



Sparse Indexing: Large-Scale, Inline Deduplication Using Sampling and Locality

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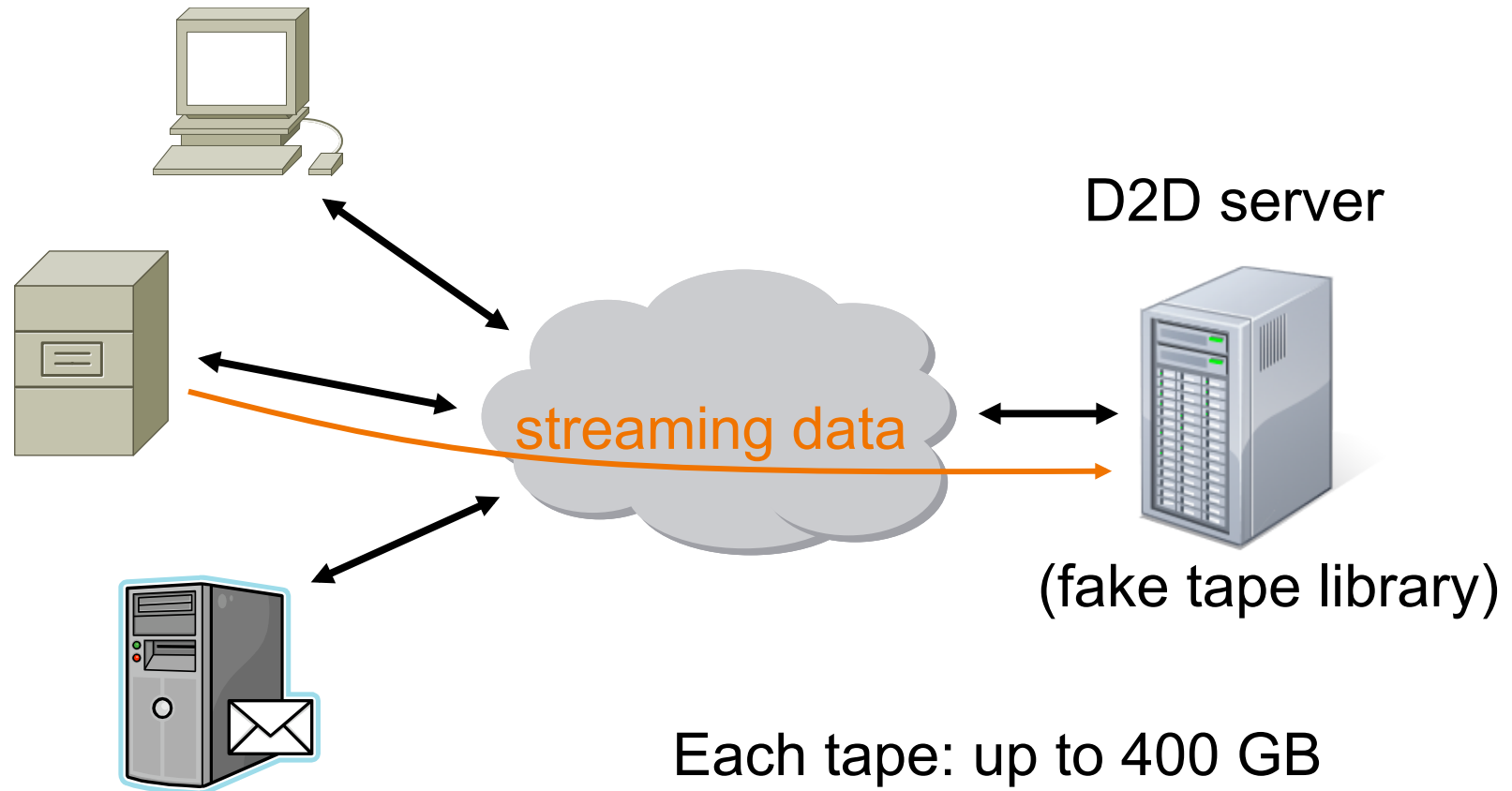
Work done at Hewlett-Packard laboratories



The Problem:

deduplication at scale
for disk-to-disk backup

A Disk-to-Disk Backup Scenario



Each tape: up to 400 GB

Total: 100 TB – 10 PB

Example backup streams

Monday:



Tuesday:



Wednesday:



Little changes from day-to-day

Monday:



Tuesday:



Wednesday:



After ideal deduplication

Monday:



Tuesday:



Wednesday:



Chunk-based deduplication

Monday:



Tuesday:



Wednesday:



Chunk-based deduplication

Monday:



Tuesday:



Wednesday:



Chunk-based deduplication

Monday:



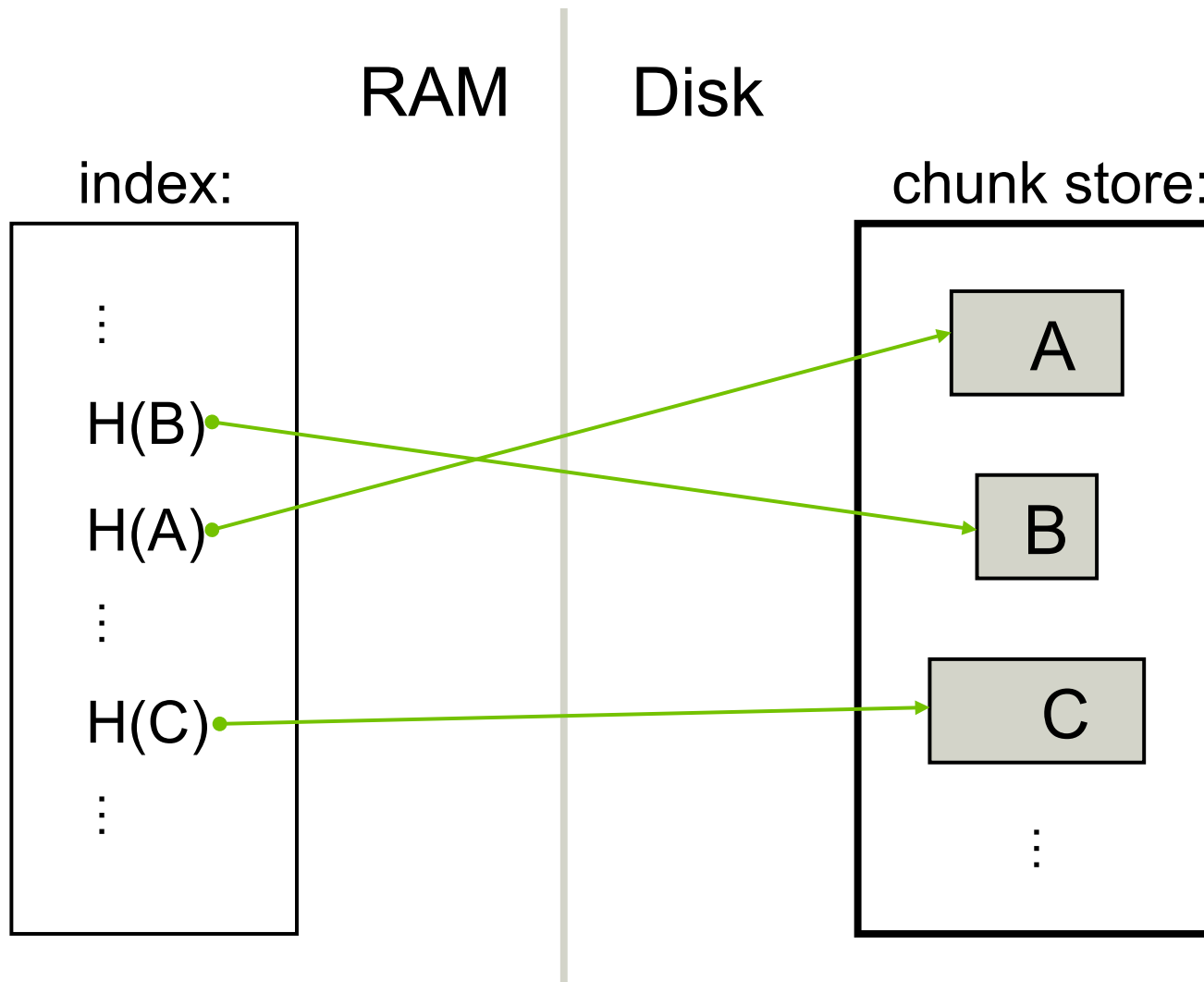
Tuesday:



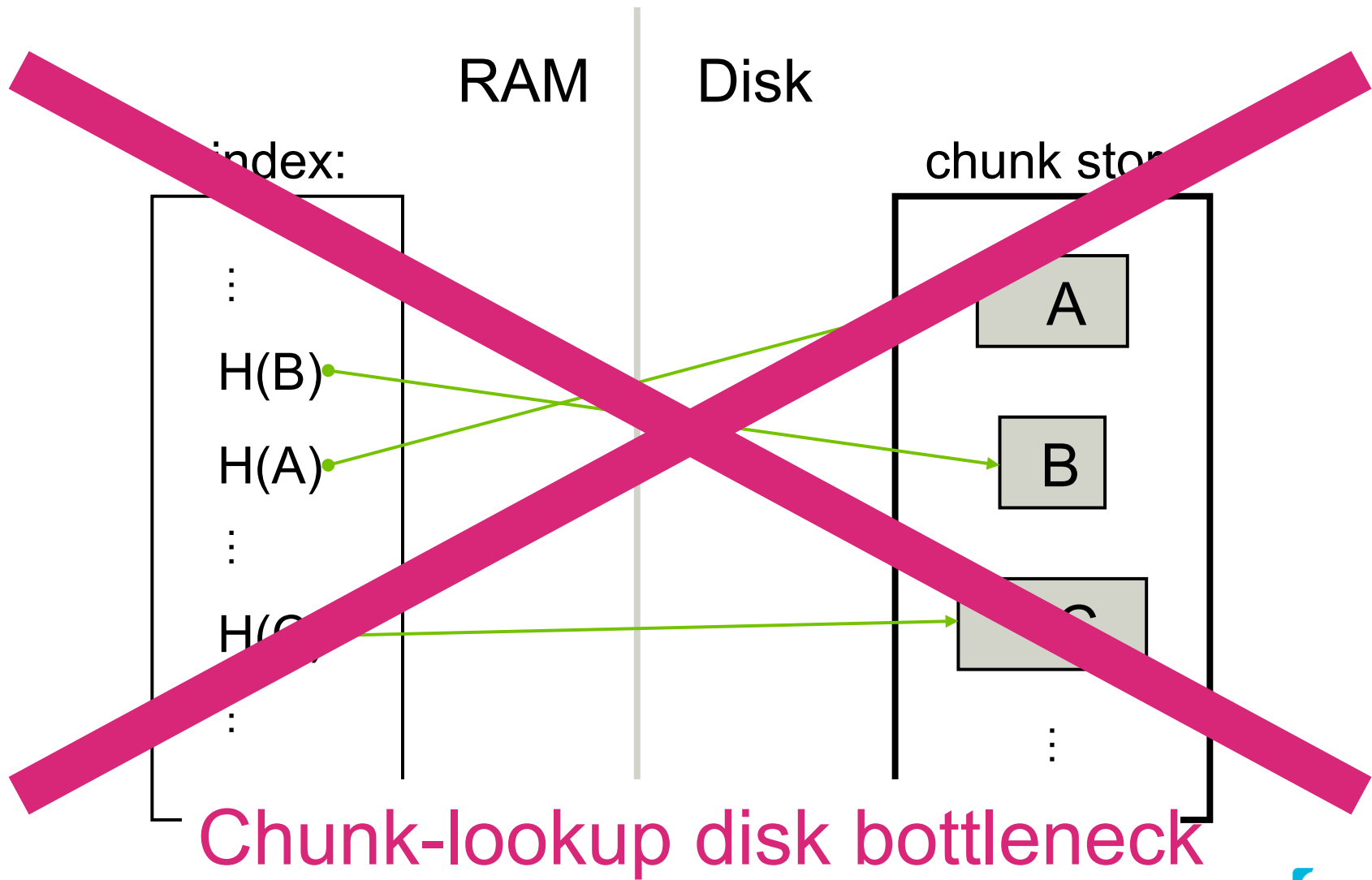
Wednesday:



The standard implementation



The standard implementation



Chunk-lookup disk bottleneck

One existing solution

- **Avoiding the Disk Bottleneck in the Data Domain Deduplication File System.** Benjamin Zhu, *Data Domain, Inc.*; Kai Li, *Data Domain, Inc.*, and Princeton University; Hugo Patterson, *Data Domain, Inc.* FAST'08.
- Today: a new approach that
 - uses significantly less RAM
 - provides a guaranteed minimum throughput

Our Approach:

Sparse indexing

Sparse indexing

- Key ideas:
 - Chunk locality
 - Sampling

No temporal locality

Monday:



Tuesday:



Wednesday:



Large sections of data reappear mostly intact → chunk locality

Monday:



Tuesday:



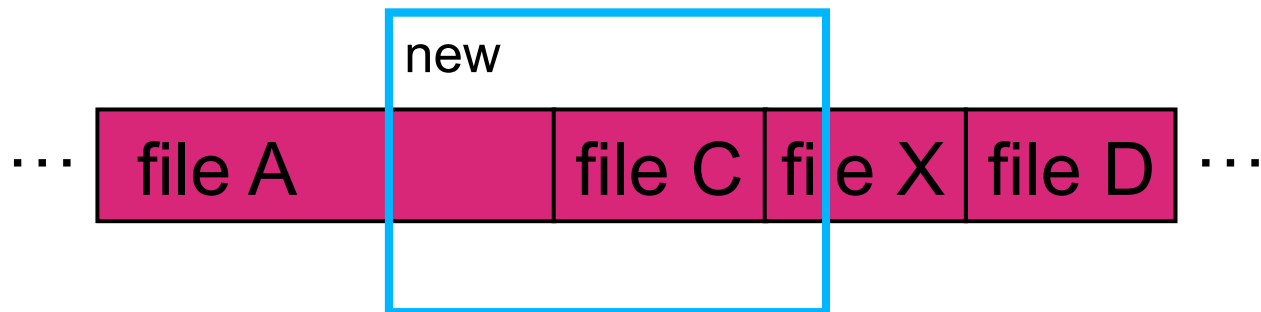
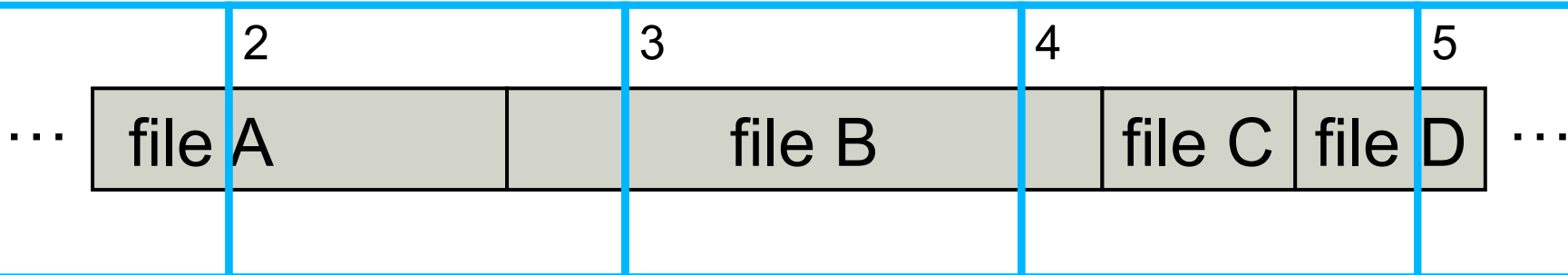
Wednesday:



Exploiting chunk locality

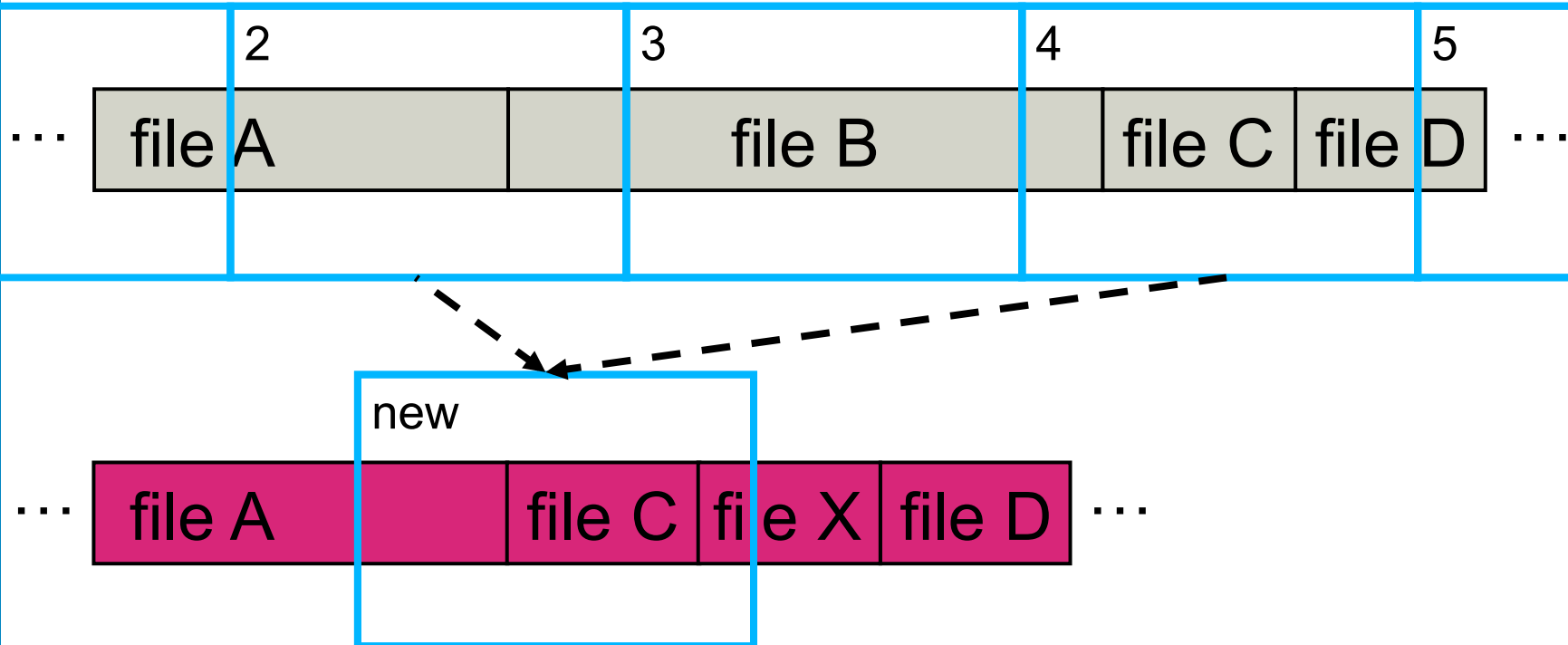


Divide into *segments*



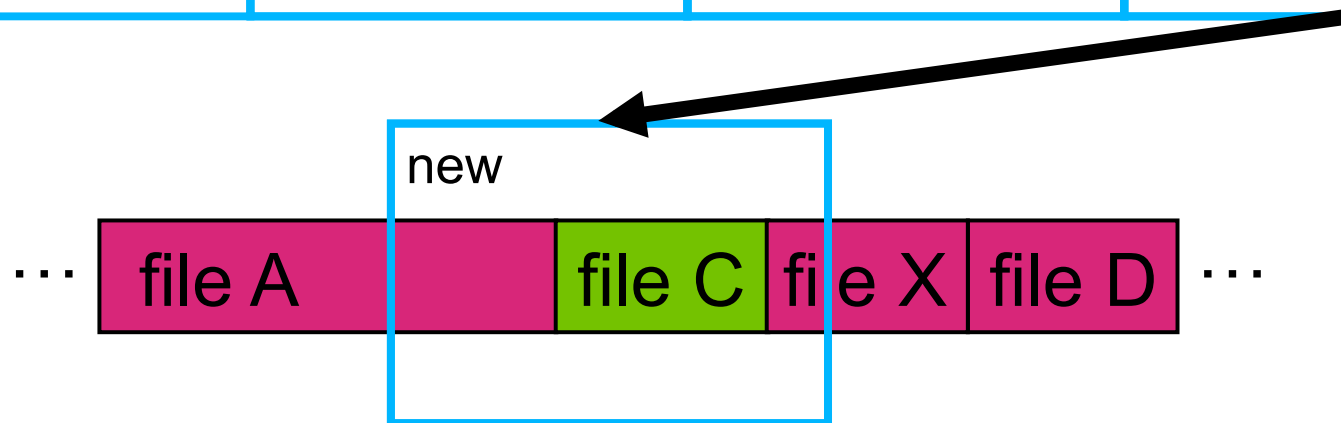
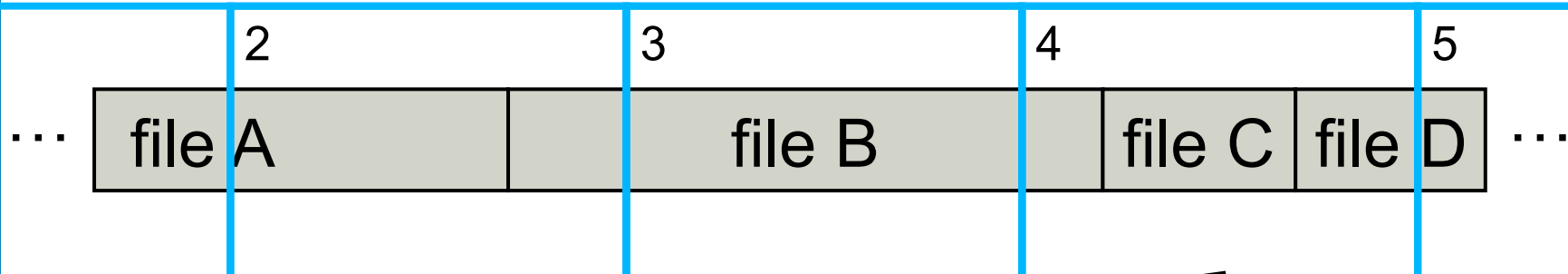
Chunks not shown, real segments much longer

