Secure Hybrid A/D Beamforming for Hardware-Efficient Large-Scale Multiple-Antenna SWIPT Systems

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Abstract—In this work, we investigate the problem of secure communications in a downlink large-scale multi-antenna assisted simultaneous wireless information and power transfer (SWIPT) system, where a base station (BS) transmits signals to serve a number of information decoding (ID) and energy harvesting (EH) users. Considering that the EH users can potentially eavesdrop the ID users' confidential information, we study the robust joint design of the hybrid analog-digital (A/D) beamforming (BF) matrices and of the artificial redundant signal (ARS) covariance matrix at the BS, where the aim is to maximize the worst-case sum secrecy rate for the ID users under a transmit power constraint, a nonlinear EH constraint and a unit-modulus constraint on the entries of the analog BF matrix. The corresponding optimization problem is very challenging due to the nonlinear and nonconvex objective function and constraints. Using innovative optimization techniques, we first transform the original problem into an equivalent but more tractable form, and then develop a novel joint iterative algorithm based on the penalty-concave-convex procedure (CCCP) for solving the resultant problem. We show that the proposed penalty-CCCP based algorithm for ARS-aided robust joint hybrid BF design converges to a Karush-Kuhn-Tucker solution of the original problem, and also analyze its computational complexity. Our simulation results verify that the resultant robust joint hybrid BF design algorithm relying on ARS significantly outperforms the conventional hybrid BF benchmark

algorithms and efficiently achieves the performance of the fully-digital BF with reduced number of radio frequency chains and energy consumption.

Index Terms—Robust hybrid A/D beamforming, secure communications, large-scale multiple-antenna, energy harvesting, joint optimization.

I. INTRODUCTION

A. Literature Review

Large-scale multiple-antenna techniques have attracted considerable interest for the development of next generation wireless networks [1]–[4], since they can dramatically increase system capacity, hence mitigating the spectrum shortage. Nevertheless, the conventional fully-digital beamforming (BF) structure imposes an excessive fabrication cost and energy consumption owing to using numerous costly radio frequency (RF) chains and analog-digital (A/D) converters. In recent years, a hybrid A/D BF structure has been proposed for circumventing this problem, which allows system developers to rely on a lower number of RF chains than the number of antenna elements [5]–[15].

The main challenges in the design of hybrid A/D BF arise owing to the transmit power constraint with coupling components, imposed by the product of analog and digital BF matrices, and by the unit modulus constraint on the elements of the RF analog BF matrix. By exploiting the structure of large-scale multiple-antenna channels, orthogonal matching pursuit (OMP) and channel matching based hybrid BF (HBF) algorithms have been proposed in [5] and [6], respectively. The authors of [7] and [8] developed joint HBF design algorithms based on manifold optimization, while a design procedure based on the maximization of the system's sum rate was proposed in [9]. The authors of [10] developed a sophisticated joint HBF algorithm based on novel penalty methods, and this algorithm was then further extended to millimeter wave multiple-input multiple-output (MIMO) relay systems in [11]. Additionally, considering the associated hardware limitations, several codebook-based HBF algorithms have been investigated in [12] and [13]. Two-timescale hybrid BF algorithms have been proposed for reducing the overhead in large-scale multiple antenna systems [14], [15], where the long-timescale analog BF matrices are designed based on the channel statistics and the short-timescale digital BF matrices are optimized by using the low-dimensional real-time effective channel state information (CSI) matrices.

Electromagnetic energy harvesting (EH) technologies enabling simultaneous wireless information and power transfer

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(SWIPT) have drawn significant attention as a means to provide cost-effective and perpetual power supply for wireless networks [16]–[20]. The authors of [21] studied a multiple-antenna wireless broadcast system consisting of three nodes, where one receiver harvests energy and another receiver decodes information separately from the signals sent by a common transmitter. Two scenarios have been examined, in which the information receiver and energy receiver are separated or co-located. A general receiver operation has been proposed in [22], which splits the received signal with adjustable power ratio for EH and information decoding (ID) separately. References [23]–[25] investigate a multiuser multiple-input single-output (MISO) broadcast SWIPT system configuration, where a multi-antenna base station (BS) sends wireless information and energy signals simultaneously via spatial multiplexing to multiple single-antenna receivers, each implementing either ID or EH. In these works, different optimization problems have been investigated for transmit beamforming design. In [26], a practical nonlinear energy harvesting model and a resource allocation algorithm have been proposed for SWIPT systems.

Considering that the same signal waveform is used for information and energy transmission, security in SWIPT is a major concern, especially within the flourishing context of heterogeneous networks and Internet of things (IoT). Artificial redundant signal (ARS) aided secure BF design algorithms have been developed as a complement to traditional encryption in order to enhance the secrecy performance for the ID users, while supplying energy for the EH users in multiple-antenna SWIPT systems. Some representative studies on this aspect can be found in [27]–[36]. Furthermore, the security problem in large-scale multiple-antenna assisted SWIPT systems has been investigated in [37]–[42]. Recently, the design of secure HBF for large-scale multiple-antenna systems has also attracted substantial interests [43]–[47]. The authors of [43] proposed a HBF algorithm for minimizing the power leaked to the eavesdroppers while guaranteeing the received signal quality of the legitimate user. In order to maximize the system's secrecy rate, the authors of [44] and [45] studied ARS aided secure HBF algorithms only relying on partial CSI. The design of secure HBF has been investigated in two-way MIMO relay networks in [46]. Moreover, the authors of [47] investigated an energy-efficient HBF algorithm to minimize the total transmit energy and meet the signal-to-noise ratio (SNR) and EH requirements of the ID and EH users, respectively.

B. Motivation

Although many studies have been conducted on beamforming for secrecy-preserving SWIPT, most of them are based on the fully-digital beamforming structure which imposes an excessive fabrication cost and energy consumption. Besides, the existing security enhancement techniques conceived for HBF typically design the digital and analog BF matrices separately, which may lead to a significant performance degradation. To the best of our knowledge, the efficient joint design of ARS-aided HBF for secrecy rate maximization in large-scale multiple-antenna assisted SWIPT systems has not been previously addressed, and it is well worth studying due to the hardware-efficient hybrid beamforming structure. Against this backdrop, we consider the problem of secure communications in a downlink large-scale multiple-antenna SWIPT system, where a BS transmits signals to serve a number of ID and EH users simultaneously. Considering that the EH users can potentially eavesdrop the ID users' confidential information, we study the joint design of the hybrid A/D BF matrices and of the ARS covariance matrix at the BS, where the aim is to maximize the worst-case sum secrecy rate for the ID users under a transmit power constraint, a nonlinear EH constraint and a unit-modulus constraint on the entries of the analog BF matrix. The corresponding optimization problem is very challenging due to the nonlinear and nonconvex objective function and constraints. By exploiting the problem structure and using avant-garde techniques, we first transform the original problem into an equivalent yet more tractable form and then develop a novel twin-loop joint iterative algorithm based on the penalty concave-convex procedure (CCCP) to solve the resultant problem. In particular, in the outer loop we adjust the penalty coefficient, while in the inner loop, through insightful approximations and linearization, we optimize the variables in a block coordinate descent fashion, where the subproblems within each block can be solved in parallel. We discuss the convergence and computational complexity of the proposed ARS-aided robust joint HBF design algorithm. Our simulation results show that the proposed algorithm significantly outperforms the existing HBF algorithms and achieves the performance of the fully-digital BF efficiently.

C. Our Contributions

The main contributions of this work are summarized as follows:

1) We consider the robust joint design of hybrid A/D BF and ARS covariance matrices to maximize the worst-case secrecy rate, by taking the CSI uncertainties into account. This problem, which has not been previously addressed in the literature, is very challenging due to the nonlinear objective function, the nonconvex transmit power and EH constraints with coupling terms, and the unit-modulus constraint on the entries of the analog BF matrix. By introducing carefully chosen auxiliary variables, we first convert the challenging constraints into simpler ones, and then transform the objective function into a more tractable form. The resultant problem is equivalent to the original problem but much easier to solve.

2) To solve the transformed problem efficiently, we further introduce a number of auxiliary variables and equality constraints and develop a novel iterative algorithm based on the penalty-CCCP. The proposed algorithm contains two loops, where in the inner loop, through insightful approximations and linearization, the resulting penalized problem is solved in a block coordinate descent fashion to update the optimization variables, while in the outer loop the penalty coefficient is adjusted. Specially, the subproblems within each block of the inner loop can be solved in parallel.

3) We show that the proposed penalty-CCCP-based algorithm for ARS-aided robust joint HBF design converges to a Karush-Kuhn-Tucker (KKT) point of the original

Fig. 1: Hardware-efficient multiple-antenna SWIPT system

where $\hat{\mathbf{h}}_j$ and $\hat{\mathbf{g}}_k$ denote the imperfect estimates while $\triangle\mathbf{h}_j$ and $\triangle\mathbf{g}_k$ capture the corresponding uncertainty. In the absence of *a priori* statistical knowledge about the channel uncertainty, the errors are simply assumed to lie within the following bounded regions

$$\mathcal{H}_j = \{\triangle\mathbf{h}_j : \|\triangle\mathbf{h}_j\|^2 \leq \eta_j^2\} \tag{9}$$

and

$$\mathcal{R}_k = \{\triangle\mathbf{g}_k : \|\triangle\mathbf{g}_k\|^2 \leq \varepsilon_k^2\} \tag{10}$$

where η_j and ε_k denote the corresponding radii of uncertainty.

The worst-case sum secrecy rate achieved by the ID users is given by (11) [38]. In this contribution we aim for jointly designing the baseband BF matrix $\mathbf{W}$, the ARS shaping matrix $\mathbf{F}$ and the analog BF matrix $\mathbf{U}$ in order to maximize the worst-case sum secrecy rate in (11). Hence, the design problem for the secure hybrid A/D BF can be formulated as (12). The constraint (12b) reflects the limited transmit power budget of the BS, while the constant modulus constraint (12c) is due to the fact that the analog beamformer is implemented using low cost phase shifters. Constraint (12d) guarantees that the energy harvested by the kth EH user meets a given energy target ψ_k, where t_k denotes the tight lower bound of $P_{in,k}$ in the presence of CSI errors. We note that the constrained optimization problem (12) is very challenging to solve due to the highly nonconvex objective function and constraints.

III. PROPOSED ARS-AIDED JOINT ROBUST HBF DESIGN ALGORITHM

In this section, we first convert the challenging problem (12) into a more easily tractable yet equivalent form, and then develop a novel penalty-CCCP based iterative algorithm to efficiently obtain its solution.

A. Problem transformation

The transformation of problem (12) into a simpler yet equivalent form relies on the following lemma whose proof is given in Appendix A.

Lemma 1. *Let* $\mathbf{G} \in \mathbb{C}^{L\times L}$ *be a semi-positive definite Hermitian matrix represented in block form as follows,*

$$\mathbf{G} \triangleq \begin{bmatrix} x & \bar{\mathbf{x}}^H \\ \bar{\mathbf{x}} & \bar{\mathbf{G}} \end{bmatrix} \succeq \mathbf{0}, \tag{13}$$

where $x \in \mathbb{C}$, $\bar{\mathbf{x}} \in \mathbb{C}^{(L-1)\times 1}$, *and* $\bar{\mathbf{G}} \in \mathbb{C}^{(L-1)\times(L-1)}$. *Then* $x \geq 0$ *and* $\bar{\mathbf{G}} \succeq \mathbf{0}$, *and if* $x = 0$, *then* $\bar{\mathbf{x}} = \mathbf{0}$.

