

Scalable Ontology-Based Information Systems

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What is an Ontology?





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A model of (some aspect of) the world

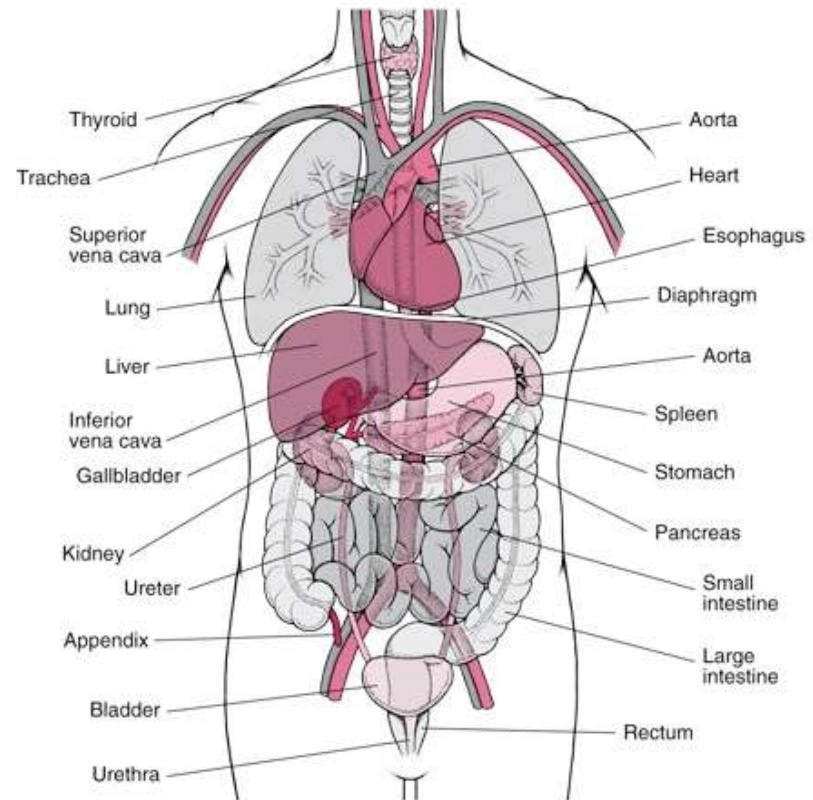




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- Introduces **vocabulary** relevant to domain, e.g.:
 - Anatomy

