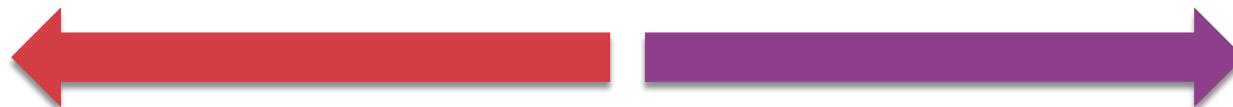


THE BEAUTY AND THE BEAST: SEPARATING DESIGN FROM ALGORITHMS

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Modularization of a program



Overdesign

Lack of modularization

- Do I extract this code into a separate method?
- Do I implement this traversal iteratively or recursively?
- Do I use a visitor pattern or do I place the computation in the structure itself? ...

A good modularization ease understandability and maintainability

Modularization of a program

```
private static int fac(int n) {  
    int f=1;  
    for (int i=1; i<=n; i++) {  
        int p = 0;  
        for (int j=1; j<=i; j++) {  
            p=p+f;  
        }  
        f=p;  
    }  
    return f ;  
}
```

```
private static int fac(int n) {  
    int f=1;  
    for (int i=1; i<=n; i++) {  
        f=mul(f, i);  
    }  
    return f ;  
}  
  
private static int mul(int a, int b) {  
    int p = 0;  
    for (int j=1; j<=a; j++) {  
        p=p+b;  
    }  
    return p ;  
}
```

Modularization of a program

```
private static int fac(int n) {  
    int f=1;  
    for (int i=1; lessOrEqual(i, n); i=addOne(i)) {f=mul(f, i);}  
    return f ;  
}  
private static int mul(int a, int b){  
    int p = 0;  
    for (int j=1; lessOrEqual(j, a); j=addOne(j)) {p=add(p, b);}  
    return p;  
}  
private static int add(int a, int b) { return a+b; }  
private static boolean lessOrEqual ( int a , int b ) { return a<=b; }  
private static int addOne(int a) { return add(a, 1); }
```

Aims of the paper

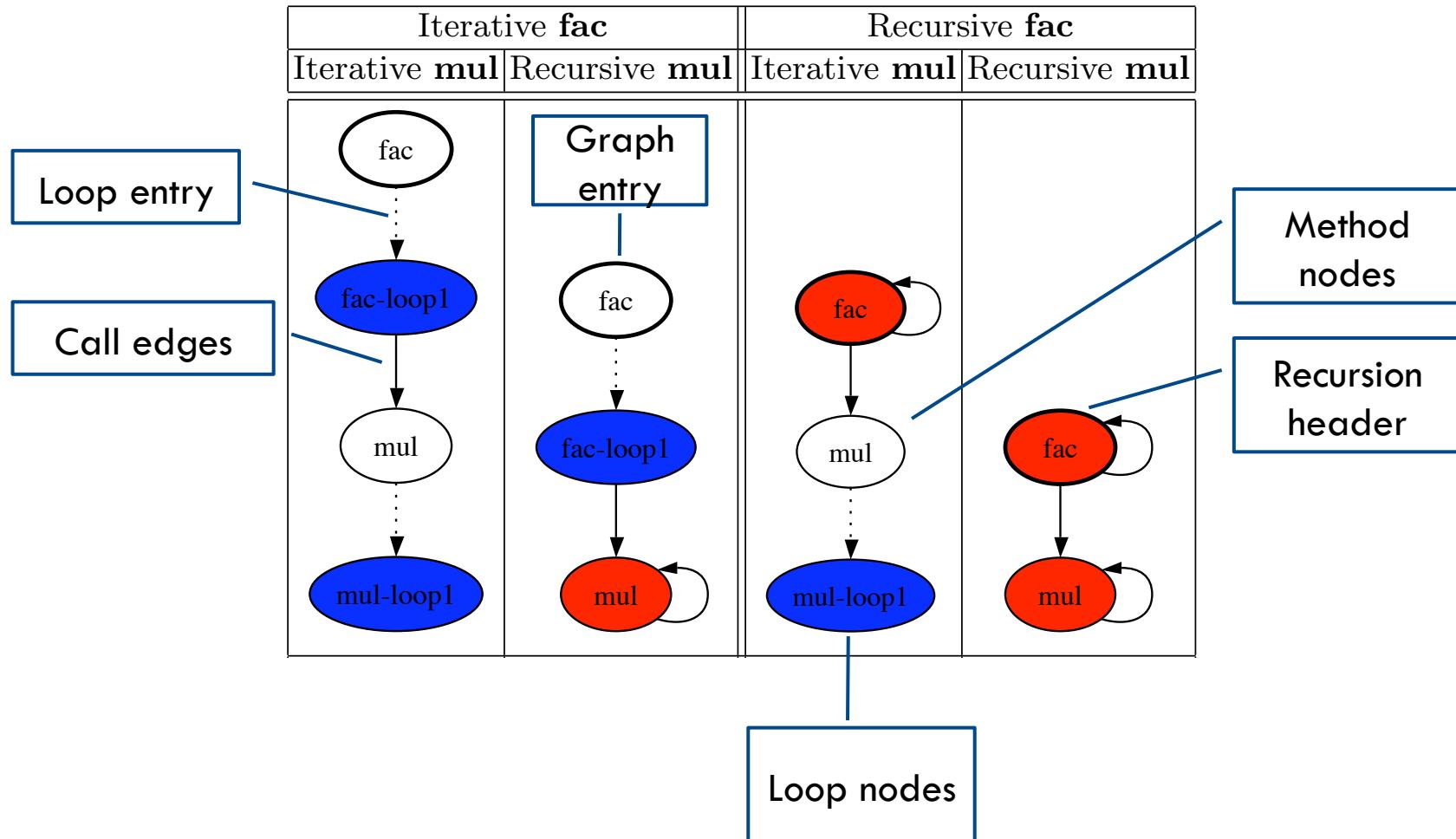
- Propose a metric to quantify the algorithmically essential parts of a program
 - ▣ Localizable
 - ▣ Intuitive
 - ▣ Stable
 - ▣ Language independent
- Consider loops and recursive calls
- The metric is computed using three different representations
 - ▣ Call graph
 - ▣ Control-flow
 - ▣ Loop call graphs