By defining $\mathbf{X} \triangleq [\mathbf{W}, \mathbf{F}] \in \mathbb{C}^{N\times(J+N)}$, we can represent $\mathbf{W}$ and $\mathbf{F}$ as $\mathbf{W} = \mathbf{X}\boldsymbol{\Omega}_1$ and $\mathbf{F} = \mathbf{X}\boldsymbol{\Omega}_2$, where $\boldsymbol{\Omega}_1 \triangleq [\mathbf{I}_J, \mathbf{0}_{J\times N}]^T$ and $\boldsymbol{\Omega}_2 \triangleq [\mathbf{0}_{N\times J}, \mathbf{I}_N]^T$ denote the selection matrices. Making use of this notation, we can now state the following theorem, which provides the desired equivalence between (12) and a more tractable problem.

Theorem 1. *Problem* (12) *is equivalent to the following problem in* (14), *in the sense that the global optimal solution* $\mathbf{X}$ *and* $\mathbf{U}$ *for the two problems are identical.*

Here $\mathbf{A} \in \mathbb{C}^{M\times M}, \mathbf{B} \in \mathbb{C}^{M\times(N+1)}, \mathbf{C} \in \mathbb{C}^{(N+1)\times(N+1)}, \alpha_j \in \mathbb{R}^+, \beta_j \in \mathbb{R}^+$ and $\gamma_j \in \mathbb{R}^+$ denote a set of auxiliary variables introduced to simplify the problem solution. $\boldsymbol{\omega}_j$ denotes a $J \times 1$ selection vector, where the jth element is one and the remaining elements are zeros.

Proof: Firstly, we show that constraints (12b) and (12c) are equivalent to constraints (14b)–(14e). By defining $\bar{\mathbf{U}} \triangleq \mathbf{U}\mathbf{U}^H$, constraint (14b) can be rewritten as

$$\tilde{\mathbf{U}} \triangleq \begin{bmatrix} \bar{\mathbf{U}} - \mathbf{A} & \mathbf{U}\mathbf{X} - \mathbf{B} \\ \mathbf{X}^H\mathbf{U}^H - \mathbf{B}^H & \mathbf{X}^H\mathbf{X} - \mathbf{C} \end{bmatrix} \succeq \mathbf{0}. \tag{15}$$

It is readily seen that $\hat{\mathbf{U}} \triangleq \bar{\mathbf{U}} - \mathbf{A} \succeq \mathbf{0}$. Hence, it follows on the one hand that $\hat{\mathbf{U}}(m,m) \geq 0$, while on the other hand, we have $\hat{\mathbf{U}}(m,m) = \bar{\mathbf{U}}(m,m) - \mathbf{A}(m,m) = \sum_{n=1}^N |\mathbf{U}(m,n)|^2 - \mathbf{A}(m,m)$. Based on constraints (14d) and (14e), we obtain $\hat{\mathbf{U}}(m,m) \leq 0$. Therefore, we must have $\hat{\mathbf{U}}(m,m) = 0$.

By considering in turn the upper M diagonal entries of the semi-positive definite matrix $\tilde{\mathbf{U}}$, which are equal to zero as shown above, and repeatedly applying *Lemma 1*, we obtain $\bar{\mathbf{U}} = \mathbf{A}$ and $\mathbf{B} = \mathbf{U}\mathbf{X}$. By substituting $\mathbf{B} = \mathbf{U}\mathbf{X}$ into (14c), we obtain the constraint (12b). Moreover, since $\sum_{n=1}^N |\mathbf{U}(m,n)|^2 = \mathbf{A}(m,m) \geq N$, while from (14e), $|\mathbf{U}(m,n)| \leq 1$, we obtain $|\mathbf{U}(m,n)| = 1$, i.e. constraint (12c) must be satisfied.

Secondly, let us turn our attention to the objective function in (12a). By introducing the auxiliary variables α_j, β_j and γ_j, where α_j denotes the upper bound of the SINR among the EH users, $\beta_j \geq \frac{1}{\alpha_j}$, and γ_j denotes the lower bound of the jth ID user's SINR, we can see that problem (12) can be equivalently

$$\sum_{j=1}^{J}\left[\log\left(1+\min_{\forall\triangle\mathbf{h}_j\in\mathcal{H}_j}\mathrm{SINR}_D^j\right)-\log\left(1+\max_{\forall k\in\mathcal{K}}\max_{\forall\triangle\mathbf{g}_k\in\mathcal{R}_k}\mathrm{SINR}_{E,k}^j\right)\right]^+. \tag{11}$$

$$\max_{\mathbf{W},\mathbf{F},\mathbf{U},t_k}\quad\sum_{j=1}^{J}\log\left(1+\min_{\forall\triangle\mathbf{h}_j\in\mathcal{H}_j}\mathrm{SINR}_D^j\right)-\sum_{j=1}^{J}\log\left(1+\max_{\forall k\in\mathcal{K}}\max_{\forall\triangle\mathbf{g}_k\in\mathcal{R}_k}\mathrm{SINR}_{E,k}^j\right) \tag{12a}$$

$$\text{s.t.}\quad \|\mathbf{UW}\|^2+\|\mathbf{UF}\|^2\leq P_t, \tag{12b}$$

$$|\mathbf{U}(m,n)|=1,\forall m,n, \tag{12c}$$

$$\phi_k(t_k)\geq\psi_k,\ P_{in,k}\geq t_k,\ \forall\triangle\mathbf{g}_k\in\mathcal{R}_k,\ \forall k. \tag{12d}$$

$$\max_{\substack{\mathbf{X},\mathbf{U},\mathbf{A},\mathbf{B},\\ \mathbf{C},\alpha_j,\beta_j,\gamma_j>0,t_k}}\quad\sum_{j=1}^{J}\log\left(1+\gamma_j\right)-\sum_{j=1}^{J}\log\left(1+\alpha_j\right) \tag{14a}$$

$$s.t.\quad\begin{bmatrix}\mathbf{U}\\\mathbf{X}^H\end{bmatrix}\begin{bmatrix}\mathbf{U}^H & \mathbf{X}\end{bmatrix}\succeq\begin{bmatrix}\mathbf{A} & \mathbf{B}\\\mathbf{B}^H & \mathbf{C}\end{bmatrix} \tag{14b}$$

$$\|\mathbf{B}\|^2\leq P_t, \tag{14c}$$

$$\mathbf{A}(m,m)\geq N, \tag{14d}$$

$$|\mathbf{U}(m,n)|\leq 1,\forall m,n, \tag{14e}$$

$$\frac{|\mathbf{h}_j^H\mathbf{B}\boldsymbol{\Omega}_1\boldsymbol{\omega}_j|^2}{\|\mathbf{h}_j^H\mathbf{B}\boldsymbol{\Omega}_2\|^2+\sum_{j'\neq j}^{J}|\mathbf{h}_j^H\mathbf{B}\boldsymbol{\Omega}_1\boldsymbol{\omega}_{j'}|^2+\sigma_{D,j}^2}\geq\gamma_j,\ \forall\triangle\mathbf{h}_j\in\mathcal{H}_j,\ \forall j, \tag{14f}$$

$$\alpha_j\beta_j\geq 1,\ \forall j, \tag{14g}$$

$$\frac{|\mathbf{g}_k^H\mathbf{B}\boldsymbol{\Omega}_1\boldsymbol{\omega}_j|^2}{\|\mathbf{g}_k^H\mathbf{B}\boldsymbol{\Omega}_2\|^2+\sum_{j'\neq j}^{J}|\mathbf{g}_k^H\mathbf{B}\boldsymbol{\Omega}_1\boldsymbol{\omega}_{j'}|^2+\sigma_{E,k}^2}\leq\frac{1}{\beta_j},\ \forall\triangle\mathbf{g}_k\in\mathcal{R}_k,\ \forall k, \tag{14h}$$

$$\phi_k(t_k)\geq\psi_k,\ \forall k, \tag{14i}$$

$$\|\mathbf{g}_k^H\mathbf{B}\|^2+\sigma_{E,k}^2\geq t_k,\ \forall\triangle\mathbf{g}_k\in\mathcal{R}_k,\ \forall k. \tag{14j}$$

converted into the following problem:

$$\max_{\substack{\mathbf{X},\mathbf{U},\mathbf{A},\mathbf{B},\\ \mathbf{C},\alpha_j,\beta_j,\gamma_j>0,t_k}}\quad\sum_{j=1}^{J}\log\left(1+\gamma_j\right)-\sum_{j=1}^{J}\log\left(1+\alpha_j\right) \tag{16a}$$

$$\text{s.t.}\quad (14b)-(14h),(12d). \tag{16b}$$

Due to the fact that the constraints (12b) and (12c) are equivalent to constraints (14b)–(14e) and $\mathbf{B}=\mathbf{UX}$, we find that problem (12) can be equivalently transformed into problem (14). This completes the proof. ∎

B. Proposed joint iterative design algorithm

In this subsection, we develop the proposed penalty-CCCP based iterative algorithm for the efficient joint design of robust secure HBF via the solution of problem (14) and subsequently discuss its properties, including convergence behavior and complexity. The proposed penalty-CCCP based algorithm contains two loops[1], where the penalty coefficient is adjusted in the outer loop, while in the inner loop the optimization variables are updated in a block coordinate descent fashion. In each block of the inner loop, we aim to decompose the

[1]The penalty-CCCP optimization framework is introduced in Appendix B.

resulting penalized problem into a number of subproblems, which can be solved easily in parallel. To this end, we introduce auxiliary variables $\mathbf{Y}\in\mathbb{C}^{(M+N+1)\times N}$, $\mathbf{D}\in\mathbb{C}^{(M+N+1)\times(M+N+1)}$, $\mathbf{P}_j\in\mathbb{C}^{M\times(N+1)}$, $\tilde{\mathbf{B}}_{k,j}\in\mathbb{C}^{M\times(N+1)}$, $\bar{\mathbf{B}}_k\in\mathbb{C}^{M\times(N+1)}$, $\tilde{\beta}_{k,j}$, $\bar{t}_k$ and $\bar{\phi}_k$, $\forall k$, $\forall j$, which are subject to the following equality constraints,

$$\mathbf{Y}=\begin{bmatrix}\mathbf{U}\\\mathbf{X}^H\end{bmatrix},\ \mathbf{D}=\begin{bmatrix}\mathbf{A} & \mathbf{B}\\\mathbf{B}^H & \mathbf{C}\end{bmatrix}, \tag{17}$$

$$\mathbf{P}_j=\mathbf{B},\ \tilde{\mathbf{B}}_{k,j}=\mathbf{B},\ \bar{\mathbf{B}}_k=\mathbf{B},\ \tilde{\beta}_{k,j}=\beta_j,\ \bar{t}_k=t_k,$$
$$\bar{\phi}_k=\phi_k(\bar{t}_k),\ \forall k,\ \forall j. \tag{18}$$

Using these variables, we conveniently recast problem (14) as the equivalent form shown in (19), where $\mathcal{S}\triangleq\{\mathbf{X},\mathbf{U},\mathbf{A},\mathbf{B},\mathbf{Y},\mathbf{P}_j,\mathbf{C},\mathbf{D},\tilde{\mathbf{B}}_{k,j},\bar{\mathbf{B}}_k,\bar{t}_k,t_k,\alpha_j,\beta_j,\tilde{\beta}_{k,j},\gamma_j,\bar{\phi}_k\}$ denotes the variable set, while $\boldsymbol{\Phi}_1=[\mathbf{I}_M,\mathbf{0}_{M\times(N+1)}]$, $\boldsymbol{\Phi}_2=[\mathbf{0}_{(N+1)\times M},\mathbf{I}_{N+1}]$, $\boldsymbol{\Psi}_1=[\mathbf{I}_M,\mathbf{0}_{M\times(N+1)}]^T$ and $\boldsymbol{\Psi}_2=[\mathbf{0}_{(N+1)\times M},\mathbf{I}_{N+1}]^T$ are selection matrices.