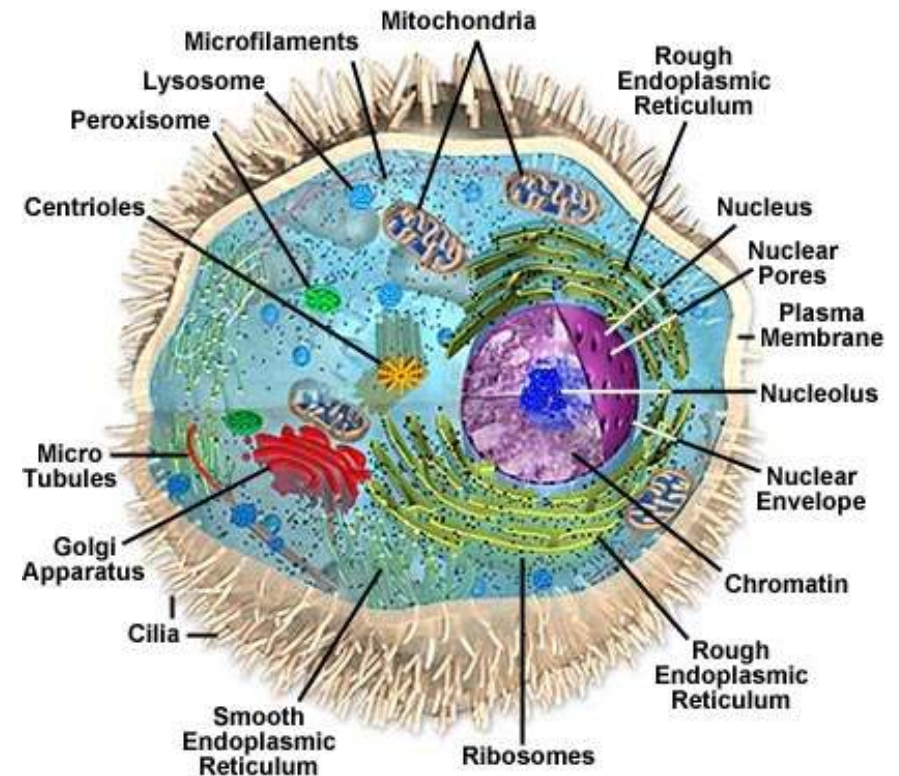




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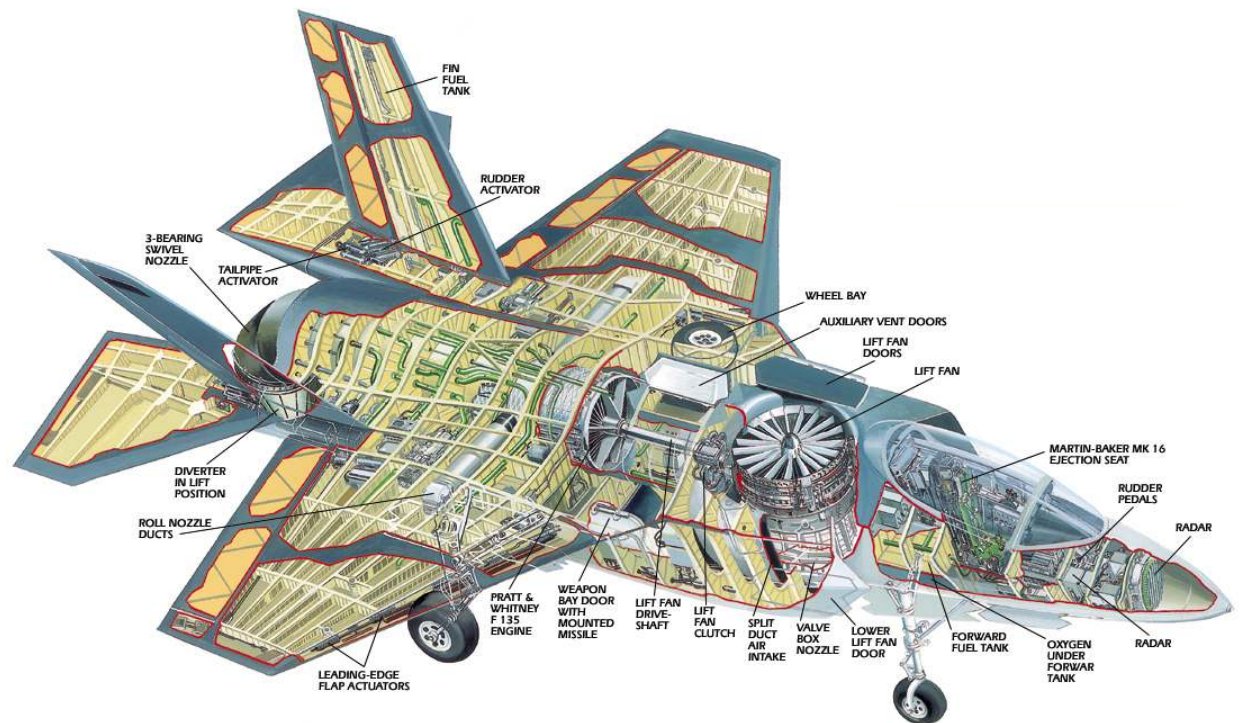




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 - Anatomy
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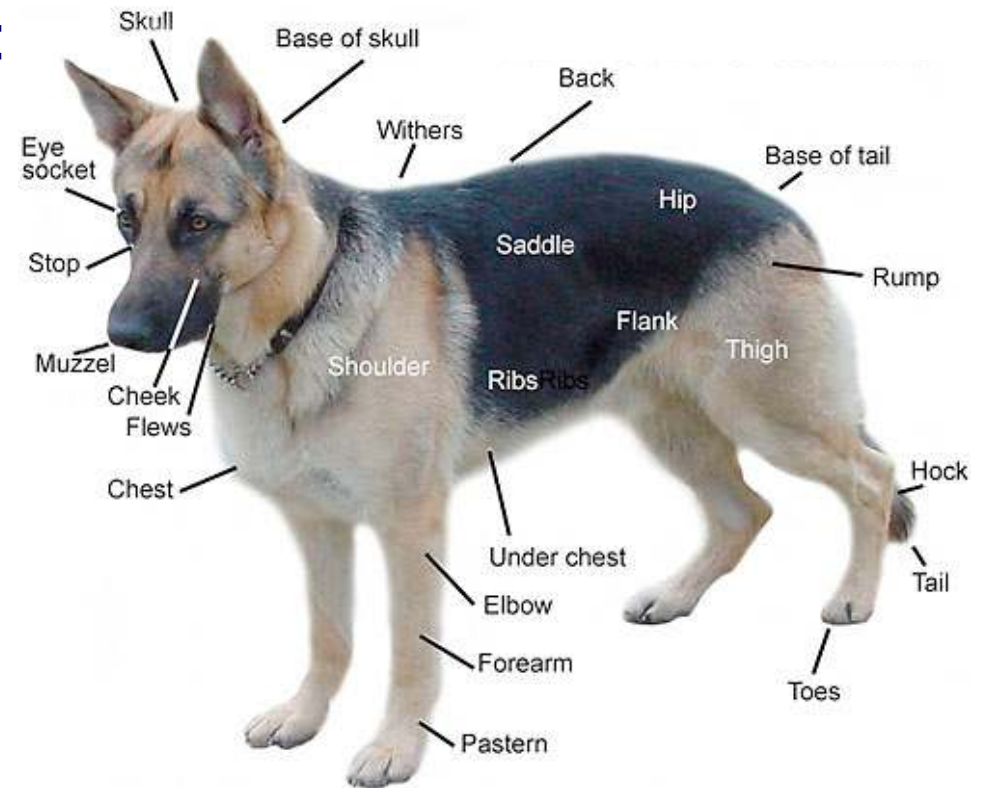




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- Anatomy
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- Dogs
- Hotdogs
- ...



