



Note

# Sparse halves in triangle-free graphs <sup>☆</sup>

Peter Keevash <sup>a</sup>, Benny Sudakov <sup>b</sup>

<sup>a</sup> Department of Mathematics, Caltech, Pasadena, CA 91125, USA

<sup>b</sup> Department of Mathematics, Princeton University, Princeton, NJ 08544, USA

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## Abstract

One of Erdős' favourite conjectures was that any triangle-free graph  $G$  on  $n$  vertices should contain a set of  $n/2$  vertices that spans at most  $n^2/50$  edges. We prove this when the number of edges in  $G$  is either at most  $n^2/12$  or at least  $n^2/5$ .

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## 1. Introduction

A fundamental result in extremal graph theory is Turán's theorem from 1941, that the unique largest graph on  $n$  vertices not containing a copy of  $K_t$  (the complete graph on  $t$  vertices) is the Turán graph  $T_{t-1}(n)$ , which is the complete  $(t-1)$ -partite graph with part sizes as equal as possible. (The case  $t=3$  was proved by Mantel in 1907.) A generalisation that takes into account edge distribution, or local density, was introduced by Erdős [2] who asked the following question. Suppose  $0 \leq \alpha, \beta \leq 1$  and that  $G$  is a  $K_t$ -free graph on  $n$  vertices in which every set of  $\alpha n$  vertices span at least  $\beta n^2$  edges. How large can  $\beta$  be as a function of  $\alpha$ ? Erdős, Faudree, Rousseau and Schelp [5] studied this problem and conjectured that there is a constant  $c_t < 1$  so that if  $c_t \leq \alpha \leq 1$  then the largest possible  $\beta$  is  $\frac{t-2}{t-1}(\alpha - 1/2)$  (which is attained by the Turán graph  $T_{t-1}(n)$ ). They proved this for triangle-free graphs ( $t=3$ ) and the general case was proved by the authors [7].

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*E-mail addresses:* keevash@caltech.edu (P. Keevash), bsudakov@math.princeton.edu (B. Sudakov).

Moreover, for triangle-free graphs and general  $\alpha$  it was conjectured in [5] that  $\beta$  is determined by a family of extremal triangle-free graphs. Besides the complete bipartite graph  $T_2(n)$  already mentioned, another important graph is  $C_5(n/5)$ , which is obtained from a 5-cycle by replacing each vertex  $i$  by an independent set  $V_i$  of size  $n/5$  (assuming for simplicity that  $n$  is divisible by 5), and each edge  $ij$  by a complete bipartite graph joining  $V_i$  and  $V_j$  (this operation is called a ‘blow-up’). It is not difficult to see that a set  $S$  of  $\alpha n$  edges spanning as few edges as possible will either contain  $V_i$  or be disjoint from  $V_i$  for all but at most one  $i$ ; indeed, if  $S$  intersects  $V_i$  and  $V_j$  non-trivially it is possible to increase one intersection and decrease the other without increasing the number of edges spanned by  $S$ . For  $2/5 \leq \alpha \leq 3/5$  it follows that every  $\alpha n$  vertices in  $C_5(n/5)$  span at least  $\frac{5\alpha-2}{25}n^2$  edges. This is larger than  $\frac{2\alpha-1}{4}n^2$  (the value for  $T_2(n)$ ) when  $\alpha < 17/30$ . Erdős et al. conjectured that above this value the largest  $\beta$  is always  $\frac{2\alpha-1}{4}$ , i.e., the constant  $c_3$  (defined above) can be taken equal to  $17/30$ . The best-known bound for this problem is due to Krivelevich [8], who showed that one can take  $c_3 < 3/5$ . They also conjectured that  $\beta$  is  $\frac{5\alpha-2}{25}$  for a certain range of  $\alpha$  below  $17/30$ , including  $\alpha = 1/2$ .

