)$ | (2370,2612,2976) | $(2419,2678,3023)$ | $(2300,2584,2967)$ |
| 29 | (2657,2960,3322) | $(2842,3118,3585)$ | (2845,3123,3512) | $(2851,3188,3586)$ | $(2790,3079,3498)$ |
| 30 | $(3340,3688,4209)$ | $(3560,3896,4436)$ | $(3526,3900,4394)$ | $(3566,3940,4498)$ | $(3513,3844,4358)$ |
| 31 | $(4465,4948,5521)$ | $(4664,5211,5950)$ | (4747,5194,5920) | $(4776,5292,5953)$ | $(4671,5163,5860)$ |
| 32 | (704,781,880) | $(775,852,1023)$ | $(683,778,901)$ | $(803,850,986)$ | $(745,810,903)$ |
| 33 | $(917,1029,1132)$ | $(990,1097,1280)$ | $(919,1039,1171)$ | $(1017,1134,1284)$ | $(970,1078,1206)$ |
| 34 | $(1139,1253,1399)$ | $(1200,1345,1603)$ | $(1115,1243,1357)$ | $(1253,1383,1543)$ | $(1184,1318,1434)$ |
| 35 | $(1337,1464,1642)$ | $(1486,1636,1921)$ | $(1324,1465,1716)$ | $(1447,1586,1848)$ | $(1401,1567,1757)$ |
| 36 | $(1579,1736,1959)$ | $(1711,1870,2196)$ | (1564,1732,1902) | $(1711,1884,2172)$ | $(1659,1826,2059)$ |
| 37 | (1811,1994,2224) | (1941,2131,2461) | $(1791,1997,2257)$ | $(1977,2181,2440)$ | (1862,2072,2351) |
| 38 | $(2007,2218,2464)$ | $(2189,2415,2767)$ | (2111,2354,2632) | $(2159,2401,2744)$ | (2106,2314,2598) |
| 39 | (2249,2490,2812) | $(2413,2647,3071)$ | $(2382,2602,2981)$ | $(2403,2658,2961)$ | $(2334,2582,2938)$ |
| 40 | (2681,2956,3278) | (2902,3178,3654) | (2858,3161,3491) | (2987,3251,3561) | (2813,3093,3462) |
| 41 | (3375,3744,4090) | (3575,3907,4512) | $(3506,3875,4419)$ | $(3542,3943,4465)$ | $(3479,3856,4334)$ |
| 42 | $(4467,4941,5533)$ | $(4743,5217,5990)$ | $(4692,5203,5884)$ | $(4709,5207,5994)$ | $(4671,5163,5809)$ |
| 43 | $(538,610,660)$ | $(598,662,799)$ | $(533,607,715)$ | $(700,761,851)$ | $(576,632,729)$ |
| 44 | $(708,766,850)$ | $(757,841,1025)$ | $(689,774,874)$ | $(793,874,1026)$ | $(750,824,926)$ |
| 45 | (849,954,1101) | $(947,1039,1235)$ | $(873,976,1056)$ | $(1089,1166,1316)$ | $(916,1011,1167)$ |
| 46 | $(1026,1136,1277)$ | $(1101,1226,1470)$ | $(1018,1115,1255)$ | $(1151,1263,1448)$ | $(1077,1196,1360)$ |
| 47 | $(1165,1316,1464)$ | $(1305,1437,1640)$ | $(1205,1337,1487)$ | $(1376,1526,1755)$ | $(1235,1386,1597)$ |
| 48 | $(1357,1501,1648)$ | $(1464,1595,1862)$ | $(1339,1478,1694)$ | $(1529,1681,1914)$ | $(1446,1577,1748)$ |
| 49 | $(1502,1673,1871)$ | $(1630,1811,2146)$ | $(1632,1813,2080)$ | $(1678,1865,2137)$ | $(1618,1777,2009)$ |
| 50 | $(1662,1859,2095)$ | $(1818,2003,2296)$ | $(1809,1988,2282)$ | $(1908,2062,2340)$ | $(1765,1948,2240)$ |

Table 10 Computational results of all algorithms on Max for instances 51-100.