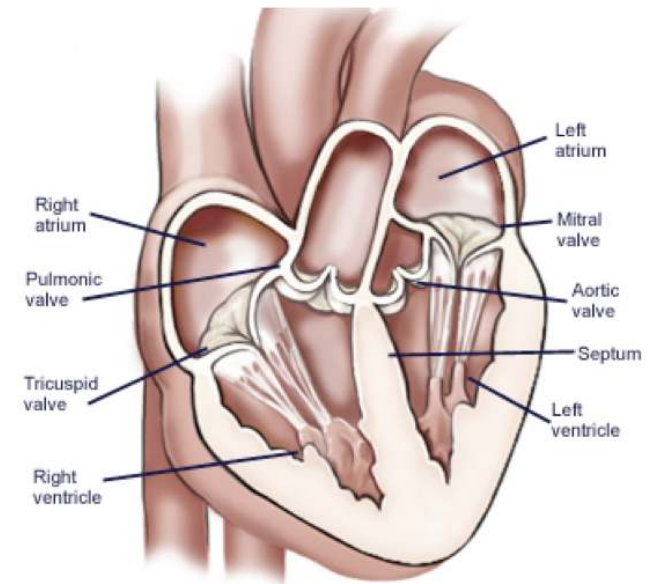


What is an Ontology?

A model of (some aspect of) the world

- Introduces **vocabulary** relevant to domain
- Specifies **meaning** (semantics) of terms

Heart **is** a muscular organ that **is part of** the circulatory system





What is an Ontology?

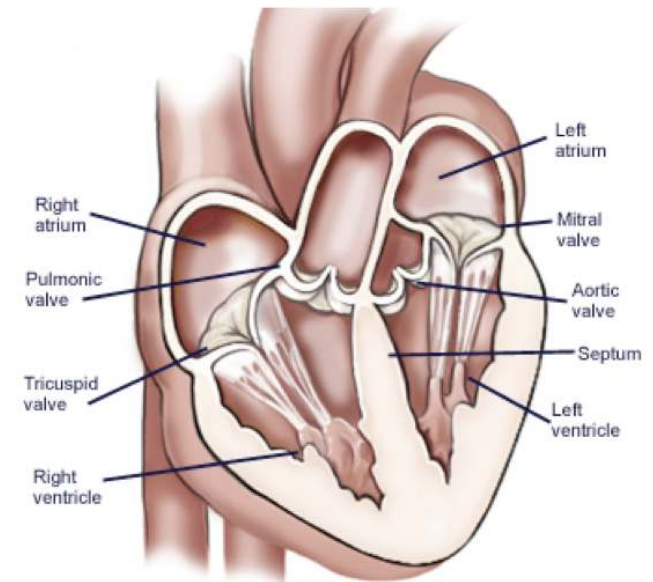
A model of (some aspect of) the world

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Heart **is a** muscular organ that **is part of** the circulatory system

- **Formalised** using suitable logic

$$\forall x. [\text{Heart}(x) \rightarrow \text{MuscularOrgan}(x) \wedge \exists y. [\text{isPartOf}(x, y) \wedge \text{CirculatorySystem}(y)]]$$





Web Ontology Language OWL (2)

- **W3C** recommendation(s)
- Motivated by **Semantic Web** activity
 - Add meaning to web content by annotating it with terms defined in ontologies
- Supported by **tools and infrastructure**
 - APIs (e.g., OWL API, Thea, OWLink)
 - Development environments (e.g., Protégé, Swoop, TopBraid Composer)
 - Reasoners & Information Systems (e.g., Pellet, Racer, HermiT, Quonto, ...)
- Based on a **Description Logics** (*SHOIN / SROIQ*)





Description Logics (DLs)

- Fragments of **first order logic** designed for KR
- Desirable computational properties
 - **Decidable** (essential)
 - Low complexity (desirable)
- Succinct and **variable free syntax**

$$\forall x. [\text{Heart}(x) \rightarrow \text{MuscularOrgan}(x) \wedge \\ \exists y. [\text{isPartOf}(x, y) \wedge \\ \text{CirculatorySystem}(y)]]$$

$$\text{Heart} \sqsubseteq \text{MuscularOrgan} \sqcap \\ \exists \text{isPartOf}. \text{CirculatorySystem}$$



Description Logics (DLs)

DL **Knowledge Base** (KB) consists of two parts:

- Ontology (aka **TBox**) axioms define terminology (schema)

$\text{Heart} \sqsubseteq \text{MuscularOrgan} \sqcap$
 $\exists \text{isPartOf}.\text{CirculatorySystem}$

$\text{HeartDisease} \equiv \text{Disease} \sqcap$
 $\exists \text{affects}.\text{Heart}$

$\text{VascularDisease} \equiv \text{Disease} \sqcap$
 $\exists \text{affects} . (\exists \text{isPartOf} . \text{CirculatorySystem})$

- Ground facts (aka **ABox**) use the terminology (data)

$\text{John} : \text{Patient} \sqcap$
 $\exists \text{suffersFrom} . \text{HeartDisease}$

