# The cohesive principle and the Bolzano-Weierstraß principle

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## The Bolzano-Weierstraß principle

A real number is a sequence of rational numbers with Cauchy-rate  $2^{-n}$ .

## Definition

(BW):

Every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  has a cluster point.

*Equivalently*, every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  contains a Cauchy-subsequence  $(y_n)$  with Cauchy-rate  $2^{-n}$ , i.e. with

$$\forall n \,\forall i,j \ge n \, \left( |y_i - y_j| < 2^{-n} \right).$$

#### Definition

 $(\mathsf{BW}_{\mathsf{weak}})$ :

Every bounded sequence  $(x_n)_n \subseteq \mathbb{R}$  contains a Cauchy-subsequence  $(y_n)_n$ ,

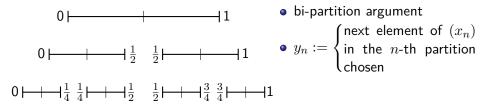
i.e.

$$\forall n \,\exists k \,\forall i, j \geq k \, \left( |y_i - y_j| < 2^{-n} \right).$$

## Computing BW

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ .

Goal: Construct a subsequence  $(y_n)$  with  $\forall n \forall i, j \ge n \ (|y_i - y_j| < 2^{-n})$ .



• The partitions form a  $\Pi_2^0$ -0/1-tree.

• This is a  $\Pi_1^0$ -0/1-tree in 0'.

• WKL relativized to 0' yields an infinite branch and therefore computes the sequence of partitions.

#### Theorem (Kohlenbach '98, Kohlenbach, Safarik '10, K. '10)

- For each computable sequence  $(x_n)$ there is a 0'-computable 0/1-tree T, such that an infinite branch of T computes a cluster point, and vice versa.
- Over RCA<sub>0</sub> the principles BW and WKL for Σ<sup>0</sup><sub>1</sub>-trees are instance-wise equivalent.

By the low basis theorem:

#### Corollary

BW has for computable instances a solution low relative to 0', *i.e.* the first Turing jump of a solution is computable in 0".

## Computing $BW_{weak}$

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ . Goal: Construct a subsequence  $(y_n)$  with  $\forall n \exists k \forall i, j \geq k \ (|y_i - y_j| < 2^{-n})$ and compute the Turing jump of  $(y_n)$ .

It is clear that

$$\Phi_e^{(y_n)_n} \downarrow \quad \text{iff} \quad \exists k \, \Phi_e^{(y_n)_{n < k}} \downarrow.$$

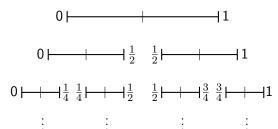
Suppose that  $(y_n)_{n < m}$  is an initial segment that has already been computed. Deciding, whether there is an extension  $(y_n)_{n < l}$ , such that

 $\Phi_e^{(y_n)_{n < l}} \downarrow$ 

can be done in 0'.

## Computing $BW_{weak}$

Assume that  $(x_n)_n \subseteq [0,1] \cap \mathbb{Q}$ . Goal: Construct a subsequence  $(y_n)$  with  $\forall n \exists k \forall i, j \geq k \ (|y_i - y_j| < 2^{-n})$ and compute the Turing jump of  $(y_n)$ .



- bi-partition argument
- now add at each step not only one element but finitely many elements of the chosen interval to (y<sub>n</sub>).
- Let  $(y_n)_{n < m}$  be the initial segment of  $(y_n)$  computed up to the k-th step. At the k-th step extend this to  $(y_n)_{n < l}$  by elements in the k-th
  - chosen interval, such that  $\Phi_k^{(y_n)_{n < l}}\downarrow$ , if possible.
- Then extend this by another element of the interval.

## Theorem (K.)

For each bounded, computable sequence  $(x_n)$ there is a Cauchy-subsequence  $(y_n)$ , such that  $(y_n)$  and  $(y_n)'$  are computable in a Turing degree that contains infinite branches of 0'-computable 0/1-trees.

#### Corollary

 $BW_{weak}$  has  $low_2$  solutions, i.e.  $(y_n)''$  is computable in 0''.