The case  $\alpha = 1/2$  is an old question of Erdős that he returned to often in his problems papers, starting with [2], through to [3] where he offered a \$ 250 prize for its solution. Here the conjecture is that any triangle-free graph on  $n$  vertices should contain a set of  $n/2$  vertices that spans at most  $n^2/50$  edges. Krivelevich [8] has shown that this holds when  $n^2/50$  is replaced by  $n^2/36$ . In this paper we prove the following result, which establishes the conjecture under an additional assumption on the total number of edges in the graph, and shows that  $C_5(n/5)$  is the unique extremal example in the range that we consider.

**Theorem 1.1.** *Let  $G$  be a triangle-free graph on  $n$  vertices with at least  $n^2/5$  edges such that every set of  $\lfloor n/2 \rfloor$  vertices of  $G$  spans at least  $n^2/50$  edges. Then  $n = 10m$  for some integer  $m$  and  $G = C_5(2m)$ .*

Also, it is not difficult to obtain an analogous result for graphs with few edges.

**Proposition 1.2.** *Let  $G$  be a triangle-free graph on  $n$  vertices with at most  $n^2/12$  edges. Then some set of  $\lfloor n/2 \rfloor$  vertices of  $G$  spans at most  $n^2/50$  edges.*

Another problem, that is similar in spirit, is to determine how many edges one may need to delete from a triangle-free graph on  $n$  vertices in order to make it bipartite. A long-standing conjecture of Erdős [2] is that at most  $n^2/25$  edges need to be deleted, and  $C_5(n/5)$  shows that this would be best possible. A related conjecture, that any  $K_4$ -free graph on  $n$  vertices can be made bipartite with the omission of at most  $n^2/9$  edges was proved recently by the second author [9]. For triangle-free graphs, the best-known bound is  $(1/18 - \epsilon)n^2$  for some calculable constant  $\epsilon > 0$ , obtained by Erdős, Faudree, Pach and Spencer [4].

Krivelevich [8] noticed that for regular graphs a bound in the local density problem implies a bound for the problem of making the graph bipartite. Indeed, suppose  $n$  is even,  $G$  is a  $d$ -regular graph on  $n$  vertices and  $S$  is a set of  $n/2$  vertices. Then  $dn/2 = \sum_{s \in S} d(s) = 2e(S) + e(S, \bar{S})$  and  $dn/2 = \sum_{s \notin S} d(s) = 2e(\bar{S}) + e(S, \bar{S})$ , i.e.,  $e(S) = e(\bar{S})$ . Deleting the  $2e(S)$  edges within  $S$  or  $\bar{S}$  makes the graph bipartite, so if we could find  $S$  spanning at most  $n^2/50$  edges we would delete at most  $n^2/25$  in making  $G$  bipartite. The converse reasoning does not work, as may be seen from considering the blow-up  $P(n/10)$  (supposing  $n$  is divisible by 10), where  $P$  is the Petersen graph (see Fig. 1). In this graph every set of  $n/2$  vertices spans at least  $n^2/50$  edges (illustrated by the black circles), but it can be made bipartite by deleting only  $3n^2/100$  edges (illustrated by

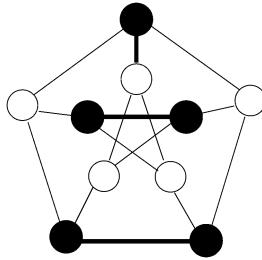


Fig. 1. The Petersen blow-up  $P(n/10)$ .

the bold lines). This seems to indicate that the local density problem may be harder, especially since it appears that there is not a unique extremal example. Noting that  $P(n/10)$  has  $3n^2/20$  edges, it seems interesting to extend our bound of  $n^2/12$  in Proposition 1.2 to  $3n^2/20$ . We expect that graphs with between  $3n^2/20$  and  $n^2/5$  edges will be the most challenging to deal with.

It is not difficult to show (see [4]) that if  $G$  is a triangle-free graph on  $n$  vertices with  $e \geq n^2/5$  edges then it can be made bipartite by deleting at most  $n^2/25$  edges. Some extensions of this result can be found in [6]. With Theorem 1.1 we solve the local density problem in the same range.