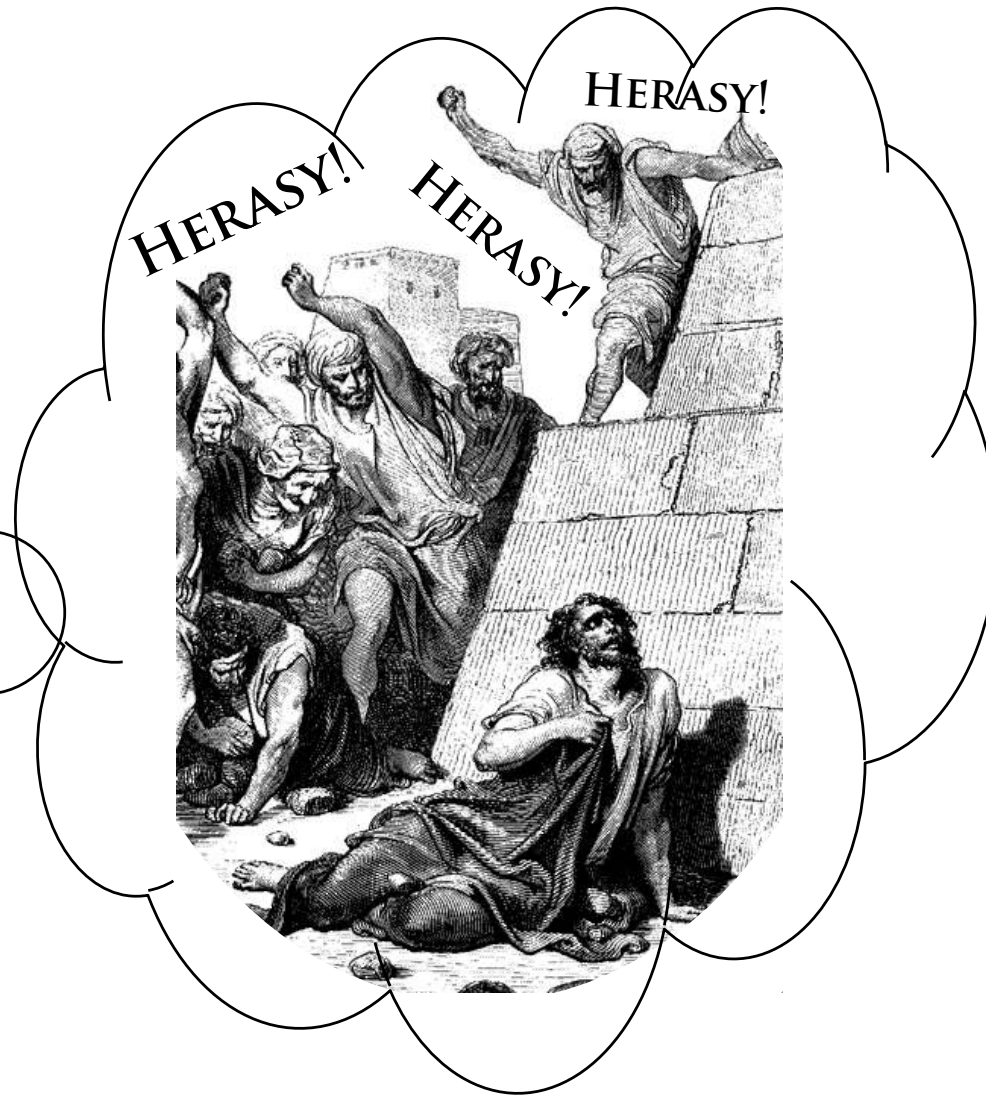


Why Care About Semantics?





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Why should I care about semantics?





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Why Care About Semantics?

Why should I care about semantics?

Well, from a philosophical POV, we need to specify the relationship between statements in the logic and the existential phenomena they describe.

That's OK, but I don't get paid for philosophy.

From a practical POV, in order to specify and test ontology-based information systems we need to precisely define their intended behaviour





Why Care About Semantics?

In FOL we define the semantics in terms of models (a model theory). A model is supposed to be an analogue of (part of) the world being modeled. FOL uses a very simple kind of model, in which “objects” in the world (not necessarily physical objects) are modeled as elements of a set, and relationships between objects are modeled as sets of tuples.





Why Care About Semantics?

In FOL we define the semantics in terms of models (a model theory). A model is supposed to be an analogue of (part of) the world being modeled. FOL uses a very simple kind of model, in which “objects” in the world (not necessarily physical objects) are modeled as elements of a set, and relationships between objects are modeled as sets of tuples.

This is exactly the same kind of model as used in a database: objects in the world are modeled as values (elements) and relationships as tables (sets of tuples).





What are Ontologies Good For?