By penalizing the equality constraints (19h)-(19k) into the objective function, we obtain the following penalized problem for (19):

$$\max_{\mathcal{S}}\quad\sum_{j=1}^{J}\log\left(1+\gamma_j\right)-\sum_{j=1}^{J}\log\left(1+\alpha_j\right)-\mathcal{L}_\rho(\mathcal{X}) \tag{20}$$

$$\text{s.t.}\quad (19b)-(19g),$$

$$\max_{\mathcal{S}} \quad \sum_{j=1}^{J} \log\left(1 + \gamma_j\right) - \sum_{j=1}^{J} \log\left(1 + \alpha_j\right) \tag{19a}$$

$$\text{s.t.} \quad \mathbf{YY}^H \succeq \mathbf{D}, \tag{19b}$$

$$(14c) - (14e), (14g), \tag{19c}$$

$$\frac{|\mathbf{h}_j^H \mathbf{P}_j \boldsymbol{\Omega}_1 \boldsymbol{\omega}_j|^2}{\|\mathbf{h}_j^H \mathbf{P}_j \boldsymbol{\Omega}_2\|^2 + \sum_{j' \neq j}^{J} |\mathbf{h}_j^H \mathbf{P}_j \boldsymbol{\Omega}_1 \boldsymbol{\omega}_{j'}|^2 + \sigma_{D,j}^2} \geq \gamma_j, \, \forall \triangle \mathbf{h}_j \in \mathcal{H}_j, \, \forall j, \tag{19d}$$

$$\frac{|\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \boldsymbol{\Omega}_1 \boldsymbol{\omega}_j|^2}{\|\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \boldsymbol{\Omega}_2\|^2 + \sum_{j' \neq j}^{J} |\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \boldsymbol{\Omega}_1 \boldsymbol{\omega}_{j'}|^2 + \sigma_{E,k}^2} \leq \frac{1}{\tilde{\beta}_{k,j}}, \, \forall \triangle \mathbf{g}_k \in \mathcal{R}_k, \, \forall k, \, \forall j, \tag{19e}$$

$$\bar{\phi}_k \geq \psi_k, \, \forall k, \tag{19f}$$

$$\|\mathbf{g}_k^H \bar{\mathbf{B}}_k\|^2 + \sigma_{E,k}^2 \geq t_k, \, \forall \triangle \mathbf{g}_k \in \mathcal{R}_k, \, \forall k, \tag{19g}$$

$$\boldsymbol{\Phi}_1 \mathbf{Y} = \mathbf{U}, \, \boldsymbol{\Phi}_2 \mathbf{Y} = \mathbf{X}^H, \tag{19h}$$

$$\boldsymbol{\Psi}_1^H \mathbf{D} \boldsymbol{\Psi}_2 = \mathbf{B}, \, \boldsymbol{\Psi}_2^H \mathbf{D} \boldsymbol{\Psi}_1 = \mathbf{B}^H, \tag{19i}$$

$$\boldsymbol{\Psi}_2^H \mathbf{D} \boldsymbol{\Psi}_2 = \mathbf{C}, \, \boldsymbol{\Psi}_1^H \mathbf{D} \boldsymbol{\Psi}_1 = \mathbf{A}, \tag{19j}$$

$$\mathbf{P}_j = \mathbf{B}, \, \tilde{\mathbf{B}}_{k,j} = \mathbf{B}, \, \bar{\mathbf{B}}_k = \mathbf{B}, \, \tilde{\beta}_{k,j} = \beta_j, \, t_k = \bar{t}_k, \, \phi_k(\bar{t}_k) = \bar{\phi}_k, \, \forall k, \, \forall j. \tag{19k}$$

$$\begin{aligned}
\mathcal{L}_\rho(\mathcal{X}) \triangleq \rho\bigg(&\|\boldsymbol{\Phi}_1 \mathbf{Y} - \mathbf{U}\|^2 + \|\boldsymbol{\Phi}_2 \mathbf{Y} - \mathbf{X}^H\|^2 + \|\boldsymbol{\Psi}_1^H \mathbf{D} \boldsymbol{\Psi}_1 - \mathbf{A}\|^2 + \|\boldsymbol{\Psi}_1^H \mathbf{D} \boldsymbol{\Psi}_2 - \mathbf{B}\|^2 \\
&+ \|\boldsymbol{\Psi}_2^H \mathbf{D} \boldsymbol{\Psi}_1 - \mathbf{B}^H\|^2 + \|\boldsymbol{\Psi}_2^H \mathbf{D} \boldsymbol{\Psi}_2 - \mathbf{C}\|^2 + \sum_{j=1}^{J} \|\mathbf{P}_j - \mathbf{B}\|^2 + \sum_{j=1}^{J} \sum_{k=1}^{K} |\beta_j - \tilde{\beta}_{k,j}|^2 \\
&+ \sum_{j=1}^{J} \sum_{k=1}^{K} \|\mathbf{B} - \tilde{\mathbf{B}}_{k,j}\|^2 + \sum_{k=1}^{K} \|\mathbf{B} - \bar{\mathbf{B}}_k\|^2 + \sum_{k=1}^{K} |t_k - \bar{t}_k|^2 + \sum_{k=1}^{K} |\phi_k(\bar{t}_k) - \bar{\phi}_k|^2 \bigg).
\end{aligned} \tag{21}$$

where $\mathcal{L}_\rho(\mathcal{X})$ is expressed as (21) which captures the penalty terms and ρ denotes the penalty coefficient. Moreover, we define $\mathcal{X} \triangleq \{\mathbf{X}, \mathbf{U}, \mathbf{A}, \mathbf{B}, \mathbf{Y}, \mathbf{P}_j, \mathbf{C}, \mathbf{D}, \tilde{\mathbf{B}}_{k,j}, \bar{\mathbf{B}}_k, \beta_j, \tilde{\beta}_{k,j}, t_k, \bar{t}_k, \bar{\phi}_k\}$. In the proposed penalty-CCCP algorithm, the outer loop adjusts the penalty coefficient ρ and the inner loop solves the penalized problem (20) iteratively based on the CCCP in a block coordinate descent fashion.

In order to handle the nonconvex objective function and the nonconvex constraints shown in (19b), (19d), (19e) and (19g), we derive a locally tight lower bound for the objective function along with linearization of the nonconvex constraints. The details are as follows.

1) Lower bound of the objective function: Due to the logarithmic term in its objective function, problem (20) remains difficult to handle. To overcome this difficulty in the local iterative search, we propose to minimize a locally tight upper bound of the objective function instead. Based on the CCCP concept [51], [52], we approximate the function $d_1(\alpha_j) = \log\left(1 + \alpha_j\right)$ in the ith iteration of the proposed algorithm by its first order Taylor expansion around the current value α_j^i, denoted as $\hat{d}_1(\alpha_j^i, \alpha_j) = d_1(\alpha_j^i) + 2\Re\{\nabla d_1^*(\alpha_j^i)(\alpha_j - \alpha_j^i)\} = \log\left(1 + \alpha_j^i\right) + \frac{(\alpha_j - \alpha_j^i)}{(1 + \alpha_j^i)}$. Thus, the locally tight lower bound

of the objective function in (20) is given by

$$\sum_{j=1}^{J} \log\left(1 + \gamma_j\right) - \sum_{j=1}^{J} \left(\log\left(1 + \alpha_j^i\right) + \frac{(\alpha_j - \alpha_j^i)}{(1 + \alpha_j^i)}\right) - \mathcal{L}_\rho(\mathcal{X}). \tag{22}$$

2) Approximation of constraint (19b): Next, let us approximate the nonconvex constraint in (19b). Note that the conventional linearization methods cannot be applied for (19b). However, the following theorem provides a useful alternative.

Theorem 2. *A tight lower bound of* $\mathbf{YY}^H$ *is given by* $\mathbf{Y}(\mathbf{Y}^i)^H + \mathbf{Y}^i \mathbf{Y}^H - \mathbf{Y}^i(\mathbf{Y}^i)^H$, *where* $\mathbf{Y}^i$ *denotes the value of* $\mathbf{Y}$ *at the ith iteration.*

Proof: Let us consider the following positive semidefinite matrix:

$$(\mathbf{Y} - \mathbf{Y}^i)(\mathbf{Y} - \mathbf{Y}^i)^H \succeq \mathbf{0}. \tag{23}$$

By expanding the matrix product in (23), we obtain $\mathbf{YY}^H \succeq \mathbf{YY}^{iH} + \mathbf{Y}^i \mathbf{Y}^H - \mathbf{Y}^i \mathbf{Y}^{iH}$. This completes the proof. ∎

Therefore, by using *Theorem* 2 constraint (19b) can be approximated as the following linear matrix inequality (LMI) constraint:

$$\mathbf{YY}^{iH} + \mathbf{Y}^i \mathbf{Y}^H - \mathbf{Y}^i \mathbf{Y}^{iH} \succeq \mathbf{D}. \tag{24}$$

3) Approximation of constraints (19d)*,* (19e) *and* (19g) *:* Let us rewrite the constraint (19e) as the following difference

of convex (DC) functions:

$$d_2(\tilde{\mathbf{B}}_{k,j}) - d_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}) \leq 0, \forall k, \forall j, \qquad (25)$$

where $d_2(\tilde{\mathbf{B}}_{k,j}) \triangleq |\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \mathbf{\Omega}_1 \boldsymbol{\omega}_j|^2$, $d_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}) \triangleq \frac{\|\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \bar{\mathbf{\Omega}}_j\|^2 + \sigma_{E,k}^2}{\tilde{\beta}_{k,j}}$,
$\bar{\mathbf{\Omega}}_j \triangleq [\mathbf{\Omega}_1 \boldsymbol{\omega}_1, \ldots, \mathbf{\Omega}_1 \boldsymbol{\omega}_{j-1}, \mathbf{\Omega}_1 \boldsymbol{\omega}_{j+1}, \ldots, \mathbf{\Omega}_1 \boldsymbol{\omega}_J, \mathbf{\Omega}_2]$. Using the first order Taylor expansion around $(\tilde{\mathbf{B}}_{k,j}^i, \tilde{\beta}_{k,j}^i)$ in the ith iteration, we linearize $d_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j})$ as shown in (26). Then, (19e) can be approximated as the following convex constraint:

$$d_2(\tilde{\mathbf{B}}_{k,j}) - \hat{d}_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}; \tilde{\mathbf{B}}_{k,j}^i, \tilde{\beta}_{k,j}^i) \leq 0, \forall k, \forall j. \qquad (27)$$