**Notation.** Suppose  $G$  is a graph. For a vertex  $v$  we let  $N(v)$  denote its neighbourhood and  $d(v) = |N(v)|$  its degree. If  $X$  is a set of vertices then  $G[X]$  is the restriction of  $G$  to  $X$ , i.e., the graph with vertex set  $X$  whose edges are edges of  $G$  with both endpoints in  $X$ . We write  $e(X) = e(G[X])$  for the number of edges in  $X$ . If  $X$  and  $Y$  are sets of vertices then  $e(X, Y)$  is the number of edges with one endpoint in  $X$  and the other in  $Y$ . For an integer  $t$ , the blow-up  $G(t)$  is the graph obtained from  $G$  by replacing each vertex  $i$  by an independent set  $V_i$  of size  $t$ , and each edge  $ij$  by a complete bipartite graph joining  $V_i$  and  $V_j$ .

## 2. Proof of Proposition 1.2

Throughout this note we will repeatedly use the two following standard averaging arguments. Suppose we have a graph containing a set  $X$  of  $x$  vertices, such that there are  $e_1$  edges with exactly one endpoint in  $X$  and  $e_2$  edges with both endpoints in  $X$ . Then for any integer  $y$  with  $0 \leq y \leq x$ , by considering a random  $y$ -subset of  $X$  we see that:

- (i) there is a subset of  $X$  of size  $y$  containing at most  $(y/x)^2 e_2$  edges, and
- (ii) there is a subset  $Y$  of  $X$  of size  $y$  so that the number of edges incident to  $Y$ , i.e., having at least one endpoint in  $Y$ , is at most  $(y/x)e_1 + (y/x)^2 e_2 \leq (y/x)(e_1 + e_2)$ .

Suppose  $G$  is a triangle-free graph on  $n$  vertices and  $\theta n^2$  edges. As a warm-up, we note that the case  $\theta \leq 2/25$  is easy to deal with: by averaging there is a set of  $n/2$  vertices spanning at most  $(1/2)^2 \theta n^2 \leq n^2/50$  edges. With a little more work we can deal with the range  $\theta \leq 1/12$  as follows. Let  $I$  be an independent set of size  $2\theta n$  (we can choose such a set inside the neighbourhood of a vertex of maximum degree). Let  $I' = V(G) - I$  and suppose that there are  $\phi n^2$  edges in  $G[I']$ . By averaging some set of  $n/2$  vertices in  $I'$  spans at most  $(\frac{1/2}{1-2\theta})^2 \phi n^2 \geq n^2/50$  edges, so we can assume  $\phi \geq \frac{2}{25}(1 - 2\theta)^2$ . Now by averaging there is a subset  $J \subseteq I'$  of  $(1/2 - 2\theta)n$  vertices so that

$$\begin{aligned}
 n^{-2}e(I \cup J) &\leq n^{-2} \left( \frac{1/2 - 2\theta}{1 - 2\theta} (e(G) - e(I')) + \left( \frac{1/2 - 2\theta}{1 - 2\theta} \right)^2 e(I') \right) \\
 &= \frac{1/2 - 2\theta}{1 - 2\theta} (\theta - \phi) + \left( \frac{1/2 - 2\theta}{1 - 2\theta} \right)^2 \phi \\
 &= \frac{1/2 - 2\theta}{1 - 2\theta} \theta - \frac{1/2 - 2\theta}{2(1 - 2\theta)^2} \phi \\
 &\leq \frac{1/2 - 2\theta}{1 - 2\theta} \theta - (1/2 - 2\theta)/25 \leq 1/50.
 \end{aligned}$$

Here the last inequality follows from the fact that the function  $f(t) = \frac{1/2 - 2t}{1 - 2t} t - (1/2 - 2t)/25$  is increasing when  $t \leq 1/12$ .

### 3. Proof of Theorem 1.1

We will need a result of Andrásfai, Erdős and Sós [1], which states that if  $G$  is a triangle-free graph on  $n$  vertices with minimum degree at least  $2n/5$  then either  $G$  is bipartite or  $G = C_5(n/5)$  is the blow-up of a 5-cycle.<sup>1</sup>

Now let  $G$  be a triangle-free graph on  $n$  vertices with at least  $n^2/5$  edges such that every set of  $\lfloor n/2 \rfloor$  vertices of  $G$  spans at least  $n^2/50$  edges. We will show that  $n$  is divisible by 10 and  $G = C_5(n/5)$ .