Deduplicate one segment at a time

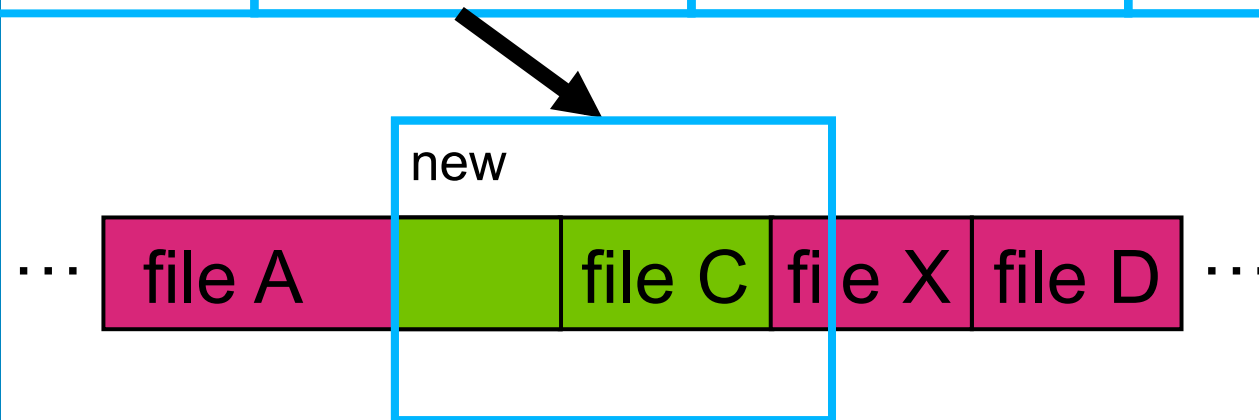
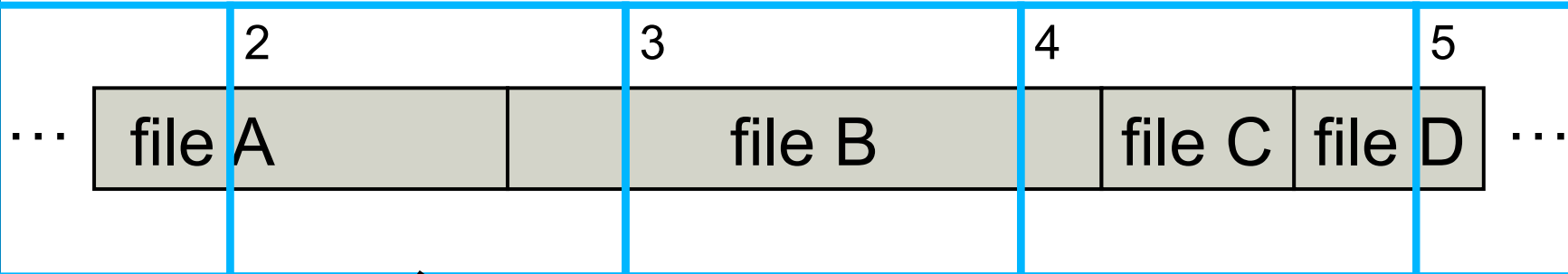