Overall approach



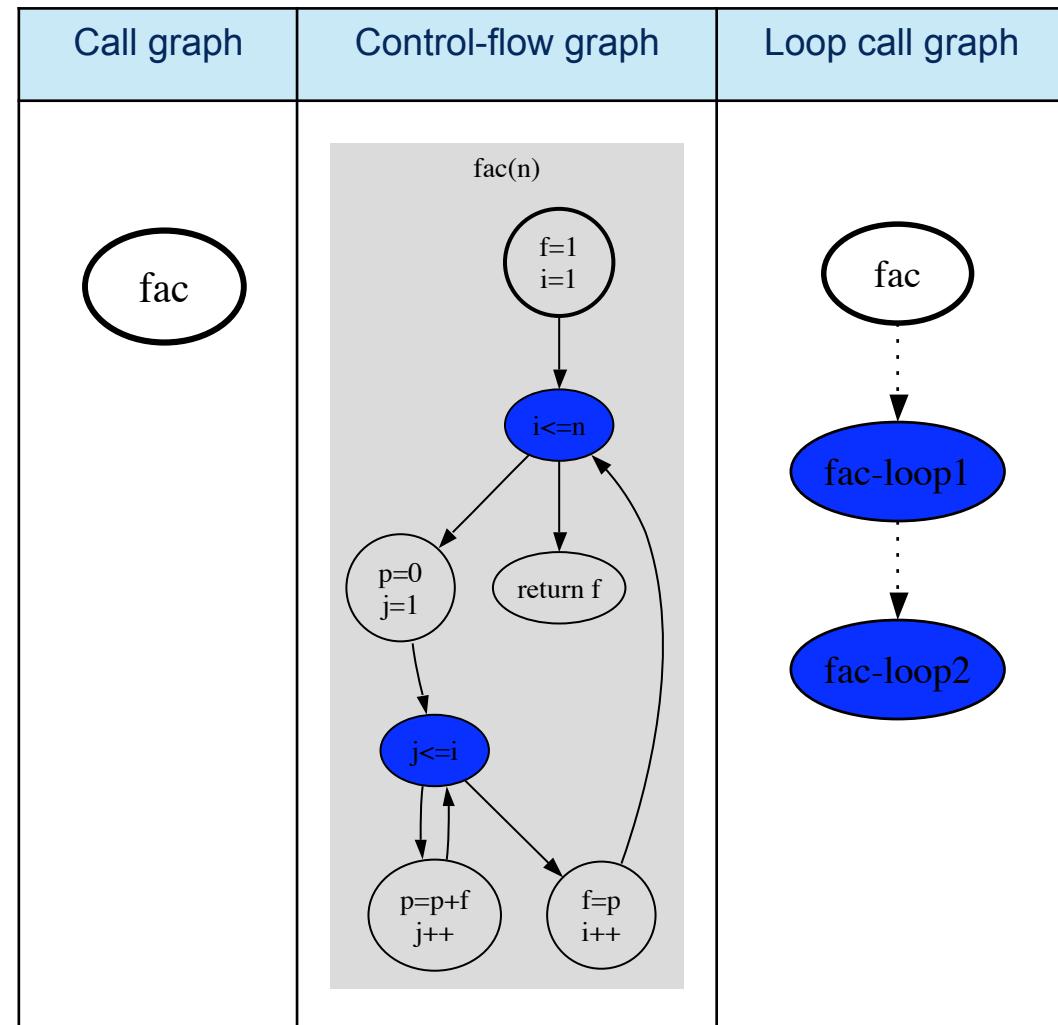
1. Build a control-flow graph
2. Identify loop forests in control-flow graphs
3. Build call graph
4. Identify recursion forests in call graph
5. Combine loop forests & call graph into loop call graph
6. Compute metrics