First suppose that we have proved the theorem under the assumption that  $n$  is divisible by 10, and consider the case when  $n$  is not divisible by 10. Then every set of  $\lfloor n/2 \rfloor$  vertices of  $G$  spans at least  $\lceil n^2/50 \rceil > n^2/50$  edges, since  $n^2/50$  is not an integer. Consider the blow-up  $H = G(10)$ , obtained by replacing each vertex  $i$  of  $G$  by a set  $V_i$  of size 10. Then  $H$  has  $h = 10n$  vertices and at least  $h^2/5$  edges. As noted in the introduction, it is not difficult to see that there is a set  $S$  of  $h/2$  vertices spanning the minimal number of edges which either contains or is disjoint from the set  $V_i$  for all but at most one index  $i$ . Then  $S$  contains  $\lfloor n/2 \rfloor$  sets  $V_i$ , so it spans at least  $100\lceil n^2/50 \rceil > h^2/50$  edges. Applying the theorem to  $H$  we see that  $H = C_5(2n)$ . But then  $H$  has a set of  $h/2$  vertices spanning  $h^2/50$  edges, a contradiction. Therefore we can assume that  $n$  is divisible by 10.

Next we mention two simple consequences of the fact that  $G$  is triangle-free. One is that for any vertex  $v$  the neighbourhood  $N(v)$  is an independent set. Another is that any set of vertices  $U$  spans at most  $|U|^2/4$  edges, by Mantel’s theorem.

We start the argument with the following useful observation. Suppose that  $G$  contains an independent set  $I$  of size  $(2/5 + t)n$  for some  $t > 0$ , and that there is a vertex  $v$  with at least  $(1/5 - t)n$  neighbours in  $I$ . Then  $v$  has less than  $(1/10 - t)n$  neighbours outside  $I$ . For otherwise we can add a set of size  $(1/10 - t)n$  from  $N(x) \setminus I$  to obtain a set of size  $n/2$  in which the only edges are those going between  $N(x) \setminus I$  and  $I \setminus N(x)$ .<sup>2</sup> The number of such edges is at most  $(1/10 - t)n \cdot (1/5 + 2t)n = (1/50 - 2t^2)n^2 < n^2/50$ , a contradiction.

Suppose that the maximum degree of  $G$  is  $(2/5 + t)n$ , where  $t \geq 0$ . We divide the proof into two cases according to the value of  $t$ .

<sup>1</sup> Note that this theorem immediately implies our theorem when restricted to regular graphs, i.e., any regular triangle-free graph on  $n$  vertices with at least  $n^2/5$  edges contains a set of  $n/2$  vertices that spans at most  $n^2/50$  edges.

<sup>2</sup> Since  $n$  is divisible by 10 and  $|I| = (2/5 + t)n$  is an integer we see that  $tn$  is an integer, and so  $(1/10 - t)n$  is an integer. Similar comments apply throughout the proof, but we will not labour the point.

**Case 1.**  $t \geq 1/135$ .

Let  $A$  be the neighbourhood of a vertex of maximum degree, so  $|A| = (2/5 + t)n$ . Since  $A$  is an independent set, we certainly have  $t < 1/10$ . Let  $B$  be the set of vertices with at least  $(1/5 - t)n$  neighbours in  $A$ , and let  $C$  be the set of vertices not in  $A$  or  $B$ . By the previous observation every vertex in  $B$  has less than  $(1/10 - t)n$  neighbours outside  $A$ .

By definition every vertex in  $C$  has less than  $(1/5 - t)n$  neighbours in  $A$ . We claim that  $|C| < (1/10 - t)n$ . Otherwise let  $X \subseteq C$  has size  $(1/10 - t)n$  and considers  $A \cup X$ , which has size  $n/2$ . Now

$$\begin{aligned} n^2/50 &\leq e(A \cup X) < |X| \cdot (1/5 - t)n + |X|^2/4 \\ &= (1/10 - t)(1/5 - t)n^2 + (1/10 - t)^2 n^2/4 \\ &= n^2/50 + (5t^2/4 - 7t/20 + 1/400)n^2, \end{aligned}$$

which gives  $5t^2/4 - 7t/20 + 1/400 > 0$ . Solving the quadratic and recalling that  $t < 1/10$  we obtain  $t < (7 - 2\sqrt{11})/50 < 1/135$ , a contradiction.