Against a few carefully chosen champion segments

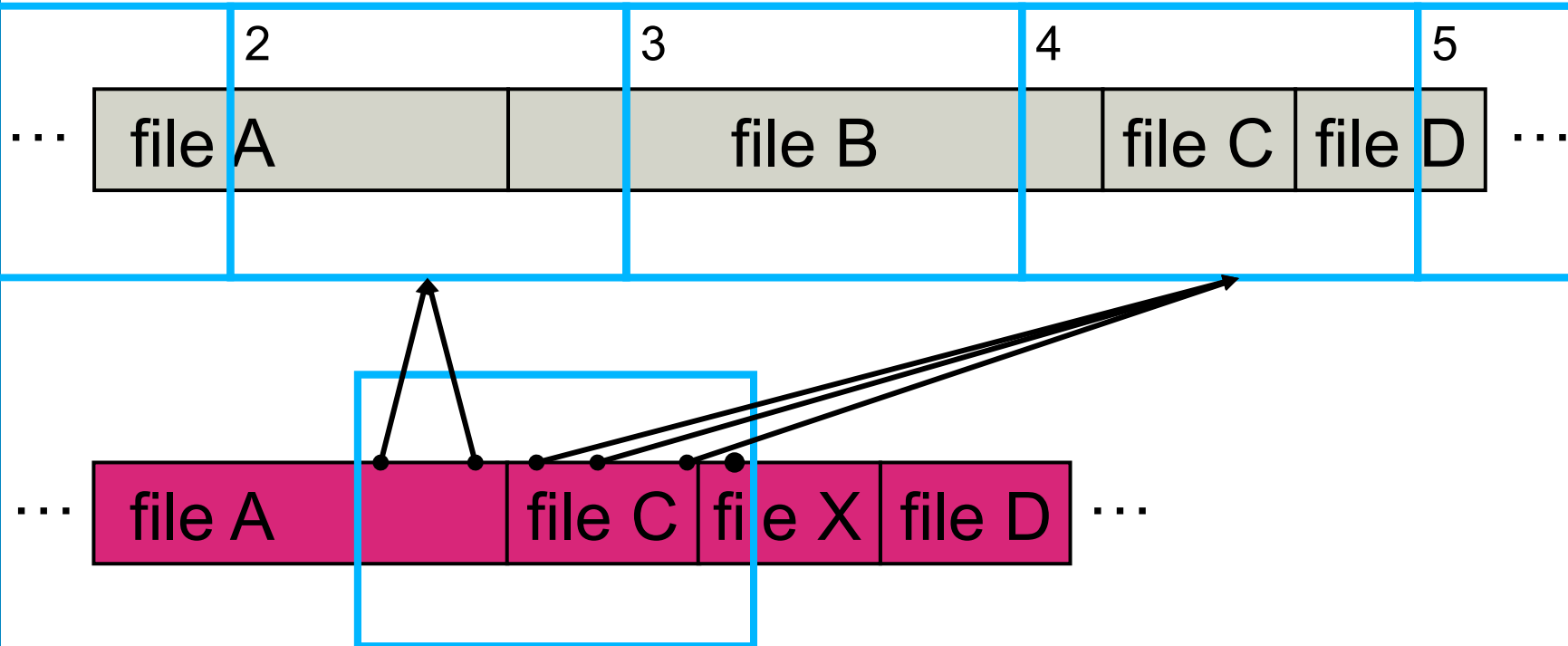
Champion #1: the most similar segment



Champion #2: most similar to remainder



Finding similar segments by sampling

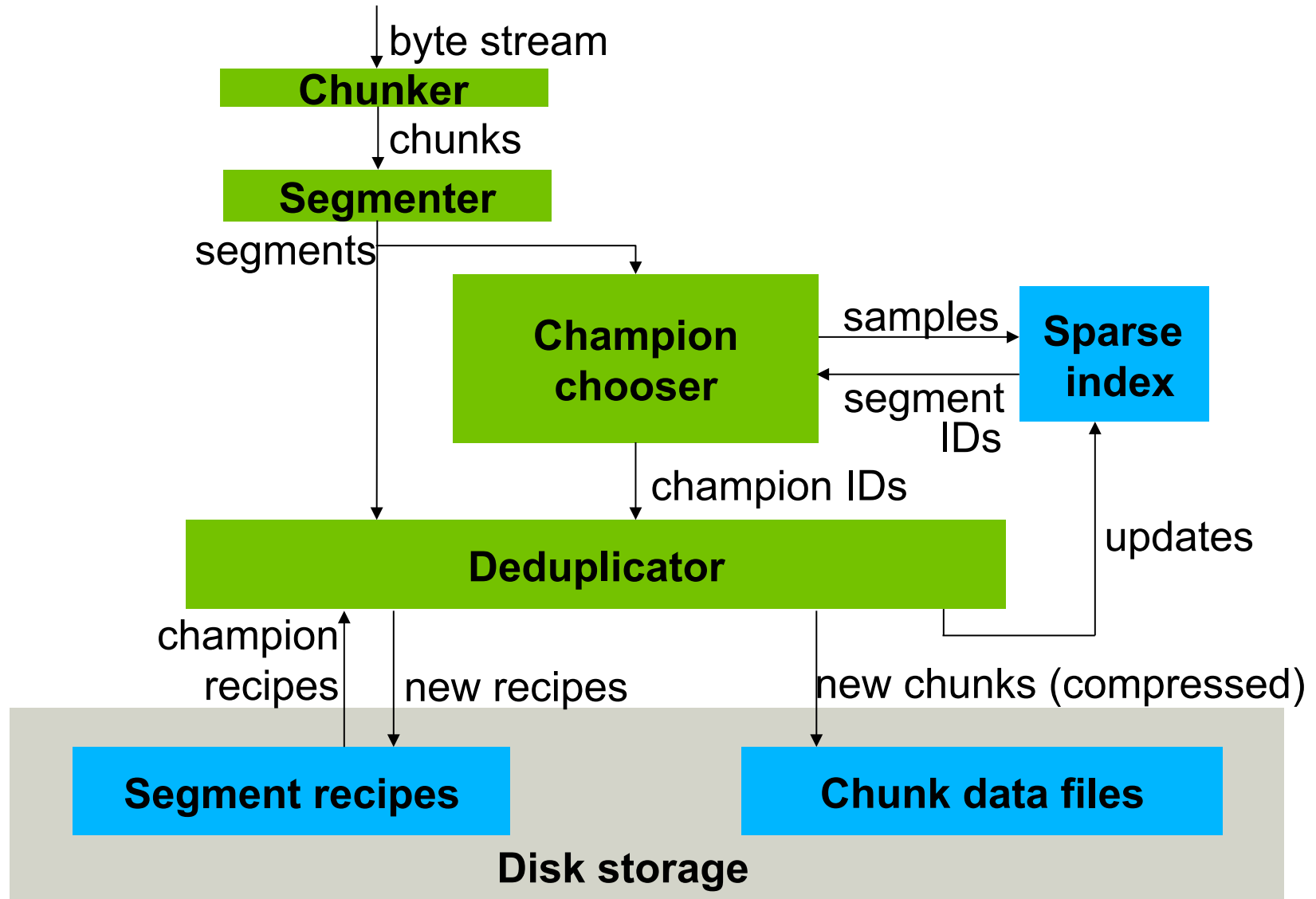


Sparse Index: samples $\bullet \rightarrow$ containing segment(s)

A few details

- Also keep segment *recipes*:
 - list of pointers to a segment's chunks
- Actually deduplicate against champion recipes
- Better with variable-sized segments
 - boundaries based on landmarks (“superchunks”)
 - reduces number of champions required

Putting it all together



Results

Methodology

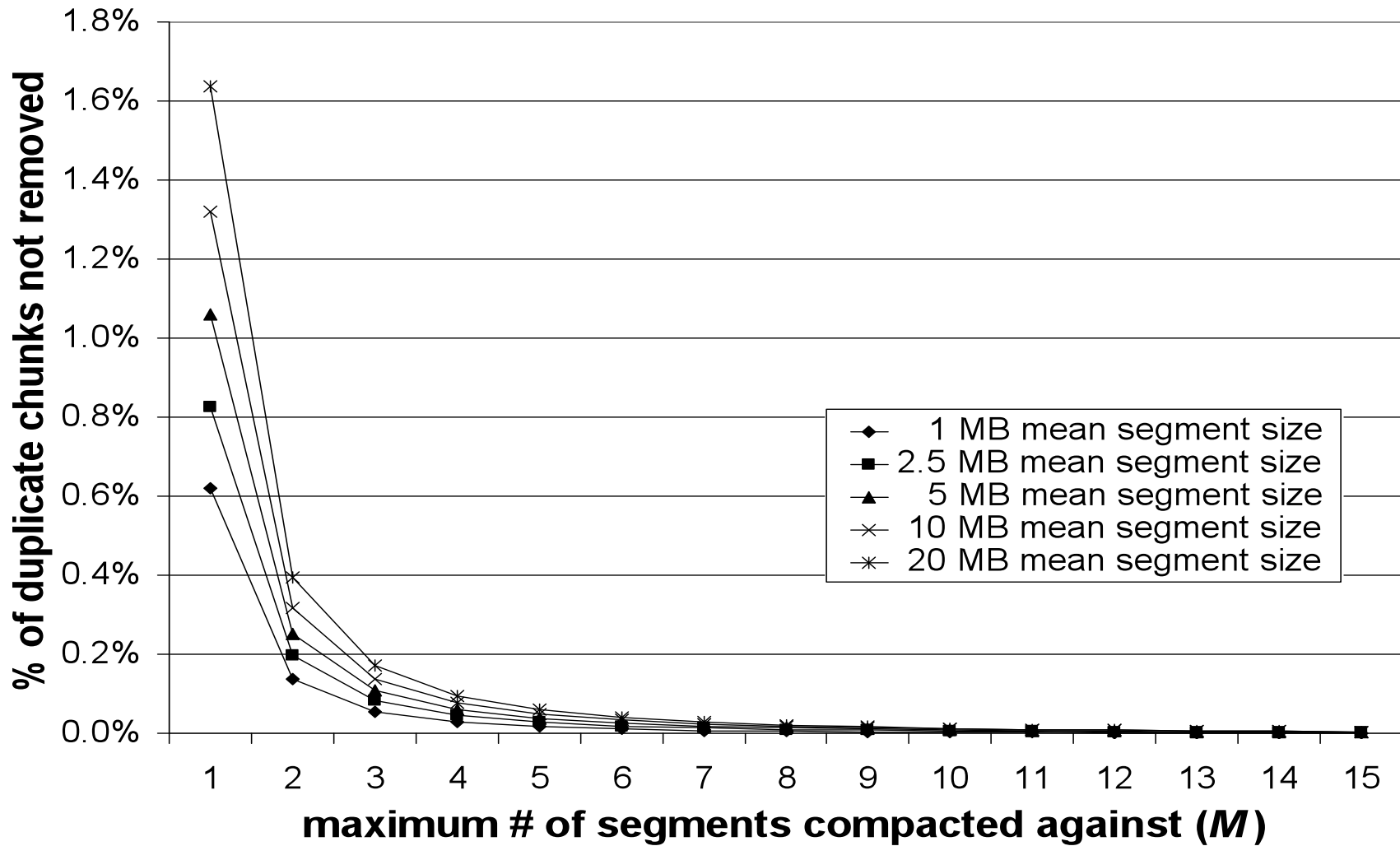
- Built a simulator
- Fixed parameters:
 - 4 KB mean chunk size
 - variable-size segments
 - maximum of 1 segment ID kept per sample
- Varying parameters:
 - mean segment size
 - sampling rate
 - maximum number of champions per segment (M)

