

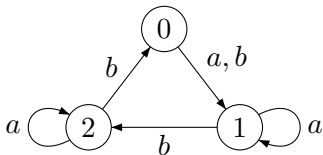
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# Reset Complexity of Ideal Languages

Marina I. Maslennikova

Ural Federal University  
Ekaterinburg, Russia

- A DFA  $\mathcal{A} = \langle Q, \Sigma, \delta \rangle$  is called *synchronizing*, if there exists a word  $w \in \Sigma^*$  which leaves the automaton in one particular state no matter which state in  $Q$  it starts at:  
$$\delta(q_1, w) = \delta(q_2, w) \text{ for all } q_1, q_2 \in Q.$$
$$|\delta(Q, w)| = 1.$$
- Any such word is said to be *synchronizing* (or *reset*) for the DFA  $\mathcal{A}$ .
- $Syn(\mathcal{A})$  – the language of all words synchronizing  $\mathcal{A}$ .



A reset word is  $abba$ : applying it at any state brings the automaton to the state 1.

## Černý conjecture

Any synchronizing  $n$ -state automaton possesses a synchronizing word of length at most  $(n - 1)^2$ .

- The *state complexity* of the language  $L$  (the number of states in the minimal DFA recognizing  $L$ ) is denoted by  $sc(L)$ .
- In what follows we consider only ideal languages:  $L = \Sigma^* L \Sigma^*$ .

## Lemma 1.

Let  $L$  be an ideal language and  $\mathcal{A}$  the minimal automaton recognizing  $L$ . Then  $\mathcal{A}$  is synchronizing and  $Syn(\mathcal{A}) = L$ .

- *Reset complexity* of an ideal language  $L$  is the minimal possible number of states in a synchronizing automaton  $\mathcal{A}$  such that  $Syn(\mathcal{A}) = L$ .

- $rc(L) \leq sc(L)$ .

## Proposition 1.

Let  $L$  be an ideal language over a unary alphabet. Then  $sc(L) = rc(L) = \ell + 1$ , where  $\ell$  is the minimum length of words in  $L$ .

- A synchronizing automaton, whose shortest reset words have length close to  $(n - 1)^2$ , is called “*slowly*” synchronizing.

## Proposition 2.

For every “slowly” synchronizing automata  $\mathcal{A}_n$  with  $n$  states,  $sc(\text{Syn}(\mathcal{A}_n)) = 2^n - n$  and  $rc(\text{Syn}(\mathcal{A}_n)) = n$ .

- For the language  $L = (a + b)^*(b^3ab^2a + a^2b^3a + abab^3a + ab^2ab^3a)(a + b)^*$  there exist two different synchronizing automata for which  $L$  serves as the language of synchronizing words.