- Coherent **user-centric view** of domain
 - Help identify and resolve disagreements
- Ontology-based **Information Systems**
 - View of data that is independent of logical/physical schema
 - Queries use terms familiar to users
 - Answers reflect schema & data, e.g.:
 - “Patients suffering from Vascular Disease”
 - Query expansion/navigation/refinement
 - Incomplete and semi-structured data
 - Integration of heterogeneous sources



Now... *that* should clear up a few things around here



Healthcare

- UK NHS **£6.2 billion** “Connecting for Health” IT programme
- Key component is **Care Records Service** (CRS)
 - “Live, interactive patient record service accessible 24/7”
 - Patient **data distributed** across local centres in 5 regional clusters, and a national DB
 - Detailed **records** held by local service providers
 - Diverse **applications** support radiology, pharmacy, etc
 - **Summaries** sent to national database
 - **SNOMED-CT** ontology provides common **vocabulary** for data
 - Clinical data uses terms drawn from ontology



SNOMED-CT

- It's **BIG** – over **400,000** concepts
- Language used is **EL** profile of **OWL 2**
- **Multiple hierarchies** and **rich definitions**





CliniClue 2006: SNOMED CT(International 0801int[Release]) [Registered user: phendler@hotmail.com]

File Edit Subsets Restrict Language Layout Tools Help

Concept Id 154283005 **TB - Pulmonary tuberculosis**

Description Id 1784750013 clinical finding

Words - any order

Find pulmonary tuber

- P pulmonary tuberculosis
- S TB - Pulmonary tuberculosis**
- P pulmonary tuberosa sclerosis
- S PTB - Pulmonary tuberculosis
- S inactive pulmonary tuberculosis

caused by Mycobacterium tuberculosis complex

Hierarchy Subtype hierarchy

- C 205237003 pneumonitis
- C 56717001 tuberculosis
- C 84353005 pulmonary disease due to Mycobacteria
 - 154283005 **pulmonary tuberculosis**
 - C 428697002 inactive tuberculosis of lung
 - C 186175002 infiltrative lung tuberculosis
 - C 186188004 isolated tracheal or bronchial tuberculosis
 - C 77688003 isolated tracheal tuberculosis
 - C 80602006 nodular tuberculosis of lung
 - C 186192006 respiratory tuberculosis, bacteriologically and histologically confirmed
 - C 186202007 respiratory tuberculosis, not confirmed bacteriologically
 - C 186177005 tuberculosis of lung with cavitation
 - C 81554001 tuberculosis of lung with involvement of bronchus
 - C 186204008 tuberculosis of lung, bacteriologically and histologically confirmed
 - C 186184007 tuberculosis of lung, confirmed by culture only
 - C 186193001 tuberculosis of lung, confirmed by sputum microscopy
 - C 186195008 tuberculosis of lung, confirmed histologically
 - C 23022004 tuberculous bronchiectasis
 - C 90117007 tuberculous fibrosis of lung

pulmonary tuberculosis - Definition
Concept Status: **Current**

Descriptions

- F pulmonary tuberculosis (disorder)
- P pulmonary tuberculosis
- S PTB - Pulmonary tuberculosis
- S TB - Pulmonary tuberculosis

Definition: Fully defined by ...

is a

- D pneumonitis
- D inflammatory disorder of lower respiratory tract
- D disorder of lung
- D inflammation of specific body organs
- D tuberculosis
- D pulmonary disease due to Mycobacteria
- D infectious disease of lung
- D bacterial lower respiratory infection
- D mycobacteriosis

causative agent

- D Mycobacterium tuberculosis complex

Group

- associated morphology
- finding site
 - D lung structure

Qualifiers

- severity
 - p severities
- episodicity
 - p episodicities
- clinical course
 - p courses

Codes

Original SnomedId : R-F46B3

Pulmonary Tuberculosis

kind of pneumonitis

kind of tuberculosis

kind of Pulmonary disease due to Mycobacteria

found in lung structure



SNOMED-CT

- Over **400,000 concepts**
- Language used is **EL fragment of OWL 2**
- **Multiple hierarchies** and **rich definitions**
- Supports, e.g., retrieving details of **all patients having pulmonary TB**





QOMCriteria - Microsoft Internet Explorer

Query Model

Save CQML | VCG/HCM Metadata | Run | Clear

Criteria Link | Export | Help

Query Details

Query Name:
Query Desc:
INI: EDG Clinical (Diagnosis)
Load CQML:

Property Based Query

Icd-9 Code	Value	300000	300010
<input type="text"/>			
<input type="text"/>			
<input type="text"/>			
<input type="text"/>			