The data sets

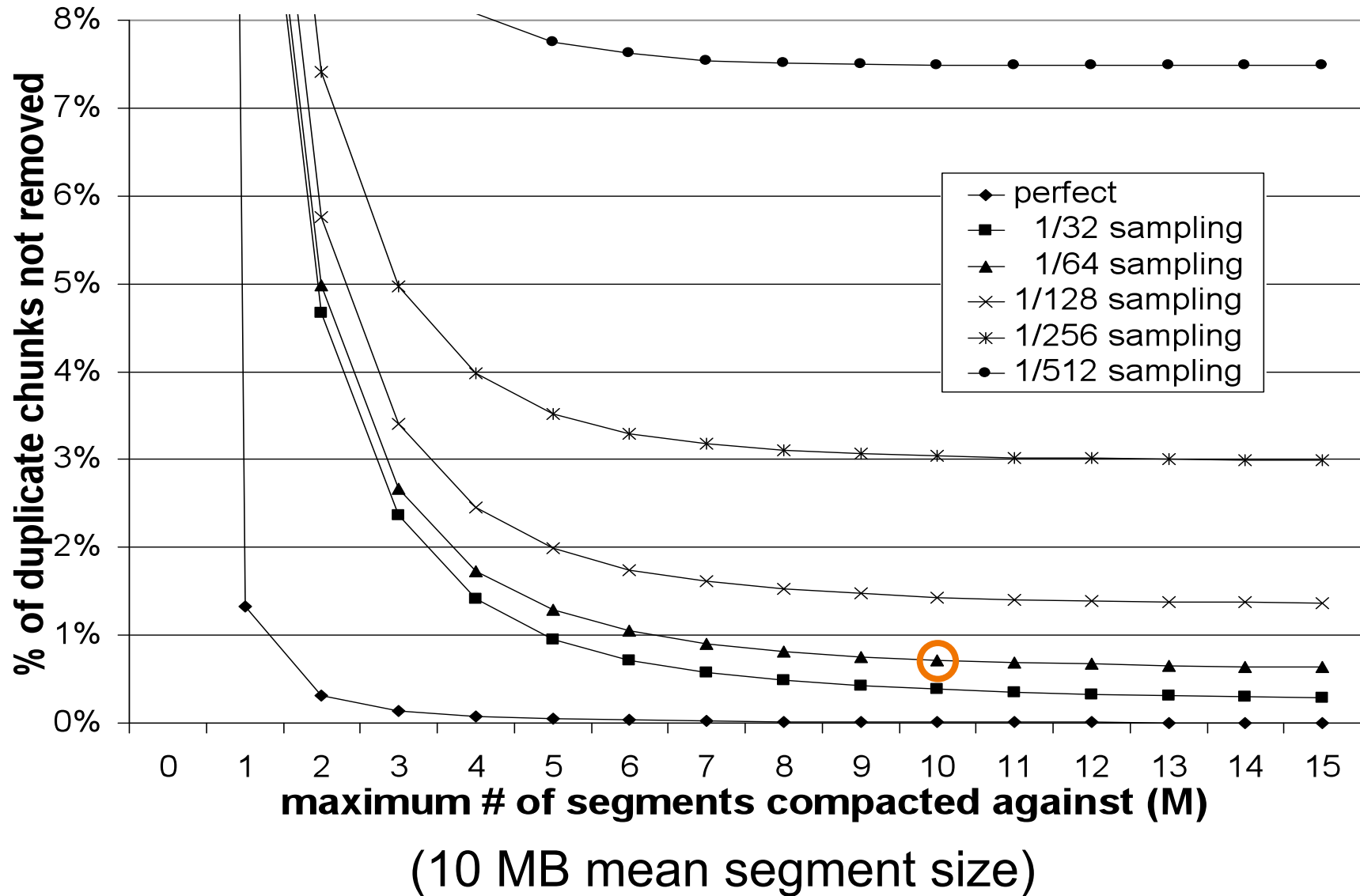
- **Workgroup** [this talk] 3.8 TB
 - backups of 20 desktop PCs belonging to engineers
 - semi-regular backups over 3 months via tar
 - 154 full backups and 392 incremental backups
 - end-of-week full backups are synthetic

- **SMB** [see paper] 0.6 TB
 - backups of a server with
 - real Oracle data
 - synthetic Microsoft Exchange data
 - two weeks