By introducing an auxiliary matrix variable $\mathbf{S}_{k,j} \in \mathbb{C}^{M \times M}$ as the upper bound of $\tilde{\mathbf{B}}_{k,j} \mathbf{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \tilde{\mathbf{B}}_{k,j}^H$, i.e.,

$$\mathbf{S}_{k,j} \succeq \tilde{\mathbf{B}}_{k,j} \mathbf{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \tilde{\mathbf{B}}_{k,j}^H, \forall k, \forall j, \qquad (28)$$

(27) can be equivalently transformed as

$$\mathbf{g}_k^H \mathbf{S}_{k,j} \mathbf{g}_k - \hat{d}_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}; \tilde{\mathbf{B}}_{k,j}^i, \tilde{\beta}_{k,j}^i) \leq 0, \forall k, \forall j. \qquad (29)$$

Note also that (29) must hold for $\forall \triangle \mathbf{g}_k \in \mathcal{R}_k$, which results in an infinite number of inequalities. By applying the S-procedure [53], (29) can be equivalently converted to the following finite number of LMI constraints shown in (30) and (31), where $\mathbf{W}_{k,j} \triangleq$
$\mathbf{S}_{k,j} - \left(\frac{\tilde{\mathbf{B}}_{k,j} \bar{\mathbf{\Omega}}_j \bar{\mathbf{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^{iH} + \tilde{\mathbf{B}}_{k,j}^i \bar{\mathbf{\Omega}}_j \bar{\mathbf{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^H}{\tilde{\beta}_{k,j}^i} - \tilde{\beta}_{k,j} \frac{\tilde{\mathbf{B}}_{k,j}^i \bar{\mathbf{\Omega}}_j \bar{\mathbf{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^{iH}}{(\tilde{\beta}_{k,j}^i)^2} \right)$
and $\tau_{k,j}$ is a slack variable.

Furthermore, by using the Schur complement [49], [53], (28) can be equivalently converted to the following LMI constraints

$$\begin{bmatrix} \mathbf{S}_{k,j} & \tilde{\mathbf{B}}_{k,j} \mathbf{\Omega}_1 \boldsymbol{\omega}_j \\ \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \tilde{\mathbf{B}}_{k,j}^H & 1 \end{bmatrix} \succeq \mathbf{0}, \forall k, \forall j. \qquad (32)$$

By following the same approach, constraint (19d) can be approximated as the convex one shown in (33) which must hold $\forall \triangle \mathbf{h}_j \in \mathcal{H}_j$. Based on the S-procedure and Schur complement, we convert (33) into the following finite number of LMI constraints:

$$\begin{bmatrix} -\bar{\mathbf{W}}_j + \bar{\tau}_j \mathbf{I} & -\bar{\mathbf{W}}_j \hat{\mathbf{h}}_j \\ -\hat{\mathbf{h}}_j^H \bar{\mathbf{W}}_j & -\sigma_{D,j}^2 - \hat{\mathbf{h}}_j^H \bar{\mathbf{W}}_j \hat{\mathbf{h}}_j - \bar{\tau}_j \eta_j^2 \end{bmatrix} \succeq \mathbf{0}, \qquad (34)$$

$$\begin{bmatrix} \bar{\mathbf{S}}_j & \mathbf{P}_j \bar{\mathbf{\Omega}}_j \\ \bar{\mathbf{\Omega}}_j^H \mathbf{P}_j^H & \mathbf{I} \end{bmatrix} \succeq \mathbf{0}, \qquad (35)$$

and

$$\bar{\tau}_j \geq 0, \qquad (36)$$

where $\bar{\mathbf{S}}_j$ is an auxiliary variable, $\bar{\mathbf{W}}_j \triangleq \bar{\mathbf{S}}_j - \left(\frac{\mathbf{P}_j \mathbf{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \mathbf{P}_j^{iH} + \mathbf{P}_j^i \mathbf{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \mathbf{P}_j^H}{\gamma^i} - \gamma_j \frac{\mathbf{P}_j^i \mathbf{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \mathbf{P}_j^{iH}}{\gamma_j^{i2}} \right)$,
and $\bar{\tau}_j$ is a slack variable.

Similarly, (19g) can be approximated as the following linear constraints:

$$\mathbf{g}_k^H \mathbf{\Pi}_k \mathbf{g}_k + \sigma_{E,k}^2 - t_k \geq 0, \forall \triangle \mathbf{g}_k \in \mathcal{R}_k, \qquad (37)$$

where $\mathbf{\Pi}_k \triangleq \bar{\mathbf{B}}_k^i \bar{\mathbf{B}}_k^H + \bar{\mathbf{B}}_k \bar{\mathbf{B}}_k^{iH} - \bar{\mathbf{B}}_k^i \bar{\mathbf{B}}_k^{iH}$. Then, by using the S-procedure, (37) can be transformed into the following finite number of LMI constraints,

$$\begin{bmatrix} \mathbf{\Pi}_k + \tilde{\tau}_k \mathbf{I} & \mathbf{\Pi}_k \hat{\mathbf{g}}_k \\ \hat{\mathbf{g}}_k^H \mathbf{\Pi}_k & -t_k + \sigma_{E,k}^2 + \hat{\mathbf{g}}_k^H \mathbf{\Pi}_k \hat{\mathbf{g}}_k - \tilde{\tau}_k \varepsilon_k^2 \end{bmatrix} \succeq \mathbf{0}, \forall k, \qquad (38)$$

$$\tilde{\tau}_k \geq 0, \forall k, \qquad (39)$$

where $\tilde{\tau}_k$ denotes a slack variable.

Moreover, we note that constraint (14g) can be represented as a second-order cone programming (SOCP) as follows,

$$\left\| \left[1, \frac{\alpha_j - \beta_j}{2} \right] \right\| \leq \frac{\alpha_j + \beta_j}{2}, \forall j. \qquad (40)$$

4) Proposed penalty-CCCP based joint design: By applying (22), (24), (30)–(32), (34)–(36), (38) and (39), we obtain the following approximated convex problem for (20) in the ith iteration of the inner loop:

$$\max_{\bar{\mathcal{S}}} \quad \sum_{j=1}^J \log(1 + \gamma_j) - \sum_{j=1}^J \left(\log(1 + \alpha_j^i) + \frac{(\alpha_j - \alpha_j^i)}{(1 + \alpha_j^i)} \right)$$
$$- \mathcal{L}_\rho(\mathcal{X})$$
$$\text{s.t.} \quad (24), (14c) - (14e), (30) - (32),$$
$$(34) - (36), (38) - (40), (19f),$$
$$\qquad (41)$$

where $\bar{\mathcal{S}} \triangleq \{\mathbf{X}, \mathbf{U}, \mathbf{A}, \mathbf{B}, \mathbf{Y}, \mathbf{P}_j, \mathbf{C}, \mathbf{D}, \tilde{\mathbf{B}}_{k,j}, \bar{\mathbf{B}}_k, \mathbf{S}_{k,j}, \bar{\mathbf{S}}_j, \alpha_j, \beta_j, \tilde{\beta}_{k,j}, \gamma_j, \tau_{k,j}, \bar{\tau}_j, \tilde{\tau}_k, t_k, \bar{t}_k, \bar{\phi}_k\}$ denotes the complete set of variables. In each iteration, the optimization variables are divided into the following four blocks, such that for each block the corresponding subproblems can be solved efficiently in parallel.

In **Block 1**, we optimize $\{\tau_{k,j}, \mathbf{S}_{k,j}, \tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}\}$, $\{\tilde{\tau}_k, \bar{\mathbf{B}}_k, t_k\}$ and $\{\bar{\tau}_j, \bar{\mathbf{S}}_j, \mathbf{P}_j, \gamma_j\}$ by fixing the other variables. In this block, the corresponding subproblems for $\forall k \in \mathcal{K}$ and $\forall j \in \mathcal{J}$ can be solved in parallel.

In **Block 2**, we optimize $\mathbf{B}$, $\{\alpha_j, \beta_j\}$ and $\{\bar{t}_k\}$ by fixing the other variables. The corresponding subproblems can be solved in closed form.

In **Block 3**, we optimize $\mathbf{A}$, $\mathbf{U}$, $\mathbf{X}$ and $\mathbf{C}$ by fixing the other variables. In this block, we have two constrained subproblems and two unconstrained subproblems, which can be solved in closed form.

In **Block 4**, we optimize $(\mathbf{X}, \mathbf{D})$ and $\{\bar{\phi}_k\}$ by fixing the other variables.

A detailed derivation for each one of the updating steps of the inner-loop iterative algorithm is provided in Appendix C. The proposed two-loop penalty-CCCP based joint iterative algorithm is summarized in **Algorithm 1**.

C. Convergence and complexity

Based on the proposed approximations and linearization in Subsection III-B and the convergence analysis for penalty-CCCP in [49], we can see that the proposed **Algorithm 1** converges to the KKT solution of problem (19). Note that (12) and (19) are equivalent, in the sense that the global optimal solution $(\mathbf{w}, \mathbf{F}, \mathbf{U})$ for the two problems are identical. We therefore conclude that any limit point $(\mathbf{w}^\star, \mathbf{F}^\star, \mathbf{U}^\star)$ in $\bar{\mathcal{S}}$ of the iterates generated by **Algorithm 1** is a KKT point of problem (12).