Loop call graph



Example 1

```
private static int fac(int n) {  
    int f=1;  
    for (int i=1; i<=n; i++) {  
        int p = 0;  
        for (int j=1; j<=i; j++) {  
            p=p+f;  
        }  
        f=p;  
    }  
    return f ;  
}
```



Example 2

```

private static int fac(int n) {
    int f=1;
    for (int i=1; i<=n; i++) {
        f=mul(f, i);
    }
    return f ;
}

private static int mul(int a, int b) {
    int p = 0;
    for (int j=1; j<=a; j++) {
        p=p+b;
    }
    return p ;
}

```

Call graph	Control-flow graph	Loop call graph
<pre> graph TD fac((fac)) --> mul((mul)) </pre>	<p>fac(n)</p> <pre> graph TD f1i1((f=1 i=1)) --> iLeqN((i<=n)) iLeqN --> fMul((f=mul(f, i) i++)) fMul --> returnF((return f)) f1i1 <--> iLeqN f1i1 <--> fMul </pre> <p>mul(a,b)</p> <pre> graph TD p0j1((p=0 j=1)) --> jLeqA((j<=a)) jLeqA --> ppb((p=p+b j++)) ppb --> returnP((return p)) p0j1 <--> jLeqA p0j1 <--> ppb </pre>	<pre> graph TD fac((fac)) -.-> facLoop1((fac-loop1)) mul((mul)) --> mulLoop1((mul-loop1)) </pre>

Example 3

```
private static int fac(int n) {  
    int f=1;  
    for (int i=1; lessOrEqual(i, n); i=addOne(i)) {f=mul(f, i);}  
    return f ;  
}  
  
private static int mul(int a, int b){  
    int p = 0;  
    for (int j=1; lessOrEqual(j, a); j=addOne(j)) {p=add(p, b);}  
    return p;  
}  
  
private static int add(int a, int b) { return a+b; }  
private static boolean lessOrEqual ( int a , int b) { return a<=b;}  
private static int addOne(int a) { return add(a, 1); }
```

Example 3

Call graph	Control-flow graph	Loop call graph
<pre> graph TD fac --> mul fac --> inc mul --> inc mul --> leq inc --> add leq --> add </pre>	<p>fac(n)</p> <pre> graph TD f1i1((f=1 i=1)) --> leqin((leq(i,n))) leqin --> fmul((f=mul(f,i) i=inc(i))) fmul --> returnf((return f)) </pre> <p>mul(a,b)</p> <pre> graph TD p0j1((p=0 j=1)) --> leqja((leq(j,a))) leqja --> pab((p=add(p,b) j=inc(j))) pab --> returnp((return p)) </pre> <p>add(a,b)</p> <pre> graph TD rab((r=a+b return r)) </pre> <p>inc(a)</p> <pre> graph TD raa1((r=add(a,1) return r)) </pre> <p>leq(a,b)</p> <pre> graph TD rab((r=a<=b return r)) </pre>	<p>fac</p> <p>fac-loop1</p> <p>mul</p> <p>mul-loop1</p> <p>inc</p> <p>leq</p> <p>add</p> <pre> graph TD fac -.-> facloop1 facloop1 --> mul mul --> mulloop1 mulloop1 --> inc mulloop1 --> leq inc --> add leq --> add </pre>

Metrics

NN = |non-recursive method nodes|

NR = |recursive method nodes|

NL = |loop nodes|

$$E = NL + NR$$

$$\text{recursiveness} = \frac{NR}{NN + NR}$$

$$\text{loopyness} = \frac{NL}{NN + NR}$$

$$e = \frac{E}{NN + NR}$$

'e' for previous examples:

- Example 1 = 2
- Example 2 = 1
- Example 3 = 0.4

Iterative fac		Recursive fac	
Iterative mul	Recursive mul	Iterative mul	Recursive mul
<pre> graph TD fac((fac)) --> facLoop1([fac-loop1]) facLoop1 --> mul1((mul)) mul1 --> mulLoop1([mul-loop1]) </pre>	<pre> graph TD fac((fac)) --> facLoop1([fac-loop1]) facLoop1 --> mul2((mul)) mul2 --> facLoop1 </pre>	<pre> graph TD fac3((fac)) --> mul3((mul)) mul3 --> mulLoop1([mul-loop1]) </pre>	<pre> graph TD fac4((fac)) --> mul4((mul)) mul4 --> fac4 </pre>
$N_N = 2$ $N_R = 0$ $N_L = 2$ $E = 2$ $e = \frac{2}{2+0} = 1$	$N_N = 1$ $N_R = 1$ $N_L = 1$ $E = 2$ $e = \frac{2}{1+1} = 1$	$N_N = 1$ $N_R = 1$ $N_L = 1$ $E = 2$ $e = \frac{2}{1+1} = 1$	$N_N = 0$ $N_R = 2$ $N_L = 0$ $E = 2$ $e = \frac{2}{0+2} = 1$

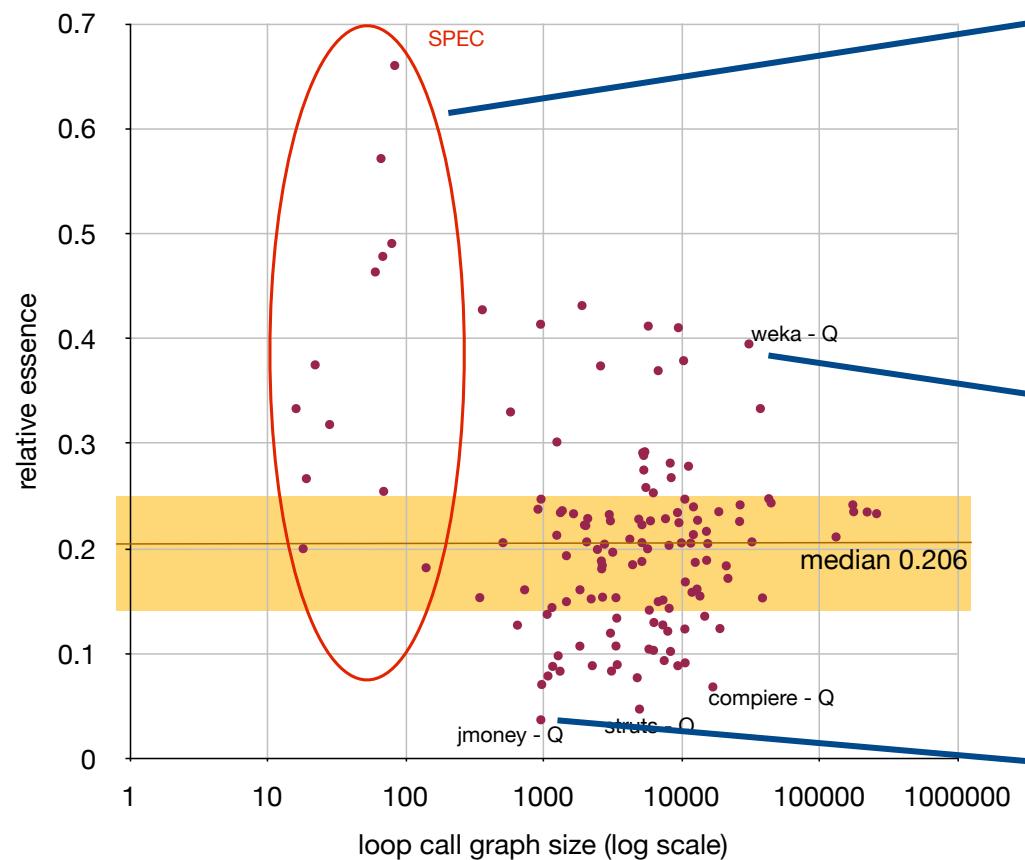
Metrics



- E
 - ▣ It is not affected by the degree of indirection
 - ▣ It does not depend on the implementation (recursive or iterative)