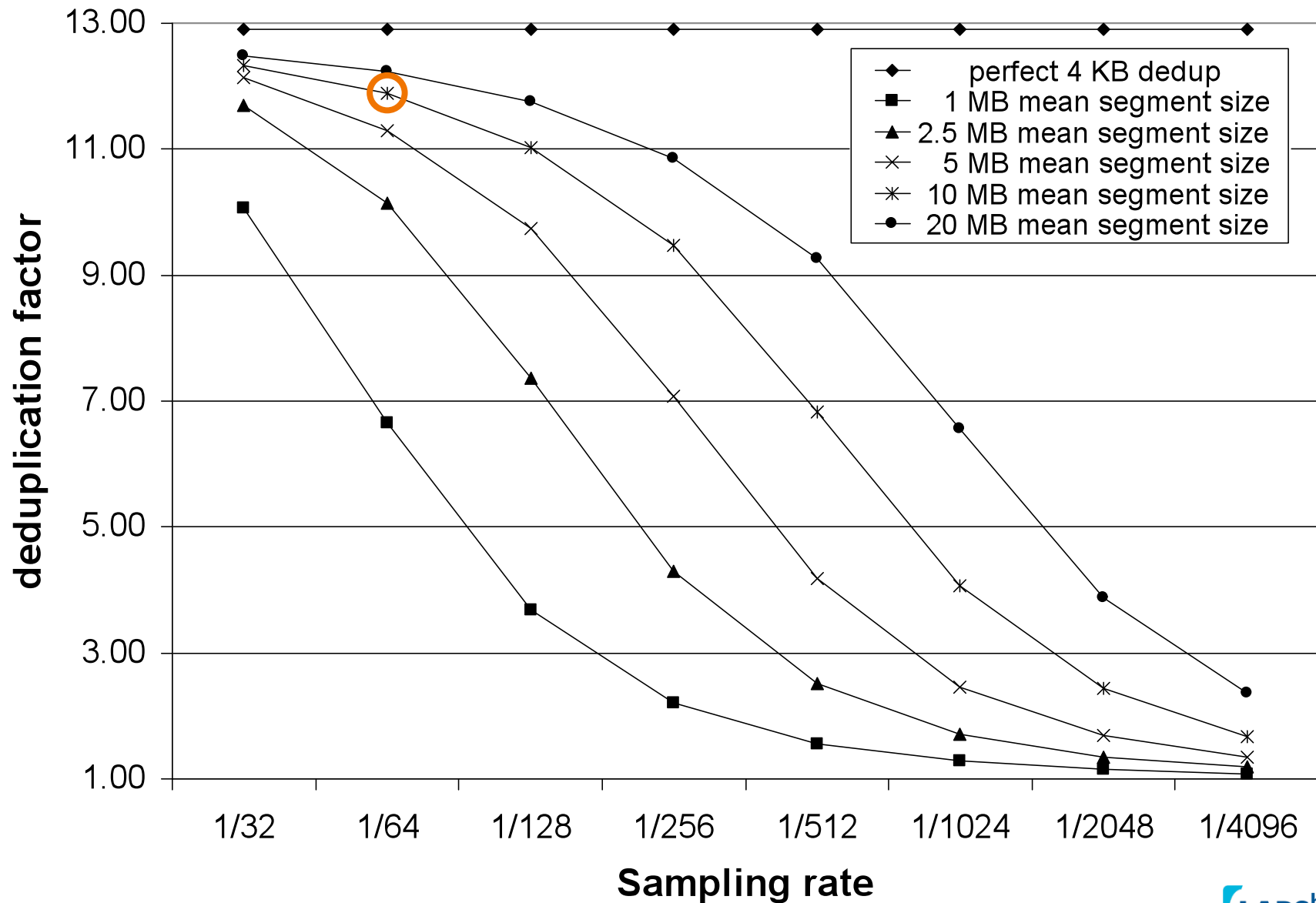
Chunk locality exists...



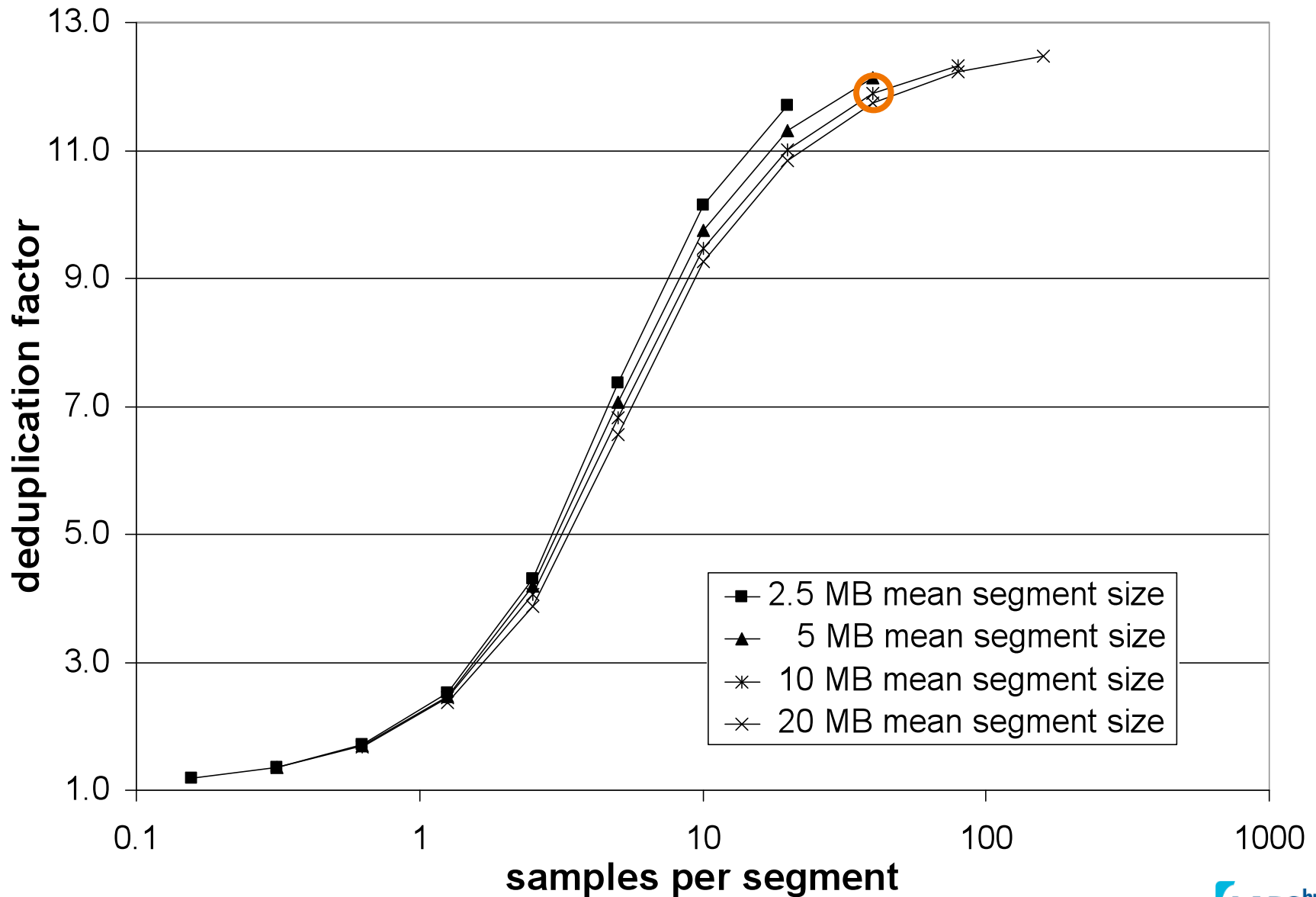
Sampling can exploit most of it...