According to the aforementioned description, the proposed joint design algorithm can be implemented in a parallel manner. We note that the overall computational complexity is dominated by the semidefinite programming (SDP) subproblem with respect to $(\bar{\tau}_j, \bar{\mathbf{S}}_j, \mathbf{P}_j, \gamma_j)$ in **Block 1** as shown

$$\hat{d}_3(\tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}; \tilde{\mathbf{B}}_{k,j}^i, \tilde{\beta}_{k,j}^i) = \frac{\text{Tr}\{\bar{\mathbf{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^{iH} \mathbf{g}_k \mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j} \bar{\mathbf{\Omega}}_j + \bar{\mathbf{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^H \mathbf{g}_k \mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j}^i \bar{\mathbf{\Omega}}_j\}}{\tilde{\beta}_{k,j}^i} - \tilde{\beta}_{k,j} \frac{\|\mathbf{g}_k^H \tilde{\mathbf{B}}_{k,j}^i \bar{\mathbf{\Omega}}_j\|^2 + \sigma_{E,k}^2}{(\tilde{\beta}_{k,j}^i)^2} + \frac{2\sigma_{E,k}^2}{\tilde{\beta}_{k,j}^i}. \quad (26)$$

$$\begin{bmatrix} -\mathbf{W}_{k,j} + \tau_{k,j}\mathbf{I} & -\mathbf{W}_{k,j}\hat{\mathbf{g}}_k \\ -\hat{\mathbf{g}}_k^H \mathbf{W}_{k,j} & \frac{2\sigma_{E,k}^2}{\tilde{\beta}_{k,j}^i} - \tilde{\beta}_{k,j}\frac{\sigma_{E,k}^2}{\tilde{\beta}_{k,j}^{i2}} - \hat{\mathbf{g}}_k^H \mathbf{W}_{k,j}\hat{\mathbf{g}}_k - \tau_{k,j}\varepsilon_k^2 \end{bmatrix} \succeq \mathbf{0}, \forall k, \forall j, \quad (30)$$

$$\tau_{k,j} \geq 0, \forall k, \forall j. \quad (31)$$

$$\|\mathbf{h}_j^H \mathbf{P}_j \bar{\mathbf{\Omega}}_j\|^2 + \sigma_{D,j}^2 - \left(\frac{\boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \mathbf{P}_j^{iH} \mathbf{h}_j \mathbf{h}_j^H \mathbf{P}_j \mathbf{\Omega}_1 \boldsymbol{\omega}_j + \boldsymbol{\omega}_j^H \mathbf{\Omega}_1^H \mathbf{P}_j^H \mathbf{h}_j \mathbf{h}_j^H \mathbf{P}_j^i \mathbf{\Omega}_1 \boldsymbol{\omega}_j}{\gamma_j^i} - \frac{\gamma_j \|\mathbf{h}_j^H \mathbf{P}_j^i \mathbf{\Omega}_1 \boldsymbol{\omega}_j\|^2}{(\gamma_j^i)^2}\right) \leq 0. \quad (33)$$

Algorithm 1 Proposed penalty-CCCP based joint iterative design algorithm

---

1. Define the tolerance of accuracy δ_1 and δ_2. Initialize the algorithm with a feasible point. Set the iteration number $i = 0$ and $n = 0$. Set $c > 1$ and $\rho^{(0)} > 0$.
2. **Repeat**
   - **Repeat**
     * Update $\{\tau_{k,j}, \mathbf{S}_{k,j}, \tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}\}$, $\{\tilde{\tau}_k, \bar{\mathbf{B}}_k, t_k\}$, $\{\bar{\tau}_j, \bar{\mathbf{S}}_j, \mathbf{P}_j, \gamma_j\}$, $\forall k \in \mathcal{K}$, $\forall j \in \mathcal{J}$, by solving the subproblems in parallel.
     * Update $\mathbf{B}$, $\{\alpha_j, \beta_j\}$ and $\{\bar{t}_k\}$ by solving the subproblems in parallel.
     * Update $\mathbf{A}$, $\mathbf{U}$, $\mathbf{X}$ and $\mathbf{C}$ by solving four subproblems in parallel.
     * Update $(\mathbf{X}, \mathbf{D})$ and $\{\bar{\phi}_k\}$ by fixing the other variables.
     * Update the iteration number : $i = i + 1$.
   - **Until** the penalty term value is less than δ_1
   - $\rho^{(n+1)} = c\rho^{(n)}$.
   - Update the iteration number : $n = n + 1$.
3. **Until** the difference of successive objective function values is less than δ_2 or the maximum number of iterations is reached, i.e., $n > N_{\max}$.

---

in Appendix C. In this work, we focus on evaluating the order of computational complexity. The complexity of the other subproblems that are solved in closed form or by the method of Lagrange multipliers is much lower than that of the SDP subproblem, hence it can be neglected. Thus, we only concentrate on the main complexity for solving the SDP subproblem which has 1 LMI constraint of size $M + 1$ and 1 LMI constraint of size $M + N + J - 1$. Hence, the overall complexity of the proposed algorithm is given by

$$\mathcal{O}(n\sqrt{2M + N + J}((M + 1)^3 + n(M + 1)^2 \\ + (M + N + J - 1)^3 + n(M + N + J - 1)^2 + n^2)), \quad (42)$$

where the number of decision variables n is on the order of $\mathcal{O}\left(M^2 + M(N + J) + 2\right)$ [54].

IV. SIMULATION RESULTS

In this section, we evaluate the performance of the proposed robust design algorithm for ARS-aided joint HBF. We assume that the BS is equipped with a uniform linear array of $M = 32$ transmit antennas and, unless specified otherwise, the transmit power, the number of RF chains, the number of ID users, the number of EH users, the saturation threshold, the energy conversion efficiency, and the EH target are set to $P_t = 40$ dBm, $N = 4$, $J = 2$, $K = 2$, $P_{sat,k} = -5$ dBm, $\xi = 0.5$, and $\psi = \psi_k = -20$ dBm, $\forall k \in \mathcal{K}$, respectively. For the large-scale multiple-antenna channel model, we employ the uniform linear antenna array and assume that there are 6 clusters and 8 rays per cluster as in [9], [10]. The large-scale path loss is modeled as (in dB) $PL = PL_0 - 10\alpha \log_{10}(d/d_0)$, where PL_0 denotes the path loss at reference distance d_0, d is the link distance in meters, and α denotes the path loss exponent. Here, we set $d_0 = 1$ and $PL_0 = -30$ dB. The distances of the ID and EH users are $d_{ID,j} = 500$, $\forall j \in \mathcal{J}$ and $d_{EH,k} = 5$, $\forall k \in \mathcal{K}$, respectively [28]. The path loss exponents regarding the ID and EH users are set as $\alpha_{ID,j} = \alpha_{EH,k} = 3.6$ [55]. Moreover, we set the channel bandwidth as 20 MHz and the noise power density as -174 dBm/Hz. For simplicity, the channel error bounds are assumed to be identical, i.e., $\eta_j = \varepsilon_k = 0.1$, $\forall j \in \mathcal{J}$, $\forall k \in \mathcal{K}$. Regarding the implementation of the proposed ARS-aided robust joint HBF design algorithm (**Algorithm 1**), we set $N_{max} = 30$, $\rho_0 = 1$, $c = 2$, $\delta_1 = 10^{-6}$ and $\delta_2 = 10^{-4}$.

Fig. 2 shows the convergence performance of the proposed **Algorithm 1**. Specifically, Fig. 2 (a) shows the evolution of the sum secrecy rate versus the number of iterations for this algorithm. We can observe that the secrecy rate increases with the number of iterations and converges within 25 iterations, which in this case demonstrates the convergence of the proposed algorithm. Moreover, Fig. 2 (b) indicates the sum of all the penalty terms versus the number of iterations, which supports our claim that the proposed penalty-CCCP based joint design algorithm can tackle the equality constraint effectively.

In the following, the performance of the proposed algorithm is compared with that of other benchmark approaches. Specifically, we consider the following HBF approaches:

- Proposed: The proposed ARS-aided robust joint HBF design algorithm.
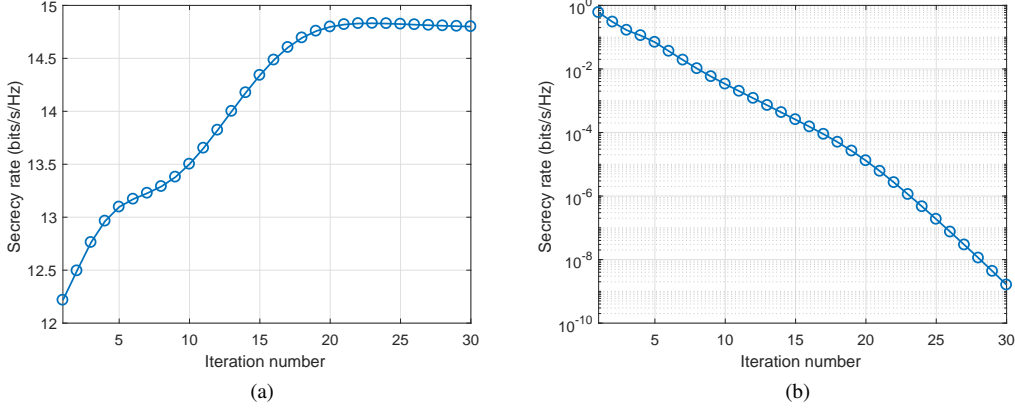- Separate: The analog BF matrix $\mathbf{U}$ is first obtained

Fig. 2: The convergence performance of the proposed ARS-aided robust joint HBF design algorithm: (a) Secrecy rate versus iteration number; (b) penalty term versus iteration number.

by using the channel matching approach [6]. The total transmit power is equally allocated to the digital BF matrix $\mathbf{W}$ and the ARS shaping matrix $\mathbf{F}$, where the former is designed to maximize the ID users' sum rate and the latter is designed in the null space of the ID users' estimated channel matrices and aligned with the EH users estimated channel matrices.

- Without ARS: The analog BF matrix $\mathbf{U}$ and digital BF matrix $\mathbf{W}$ are obtained by the channel matching approach [6] for ID users without the consideration of EH users.
- Fully digital: The ARS-aided robust BF is optimized assuming a fully-digital structure.
- Nonrobust: Simplified version of the proposed ARS-aided joint HBF design algorithm based on estimated CSI only (that is, knowledge of CSI errors is not used).

We first compare the proposed ARS-aided robust joint HBF design algorithm to its nonrobust counterpart. Fig. 3 (a) and (b) demonstrate the secrecy rate and feasible rate versus different values of ψ, respectively. In this simulation, the solution produced by a design algorithm is considered *feasible* for a given channel realization if all the associated constraints are satisfied. The feasible rate represents the proportion (fraction) of simulations experiments where the design algorithm produces a feasible solution in this sense. From the results, the robust algorithm significantly outperforms the nonrobust algorithm for $\psi \leq -15$ dBm. Although the secrecy rate of the nonrobust algorithm is higher than that of the robust algorithm for $\psi = -10$ dBm, the proposed robust design algorithm provides 100% feasible rate which is significantly better than that of its nonrobust counterpart. It shows the effectiveness of the proposed approach against the CSI uncertainties. The improvement in feasible rate comes from the robust operation which takes the CSI error into account.

In Fig. 4, we show the secrecy rate versus the transmit power for the algorithms under comparison. It is seen that the proposed ARS-aided robust joint HBF design algorithm achieves the performance of the fully-digital robust BF algorithm with reduced RF chains and energy consumption. Moreover, the proposed design algorithm significantly outperforms the conventional separate HBF design algorithm and the HBF

algorithm without ARS. In particular, for a given secrecy rate of 13 bits/s/Hz, the proposed algorithm can save almost 3 dBm in transmit power compared to the conventional separate HBF design method. The secrecy rate of the HBF algorithm without ARS is higher than that of the separate HBF design algorithm for 30 dBm transmit power. That is because the separate HBF design algorithm has to allocate more power to serve the EH users in the case of low transmit power, which degrades the secrecy rate performance. However, the HBF algorithm without ARS does not consider the EH constraints. Fig. 5 shows the secrecy rate performance versus the value of EH target ψ. The proposed HBF design algorithm outperforms the separate HBF design algorithm and almost achieves the performance of the fully-digital BF algorithm. The difference between the secrecy rate performance of the proposed ARS-aided robust joint HBF design algorithm and that of the fully-digital BF algorithm increases with the increase of the EH target value, since more energy is used for wireless power transfer in order to meet the EH constraint in the hybrid beamforming structure. The number of RF chains can be increased to guarantee the secrecy rate performance for large values of EH target. Due to the heuristic design without the consideration of EH users, the HBF algorithm without ARS cannot satisfy the EH constraint and its performance does not change with ψ.