| Instance | QTLBO | TLBO | IABC | SIG | IBSO |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 51 | (2034,2243,2564) | (2167,2397,2782) | $(2148,2403,2656)$ | (2307,2582,2921) | $(2125,2352,2624)$ |
| 52 | (2487,2789,3160) | $(2715,2975,3412)$ | $(2691,2983,3323)$ | (2932,3208,3631) | $(2666,2943,3230)$ |
| 53 | $(3359,3735,4174)$ | $(3583,3939,4515)$ | $(3557,3919,4496)$ | $(3635,3988,4543)$ | $(3490,3900,4367)$ |
| 54 | $(541,606,669)$ | $(617,691,815)$ | $(554,617,682)$ | (629,701,771) | $(581,625,708)$ |
| 55 | $(709,765,850)$ | $(756,840,990)$ | $(709,766,860)$ | $(835,928,1061)$ | $(740,810,929)$ |
| 56 | $(860,953,1084)$ | $(967,1053,1246)$ | (872,949,1093) | $(1000,1117,1235)$ | $(904,1017,1146)$ |
| 57 | $(1033,1127,1254)$ | $(1120,1236,1474)$ | $(1008,1128,1258)$ | $(1145,1256,1401)$ | $(1080,1185,1372)$ |
| 58 | $(1178,1310,1449)$ | (1290,1434,1673) | $(1170,1315,1471)$ | $(1392,1525,1710)$ | $(1281,1407,1532)$ |
| 59 | $(1359,1494,1659)$ | $(1480,1627,1878)$ | $(1378,1503,1665)$ | $(1534,1682,1867)$ | $(1435,1567,1772)$ |
| 60 | $(1525,1686,1830)$ | $(1628,1812,2169)$ | $(1614,1800,2058)$ | (1677,1849,2083) | $(1602,1761,2001)$ |
| 61 | $(1706,1864,2058)$ | $(1800,1999,2347)$ | $(1808,1988,2273)$ | (1982,2210,2589) | $(1753,1952,2228)$ |
| 62 | (2027,2267,2506) | (2170,2404,2785) | $(2158,2388,2702)$ | $(2353,2571,2932)$ | $(2125,2346,2658)$ |
| 63 | $(2524,2802,3134)$ | $(2709,2988,3425)$ | (2690,2987,3344) | (2710,3017,3414) | $(2639,2913,3344)$ |
| 64 | $(3428,3747,4129)$ | (3609,3972,4555) | $(3609,3982,4443)$ | $(3889,4306,4818)$ | $(3547,3903,4420)$ |
| 65 | $(856,956,1056)$ | $(945,1027,1262)$ | $(881,964,1037)$ | $(1031,1142,1240)$ | $(903,1010,1128)$ |
| 66 | $(1021,1121,1272)$ | $(1131,1230,1475)$ | $(1007,1126,1230)$ | $(1152,1253,1398)$ | $(1096,1199,1338)$ |
| 67 | $(1178,1307,1450)$ | $(1310,1456,1696)$ | $(1181,1313,1445)$ | $(1319,1453,1633)$ | $(1242,1386,1568)$ |
| 68 | $(1348,1505,1639)$ | $(1469,1608,1868)$ | $(1379,1532,1732)$ | $(1597,1778,2055)$ | $(1432,1587,1802)$ |
| 69 | $(1501,1676,1877)$ | (1641,1811,2161) | $(1645,1802,2043)$ | $(1723,1921,2182)$ | $(1598,1771,2015)$ |
| 70 | $(1656,1853,2065)$ | $(1821,2032,2384)$ | $(1800,1997,2249)$ | $(1887,2075,2369)$ | $(1778,1964,2198)$ |
| 71 | (2044,2252,2546) | (2194,2422,2787) | (2164,2392,2655) | $(2295,2535,2923)$ | $(2130,2357,2620)$ |
| 72 | (2509,2787,3095) | $(2710,2973,3440)$ | $(2691,2971,3424)$ | (2824,3133,3581) | (2637,2930,3321) |
| 73 | $(3365,3726,4152)$ | (3608,3994,4676) | $(3567,3937,4464)$ | $(3674,4045,4591)$ | $(3531,3900,4379)$ |
| 74 | $(702,769,840)$ | $(762,845,1029)$ | $(687,760,854)$ | $(778,861,1013)$ | $(738,821,920)$ |
| 75 | $(831,919,982)$ | $(923,1014,1220)$ | $(820,909,1036)$ | $(970,1076,1220)$ | $(879,985,1092)$ |
| 76 | $(946,1053,1221)$ | $(1055,1158,1402)$ | $(956,1064,1194)$ | $(1130,1255,1429)$ | $(1014,1138,1262)$ |
| 77 | $(1096,1211,1297)$ | $(1212,1318,1605)$ | $(1143,1271,1389)$ | $(1330,1445,1650)$ | $(1173,1272,1428)$ |
| 78 | $(1217,1344,1531)$ | $(1319,1460,1753)$ | $(1314,1459,1703)$ | $(1439,1588,1771)$ | $(1300,1441,1597)$ |