Add more rows +

Hierarchical Based Query

SCTID	<input type="text"/>
SCTID	<input type="text"/>
SCTID	<input type="text"/>

Add more rows +

Subsumption Based Query

Add concept with +

Is A: +

Role: +

1 Total Results = 15

Exclude	CSM_ID	item 2	item 1	item 11	item 40	item 2000	
<input type="checkbox"/>	105302	TB LUNG, BRONCHIECTASIS	26224	12126224	011.50A	011.50	Ti bi (c
<input type="checkbox"/>	105303	TB LUNG, FIBROSIS	26225	12126225	011.40A	011.40	Ti fil lu (c
<input type="checkbox"/>	105825	TB LUNG, CAVITARY	26747	12126747	011.20A	011.20	Ti of ca (c
<input type="checkbox"/>	124660	TB LUNG, CULTURE DIAGNOSIS	27295	12127295	011.94A	011.94	Ti of cu ct (c
<input type="checkbox"/>	127809	MYCOBACTERIAL PNEUMONIA, NON TB	28774	12128774	500725	500725	N tu m pi (c
<input type="checkbox"/>	132112	TUBERCULOSIS (TB), PULMONARY, ACTIVE.	31139	12131139	011.90D	011.90	Pi tu (c
<input type="checkbox"/>	171105	PULMONARY MYCOBACTERIAL AVIUM	37507	12137507	031.0D	031.0	B di (c
<input type="checkbox"/>	184096	TB LUNG, CONFIRMED HISTOLOGICALLY	39390	12139390	011.95A	011.95	Ti of ca hi (c
<input type="checkbox"/>	189147	TB OF LUNG, INFILTRATIVE	39861	12139861	011.00A	011.00	Ir lu tu (c
<input type="checkbox"/>	189149	TB PNEUMONIA	39863	12139863	011.60A	011.60	Ti di





SNOMED-CT

- Over **400,000 concepts**
- Language used is **EL fragment of OWL 2**
- **Multiple hierarchies** and **rich definitions**
- Supports, e.g., retrieving details of **all patients having pulmonary TB**
 - information used e.g., to improve Quality of Care, for Reporting, in epidemiological research, in Decision Support, ...
- **Building and maintenance** is a huge task
 - supported by reasoning tools, e.g., to enrich hierarchies





What About Scalability?

- Only **useful in practice** if we can deal with large ontologies and/or large data sets
- Unfortunately, many ontology languages are highly intractable
 - Satisfiability for OWL 2 ontologies is **2NEXPTIME-complete**
- Problem addressed in practice by
 - Algorithms that work well in **typical cases**
 - Highly **optimised implementations**
 - Use of tractable fragments (aka **profiles**)





Tableau Reasoning Algorithms





Tableau Reasoning Algorithms

Standard technique based on (hyper-) **tableau**

- Reasoning tasks reducible to (un)**satisfiability**
 - E.g., $\text{KB} \models \text{HeartDisease} \sqsubseteq \text{VascularDisease}$ iff
 $\text{KB} \cup \{x:(\text{HeartDisease} \sqcap \neg \text{VascularDisease})\}$ is *not* satisfiable





Tableau Reasoning Algorithms

Standard technique based on (hyper-) **tableau**

- Reasoning tasks reducible to (un)**satisfiability**
 - E.g., $\text{KB} \models \text{HeartDisease} \sqsubseteq \text{VascularDisease}$ iff $\text{KB} \cup \{x:(\text{HeartDisease} \sqcap \neg \text{VascularDisease})\}$ is *not* satisfiable
- Algorithm tries to construct (an abstraction of) a model in which some individual (x) is an instance of **HeartDisease** and not an instance of **VascularDisease**
 - such a model is a counter-example for postulated subsumption



Highly Optimised Implementations

- Lazy unfolding
- Simplification and rewriting,
e.g., $A \sqcap B \sqsubseteq C \longrightarrow A \sqsubseteq C \sqcup \neg B$
- HyperTableau (reduces non-determinism)
- Fast semi-decision procedures
- Search optimisations
- Reuse of previous computations
- Heuristics

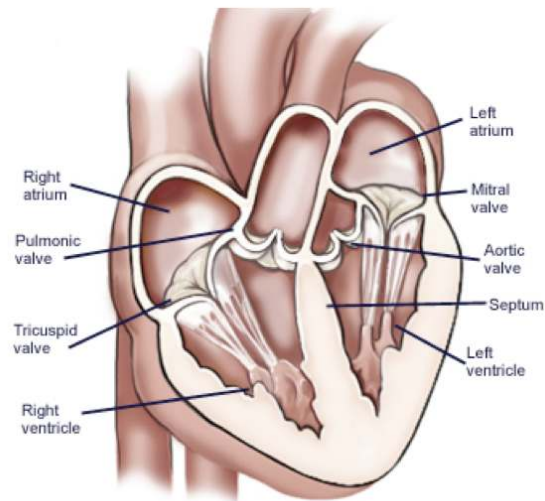
**Not computationally optimal,
but effective with many realistic ontologies**





Scalability Issues

- Problems with very **large and/or cyclical ontologies**



LeftSide $\sqsubseteq \exists$ hasComponent.AorticValve
LeftSide $\sqsubseteq \exists$ hasComponent.MitralValve
AorticValve $\sqsubseteq \exists$ hasConnection.LeftVentricle
MitralValve $\sqsubseteq \exists$ hasConnection.LeftVentricle
LeftVentricle $\sqsubseteq \exists$ isDivisionOf.LeftSide

- Ontologies may define 10s/100s of thousands of terms
 - can lead to construction of *very* large models
 - requires many (worst case n^2) tests to construct taxonomy





Scalability Issues

- Problems with **large data sets** (ABoxes)
 - Main reasoning problem is (conjunctive) query answering, e.g., retrieve all patients suffering from vascular disease:
 $Q(x) \leftarrow \text{Patient}(x) \wedge \text{suffersFrom}(x, y) \wedge \text{VascularDisease}(y)$
 - Decidability still open for OWL, although minor restrictions (on cycles in non-distinguished variables) restore decidability
 - Query answering reduced to standard decision problem, e.g., by checking for each individual x if $\text{KB} \models Q(x)$
 - Model construction starts with *all* ground facts (data)
- Typical applications may use data sets with **10s/100s of millions** of individuals (or more)