Subsequently, we investigate the power consumption and energy efficiency performance of the proposed and conventional algorithms. We assume that the levels of power consumption for each low noise amplifier (LNA), RF chain and A/D converter are given by $20mW$, $40mW$ and $350mW$, respectively [56]. We also consider that the power consumption of the baseband processing is given by $200mW$. Thus, the power consumption of the fully-digital BF is given by $M \times (20mW + 40mW + 350mW) + 200mW$, while the one of the A/D HBF is given by $M \times 20mW + N \times 40mW + N \times 350mW + 200mW$. By fixing the number of RF chains to $N = 8$, the results in Table I show that the HBF structure can significantly reduce the power consumption compared to the fully-digital BF structure, especially as the number of transmit antennas increases. Then, let us define the energy efficiency as $\frac{\tilde{R}}{\tilde{P}}$,
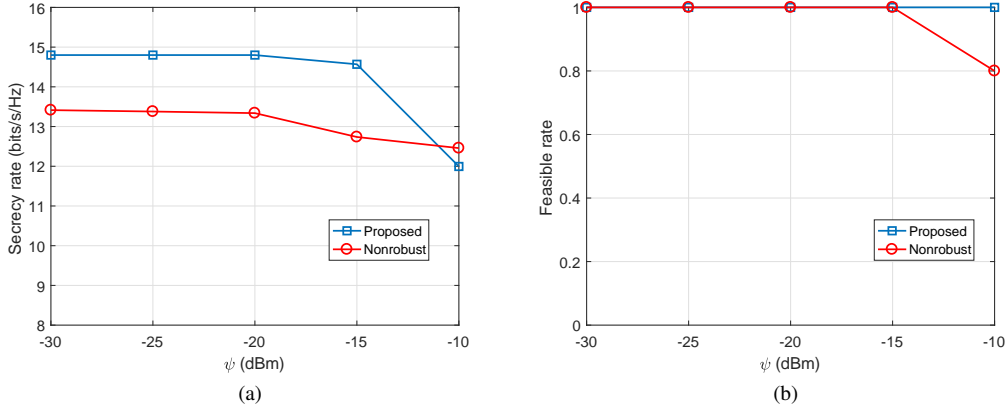
Fig. 3: Comparison of the proposed ARS-aided robust joint HBF design algorithm and its nonrobust counterpart: (a) Secrecy rate versus ψ; (b) feasible rate versus ψ.

Fig. 4: Comparison of the proposed ARS-aided robust joint HBF design algorithm and the conventional algorithms: Secrecy rate versus transmit power.
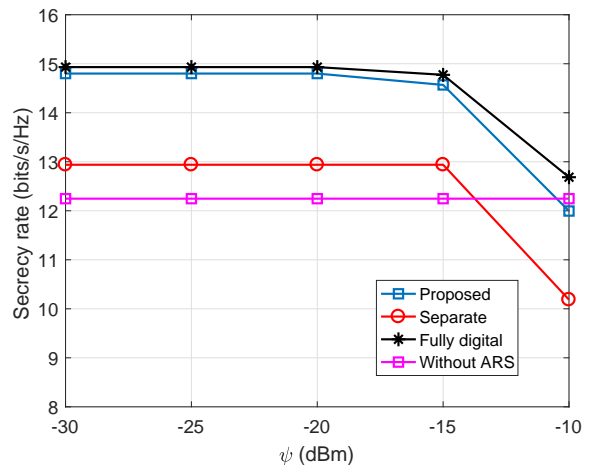
Fig. 5: Comparison of the proposed ARS-aided robust joint HBF design algorithm and the conventional algorithms: Secrecy rate versus the EH target.

where $\tilde{R}$ denotes the secrecy rate of the ID user and $\tilde{P}$ denotes the total power consumption. Fig. 6 shows the energy efficiency for the various algorithms under study versus the number of RF chains. It is seen that the energy efficiency performance of all the HBF approaches decreases with the number of RF chains, and that the proposed ARS-aided robust joint HBF design algorithm outperforms the other compared algorithms and provides the best performance. This is because the proposed HBF design algorithm based on the innovative optimization technique can achieve the same secrecy rate performance as the fully-digital BF with greatly reduced power consumption.

In Fig. 7, we show the secrecy rate performance of the proposed joint hybrid beamforming design algorithm and existing beamforming algorithms versus the number of EH users. From the results, the performance of all the analyzed algorithms degrades with the increase of the number of EH users. The proposed ARS-aided robust joint hybrid beamforming design algorithm outperforms the existing hybrid beamforming al-

gorithms and approaches the performance of the fully-digital beamforming algorithm.

Finally, we investigate the impact of finite resolution phase shifters on the proposed algorithm. Specifically, in Fig. 8, we plot its secrecy rate versus the transmit power for different number N_b of phase shifter quantization bits, in the case of $N = 4$ RF chains and $K = 2$ EH users. It is seen that the performance of the proposed algorithm improves with the number of quantization bits N_b as expected. In particular, the performance with $N_b = 8$ can approach the performance with infinite resolution phase shifters.

V. CONCLUSION

We have designed ARS-aided robust joint hybrid A/D BF for the secrecy rate maximization of a large-scale multiple-antenna assisted SWIPT system. To tackle this challenging multi-dimensional optimization problem, featuring nonlinear objective and nonconvex constraints, we first transformed it into a more tractable, yet equivalent form. We then proposed

TABLE I: Power consumption for fully-digital and hybrid beamforming structures ($N = 8$)

	$M = 32$	$M = 64$	$M = 128$	$M = 256$
Fully Digital	$13.32W$	$26.44W$	$52.68W$	$105.16W$
Hybrid A/D	$3.96W$	$4.60W$	$5.88W$	$8.44W$

Fig. 6: Comparison of the proposed ARS-aided robust joint HBF design algorithm and the conventional algorithms: Energy efficiency versus number of RF chains.

Fig. 8: Secrecy rate of the proposed ARS-aided robust joint HBF design versus the transmit power with finite resolution phase shifters ($N = 4$ and $K = 2$).
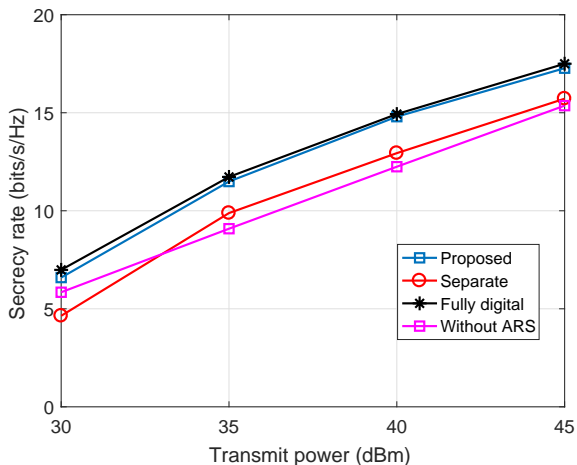

Fig. 7: Comparison of the proposed ARS-aided robust joint HBF design algorithm and the conventional algorithms: Secrecy rate versus number of EH users.
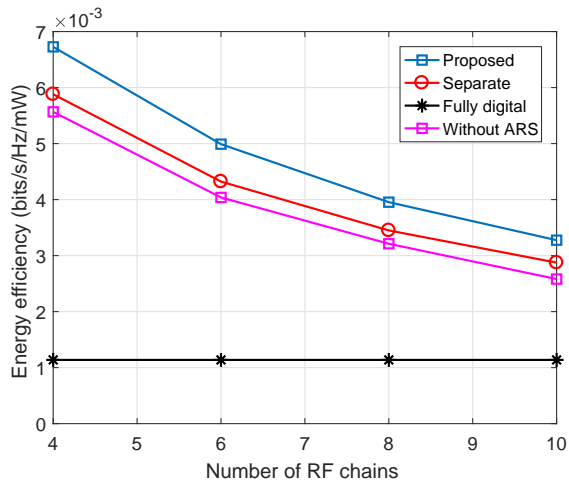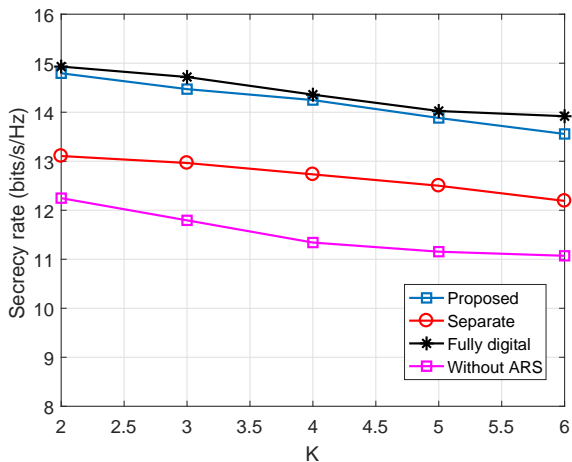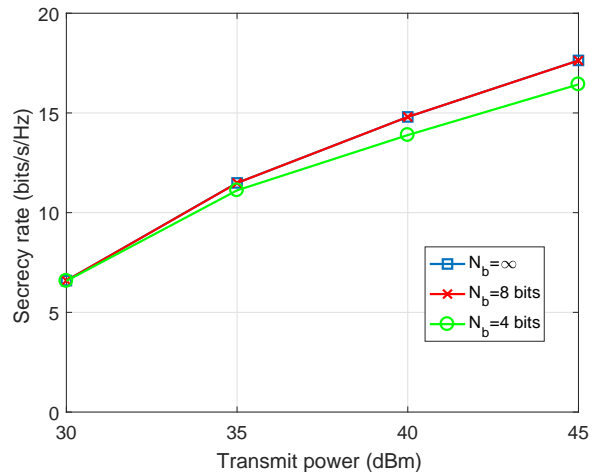
a penalty-CCCP based twin-loop joint iterative algorithm to solve the resultant problem. Specifically, the penalty coefficient is adjusted in the outer loop, while in the inner loop the optimization variables are updated in a block coordinate descent fashion, where the subproblems within each block can be solved in parallel. We discussed the convergence and computational complexity of the proposed algorithm. Our simulation results have shown that the proposed ARS-aided robust joint HBF design algorithm can outperform the existing HBF benchmark algorithms and efficiently achieves the performance of the fully-digital BF with reduced RF chains and energy consumption.

APPENDIX A
PROOF OF LEMMA 1

For any $\mathbf{d} \triangleq [\tilde{d}, \bar{\mathbf{d}}^T]^T \in \mathbb{C}^{L \times 1}$, where $\bar{\mathbf{d}} \in \mathbb{C}^{(L-1) \times 1}$, we have

$$\mathbf{d}^H \begin{bmatrix} x & \bar{\mathbf{x}}^H \\ \bar{\mathbf{x}} & \bar{\mathbf{G}} \end{bmatrix} \mathbf{d} = x|\tilde{d}|^2 + \tilde{d}\bar{\mathbf{d}}^H\bar{\mathbf{x}} + \tilde{d}^*\bar{\mathbf{x}}^H\bar{\mathbf{d}} + \bar{\mathbf{d}}^H\bar{\mathbf{G}}\bar{\mathbf{d}} \geq 0.$$

(43)

Hence, by using $\mathbf{d} = [0, \bar{\mathbf{d}}^T]^T$, we obtain $\bar{\mathbf{G}} \succeq \mathbf{0}$. Next, we show that when $x = 0$ we have $\bar{\mathbf{x}} = \mathbf{0}$ by contradiction.