## Proof.

$$(y_n)' \leq_T 0' + \mathsf{WKL} \implies (y_n)'' \leq_T 0''$$

 $\mathsf{BW}_{\mathsf{weak}}$  does not compute 0' and is therefore strictly weaker than BW.

## The cohesive principle

Write  $X \subseteq^* Y$  if  $X \setminus Y$  is finite.

## Definition

• A set X is *cohesive* for a sequence of set  $(R_n)_n \subseteq 2^{\mathbb{N}}$  if

$$X \subseteq^* R_n \lor X \subseteq^* \overline{R_n}$$
 for each  $n$ .

• The *cohesive principle* (COH) states that for each  $(R_n)_n$  there is an infinite cohesive set X.

## Theorem (K.)

- For each sequence (x<sub>n</sub>)<sub>n</sub> ⊆ ℝ there exists (R<sub>n</sub>)<sub>n</sub> ⊆ 2<sup>N</sup>, such that from an infinite cohesive set for (R<sub>n</sub>) one can compute a Cauchy-subsequence of (x<sub>n</sub>) and vice versa.
- RCA<sub>0</sub> ⊢ BW<sub>weak</sub> ↔ COH ∧ BΣ<sub>2</sub><sup>0</sup> Moreover, this equivalence also holds instance-wise.

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#### Theorem

- COH and hence also BW<sub>weak</sub> do not compute solutions to WKL in general. (Cholak, Jockusch, Slaman '01)
- There are instance of these principle which have no low solutions. (Jockusch, Stephan '93)

Proof of the  $low_2$ -ness of BW<sub>weak</sub> is a streamlined version of the  $low_2$ -ness of COH (Jockusch, Stephan '93).

## Theorem (Chong, Slaman, Yang '10)

 $\mathsf{RCA}_0 + \mathsf{COH} + B\Sigma_2^0$  and thus  $\mathsf{RCA}_0 + \mathsf{BW}_{\mathsf{weak}}$ are  $\Pi_1^1$ -conservative over  $\mathsf{RCA}_0 + B\Sigma_2^0$ .

#### Theorem (K., Kohlenbach '10)

If WKL<sub>0</sub> + BW<sub>weak</sub>  $\vdash \forall f \exists y \phi(f, y)$ for quantifier free  $\phi$ , then one can extract from a given proof a **primitive recursive** function(al) t such that  $\forall f \phi(f, t(f))$ .

"Proof mining"

## Bolzano-Weierstraß in the weak topology

We consider the Hilbert space  $\ell_2 = (\mathbb{R}^{\mathbb{N}}, \langle \cdot, \cdot \rangle)$ . An element of  $\ell_2$  is given by a Cauchy-sequence  $(w_n)_n$  of finite dimensional and rational approximations, i.e.  $w_n \in \mathbb{Q}^{<\mathbb{N}}$ , with Cauchy-rate  $2^{-n}$  with respect to  $\|\cdot\|$ .

### Definition

(weak-BW): Every  $\|\cdot\|$ -bounded sequence  $(x_n) \subseteq \ell_2$  has a weak cluster point x, i.e.  $\forall y \in \ell_2 \lim_{n \to \infty} \langle y, x_n \rangle = \langle y, x \rangle$ .

## Theorem (K.)

- For each bounded sequence (x<sub>n</sub>) ⊆ ℓ<sub>2</sub> there is a weak cluster point x computable in 0".
- There is a bounded and computable sequence (x<sub>n</sub>) ⊆ ℓ<sub>2</sub>, such that each weak cluster point of it computes 0".
- Over RCA<sub>0</sub> the principles Π<sup>0</sup><sub>2</sub>-CA and weak-BW are instance-wise equivalent.

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- $\bullet~$  BW is equivalent to WKL for  $0'\mbox{-}computable$  trees.
- BW<sub>weak</sub> is equivalent to COH.
  - Hence, it does not imply 0'.
  - It admits extraction of primitive recursive terms.
- weak-BW is equivalent to 0''.

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The cohesive principle and the Bolzano-Weierstraß principle, Math. Log. Quart. **57** (2011), no. 3, 292–298.

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