It follows that  $|B| = n - |A| - |C| > n/2$ . Consider  $Y \subseteq B$  of size  $n/2$ . Recalling that every vertex in  $B$  has less than  $(1/10 - t)n$  neighbours outside  $A$  we have  $n^2/50 \leq e(Y) < \frac{1}{2} \cdot n/2 \cdot (1/10 - t)n$ , so  $t < 1/50$ .

Let  $D$  be the set of vertices  $v \in B$  with at most  $(1/5 + 2t)n$  neighbours in  $A$ . Vertices in  $B$  have less than  $(1/10 - t)n$  neighbours outside  $A$ , so vertices in  $D$  have degree less than  $(3/10 + t)n$ . Recall that the maximum degree of  $G$  is  $(2/5 + t)n$  and let  $|D| = pn$ . Now we have

$$\begin{aligned} 2/5 &\leq 2e(G)/n^2 = n^{-2} \sum_v d(v) \leq n^{-2} (|D|(3/10 + t)n + (n - |D|)(2/5 + t)n) \\ &= p(3/10 + t) + (1 - p)(2/5 + t) = 2/5 + t - p/10, \end{aligned}$$

so  $p \leq 10t$ . Now we note that the only edges of  $G$  within  $B$  are those internal to  $D$ . Indeed, in any other pair of vertices, both have at least  $(1/5 - t)n$  neighbours in  $A$ , and one vertex has more than  $(1/5 + 2t)$  neighbours in  $A$ . Since  $|A| = (2/5 + t)n$  they must have a common neighbour in  $A$ . Therefore these vertices are non-adjacent, as  $G$  is triangle-free. This implies that any set of  $n/2$  vertices in  $B$  spans at most  $e(D) \leq |D|^2/4 \leq 25t^2 n^2$  edges. Since  $t \leq 1/50$  this is less than  $n^2/50$ , so we have a contradiction. This completes the analysis of the first case.

**Case 2.**  $t < 1/135$ .

Let  $e \geq n^2/5$  be the number of edges in  $G$ . By the Cauchy–Schwartz inequality, we have

$$\frac{1}{n} \sum_v \sum_{u \in N(v)} d(u) = \frac{1}{n} \sum_v d(v)^2 \geq \left( \frac{\sum_v d(v)}{n} \right)^2 = 4e^2/n^2.$$

Moreover equality holds only if all degrees in  $G$  are exactly  $2e/n \geq 2n/5$ . In this case we can apply the theorem of Andrásfai, Erdős and Sós mentioned earlier.  $G$  cannot be bipartite, as then one of its parts would have size at least  $n/2$  and contain no edges at all! It follows that  $G = C_5(n/5)$ , as required. We can assume then that equality does not hold, so there is a vertex  $v$  for which  $\sum_{u \in N(v)} d(u) > 4e^2/n^2$ . Let  $A = N(v)$  (an independent set) and suppose  $|A| = (2/5 + s)n$ , where  $s \leq t < 1/135$ .

There is some vertex  $a \in A$  with  $d(a) \geq (2/5 - s)n$ . Otherwise we would have

$$\sum_{a \in A} d(a) < (2/5 + s)n \cdot (2/5 - s)n = (4/25 - s^2)n^2 \leq 4e^2/n^2,$$

a contradiction. Let  $B$  be a subset of  $N(a)$  of size  $(2/5 - s)n$  and let  $C$  be the set of vertices not in  $A \cup B$ . By definition,  $B$  is an independent set disjoint from  $A$  and  $|C| = n - |A| - |B| = n/5$ .