Assume that when $x = 0$ we have $\bar{\mathbf{x}} \neq \mathbf{0}$, then for any $\mathbf{d}$ with $\bar{\mathbf{x}}^H\bar{\mathbf{d}} \neq 0$ (43) can be recast as

$$\tilde{d}\bar{\mathbf{d}}^H\bar{\mathbf{x}} + \tilde{d}^*\bar{\mathbf{x}}^H\bar{\mathbf{d}} + \bar{\mathbf{d}}^H\bar{\mathbf{G}}\bar{\mathbf{d}} \geq 0. \tag{44}$$

However, using $\tilde{d} = -\frac{\bar{\mathbf{d}}^H\bar{\mathbf{G}}\bar{\mathbf{d}}\bar{\mathbf{x}}^H\bar{\mathbf{d}}}{|\bar{\mathbf{x}}^H\bar{\mathbf{d}}|^2}$ we have

$$\tilde{d}\bar{\mathbf{d}}^H\bar{\mathbf{x}} + \tilde{d}^*\bar{\mathbf{x}}^H\bar{\mathbf{d}} + \bar{\mathbf{d}}^H\bar{\mathbf{G}}\bar{\mathbf{d}} < 0. \tag{45}$$

This contradicts (43). Therefore, when $x = 0$ we have $\bar{\mathbf{x}} = \mathbf{0}$. Lemma 1 is thus proved.

APPENDIX B
PENALTY-CCCP METHOD

In this appendix, we briefly introduce the proposed penalty-CCCP method in a general framework. Let us consider a general problem:

$$(P) \quad \min_{\boldsymbol{x} \in \mathcal{X}} f(\boldsymbol{x})$$
$$\text{s.t.} \quad \boldsymbol{h}(\boldsymbol{x}) = \mathbf{0}, \tag{46}$$
$$\boldsymbol{g}(\boldsymbol{x}) \leq \mathbf{0}$$

where $\mathcal{X} \subseteq \mathbb{R}^n$ denotes a closed convex set, $f(\boldsymbol{x})$ is a scalar continuously differentiable function, $\boldsymbol{h}(\boldsymbol{x}) \in \mathbb{R}^p$ is a vector of

TABLE II: Penalty-CCCP method for problem (46)

```
0.  initialize x⁰ ∈ 𝒳, ϱ₀ > 0, and set c > 1, k = 0
1.  repeat
2.      x^{k+1} = CCCP(P_{ϱ_k}, x^k)
3.      ϱ_{k+1} = cϱ_k
4.      k = k + 1
5.  until some termination criterion is met
```

p continuously differentiable functions, and $\boldsymbol{g}(\boldsymbol{x}) \in \mathbb{R}^q$ is a vector of differentiable but *possibly nonconvex* functions.

We can tackle problem (P) by using a penalty method [57], i.e., by solving the penalized problem:

$$(P_\varrho) \quad \min_{\boldsymbol{x} \in \mathcal{X}} f_\rho(\boldsymbol{x}) \triangleq f(\boldsymbol{x}) + \frac{\varrho}{2}\|\boldsymbol{h}(\boldsymbol{x})\|^2 \tag{47}$$
$$\text{s.t. } \boldsymbol{g}(\boldsymbol{x}) \leq \boldsymbol{0}$$

where $\varrho > 0$ is a scalar penalty parameter that prescribes a high cost for the violation of the equality constraints. In particular, when $\varrho \to \infty$, solving the above problem yields an identical solution to problem (P) [57]. However, it is still difficult to globally solve problem (P_ϱ) when it is a nonconvex problem. Thus, we propose the penalty-CCCP method summarized in Table II to solve problem (P), where in each iteration problem (P_ϱ) is approximately solved using the CCCP method. The resulting solution is denoted simply as $\boldsymbol{x}^{k+1} = CCCP(P_{\varrho_k}, \boldsymbol{x}^k)$ in Step 2 of Table II.

Regarding the convergence of penalty-CCCP, this method converges to the KKT solution of the original problem (P). The detailed proof is available from [49].

APPENDIX C
DERIVATION OF UPDATING STEPS FOR THE INNER-LOOP ALGORITHM

A. Variables Updating in Block 1

In this block, problem (41) can be decomposed into a number of independent subproblems, which can be solved in parallel.

1) Subproblems with respect to $\{\tau_{k,j}, \mathbf{S}_{k,j}, \tilde{\mathbf{B}}_{k,j}, \tilde{\beta}_{k,j}\}$, $\forall k \in \mathcal{K}, \forall j \in \mathcal{J}$: Each subproblem can be expressed as (48), where $\mathbf{W}_{k,j} \triangleq \mathbf{S}_{k,j} - \left(\frac{\tilde{\mathbf{B}}_{k,j} \bar{\boldsymbol{\Omega}}_j \bar{\boldsymbol{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^{iH} + \tilde{\mathbf{B}}_{k,j}^i \bar{\boldsymbol{\Omega}}_j \bar{\boldsymbol{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^H}{\tilde{\beta}_{k,j}^i} - \tilde{\beta}_{k,j} \frac{\tilde{\mathbf{B}}_{k,j}^i \bar{\boldsymbol{\Omega}}_j \bar{\boldsymbol{\Omega}}_j^H \tilde{\mathbf{B}}_{k,j}^{iH}}{(\tilde{\beta}_{k,j}^i)^2} \right)$. This is a standard SDP problem, which can be solved by using the convex programming toolbox CVX [58].

2) Subproblems with respect to $\{\tilde{\tau}_k, \bar{\mathbf{B}}_k, t_k\}$, $\forall k \in \mathcal{K}$: Each subproblem can be expressed as

$$\min_{\tilde{\tau}_k \geq 0, \bar{\mathbf{B}}_k, t_k} \quad \|\mathbf{B} - \bar{\mathbf{B}}_k\|^2$$
$$\text{s.t. } \begin{bmatrix} \boldsymbol{\Pi}_k + \tilde{\tau}_k \mathbf{I} & \boldsymbol{\Pi}_k \hat{\mathbf{g}}_k \\ \hat{\mathbf{g}}_k^H \boldsymbol{\Pi}_k & -t_k + \sigma_{E,k}^2 + \hat{\mathbf{g}}_k^H \boldsymbol{\Pi}_k \hat{\mathbf{g}}_k - \tilde{\tau}_k \varepsilon_k^2 \end{bmatrix} \succeq \mathbf{0}, \tag{49}$$

where $\boldsymbol{\Pi}_k \triangleq \bar{\mathbf{B}}_k^i \bar{\mathbf{B}}_k^H + \bar{\mathbf{B}}_k \bar{\mathbf{B}}_k^{iH} - \bar{\mathbf{B}}_k^i \bar{\mathbf{B}}_k^{iH}$. This SDP problem can be solved by using the toolbox CVX.

3) Subproblems with respect to $\{\bar{\tau}_j, \bar{\mathbf{S}}_j, \mathbf{P}_j, \gamma_j\}$: This subproblem can be expressed as

$$\min_{\bar{\tau}_j \geq 0, \bar{\mathbf{S}}_j, \mathbf{P}_j, \gamma_j} \quad \rho\|\mathbf{B} - \mathbf{P}_j\|^2 - \log(1 + \gamma_j)$$
$$\text{s.t. } \begin{bmatrix} -\bar{\mathbf{W}}_j + \bar{\tau}_j \mathbf{I} & -\bar{\mathbf{W}}_j \hat{\mathbf{h}}_j \\ -\hat{\mathbf{h}}_j^H \bar{\mathbf{W}}_j & -\sigma_{D,j}^2 - \hat{\mathbf{h}}_j^H \bar{\mathbf{W}}_j \hat{\mathbf{h}}_j - \bar{\tau}_j \eta_j^2 \end{bmatrix} \succeq \mathbf{0},$$
$$\begin{bmatrix} \bar{\mathbf{S}}_j & \mathbf{P}_j \bar{\boldsymbol{\Omega}}_j \\ \bar{\boldsymbol{\Omega}}_j^H \mathbf{P}_j^H & \mathbf{I} \end{bmatrix} \succeq \mathbf{0}, \tag{50}$$

where $\bar{\mathbf{W}}_j \triangleq \bar{\mathbf{S}}_j - \left(\frac{\mathbf{P}_j \boldsymbol{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \boldsymbol{\Omega}_1^H \mathbf{P}_j^{iH} + \mathbf{P}_j^i \boldsymbol{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \boldsymbol{\Omega}_1^H \mathbf{P}_j^H}{\gamma^i} - \gamma_j \frac{\mathbf{P}_j^i \boldsymbol{\Omega}_1 \boldsymbol{\omega}_j \boldsymbol{\omega}_j^H \boldsymbol{\Omega}_1^H \mathbf{P}_j^{iH}}{\gamma_j^{i2}} \right)$. This SDP problem can be solved by using the toolbox CVX.

B. Variables Updating in Block 2

In this block, by fixing the other variables, problem (41) with respect to $\mathbf{B}$, $\{\alpha_j, \beta_j\}$ and $\{\bar{t}_k\}$, $\forall k \in \mathcal{K}, \forall j \in \mathcal{J}$, can be decomposed into several independent subproblems, so that the variables can be optimized in parallel.

1) Subproblems with respect to $\{\alpha_j, \beta_j\}$: Each subproblem is given by

$$\min_{\alpha_j, \beta_j} \quad \frac{(\alpha_j - \alpha_j^i)}{(1 + \alpha_j^i)} + \rho \sum_{k=1}^{K} |\beta_j - \tilde{\beta}_{k,j}|^2 \tag{51a}$$
$$\text{s.t. } \sqrt{1 + \frac{(\alpha_j - \beta_j)^2}{4}} \leq \frac{\alpha_j + \beta_j}{2}. \tag{51b}$$

By attaching a Lagrange multiplier μ to the constraint, we obtain the following Lagrange function

$$\mathcal{L}(\alpha_j, \beta_j, \mu) \triangleq \frac{(\alpha_j - \alpha_j^i)}{(1 + \alpha_j^i)} + \rho \sum_{k=1}^{K} |\beta_j - \tilde{\beta}_{k,j}|^2$$
$$+ \mu \left(\sqrt{1 + \frac{(\alpha_j - \beta_j)^2}{4}} - \frac{\alpha_j + \beta_j}{2} \right).$$

By examining the first order optimality condition of $\mathcal{L}(\alpha_j, \beta_j, \mu)$ with respect to α_j, we obtain

$$\alpha_j(\mu) = \frac{(1 + \alpha_j^i)\mu - 2}{\sqrt{(1 + \alpha_j^i)\mu - 1}} + \frac{2\rho \sum_{k=1}^{K} \tilde{\beta}_{k,j} - \frac{1}{1 + \alpha_j^i} + \mu}{2\rho K}. \tag{52}$$

Similarly, by examining the first order optimality condition of $\mathcal{L}(\alpha_j, \beta_j, \mu)$ with respect to β_j, we have

$$\beta_j(\mu) = \frac{2\rho \sum_{k=1}^{K} \tilde{\beta}_{k,j} - \frac{1}{1 + \alpha_j^i} + \mu}{2\rho K}. \tag{53}$$