OWL 2 Profiles

- OWL recommendation now updated to **OWL 2**
- OWL 2 defines several **profiles** – fragments with desirable computational properties
 - **OWL 2 EL** targeted at very large ontologies
 - **OWL 2 QL** targeted at very large data sets





OWL 2 EL

- A (near maximal) fragment of OWL 2 such that
 - Satisfiability checking is in PTime (**PTime-Complete**)
 - Data complexity of query answering also PTime-Complete
- Based on \mathcal{EL} family of description logics
- Can exploit **saturation** based reasoning techniques
 - Computes classification in “one pass”
 - Computationally optimal
 - Can be extended to Horn fragment of OWL DL



Saturation-based Technique (basics)

- Normalise ontology axioms to standard form:

$$A \sqsubseteq B \quad A \sqcap B \sqsubseteq C \quad A \sqsubseteq \exists R.B \quad \exists R.B \sqsubseteq C$$

- Saturate using inference rules:

$$\frac{A \sqsubseteq B \quad B \sqsubseteq C}{A \sqsubseteq C}$$

$$\frac{A \sqsubseteq B \quad A \sqsubseteq C \quad B \sqcap C \sqsubseteq D}{A \sqsubseteq D}$$

$$\frac{A \sqsubseteq \exists R.B \quad B \sqsubseteq C \quad \exists R.C \sqsubseteq D}{A \sqsubseteq D}$$

- Extension to Horn fragment requires (many) more rules



Saturation-based Technique (basics)

Example:

OrganTransplant \equiv Transplant $\sqcap \exists \text{site}.Organ$

HeartTransplant \equiv Transplant $\sqcap \exists \text{site}.Heart$

Heart \sqsubseteq Organ





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OrganTransplant \sqsubseteq Transplant

OrganTransplant \sqsubseteq \exists site.Organ

\exists site.Organ \sqsubseteq SO

Transplant \sqcap SO \sqsubseteq OrganTransplant





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$\text{HeartTransplant} \sqsubseteq \text{SO}$
 $\text{HeartTransplant} \sqsubseteq \text{OrganTransplant}$





Saturation-based Technique

Performance with large bio-medical ontologies:

	GO	NCI	Galen v.0	Galen v.7	SNOMED
Concepts:	20465	27652	2748	23136	389472
FACT++	15.24	6.05	465.35	—	650.37
HERMIT	199.52	169.47	45.72	—	—
PELLET	72.02	26.47	—	—	—
CEL	1.84	5.76	—	—	1185.70
CB	1.17	3.57	0.32	9.58	49.44
Speed-Up:	1.57X	1.61X	143X	∞	13.15X



OWL 2 QL

- A (near maximal) fragment of OWL 2 such that
 - Data complexity of conjunctive query answering in **AC⁰**, i.e., query answering is *first order reducible*
- Based on **DL-Lite** family of description logics
- Can exploit **query rewriting** based reasoning technique
 - Computationally optimal
 - Data storage and query evaluation can be delegated to standard RDBMS
 - Can be extended to more expressive languages (beyond AC⁰) by delegating query answering to a Datalog engine



Query Rewriting Technique (basics)

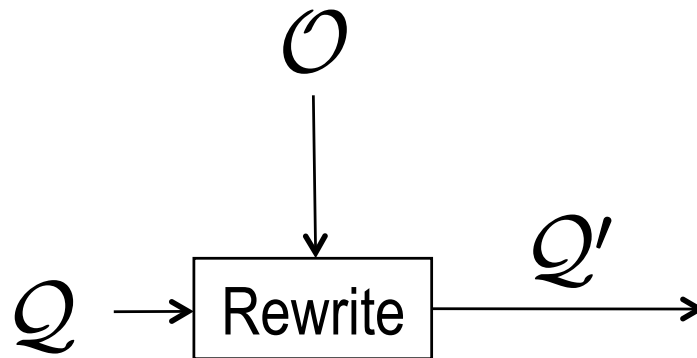
- Given ontology \mathcal{O} and query Q , use \mathcal{O} to rewrite Q as Q' such that, for any set of ground facts \mathcal{A} :
 - $\text{ans}(Q, \mathcal{O}, \mathcal{A}) = \text{ans}(Q', \emptyset, \mathcal{A})$





Query Rewriting Technique (basics)