By our construction the number of edges between  $A$  and  $B \cup C$  equals  $\sum_{a \in A} d(a) > 4e^2/n^2$  and therefore the set  $B \cup C$  spans less than  $e - 4e^2/n^2$  edges. Since  $e \geq n^2/5$  and  $f(t) = t - 4t^2$  is a decreasing function for  $t \geq 1/5$  we have that the number of edges of  $G[B \cup C]$  is less than  $n^2/25$ . Then, by the second averaging argument mentioned in the previous section, there is some subset  $X \subseteq C$  of size  $n/2 - |B| = (1/10 + s)n$  incident to less than  $(n/2 - |B|)/|C| \cdot n^2/25 = (1/2 + 5s)n^2/25$  edges of  $G[B \cup C]$ . Since  $B$  is an independent set, we have that  $1/50 \leq n^{-2}e(B \cup X) < 1/50 + s/5$  and so  $s > 0$ .

Suppose that there are  $\beta n^2$  edges with both endpoints in  $C$ . Then  $e(C, B) \leq (1/25 - \beta)n^2$  and taking  $X$  to be a random subset of  $C$  of size  $(1/10 + s)n$  we can improve the previous computation as follows:

$$\begin{aligned} 1/50 \leq n^{-2}e(B \cup X) &\leq ((1/2 + 5s)e(C, B) + (1/2 + 5s)^2e(C))/n^2 \\ &\leq (1/2 + 5s)(1/25 - \beta) + (1/2 + 5s)^2\beta \\ &= (1/2 + 5s)/25 - (1/4 - 25s^2)\beta. \end{aligned}$$

Since  $s < 1/135$ , this gives  $\beta \leq (4s/5)/(1 - 100s^2) \leq 81s/100$ .

Let  $K$  be the set of vertices in  $C$  with at least  $(1/5 - s)n$  neighbours in  $A$ , and let  $L = C \setminus K$ . Recalling the observation made in the third paragraph of the proof, we see that every vertex in  $K$  has at most  $(1/10 - s)n$  neighbours outside  $A$ .

Suppose that  $|K| \geq 4sn$ . Let  $K' \subseteq K$  be a set of size  $4sn$ . Suppose there are  $\alpha n^2$  edges of  $G[B \cup C]$  incident to  $K'$ , where we have  $\alpha \leq 4s(1/10 - s)$ . There is some set  $K'' \subseteq C \setminus K'$  of size  $n/2 - |B \cup K'|$  incident to at most  $(n/2 - |B \cup K'|)/|C \setminus K'| \cdot (1/25 - \alpha)n^2$  edges of  $G[B \cup C]$ . Then

$$\begin{aligned} 1/50 \leq n^{-2}e(B \cup K' \cup K'') &\leq \alpha + \frac{1/10 - 3s}{1/5 - 4s}(1/25 - \alpha) \\ &\leq \alpha + (1/2 - 5s)(1/25 - \alpha) = 1/50 - s/5 + (1/2 + 5s)\alpha \\ &\leq 1/50 - s/5 + (1/2 + 5s) \cdot 4s(1/10 - s) \\ &= 1/50 - 20s^3. \end{aligned}$$

This contradiction shows that  $|K| < 4sn$ .

Recall that, by definition of  $K$ , each vertex in  $L = C \setminus K$  has less than  $(1/5 - s)n$  neighbours in  $A$ . Since  $s < 1/135$ , we also have that  $|L| = |C| - |K| > (1/5 - 4s)n > (1/10 - s)n$ . By averaging there is some  $L' \subseteq L$  of size  $(1/10 - s)n$  so that  $e(A \cup L') \leq |L'|(1/5 - s)n + (|L'|/|L|)^2e(C)$ . Now  $|L'|/|L| < (1/10 - s)/(1/5 - 4s) < 1/2 + 10s$  and  $(1/2 + 10s)^2 < 1/4 + 11s$ . Recalling that  $n^{-2}e(C) = \beta < 81s/100$  we have

$$\begin{aligned} 1/50 \leq n^{-2}e(A \cup L') &\leq (1/10 - s)(1/5 - s) + (1/4 + 11s) \cdot 81s/100 \\ &\leq 1/50 - 39s/400 + 10s^2. \end{aligned}$$

This gives  $10s^2 > 39s/400 > s/11$ , i.e.,  $s > 1/110$ , a contradiction that completes the proof.

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