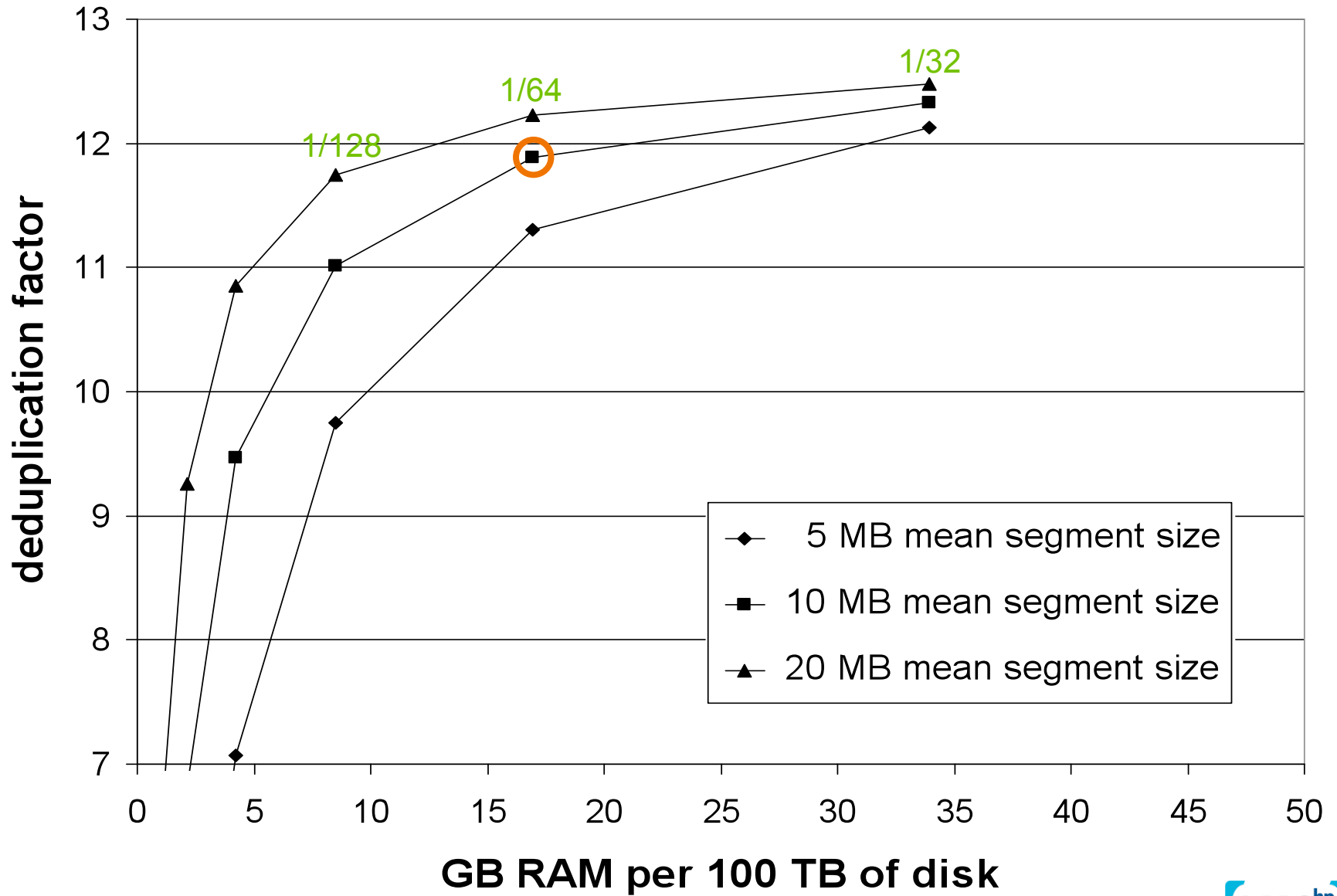
Deduplication with at most 10 champions



Deduplication depends primarily on...



Index RAM usage



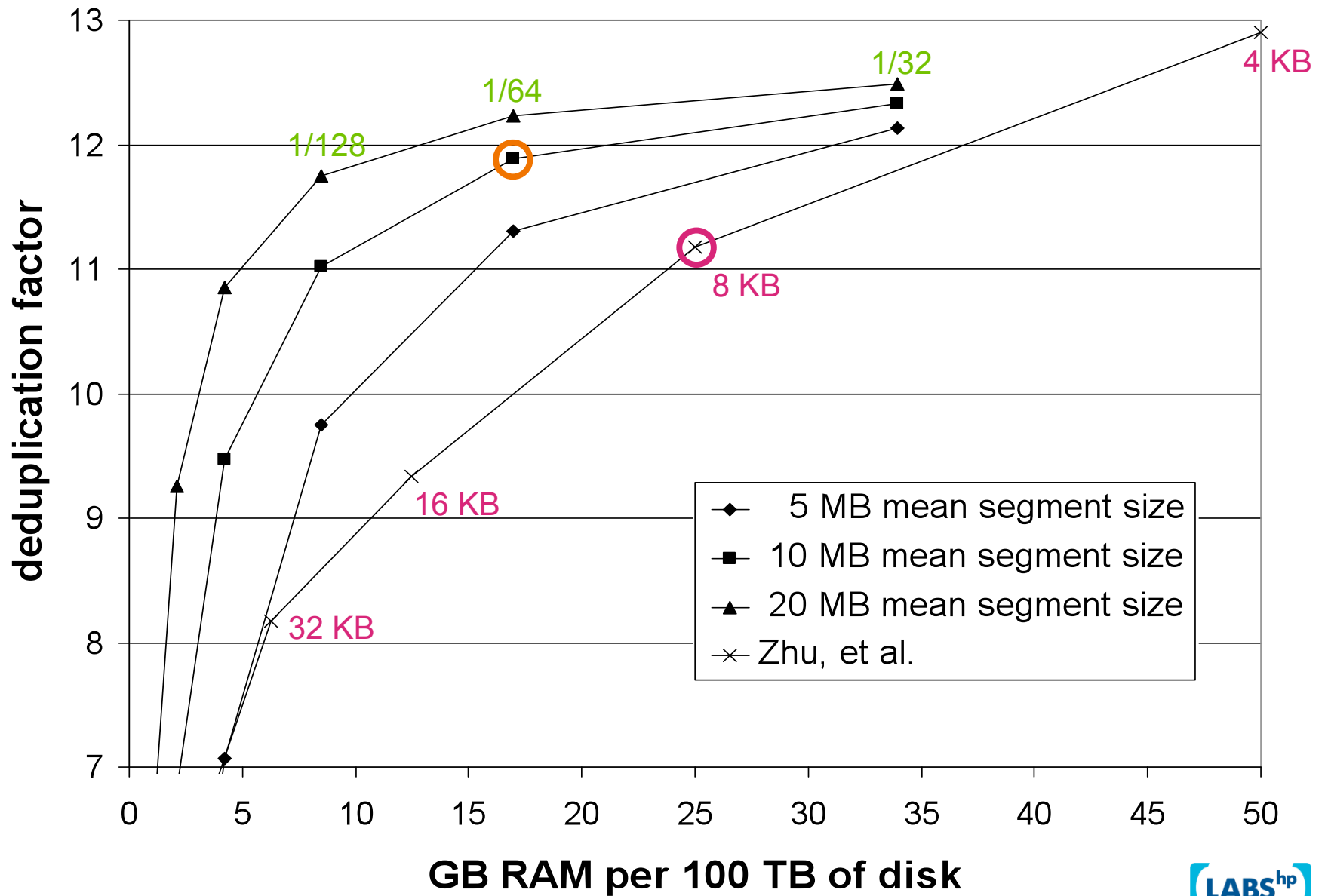
Comparison with Zhu, et al.

- Their chunk lookup:
 - bloom filter: might the store have a copy?
 - cache of chunk container indexes
 - full on disk index

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- Their chunk lookup:
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- When chunk locality is poor,
 - deduplication quality remains constant
 - but throughput degrades
- Find all duplicate chunks
 - but larger chunk size

Ram usage comparison



What about all those disk accesses?

- Infrequent due to batch processing
- **Example:**
 - load at most **10** champions per **10 MB** segment
 - average of 1.7 champions per 10 MB segment
 - = 0.17 champions/MB
 - = 1 seek per 5 MB
- I/O burden:
 - 20 ms to load a champion recipe (~100 KB)
 - → 1 drive can handle > 250 MB/s ingestion rate

Thank You