- NN is correlated with the level of indirection

Essence



Small loop-driven algorithm implementations

Machine learning algorithms

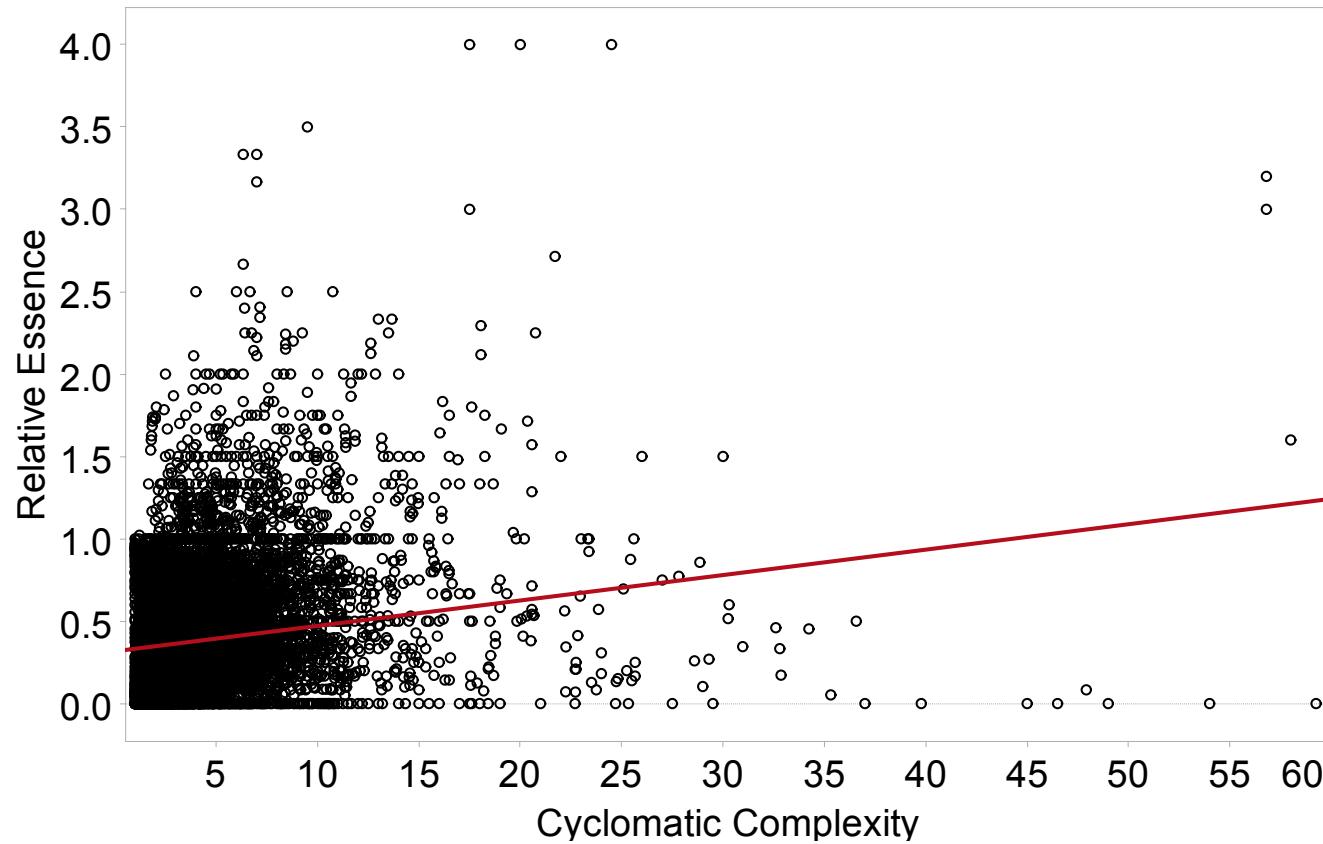
Connects eclipse plugins

Essence & design



- Code smells
 - ▣ In some cases, the essence is too low but the amount of indirection too high
- Refactoring
- Design patterns
 - ▣ This technique successfully identifies that the visitor pattern introduces an extra level indirection for each method

Cyclomatic complexity



Usage scenarios

- Deviation from reference
 - ▣ Compare essence with a reference value of a system of high design quality
- Problem localization
 - ▣ Generated graphs can help to spot nodes with high or low essence
- Refactoring recommendation
 - ▣ Guide modularization of a system
- Quality and process attribute prediction
 - ▣ It might be possible to predict process attributes, such as error rates or times to fix an error

Conclusions



- The presented metric (relative essence) provides hints on which parts of the system to remove, and where to add extra indirections
- Not all recursion and loops are necessarily required for solving the problem the program needs to solve
- This metric relates with design patterns, code smells and refactoring
- None of the existing metrics correlates with relative essence