The value of the multiplier μ must be chosen to ensure the complementarity slackness condition of the constraint (51b). In this regard, it is convenient to define the function

$$\mathcal{C}(\alpha_j, \beta_j) \triangleq \sqrt{1 + \frac{(\alpha_j - \beta_j)^2}{4}} - \frac{\alpha_j + \beta_j}{2}.$$

$$\min_{\tau_{k,j}\geq 0,\mathbf{S}_{k,j},\tilde{\mathbf{B}}_{k,j},\tilde{\beta}_{k,j}} \quad |\beta_j - \tilde{\beta}_{k,j}|^2 + \|\mathbf{B} - \tilde{\mathbf{B}}_{k,j}\|^2$$

$$\text{s.t.} \quad \begin{bmatrix} -\mathbf{W}_{k,j} + \tau_{k,j}\mathbf{I} & -\mathbf{W}_{k,j}\hat{\mathbf{g}}_k \\ -\hat{\mathbf{g}}_k^H\mathbf{W}_{k,j} & \frac{2\sigma_{E,k}^2}{\tilde{\beta}_{k,j}^i} - \tilde{\beta}_{k,j}\frac{\sigma_{E,k}^2}{\tilde{\beta}_{k,j}^{i2}} - \hat{\mathbf{g}}_k^H\mathbf{W}_{k,j}\hat{\mathbf{g}}_k - \tau_{k,j}\varepsilon_k^2 \end{bmatrix} \succeq \mathbf{0}, \tag{48}$$

$$\begin{bmatrix} \mathbf{S}_{k,j} & \tilde{\mathbf{B}}_{k,j}\boldsymbol{\Omega}_1\boldsymbol{\omega}_j \\ \boldsymbol{\omega}_j^H\boldsymbol{\Omega}_1^H\tilde{\mathbf{B}}_{k,j}^H & 1 \end{bmatrix} \succeq \mathbf{0}.$$

The optimal μ should be positive and satisfy $\mathcal{C}(\alpha_j(\mu), \beta_j(\mu)) = 0$, which is equivalent to

$$\frac{1}{\sqrt{(1+\alpha_j^i)\mu - 1}} - \left(\frac{2\rho\sum_{k=1}^K \tilde{\beta}_{k,j} - \frac{1}{1+\alpha_j^i} + \mu}{2\rho K} \right) = 0. \tag{54}$$

We can obtain the optimal μ by using a one dimensional search to solve (54). By substituting the resulting value of μ in (52) and (53), we can obtain the optimal values of α_j and β_j.

2) Subproblem with respect to $\mathbf{B}$: The subproblem is given by

$$\min_{\mathbf{B}} \quad \|\boldsymbol{\Psi}_1^H\mathbf{D}\boldsymbol{\Psi}_2 - \mathbf{B}\|^2 + \|\boldsymbol{\Psi}_2^H\mathbf{D}\boldsymbol{\Psi}_1 - \mathbf{B}^H\|^2 + \sum_{j=1}^J \|\mathbf{P}_j - \mathbf{B}\|^2$$

$$+ \sum_{j=1}^J\sum_{k=1}^K \|\mathbf{B} - \tilde{\mathbf{B}}_{k,j}\|^2 + \sum_{k=1}^K \|\mathbf{B} - \bar{\mathbf{B}}_k\|^2$$

$$\text{s.t.} \quad \|\mathbf{B}\|^2 \leq P_t. \tag{55}$$

The objective function can be reformulated as $(2K+3)\|\mathbf{B} - \mathbf{B}_0\|^2 + Z$, where

$$\mathbf{B}_0 = \frac{1}{JK + K + J + 2}\left(\boldsymbol{\Psi}_1^H\mathbf{D}\boldsymbol{\Psi}_2 + \boldsymbol{\Psi}_1^H\mathbf{D}^H\boldsymbol{\Psi}_2 + \sum_{j=1}^J \mathbf{P}_j \right.$$

$$\left. + \sum_{j=1}^J\sum_{k=1}^K \tilde{\mathbf{B}}_{k,j} + \sum_{k=1}^K \bar{\mathbf{B}}_k \right), \tag{56}$$

and

$$Z = \|\boldsymbol{\Psi}_1^H\mathbf{D}\boldsymbol{\Psi}_2\|^2 + \|\boldsymbol{\Psi}_1^H\mathbf{D}^H\boldsymbol{\Psi}_2\|^2 + \sum_{j=1}^J \|\mathbf{P}_j\|^2$$

$$+ \sum_{j=1}^J\sum_{k=1}^K \|\tilde{\mathbf{B}}_{k,j}\|^2 + \sum_{k=1}^K \|\bar{\mathbf{B}}_k\|^2 - (JK+K+J+2)\|\mathbf{B}_0\|^2. \tag{57}$$

Therefore, the subproblem (55) is equivalent to

$$\min_{\mathbf{B}} \quad \|\mathbf{B} - \mathbf{B}_0\|^2$$

$$\text{s.t.} \quad \|\mathbf{B}\|^2 \leq P_t. \tag{58}$$

It is readily seen that (58) is equivalent to the projection of $\mathbf{B}_0$ onto a hypersphere $\|\mathbf{B}\|^2 \leq P_t$. Thus, the solution is

$$\mathbf{B} = P_t\frac{\mathbf{B}_0}{\|\mathbf{B}_0\| + \max(0, P_t - \|\mathbf{B}_0\|)}. \tag{59}$$

3) Subproblems with respect to $\{\bar{t}_k\}$: Each subproblem is given by

$$\min_{\bar{t}_k} \quad |t_k - \bar{t}_k|^2 + |\phi_k(\bar{t}_k) - \bar{\phi}_k|^2. \tag{60}$$

According to the first order optimality condition and the piecewise nature of the function $\phi_k(\bar{t}_k)$, the solution for $\bar{t}_k$ is given by

$$\bar{t}_k = \begin{cases} t_k & \bar{\phi}_k > \xi P_{sat,k} \\ \frac{\xi\bar{\phi}_k + t_k}{1+\xi^2} & \bar{\phi}_k \leq \xi P_{sat,k}. \end{cases} \tag{61}$$

C. Variables Updating in Block 3

In this block, by fixing the other variables, problem (41) with respect to $\mathbf{A}$, $\mathbf{U}$, $\mathbf{X}$, $\mathbf{C}$ can be decomposed into four independent subproblems, so that the variables in this block can be optimized in parallel.

1) Subproblem 1 with respect to $\mathbf{A}$: The subproblem is formulated as

$$\min_{\mathbf{A}} \quad \|\boldsymbol{\Psi}_1^H\mathbf{D}\boldsymbol{\Psi}_1 - \mathbf{A}\|^2$$

$$\text{s.t.} \quad \mathbf{A}(m,m) \geq N, \forall m.$$

Based on the first order optimality condition, we obtain $\mathbf{A}_0 = \boldsymbol{\Psi}_1^H\mathbf{D}\boldsymbol{\Psi}_1$. For the diagonal elements in $\mathbf{A}$, we set

$$\mathbf{A}(m,m) = \max(\mathbf{A}_0(m,m), N), \forall m.$$

For the elements that are not on the diagonal, we set

$$\mathbf{A}(m,n) = \mathbf{A}_0(m,n), \forall m \neq n.$$

2) Subproblem 2 with respect to $\mathbf{C}$: The subproblem is given by

$$\min_{\mathbf{C}} \quad \|\boldsymbol{\Psi}_2^H\mathbf{D}\boldsymbol{\Psi}_2 - \mathbf{C}\|^2.$$

By applying the first order optimality condition, the optimal $\mathbf{C}$ can be obtained as

$$\mathbf{C} = \boldsymbol{\Psi}_2^H\mathbf{D}\boldsymbol{\Psi}_2.$$

3) Subproblem 3 with respect to $\mathbf{U}$: The subproblem is given as

$$\min_{\mathbf{U}} \quad \|\mathbf{U} - \mathbf{U}_0\|^2$$

$$\text{s.t.} \quad |\mathbf{U}(m,n)| \leq 1, \forall m,n, \tag{62}$$

where we define $\mathbf{U}_0 \triangleq \boldsymbol{\Phi}_1\mathbf{Y}$. The subproblem (62) can be further divided into MN parallel subproblems. For each element $\mathbf{U}(m,n)$, we have the following subproblem:

$$\min_{\mathbf{U}(m,n)} \quad \|\mathbf{U}(m,n) - \mathbf{U}_0(m,n)\|^2$$

$$\text{s.t.} \quad |\mathbf{U}(m,n)| \leq 1. \tag{63}$$

Similar to subproblem (55), the subproblem (63) is equivalent to a projection. Therefore, we can obtain

$$\mathbf{U}(m,n) = \frac{\mathbf{U}_0(m,n)}{|\mathbf{U}_0(m,n)| + \max(0, 1 - |\mathbf{U}_0(m,n)|)}, \forall m,n.$$

4) Subproblem 4 with respect to $\mathbf{X}$*:* The subproblem is given as

$$\min_{\mathbf{X}} \quad \|\mathbf{\Phi}_2 \mathbf{Y} - \mathbf{X}^H\|^2.$$

By applying the first order optimality condition, the optimal $\mathbf{X}$ can be obtained as

$$\mathbf{X} = \mathbf{Y}^H \mathbf{\Phi}_2^H.$$

D. Variables Updating in Block 4

In this block, problem (41) can be decomposed into the following subproblems, which can be solved in parallel.

1) Subproblem with respect to $(\mathbf{Y}, \mathbf{D})$*:* This subproblem is given by

$$\min_{\mathbf{Y},\mathbf{D}} \quad \|\mathbf{\Phi}_1 \mathbf{Y} - \mathbf{U}\|^2 + \|\mathbf{\Phi}_2 \mathbf{Y} - \mathbf{X}^H\|^2 + \|\mathbf{\Psi}_1^H \mathbf{D} \mathbf{\Psi}_1 - \mathbf{A}\|^2$$
$$+ \|\mathbf{\Psi}_1^H \mathbf{D} \mathbf{\Psi}_2 - \mathbf{B}\|^2 + \|\mathbf{\Psi}_2^H \mathbf{D} \mathbf{\Psi}_1 - \mathbf{B}^H\|^2$$
$$+ \|\mathbf{\Psi}_2^H \mathbf{D} \mathbf{\Psi}_2 - \mathbf{C}\|^2$$
$$\text{s.t.} \quad \mathbf{Y}\mathbf{Y}^{iH} + \mathbf{Y}^i \mathbf{Y}^H - \mathbf{Y}^i \mathbf{Y}^{iH} \succeq \mathbf{D}. \tag{64}$$

This is a standard SDP problem, which can be solved by using the toolbox CVX.

2) Subproblems with respect to $\{\bar{\phi}_k\}$*:* Each subproblem is given by

$$\min_{\bar{\phi}_k} \quad |\phi_k(\bar{t}_k) - \bar{\phi}_k|^2 \tag{65}$$
$$\text{s.t.} \quad \psi_k \leq \bar{\phi}_k \leq \xi P_{sat,k}.$$

By applying the first order optimality condition, the optimal $\bar{\phi}_k$ can be obtained as

$$\bar{\phi}_k = \min(\xi P_{sat,k}, \max(\psi_k, \phi_k(\bar{t}_k))). \tag{66}$$

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