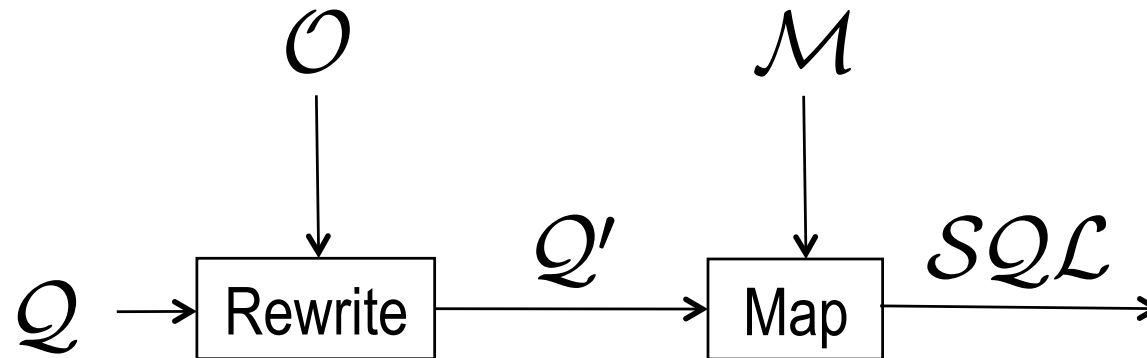
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Query Rewriting Technique (basics)

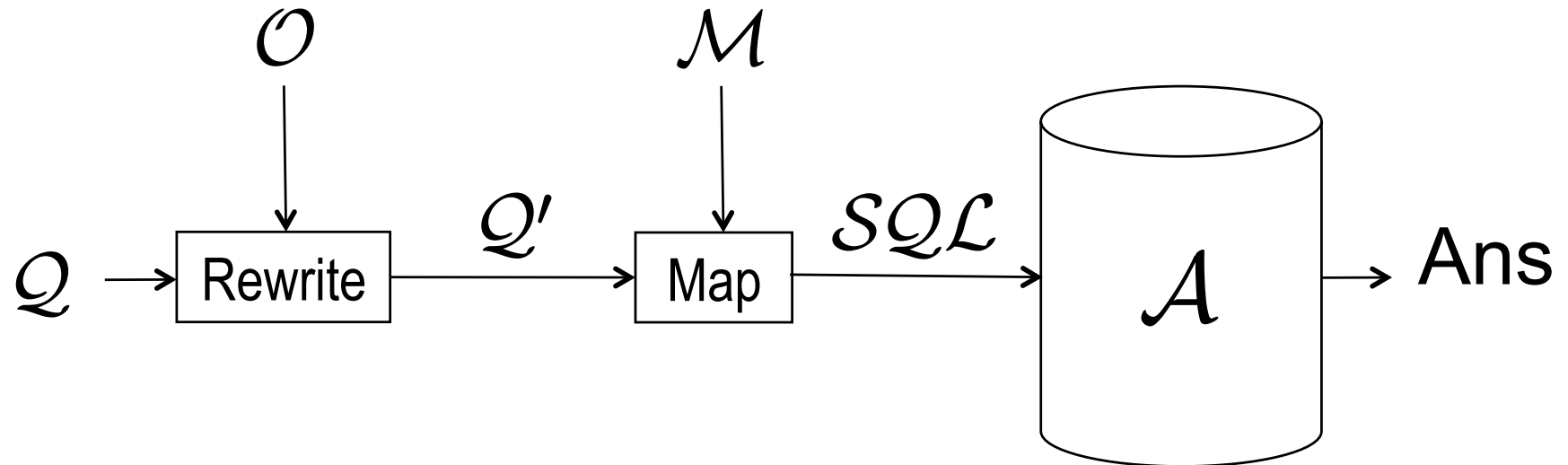
- Given ontology \mathcal{O} and query Q , use \mathcal{O} to rewrite Q as Q' such that, for any set of ground facts \mathcal{A} :
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 - $\text{ans}(Q, \mathcal{O}, \mathcal{A}) = \text{ans}(Q', \emptyset, \mathcal{A})$
- Resolution based query rewriting
 - **Clausify** ontology axioms
 - **Saturate** (clausified) ontology and query using resolution
 - **Prune** redundant query clauses



Query Rewriting Technique (basics)

- Example:

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Consultant $\sqsubseteq \text{Doctor}$

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- For DL-Lite, result is a union of conjunctive queries



Query Rewriting Technique (basics)

- Data can be stored/left in **RDBMS**
- Relationship between ontology and DB defined by **mappings**, e.g.:

Doctor \mapsto SELECT Name FROM Doctor

Patient \mapsto SELECT Name FROM Patient

treats \mapsto SELECT DName, PName FROM Treats

- UCQ translated into **SQL query**:

```
SELECT Name FROM Doctor UNION
```

```
SELECT DName FROM Treats, Patient WHERE PName=Name
```




Some Research Challenges

- Extend saturation-based techniques to non-Horn fragments
 - SNOMED users want negation and/or disjunction
 - Non infectious Pneumonia
 - Infectious or Malignant disorder of lung
 - Burn injury of face neck or scalp
- Extend reasoning support
 - Modularity
 - Explanation
 - ...





Some (more) Research Challenges

- Open questions w.r.t. query rewriting
 - FO rewritability (AC^0) only for very weak ontology languages
 - Even for AC^0 languages, queries can get very large (order $(|\mathcal{O}| \cdot |\mathcal{Q}|)^{|\mathcal{Q}|}$), and existing RDBMSs may behave poorly
 - Larger fragments require (at least) Datalog engines and/or extension to technique (e.g., partial materialisation)
- Integrating DL/DB research
 - Ontologies -v- dependencies
 - Open world -v- closed world





Thanks To

- Boris Motik
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