| 79 | $(1362,1502,1677)$ | $(1493,1632,1923)$ | $(1480,1648,1891)$ | $(1514,1662,1917)$ | $(1451,1589,1789)$ |
| 80 | $(1627,1799,2002)$ | $(1764,1947,2219)$ | $(1767,1945,2217)$ | $(1972,2181,2425)$ | $(1701,1896,2188)$ |
| 81 | $(2053,2259,2530)$ | $(2183,2392,2800)$ | $(2197,2435,2752)$ | (2384,2597,2970) | $(2154,2369,2646)$ |
| 82 | (2724,3030,3337) | (2855,3152,3667) | $(3072,3373,3849)$ | (2959,3246,3687) | (2835,3138,3571) |
| 83 | $(687,765,876)$ | $(778,864,1035)$ | $(701,768,842)$ | $(874,993,1128)$ | $(738,813,937)$ |
| 84 | $(823,911,992)$ | $(941,1017,1162)$ | $(820,912,1013)$ | $(1071,1173,1346)$ | $(876,980,1102)$ |
| 85 | $(961,1053,1169)$ | $(1050,1173,1423)$ | $(953,1045,1186)$ | $(1199,1318,1525)$ | $(1027,1112,1281)$ |
| 86 | $(1079,1194,1351)$ | $(1198,1329,1565)$ | $(1151,1280,1409)$ | $(1296,1407,1569)$ | $(1161,1292,1416)$ |
| 87 | $(1227,1343,1510)$ | $(1337,1482,1756)$ | $(1347,1483,1663)$ | $(1470,1599,1819)$ | (1294,1436,1595) |
| 88 | $(1345,1515,1687)$ | $(1457,1624,1943)$ | $(1480,1653,1906)$ | $(1555,1718,1967)$ | $(1425,1591,1836)$ |
| 89 | (1612,1794,2007) | (1780,1934,2280) | $(1804,1970,2256)$ | (1880,2039,2315) | $(1722,1896,2165)$ |
| 90 | $(2025,2248,2508)$ | $(2215,2400,2825)$ | $(2189,2407,2714)$ | (2419,2660,3018) | $(2148,2352,2659)$ |
| 91 | (2720,3002,3395) | (2855,3188,3673) | (3043,3347,3743) | (3070,3418,3860) | (2834,3132,3575) |
| 92 | (709,772,813) | $(795,853,1045)$ | $(678,754,847)$ | $(917,1009,1137)$ | $(728,805,943)$ |
| 93 | (812,914,999) | $(904,1017,1199)$ | $(815,914,1013)$ | $(993,1087,1225)$ | $(900,977,1092)$ |
| 94 | $(937,1056,1196)$ | $(1086,1172,1412)$ | $(946,1038,1177)$ | $(1098,1207,1386)$ | $(1008,1118,1266)$ |
| 95 | $(1101,1195,1347)$ | $(1235,1364,1600)$ | $(1169,1283,1426)$ | $(1392,1522,1722)$ | $(1162,1272,1466)$ |
| 96 | $(1240,1369,1507)$ | $(1354,1486,1771)$ | $(1325,1479,1667)$ | $(1490,1638,1898)$ | $(1303,1424,1626)$ |
| 97 | $(1352,1497,1653)$ | $(1473,1656,1966)$ | $(1471,1619,1842)$ | $(1575,1770,2007)$ | $(1439,1580,1808)$ |
| 98 | $(1627,1812,2004)$ | $(1798,1991,2287)$ | $(1813,1981,2244)$ | $(1931,2118,2436)$ | $(1703,1898,2220)$ |
| 99 | $(2033,2241,2502)$ | $(2229,2461,2817)$ | $(2174,2427,2690)$ | (2207,2455,2770) | $(2131,2351,2673)$ |
| 100 | (2748,3019,3402) | -305,033,403,835 | (3083,3405,3881) | (3006,3294,3791) | $(2823,3125,3570)$ |

Table 11 Results of the paired sample t-test.

| t-test | $p$-value $(A v g)$ | $p$-value $($ Max $)$ | $p$-value $($ Min $)$ |
| :---: | :---: | :---: | :---: |
| t-test(QTLBO,TLBO) | 0.0000 | 0.0000 | 0.0000 |
| t-test(QTLBO,IABC) | 0.0000 | 0.0000 | 0.0000 |
| t-test(QTLBO,SIG) | 0.0000 | 0.0000 | 0.0000 |
| t-test(QTLBO,IBSO) | 0.0000 | 0.0000 | 0.0000 |


(a) Instance 12

(c) Instance 70

(b) Instance 35

(d) Instance 90

Fig. 5 Convergence curves of all algorithms on four instances.


Fig. 6 Gantt chart of the best solution of instance 1 obtained by Q-learning-based teaching-learning based optimization (QTLBO).

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